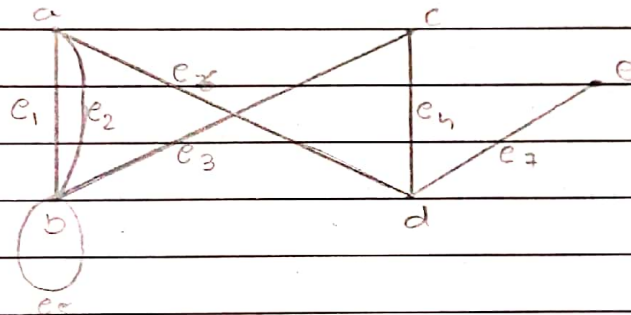


# TUTORIAL #3

Date: \_\_\_\_\_

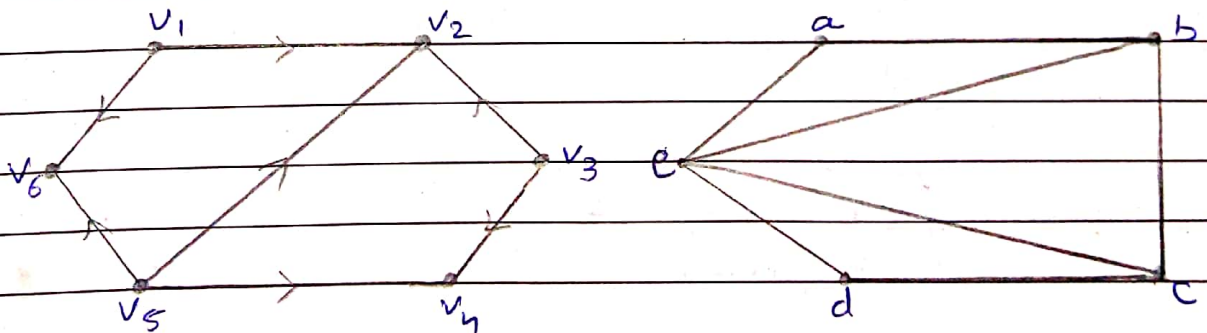
## GRAPHS

Q.1] Draw the undirected graph  $G = (V, E)$ , where  $V = \{a, b, c, d, e\}$  and  $E = \{e_1, e_2, e_3, e_4, e_5, e_6, e_7\}$  and its incidence relation given as:  
 $e_1 = (a, b)$ ,  $e_2 = (a, b)$ ,  $e_3 = (b, c)$ ,  $e_4 = (c, d)$ ,  $e_5 = (b, b)$ ,  
 $e_6 = (a, d)$ ,  $e_7 = (e, d)$



(Undirected graph)

Q.2] Consider the following graphs: Determine the degree of each node and verify handshaking lemma.



- As we can see it is directed graph  
 so,

Date: 

(1) In-degree

(2) out degree

Now, In-degree

$$v_1 = 0$$

$$v_2 = 3$$

$$v_3 = 0$$

$$v_4 = 2$$

$$v_5 = 0$$

$$v_6 = 2$$

$$\underline{7}$$

Out-degree

$$v_1 = 2$$

$$v_2 = 0$$

$$v_3 = 2$$

$$v_4 = 0$$

$$v_5 = 3$$

$$v_6 = 0$$

$$\underline{7}$$

$$\begin{aligned} \text{Total number of degree} &= \text{In-degree} + \text{out-degree} \\ &= 7 + 7 \\ &= 14 \end{aligned}$$

Now, Handshaking lemma

$$\therefore \sum_{i=1}^n \deg(v) = 2 |E|$$

$$\therefore \sum_{i=1}^6 \deg(v) = 2 |7|$$

$$\therefore \boxed{\deg(v) = 14}$$

Now, for undirected graph

Date:

degree of  $a = 2$  $b = 3$  $c = 3$  $d = 2$  $e = 4$ 

Now,

$$\begin{aligned}\text{Total Degree} &= a + b + c + d + e \\ &= 2 + 3 + 3 + 2 + 4 \\ &= 14\end{aligned}$$

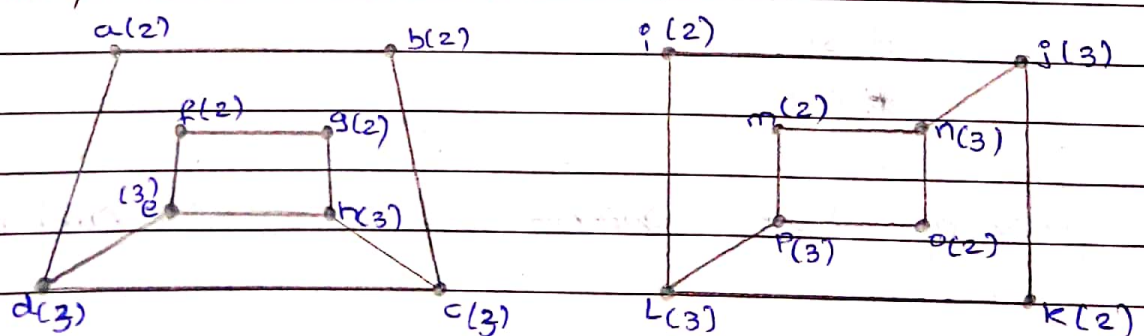
$$\text{Now, } \sum \deg(v) = 2 |E|$$

$$\Rightarrow \deg(v) = 14$$

$$\text{Hence, } 2 |E| = 14$$

$\therefore$  The given graph is an Handshaking lemma.

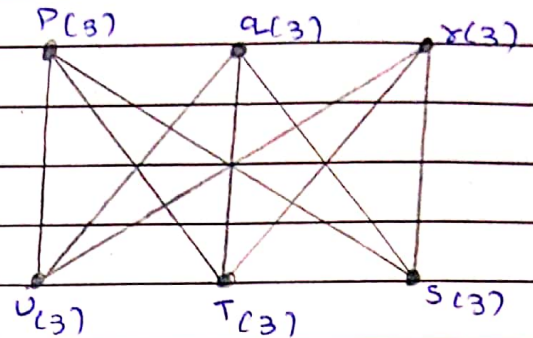
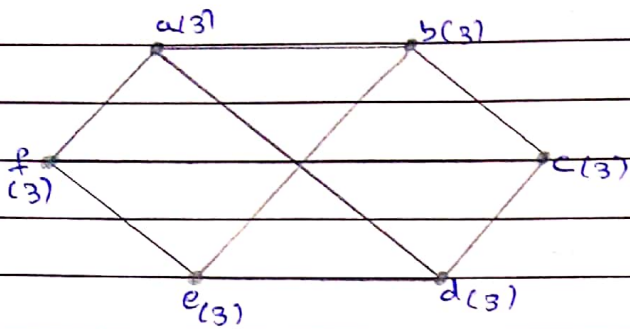
Q.3] check whether the following graphs are isomorphic or not.

 $\Rightarrow (A)$ 

Here,  $a \rightarrow j$  &  $b \nrightarrow j$  so graph is not ISOMORPHIC



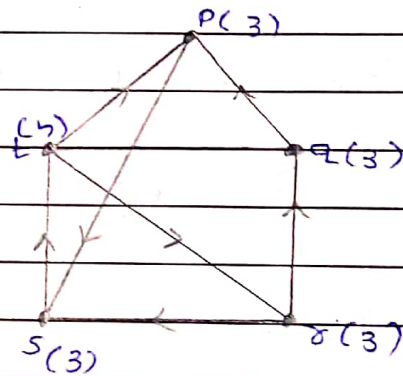
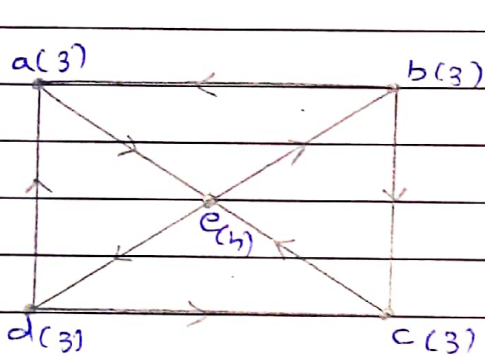
(B)



Here,  $a \rightarrow P$ ,  $b \rightarrow Q$ ,  $c \rightarrow R$ ,  $d \rightarrow S$ ,  $e \rightarrow T$ ,  $f \rightarrow U$

Hence, this graph is ISOMORPHIC.

(C)



Here,

$a \rightarrow P$

$b \rightarrow Q$

$c \rightarrow R$

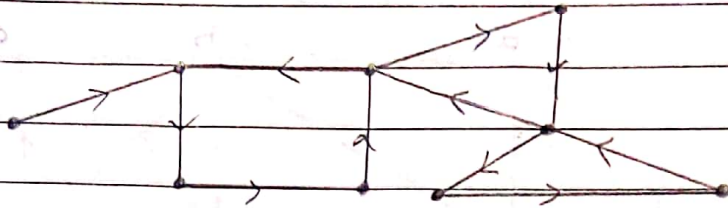
$d \rightarrow S$

$e \rightarrow T$

Hence, This graph is ISOMORPHIC

Q.5] Find converse of a given digraph.

⇒ (a)



(P)

Now, Transpose of P will be  $P^T$



( $P^T$ )

(b)



(Q)

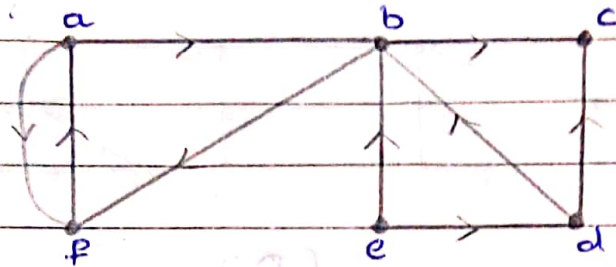
Now, Transpose of Q will be  $Q^T$



( $Q^T$ )



Q.5) Find reachability matrix for the given graph.



$\Rightarrow$  Step-1 Path of length is 1

$$R^1 = \begin{matrix} & \begin{matrix} a & b & c & d & e & f \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Step-2 Path of length is 2

$$R^2 = \begin{matrix} & \begin{matrix} a & b & c & d & e & f \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \end{matrix} & \begin{bmatrix} 0 & 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Step-3 Path of length is 3

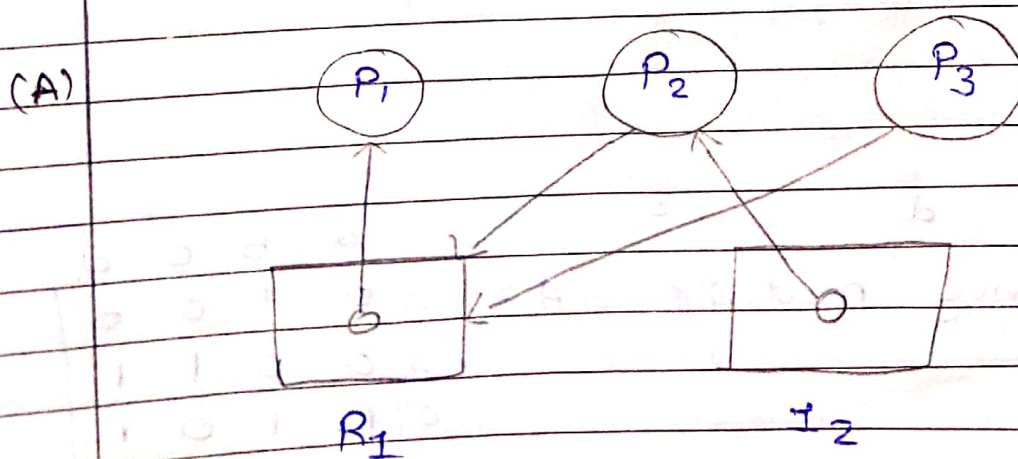
$$R^3 = \begin{matrix} & a & b & c & d & e & f \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \end{matrix} & \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

Now, Reachability matrix is  $R = R^1 + R^2 + R^3$

$$R = \begin{matrix} & a & b & c & d & e & f \\ \begin{matrix} a \\ b \\ c \\ d \\ e \\ f \end{matrix} & \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 2 \\ 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 2 & 0 & 0 & 1 \\ 1 & 2 & 2 & 1 & 0 & 2 \\ 1 & 1 & 1 & 0 & 0 & 0 \end{bmatrix} \end{matrix}$$

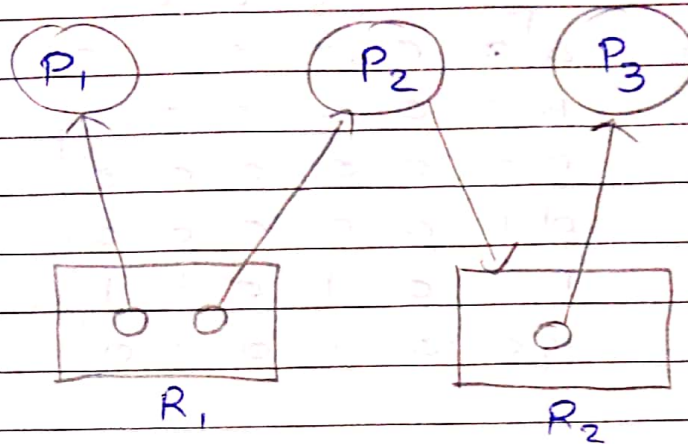
Q.6 Give three examples of resource allocation graph which does not have any deadlock

⇒ The example of resource allocation graph without deadlock.

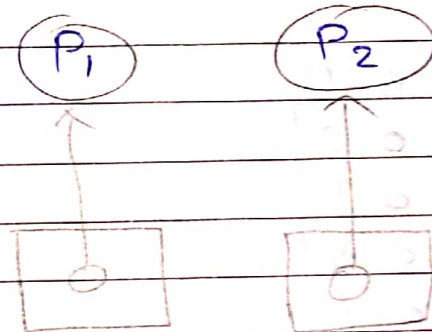




(B)

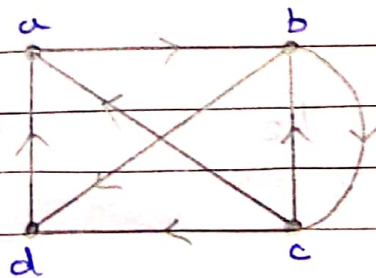


(c)



Hence, these are the three examples of Resource Allocation Graph.

Q.7] Obtain adjacency matrix  $A$  for the digraph in the figure. Find  $A$ ,  $A^2$  without matrix multiplication.



$\Rightarrow$  Adjacency matrix  $A =$

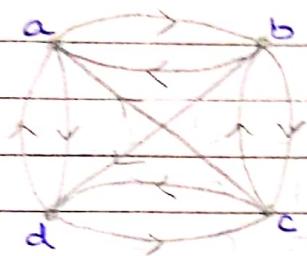
	a	b	c	d
a	0	1	0	0
b	0	0	1	1
c	1	1	0	1
d	1	0	0	0



Now,

$$A^2 = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 0 & 2 & 2 \\ 2 & 0 & 0 & 2 \\ 2 & 2 & 0 & 2 \\ 0 & 2 & 0 & 0 \end{bmatrix} \end{matrix}$$

Q.8] Find path of length 2 from the graph (using matrix)



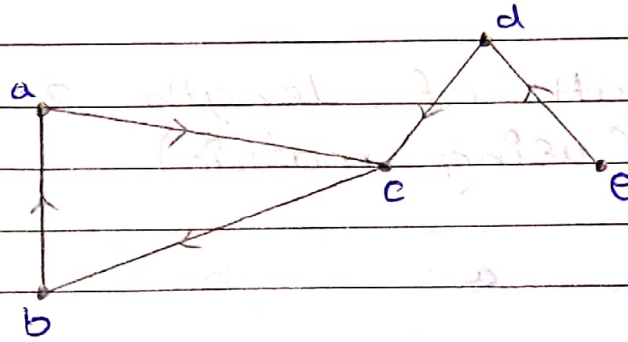
$$\Rightarrow A = \begin{matrix} & \begin{matrix} a & b & c & d \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \end{matrix} & \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

Now, path of length 2 using matrix  
 $A^2 = A \times A$

$$\therefore A^2 = \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix} \times \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{bmatrix}$$

$$\therefore A^2 = \begin{bmatrix} 2 & 0 & 2 & 1 \\ 2 & 2 & 0 & 2 \\ 2 & 1 & 2 & 2 \\ 1 & 2 & 0 & 2 \end{bmatrix}$$

Q.9] Find the shortest path between each pair of vertices for a simple digraph using warshall's algorithm. Produce Path matrix.



$$\Rightarrow A = \begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

$$\text{Now, } D_0 = \begin{matrix} & \begin{matrix} a & b & c & d & e \end{matrix} \\ \begin{matrix} a \\ b \\ c \\ d \\ e \end{matrix} & \begin{bmatrix} 0 & 0 & 1 & \infty & \infty \\ 1 & 0 & \infty & \infty & \infty \\ 0 & 1 & 0 & \infty & \infty \\ 0 & \infty & 1 & 0 & \infty \\ 0 & \infty & \infty & 1 & 0 \end{bmatrix} \end{matrix}$$

- $D_1[b, c] = \min[D_0(b, c), D_0(b, a) + D_0(a, c)]$   
 $D_1[b, c] = \min[\infty, 1 + 1]$   
 $D_1[b, c] = 2$
- $D_1[b, d] = \min[D_0(b, d), D_0(b, a) + D_0(a, d)]$   
 $= \min[\infty, 1 + 0]$   
 $= 1$



$$\bullet D_1[b, e] = \min [D_0(b, e), D_0(b, i) + D_0(i, e)]$$

$$= 1$$

$$\bullet D_1[c, d] = \min [D_0(c, d), D_0(c, i) + D_0(i, d)]$$

$$= 0$$

$$\bullet D_1[c, e] = \min [D_0(c, e), D_0(c, i) + D_0(i, e)]$$

$$= 0$$

$$\bullet D_1[d, b] = \min [D_0(d, b), D_0(d, i) + D_0(i, b)]$$

$$= 0$$

$$\bullet D_1[e, b] = \min [D_0(e, b), D_0(e, i) + D_0(i, b)]$$

$$= 0$$

$$\bullet D_1[e, c] = \min [D_0(e, c), D_0(e, i) + D_0(i, c)]$$

$$= 1$$

$$D_1 =$$

	a	b	c	d	e
a	0	0	1	$\infty$	$\infty$
b	1	0	2	1	1
c	$\infty$	1	0	$\infty$	$\infty$
d	$\infty$	0	1	0	$\infty$
e	$\infty$	0	1	1	0

$$\bullet D_2[a, d] = \min [D_1(a, d), D_1(a, 2) + D_1(2, d)]$$

$$= 1$$

$$\bullet D_2[a, e] = \min [D_1(a, e), D_1(a, 2) + D_1(2, e)]$$

$$= 1$$

$$\bullet D_2[c, a] = \min [D_1(c, a), D_1(c, 2) + D_1(2, a)]$$

$$= 2$$

$$\bullet D_2[c, d] = \min [D_1(c, d), D_1(c, 2) + D_1(2, d)]$$

$$= 2$$

$$\bullet D_2[c, e] = \min [D_1(c, e), D_1(c, 2) + D_1(2, e)]$$

$$= 2$$

$$\bullet D_2[d, a] = \min [D_1(d, a), D_1(d, 2) + D_1(2, a)]$$

$$= 1$$

$$D_2[d, e] = \min [D_1(d, e), D_1(d, 2) + D_1(2, e)]$$

$$= 1$$

$$D_2[e, a] = \min [D_1(e, a), D_1(e, 2) + D_1(2, a)]$$

$$= 0$$

$$D_2 =$$

	a	b	c	d	e
a	0	$\infty$	1	1	1
b	1	0	2	1	1
c	2	1	0	2	2
d	1	$\infty$	1	0	1
e	$\infty$	$\infty$	1	1	0

$$D_3[a, b] = \min [D_2(a, b), D_2(a, 3) + D_2(3, b)]$$

$$= 2$$

$$D_3[d, b] = \min [D_2(d, b), D_2(d, 3) + D_2(3, b)]$$

$$= 2$$

$$D_3[e, b] = \min [D_2(e, b), D_2(e, 3) + D_2(3, b)]$$

$$= 2$$

$$D_4 =$$

	a	b	c	d	e
a	0	2	1	1	1
b	1	0	2	1	1
c	2	1	0	2	2
d	1	2	1	0	1
e	2	2	1	1	0

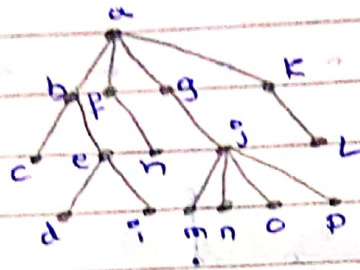
$$D_5 =$$

	a	b	c	d	e
a	0	2	1	1	1
b	1	0	2	1	1
c	2	1	0	2	2
d	1	2	1	0	1
e	2	2	1	1	0

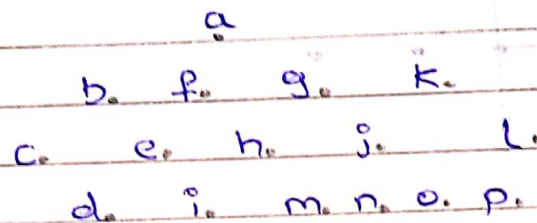
Hence,  
 $D_5$  is the  
 path matrix



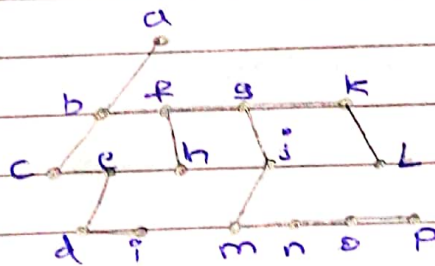
Q.10] Convert following m-ary trees into binary trees.



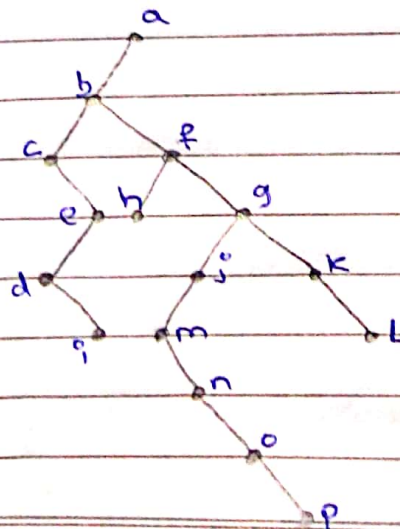
⇒ Step-1 Removing all edges.



Step-2 connecting left side vertices.



Step-3 Representation in proper manner

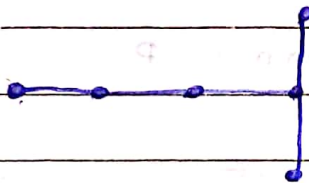


Q.11] Draw three non isomorphic trees with 6 vertices.

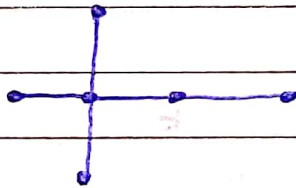
=> (A)



(B)



(C)



=> Hence, these are the required non-isomorphic trees with 6-vertices.