

## Analysis of Algorithms | Set 3 (Asymptotic Notations)

We have discussed [Asymptotic Analysis](#), and [Worst, Average and Best Cases of Algorithms](#). The main idea of asymptotic analysis is to have a measure of efficiency of algorithms that doesn't depend on machine specific constants, and doesn't require algorithms to be implemented and time taken by programs to be compared. Asymptotic notations are mathematical tools to represent time complexity of algorithms for asymptotic analysis. The following 3 asymptotic notations are mostly used to represent time complexity of algorithms.

**1)  $\Theta$  Notation:** The theta notation bounds a functions from above and below, so it defines exact asymptotic behavior. A simple way to get Theta notation of an expression is to drop low order terms and ignore leading constants. For example, consider the following expression.

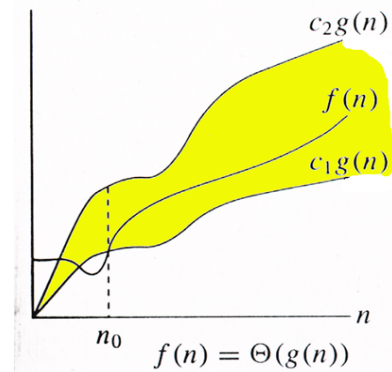
$$3n^3 + 6n^2 + 6000 = \Theta(n^3)$$

Dropping lower order terms is always fine because there will always be a  $n_0$  after which  $\Theta(n^3)$  beats  $\Theta(n^2)$  irrespective of the constants involved.

For a given function  $g(n)$ , we denote  $\Theta(g(n))$  is following set of functions.

$$\Theta(g(n)) = \{f(n): \text{there exist positive constants } c_1, c_2 \text{ and } n_0 \text{ such that } 0 \leq c_1 \cdot g(n) \leq f(n) \leq c_2 \cdot g(n) \text{ for all } n \geq n_0\}$$

The above definition means, if  $f(n)$  is theta of  $g(n)$ , then the value  $f(n)$  is always between  $c_1 \cdot g(n)$  and  $c_2 \cdot g(n)$  for large values of  $n$  ( $n \geq n_0$ ). The definition of theta also requires that  $f(n)$  must be non-negative for values of  $n$  greater than  $n_0$ .



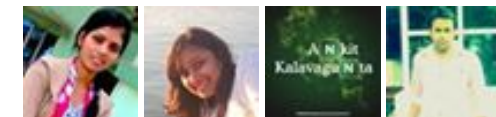
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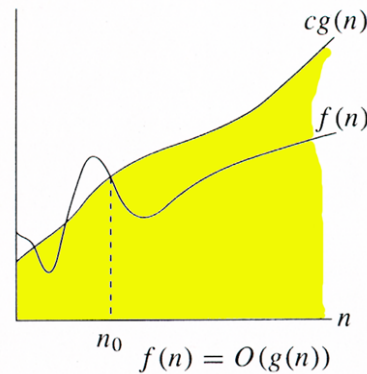
**2) Big O Notation:** The Big O notation defines an upper bound of an algorithm, it bounds a function only from above. For example, consider the case of Insertion Sort. It takes linear time in best case and quadratic time in worst case. We can safely say that the time complexity of Insertion sort is  $O(n^2)$ . Note that  $O(n^2)$  also covers linear time.

If we use  $\Theta$  notation to represent time complexity of Insertion sort, we have to use two statements for best and worst cases:

1. The worst case time complexity of Insertion Sort is  $\Theta(n^2)$ .
2. The best case time complexity of Insertion Sort is  $\Theta(n)$ .

The Big O notation is useful when we only have upper bound on time complexity of an algorithm. Many times we easily find an upper bound by simply looking at the algorithm.

$$O(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq cg(n) \text{ for all } n \geq n_0 \}$$



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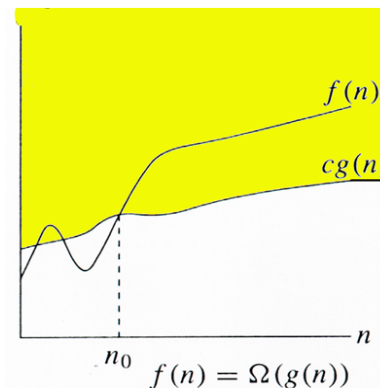
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**3)  $\Omega$  Notation:** Just as Big O notation provides an asymptotic upper bound on a function,  $\Omega$  notation provides an asymptotic lower bound.

$\Omega$  Notation can be useful when we have lower bound on time complexity of an algorithm. As discussed in the previous post, the **best case performance of an algorithm is generally not useful**, the Omega notation is the least used notation among all three.

For a given function  $g(n)$ , we denote by  $\Omega(g(n))$  the set of functions.

$$\Omega(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq cg(n) \leq f(n) \text{ for all } n \geq n_0 \}.$$



Let us consider the same Insertion sort example here. The time complexity of Insertion Sort can be written as  $\Omega(n)$ , but it is not a very useful information about insertion sort, as we are generally interested in worst case and sometimes in average case.

### Exercise:

Which of the following statements is/are valid?

1. Time Complexity of QuickSort is  $\Theta(n^2)$
2. Time Complexity of QuickSort is  $O(n^2)$
3. For any two functions  $f(n)$  and  $g(n)$ , we have  $f(n) = \Theta(g(n))$  if and only if  $f(n) = O(g(n))$  and  $f(n) = \Omega(g(n))$ .
4. Time complexity of all computer algorithms can be written as  $\Omega(1)$

### References:

Lec 1 | MIT (Introduction to Algorithms)

Introduction to Algorithms 3rd Edition by Clifford Stein, Thomas H. Cormen, Charles E. Leiserson, Ronald L. Rivest

This article is contributed by **Abhay Rathi**. Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

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**Marsha Donna** · 4 months ago

so does this mean that  
big O notation corresponds to worst case time complexity,  
big omega notation corresponds to best case time complexity,  
big theta notation corresponds to average case time complexity???

9 ^ | v .



**darubramha** → Marsha Donna · a month ago

Somewhat correct.

big O notation corresponds to worst case time complexity, generally it  
 $T(n)$  [time complexity] =  $O(n)$  then it means,  $T(n)$  can have maximum c

big omega notation corresponds to best case complexity, it is similar to  
case.

705



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But, big theta notation, tells us that the time complexity of the algorithm given function, which means if  $T(n) = \theta(n)$ , the time complexity can be better than  $c_2n$ , where  $c_1, c_2$  are constants. It does not mean the approach towards "=", i.e.  $c_1n \leq T(n) \leq c_2n$

1 ^ | v .



**Ankur Teotia** → Marsha Donna · 2 months ago

i have the same doubt , someone please clarify.  
it seems to be the case but no where i see this explicitly mentioned.

1 ^ | v .



**noob** · 7 months ago

1. Time Complexity of QuickSort is  $(n^2)$  -----> NO
2. Time Complexity of QuickSort is  $O(n^2)$  -----> YES
3. For any two functions  $f(n)$  and  $g(n)$ , we have  $f(n) = (g(n))$  if and only if  $f(n) =$
4. Time complexity of all computer algorithms can be written as (1) -----> YES

8 ^ | v .



**Chaithanya Kanumolu** → noob · 4 months ago

I am confused with 4th point. Since best case for insertion sort is  $\Theta(n^2)$ , write it as  $\Theta(1)$ ?

1 ^ | v .



**ibn** → noob · 5 months ago

Why was question 4 Yes?

^ | v .



**Utkarsh Gupta** → ibn · 3 months ago

If  $f(n) = O(n)$  then it means  $f(n)$  can be less than or equal to order

If  $f(n) = \Theta(n)$  then it means  $f(n)$  is equal to the order of  $n$ .

If  $f(n) = \Omega(n)$  then it means  $f(n)$  is greater than or equal to order

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So if you see  $\Omega(1)$ , it means that this is the smallest lower bound algorithm  $f(n) = \Omega(1)$ . But it can be even more precise that is  $f(n)$  under different conditions.

^ | v .



**mog** → ibn · 4 months ago

bcuz omega always provide a lower bound on running time of an algorithm. It would be  $\Omega(1)$ , it can't be less than  $\Omega(1)$ .....so we can say as  $\Omega(1)$ .

^ | v .



**Sourabh** → mog · 3 months ago

For a sorting algorithm, best case can not be better than  $\Omega(n)$  elements before saying that it is sorted, so I guess answer is

1 ^ | v .



**rohit** → Sourabh · a month ago

$1 < n \Rightarrow$  yes

^ | v .



**Sourabh** → rohit · 17 days ago

Well you got me there, I saw this answer again today and IMO  $\Omega(1)$  is true for any algorithm (it's a lower bound).

^ | v .



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