

Today's Content :-

- Modular operators
- Modular arithmetic
- 1 Hard problem.

int range → $(-2 * 10^9 \text{ to } 2 * 10^9)$

long rang → $(-8 * 10^8 \text{ to } 8 * 10^8)$

∴ → modules / remainder

Dividend = divisor * quotient + remainder

$$10 \div 4 = 2 \Rightarrow 4 * \left(\frac{10}{4}\right) + r \Rightarrow 8 + r = 10, \underline{r=2}$$

$$13 \div 5 = 2 \Rightarrow 5 * \left(\frac{13}{5}\right) + r \Rightarrow 10 + r = 13, \underline{r=3}$$

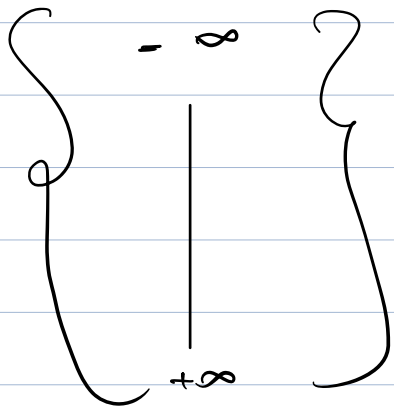
$$100 \div 7 \Rightarrow 7 * \left(\frac{100}{7}\right) + r \Rightarrow 100 \Rightarrow 98 + r = 100, \underline{r=2}$$

$$150 \div 7 \Rightarrow 7 * \left(\frac{150}{7}\right) + r \Rightarrow 147 + r = 150, \underline{r=3}$$

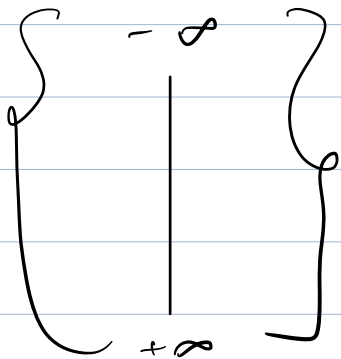
$$-60 \div 9 \Rightarrow 9 * \left(\frac{-60}{9}\right) + r = -60,$$

$$-54 + r = -60$$

$$\underline{r = -6}$$



$$\text{1..10} \rightarrow \begin{matrix} \text{min} & & \text{max} \\ \{ 0, & 9 \} \end{matrix}$$



$$\rightarrow \text{1..M} = \begin{matrix} \text{min} & & \text{max} \\ \{ 0 & M-1 \} \end{matrix}$$

\downarrow
 limit range { consistent hashing
 hashmap/dict
 cryptography

Conceptually

$$\text{Remainder} = \text{Dividend} - (\text{divisor} \times \text{Quotient})$$

greatest multiple of
divisor \leq dividend.

$$10 \div 4 = 10 - 8 \Rightarrow \underline{2}$$

$$13 \div 5 = 13 - 10 \Rightarrow \underline{3}$$

$$100 \div 7 = 100 - 98 \Rightarrow \underline{2}$$

$$150 \div 7 \Rightarrow 150 - 147 \Rightarrow \underline{3}$$

$$-60 \div 9 \Rightarrow -60 - (\text{greatest multiple of } 9 \leq -60)$$

$$-60 - (-63) \Rightarrow -60 + 63 \Rightarrow \underline{3}$$

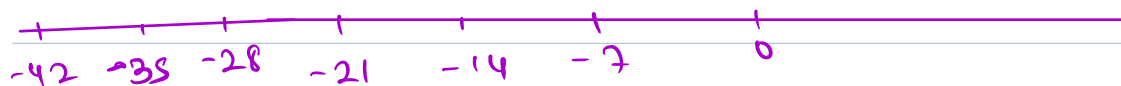


$$-63 \quad -54$$



1

$$-40 \div 7 \Rightarrow -40 - \{ \text{greatest multiple of } 7 \leq -40 \}$$



$$-40 - \{-42\} \Rightarrow -40 + 42$$

$$\Rightarrow 42 - 40 \Rightarrow \underline{2}$$

As per your language :-

$$-80 \% 9 \begin{cases} \rightarrow \text{language} \rightarrow -8 \\ \rightarrow \text{concept} \rightarrow -80 - \{ \text{greatest multiple of } 9 \leq -80 \} \end{cases}$$

$$\Rightarrow -80 - (-81) \Rightarrow \underline{1}$$

$$-40 \% 9 \begin{cases} \rightarrow \text{language} \rightarrow -4 \\ \rightarrow \text{concept} \rightarrow -40 - \{ \text{greatest multiple of } 9 \leq -40 \} \end{cases}$$

$$-40 - (-45) \Rightarrow \underline{5}$$

$$-60 \% 9 \begin{cases} \rightarrow \text{language} \rightarrow -6 \\ \rightarrow \text{concept} \rightarrow -60 - \{ \text{greatest multiple of } 9 \leq -60 \} \end{cases}$$

$$-60 - (-63) \Rightarrow \underline{3}$$

if ($x < 0$) {
 $\{ x \% p + p \}$ \rightarrow only by adding p we can get expected no.
}

}

modular arithmetic

$$(a + b) \% m = (\underbrace{a \% m}_{(0, m-1)} + \underbrace{b \% m}_{(0, m-1)}) \% m$$

$$\begin{aligned} a = 6, b = 13, m = 7 &\Rightarrow (6 \% 7 + 13 \% 7) \% 7 \\ &\Rightarrow (6 + 6) \% 7 \\ &\Rightarrow 12 \% 7 \\ &\Rightarrow \underline{5} \\ (19 \% 7) &\Rightarrow \underline{5} \end{aligned}$$

$$\begin{aligned} a = 4, b = 5, m = 6 &\Rightarrow (4 \% 6 + 5 \% 6) \% 6 \\ &\Rightarrow (4 + 5) \% 6 \Rightarrow \underline{3} \\ (9 \% 6) &= 3 \end{aligned}$$

$$(a * b) \% m \Rightarrow \{ \underbrace{a \% m}_{(0, m-1)} * \underbrace{b \% m}_{(0, m-1)} \} \% m$$

$$\begin{aligned} a = 6, b = 7, m = 4 &\Rightarrow (6 \% 4 * 7 \% 4) \% 4 \\ &\Rightarrow (2 * 3) \% 4 \\ &\Rightarrow \underline{2} \\ 42 \% 4 &= 2 \end{aligned}$$

$(a-b) \% m$
 $(a/b) \% m$ } advance batch.

Problems :-

given $a, m, p \rightarrow$ calculate $a^m \% p$.

$$\frac{a}{3} \quad \frac{m}{4} \quad \frac{p}{2} \Rightarrow (3^4) \% 2 = 81 \% 2 = 1$$

given $a, m=5,$
 $a^5 \% p$

power (int a, int m, int p)

// no overflow

for (int i=1; i<=m; i++) {

$a = a * a$

return $a \% p$

wrong

i	// a	a
1	$a = a * a$	a^2
2	$a = a * a$	a^4
3	$a = a * a$	a^8

```
power (int a, int n, int p) { // a^n % p
```

// assume no overflows.

```
int ans = 1
```

```
for (i = 1; i <= n; i++) {
```

```
    |   ans = a * a
```

```
    }
```

```
return a % p
```

3

→ wrong

```
power (int a, int n, int p) { // a^n % p
```

```
long ans = 1
```

```
for (i = 1; i <= n; i++) {
```

```
    |   ans = (ans * a) % p
```

```
    }
```

$ans = (ans \% p * a \% p) \% p$

```
return ans % p
```

3

Constraints :-

$1 \leq n \leq 10^5$

$1 \leq a \leq 10^9$

$1 \leq p \leq 10^9$

$a = 10, n = 14, p = 25$

$(10^{14}) \% 25$

i	ans
1	10

$$1.C \rightarrow 000$$

$$8.C \rightarrow 0011$$

$$2 \quad 0$$

$$3$$

Break \rightarrow 10:10pm to 10:20pm.

Divisibility of 3 :- Sum of digits are divisible by 3.

$$= (789) \div 3 \quad (8215) \div 3 \quad (3458) \div 3$$

$\checkmark \quad \quad \quad \times \quad \quad \quad \times$

$$10 \div 3 = 1$$

$$10^2 \div 3 = 1$$

$$10^3 \div 3 = 1$$

$$10^x \div 3 = 1$$

$$(3458)$$

$$= 3 \times 10^3 + 4 \times 10^2 + 5 \times 10^1 + 8 \times 10^0$$

$$(3 \times 10^3 + 4 \times 10^2 + 5 \times 10^1 + 8 \times 10^0) \div 3$$

$$\left(\underbrace{(3 \times 10^3)}_1 \div 3 + \underbrace{(4 \times 10^2)}_1 \div 3 + \underbrace{(5 \times 10^1)}_1 \div 3 + \underbrace{(8 \times 10^0)}_1 \div 3 \right) \div 3$$

$$(3 \div 3 + 4 \div 3 + 5 \div 3 + 8 \div 3) \div 3$$

$$\rightarrow (3 + 4 + 5 + 8) \div 3$$

$$\rightarrow (20) \div 3 = \underline{2}$$

* Divisibility of 4 :- (last 2 digits $\div 4$)

$$\{ a_3 a_2 a_1 a_0 \} \div 4 \rightarrow 5432$$

\downarrow

$$5 \times 1000 + 4 \times 100 + 3 \times 10 + 2$$

$$\{ a_3 \times 10^3 + a_2 \times 10^2 + a_1 a_0 \} \div 4$$

$$((a_3 \times 10^3) \div 4) + ((a_2 \times 10^2) \div 4) + (a_1 a_0 \div 4) \div 4$$

$\downarrow \quad \quad \downarrow \quad \quad \swarrow$

Last 2 digits are divisible by 4
or not.

Divisible by 9 \Rightarrow { sum of digits $\div 9$ }

\downarrow

Todo .

Ques) Given n numbers in $arr[]$, calculate number $\% p$.

Each $arr[i]$ contains single digit of a no.

$N = 5$
 $arr[5] = 7 \ 8 \ 9 \ 6 \ 2$

$$0 \leq arr[i] \leq 9$$

Constraints :-

$$1 \leq N \leq 10^5$$

$$1 \leq p \leq 10^9$$

$$N = 3, \quad 999 : 10^3 - 1$$

$$N = 4, \quad 9999 : 10^4 - 1$$

$$N = 10, \quad 10^{10} - 1$$

$$N = 20, \quad 10^{20} - 1$$

$$N = 10^5, \quad 10^{10^5} - 1$$

$arr[5] =$

0	1	2	3	4
3	2	6	4	9

Diagram showing the conversion of the array to a number:

- 3 $\rightarrow 3 \times 10^4$
- 2 $\rightarrow 2 \times 10^3$
- 6 $\rightarrow 6 \times 10^2$
- 4 $\rightarrow 4 \times 10^1$
- 9 $\rightarrow 9 \times 10^0$

$$(3 \times 10^4 + 2 \times 10^3 + 6 \times 10^2 + 4 \times 10^1 + 9 \times 10^0) \% p.$$

$$(3 \times 10^4 + 2 \times 10^3 + 6 \times 10^2 + 4 \times 10^1 + 9 \times 10^0) \cdot 10^1$$

$$\left((3 \times 10^4) \cdot 10^1 + (2 \times 10^3) \cdot 10^1 + (6 \times 10^2) \cdot 10^1 + (4 \times 10^1) \cdot 10^1 + (9 \times 10^0) \cdot 10^1 \right) \cdot 10^1$$

$$(3 \cdot 10^1 \times 10^4 \cdot 10^1) \cdot 10^1$$

+

$$(2 \cdot 10^1 \times 10^3 \cdot 10^1) \cdot 10^1$$

+

$$(6 \cdot 10^1 \times 10^2 \cdot 10^1) \cdot 10^1$$

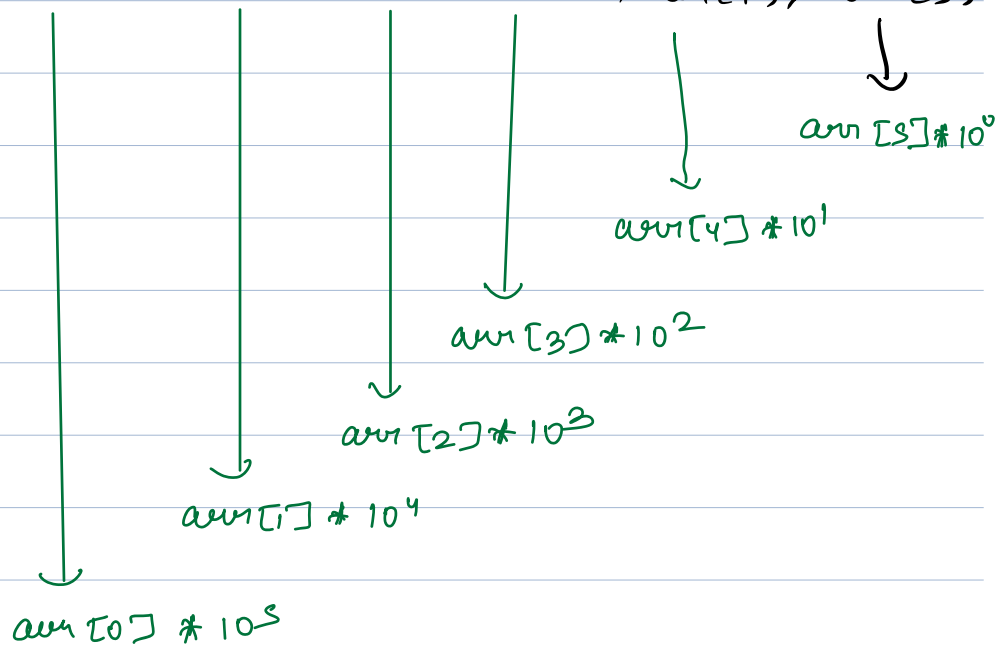
+

$$(4 \cdot 10^1 \times 10^1 \cdot 10^1) \cdot 10^1$$

+

$$(9 \cdot 10^1 \times 10^0 \cdot 10^1) \cdot 10^1$$

// arr[6] = arr[0], arr[1], arr[2], arr[3], arr[4], arr[5]



$$(ans[0] * 10^5 + ans[1] * 10^4 + ans[2] * 10^3 + ans[3] * 10^2 + ans[4] * 10^1 + ans[5] * 10^0) \% p$$

$$\downarrow$$

$$(ans[0] \% p * 10^5 \% p) \% p$$

$$\downarrow$$

$$(ans[1] \% p * 10^4 \% p) \% p$$

$$\downarrow$$

$$(ans[2] \% p * 10^3 \% p) \% p$$

$$\downarrow$$

$$10^2 \% p$$

$$\downarrow$$

$$10 \% p$$

$$\downarrow$$

$$(ans[5] \% p * 1 \% p) \% p$$

Same 10 :-

$$(p = 3)$$

$$ans = 1$$

$$ans = (ans * 10) \% p = 10 \% p \Rightarrow 1$$

$$ans = (ans * 10) \% p = (10^2) \% p \Rightarrow 1$$

$$ans = (ans * 10) \% p = (10^3) \% p$$

Pseudo Code :-

// given arr[n] & p

long fans = 0

long ans = 1

for (i = N-1; i >= 0; i--)

 fans = fans + { (arr[i] % p) * (ans % p) } % p

 fans = fans % p.

 ans = (ans * 10) % p

}

return fans;

$10 \% p \rightarrow 10^2 \% p \rightarrow 10^3 \% p \rightarrow 10^4 \% p$

↓
x

↓
(10 * 10) % p

↓
(10 * 10 + 10 * 10) % p

(x * 10) % p