

Z Specification:

Z specification is formal language used for modelling and describing system. It is based on mathematical notations, set theory, and lambda calculus

Formal specifications use mathematical notation to describe in a precise way the properties which an information system must have, without unduly constraining the way in which these properties are achieved. It describes what the system must do rather than how?

This abstraction makes formal specifications useful in the process of developing a computer system. It is independent of program code. A formal specification of a system can be completed early in its development. The other main ingredient in Z is a way of decomposing a specification into small pieces called schemas. By splitting the specification into schemas, we can present it piece by piece. Each piece can be linked with a commentary which explains informally the significance of the formal mathematics. In Z, schemas are used to describe both static and dynamic aspects of a system. The static aspects include.

- state it can occupy
- invariant relationship that maintains system moves from state to state.

Dynamic aspects include.

- operations that are possible
- relationship between input and output
- the changes of state that happen

This model consists of description of state space and system operations. System state space is a set of all states that the system could be in. It describes the value of each variable.

We represent state space of a system as collection of functions, sets, relations, sequences, bags etc.

Together with collection of invariant properties on these objects.

It describes regularities between state changes.

Operations are generally defined in terms of pre and post conditions.

Operations define acceptable state transitions.

Z schemas.

It is a 2 dimensional graphical notation for describing.

1) state space and 2) operations.

Schema is either of the form

Schema name
Declararations
Perdicates

or

Schema name
Declarations

In the later case predicate part is assumed to be true

after this, schema name will be associated with schema which is content of box

declaration part of schema contain
list of variable declaration.

And reference to other schemas

variable declaration has usual form

$x_1, x_2, x_3, \dots, x_n.T$

predicate part of schema contaion list of predicates.separated by
semicolens.

Use case formalization.

set of all unique identifiers throughout the specification is denoted by
given set of classifier

there are several relationships between use cases. This include relationship,
and generalization.

Include is a relationship between two use cases which is used to show
behavior of the included use case is inserted into behavior of inluding use
cases.

Exteend relationship is that specifies how and when behavior defined in
usually operations extending use case can be inserted into behavior defined
in the extended use case

The relationships include and extend are basically represented as a Z
relation that relates two use cases. The relation include is used to record
which use case is directly included by the original use case, whereas the
relation extend is used to record which use case is directly extending the
original use case.

The functions `includedUsecases` and `extendingUsecases` are respectively obtained with the transitive closure of `include` and `extend`. They are used to formally specify the function that returns the set of use cases that are directly or indirectly related with the relations `include` or `extend`.

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Actor formalization.

An actor specifies a role played by a user or any other system that interacts with the subject

given set `CLASSIFIER` is used to define the set of all unique identifiers throughout the specification. The variable `ACTOR` is introduced as a classifier to specify the set of actors' identifiers.

The actor can be human user, some internal application or may be some external application. The actor type is introduced as an enumerated set representing two types of actors.

The relationship between actor and use case is basically represented as a Z relation that relates an actor to a use case. The relation `association`. `AU` is used to record which use case can be used by the actor, whereas the function `relatedUsecases` returns the set of use cases that are associated to the actor.

Actor generalization is used to factor out and reuse similarities between actors. This relationship shows that one actor inherits the role and properties of another actor. To represent generalization, we chose to add a state variable to declaration part of the schema `Actor` as shown below. This state variable is called `parents` and specified as a set of actors. The generalization relationship also implies that the descendant actor can use all the use cases that have been defined for its ancestor.

The specification of requirements is a key activity for achieving the goals of any software project and it has long been established and recognized by researchers and practitioners

The objectives of our work are providing concise and unambiguous specification of use case diagrams using Z notation and consistency checking by means of Z theorems

Formal specification of use case diagrams has been done in the past using different languages but we believe that the expressiveness and the mathematical basis of Z notation provide a richer framework for such formalism. fact that the use of formal specification can greatly enhance the quality of the produced software as well as to contribute significantly to cost savings in the software development process.

These are the connectives of propositional logic. they allow complex predicates to be built from simpler ones in such a way that the truth or falsity of the compound in some binding depends only on the truth or falsity of the parts in the same binding

For example, the predicate $P1 \wedge P2$ is true in any binding if and only if both P1 and P2 are true in that binding.

Operations schemas

In specifying a system operation, we must consider the objects that are accessed by the operation, and of these.

the objects that are known to remain unchanged by the operation (value parameters)

the objects that may be altered by the operation

the pre-conditions of the operation, i.e., the things that must be true for the operation to succeed.

the post-conditions the things that will be true after the operation, if the pre-condition was satisfied before the operation

Example telephone book.

Consider lookup operation

we put a name in, and get a phone number out

- this operation accesses the PhoneBook schema.

it does not change it.

- it takes a single 'input' a name for which we want to find a phone number.

it produces a single output a phone number

it has the pre-condition that the name is known to the database

Procedural Abstraction: Schemas only describe what is to be done; issues relating to "how" are ignored