

COL216 Assignment 8

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To Run:

(A) To give instructions in form of MIPS assembly code:

1. Set `syntax_input = true` in `main.cpp` (It has been set to true already while submitting)
2. Run `'make'` to create an executable `'main'`
3. Enter the required instructions in `'input.txt'` (one test case already filled in while submitting). The registers could be like `$8`, `$9`, etc or `$t1`, `$t2`, etc
4. Run `'./main'`

(B) To give instructions in form of 32-bit instructions directly:

1. Set `syntax_input = false` in `main.cpp`
2. Run `'make'` to create an executable `'main'`
3. Enter the required instructions in `'binary_input.txt'` (one test case already filled in while submitting)
4. Run `'./main'`

Output:

1. The state for each clock cycle is printed in the console. For each clock cycle, following is shown:
 - a. Clock cycle no.
 - b. Current instruction in each stage of the pipeline (Here instruction nos. refer to the line no in `input_clean.txt`, read more below)
 - c. Details about any read/write from data memory (Memory address and value)
 - d. Details about any write to the register file (Register no. and value)
 - e. The status of the register file, i.e. values in each register
2. At the end, the following 2 statistics are printed:
 - a. Total clock cycles
 - b. IPC

Documentation

Implementation structure

1. If the input is in the form of assembly code, we first 'clean' the file by removing unnecessary spaces and newlines. We parse instructions into tokens (delimited by space, bracket, comma) and create a 'clean' input file `input_clean.txt`. **All the instruction nos. in the output refer to corresponding line nos. in `input_clean.txt`**
2. We then convert each instruction into 32-bit binary instructions, and output them to '`binary_input.txt`'. We also take care of labels for jump/branch instructions, by mapping each label to corresponding instruction no first and then encoding the required jump/branch value in the instruction.
3. We have represented each instruction as an object of a class 'Instruction'.
 - a. This object stores the instruction string, rd/rs/rt register nos, values read while decoding, result of alu operation, etc. Thus it acts as a proxy for the pipeline registers, while making it easy to organise and print data during the simulation.
 - b. This object also has functions: `instruction_decode()`, `execute()`, `memory_access()` and `write_back()` which are run in the corresponding stage of the pipeline the instruction is in. Each of these functions is implemented in the form of switch-case, with the action depending on the type of instruction.
4. We have implemented stalling by creating a function `stall()`, which stalls instruction $\geq i$ (the i th instruction and instructions after it) for n cycles. This function is called in case of data hazards and for branch hazards accordingly, during the instruction decode stage.
 - a. Data hazards is checked using a function '`is_data_hazard()`', which checks register nos of instructions already in the pipeline. We check if the register which we want to read data from (is **rs** or **rt** of current instruction) is equal to the **rd** of an `add/sub/sll/srl` or **rt** of an `lw` instruction already in the pipeline. If yes, then there's a data hazard.
 - b. In case of jump instructions we stall for 1 cycle, and for branch instructions stall for 2 cycles
5. For each clock cycle, a new instruction is fetched (if not stalled), and the instructions in the pipeline run their corresponding functions (depending on the pipeline stage they are in) (if not stalled)

Test cases:

For all test cases, we have assumed the following initialisation of the register file:

reg_array[0] = 2

reg_array[1] = 1

reg_array[29] = 4095

Others: 0

S. No	Instructions	Details, Final register file	No of clock cycles (without forwarding)
1	sw \$0 , 1024(\$1) add \$0,\$0,\$0 sub \$1 , \$0, \$1 lw \$2, 1022(\$1)	4 3 2 0 4095 0 0	12
2	bgtz \$0, second first: add \$1, \$1, \$1 add \$2, \$1, \$0 j end second: bgtz \$0, first end:	2 2 4 0 4095 0 0	15
3	sw \$0 , 1024(\$1) sll \$0,\$0,2 srl \$0,\$0,1 sub \$1 , \$0, \$1	4 3 2 0 4095 0 0	15

	lw \$2, 1022(\$1)		
4	sw \$0 , 1000(\$0) sw \$1 , 1001(\$0) add \$2, \$1, \$0 sub \$3, \$1, \$0 sw \$2 , 1003(\$1) sw \$3 , 1004(\$1) lw \$4 , 1000(\$0) lw \$5 , 1001(\$0) lw \$6 , 1002(\$0) lw \$7 , 1003(\$0) add \$6, \$6, \$7 add \$5, \$5, \$6 add \$4, \$5, \$4	2 1 3 -1 5 3 2 -1 0 4095 0 0	24
5	add \$2 , \$0 , \$1 sub \$3 , \$2 , \$0 sw \$2 , 1024(\$1) sw \$3 , 1024(\$0) add \$0 , \$0 , \$0 add \$4 , \$0 , \$3 sw \$4 , 1024(\$0) sub \$5 , \$4 , \$3	4 1 3 1 5 4 16 0 4095 0 0	22

	sll \$6 , \$5 , 2		
6	sll \$4 , \$0 , 1 sll \$5 , \$0 , 2 jal one add \$17 , \$2 , \$16 j end one: sub \$29 , \$29 , \$0 sw \$31 , 1(\$29) sw \$4 , 0(\$29) bgtz \$4 , two add \$29 , \$29 , \$0 jr \$31 two: sub \$4 , \$4 , \$1 jal one lw \$4 , 0(\$29) lw \$31 , 1(\$29) add \$29 , \$29 , \$0 add \$2 , \$2 , \$5 jr \$31 end:	2 1 32 0 4 8 0 0 0 0 0 0 0 0 0 0 32 0 0 0 0 0 0 0 0 0 0 0 4095 0 3	89

7	sll \$4 , \$0 , 3 jal acc add \$17 , \$2 , \$16 j end acc: sub \$sp , \$sp , \$0 sw \$ra , 1(\$sp) sw \$4 , 0(\$sp) bgtz \$4 , one add \$sp , \$sp , \$0 jr \$ra one: sub \$4 , \$4 , \$1 jal acc lw \$4 , 0(\$sp) lw \$ra , 1(\$sp) add \$sp , \$sp , \$0 add \$2 , \$2 , \$4 jr \$ra end:	2 1 136 0 16 0 0 0 0 0 0 0 0 0 0 0 136 0 0 0 0 0 0 0 0 0 0 0 4095 0 2	292
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More test cases:

S.N o	Instructions	Details, Final register file
1	sw \$0 , 1024(\$1) add \$0,\$0,\$0 sub \$1 , \$0, \$1 lw \$2, 1022(\$1)	4 3 2 0 4095 0 0
2	sw \$0 , 1024(\$1) sll \$0,\$0,1 sub \$1 , \$0, \$1 lw \$2, 1022(\$1)	4 3 2 0 4095 0 0
3	sw \$0 , 1024(\$1) sll \$0,\$0,2 srl \$0,\$0,1 sub \$1 , \$0, \$1 lw \$2, 1022(\$1)	4 3 2 0 4095 0 0
4	sw \$0 , 1000(\$0) sw \$1 , 1001(\$0)	2 1 3 -1 5 3 2 -1 0 4095 0 0

	jr \$31 end:	
11	sll \$4 , \$0 , 3 jal acc add \$17 , \$2 , \$16 j end acc: sub \$sp , \$sp , \$0 sw \$ra , 1(\$sp) sw \$4 , 0(\$sp) bgtz \$4 , one add \$sp , \$sp , \$0 jr \$ra one: sub \$4 , \$4 , \$1 jal acc lw \$4 , 0(\$sp) lw \$ra , 1(\$sp) add \$sp , \$sp , \$0 add \$2 , \$2 , \$4 jr \$ra end:	Non leaf procedure Accumulator with n = 16 Argument in \$4 \$2 is \$v0 \$17 contains final answer 2 1 136 0 16 0 0 0 0 0 0 0 0 0 0 0 136 0 0 0 0 0 0 0 0 0 0 0 4095 0 2

	lw \$31 , 1(\$29) add \$29 , \$29 , \$0 add \$2 , \$2 , \$5 jr \$31 end:	
14	bgtz \$0, second first: add \$1, \$1, \$1 add \$2, \$1, \$0 j end second: bgtz \$0, first end:	Negative offset on branch Register file: 2 2 4 0 4095 0 0