

Today's Content :-

- Comparing 2 Algs
- why Big O needed
- Space Complexity
- TLE
- TC naming conventions
- 1 advance problem in calculating iterations.

Context :- Given n elements sort them in
as ascending order.

$$arr[S] = \{3, 2, 6, 8, 13\} \rightarrow \{1, 2, 3, 6, 8\}$$

$$\begin{cases} & \\ \rightarrow & N = 10^4 \end{cases}$$

Vishal (Algo 1)



15 sec

↓
(window 1 x 1)



(Macbook M2)



7 sec



C++



(Top of mt. volcano)



mt. Everest



5 sec

Rohan (Algo 2)



10 sec

↓
(Macbook M2)



10 sec



Python



C++



5 sec



(mt. Everest)

Execution time :- It depends on so many external factors, hence we generally don't compare execution time b/w 2 algorithms.

Iterations, $\text{for } (i=1; i \leq n; i++) \{ \dots \}$

generations
 n

↓
3
print(i)

Con1:- To compare 2 algorithms, calculate their execution, based on input size and compare them.

$\Theta \rightarrow$

Algo 1 (Taw)

$\log_2 n$

Algo 2 (Saket)

$n^{\frac{1}{10}}$

Obs :-

Till $n \leq 350$

$n > 350$

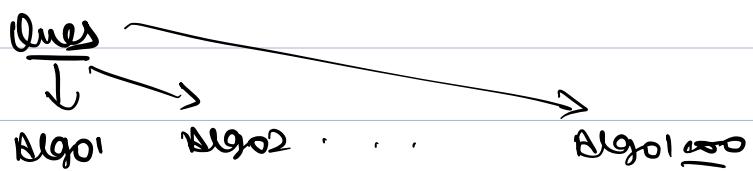
Generations

Taw > Saket

Saket > Taw

Find us pat : 10 - 11 million
Google result : millions results.

Pick algorithm, which works better for very large input.
cel2 :- ↑



Asymptotic Analysis of Algorithms

- ↳ Analyse performance of an algo for very large inputs.
- ↳ Big(O) → Notation

Algo 1 (Taru)

$\sim \log_2 N$

Algo 2 (Saket)

$N^{1/2}$

$$\Theta(\log_2 N) < O(N)$$

To calculate Big(O)

→ Calculate Iterations

→ Take higher order terms (Ignore lower order)

→ Ignore constant co-efficients.

Neglect lower order terms?

0 → Algo (Priyanka)

Generations

$$N^2 + 10N$$

Input

$$N = 10$$

Total Generations

$$200$$

% of lower order terms
in total Generations.

$$\frac{10}{200} \times 100 = 5\%$$

$$N = 100$$

$$10^4 + 10^3$$

$$\frac{10^3}{10^4 + 10^3} \times 100 \approx 9\%$$

$$N = 10^4$$

$$10^8 + 10^5$$

$$\frac{10^5}{10^8 + 10^5} \times 100 \approx 0.1\%$$

Obs:- as Input size increases, contribution
of lower order term decreases.

Neglect Constant Co-efficients

$\Theta \rightarrow$

Algo1 (howind) Algo2 (hima) for large Inputs

$10 \log_2 N$

N

howind

$10 \log_2 N$

N

howind

$10^3 \log_2 N$

N_{10}

howind

$10 N$

$\frac{N^2}{10}$

howind

$O \log N$

$10 N$

hima

Issues in Big O

$\Theta \rightarrow$ Aminash (Alg1)	Path (Alg2)	more optimized.
$10^3 N$	N^2	
$O(N)$	$O(N^2)$	
BigO		Aminash always better?

Input

$$N=10 \rightarrow 10^4$$

$$N=10 \rightarrow 10^5$$

$$N=10^3 \rightarrow 10^6$$

$$10^2 \rightarrow (\text{Path})$$

$$10^4 \rightarrow (\text{Path})$$

$$10^6 \rightarrow (\text{Both are same})$$

$$N=10^3+1 \rightarrow 10^3(10^3+1) \quad (10^3+1) * (10^3+1) \quad \underline{\text{Aminash}}$$

$$N=10^4 \rightarrow 10^7 \quad 10^8 \rightarrow \underline{\text{Aminash}}$$

claim :- for all Input $\geq 10^6$, Aminash is better.

final claim :-

when we compare 2 algo using Big O, Alg1 will be better than alg2, for all Input value above a certain point.

↳ threshold point

- 1) After threshold Big(O) holds.
- 2) Don't worry about threshold value.

values in Big(O)

$O \rightarrow$	Algo1 (Amritesh) $2n^2 + 4n$	Algo2 (Chetan) $3n^2$
-----------------	---------------------------------	--------------------------

$\text{Big}(O) \rightarrow$	$O(n^2)$	$O(n^2)$
-----------------------------	----------	----------

(Both are same acc. to
Big O)

Err2: $5n^3 + 6n^2$	$4n^3 + 10n$
$\text{Big}(O)$ $O(n^3)$	$O(n^3)$

(Both are same
according to Big O)

Obs 2: If 2 algos have same Big(O),
we can't really compare with
Big O, we need iteration to
compare. (Co-efficients of higher
order seems in iterations)

Code :- Searching for an element = k

```
bool search (int arr[], int k){  
    int n = arr.length;  
    for (i=0; i<n; i++) {  
        if (arr[i] == k) { return true; }  
    }  
    return false;  
}
```

T.C \rightarrow $O(n)$

generation
↓
Best Worst
i n

Note:- while writing Big(O) we consider worst case iterations

Break :- 10:12 \rightarrow 10:20 pm

$$\left(\frac{3}{2}\right)^n$$

Code → Time Complexity
→ Space Complexity

Space Complexity

int → 4 B
long → 8 B.

— func (int n) {

 int m = n + 4 → 4

 int y = m * m + 4

 long z = my + 1

Total memory
= 20 B.

3

S.C. ⇒ Independent of Imp.
 ↳ O(1)

— func (int n) {

 int m = n → 4

 int y = m * m → 4

 long z = m + y → 8

// Declare an array.

 int [] arr = new int [n].

 4N

2

Total memory = 4N + 20

S.C. ⇒ O(N)

func `(int n)` {

$$\text{int } x = n \rightarrow 4$$

$$\text{int } y = x^2 \rightarrow 4$$

$$\text{long } z = x + y \rightarrow 8$$

`int [] arr = new int [n]; 4n`

`long [] [] l = new long [n] [n]`

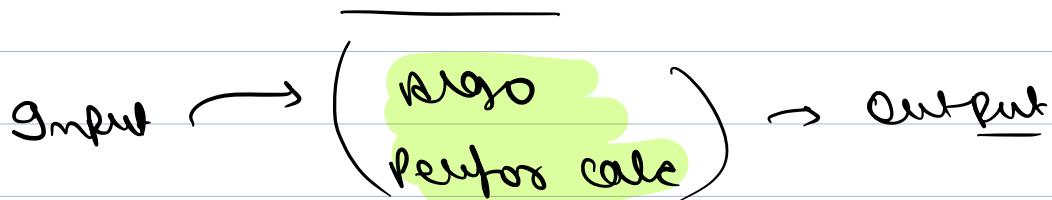
$$8n^2$$

3

$$\text{S.C} \rightarrow 8n^2 + 4n + 20$$

$$\Rightarrow O(n^2)$$

→ Algo we developed →



S.C. of an algorithm :- It is amt. of space additionally used by algorithm other than input space to perform necessary computation.
auxiliary space

// Given a program to find max.

```
int max arr (int arr[], int n) {
```

```
    int +4B ans = arr[0];
```

```
    for (i=1; i<n; i++) {
```

```
        | ans = max (ans, arr[i])
```

```
    }
```

```
3
```

T.C $\rightarrow O(N)$

S.C $\rightarrow O(1)$

// Given an array to do some task.

```
int Task (int arr[]) {
```

```
+4B int m = arr.length
```

```
int pf[n]  $\rightarrow$  4N
```

```
pf[0] = arr[0]  $\rightarrow$  4
```

```
for (i=1; i<m; i++) {
```

```
    | pf[i] = pf[i-1] + arr[i]
```

3 // perform something.

T.C \rightarrow

Iteration $\rightarrow \underline{n-1} \rightarrow O(\underline{m})$

$I.C \rightarrow (4N^2)$ $\rightarrow O(n^2)$

$T.C > S.C$ { gm most cases you'll
be asked to improve
Time first and then Space}

TLE :- (Time Limit Exceeded)



Mehta \rightarrow (Amazon) \rightarrow Hiring challenge
 $\rightarrow 20 \rightarrow$ Time.

Optimize code

$O \rightarrow$ Idea \rightarrow Code \rightarrow Submit \rightarrow TLE

\rightarrow without even writing a single line
code, we can say whether logic
will work or not.

Online Editors → code server

↳ Code time

exceeds 1 sec incⁿ

2 sec Java,
4 sec Python

$\hookrightarrow P \rightarrow 1 \text{ GHz}$

\downarrow

10^9 instructions / sec.

Obs 1:- At max our code can have 10^9

instructions

variable =
operator
function calling

Pseudo Code

but count factor (n)?

int c = 0 + 1
 for (i=1; i<=N; i++) { > few iterations
 | if (N/i == 0) 6 or 7
 | c = c + 1 + 1 instructions
 3 Total instructions
 reference c + 1 6 N or 7 N.
 3

App 1 :-

↳ In our code 1 iteration = 10¹⁰ instruction

→ Our code can almost
contains = 10^9 instructions
= $10^8 \times 10$ instruction

Our code can have at most 10^8 iterations

App 2 :-

In our code 1 iteration =
100 instructions

→ Our code can contain only
 10^9 instructions

$$= 10^7 \times 10$$

We can have only 10^7 iterations,

In general \rightarrow Code Iterations =

$[10^7 - 10^8]$ iterations

Solve a Question

Read Q

Understand Q

Logic

Correctness of logic

Code

you need
to optimize

Constraints

Q :- Read the question

Constraints

$$1 \leq N \leq 10^3$$

idea \rightarrow pseudocode

↓

T.C

↓

$$\left\{ \begin{array}{l} O(n^2) \\ \cup \end{array} \right.$$

no need to code

Q Read it

Constraints

$$1 \leq N \leq 10^3$$

idea $\rightarrow O(n^2)$

yes

Q Read it

Constraints

$$1 \leq N \leq 10^4$$

idea \rightarrow pseudocode $\rightarrow O(n^2)$

↓

WA / TLE

→ Dream Companies :- (franchises)

→ Strong hold → scales → service it
→ Affiliation
→ Time

→ resume

→ Tell me about yourself.

→ Career C.I. → ED
→ SDM
→ System Design