

## MA 251 Data Structures

### Laboratory Assignment 6

25-09-2019

Note: Upload your programs to the server (deadline: 4:30 pm)

1. **Merge sort with arrays:** Write a program to implement merge sort using arrays. As discussed in class, the procedure MERGE-SORT( $A, p, r$ ) sorts the elements in the subarray  $A[p \dots r]$ . The pseudocode of Merge sort is given below.

```
MERGE-SORT( $A, p, r$ )
1  if  $p < r$ 
2     $q = \lfloor (p + r)/2 \rfloor$ 
3    MERGE-SORT( $A, p, q$ )
4    MERGE-SORT( $A, q + 1, r$ )
5    MERGE( $A, p, q, r$ )
```

The procedure of MERGE() is given overleaf. Your code should be implementation of these pseudocodes.

In the main function,

- populate an array **A** with **n** random integers in the range  $[0, 10^3-1]$ . Take **n=10<sup>k</sup>**, where  $k = 1, 2$  and  $3$ .
- Call Merge sort for the three different values of  $k$  and compute the time taken by each call.
- For each value of  $k$ , print the unsorted array (row wise in line 1), sorted array (row wise in line 2) and the time taken (in third line).

#### Hint:

The *random* nos. can be generated by the *rand()* function. Do not forget to *seed* the *rand()* function with the current time.

The predefined *clock()* function can be used to measure the running time. An example code segment is

```
c1 = clock();
.... Call Merge sort ...
c2 = clock();
runtime = c2 - c1;
```

For small values of **n**, it is likely that you will get the runtime to be 0. In order to ensure a more accurate timing measurement, for a particular value of **n** (say 10), perform the sort experiment 1000 times in a row and then divide the total elapsed time by 1000. This strategy gives you a timing measurement that is much more accurate.

2. **Merge sort with linked list:** Repeat the above assignment using linked lists.

### Pseudocode of MERGE

MERGE( $A, p, q, r$ )

```
1   $n_1 = q - p + 1$ 
2   $n_2 = r - q$ 
3  let  $L[1 \dots n_1 + 1]$  and  $R[1 \dots n_2 + 1]$  be new arrays
4  for  $i = 1$  to  $n_1$ 
5       $L[i] = A[p + i - 1]$ 
6  for  $j = 1$  to  $n_2$ 
7       $R[j] = A[q + j]$ 
8   $L[n_1 + 1] = \infty$ 
9   $R[n_2 + 1] = \infty$ 
10  $i = 1$ 
11  $j = 1$ 
12 for  $k = p$  to  $r$ 
13     if  $L[i] \leq R[j]$ 
14          $A[k] = L[i]$ 
15          $i = i + 1$ 
16     else  $A[k] = R[j]$ 
17          $j = j + 1$ 
```