

CS674A Assignment

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The Python code in the Assignment.py file is implementing a fast convolution algorithm(negative wrapped convolution) using the Number Theoretic Transform (NTT) and comparing the results with a regular polynomial convolution.

Here's a summary of what the code does:

- ➔ The code defines a few constants, including n (the size of the input arrays/polynomials), q (a prime number) and γ .
- ➔ It defines helper functions for bit-reversal, generating NTT twiddle factors, and performing the Cooley-Tukey NTT and Gentleman-Sande Inverse NTT operations.
- ➔ In the main function, it generates NTT twiddle factors and initializes two random arrays a and b of size n , representing polynomial coefficients. It also creates a polynomial f , which is used to perform polynomial division to obtain the convolution result.
- ➔ The code calculates the convolution using the regular polynomial multiplication and division operations, resulting in the polynomial p .
- ➔ It then performs the Cooley-Tukey NTT on arrays a and b , element-wise multiplication of the NTT-transformed arrays to get the convolution result c , and finally applies the Gentleman-Sande Inverse NTT to obtain the inverse transform.
- ➔ The code prints the NTT-convolved result c and checks if it is equal to the polynomial-convolved result p .

The code essentially demonstrates how to perform polynomial convolution using the NTT and validates that the NTT-based convolution matches the result obtained using traditional polynomial multiplication and division.

Time Complexity

Traditional Approach: $O(n^2)$

```
91 a = np.random.randint(0, q, n)
92 b = np.random.randint(0, q, n)
93
94 p = np remainder(np.polydiv(np.polymul(a[::-1], b[::-1]), f)[1], q).astype(int)[::-1]
95
96 print("Convolution result (p):", p)
```

- ➔ First of all, we are generating random arrays *a* and *b* of size *n* each using *np.random.randint(0, q, n)* takes $O(n)$ time.
- ➔ Now, we are performing polynomial multiplication using *np.polymul(a[::-1], b[::-1])* which takes $O(n^2)$ time because it involves multiplying each term of the first polynomial by each term of the second polynomial.
- ➔ Then, we are doing polynomial division using *np.polydiv* essentially has a time complexity of $O(n)$ because it involves dividing the two polynomials of degree *n* and yields both quotient and remainder polynomials.
- ➔ Now, the modulo operation *np.remainder* applied to the result of the polynomial division which will take $O(n)$ time, as it iterates through the coefficients.
- ➔ Then, we are converting the result to an integer array using *astype(int)* is a constant-time operation and can be considered $O(1)$.
- ➔ Finally, we are reversing the result array using *[::-1]* which is also a constant-time operation and can be considered $O(1)$.

Overall, the dominant time complexity for this code segment is the polynomial multiplication, which is **$O(n^2)$** . The other operations, such as array generation, polynomial division, and modification, have linear or constant time complexities relative to the polynomial multiplication, so they do not significantly impact the overall time complexity.

Fast Approach: $O(n \cdot \log(n))$

The code implements a fast polynomial multiplication algorithm based on the Number Theoretic Transform (NTT) using the Cooley-Tukey and Gentleman-Sande algorithms.

To discuss the time complexity in detail, let's break down the main components and their time complexities:

1) Twiddle Factor Generation:

```
14 def generate_twiddleFactor(n, gamma, q):
15     alpha = 1
16     tmp = []
17     tFactor = []
18     inv_tFactor = []
19
20     for x in range(0, n):
21         tmp.append(alpha)
22         alpha = alpha * gamma % q
23
24     positions = generate_positions(n)
25
26     for x in range(0, n):
27         val = tmp[positions[x]]
28         inv_val = pow(val, -1, q)
29         tFactor.append(val)
30         inv_tFactor.append(inv_val)
31
32     return tFactor, inv_tFactor
```

- ➔ The generate_twiddleFactor() function generates twiddle factors for the NTT.
- ➔ It computes the n twiddle factors using a loop with $O(n)$ complexity.

Therefore, the time complexity is $O(n)$ for generating twiddle factors.

2) Cooley-Tukey NTT:

```
34 def cooley_tukey_ntt(a, tFactor, q):
35     n = len(a)
36     t = n
37     m = 1
38
39     while m < n:
40         t = t // 2
41
42         for i in range(0, m):
43             j1 = 2 * i * t
44             j2 = j1 + t - 1
45             S = tFactor[m + i]
46
47             for j in range(j1, j2 + 1):
48                 U = a[j]
49                 V = a[j + t] * S
50                 a[j] = (U + V) % q
51                 a[j + t] = (U - V) % q
52
53         m = 2 * m
```

- ➔ The `cooley_tukey_ntt()` function performs the Cooley-Tukey NTT on two input arrays/polynomials `a` and `b`.
- ➔ It consists of two nested loops. The outer loop runs $\log_2(n)$ times, where n is the size of the input array, and each iteration halves the value of `t` until it reaches 1.
- ➔ The inner loop performs arithmetic operations on the input arrays, and it runs n times for each outer loop iteration.

Therefore, the overall time complexity of Cooley-Tukey NTT is $O(n * \log(n))$, where n is the size of the input arrays.

3) Gentleman-Sande Inverse NTT:

```
55 def gentleman_sande_inv_ntt(a, inv_tFactor, q):
56     n = len(a)
57     t = 1
58     m = n
59
60     while m > 1:
61         j1 = 0
62         h = m // 2
63
64         for i in range(0, h):
65             j2 = j1 + t - 1
66             S = inv_tFactor[h + i]
67
68             for j in range(j1, j2 + 1):
69                 U = a[j]
70                 V = a[j + t]
71                 a[j] = (U + V) % q
72                 a[j + t] = (U - V) * S % q
73
74             j1 = j1 + 2 * t
75
76         t = 2 * t
77         m = m // 2
78
79     n_inv = pow(n, -1, q)
80     for i in range(0, n):
81         a[i] = a[i] * n_inv % q
```

- ➔ The `gentleman_sande_inv_ntt()` function performs the inverse NTT using the Gentleman-Sande algorithm.
- ➔ Similar to Cooley-Tukey, it consists of two nested loops, with the outer loop running $\log_2(n)$ times, and each iteration halves the value of m until it reaches 1.
- ➔ The inner loop performs arithmetic operations on the input array and runs n times for each outer loop iteration.

Therefore, the overall time complexity of Gentleman-Sande Inverse NTT is $O(n * \log(n))$, where n is the size of the input array.

4) Polynomial Multiplication:

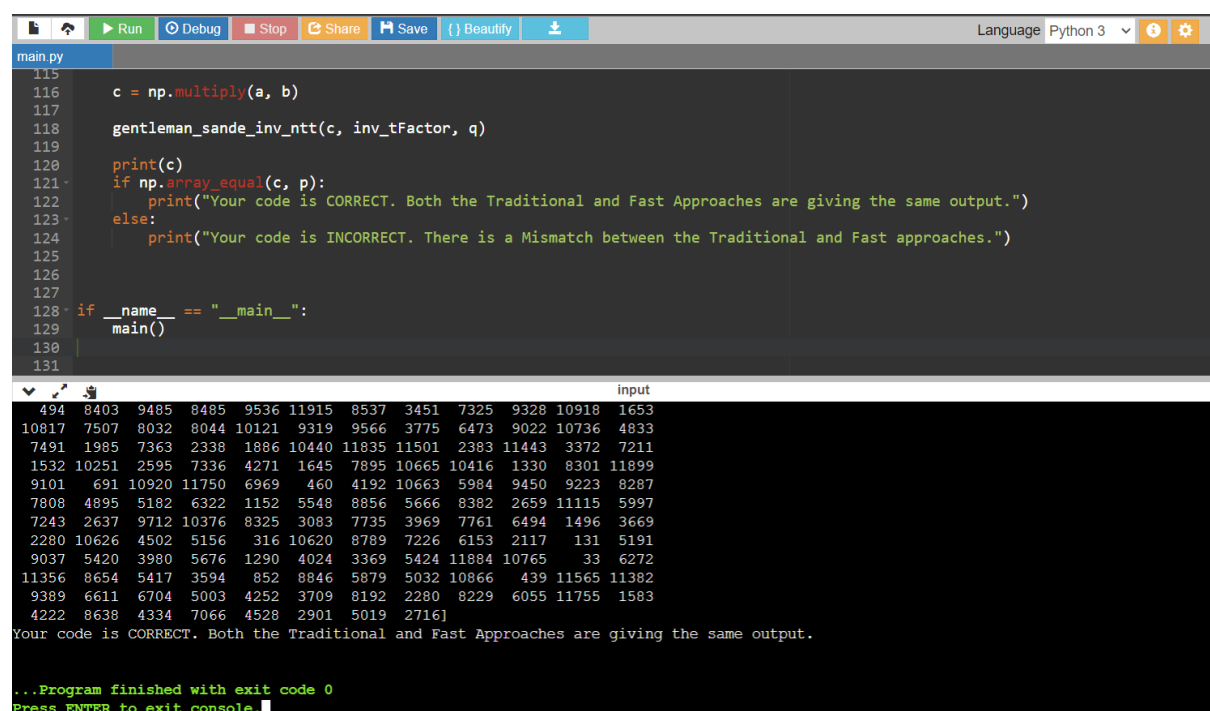
```
101 c = np.multiply(a, b)
```

The polynomial multiplication is performed using NumPy's `np.multiply()` function, which has a time complexity of $O(n)$ since it computes the element-wise product of two arrays.

5) Overall Time Complexity:

- ➔ The overall time complexity is dominated by the Cooley-Tukey NTT and Gentleman-Sande Inverse NTT, both of which have a time complexity of $O(n * \log(n))$.
- ➔ The generation of twiddle factors and polynomial multiplication contribute additional $O(n)$ time complexity.

OUTPUT:



The screenshot shows a Python IDE with a dark theme. The top toolbar includes buttons for Run, Debug, Stop, Share, Save, and Beautify. The language is set to Python 3. The code in the editor is as follows:

```
115
116     c = np.multiply(a, b)
117
118     gentleman_sande_inv_ntt(c, inv_tFactor, q)
119
120     print(c)
121     if np.array_equal(c, p):
122         print("Your code is CORRECT. Both the Traditional and Fast Approaches are giving the same output.")
123     else:
124         print("Your code is INCORRECT. There is a Mismatch between the Traditional and Fast approaches.")
125
126
127
128 if __name__ == "__main__":
129     main()
130
131
```

The output window shows a large grid of numbers, with the last line indicating the code is correct:

```
494 8403 9485 8485 9536 11915 8537 3451 7325 9328 10918 1653
10817 7507 8032 8044 10121 9319 9566 3775 6473 9022 10736 4833
7491 1985 7363 2338 1886 10440 11835 11501 2383 11443 3372 7211
1532 10251 2595 7336 4271 1645 7895 10665 10416 1330 8301 11899
9101 691 10920 11750 6969 460 4192 10663 5984 9450 9223 8287
7808 4895 5182 6322 1152 5548 8856 5666 8382 2659 11115 5997
7243 2637 9712 10376 8325 3083 7735 3969 7761 6494 1496 3669
2280 10626 4502 5156 316 10620 8789 7226 6153 2117 131 5191
9037 5420 3980 5676 1290 4024 3369 5424 11884 10765 33 6272
11356 8654 5417 3594 852 8846 5879 5032 10866 439 11565 11382
9389 6611 6704 5003 4252 3709 8192 2280 8229 6055 11755 1583
4222 8638 4334 7066 4528 2901 5019 2716]
Your code is CORRECT. Both the Traditional and Fast Approaches are giving the same output.

...Program finished with exit code 0
Press ENTER to exit console.
```

Warning: The Python code is running without issues on online compilers and macOS, but it encounters overflow errors on Windows computers.

In summary, the provided code has an overall time complexity of $O(n * \log(n))$ due to the Cooley-Tukey and Gentleman-Sande NTT algorithms, with some additional $O(n)$ complexity for twiddle factor generation and polynomial operations.

Thank you 😊