

Ques) Given N array elements, find no. of distinct elements.

$$arr[7] = \{ \underline{6}, \underline{3}, \underline{2}, \underline{3}, \underline{8}, \underline{6}, \underline{9} \} \rightarrow 5.$$

$$arr[8] = \{ \underline{2}, \underline{9}, \underline{8}, \underline{7}, \underline{6}, \underline{8}, \underline{2}, \underline{7} \} \rightarrow 5.$$

 * *

<2,1>
<9,1>
<8,2>
<7,2>
<6,1>

→ hashmap size.

T.C → O(n)

2
9
8
7
6

T.C → O(n)

→ set.size.

```
Hashset<integer> hs;
```

```
for (i=0; i<m; i++) {
```

```
    hs.insert(arr[i]);
```

```
}
```

```
cout << hs.size();
```

(If we store same keys again in hashset, it will only store a single occurrence).

Ques) :- Given N array elements, check if there exists a subarray with $\text{sum} = 0$.

arr :- 0 1 2 3 4 5 6 7 8 9
 2 2 1 -3 4 3 1 -2 -3 2

PF :- 2 4 5 2 6 9 10 8 5 7.

5

$$PF[2] = \text{sum}[0-2] = 5$$

5

$$PF[8] = \text{sum}[0-8] = 5.$$

These elements sum is 0.

$$PF[8] = \text{sum}[0-2] + \text{sum}[3-8]. \rightarrow PF[8] = PF[2] + \text{sum}[3-8].$$

Brute force:-

- 1) Generate all subarray

Time \downarrow $O(n^3)$ \rightarrow carry forward \downarrow $O(n^2)$

$$\text{Sum}[3-8] = \text{PF}[3-8] - \text{PF}[2]$$

$$\text{Sum}[3-8] = 0,$$

Obs:- ① If your $\text{PF}()$, if an element repeats, subarray sum = 0.

② Check if 0 is in $\text{PF}[]$.

	0	1	2	3
En:-	3	-1	-2	4
PF:-	3	2	0	4

Pseudocode:-

Steps:-

- 1) calculate $\text{PF}[]$.
- 2) check if 0 is present in $\text{PF}[]$.
- 3) check if element repeats in $\text{PF}[]$.

→ do it with prev ques.



In the last ques and check:-

($\text{iset.size}() == n$) or not.

Ques) find the length of longest subarray with sum = 0.

	0	1	2	3	4	5	6	7	8	9	10	11	12
arr [] =	3	3	4	-5	-2	2	1	-3	3	-1	5	-4	-1
pf [] =	3	6	10	5	3	5	6	3	6	5	10	6	5

Obs:- for every element in pf(), store its 1st and last occurrence.

2 hashmaps :-



1 hashmap



for every element
will store
first occurrence

2 hashmap



for every element
will store
last occurrence.

$$PF[7] = \frac{0}{3} \frac{1}{6} \frac{2}{10} \frac{3}{5} \frac{4}{3} \frac{5}{5} \frac{6}{6} \frac{7}{3} \frac{8}{6} \frac{9}{5} \frac{10}{10} \frac{11}{6} \frac{12}{5}$$

	firstIndex	len	ans
<3, 0>	0	4	4
<6, 1>	3	2	4
<10, 2>	1	6	5
<5, 3>	0	7	7
	1	7	7
	3	6	7
	2	8	8
	1	10	10
	3	9	10

firstIndex	len	ans
0	4	4
3	2	4
1	6	5
0	7	7
1	7	7
3	6	7
2	8	8
1	10	10
3	9	10

hm <int, int>,

hm.insert <0, -1> → Edge cons.

for (i=0; i<n; i++) {

if (PF[i] is in hm) {
int len = i - hm.get(PF[i]);

ans = max(ans, len);

3

else {

hm.insert (PF[i], i);

3

Edge Case :-

	0	1	2	3	4	5	
arr \rightarrow	4	-3	-1	2	3	-5	
pf \rightarrow	4	1	0	2	5	0	

$\langle 0, -1 \rangle$

Ques) Given array only contains 0's and 1's,
find max length subarray
which contains equal 0's and 1's.

Soln:- replace ^{all} 0 \rightarrow with -1.

\rightarrow 0 0 1 1 0 0 1 1 1



-1 -1 1 1 -1 -1 1 1 1

Ques) Count the Subarrays with sum = 0.



$$5 \rightarrow 4 \rightarrow {}^4C_2 \rightarrow 6$$

$$6 \rightarrow 3 \rightarrow {}^4C_2 \rightarrow 6 \dots$$

* Count no. of times a pf sum [] is repeating.

Pf sum freq.

$$3 \rightarrow 3 \rightarrow {}^3C_2 +$$

$$6 \rightarrow 4 \rightarrow {}^4C_2 + \Rightarrow 16$$

$$10 \rightarrow 2 \rightarrow {}^2C_2 +$$

$$5 \rightarrow 4 \rightarrow {}^4C_2 +$$

$$9 \rightarrow 1 \rightarrow {}^1C_2 \rightarrow 1 \left(\frac{1-1}{2}\right) = 0.$$

$$mC_2 \Rightarrow \frac{m(m-1)}{2}.$$

↳

$$\frac{m}{2! (m-2)!} \Rightarrow \frac{m(m-1)(m-2)!}{2 \times (m-2)!}$$

~~# Pseudo code:~~

① `HashMap<int, int> fmap;`

② Create a pf array.

③ Store freq of pf array in fmap.

add <0, 17>

int ans = 0;

for each root

for (int val : fmap.keySet()) {

int freq = fmap.get(val);

ans = ans + (freq * $\frac{freq-1}{2}$)

3

return ans;

Edge Care:

arr1 →	4	-3	-1	2	3	-5
pf →	4	1	0	2	5	0

0 → 2 → ${}^a C_2 \rightarrow 1$.

0 → 3 → ${}^3 C_2 \rightarrow 3$.

$\langle 0, 1 \rangle$

add it to start.

There is a sum 0 with

freq 1.

④ Arrays and Maths

* Majority Element ($\frac{n}{3}$) → freq should be greater than $\frac{n}{3}$.

$$(10, 10, 20, 30, 10, 10) \rightarrow \frac{6}{3} \approx 2 \rightarrow \underline{\underline{3}}$$

→ How many majority elements an array can have?



At most two majority elements can be there.

* Here, if you remove 3 distinct elements majority will not change.

10, 10, 20, 30, 10, 10, 80

$$\text{may 1} = 10 \quad \text{may 2} = \cancel{20}$$

$$\text{freq 1} = 8 \times 2 \times 2 \times 3 \quad \text{freq 2} = 8 \times 0$$

Pseudo Code :-

int may1 = -∞
for (int i=0; i<n; i++) { int val = arr[i]; }
int may2 = ∞
for each loop,
for (int val: arr) { // val: arr[i]; }

if (may1 == val) {
freq1++;

else if (may2 == val) {
freq2++;

else if (freq1 == 0) {
freq1 = 1;
may1 = val;

else if (freq2 == 0) {
freq2 = 1;
may2 = val;

else {
freq1--;
freq2--;

}

3
}

$$N = \frac{12}{3} \rightarrow \min \text{ freq} \Rightarrow \frac{12^4}{3} \rightarrow \underline{\underline{S}}$$

4 S 6 6 S 5 6 S 6 4 S 6

$$\text{may}_1 = 4^6$$

$$\text{may}_2 = \cancel{6}^5$$

$$\text{freq}_1 = 0 \times 0 \times 2 \times 2 \times 2 \quad \text{freq}_2 = 0 \times 0 \times 2 \times 1 \times 2$$

$\frac{N}{2}$ majority

Once you count that,

$$\frac{N}{2} \rightarrow 1 \text{ Max} \rightarrow 2$$

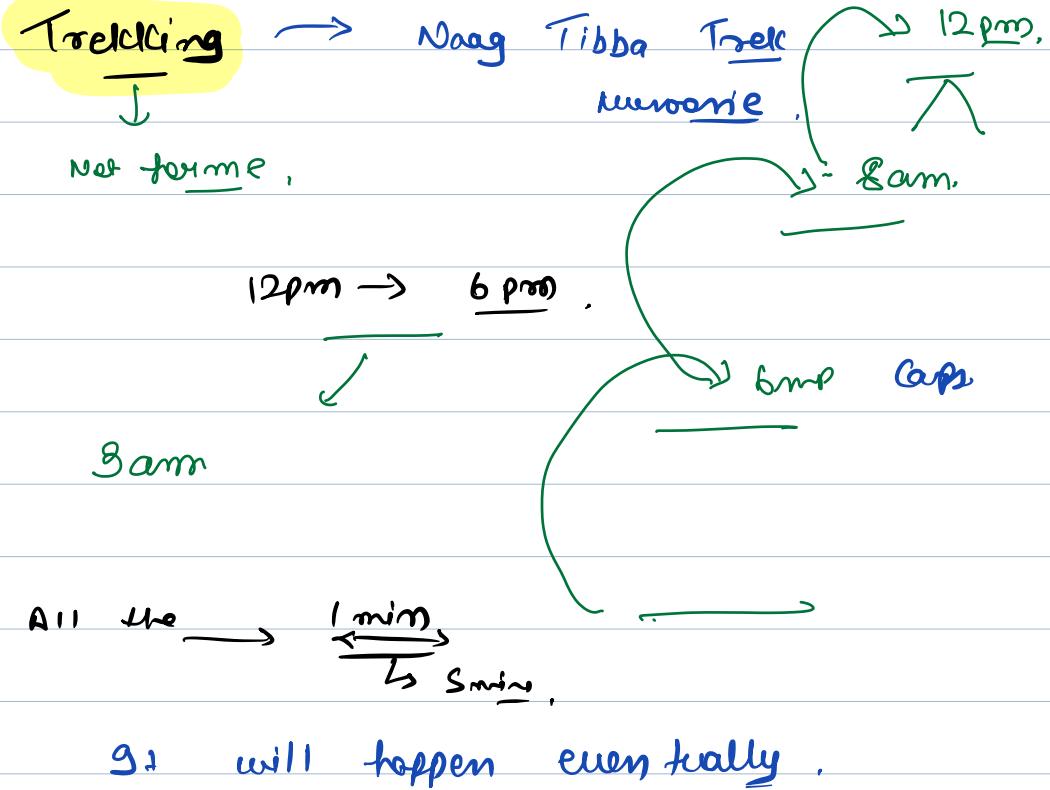
$$\frac{N}{3} \rightarrow 2 \rightarrow 3$$

$$\frac{N}{4} \rightarrow 3 \rightarrow 4$$

$$\frac{N}{k} \rightarrow k-1 \rightarrow k$$

2 : 04 - 2 : 14 pm.

Top
T



Trelding → Once.

← Comparators →

Ques) Given n array elements, sort them in Increasing order of their no. of factors.

If two elements have same factors, ele with less value should come first.

# factors:	9	3	4	8	16	37	6	13	15
	3	2	3	4	5	2	4	2	9

→ 8, 18, 32, 4, 9, 6, 8, 15, 16.

Given:-
→ { 1, 21, 6, 23, 10, 14, 25 }
factors
1 4 4 2 4 4 3

→ { 1, 23, 25, 6, 10, 14, 21, }

→ Arrays.sort (array) *

int array[] = { 1, 2, 3, 4, 5, -13 };

Arrays.sort (array, new Comparator<Integer>){

// for descending .

@Override

public int compare (Integer a, Integer b) {

if (a > b) {

// you want a to come first
return -1

3 else {

// you want b to come first
return 1

3

3);

Arrays.sort (array, new Comparator<Integer>){

@Override

public int compare(Integer a, Integer b)

if (a should come before b
return -1

if a = b return 0
else return 1.

3

3);

Sorting on factors

Arrays.sort (arr, new Comparator<Integer>){

@Override

public int compare(Integer a, Integer b)

int f1 = factors(a)

int f2 = factors(b)

if (f1 < f2) {

return -1

3 else if (f1 == f2 and a < b) {

return -1

3 else {

return 1

3

3

3),

```
Arrays.sort(arr, new Comparator<Integer>(){  
    @Override  
    public int compare(Integer o1, Integer o2) {  
        return o1.compareTo(o2); → write your logic  
    }  
});  
  
for(int str:list){  
    System.out.println(str);  
}
```

→ above logic will not work on,
int[] arr,
Integer[] arr.

```

Integer[] arr = {1,2,3,4,5};
ArrayList<Integer> list = new ArrayList<>();
list.add(1);
list.add(2);
list.add(3);
list.add(4);

Collections.sort(list,new Comparator<Integer>(){

    @Override
    public int compare(Integer a, Integer b) {
        if(a<b){
            return -1;
        }else{
            return 1;
        }
    }
});

Arrays.sort(arr,new Comparator<Integer>(){

    @Override
    public int compare(Integer o1, Integer o2) {
        return o1.compareTo(o2);
    }
});

```

Ques) :- Given N array ele, sort in increasing order of no. of digits, if two ele have same no. of digits, ele with more value should come first.

98, 2, 37, 639, 8, 100, 345, 49
 ↓ ↓ ↓ ↓ ↓ ↓ ↓
 2 1 2 3 1 3 3 2

$\rightarrow 8 \ 2 \ 93 \ 49 \ 37 \ 639 \ 845 \ 100.$

Arrays.sort (arr, new Comparator<Integer>){

@Override

public int compare (Integer a, Integer b){

int d1 = digit (a)

int d2 = digit (b)

if (d1 < d2) {

return -1;

3

if (d1 == d2 and a > b) {

return -1

3 else {

3 return 1

3};

Tuesday \rightarrow Last class.

→ Pending problems

Problem Solving Session



↳ Timing,

7 days gap you'll get.

Advance → ?

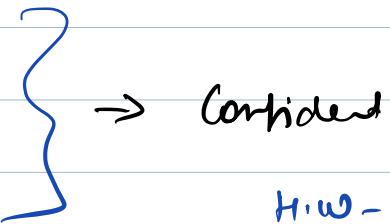
Arrays

Bits

recursion

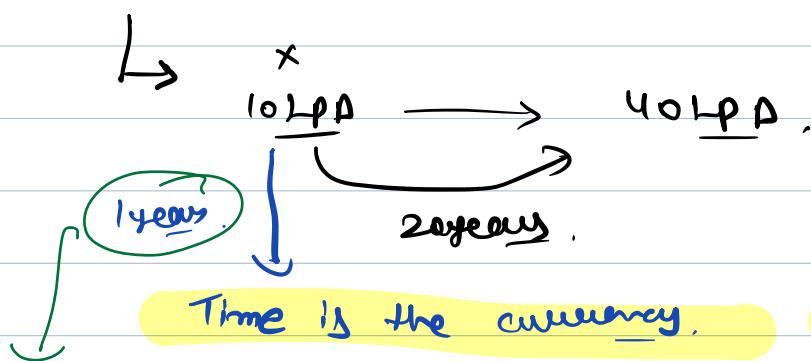
HashMap

Strings



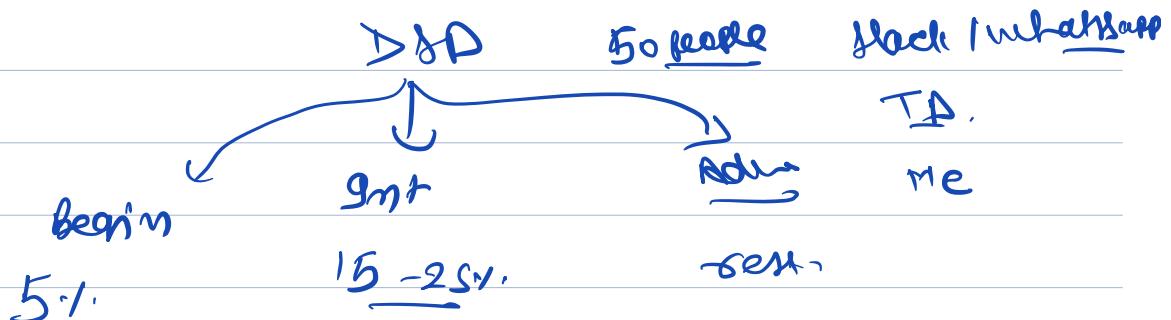
H.W - 80-60%.

7 days break



1 year

$\alpha \rightarrow \beta \rightarrow \gamma \rightarrow \delta$



* Rest in the end met in the middle

Advanced → 4-5 months

kenmai.aman@scaleer.com.