

## Divide and Conquer | Set 4 (Karatsuba algorithm for fast multiplication)

Given two binary strings that represent value of two integers, find the product of two strings. For example, if the first bit string is "1100" and second bit string is "1010", output should be 120.

For simplicity, let the length of two strings be same and be  $n$ .

A **Naive Approach** is to follow the process we study in school. One by one take all bits of second number and multiply it with all bits of first number. Finally add all multiplications. This algorithm takes  $O(n^2)$  time.

```

X = 101001 = 41
Y = 101010 = 42
-----
      1010010
     101001
+  101001
-----
11010111010 = 1722
    
```

Using **Divide and Conquer**, we can multiply two integers in less time complexity. We divide the given numbers in two halves. Let the given numbers be X and Y.

For simplicity let us assume that  $n$  is even

```

X = X1*2n/2 + Xr    [X1 and Xr contain leftmost and rightmost n/2 bits of X]
Y = Y1*2n/2 + Yr    [Y1 and Yr contain leftmost and rightmost n/2 bits of Y]
    
```

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The product XY can be written as following.

$$\begin{aligned}XY &= (X_l * 2^{n/2} + X_r)(Y_l * 2^{n/2} + Y_r) \\ &= 2^n X_l Y_l + 2^{n/2}(X_l Y_r + X_r Y_l) + X_r Y_r\end{aligned}$$

If we take a look at the above formula, there are four multiplications of size  $n/2$ , so we basically divided the problem of size  $n$  into four sub-problems of size  $n/2$ . But that doesn't help because solution of recurrence  $T(n) = 4T(n/2) + O(n)$  is  $O(n^2)$ . The tricky part of this algorithm is to change the middle two terms to some other form so that only one extra multiplication would be sufficient. The following is tricky expression for middle two terms.

$$X_l Y_r + X_r Y_l = (X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r$$

So the final value of XY becomes

$$XY = 2^n X_l Y_l + 2^{n/2} * [(X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r] + X_r Y_r$$

With above trick, the recurrence becomes  $T(n) = 3T(n/2) + O(n)$  and solution of this recurrence is  $O(n^{1.59})$ .

*What if the lengths of input strings are different and are not even?* To handle the different length case, we append 0's in the beginning. To handle odd length, we put  $\text{floor}(n/2)$  bits in left half and  $\text{ceil}(n/2)$  bits in right half. So the expression for XY changes to following.

$$XY = 2^{2\text{ceil}(n/2)} X_l Y_l + 2^{\text{ceil}(n/2)} * [(X_l + X_r)(Y_l + Y_r) - X_l Y_l - X_r Y_r] + X_r Y_r$$

The above algorithm is called Karatsuba algorithm and it can be used for any base.

Following is C++ implementation of above algorithm.

```
// C++ implementation of Karatsuba algorithm for bit string multiplication
#include<iostream>
#include<stdio.h>

using namespace std;

// FOLLOWING TWO FUNCTIONS ARE COPIED FROM http://goo.gl/q00hZ
// Helper method: given two unequal sized bit strings, converts them to
// same length by adding leading 0s in the smaller string. Returns the
// the new length
```



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```

int makeEqualLength(string &str1, string &str2)
{
    int len1 = str1.size();
    int len2 = str2.size();
    if (len1 < len2)
    {
        for (int i = 0 ; i < len2 - len1 ; i++)
            str1 = '0' + str1;
        return len2;
    }
    else if (len1 > len2)
    {
        for (int i = 0 ; i < len1 - len2 ; i++)
            str2 = '0' + str2;
    }
    return len1; // If len1 >= len2
}

// The main function that adds two bit sequences and returns the addit
string addBitStrings( string first, string second )
{
    string result; // To store the sum bits

    // make the lengths same before adding
    int length = makeEqualLength(first, second);
    int carry = 0; // Initialize carry

    // Add all bits one by one
    for (int i = length-1 ; i >= 0 ; i--)
    {
        int firstBit = first.at(i) - '0';
        int secondBit = second.at(i) - '0';

        // boolean expression for sum of 3 bits
        int sum = (firstBit ^ secondBit ^ carry)+'0';

        result = (char)sum + result;

        // boolean expression for 3-bit addition
        carry = (firstBit&secondBit) | (secondBit&carry) | (firstBit&c.
    }

    // if overflow, then add a leading 1
    if (carry) result = '1' + result;

    return result;
}

```



```
// A utility function to multiply single bits of strings a and b
int multiplySingleBit(string a, string b)
{   return (a[0] - '0')*(b[0] - '0');   }

// The main function that multiplies two bit strings X and Y and returns
// result as long integer
long int multiply(string X, string Y)
{
    // Find the maximum of lengths of x and Y and make length
    // of smaller string same as that of larger string
    int n = makeEqualLength(X, Y);

    // Base cases
    if (n == 0) return 0;
    if (n == 1) return multiplySingleBit(X, Y);

    int fh = n/2;    // First half of string, floor(n/2)
    int sh = (n-fh); // Second half of string, ceil(n/2)

    // Find the first half and second half of first string.
    // Refer http://goo.gl/lLmgn for substr method
    string Xl = X.substr(0, fh);
    string Xr = X.substr(fh, sh);

    // Find the first half and second half of second string
    string Yl = Y.substr(0, fh);
    string Yr = Y.substr(fh, sh);

    // Recursively calculate the three products of inputs of size n/2
    long int P1 = multiply(Xl, Yl);
    long int P2 = multiply(Xr, Yr);
    long int P3 = multiply(addBitStrings(Xl, Xr), addBitStrings(Yl, Yr));

    // Combine the three products to get the final result.
    return P1*(1<<(2*sh)) + (P3 - P1 - P2)*(1<<sh) + P2;
}
```

```
// Driver program to test above functions
int main()
{
    printf ("%ld\n", multiply("1100", "1010"));
    printf ("%ld\n", multiply("110", "1010"));
    printf ("%ld\n", multiply("11", "1010"));
    printf ("%ld\n", multiply("1", "1010"));
    printf ("%ld\n", multiply("0", "1010"));
    printf ("%ld\n", multiply("111", "111"));
}
```

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```
    printf ("%ld\n", multiply("11", "11"));  
}
```

Output:

```
120  
60  
30  
10  
0  
49  
9
```

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**Time Complexity:** Time complexity of the above solution is  $O(n^{1.59})$ .

Time complexity of multiplication can be further improved using another Divide and Conquer algorithm, fast Fourier transform. We will soon be discussing fast Fourier transform as a separate post.

### Exercise

The above program returns a long int value and will not work for big strings. Extend the above program to return a string instead of a long int value.

### References:

[Wikipedia page for Karatsuba algorithm](#)

[Algorithms 1st Edition by Sanjoy Dasgupta, Christos Papadimitriou and Umesh Vazirani](#)

<http://courses.csail.mit.edu/6.006/spring11/exams/notes3-karatsuba>

<http://www.cc.gatech.edu/~ninamf/Algos11/lectures/lect0131.pdf>

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can the complexity...



**Amit** • 6 days ago

For normal high school multiplication, the complexity is  $O(n^2)$  right?

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**Spanky** • a month ago

Mistake :  $O(n^{159}) \rightarrow O(n^{158}) \approx \log_2(3)$

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**anvesh** • 2 months ago

can any one write code for complex numbers multiplication using FFT

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**frenzydude** • 2 months ago

This can be done even faster in  $O(n \log n)$  using FFT.

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**sudhanshu** • 3 months ago

can any1 pls tell me wht are floor n ceil values ??

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**Himanshu Dagar** → sudhanshu • 3 months ago

floor means greatest integer less than or equal to that no

and ceil is least integer greater than or equal to that no

e.g floor(2.6)=2;

ceil(2.6)=3;

defined in math.h

^ | v • Reply • Share ›



**Sanjeet** • 5 months ago

/\*\*\*\*\* Sanjeet from NIT allahabad \*\*\*\*\*/ you can simply adjust the value of i

strings .

```
#include<stdio.h>
```

```
#include<math.h>
```

```
int convertbin(char a[]);
```

```
int main()
```

```
{
```

```
char string1[4]="0111";
```

```
char string2[4]="0111";
```

```
int DecimalNum1=0;
```

```
int DecimalNum2=0;
```

```
int Product=0;
```

```
DecimalNum1=convertbin(string1);
```

```
DecimalNum2=convertbin(string2);
```

```
Product=DecimalNum1*DecimalNum2;
```

```
printf("Product of two string =%d\n",Product);
```

```
return 0;
```

```
}
```

[see more](#)

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**Shabaz Ahmed** → Sanjeet • 2 months ago

The point here is to do the multiplication in less than  $O(n^2)$ ..!

^ | v • Reply • Share ›



**Raunak Lakhwani** • 6 months ago

```
public class Main {
```

```
static String first = "1", second = "00110";
```

```
/**
```

```
* @param args
```



@param args

\*/

```
public static void main(String[] args) {  
  
    //int n = makeEqualString(first, second);  
  
    System.out.println(multiplication(first, second));  
  
}  
  
public static int makeEqualString(StringBuffer a, StringBuffer b) {  
  
    if (a.length() > b.length()) {
```

[see more](#)

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**Trilok** • 7 months ago

Typo mistake ..

$XY = 2n \cdot XYI + 2\text{ceil}(n/2) * [(XI + Xr)(YI + Yr) - XYI - XrYr] + XrYr$

^ | v • Reply • Share ›



**komal** • a year ago

#include

#include

#include

using namespace std;

int pow(int x,int n)

{

if(x==0) return 0;

if(n==0) return 1;

if(n==1) return x;

```

if(n==1) return (x);
return (x*pow(x,n-1));
}
void print(string s)
{
for(int i=0;i<=s.size();i++)
cout<<s[i]<<" ";
cout<<endl;
}

```

[see more](#)

^ | v • Reply • Share ›



**akshat gupta** → komal • a year ago

- 1.you are entering a number in binary representation(serves no purpos
- 2.converting it to decimal,
- 3.printing the multiplication...

PROBLEM:

question asks you to find an algorithm for "fast multiplication of 2 numb

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**akshat gupta** • a year ago

```

int karatsuba(int n1,int n2)
{
int l1,l2,l;
int i,mask1=1;
int a,b,c,d;
int res1,res2,res3;

```

l1=dig(n1);//returns number of digits

l2=dig(n2);

```

if(n1==0 || n2==0)

```

```
return(0);

l=min(l1,l2);

if(l>1)
{
for(i=1;i<=l/2;i++)
mask1=mask1*10;
```

[see more](#)

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**akshat gupta** • a year ago

C IMPLEMENTATION:

```
int karatsuba(int n1,int n2)
{
int l1,l2,l;
int i,mask1=1;
int a,b,c,d;
int res1,res2,res3;
l1=dig(n1);//returns number of digits
l2=dig(n2);
if(n1==0 || n2==0)
return(0);

l=min(l1,l2);

if(l>1)
{
for(i=1;i<=l/2;i++)
mask1=mask1*10;
```

[see more](#)

^ | v · Reply · Share ›



**Sandeep Jain** · a year ago

Thanks for pointing this out. We have corrected it.

^ | v · Reply · Share ›



**Anwit Roy** · a year ago

There is a typo: it should be  $XIYr + XrYl = (XI + Xr)(Yl + Yr) - XIYl - XrYr$ .

^ | v · Reply · Share ›



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