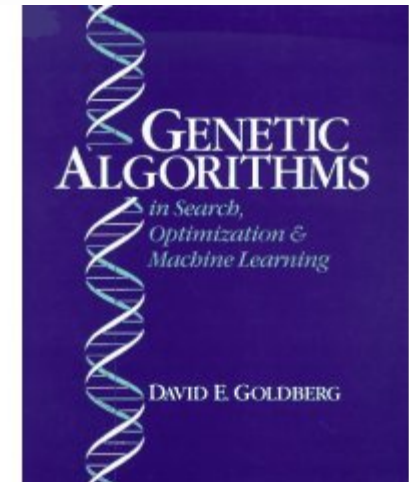




Lecture 4: Schemata & Schema Theorem

Genetic Algorithms and Other Evolutionary Techniques

- David E. Goldberg “Genetic Algorithms in Search, Optimization and Machine Learning” Addison Wesley (1 Jan 1989)
- David E. Goldberg “The Design of Innovation” Springer-Verlag 1992



- What does a genetic algorithm process?
- Schemata, the similarity templates
- Schema Theorem
- Building Block Hypothesis
- Implicit Parallelism

x	$fitness$
10110111	1.00
00110100	4.00
11011011	1.00
00101000	5.00
11011111	1.00

Can you suggest a better individual?

x	$fitness$
10110111	6.00
00110100	4.00
11011011	6.00
00101000	2.00
11011111	7.00

Can you suggest a better individual?

What approach?

Define similarity among individuals

Study evolution of similarities
between candidate solutions

Study evolution of groups of
similar candidate solutions

- To describe similarity among individuals in the population we use the concept of schema (schemata)
- What is a schema?
 - Similarity template describing a subset of strings with similarities at certain positions
 - String over the alphabet $\{0,1,*\}$
 - $*$ is the don't care symbol
- The don't care symbol $*$ is a metasympol:
it is never explicitly process by the genetic algorithm

- Schema $***II$ matches,
 - 000II
 - 00III
 - 0I0II
 - ...
 - IIIII
- Schema $0***I$ matches,
 - 0000I
 - 000II
 - 00I0I
 - ...
 - 0IIII

- Order of schema H , $o(H)$
the number of specific (non-*) positions in the schema
- Examples
 - $o(**|*0) = 2$
 - $o(|**||) = 3$
- Defining length of schema H , $\delta(H)$
distance between first and last non-* in a schema
- Examples
 - $\delta(**|*0) = 2$
 - $\delta(|**||) = 4$

- Schema $H1$ is a competitor of schema $H2$ if,
 - $H1$ and $H2$ are specified in the same positions.
 - $H1$ and $H2$ differ in at least one specified position.
- Example
 - $*| | |*$ is a competitor of $*|0|*$

- A string of length l is an instance of $2^l - 1$ schemata
- But how many schemata are there in the whole search space? (how many choices each locus?)
- Since one string instances $2^l - 1$ schemata, how much does a population tell us about schemata of various orders?
- Implicit parallelism
 - One string's fitness tells us something about the relative fitness of more than one schema

- We evaluate solutions at the level of full strings, solutions, or structures
- We modify a structure at the level of modules, substructures, components, etc.
- Trick of genetic optimization: Only get evaluation at the highest level.
- How can we get best substructures without explicit evaluation of them?

How a Simple Genetic Algorithm Works?

- A simple genetic algorithm
 - Population size N
 - Fixed-length binary strings of length l
 - Fitness-proportionate selection
 - One-point crossover with probability p_c
 - Bit-flip mutation
 - Fitness is positive
- Notation
 - $A(t)$ population at time t
 - A_i i -th individual in $A(t)$
 - $f(H,t)$ fitness of schema H at time t
 - $m(H,t)$ number of examples of H in $A(t)$
 - f_A average fitness in $A(t)$
- How selection, crossover, and mutation work?

How a Simple Genetic Algorithm Works? Proportionate Selection

- During reproduction, A_i is selected with probability p_i ,

$$p_i = \frac{f_i}{\sum f_j}$$

- After reproduction, we expect,

$$m(H, t + 1) = m(H, t) \frac{f(H, t)}{f_A}$$

- The equation,

$$m(H, t + 1) = m(H, t) \frac{f(H, t)}{f_A}$$

- can be rewritten as,

$$\begin{aligned} m(H, t + 1) &= m(H, t) \frac{f_A + cf_A}{f_A} \\ &= m(H, t)(1 + c) \\ &= m(H, 0)(1 + c)^t \end{aligned}$$

Reproduction allocates exponentially increasing/decreasing number of trials of above/below average schemata

How a Simple Genetic Algorithm Works? Crossover

- Schema H survives crossover with probability p_s ,

$$p_s = 1 - \delta(H)/(l - 1)$$

- When crossover probability is taken into account,

$$p_s \geq 1 - p_c \delta(H)/(l - 1)$$

- So that,

$$m(H, t + 1) \geq m(H, t) \frac{f(H, t)}{f_A} \left[1 - p_c \frac{\delta(H)}{l - 1} \right]$$

How a Simple Genetic Algorithm Works? Mutation

- Each allele survives with probability

$$(1 - p_m)$$

- The schema H survives when all the $o(H)$ alleles survive

$$(1 - p_m)^{o(H)}$$

- When p_m is very small,

$$(1 - p_m)^{o(H)} \approx 1 - o(H)p_m$$

- So that,

$$m(H, t + 1) \geq m(H, t) \frac{f(H, t)}{f_A} \left[1 - p_c \frac{\delta(H)}{l - 1} - p_m o(H) \right]$$

- Define the probability of survival $P_s(H,t)$,

$$P_s(H, t) = \left[1 - p_c \frac{\delta(H)}{l-1} - p_m o(H) \right]$$

- H gets more copies when $P_s(H, t) \frac{f(H,t)}{f_A} > 1$
- H gets less copies when $P_s(H, t) \frac{f(H,t)}{f_A} < 1$
- Expected number of representatives of H grows with $\frac{f(H,t)}{f_A}$
- $P_s(H,t)$ decreases with increasing $o(H)$ and $\delta(H)$

- Two factors influence the number of representatives of H
 - Over the average schema fitness, leads to higher $m(H, t+1)$
 - Greater $o(H)$ and $\delta(H)$, lead to a smaller $m(H, t)$
- Short, low order, above-average schemata receive exponentially increasing trials in subsequent generations.

$$m(H, t + 1) \geq m(H, t) \frac{f(H, t)}{f_A} \left[1 - p_c \frac{\delta(H)}{l - 1} - p_m o(H) \right]$$

- Generalization of the Schema Theorem

$$m(H, t + 1) \geq m(H, t) \phi(H, t) [1 - \varepsilon(H, t)]$$

- Reproductive ratio $\phi(H, t)$
- Growth ratio $\gamma = \phi(H, t) [1 - \varepsilon(H, t)]$
- We want $\gamma \geq 1$

Building blocks

Low order, low defining-length schemata with above average fitness.

Building Block Hypothesis

“Short, low-order, and highly fit schemata are sampled, recombined, and resampled to form strings of potentially higher fitness [...] we construct better and better strings from the best partial solutions of the past samplings.”

David Goldberg, 1989

- A population with length l and size n , processes between 2^l to $n2^l$ schemata
- Consider only schemata that survive with a certain probability
- Count how many schemata are effectively processed
- Effective processing = selection + mixing (recombination).

Implicit Parallelism

A genetic algorithm with a population size n can usefully process on the order of $O(n^3)$ schemata in a generation despite the processing of only n solutions

No.	Old	$f(x)$	New	$f(x)$
1	11000	2	10010	2
2	10111	4	11011	4
3	01011	3	10111	4
4	00100	1	11001	3
avg		2.5		3.25

Example 1

Expected number of copies of I****

- $m(I^{****}, 0) = 2$
- $f(I^{****}, 0) = (2+4)/2 = 3 \ (3 > 2.5)$
- Survival after mutation $1 - p_m o(H) = 1 - 0.2 \times 1 = 0.8$
- Survival after crossover and mutation

$$1 - p_m o(H) - p_c \frac{\delta(H)}{5-1} = 1 - 0.2 \times 1 - 1 \times 0/4 = 0.8$$

- After selection $m(I^{****}, 1) = 2 \times 3 / 2.5 = 2.4$
- After selection, crossover, and mutation
 $m(I^{****}, 1) , 2.4 \times 0.8 > 1.92$

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3	01011	3	10111	4
4	00100	1	11001	3
avg		2.5		3.25

- $m(***|*, 0) = 2$
- $f(***|*, 0) = (3 + 4)/2 = 3.5 \ (3.5 > 2.5)$
- $m(***|*, 1) = 3 \ (+1)$
- $f(***|*, 1) = (2 + 4 + 4)/3 = 3.33 \ (-0.17)$

No.	Old	$f(x)$	New	$f(x)$
1	11000	2	10010	2
2	10111	4	11011	4
3	01011	3	10111	4
4	00100	1	11001	3
avg		2.5		3.25

- $m(\text{****}0, 0) = 2$
- $f(\text{****}0, 0) = (2 + 1)/2 = 1.5 \ (1.5 < 2.5)$
- $m(\text{****}0, 1) = 1 \ (-1)$
- $f(\text{****}0, 1) = 2 \ (+0.5)$

- Schemata are similarity templates
- High performance, short, low-order schemata receive at least exponentially increasing number of trials in successive generations
- By processing similarities, a genetic algorithm reduces the complexity of an arbitrary problems
- Highly fit, short, low order schemata become the partial solutions to a problem (these are called building-blocks)