



# **Sunbeam Institute of Information Technology**

## **Pune and Karad**

### **Module – Operating System Concepts**

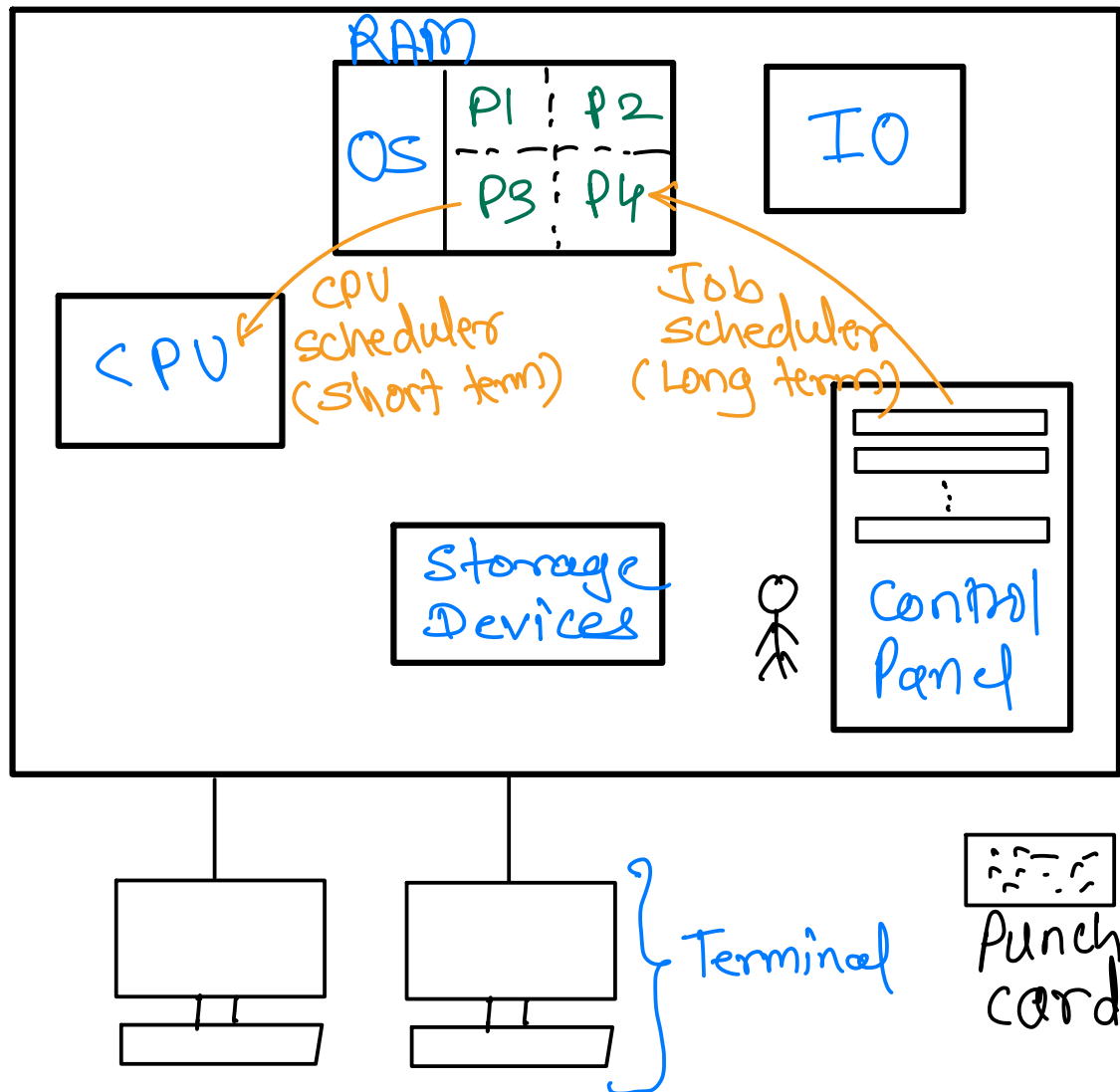
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# Types of Operating System

- OSs are classified by looking at targeted hardware
    - 1) Desktop OS (GPOS) - responsiveness
    - 2) Server OS - throughput
    - 3) Handheld OS
    - 4) Embedded OS
    - 5) Real Time OS
    - 6) Distributed OS - load balancing
- } small footprint (size)

# Types of Operating System



1) Resident Monitor System

2) Batch Systems

3) Multi-programming system

- multiple programs are loaded into memory (RAM)

Degree of multiprogramming:

↳ no. of programs loaded in RAM

CPU time/burst - time spent on CPU

IO time/burst - time spent for IO

$\text{CPU burst} > \text{IO burst}$  } CPU bound process

$\text{IO burst} > \text{CPU burst}$  } IO bound process

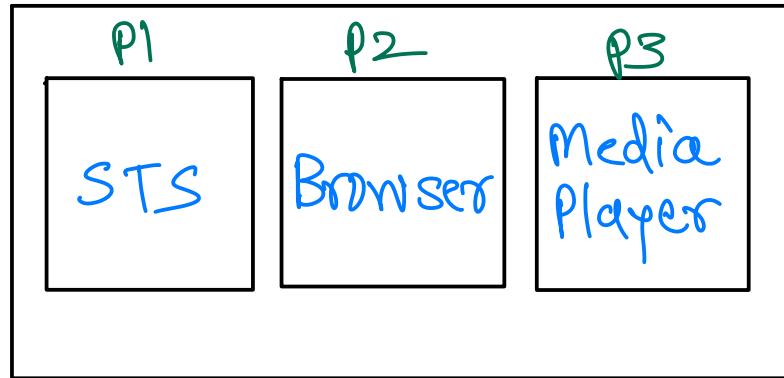
- mixture of CPU bound & IO bound processes is loaded into RAM

# Types of Operating System

## 4) Time sharing system / Multitasking System

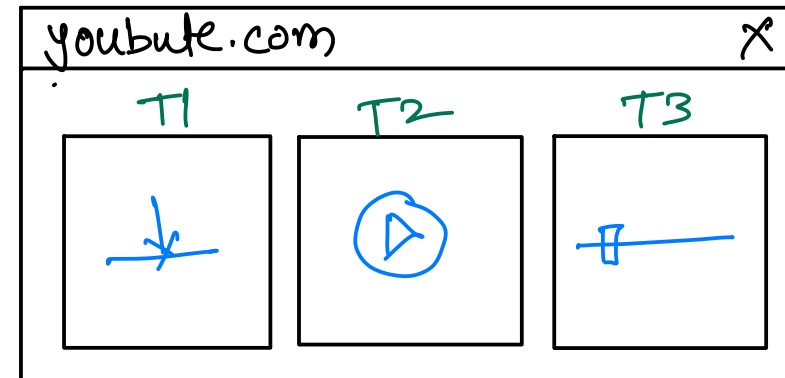
- CPU Time is shared in all the processes of RAM
- Response time  $< 1\text{sec}$

### i) Process based multitasking



- system wide multitasking

### ii) Thread based Multitasking (Multithreading)



- multitasking within process

## 5) Multi User System

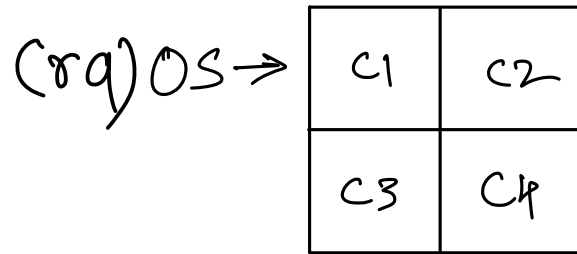
- multiple terminals (monitor + keyboard) are connected to single system
- due to this multiple users can operate single system
- whoami, who, w, tty

# Types of Operating System

## 6) Multiprocessing System

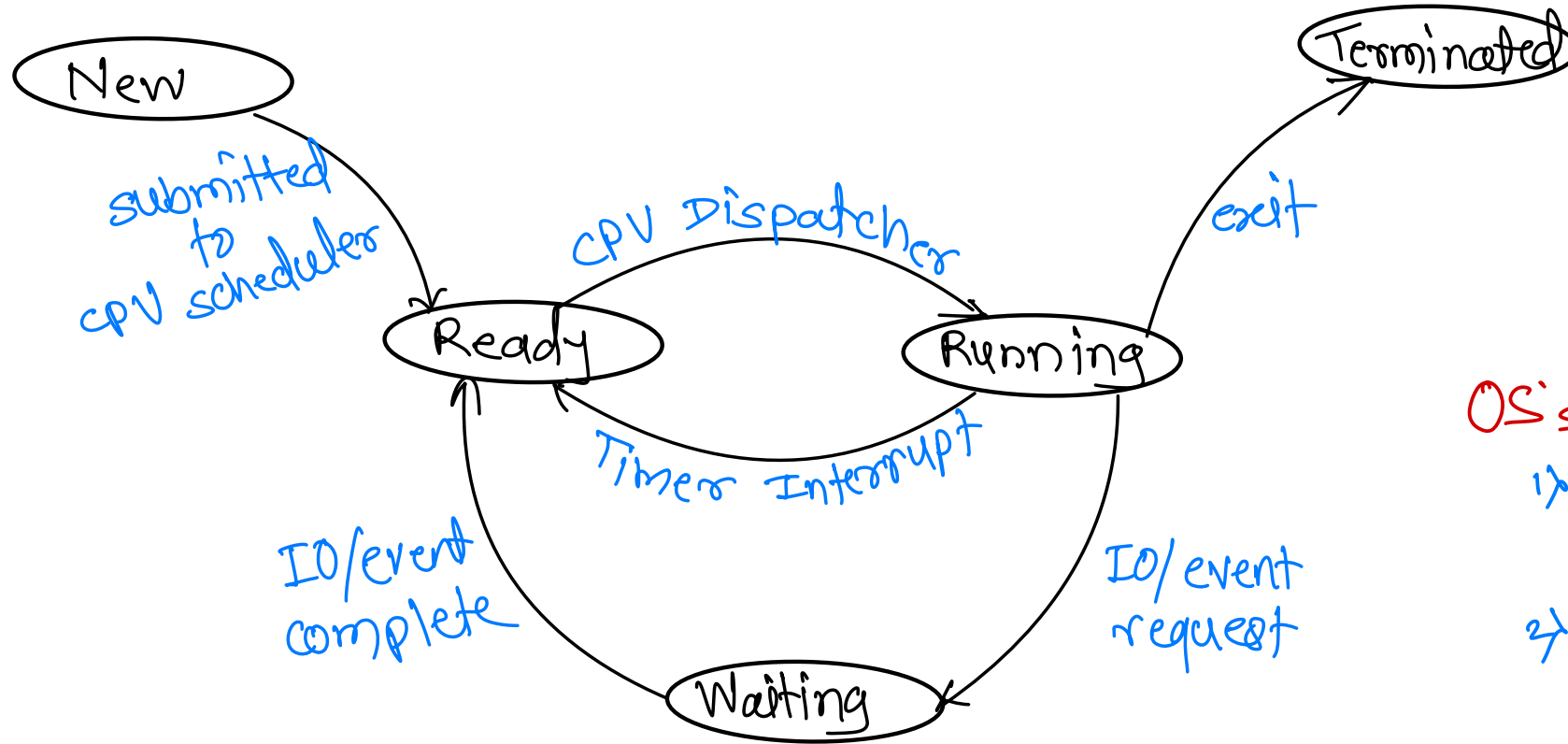
- multiple CPUs are putted on single chip, such processor is known as "multiprocessor" / "multicore".
- OS can schedule multiple processes for those multiple cores
- multiple processes can execute simultaneously, due to this it is also known as parallel system.

i) Symmetric multiprocessing      ii) Asymmetric multiprocessing



Windows Vista } multiprocessing was  
 linux 2.6+      supported from these versions

# Process Life Cycle



OS's data structure:

1) Job queue/process list  
- all processes of system

2) Ready queue  
- all ready processes

3) Waiting queue  
- processes waiting for IO/event  
- multiple in number

# CPU Scheduling – Types and Algorithms

- 1) Running → Terminated
  - 2) Running → Waiting
  - 3) Running → Ready
  - 4) Waiting → Ready
- } voluntarily
- } forcefully

## Algorithms:

- 1) FCFS
- 2) SJF
- 3) Priority
- 4) RR
- 5) Fair Share

## Types:

- 1) Preemptive scheduling
  - CPU access is given to another process forcefully
- 2) Non-preemptive scheduling
  - CPU access is given to another process voluntarily.

# CPU Scheduling – Criteria's

## 1) CPU Utilization (Max)

- Desktop OS - 70%    - Server OS - 90%

## 2) Throughput (Max)

- Amount of work done in unit time

## 3) Waiting time (min)

- time spent by process into ready queue to get CPU access

## 4) Response time (min)

- time from arrival of process into ready queue upto first time getting scheduled.

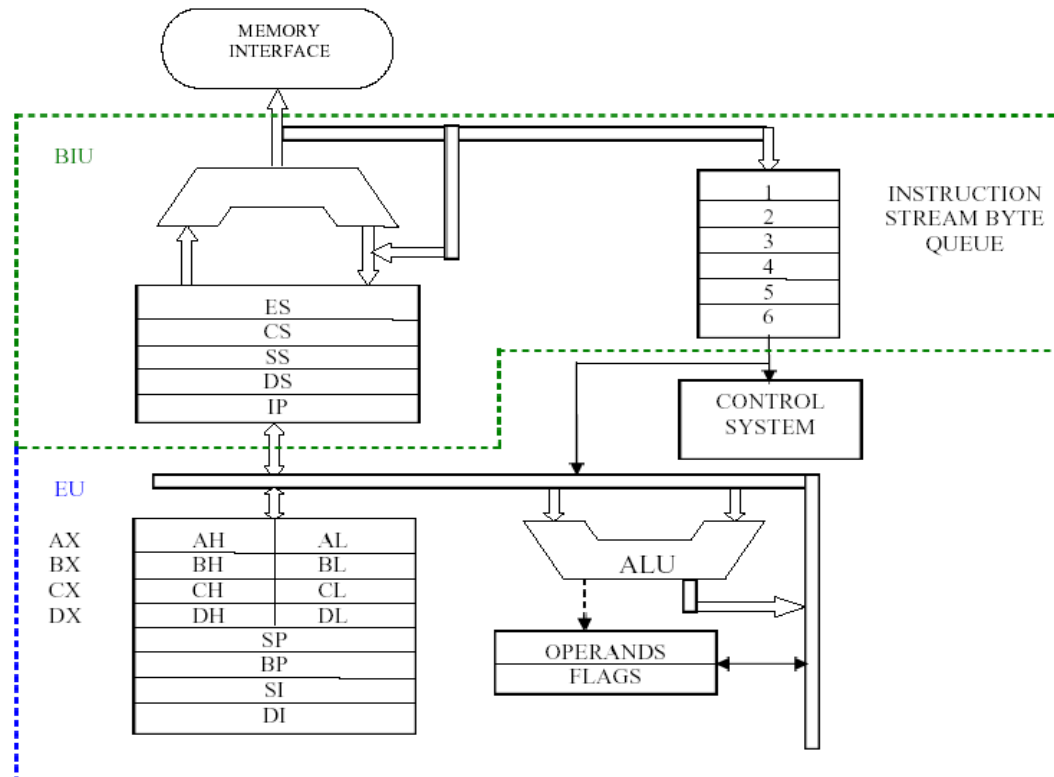
## 5) Turn Around time (min)

- Total time spent by process into memory (RAM)

$$TAT = \text{CPU waiting} + \text{CPU burst} + \text{IO waiting} + \text{IO burst}$$



# Execution Context and Context Switching



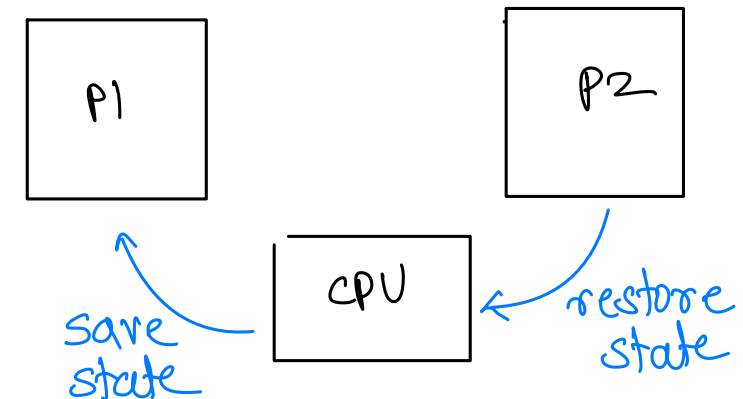
8086 Architecture

Execution context:

- values of CPU registers

Context switching:

- changing the process of CPU
- CPU dispatcher is responsible for context switching.



# FCFS (First Come First Serve) (non preemptive)

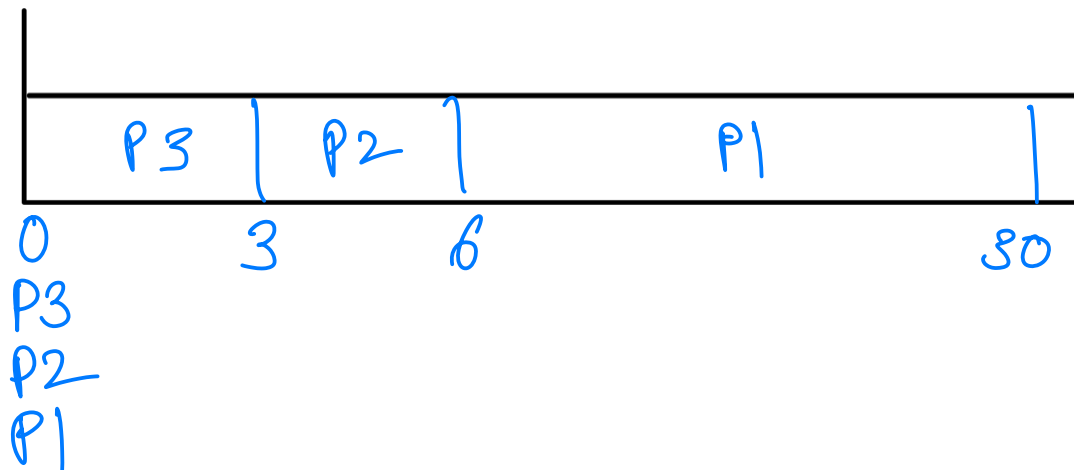
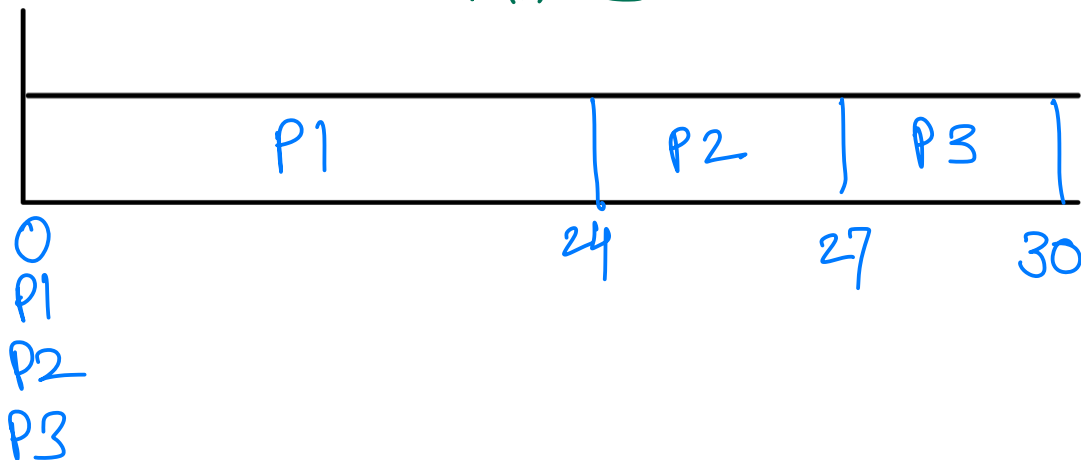
Process	Arrival	CPU Burst
P1	0	24
P2	0	3
P3	0	3

WT RT TAT  
 0 0 24  
 24 24 27  
 27 27 30

Process	Arrival	CPU Burst
P3	0	3
P2	0	3
P1	0	24

WT RT TAT  
 0 0 3  
 3 3 6  
 6 6 30

Gantt's chart



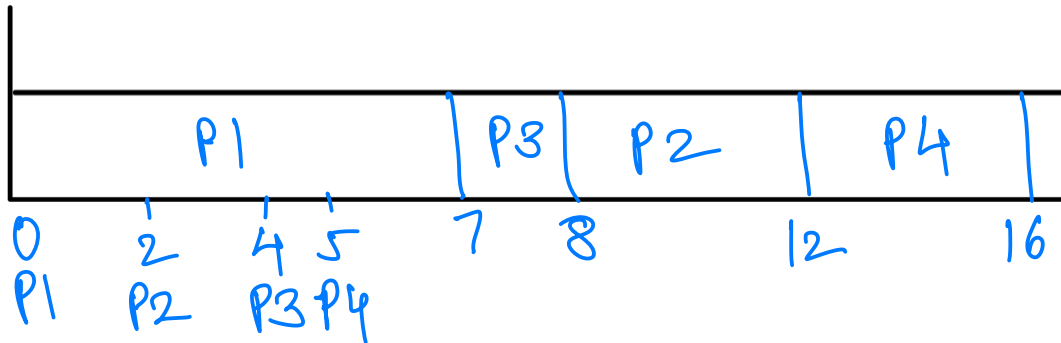
**Conroy effect:** due to arrival of longer process early all other processes need to wait for longer time

# SJF (Shortest Job First)

(Non Preemptive)

Process	Arrival	CPU Burst
P1	0	7
P2	2	4
P3	4	1
P4	5	4

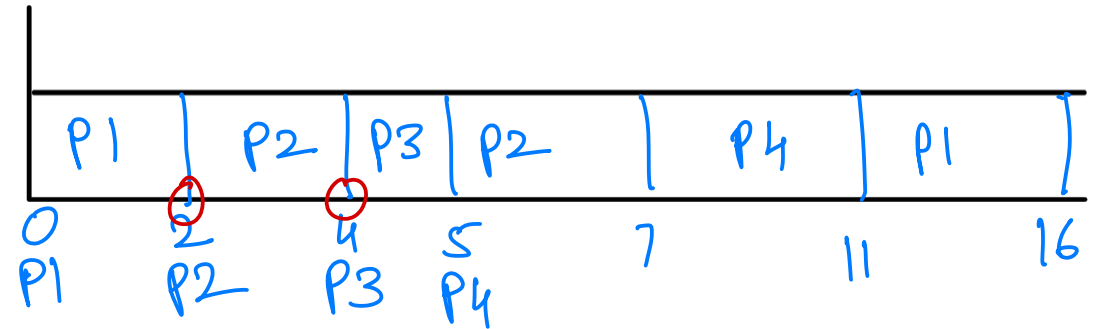
WT	RT	TAT
0	0	7
6	6	10
3	3	4
7	7	11



Shortest Remaining Time First  
(Preemptive)

Process	Arrival	CPU Burst
P1	0	7
P2	2	4
P3	4	1
P4	5	4

Remain time	WT	RT	TAT
5	9	0	16
2	1	0	5
	0	0	1
	2	2	6
		0.5	



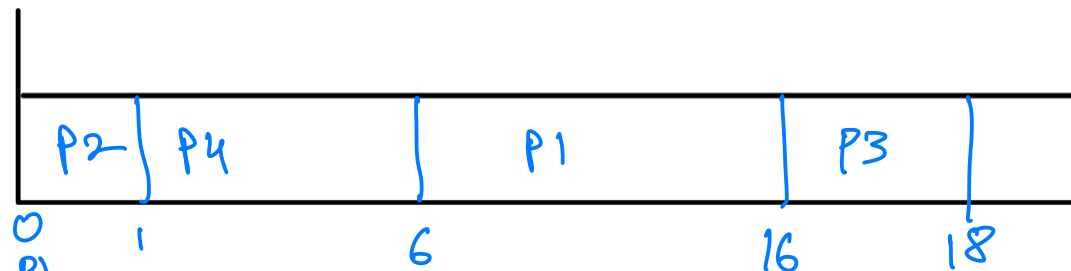
**Starvation:**  
due to longer CPU burst, process doesn't get scheduled for duration.

# Priority

(Non-Preemptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	0	1	1 (H)
P3	0	2	4 (L)
P4	0	5	2

WT RT TAT  
 6 6 16  
 0 0 1  
 16 16 18  
 1 1 6



0  
P1  
P2  
P3  
P4

P1 (7)  
P2 (4)  
P3 (5)  
P4 (6)  
P5 (5)  
P6 (5)

P2  
P3  
P5  
P4  
P6  
P4  
P1

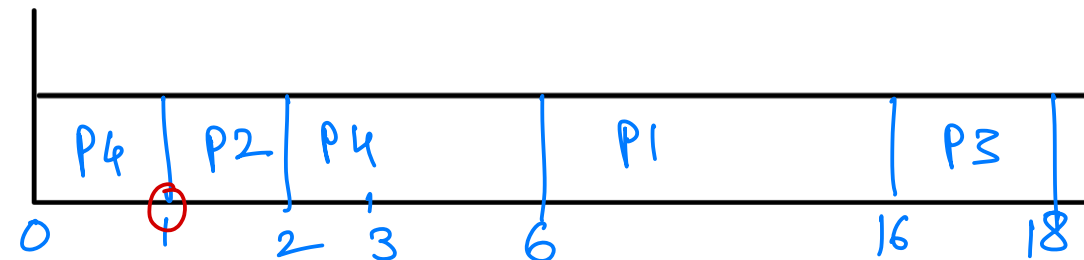
**Starvation:**  
 due to less priority,  
 process don't get enough  
 time to execute.

**Aging:**  
 increase the priority of  
 starved process gradually

(Pre-emptive)

Process	Arrival	CPU Burst	Priority
P1	0	10	3
P2	1	1	1 (H)
P3	3	2	4 (L)
P4	0	5	2

WT RT TAT  
 6 6 16  
 0 0 1  
 13 13 15  
 1 0 6



0  
P1  
P4

P2  
P3

# RR (Round Robin)

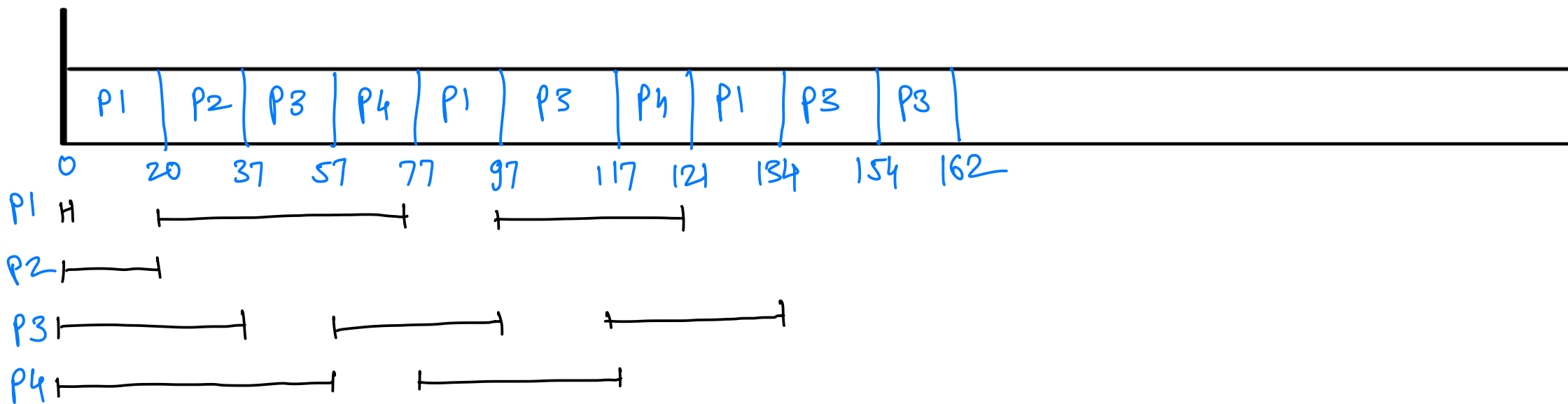
Process	CPU Burst
P1	53
P2	17
P3	68
P4	24

	WT	RT
P1	<u>0</u> + 57 + 24	0
P2	<u>20</u>	20
P3	<u>37</u> + 40 + 17	37
P4	<u>57</u> + 40	57

Time quantum = 20

TG = 100  
↳ behave like  
FCFS

TG = 4  
↳ CPU overhead  
will increase



# Fair Share

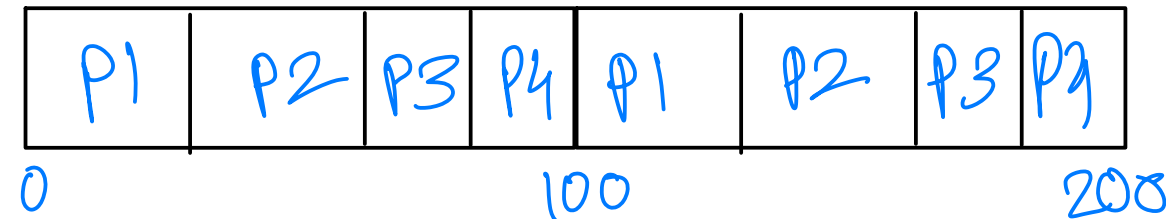
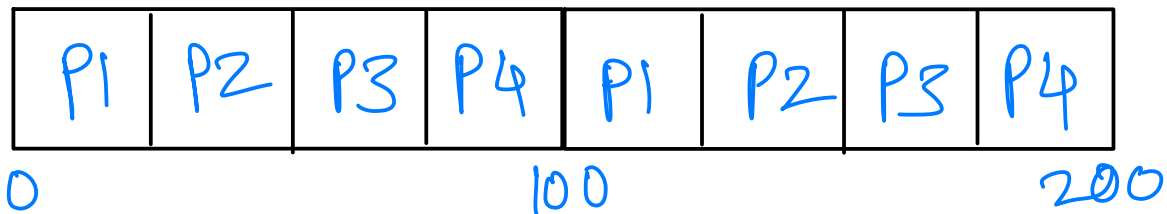
- CPU time is divided into time slices (epoch)
- some share of each epoch is given to the processes which are in ready queue.
- share is given to the process on the basis of their priority
- priority of every process is decided by its nice value
- nice values range ---> -20 to +19 (40 values)
  - \* -20 - highest priority
  - \* +19 - lowest priority

Linux scheduler:  
completely fair  
scheduler (CFS)

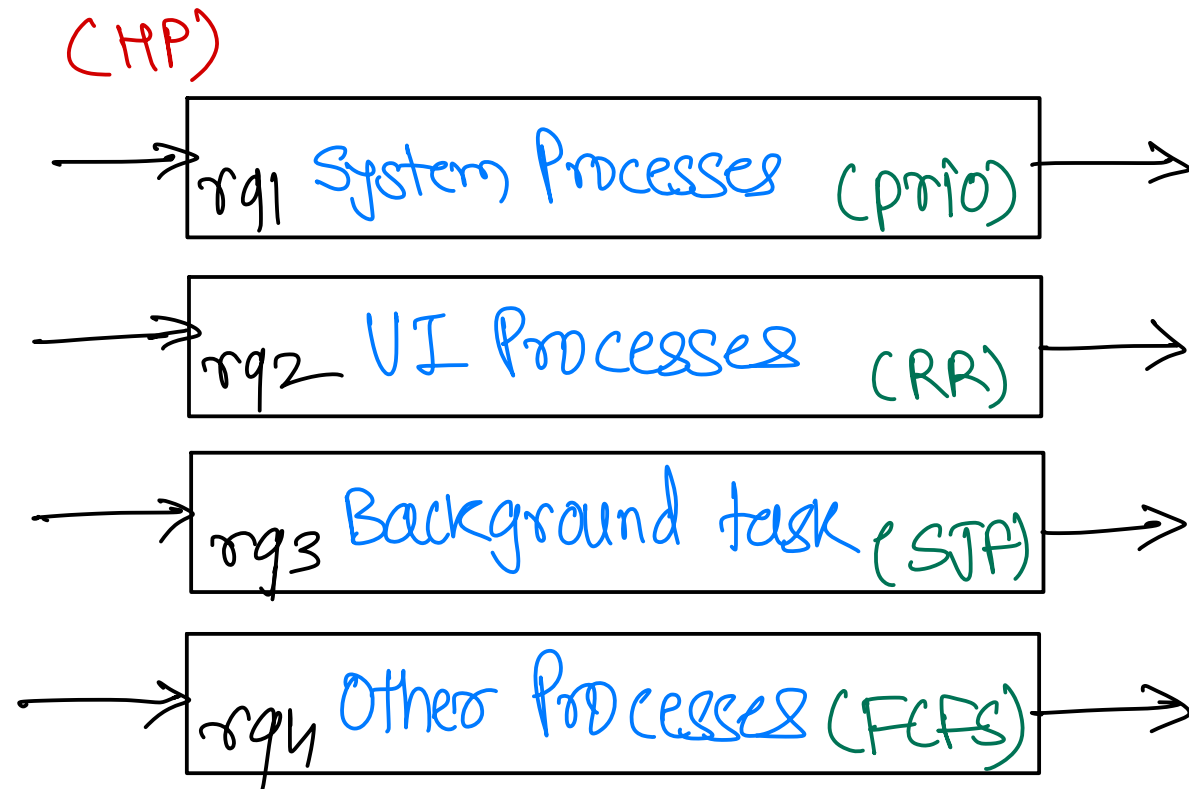
Process	Nice Value
P1	10
P2	10
P3	10
P4	10

Epoch - 100

Process	Nice Value
P1	5
P2	5
P3	10
P4	10

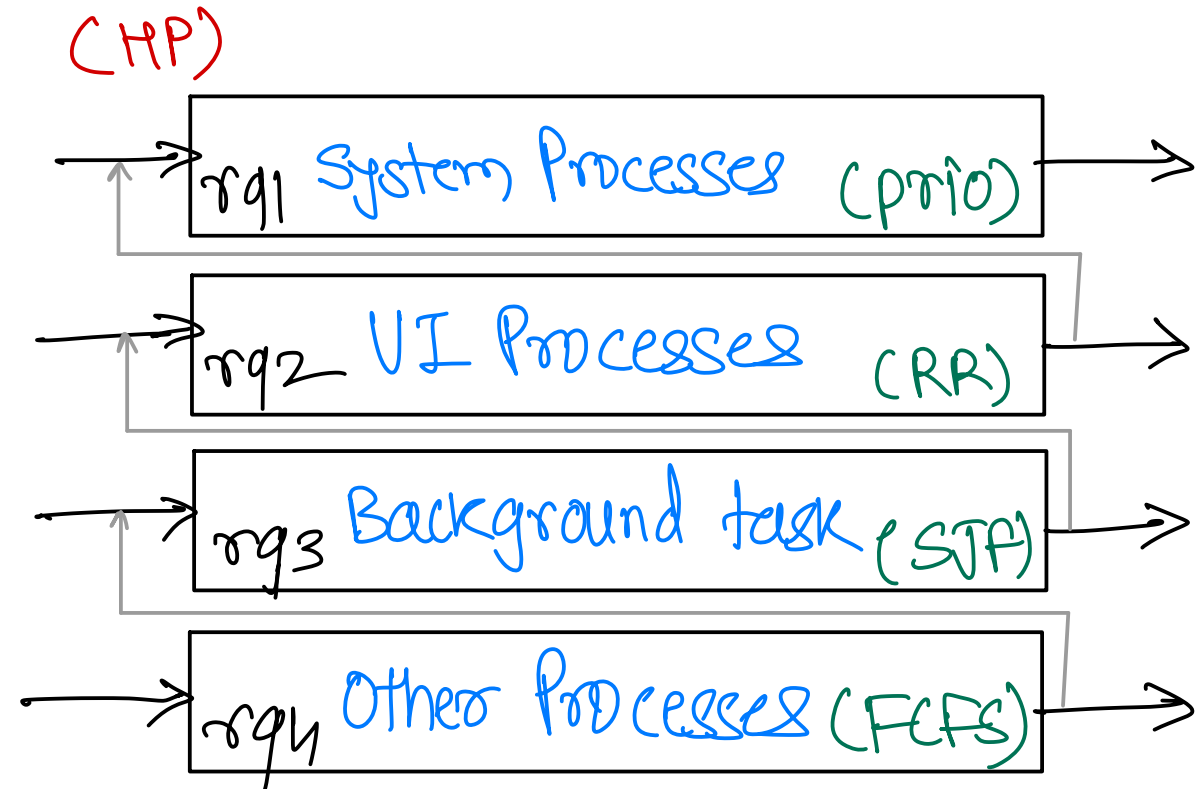


# Multi Level Ready Queue



(LP)

Multi level



(LP)

Multi Level Feedback



Thank you!!!

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