

② Symbolic Constants.  
Preprocessor Directives / macros.

#define — —

`#define` — —

$$pi \rightarrow 2.14 + 1$$

cont `<< pi << y.`

$$<< 2.14 + 1 \times y \times y \\ y \times y = b + 2.14$$

$$\underline{18.14}$$

### operators

Unary:  $\rightarrow$  It requires only 1 operand to fulfill its task.

int  $a = +10;$ , int  $b = -10;$   
 $++a;$      $a++;$      $-b;$      $b--;$   
size} (1arity)

Arithmetic  $\frac{a}{b} \rightarrow$  requires 2 operands  
and only works when  
both sides are numbers / easily  
convertible to numbers  
 $5 + 9,$      $13.0 + 6.0 \checkmark$

$65 \quad 'A' + 1 \Rightarrow 66.$   
implicit type conversion  
to integer value.

$+, -, *, /, ^.$

$\sim \rightarrow$  condition check, ...

( ) → condition check,  
to give priority to the  
expression inside brackets.

Assignment ( $=$ ) → to assign values

Relational → to compare / comparison.

$'A' \geq 1$   
 $\neg A \neq 1$  value. (same rules  
as arithmetic)

it has a return type  
it returns either true/false.

logical operators. ∵ conditions to check  
true & vice versa.

8l → AND.  
= → If all true/1 true  
1, else false.

$v_1$	$v_2$	Result
0/F	0/F	0/F
1/T	0/F	0/F
0/F	1/T	0/F
1/T	1/T	1/T

If AND encounters any false  
value, it will move out of that  
condition instantly without further  
evaluating rest of the expressions  
conditions.

11 → logical OR. → Any 1 → 1.

$v_1$	$v_2$	Result
0	0	0
0	1	1
1	0	1
1	1	1

0      0      |  
|      |      |  
|

! → value      result  
v1  
0      |  
|      0.  
|

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true in C++ / C

∴ / any non zero value is  
always considered true.

if (true) → execute

if (!false) → execute

if (0) → not execute

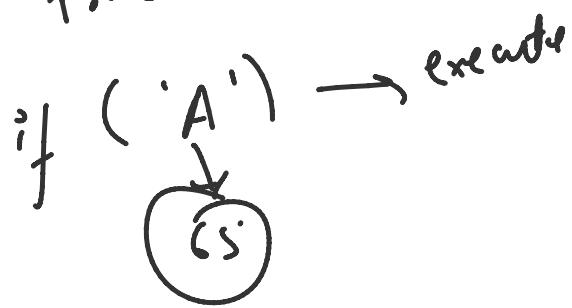
if (!true) → —;

if (0.1) → execute

if (-10) → execute.

# only false values are false and 0.  
all are true.

# only false now  
rest all are true.



Increment

2 types

int a = 10;

(i)  $a++$ .  
pre increment.

↳ pebble increment nops.  
for task nops.

(ii)  $a++;$   
post increment  
Pebble task for  
increment.

Decrement

2 types.

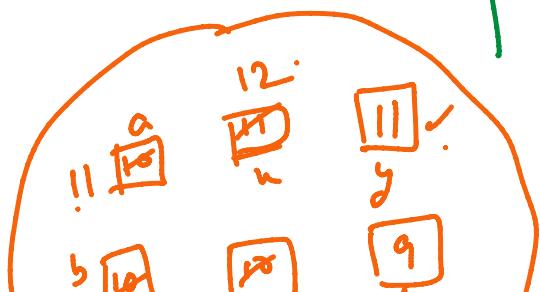
int b = 10;

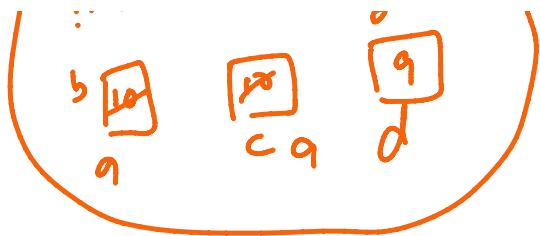
(i)  $--b$   
pebble decrement  
for task

(ii)  $b--$   
pure task  
for decrement.

int a = 10;  
int n =  $++a$ ; pre  
int y =  $n++$ ; post

int b = 10;  
int c =  $--b$ ; post  
int d =  $--c$ ; pre.





## writing blocks

Anik / Gayatri	Kavi
$a = 0$	0
$b = 1$	0
$c = 0$	6
$d = 1$	1

infra  
1, 1, 0 2.

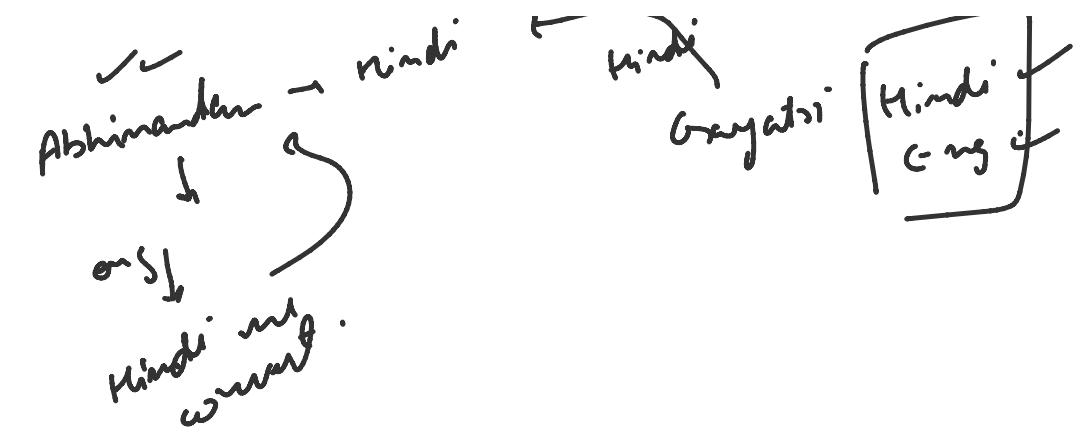
Anjali 1, 0, -1, 0.

Bitwise operators.  
operation  
right  
left  
Binary works on current.

$6 + 7 = 13$   
Binary      DNS      Binary  
6.23 + 7.23      6 + 7  
Arithmetic operators  
(needs) works on decimal  
numbers.

Bitwise operators work on BNS.

task.  
findi  
findi  
2 choice  
a: Hindi  
b: Hindi



Binary first  
It will understand quickly and also act quickly.

BNS

generally DNS.

In Bitwise operators, while they are performing their task, both numbers are present in their binary form.

Bitwise operators → &, |, ~, ^, <<, >>

Bitwise AND ' $\&$ '

All 1 → 1, else 0.

3 6 5

$$\begin{array}{r} 1 \\ \times \\ 1 \\ \hline 1 \end{array}$$

$$\begin{array}{r} 1 \\ \times \\ 0 \\ \hline 0 \end{array}$$

$$1 + 0 = 1$$

$$1 + 0 = 1$$

$$1 + 2 + 0 \times 2 + 1 \times 2^1 = 3$$

$$\begin{array}{r} 5 \\ \hline 2 \\ 2 \\ 1 \end{array}$$

$$10.1 = 5$$

$$\begin{array}{r}
 & 1 & 0 & 0 & 0 & 1 \\
 & \downarrow & & & & \downarrow \\
 2 & 1 & 7 & 0 & 0 & 1 \\
 & \downarrow & & & & \downarrow \\
 2 & 8 & 0 & 0 & 1 \\
 & \downarrow & & & \downarrow \\
 2 & 4 & 0 & 0 & 1 \\
 & \downarrow & & & \downarrow \\
 2 & 2 & 0 & 0 & 1 \\
 & \downarrow & & & \downarrow \\
 2 & 1 & 0 & 0 & 1
 \end{array}$$

$10001 = 47$

$10001$

$$\begin{array}{r}
 & 1 & 0 & 1 \\
 & \downarrow & & \downarrow \\
 2 & 1 & 3 & 1 \\
 & \downarrow & & \downarrow \\
 2 & 6 & 1 & 1 \\
 & \downarrow & & \downarrow \\
 2 & 3 & 1 & 1 \\
 & \downarrow & & \downarrow \\
 1 & 1 & 1 & 1
 \end{array}$$

$(1101) = (13)_2$

$\frac{1 \times 2^3}{8} + \frac{1 \times 2^2}{4} + \frac{0 \times 2^1}{2} + \frac{1 \times 2^0}{1} = 13$

$$13 - 8 = 5$$

$$\begin{array}{r}
 & 2 & 1 & 0 & 1 \\
 & \underline{-} & 8 & \underline{4} & \underline{2} \\
 & 4 & & & 
 \end{array}$$

$$\begin{array}{r}
 & 1 & 0 & 1 & 0 & 1 \\
 & \underline{-} & 8 & \underline{4} & \underline{2} & \underline{1} \\
 & 4 & & & & 
 \end{array}$$

$12 - 8 = 4$

$$\begin{array}{r}
 385 \\
 385 = 0. \\
 \hline
 8 \text{ rem.} \quad \begin{array}{r}
 0 & 1 & 1 \\
 \underline{-} & 1 & 0 & 1 \\
 \hline
 0 & 0 & 1
 \end{array}
 \end{array}$$

$$\begin{array}{r}
 9813 \\
 - (9) 5 \\
 \hline
 0
 \end{array}
 \quad \begin{array}{r}
 1 & 0 & 0 & 1 \\
 \underline{-} & 1 & 1 & 0 \\
 \hline
 0 & 0 & 1
 \end{array}$$

$$= \textcircled{9} \textcircled{5}$$

9:  $\begin{array}{r} 1 \\ - \\ 1 \\ - \\ 0 \\ - \\ 1 \end{array}$   
 13:  $\begin{array}{r} 1 \\ - \\ 1 \\ - \\ 0 \\ - \\ 1 \end{array}$   
 6's comp =  $\begin{array}{r} 1 \\ 0 \\ 0 \\ 1 \end{array}$

$$\begin{array}{r} 17, 13, 7 \\ \underline{17 \ 6 \ 13 \ 8 \ 7} = \textcircled{1} \\ 17 \quad \begin{array}{r} 1 \\ 0 \\ 0 \\ 0 \\ 1 \end{array} \\ 13 \quad \begin{array}{r} 0 \\ 1 \\ 1 \\ 0 \\ 1 \end{array} \\ + \quad \begin{array}{r} 0 \\ 0 \\ 1 \\ 1 \\ 1 \end{array} \\ \hline 8 \quad \begin{array}{r} 0 \\ 0 \\ 0 \\ 0 \\ 1 \end{array} \end{array}$$

Bitwise OR '1'  
 - Any 1 → 1. / All zero → 0.

$$3 | 5 = \textcircled{7}$$

$\begin{array}{r} 0 \\ 1 \\ 0 \\ 1 \end{array}$   
 4's comp  $\begin{array}{r} 1 \\ 1 \\ 1 \end{array}$

$$9 | 13 \rightarrow 9 \quad \begin{array}{r} 1 \\ 0 \\ 0 \\ 0 \\ 1 \end{array}$$

$\begin{array}{r} 1 \\ 1 \\ 0 \\ 1 \end{array}$   
 6's comp.  $\begin{array}{r} 1 \\ 1 \\ 0 \\ 1 \end{array}$

$$17 | 13 | 7$$

$\begin{array}{r} 1 \\ 0 \\ 0 \\ 0 \\ 1 \end{array}$   
 $\begin{array}{r} 0 \\ 1 \\ 1 \\ 0 \\ 1 \end{array}$   
 $\begin{array}{r} 0 \\ 0 \\ 1 \\ 1 \\ 1 \end{array}$   
 $\hline 1 \\ 1 \\ 1 \\ 1 \end{array}$

$$\text{Ques} \quad \begin{array}{r} 2 \\ \hline 1 & 1 & 1 & 1 & 1 \\ \hline 16 & 8 & 4 & 2 & 1 \end{array}$$

(31)

Bitwise Not  $\sim$  (Negation)

$$0 \rightarrow 1$$

$$1 \rightarrow 0$$

$$\begin{array}{l} \sim 7 = \\ \sim 7 = 0 \end{array}$$

$$\begin{array}{r} 000 \\ \sim 13 = 1101 \\ \hline 0010 \end{array}$$

Bitwise  $XOR$ .  
 if odd no. of  $1 \rightarrow 1$   
 else 0.

$$7 \wedge 13 = \begin{array}{r} 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ \hline 1 & 0 & 1 & 0 \end{array} \quad \text{Ans} \quad = (10)$$

$$17 \wedge 13 \wedge 7 = \begin{array}{r} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ \hline 1 & 1 & 0 & 1 \end{array} \quad ? = 1$$

= 27

$$\begin{array}{r} \underline{1} \quad \underline{1} \quad \underline{0} \quad \underline{1} \\ \underline{1} \quad \underline{0} \quad \underline{1} \quad \underline{0} \end{array} = (27)$$

# << left shift.

\*  $\ll y \rightarrow$  Shift all bits  $0\}$   
\* by  $y$  places.

$$S = 101$$

$$y = 2$$

$$S \ll 2$$

$$\begin{array}{r} \text{---} \quad \text{---} \quad \underline{0} \quad \underline{1} \\ \text{---} \quad \text{---} \quad \underline{0} \quad \underline{1} \end{array} \ll 2 = \boxed{26} \quad S \times 2^2$$

$$\begin{array}{r} \underline{1} \quad \underline{0} \quad \underline{1} \quad \underline{0} \quad \underline{1} \\ \underline{1} \quad \underline{0} \quad \underline{1} \quad \underline{0} \quad \underline{1} \end{array} = \boxed{26}$$

$$7 \rightarrow \begin{array}{r} \text{---} \quad \text{---} \quad \underline{1} \quad \underline{1} \\ \text{---} \quad \text{---} \quad \underline{1} \quad \underline{1} \end{array} \ll 3 = 7 \times 2^2$$

$$\begin{array}{r} \underline{1} \quad \underline{1} \quad \underline{1} \quad \underline{0} \quad \underline{0} \quad \underline{0} \\ 32 \quad 16 \quad 8 \quad 4 \quad 2 \quad 1 \end{array}$$

$$\boxed{56}$$

int 4 bytes  $\rightarrow$  32 bits.

$$\begin{array}{r} \text{---} \quad \text{---} \quad \underline{0} \quad \underline{0} \quad \dots \quad \dots \quad \underline{1} \quad \underline{1} \quad \underline{1} \\ \text{---} \quad \text{---} \quad \underline{0} \quad \underline{0} \quad \dots \quad \dots \quad \underline{1} \quad \underline{1} \quad \underline{1} \end{array} = 7 \cdot$$

fall off.

29 bits

If I say  $n$

$$\Rightarrow \boxed{n = n \times 2^0}$$

$\Rightarrow$  right shift operator:  
shift y bits to the right.  
 $x \gg y$

$\Rightarrow$   $x \gg 2$  

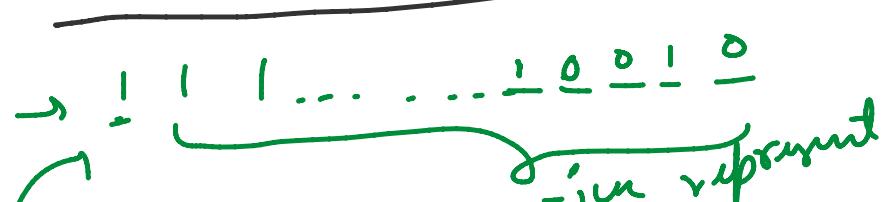
$$\begin{array}{r} 0 \\ \underline{\quad} \\ 0 \end{array} = 1$$

$$\text{If } x \gg y \Rightarrow n = \frac{x}{2^y}$$

(signed) remain ( $\sim 13$ )  
in memory  
32 bits

$$\begin{array}{r} 0 \\ \underline{\quad} \\ 0 \end{array} \dots \dots \begin{array}{r} 1 \\ \underline{\quad} \\ 1 \end{array} \begin{array}{r} 0 \\ \underline{\quad} \\ 1 \end{array}$$

$\sim$

o/p  $\rightarrow$    
 $\sim$  complement of (14).

$$\sim 1 \rightarrow -2$$

$$\sim (-2) \rightarrow 1$$

$$\gg 1$$

---

0 0 0 - .

---

If signed int , and we are doing  
Bitwise NOT on it  $\rightarrow -(num+1)$   
 $\sim num \rightarrow -(num+1)$

---

Bitmasking  $\rightarrow$  concept / questions solved.

$$7 \wedge 7 \rightarrow \begin{array}{r} 111 \\ ,\, 111 \\ \hline 111 \end{array}$$

$$5 \wedge 5 \rightarrow \begin{array}{r} 101 \\ ,\, 101 \\ \hline 000 \end{array}$$

Any number XORed with itself will  
always be 0.

int a = 10;

(Increment value of

a++ / ++a)

(a by 1)

m = a + 1000;

(Increment value of a  
by 1000)

$a = a + 1000;$        $\sim$       by 1000  
compound Assignment operators.

$a += 1000$   
 $a \underline{=} \underline{1000}$   
assignment  
original operator to be performed

$c = 10 \quad \times 10$   
equal.  $(a = a \times 10;$   
 $a *= 10;)$        $a = 10, c = 2.$

$a \ll= 2;$       0

---

ques      b, l, ~, ^, <<, >>

you are given some numbers  
in which every no. appears twice  
except 1 number. find out that  
number.

n=5      1, 2, 1, 3, 2.      I time.      ?

we will XOR each no. with

$$f_{-n^0} = 1$$

${}^0_1 \wedge {}^1_2 \wedge {}^0_3 \wedge {}^1_2 = ?$  task.

$$\begin{array}{r} 011 \\ 000 \\ \hline 011 \end{array}$$

$${}^0_1 3 - \textcircled{3}$$

$${}^1_2 2 = {}^0_?$$

$$\begin{array}{r} 010 \\ 010 \\ \hline 010 \end{array}$$

$$13 = (1101) = \textcircled{3}$$

?

you are given a number  
and you have to find out no. of  
1's in its binary form.

$$17 = 10001 = \textcircled{2}$$