

Ques) Given an array elements, check if there exists a pair (i, j) s.t.

$$\text{arr}[i] + \text{arr}[j] == k \quad \&\& (i \neq j)$$

arr[] = 0 1 2 3 4 5 6 7 8 9
 8 9 1 -2 4 5 11 -6 7 5

 i j
 k = 11, 4 8

k = 6, 2 9

k = 22, 6 6 *

$$a + b = k$$

idea :-

check all pairs.

T.C $\rightarrow O(n^2)$ S.C $\rightarrow O(1)$.

- 1) Create hashset
- 2) Insert everything in it.

H.S { 8, 9, 1, -2, 4, 5, 11, -6, 7 }

```
for (i = 0; i < n; i++) {
    a = arr[i]
    b = k - a
    for (j = i+1; j < n; j++) {
        if (arr[j] == b) {
            return true
        }
    }
}
```

3

return false

k = 11

a	b	in arr.
8	3	X
9	2	X
1	10	X
-2	13	X
4	7	✓

arr =

	0	1	2	3	4	5	6	7	8	9
	8	9	1	-2	4	5	11	-6	7	5

H.S { 8, 9, 1, -2, 4, 5, 11, -6, 7 }

k = 22

a	b	
8	14	x
9	13	x
1	21	x
-2	24	x
4	18	x
5	17	x
11	11	✓

Obs:- If we know freq. we can solve problem.

1) Hash Map < int, int > hm

2) Store & insert in hm
hm

{ 8:1 4:1 7:1
 9:1 5:2 -6:1
 1:1 11:1 -2:1 }

arr[] = 0 1 2 3 4 5 6 7 8 9
 8 9 1 -2 4 5 11 -6 7 5

k = 22,

a	b	
8	14	*
9	13	*
1	21	*
-2	24	*
4	18	*
5	17	*
11	11	*

⋮

Ex:- 7 5 2 5 9 5 r = 10

Hm { 7:1 9:1
 5:3 2:1 }

a	b	
7	3	*
5	5	✓

Pseudo Code :-

1) Create HashMap and insert arr[] (0 to n)

$T.C \rightarrow O(N^2)$

```

for (i = 0; i < N; i++) {
    a = arr[i]
    b = K - a
    if (a != b) {
        if (b is in map) {
            return True
        }
    }
    else { // a == b
        if (hm[b] > 1) {
            return True
        }
    }
}
return false

```

(hm contains key (b))

$T.C \rightarrow O(N^2)$

$S.C \rightarrow O(N)$

* Approach :- using hashset

	0	1	2	3	4	5	6	7	8	9
arr[] :	8	9	1	-2	4	5	11	-6	7	5

$K = 22$

by 

a	b	hs
8	14	{ 3 }
9	13	{ 8, 3 }
1	21	{ 8, 9, 3 }
-2	24	{ 8, 9, 1, 3 }
4	18	{ 8, 9, 1, -2, 3 }
5	17	{ 8, 9, 1, -2, 4, 3 }
11	11	{ 8, 9, 1, -2, 4, 5, 3 }
.	.	{ 8, 9, 1, -2, 4, 5, 11, 3 }

Pseudo code :-

```

hashset <int> hs
for (i=0; i<n; i++) {
    a = arr[i]
    b = k - a
    if (b is in hs) {
        return True
    }
    else {
        hs.insert(a)
    }
}
return False

```

T.C $\rightarrow O(n)$
S.C $\rightarrow O(n)$

Ques) Calculate no. of i, j s.t. that,
 $arr[i] + arr[j] = \underline{k}$.

5 5 5 5

$$i! = j$$

$$k = 10.$$

→ Trick :-

Ques) Check if there exists a pair
s.t. that, $arr[i] - arr[j] = \underline{k}$ $i! = j$.

2 3 5 7 9 11, $k = \underline{2}$.

$$a, b = \underline{k + a}.$$

Trick.

10:02 → 10:07

Ques) Given n array elements, calculate no. of distinct elements in every subarray of size k .

$c = 4$

$n = 10, k = 4 \quad [6 \quad 9]$

arr[]: 2 4 3 8 3 9 4 9 4 10

0	3	4
1	4	3
2	5	3
3	6	4
4	7	3
5	8	2
6	9	3

1 deg.

for every window, get no.
of distinct using hashset,

last Subway $[v-k, v-1]$

$$\underline{b-a+1}$$

$$\cancel{N} - \cancel{x} - \cancel{N} + x + \cancel{x} = \underline{x}$$

for ($i=0$; $i \leq n-k$; $i++$) { $[0, \frac{n-k}{b}]$

HashSet < int > hs;

```
for (j = i; j < (i + r); j++) {
```

to insert $\text{arr}[j]$)

3

2

$$T.C \rightarrow O(n - k + 1) + k$$

$$S.C \rightarrow O(1)$$

arr : ⁰2 ¹4 ²3 ³8 ⁴3 ⁵9 ⁶4 ⁷9 ⁸4 ⁹10

└──────────┘

$$k = 4$$

Initially : - remove add by Ans

[0 - 3] { ~~2~~, 4, 3, 8 } 4

[1 - 4] 0 4 { ~~4~~, 3, 8 } 3

[2 - 5] 1 5 { ~~3~~, 8, 9 } 3

[3 - 6] 2 6 { 8, 9, 4 } 3



Ans → 4

issue:- If we remove an element, indirectly,
all occurrence of same element get
removed.

11 HashMap with Sliding window :-

arr [] : 0 1 2 3 4 5 6 7 8 9
 2 4 3 8 3 9 4 9 4 10

Subarray :-	remove	add	HashMap	
[0 - 3]			{ 2:1, 4:1, 8:1, 3:1 }	4
[1 - 4]	0	4	{ 2:0 , 4:1, 8:1, 3:2 }	3
[2 - 5]	1	5	{ 4:0 , 8:1, 3:2, 9:1 }	3
[3 - 6]	2	6	{ 8:1, 3: 1 , 9:1, 4:1 }	4
[4 - 7]	3	7	{ 8:0 , 3:1, 9:2, 4:1 }	3
[5 - 8]	4	8	{ 3:0 , 9:2, 4:2 }	2
[6 - 9]	5	9	{ 9:1, 4:2, 10:1 }	3

✓

In hashMap, for our above

problem, if freq == 0, remove from map

HashMap < Integer, Integer > hm =

new HashMap < > ();



H Pseudo code:-

```
HashMap<int, int> hm = new HashMap<>();
```

```
for (i=0; i<k; i++) {
```

```
    if (!hm.containsKey(arr[i])) {
```

```
        hm.put(arr[i], hm.get(arr[i]) + 1);
```

```
    }
```

```
    else {
```

```
        hm.put(arr[i], 1);
```

```
    }
```

```
}
```

k=5

0 1 2 3 4 5 6

```
Print (hm.size());
```

```
for (i=1, j=k; i<=(n-k); i++, j++)
```

```
    hm.put(arr[i-1], hm.get(arr[i-1]) - 1)
```

```
    if (hm.get(arr[i-1]) == 0) {
```

```
        hm.remove(arr[i-1])
```

```
    }
```

```
    if (hm.containsKey(arr[j])) {
```

```
        hm.put(arr[j], hm.get(arr[j]) + 1)
```

```
    } else {
```

```
        hm.put(arr[j], 1)
```

```
    }
```

```

    |
    |_ print (hm.size());
    3

```

T.C $\rightarrow O(N)$

S.C $\rightarrow O(\underline{K})$.

Doubt

HashMap <int, int> hm;

int freq = hm.get(arr[i])

for (i=0; i<n; i++) {

hm.update(arr[i], freq+1) if (arr[i] is in hm) {

hm[arr[i]]++

hm.put(arr[i], hm.get(arr[i])+1)

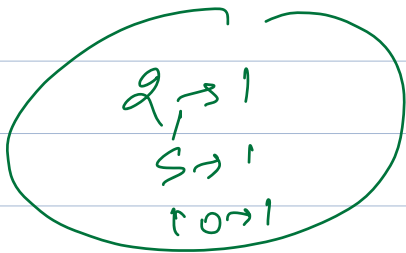
~~Kreeti Kreeti~~

K=3

0	1	2	3	4	5
2	5	10	15	20	25

i → 3
 ↑
 5

i + r



2, 2, 2, 3, 3, 5, 6

\rightarrow

~~2 \rightarrow 2 2 1~~
~~3 \rightarrow 2 1 0~~
~~5 \rightarrow 1 0~~ \rightarrow x
6 \rightarrow 1

2, 2, 3, 3, 3, 4, 5, 6, 5

\downarrow

2 2 3 3 5 2, 2, 3, 3, 5, 6

