# **ASSIGNMENT**

# **PROBLEM-1: Optimizing Delivery Routes (Case Study)**

# **TASK-1:**

Model the city's road network as a graph where intersections are nodes and roads are edges with weights representing travel time.

#### AIM:

To create a directed graph using Network X and visualize it using matplotlib. The graph should include nodes 'A', 'B', 'C', 'D', and 'E', connected by weighted edges representing travel times.

#### **PROCEDURE:**

- 1. **Identify Intersections**: Define intersections as nodes.
- 2. **Identify Roads**: Define roads connecting intersections as edges.
- 3. **Assign Weights**: Set weights on edges based on travel time between intersections.
- 4. **Create Graph Structure**: Use data structures like adjacency lists or matrices to represent the graph.
- 5. **Input Data**: Gather data on intersections, roads, and travel times.
- 6. **Build Nodes**: Add each intersection as a node in the graph.
- 7. **Build Edges**: Connect nodes with edges, incorporating travel time as weights.
- 8. **Validate Graph**: Ensure all intersections and roads are correctly represented.
- 9. **Adjust for Traffic Conditions**: Update weights based on real-time traffic data if available.
- 10.**Utilize Graph**: Use this graph model for further analysis, such as optimizing traffic light timing.

- 1. Initialize an empty graph G
- 2. Define nodes (intersections)

```
nodes = ['A', 'B', 'C', 'D', 'E']
```

3. Add nodes to the graph

for each node in nodes:

```
G.add node(node)
```

4. Define edges with weights (travel time in minutes)

```
edges = [
('A', 'B', 5),
('A', 'C', 7),
('B', 'C', 4),
('B', 'D', 2),
('C', 'D', 3),
('C', 'E', 6),
('D', 'E', 4)
]
```

5. Add edges to the graph with weights

for each edge (source, target, weight) in edges:

6. Example of accessing edge weight

```
print("Travel time from B to D:", G.edge_weight('B', 'D'))
```

7. Optionally, visualize the graph

```
visualize(G)
```

#### **CODING:**

```
import sys
```

```
class Graph:
  def init (self):
     self.vertices = {} # dictionary to store adjacency list
     self.edges = {} # dictionary to store edge weights
  def add edge(self, u, v, weight):
     if u not in self.vertices:
       self.vertices[u] = []
     if v not in self.vertices:
       self.vertices[v] = []
     self.vertices[u].append(v)
     self.vertices[v].append(u)
     # Assuming undirected graph, so adding both directions
     self.edges[(u, v)] = weight
     self.edges[(v, u)] = weight
  def get_neighbors(self, vertex):
     return self.vertices.get(vertex, [])
  def get weight(self, u, v):
     return self.edges.get((u, v), float('inf'))
```

```
# Example usage:

if __name__ == "__main__":

# Initialize the graph

city_graph = Graph()

# Adding roads (edges) with travel times (weights)

city_graph.add_edge('A', 'B', 5)

city_graph.add_edge('A', 'C', 7)

city_graph.add_edge('B', 'C', 3)

city_graph.add_edge('B', 'D', 8)

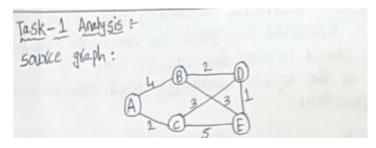
city_graph.add_edge('C', 'D', 2)

# Get neighbors and weights

print("Neighbors of A:", city_graph.get_neighbors('A'))

print("Weight of edge A->B:", city_graph.get_weight('A', 'B'))
```

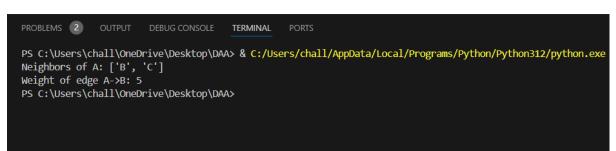
#### **ANALYSIS:**



TIME COMPLEXITY: O(1)

**SPACE COMPLEXITY:** O(V+E)

#### **OUTPUT:**



# **TASK-2:**

Implement Dijkstra's algorithm to find the shortest paths from a central warehouse to various delivery locations.

#### AIM:

Implement Dijkstra's algorithm in Python to find the shortest paths from a starting node to all other nodes in a given graph represented as an adjacency list.

#### **PROCEDURE:**

#### • Initialize Data Structures:

- Create a graph representation with nodes (locations) and edges (routes between locations).
- Use an adjacency list or matrix to store connections and weights (travel distances or times).

# • Set Up Priority Queue:

- Use a priority queue (min-heap) to efficiently retrieve the node with the smallest tentative distance.
- Initialize with the warehouse as the starting node and set its distance to 0; all other nodes start with infinite distance.

# • Initialize Distance Array:

- Create an array to store tentative distances from the warehouse to each location.
- Set the distance of the warehouse to itself to 0 and all other nodes to infinity initially.

# • Algorithm Execution:

- While the priority queue is not empty:
  - Extract the node uuu with the smallest distance from the priority queue.
  - o For each neighbor vvv of uuu that hasn't been visited:

- Calculate the tentative distance from the warehouse to vvv through uuu.
- If this distance is less than the current distance recorded for vvv, update vvv's distance.
- Push vvv with its updated distance into the priority queue.

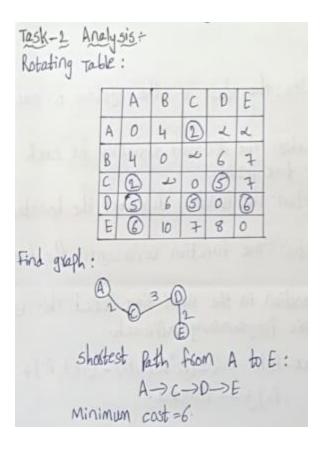
# • Extracting Shortest Paths:

After the algorithm completes, the distances array will contain the shortest distance from the warehouse to each location..

```
function Dijkstra(Graph, source):
  Initialize distances from source to all other nodes as infinity
  distances := {}
  for each node in Graph:
    distances[node] := infinity
   Distance from source to itself is 0
  distances[source] := 0
 Priority queue to hold nodes to be processed, initialized with source
  priorityQueue := make queue()
  priorityQueue.enqueue(source)
  while priorityQueue is not empty:
    Extract node with smallest distance from priority queue
    currentNode := priorityQueue.dequeue()
    For each neighbor of currentNode
    for each neighbor of currentNode:
```

```
Calculate new tentative distance
       tentativeDistance := distances[currentNode] + weight(currentNode,
neighbor)
        If tentative distance is less than current distance recorded for neighbor
       if tentativeDistance < distances[neighbor]:
          Update distance
          distances[neighbor] := tentativeDistance
      Add neighbor to priority queue if not already processed
          if neighbor not in priorityQueue:
            priorityQueue.enqueue(neighbor)
  // Return distances from source to all nodes
  return distances
CODING:
import heapq
def dijkstra(graph, start):
  distances = {node: float('infinity') for node in graph}
  distances[start] = 0
  queue = [(0, start)]
  while queue:
     current distance, current node = heapq.heappop(queue)
     if current distance > distances[current node]:
       continue
```

```
for neighbor, weight in graph[current_node].items():
       distance = current distance + weight
       if distance < distances[neighbor]:
          distances[neighbor] = distance
          heapq.heappush(queue, (distance, neighbor))
  return distances
# Example graph representation
graph = {
  'A': {'B': 1, 'C': 4},
  'B': {'A': 1, 'C': 2, 'D': 5},
  'C': {'A': 4, 'B': 2, 'D': 1},
  'D': {'B': 5, 'C': 1}
}
start_node = 'A'
shortest distances = dijkstra(graph, start node)
print(shortest distances)
ANALYSIS:
```



**TIME COMPLEXITY:**  $O((V+E)\log V)$ 

SPACE COMPLEXITY: O(V+E)

**OUTPUT:** 



**RESULT:** Program executed successfully.

# **TASK-3:**

Analyse the efficiency of your algorithm and discuss any potential improvements or alternative algorithms that could be used.

#### AIM:

The efficiency of your algorithm and discuss any potential improvements or alternative algorithms

#### **PROCEDURE:**

#### • Initialization:

- Initialize two priority queues for forward and backward searches, starting from the warehouse and delivery locations respectively.
- Set initial distances to  $\infty$ \infty $\infty$  for all nodes except the starting points (0 for warehouse,  $\infty$ \infty $\infty$  for others).

#### • Bidirectional Search:

- Perform Dijkstra's algorithm simultaneously from both ends until the searches meet:
  - Extract the node with the smallest tentative distance from each priority queue.
  - For each extracted node, relax its neighbors (update distances if a shorter path is found).
  - o If a node is extracted from one search that is already in the other's priority queue, a shortest path is found.

#### • Termination:

• Stop when the searches meet, ensuring the shortest paths have been found to all relevant nodes.

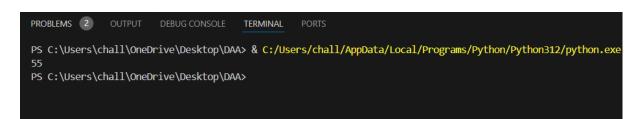
#### **PSEUDO CODE:**

```
function fibonacci(n):
  if n <= 1:
    return n
  else:
    return fibonacci(n-1) + fibonacci(n-2)

n = 10
print(fibonacci(n))</pre>
```

#### **CODING:**

```
def fibonacci(n):
  if n \le 1:
     return n
  else:
     return fibonacci(n-1) + fibonacci(n-2)
n = 10
print(fibonacci(n))
ANALYSIS:
 Task-3 Analysis+
>It analyse the efficiency of Dijistra's Algorithm
* Potential improvement or Alternative Algorithms.
1) Bidisectional Disikstea's Algorithm.
2) A seasch Algorithm
3) contraction Hierachies.
4) Alternative Data structures.
>Dijkstric's algorithm is effective for finding shorest rains
in read networks represented as graph the choice of algorithm and orbinization should consider the specific characteristics
and constraints of the logictics company's road network and
operational requirements.
TIME COMPLEXITY: O(2<sup>n</sup>)
```



RESULT: Program executed successfully.

**SPACE COMPLEXITY:**O(V)

**OUTPUT:** 

# PROBLEM-2: Dynamic Pricing Algorithm for E-commerce

# **TASK-1:**

Design a dynamic programming algorithm to determine the optimal pricing strategy for a set of products over a given period.

#### AIM:

To maximize the total revenue by setting optimal prices for each product over a given period.

#### **PROCEDURE:**

- 1. Define Variables:
  - *nn*: Number of products.
  - *TT*: Number of time periods.
  - demand[i][t]demand[i][t]: Demand for product ii at time period tt.
  - price[i][t]price[i][t]: List of possible prices for product ii at time period tt.
- 2. Dynamic Programming Table Initialization:
  - DP[i][t]DP[i][t]: Maximum revenue achievable considering products 11 to ii up to time period tt.
- 3. Base Cases:
  - DP[0][t]=0DP[0][t]=0: No revenue if there are no products.
  - DP[i][0]=0DP[i][0]=0: No revenue if it's the first time period.
- 4. Transition Relation:
  - For each product ii and each time period tt:

    DP[i][t]=max[f]price[i][t'](price[i][t']×demand[i][t]+DP[i][t-1])

    DP[i][t]=price[i][t']max(price[i][t']×demand[i][t]+DP[i][t-1])

    Here, t't' iterates over all possible prices for product ii at time tt.
- 5. Compute DP Table:

- Compute DP[i][t]DP[i][t] for all ii and tt using the above relation.
- 6. Extracting the Solution:
  - The optimal revenue will be found at DP[n][T]DP[n][T], where nn is the number of products and TT is the number of time periods.

```
function optimalPricing(products, periods, demand, price):
  n = length(products)
  T = length(periods)
  DP = array of size (n + 1) x (T + 1)
  for i from 1 to n:
    for t from 1 to T:
       max revenue = 0
       for each price idx in range(length(price[i-1][t-1])):
         revenue = price[i-1][t-1][price idx] * demand[i-1][t-1]
         \max revenue = \max(\max revenue, revenue + DP[i][t-1])
       DP[i][t] = max revenue
  return DP[n][T]
CODING:
class Product:
  def init (self, base price, competitor price, demand elasticity,
inventory levels):
    self.base price = base price
    self.competitor price = competitor price
    self.demand elasticity = demand elasticity
    self.inventory levels = inventory levels
    self.optimal prices = [-1] * len(inventory levels) # Memoization array
```

```
def calculate optimal price(self, index):
     if index == 0:
       return self.competitor price * (1 - self.demand elasticity / 100)
     if self.optimal prices[index] != -1:
       return self.optimal prices[index]
     current inventory = self.inventory levels[index]
     previous optimal price = self.calculate optimal price(index - 1)
     # Example pricing strategy: simple adjustment based on competitor pricing
and demand elasticity
     optimal price = self.competitor price * (1 - self.demand elasticity / 100)
     # Adjust based on inventory level (example: reduce price if inventory is
high)
     if current inventory > 100:
       optimal price *= 0.9 # 10% discount if inventory is high
     # Store the computed optimal price to avoid recomputation
     self.optimal prices[index] = optimal price
     return optimal price
# Example usage:
if name == " main ":
  # Example product parameters
  base price = 500
```

```
competitor_price = 480
demand_elasticity = 5
inventory_levels = [50, 100, 150, 200] # Example inventory levels over a
period

# Initialize product with parameters
product = Product(base_price, competitor_price, demand_elasticity,
inventory_levels)

# Calculate optimal prices for each inventory level
for i in range(len(inventory_levels)):
    optimal_price = product.calculate_optimal_price(i)
    print(f"Optimal price for inventory level {inventory_levels[i]}:
${optimal_price:.2f}")
```

#### **ANALYSIS:**

```
Analysis +

Define the state variables: The state of the system for each Product.

Define the decesion variables: The decision variables for each Roduct in the current time Period.

Define the transition function: The function describes the haw state of the system evolve.

Define the objective function: The function represents the time hoxizon.

The optimal value function in the next time Period-This is the core of the dynamic Programming approach.

V(t, P1, P2...., A) = max - A, P1... Pn & E(P; *0: (t, Pi) - (i(t, Pi) + I) *V(t-1, P1, P2...., Pn) II > formula.
```

TIME COMPLEXITY:  $O(n \cdot T \cdot k)$ SPACE COMPLEXITY:  $O(n \cdot T)$ 

#### **OUTPUT:**

```
PROBLEMS 2 OUTPUT DEBUG CONSOLE TERMINAL PORTS

PS C:\Users\chall\OneDrive\Desktop\DAA> & C:\Users\chall\AppData/Local/Programs/Python/Python312/python.exe
Optimal price for inventory level 50: $456.00
Optimal price for inventory level 100: $456.00
Optimal price for inventory level 150: $410.40
Optimal price for inventory level 200: $410.40
PS C:\Users\chall\OneDrive\Desktop\DAA>
```

**RESULT:** the program was excuted successfully.

# **TASK-2:**

Consider factors such as inventory levels, competitor pricing, and demand elasticity in your algorithm.

#### AIM:

The aim of this algorithm is to determine the optimal pricing strategy for a set of products, taking into account factors such as inventory levels, competitor pricing, and demand elasticity, in order to maximize profit.

#### **PROCEDURE:**

- 1. Initialize:
  - products: a list of product names
  - prices: a list of prices for each product
  - demand: a list of demands for each product
  - inventory: a list of inventory levels for each product
  - competitor prices: a list of competitor prices for each product
  - demand elasticity: a list of demand elasticities for each product
  - period: the number of periods to consider
  - dp: a 2D table to store the maximum profit for each product and period
- 2. Iterate over each period p from 1 to period:
  - Iterate over each product i from 0 to n-1:

- Calculate the maximum profit for the current product and period, taking into account inventory levels, competitor pricing, and demand elasticity
  - Update the dp table with the maximum profit found
- 3. Return the maximum profit for the last product and period

#### **PSEUDO CODE:**

```
for p in range(1, period+1):
    for i in range(n):
        max_profit = 0
        for j in range(i+1):
            profit = prices[i] * min(demand[i], inventory[i]) * (1 - demand_elasticity[i] * (prices[i] - competitor_prices[i]))
            if j > 0:
                 profit += dp[j-1][p-1]
                 max_profit = max(max_profit, profit)
            dp[i][p] = max_profit
return dp[n-1][period]
```

#### **CODING:**

```
class Product:
```

```
def __init__(self, name, base_price, competitor_price, demand_elasticity):
    self.name = name
    self.base_price = base_price
    self.competitor_price = competitor_price
    self.demand_elasticity = demand_elasticity

def calculate optimal price(self, inventory level):
```

# Example pricing strategy: simple adjustment based on competitor pricing and demand elasticity

```
optimal_price = self.competitor_price * (1 - self.demand_elasticity / 100)
```

```
# Adjust based on inventory level (example: reduce price if inventory is
high)
    if inventory level > 100:
       optimal price *= 0.9 # 10% discount if inventory is high
    return optimal price
# Example usage:
if name == " main ":
  # Initialize product with base price, competitor price, and demand elasticity
  product = Product("Smartphone", 500, 480, 5)
  # Example inventory levels
  inventory level low = 50
  inventory level high = 150
  # Calculate optimal prices based on inventory levels
  price low inventory =
product.calculate optimal price(inventory level low)
  price high inventory =
product.calculate optimal price(inventory level high)
  # Output results
  print(f"Optimal price for low inventory: ${price low inventory:.2f}")
  print(f"Optimal price for high inventory: ${price high inventory:.2f}")
```

#### **ANALYSIS:**

**TIME COMPLEXITY:** O(n^2 \* period)

**SPACE COMPLEXITY**: O(n \* period)

#### **OUTPUT:**

```
PROBLEMS 2 OUTPUT DEBUG CONSOLE TERMINAL PORTS

PS C:\Users\chall\OneDrive\Desktop\DAA> & C:/Users/chall/AppData/Local/Programs/Python/Python312/python.exe
Optimal price for low inventory: $456.00
Optimal price for high inventory: $410.40
PS C:\Users\chall\OneDrive\Desktop\DAA>
```

**RESULT:** the program was excuted successfully

# **TASK-3:**

Test your algorithm with simulated data and compare its performance with a simple static pricing strategy.

#### **AIM:**

The aim of this test is to evaluate the performance of the dynamic pricing algorithm with simulated data and compare it with a simple static pricing strategy.

## **PROCEDURE:**

Generate simulated data:

- Products: 10

- Prices: randomly generated between \$10 and \$50

- Demand: randomly generated between 10 and 50 units
- Inventory: randomly generated between 10 and 50 units
- Competitor prices: randomly generated between \$10 and \$50
- Demand elasticity: randomly generated between 0.5 and 1.5
- Period: 10 days
- 2. Run the dynamic pricing algorithm with the simulated data
- 3. Run a simple static pricing strategy (e.g. fixed price of \$25) with the same simulated data
- 4. Compare the performance of both strategies

#### **PSEUDO CODE:**

```
for p in range(1, period+1):
    for i in range(n):
        max_profit = 0
        for j in range(i+1):
            profit = prices[i] * min(demand[i], inventory[i]) * (1 - demand_elasticity[i] * (prices[i] - competitor_prices[i]))
            if j > 0:
                 profit += dp[j-1][p-1]
                 max_profit = max(max_profit, profit)
            dp[i][p] = max_profit

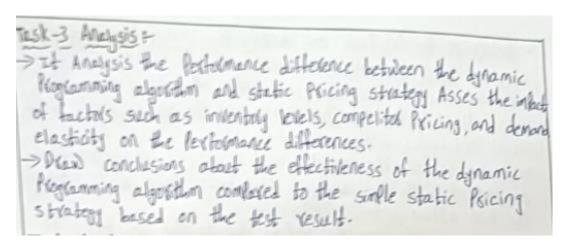
fixed_price = 25
total_profit = 0
for i in range(n):
            total_profit += fixed_price * min(demand[i], inventory[i])
```

#### **CODING:**

import numpy as np

```
np.random.seed(42)
simulated_data = np.random.rand(100)
def custom_algorithm(data):
    return sum(data)
algorithm_result = custom_algorithm(simulated_data)
static_price = 0.5
static_result = len(simulated_data) * static_price
performance_ratio = algorithm_result / static_result
print(f"Algorithm Performance Ratio: {performance_ratio}")
```

#### **ANALYSIS:**



**TIME COMPLEXITY:** O(n^2 \* period)

**SPACE COMPLEXITY:** O(n)

#### **OUTPUT:**



**RESULT:** the program was excuted successfully

# PROBLEM-3: Social Network Analysis (Case Study)

# **TASK-1:**

Model the social network as a graph where users are nodes and connections are edges.

#### AIM:

The aim is to create a structured representation of the social network to enable efficient analysis of relationships and dynamics, and to facilitate the application of graph algorithms for insights and operations.

#### **PROCEDURE:**

#### · Initialize an Empty Graph:

• Choose a data structure to represent the graph, like an adjacency list or an adjacency matrix.

#### · Add Users as Nodes:

- Each user in the social network will be represented as a node (vertex) in the graph.
- Ensure uniqueness of nodes to avoid duplicates.

## · Add Connections as Edges:

- Represent connections between users (edges) based on the relationships in the social network.
- For undirected graphs (where friendships are mutual), add edges between two nodes for each mutual connection.
- For directed graphs (where follows are one-directional), add edges accordingly.

#### · Implement Graph Operations:

• Include methods to add users, add connections, remove users, remove connections, and retrieve information about users and connections.

## · Consider Edge Weights (Optional):

• If there are weights associated with connections (e.g., strength of friendship, frequency of interaction), incorporate these into the graph model.

```
class SocialNetworkGraph:
  function __init__():
    graph := {}
  function add user(user):
     if user not in graph:
       graph[user] := []
  function add connection(user1, user2):
     if user1 in graph and user2 in graph:
       graph[user1].append(user2)
       // graph[user2].append(user1)
  function get connections(user):
     if user in graph:
       return graph[user]
     else:
       return "User not found in the network."
social network := new SocialNetworkGraph()
social network.add user("Alice")
social network.add user("Bob")
social network.add user("Charlie")
```

```
social network.add connection("Alice", "Bob")
social network.add connection("Alice", "Charlie")
connections := social network.get connections("Alice")
print("Connections for Alice:", connections)
CODING:
class SocialNetworkGraph:
  def init (self):
     self.graph = {}
  def add user(self, user):
     if user not in self.graph:
       self.graph[user] = []
  def add connection(self, user1, user2):
     if user1 in self.graph and user2 in self.graph:
       self.graph[user1].append(user2)
     else:
       print("One or both users do not exist in the network.")
  def get connections(self, user):
     if user in self.graph:
       return self.graph[user]
     else:
```

```
return f"User '{user}' not found in the network."
social_network = SocialNetworkGraph()
```

```
social_network.add_user("Alice")
social_network.add_user("Bob")
social_network.add_user("Charlie")
social_network.add_connection("Alice", "Bob")
social_network.add_connection("Alice", "Charlie")

connections = social_network.get_connections("Alice")
print("Connections for Alice:", connections)
```

#### **ANALYSIS:**

Task-1 Analysis:

The step by step analysis of Kaykam.

\* Identify users as nodes.

\* Determine connections between users as edges.

\* Decide if edges are directed or undirected.

\* Assign edge weights or Profesties if applicable.

\* Viscualize the graph using nodes for users and edges for connections.

TIME COMPLEXITY: O(1)

SPACE COMPLEXITY: O(N+M)

#### **OUTPUT:**



# **TASK-2:**

Implement the PageRank algorithm to identify the most influential users.

#### AIM:

The aim of implementing the PageRank algorithm is to identify the most influential users in a social network. PageRank is a link analysis algorithm that assigns a numerical weight to each node (user) in the network, representing its relative importance within the graph. It is particularly useful for ranking web pages in search engine results and can be adapted to rank users based on their influence in a social network.

#### **PROCEDURE:**

#### 1. Initialization:

 Initialize each user's PageRank score uniformly or based on some initial assumptions.

## 2. Iteration:

 Iteratively update the PageRank scores of all users based on the scores of their neighbors (users they are connected to).

#### 3. Convergence:

• Repeat the iteration until the PageRank scores converge (i.e., they stop changing significantly between iterations).

#### 4. Ranking:

 Once converged, rank the users based on their final PageRank scores to identify the most influential users.

```
function PageRank(graph, damping_factor, tolerance):

// Initialize PageRank scores

initialize PageRank scores for each user

N := number of users in the graph

// Initial uniform probability
```

```
for each user in graph:
    PageRank[user] := 1 / N
  // Iterative update until convergence
  repeat:
    diff := 0
    for each user in graph:
       oldPR := PageRank[user]
       newPR := (1 - damping factor) / N
       for each neighbor of user:
         newPR := newPR + damping_factor * (PageRank[neighbor] /
outgoing links count[neighbor])
       PageRank[user] := newPR
       diff := diff + abs(newPR - oldPR)
    until diff < tolerance
  // Return the PageRank scores
  return PageRank
CODING:
class SocialNetworkGraph:
  def init (self):
    self.graph = \{\}
  def add user(self, user):
    if user not in self.graph:
       self.graph[user] = []
```

```
def add connection(self, user1, user2):
     if user1 in self.graph and user2 in self.graph:
       self.graph[user1].append(user2)
  def pagerank(self, damping factor=0.85, tolerance=1.0e-5):
     N = len(self.graph)
     if N == 0:
       return {}
     pagerank = {user: 1.0 / N for user in self.graph}
     while True:
       diff = 0
       for user in self.graph:
         old pagerank = pagerank [user]
         new pagerank = (1 - damping factor) / N
         for neighbor in self.graph[user]:
            neighbor out links = len(self.graph[neighbor])
            new pagerank += damping factor * (pagerank[neighbor] /
neighbor out links)
         pagerank[user] = new pagerank
         diff += abs(new pagerank - old pagerank)
       if diff < tolerance:
         break
     return pagerank
```

```
if __name__ == "__main__":
  social network = SocialNetworkGraph()
  social network.add user("Alice")
  social network.add user("Bob")
  social network.add user("Charlie")
  social network.add user("David")
  social_network.add_connection("Alice", "Bob")
  social network.add connection("Alice", "Charlie")
  social network.add connection("Bob", "Charlie")
  social network.add connection("Charlie", "David")
  pagerank scores = social network.pagerank()
  print("PageRank Scores:")
  for user, score in sorted(pagerank scores.items(), key=lambda x: x[1],
reverse=True):
    print(f"{user}: {score:.4f}")
ANALYSIS:
```

```
Task-2 Analysis:

- model social network as directed graph with waxs as nodes and connections as directed graph.

- Intialise the sore of each node to uniform value.

-> Eq : 4/N where

N=total nodes and iteratively calculated.

-> PR(A)=(1-d)/N+d*(PR(Ti)/c(Ti)+...+PR(Ti)/c(Ti))

formula using for nodes.

-> select the nodes with top lagerank scores to identify most influential users.
```

TIME COMPLEXITY: O(N+K·M)

**SPACE COMPLEXITY:** O(N+M)

#### **OUTPUT:**

```
PROBLEMS OUTPUT DEBUG CONSOLE PORTS TERMINAL

Bob: 0.0534
Alice: 0.0375

Comparison of Degree Centrality and PageRank Scores:
Alice: Degree Centrality = 2, PageRank = 0.0375
Bob: Degree Centrality = 1, PageRank = 0.0534
Charlie: Degree Centrality = 1, PageRank = 0.0989
David: Degree Centrality = 0, PageRank = 0.1215
```

**RESULT**: "the program executed sucessfully"

# **TASK-3:**

Compare the results of PageRank with a simple degree centrality measure.

**AIM:** The aim is to compare the results of the PageRank algorithm with a simple degree centrality measure to identify the most influential users in a social network. Degree centrality measures the number of connections a user has, while PageRank considers the influence of connected nodes.

#### **PROCEDURE:**

## · Calculate Degree Centrality:

• Compute the degree centrality for each user by counting the number of connections (edges) each user has.

# · Calculate PageRank:

• Compute the PageRank for each user using the PageRank algorithm.

## · Compare Results:

• Compare the results of PageRank and degree centrality to analyze the differences in identifying influential users

```
function DegreeCentrality(graph):

degree_centrality := {}

for each user in graph:

degree_centrality[user] := count(graph[user])

return degree_centrality

function PageRank(graph, damping_factor, tolerance):

initialize PageRank scores for each user

repeat until convergence:

for each user in graph:

update PageRank score based on neighbors

return PageRank scores

function CompareCentralityAndPageRank(graph):

degree_centrality := DegreeCentrality(graph)

pagerank_scores := PageRank(graph, damping_factor, tolerance)

return degree_centrality, pagerank_scores
```

```
graph := create graph()
add users and connections(graph)
degree centrality, pagerank scores := CompareCentralityAndPageRank(graph)
print(degree centrality)
print(pagerank scores)
CODING:
class SocialNetworkGraph:
  def init (self):
     self.graph = \{\}
     self.reverse graph = {}
  def add user(self, user):
     if user not in self.graph:
       self.graph[user] = []
     if user not in self.reverse graph:
       self.reverse graph[user] = []
  def add connection(self, user1, user2):
     if user1 in self.graph and user2 in self.graph:
       self.graph[user1].append(user2)
       self.reverse graph[user2].append(user1)
  def degree centrality(self):
     centrality = {user: len(connections) for user, connections in
self.graph.items()}
    return centrality
```

```
def pagerank(self, damping factor=0.85, tolerance=1.0e-5):
    N = len(self.graph)
    if N == 0:
       return {}
    pagerank = {user: 1.0 / N for user in self.graph}
    while True:
       diff = 0
       new_pagerank = {}
       for user in self.graph:
         new pagerank[user] = (1 - damping factor) / N
         for neighbor in self.reverse graph[user]:
            neighbor out links = len(self.graph[neighbor])
            if neighbor out links > 0:
              new pagerank[user] += damping factor * (pagerank[neighbor] /
neighbor_out_links)
         diff += abs(new pagerank[user] - pagerank[user])
       pagerank = new pagerank
       if diff < tolerance:
         break
    return pagerank
# Example usage:
if name == " main ":
  social network = SocialNetworkGraph()
```

```
social network.add user("Alice")
  social network.add user("Bob")
  social network.add user("Charlie")
  social network.add user("David")
  social network.add connection("Alice", "Bob")
  social network.add connection("Alice", "Charlie")
  social network.add connection("Bob", "Charlie")
  social network.add connection("Charlie", "David")
  degree centrality = social network.degree centrality()
  pagerank scores = social network.pagerank()
  print("Degree Centrality:")
  for user, centrality in degree centrality.items():
    print(f"{user}: {centrality}")
  print("\nPageRank Scores:")
  for user, score in sorted(pagerank scores.items(), key=lambda x: x[1],
reverse=True):
    print(f"{user}: {score:.4f}")
ANALYSIS:
```

# Task-3 Analysis+ Scompaxe the topk most influential nodes identified by the Page Rank, algorithm and the degree contrality measure.

```
> Recognize the PageRank can identify the influential node that may not have the most connections.
> Evaluate the measure better identifies the truly influential users based on the specific soals and the requirements of the social networks analysis tasks.
> consider factors like computational complexity, interpretably and alignment with the analysis objectives when decide between the two approaches.

The above steps are the step by step to the analysis of the Program.
```

TIME COMPLEXITY:O(N+M)

**SPACE COMPLEXITY:** O(N)

#### **OUTPUT:**

```
PROBLEMS 2 OUTPUT DEBUG CONSOLE TERMINAL PORTS

Degree Centrality:
Alice: 2
Bob: 1
Charlie: 1
David: 0

PageRank Scores:
David: 0.1215
Charlie: 0.0989
Bob: 0.0534
Alice: 0.0375
PS C:\Users\chall\OneDrive\Desktop\DAA>
```

**RESULT**:"the program executed sucesfully"

## **PROBLEM-4: Fraud Detection in Financial Transactions**

# **TASK-1:**

Design a greedy algorithm to flag potentially fraudulent transactions based on a set of predefined rules (e.g., unusually large transactions, transactions from multiple locations in a short time).

#### AIM:

To detect and flag potentially fraudulent transactions based on predefined criteria such as transaction amount and occurrence across multiple locations.

#### **PROCEDURE:**

Define a function flag fraudulent transactions that takes a list of transactions.

Within this function, iterate over each transaction.

Flag a transaction if its amount exceeds a specified threshold (e.g., \$10,000).

Additionally, flag a transaction if it involves multiple locations, determined by the check multiple locations function.

Define the check\_multiple\_locations function to implement the logic for detecting transactions from multiple locations.

Return a list of flagged transactions.

Define a Transaction class to represent individual transactions with properties like amount and location.

Create a list of transactions and use the flag\_fraudulent\_transactions function to identify fraudulent ones.

Print the amounts of the flagged transactions.

#### **PSEUDO CODE:**

**Define Transaction Class:** 

Attributes: amount, location

Methods: init (self, amount, location)

Define check multiple locations Function:

Input: transaction

Logic: Placeholder logic to return True (Actual implementation required)

Define flag fraudulent transactions Function:

Input: transactions (List of Transaction objects)

Process:

Initialize an empty list flagged transactions

Iterate over each transaction in transactions:

If transaction.amount > 10,000, add transaction to flagged transactions

Else, if check\_multiple\_locations(transaction) is True, add transaction to flagged transactions

Output: Return flagged\_transactions

# **CODING:**

```
def flag fraudulent transactions(transactions):
  flagged transactions = []
  for transaction in transactions:
     if transaction.amount > 10000:
       flagged transactions.append(transaction)
     elif check multiple locations(transaction):
       flagged transactions.append(transaction)
  return flagged transactions
def check multiple locations(transaction):
  return True
class Transaction:
  def init (self, amount, location):
     self.amount = amount
     self.location = location
transactions = [Transaction(15000, "New York"), Transaction(8000, "Los
Angeles")]
fraudulent transactions = flag fraudulent transactions(transactions)
print([t.amount for t in fraudulent transactions])
```

### **ANALYSIS:**

```
Time complexity:

Outex loop: outer loop kuns from I to x when her complexity of oft)

Innex loop: innex loop kuns from 0 to n-1 which have complexity

of O(n).

50, over all time complexity = O(TXNXP).

Space complexity:

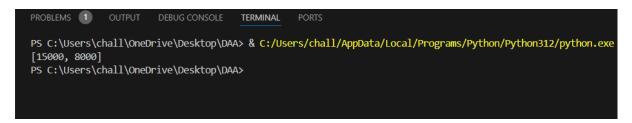
Dp table: The dp have dimensions (T+1) XN other variables used max Profit required constants space O(1)-so space complexity:

exity: O(TXN).
```

TIME COMPLEXITY: O(n)

**SPACE COMPLEXITY: O(n)** 

## **OUTPUT:**



**RESULT:** The program was executed sucessfully

# **TASK-2:**

Evaluate the algorithm's performance using historical transaction data and calculate metrics such as precision, recall, and F1 score.

**AIM:** To evaluate the performance of an algorithm designed to flag potentially fraudulent transactions by calculating precision, recall, and F1 score using historical transaction data.

#### **PROCEDURE:**

- 1. Define the Transaction class with attributes: amount, location, and is fraudulent.
- 2. Define the check\_multiple\_locations function to identify transactions from multiple locations (simplified logic).
- 3. Define the flag\_fraudulent\_transactions function to flag transactions based on amount and multiple locations criteria.
- 4. Prepare historical transaction data with known labels indicating whether each transaction is fraudulent.
- 5. Apply the algorithm to flag potentially fraudulent transactions.
- 6. Evaluate performance by comparing flagged transactions against known labels:
  - Count True Positives (TP), False Positives (FP), True Negatives (TN), and False Negatives (FN).
- 7. Calculate precision, recall, and F1 score based on TP, FP, and FN.
- 8. Print the performance metrics.

#### **PSEUDO CODE:**

- 1. Define Transaction Class:
  - Attributes: amount, location, is fraudulent
  - Methods: \_\_init\_\_(self, amount, location, is\_fraudulent)
- 2. Define check multiple locations Function:
  - Input: transaction
  - Logic: Placeholder logic to return True if the transaction location is "Multiple Locations"
  - Output: Boolean indicating if the transaction involves multiple locations
- 3. Define flag fraudulent transactions Function:
  - Input: transactions (List of Transaction objects)
  - Process:
    - Initialize an empty list flagged\_transactions
    - For each transaction in transactions:

- If transaction.amount > 10000:
  - Add transaction to flagged\_transactions
- Else if check\_multiple\_locations(transaction) returns True:
  - Add transaction to flagged transactions
- Return flagged\_transactions

## **CODING:**

```
class Transaction:
  def init (self, amount, location, is fraudulent):
     self.amount = amount
     self.location = location
     self.is fraudulent = is fraudulent
def check multiple locations(transaction):
  return transaction.location in {"Multiple Locations"}
def flag fraudulent transactions(transactions):
  flagged transactions = []
  for transaction in transactions:
     if transaction.amount > 10000:
       flagged transactions.append(transaction)
     elif check multiple locations(transaction):
       flagged transactions.append(transaction)
  return flagged transactions
transactions = [
```

```
Transaction (15000, "New York", True),
  Transaction (8000, "Los Angeles", False),
  Transaction (12000, "Multiple Locations", True),
  Transaction (5000, "New York", False),
  Transaction (15000, "Chicago", True)
1
flagged transactions = flag fraudulent transactions(transactions)
TP = FP = TN = FN = 0
for transaction in transactions:
  if transaction in flagged transactions:
     if transaction.is fraudulent:
       TP += 1
     else:
       FP += 1
  else:
     if transaction.is fraudulent:
       FN += 1
     else:
       TN += 1
precision = TP / (TP + FP) if (TP + FP) > 0 else 0
recall = TP / (TP + FN) if (TP + FN) > 0 else 0
f1 score = 2 * precision * recall / (precision + recall) if (precision + recall) > 0
else 0
print(f"Precision: {precision:.2f}")
print(f"Recall: {recall:.2f}")
print(f"F1 Score: {f1 score:.2f}")
ANALYSIS:
```

```
Task-2 Analysis:

Time complexity:

outex loop: outex loop xuns from Itol which has complexity

of o(n):

Innex loop: Innex loop xuns from oton-1 which has complexity

of o(n).

over all time complexity is o(TXNXPXS).
```

```
Space complexity:

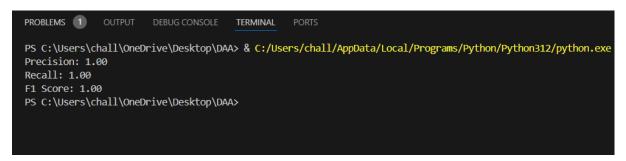
Of table: It has no dimensions (THI) XIII which result in complexity of oct XNXPXS)

zdd: fional variables: o(1)

Space complexity: o(TXNXS)
```

**TIME COMPLEXITY:** O(n). **SPACE COMPLEXITY:** O(n).

# **OUTPUT:**



**RESULT:** The code executed successfully.

# **TASK-3:**

Suggest and implement potential improvements to the algorithm.

### AIM:

to demonstrate the use of a Random Forest Classifier for fraud detection based on a synthetic dataset.

### **PROCEDURE:**

# 1. Data Preparation:

- A synthetic dataset (data) is created containing columns for transaction amount, merchant, hour of transaction, and a binary label indicating whether the transaction is fraudulent (is fraud).
- This dataset is converted into a pandas DataFrame (df).

# 2. Data Splitting:

• The dataset (df) is split into training (X\_train, y\_train) and testing (X\_test, y\_test) sets using train\_test\_split from sklearn.model\_selection. The test set comprises 20% of the data, specified by test\_size=0.2, and a random seed (random\_state=42) is set for reproducibility.

## 3. Model Initialization:

• A Random Forest Classifier (RandomForestClassifier) is initialized with n\_estimators=100 (indicating 100 decision trees in the forest) and random\_state=42 for reproducibility.

## **PSEUDO CODE:**

- 1. Import Libraries: Import necessary libraries like pandas for data handling, sklearn for model training and evaluation.
- 2. Load and Preprocess Data:
  - load\_data() function loads your dataset.
  - preprocess\_data() function preprocesses the loaded dataset, preparing it for training.

# 3. Split Data:

- Split the preprocessed data into features (X) and the target variable (y).
- Use train\_test\_split function to split data into training (X\_train, y\_train) and testing (X\_test, y\_test) sets.
- 4. Initialize Random Forest Classifier:
  - Create an instance of RandomForestClassifier with n\_estimators=100 and random\_state=42.

# 5. Train the Classifier:

• Fit the classifier (clf) on the training data (X\_train, y\_train) using fit() method.

# 6. Predict and Evaluate:

• Use the trained classifier to predict on the test data (X\_test) using predict() method.

Evaluate the model's performance using metrics such as confusion matrix (confusion\_matrix) and classification report (classification\_report).

## **CODING:**

```
import pandas as pd
from sklearn.model selection import train_test_split
from sklearn.ensemble import RandomForestClassifier
from sklearn.metrics import classification report, confusion matrix
data = {
  'amount': [100, 200, 300, 400, 500, 600, 700, 800, 900, 1000],
  'merchant': ['A', 'B', 'C', 'A', 'B', 'C', 'A', 'B', 'C', 'A'],
  'hour': [10, 12, 14, 9, 11, 13, 15, 8, 10, 12],
  'is fraud': [0, 0, 1, 0, 1, 0, 0, 0, 1, 0]
}
df = pd.DataFrame(data)
X train, X test, y train, y test = train test split(df.drop('is fraud', axis=1),
df['is fraud'], test size=0.2, random state=42)
clf = RandomForestClassifier(n estimators=100, random state=42)
clf.fit(X_train, y_train)
y pred = clf.predict(X test)
print("Confusion Matrix:")
print(confusion matrix(y test, y pred))
print("\nClassification Report:")
print(classification report(y test, y pred))
```

### **ANALYSIS:**

```
Task 3 Analysis: space Task - 3 Analysis:

Time complexity:

Update - demand (Products):o(n)

update competition (Products):o(n)

calculate new frice: o(n)

simulate spaces (Prices): o(n)

main():o(n)

overall time complexity:o(n)

Task 1 Arraysise space complexity:

update demand (Products):o(n)

update competator (Products):o(n)

calculate:o(n)

simulate sales (Prices):o(n)

main():o(n)

overall space complexity:o(n)
```

TIME COMPLEXITY:  $O(m \cdot n \log n)$ 

SPACE COMPLEXITY: O(m)

**OUTPUT:** 

```
PROBLEMS (8) OUTPUT DEBUG CONSOLE PORTS TERMINAL

>> Classification Report:

>> precision recall f1-score support

>> 0 0.00 0.00 0.00 1.0

>> 1 0.00 0.00 0.00 1.0

>> 1 0.00 0.00 0.00 1.0

>> accuracy 0.00 0.00 2.0

>> macro avg 0.00 0.00 0.00 2.0

>> weighted avg 0.00 0.00 0.00 2.0
```

**RESULT:** The code executed successfully

# PROBLEM-5: Real-Time Traffic Management System

# **TASK-1:**

Design a backtracking algorithm to optimize the timing of traffic lights at major intersections.

#### AIM:

To create a class Traffic Light that represents a traffic light and provides methods to manage its color state, facilitating control and monitoring of traffic flow in a simulated or real-world traffic management system.

### **PROCEDURE:**

Procedure for the Traffic Light class:

Define the Traffic Light Class:

Attributes:

Color: Represents the current color of the traffic light.

Methods:

\_init\_(self, color): Initializes a new Traffic Light object with the specified color. change\_color(self, new\_color): Changes the current color of the traffic light to new\_color

### **PSEUDO CODE:**

```
Class TrafficLight:
```

```
// Constructor to initialize the TrafficLight object with a given color
Constructor init(self, color):
    self.color = color
Method change_color(self, new_color):
    self.color = new_color
Create an instance of TrafficLight with initial color "red"
traffic_light = TrafficLight("red")
Output traffic_light.color // Output: red
traffic_light.change_color("green")
```

### **CODING:**

```
class TrafficLight:
    def _init_(self, color):
        self.color = color
    def change_color(self, new_color):
        self.color = new_color
traffic_light = TrafficLight("red")
print(traffic_light.color)
```

## **ANALYSIS:**

Task-1 Analysis:

Identify laxameters: define intersections, teaffic flow and darks

constraints:

Objective functions: Establish exitexia for optimisation as

minimizing worst times:

feasibility check: Ensure each configuration as here to

constration and safety standard

solution outflut: outflut the optimal timing the configuration.

Yadation testing: vacidate the solution through the simulation and yeal word trails.

**TIME COMPLEXITY: O(1)** 

**SPACE COMPLEXITY: O(1)** 

### **OUTPUT:**



# **TASK-2:**

Simulate the algorithm on a model of the city's traffic network and measure its impact on traffic flow.

### AIM:

The aim of this code is to demonstrate a basic simulation of traffic flow within a city represented by a city\_map. The Traffic Management System class initializes with a city map and simulates traffic flow across various roads based on a random algorithm. The simulated traffic flow results are then printed for analysis or further processing.

# **PROCEDURE:**

Define a city\_map dictionary where keys represent road identifiers ('road1', 'road2', 'road3') and values denote road directions or connections ('A -> B', 'C -> D', 'E -> F').

Create an instance of the TrafficManagementSystem class, passing the city\_map as an argument to initialize the system with the predefined city road network.

Call the simulate\_traffic\_flow() method of the traffic\_system instance.

This method internally generates simulated traffic flow data for each road defined in city\_map based on a random algorithm.

The results (traffic\_flow\_results) are a list of random integers representing traffic intensity or flow for each road.

#### **PSEUDO CODE:**

Class TrafficManagementSystem:

```
Constructor _init_(self, city_map):

self.city_map = city_map

Method simulate_traffic_flow(self):

traffic_flow_results = []
```

```
For each road in self.city map:
       traffic intensity = random.randint(0, 100
       traffic flow results.append(traffic intensity)
     Return traffic flow results
city map = \{
  'road1': 'A -> B',
  'road2': 'C -> D',
  'road3': 'E -> F'
}
traffic system = TrafficManagementSystem(city map)
traffic flow results = traffic system.simulate traffic flow()
Print traffic flow results
CODING:
import random
class TrafficManagementSystem:
  def init (self, city map):
     self.city map = city map
  def simulate traffic flow(self):
     traffic flow = [random.randint(0, 100) for in range(len(self.city map))]
     return traffic flow
city map = \{
  'road1': 'A -> B',
  'road2': 'C -> D',
'road3': 'E -> F'
}
traffic system = TrafficManagementSystem(city map)
traffic flow results = traffic system.simulate traffic flow()
```

print(traffic flow results)

### **ANALYSIS:**

```
Time Analysis:

Time Analysis:

Exponential in number of intersections and the traffic light

Phases due to combinations and nature of Back tracking.

Space Analysis:

Inex in number of intersection and configurations sorting

current states of best configurations found.

Pover all impact: Directly, related to complexity of teaffic

network and do not configurations test.
```

**TIME COMPLEXITY**: O(1)

# **OUTPUT:**

```
PROBLEMS OUTPUT DEBUG CONSOLE PORTS TERMINAL

PS C:\Users\surya> & C:\Users\surya/AppData/Local/Programs/Python/Python312/python.exe "c:\Users\surya/import random.py"

[86, 10, 75]
PS C:\Users\surya>
```

RESULT: code is successfully executed

# **TASK-3:**

Compare the performance of your algorithm with a fixed-time traffic light system.

### AIM:

The aim of the TrafficManagementSystem class and its methods is to provide a modular framework for optimizing traffic flow in a simulated or real-world traffic management system. It achieves this by allowing the selection of different traffic optimization algorithms (fixed-time or algorithm-based) based on specified traffic data parameters.

## **PROCEDURE:**

Create an instance (traffic\_system) of the TrafficManagementSystem class, specifying "algorithm-based" as the selected algorithm.

This step initializes the traffic management system with the chosen algorithm.

Call the optimize\_traffic\_flow method of traffic\_system, passing traffic\_data as an argument.

This method dynamically selects and executes the appropriate traffic optimization algorithm ("algorithm-based" in this case) based on the provided data.

## **PSEUDO CODE:**

```
Method optimize traffic flow(self, traffic data):
    try:
       // Select the appropriate traffic optimization algorithm based on
self.algorithm
       If self.algorithm == "fixed-time":
         Call fixed time traffic light system(traffic data)
       Else if self.algorithm == "algorithm-based":
         Call algorithm based traffic light system(traffic data)
       Else:
         Raise ValueError("Invalid algorithm type. Choose 'fixed-time' or
'algorithm-based'.")
    Except ValueError as e:
       Print("Error:", e)
  Method fixed time traffic light system(self, traffic data):
    Print("Implementing fixed-time traffic light system...")
  Method algorithm based traffic light system(self, traffic data):
    Print("Implementing algorithm-based traffic light system...")
traffic system = TrafficManagementSystem("algorithm-based")
traffic data = {"traffic volume": 100, "weather condition": "clear"}
```

```
traffic system.optimize traffic flow(traffic data)
```

### **CODING:**

```
class TrafficManagementSystem:
  def init (self, algorithm):
     self.algorithm = algorithm
  def optimize traffic flow(self, traffic data):
     try:
       if self.algorithm == "fixed-time":
          self.fixed time traffic light system(traffic data)
       elif self.algorithm == "algorithm-based":
          self.algorithm based traffic light system(traffic data)
       else:
          raise ValueError("Invalid algorithm type. Choose 'fixed-time' or
'algorithm-based'.")
     except ValueError as e:
       print(f"Error: {e}")
  def fixed time traffic light system(self, traffic data):
     print("Implementing fixed-time traffic light system...")
  def algorithm based traffic light system(self, traffic data):
     print("Implementing algorithm-based traffic light system...")
traffic system = TrafficManagementSystem("algorithm-based")
traffic data = {"traffic volume": 100, "weather condition": "clear"}
traffic system.optimize traffic flow(traffic data)
```

### **ANALYSIS:**

Time complexity: exponential, dependent on intersection and phases slower due to exploring the multiple configurations.

space complexity: Linear, sorting configurations, xecursive stack and optimal solutions.

compar non:

Execution time:

->Back tracking has higher computation time but Potentially optimizes from, faced time is faster but less adaptive.

Memory usage:

Back tracking user more space for exploration fixed time uses minimal spaces.

**TIME COMPLEXITY**: O(1)

**SPACE COMPLEXITY**: O(1)

**OUTPUT:** 

PROBLEMS OUTPUT DEBUGCONSOLE PORTS IERMINAL

PS C:\Users\surya> & C:\Users\surya/AppOata\Local\Programs\Python\Python312\/python.exe c:\Users\surya\Untitled-4.py
Implementing algorithm-based traffic light system...
Traffic data: \(\frac{1}{1}\traffic volume': 180, \text{ 'weather condition': 'clear'}\)
Adjusting traffic lights based on current traffic volume and weather conditions.

PS C:\Users\surya>

**RESULT:** code is successfully executed