

CH 4 : Queue

- It is a linear DS in which the elements that is inserted first is the one that will come out first.
- It works in FIFO (first in first out) manner.

* Types Of Queue

→ Simple Queue

→ Circular Queue

→ DeQueue [Double Ended Queue]

→ Priority Queue

→ Queue can be implemented using array & linked list.

* Basic Operation On Queue

- enqueue () - for insertion
- dequeue () - for deletion

QINSERT (Q, F, R, N, Y)

Name of Front Rear Size of Element to be
Queue Queue inserted

S. 1 [Check for Overflow]

if $R \geq N$

then write ("Overflow")

Exit ()

S. 2 [Increment Rear pointer]

$R \leftarrow R + 1$

S. 3 [Insert Element]

$Q[R] \leftarrow Y$

S. 4 [Is Front Pointer Properly Set?]

if $F = 0$

then $F \leftarrow 1$

Return

QDELETE (Q, F, R)

S.1 [Underflow?] not - () underflow
not - () underflow

If $F = 0$

then write ("Underflow")

Return (0)

S.2 [Deleted Element]

by $y \leftarrow Q[F]$

S.3 [Queue Empty?]

("if $F = R$ ") since next

then $F \leftarrow R \leftarrow 0$

Else

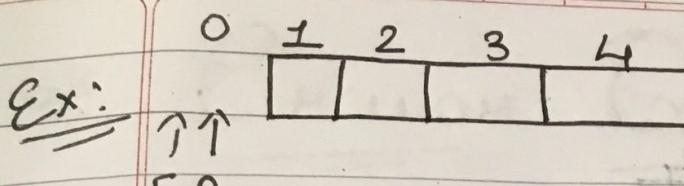
$F \leftarrow F + 1$

S.4 [Return Element]

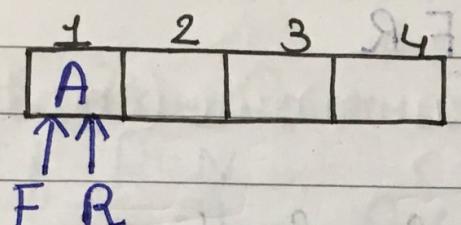
Return (Y)

$K \rightarrow [R]$

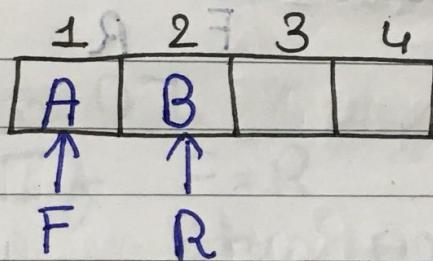
$O = 7$ fi



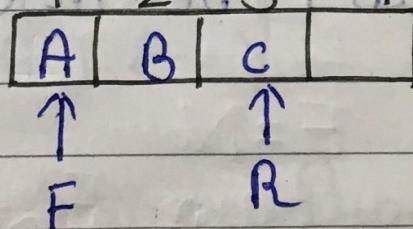
(i) Insert A



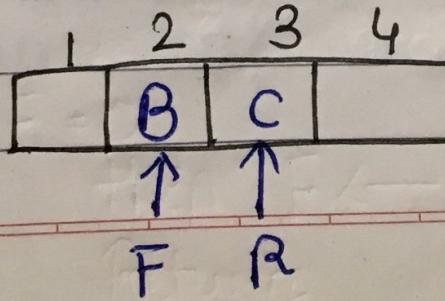
(ii) Insert B



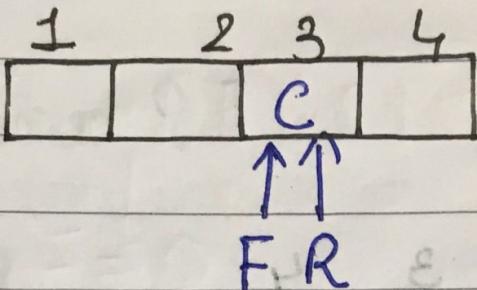
(iii) Insert C



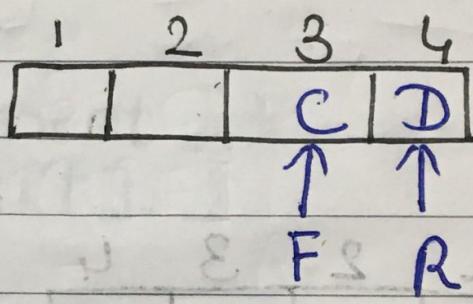
(iv) Delete



(v) Delete

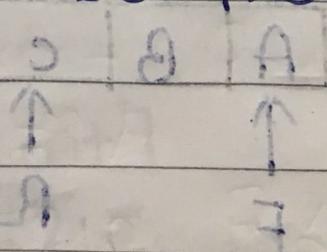


(vi) Insert D



(vii) Insert E

$R \geq N$, therefore
it will be overflow
∴ E can't be inserted



* Circular Queue

CQINSERT (Q, F, R, N, Y)

5.1 [Reset Rear Pointer ?]

If $R = N$ (0) ~~return~~
 then $R \leftarrow 1$

Else

$R \leftarrow R + 1$ (1) ~~return~~

5.2 [Overflow ?]

If $F = R$
 then write ("Overflow")

$0 \rightarrow R$ (0) ~~return~~

(1) ~~return~~ ~~sub~~

5.3 [Insert Element]

$Q[R] \leftarrow Y$

$N = T + 1$

5.4 [Is Front Pointer Properly Set ?]

if $F = 0$
 then $F \leftarrow 1$

return (1) ~~return~~

CQDELETE (Q, F, R, N)

S.1 [Underflow ?]

If $F = 0$

then write ('underflow').

return (0)

$\leftarrow \rightarrow R$ null

S.2 [Delete Element]

$y \leftarrow Q[F] \rightarrow R$

S.3 [Queue Empty ?]

("empty" If $F = R$)

then $F \leftarrow R + 1 - 0$

else return (y)

S.4 [Increment front pointer]

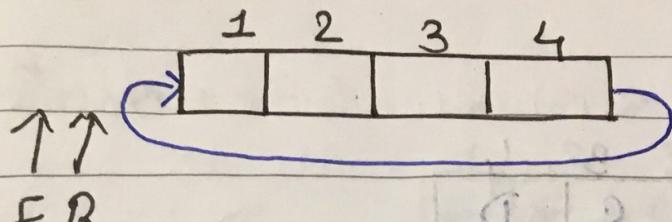
If $F = N$

then $F \leftarrow 1$

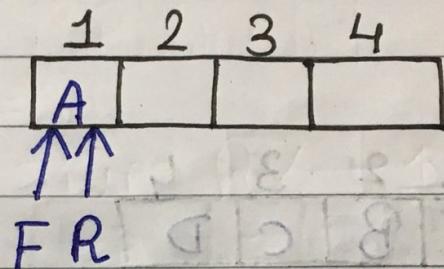
Else

$\leftarrow F \leftarrow F + 1$

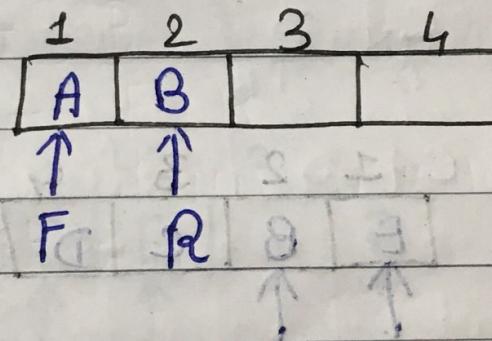
return (y) : move

Ex:

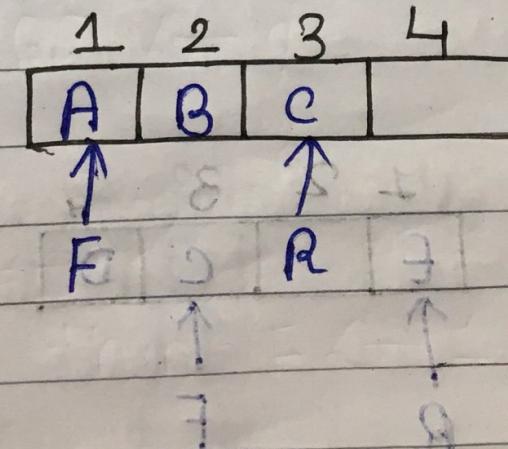
(i) Insert A



(ii) Insert B



(iii) Insert C



(iv) Insert D

1	2	3	4
A	B	C	D

↑ ↑
F R

(v) Delete

1	2	3	4
	B	C	D

↑ ↑
F R

(vi) Insert E

1	2	3	4
E	B	C	D

↑ ↑
R F

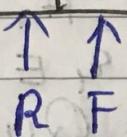
(vii) Delete

1	2	3	4
E	R	C	D

↑ ↑
R F

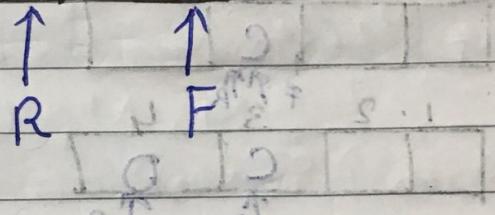
(viii) Insert F

1	2	3	4
E	F	C	D



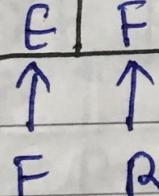
(ix) Delete

1	2	3	4	5	6	7	8	9	A
E	F		D						



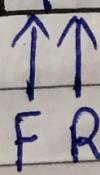
(x) Delete

1	2	3	4
E	F		



(xi) Delete

1	2	3	4
	F		



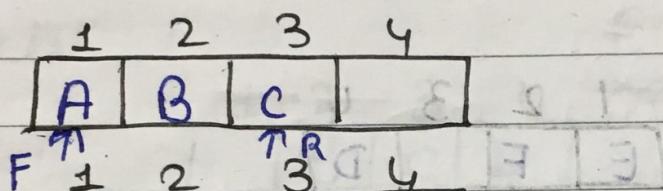
(xii) Delete

1	2	3	4

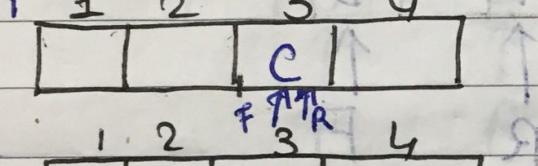
$$F = R = 0$$

Ex: Consider an Ex where Size of Queue
 Initially it is empty. It is req. to
 insert symbols A, B & C. Delete A & B
 And Insert D & E. Show the trace
 Of Content of Queue

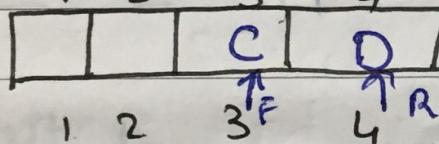
1.



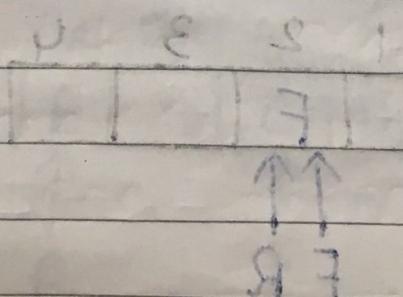
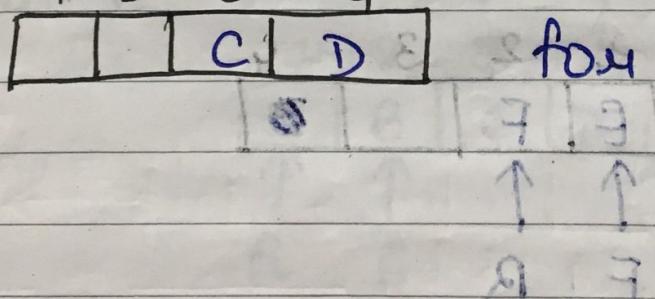
2.



3.



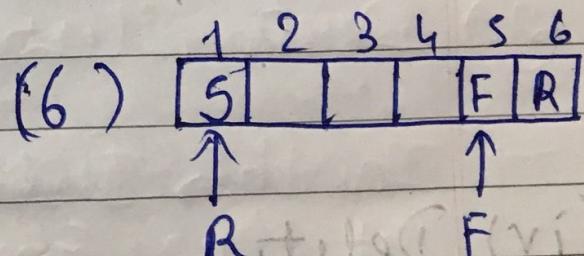
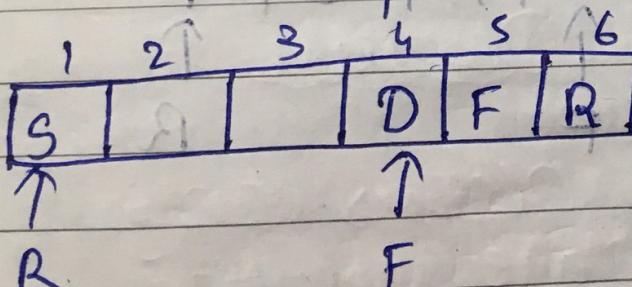
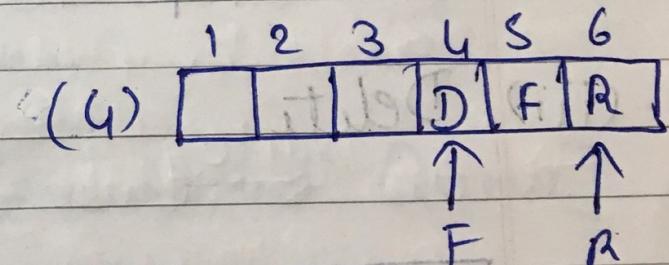
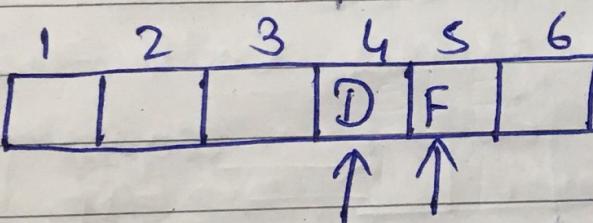
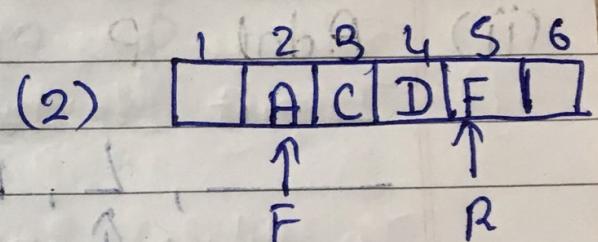
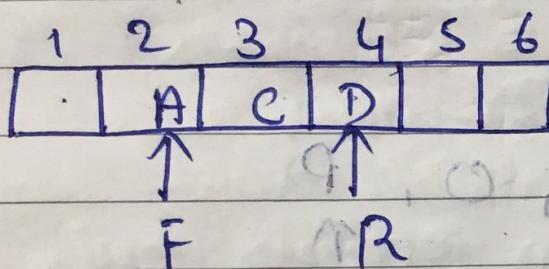
4.

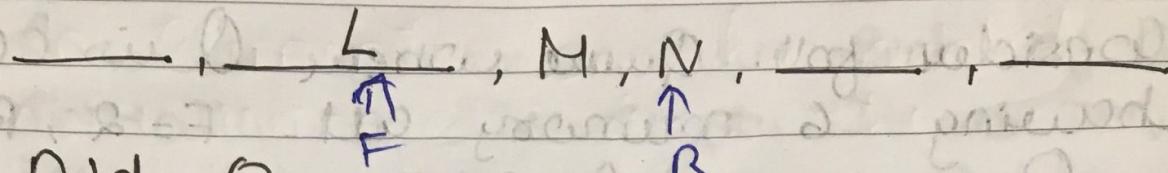


Ex: Consider foll. Queue where Q is CQ having 6 memory cell. $F = 2$, $R = 4$

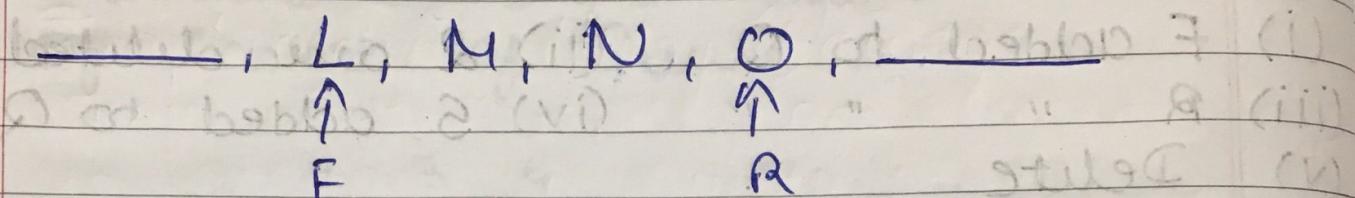
Queue: , A, C, D, ,

- (i) F added to Q
- (ii) 2 are deleted
- (iii) R " "
- (iv) S added to Q
- (v) Delete

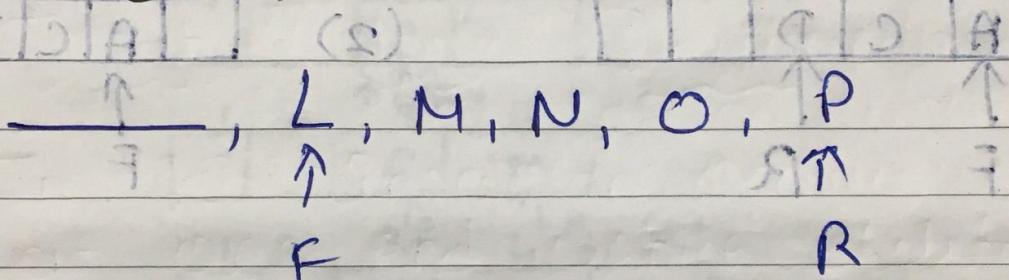


Ex:

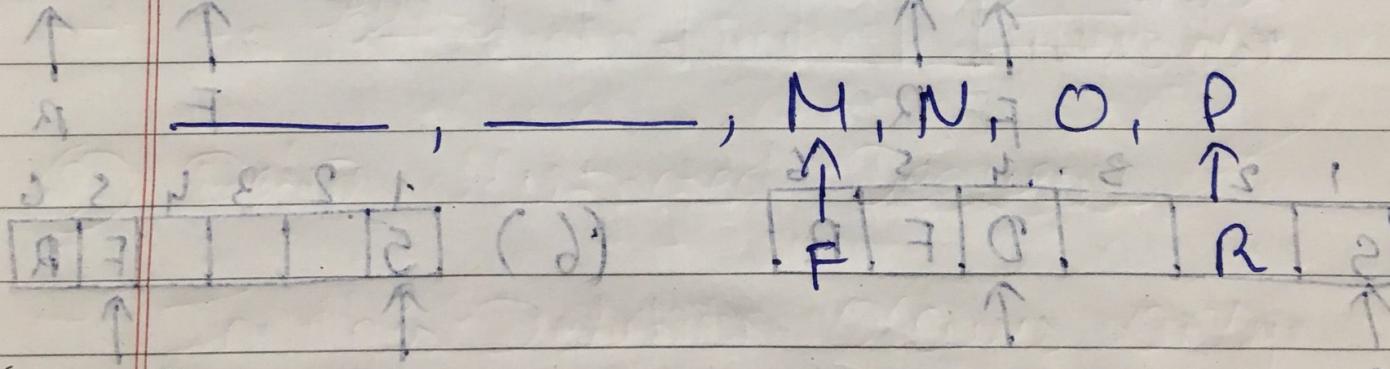
(i) Add O



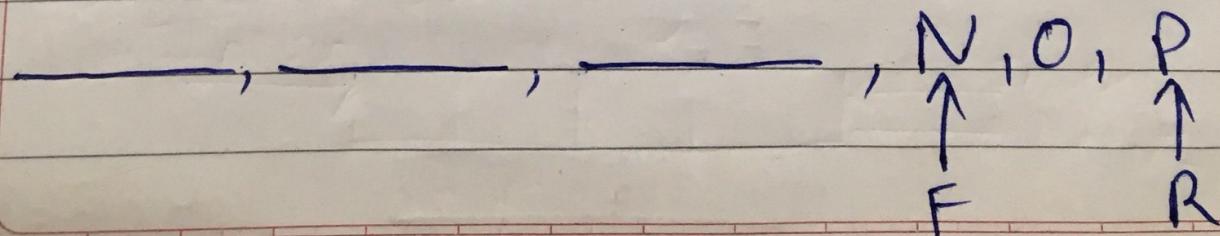
a(iiP) Add P



a(iii) Delete



(iv) Delete



(vi) Add Q, R, S and elements on next page beginning no

Q, R, S, N, O, P 2nd line
 B ↑ ↑
 P R parallel between point on line below

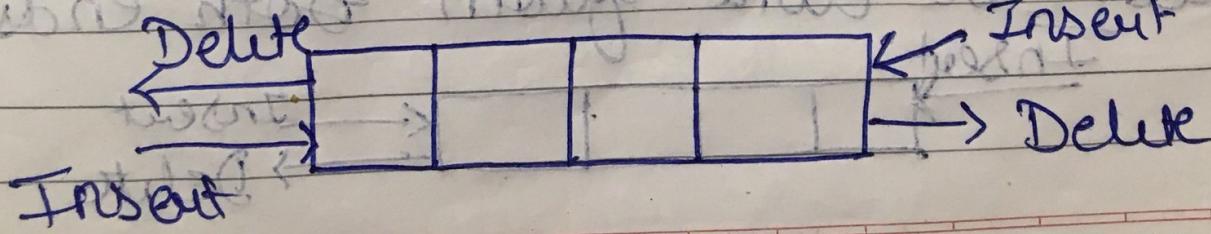
(vii) Delete

Q, R, S, —, O, P
 P R parallel between point on line below F

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* DeQueue [Double Ended Queue]

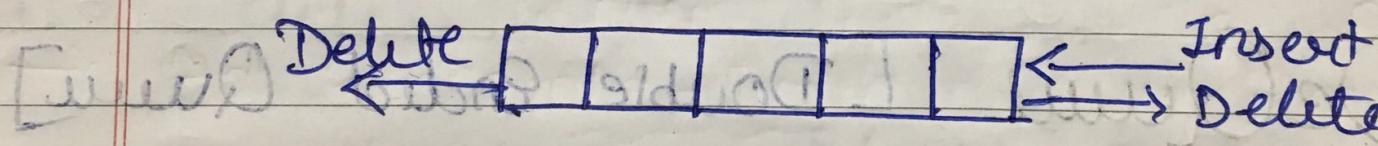
DeQueue is a list in which elements can be inserted and deleted at either end.



- (iv) No Elements Can be added or removed from the middle.
- (v) Implemented Using Doubly linked list.
- (vi) Two Variation Of deque

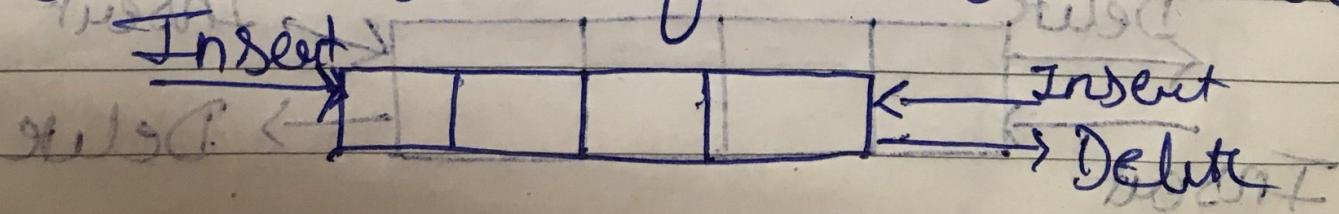
① Input Restricted Dequeue

In this insertion occurs only at one end while deletion can be done from both ends.



② Delete Restricted Dequeue

In this Deletion occurs only at One End while insertion can be done from both ends.



Ex: —, A, B, C, D, E, —, —, —, —

(i) Add f on left

E, A, B, C, D, E, —, f, —

(ii) Add G on right

E, A, B, C, D, E, G, —, —

(iii) Add H on right

F, A, B, C, D, E, G, H, —

(iv) Delete 2 alphabets from left

—, —, B, C, D, E, F, G, H, —

(v) Add I on right

—, —, B, C, D, E, F, G, H, I, —

* Priority Queue

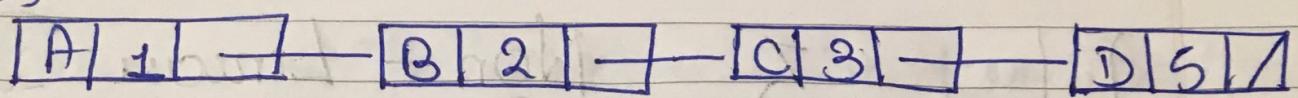
(a) Priority Queue is Data Structure in which each element is assigned a priority.

(b) This priority of elements will be used to determine the order in which the elements will be processed.

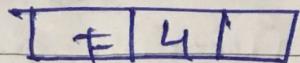
(c) Rules for Priority Queue

- An element with higher priority is processed before an element with lower priority.
- Two Elements with same priority, have proceeded based on FCFS.

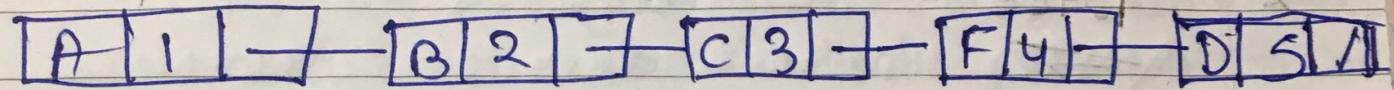
Initially



new Process



now



* Application Of Queue

- ① Job Scheduling
- ② Categorizing Data
- ④ Josephus Problem