

Backtracking | Set 5 (m Coloring Problem)

Given an undirected graph and a number m , determine if the graph can be colored with at most m colors such that no two adjacent vertices of the graph are colored with same color. Here coloring of a graph means assignment of colors to all vertices.

Input:

- 1) A 2D array $graph[V][V]$ where V is the number of vertices in graph and $graph[V][V]$ is adjacency matrix representation of the graph. A value $graph[i][j]$ is 1 if there is a direct edge from i to j , otherwise $graph[i][j]$ is 0.
- 2) An integer m which is maximum number of colors that can be used.

Output:

An array $color[V]$ that should have numbers from 1 to m . $color[i]$ should represent the color assigned to the i th vertex. The code should also return false if the graph cannot be colored with m colors.

Following is an example graph (from [Wiki page](#)) that can be colored with 3 colors.

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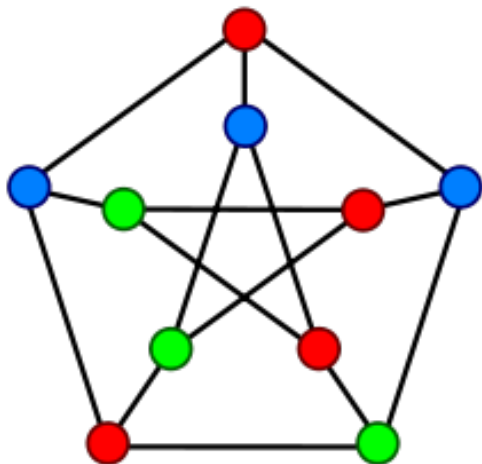
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Naive Algorithm

Generate all possible configurations of colors and print a configuration that satisfies the given constraints.

```
while there are untried configurations
{
    generate the next configuration
    if no adjacent vertices are colored with same color
    {
        print this configuration;
    }
}
```

There will be V^m configurations of colors.

Backtracking Algorithm

The idea is to assign colors one by one to different vertices, starting from the vertex 0. Before assigning a color, we check for safety by considering already assigned colors to the adjacent vertices. If we find a color assignment which is safe, we mark the color assignment as part of solution. If we do not find a color due to clashes then we backtrack and return false.

Implementation of Backtracking solution

```
#include<stdio.h>
```

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```
// Number of vertices in the graph
#define V 4

void printSolution(int color[]);

/* A utility function to check if the current color assignment
is safe for vertex v */
bool isSafe (int v, bool graph[V][V], int color[], int c)
{
    for (int i = 0; i < V; i++)
        if (graph[v][i] && c == color[i])
            return false;
    return true;
}

/* A recursive utility function to solve m coloring problem */
bool graphColoringUtil(bool graph[V][V], int m, int color[], int v)
{
    /* base case: If all vertices are assigned a color then
    return true */
    if (v == V)
        return true;

    /* Consider this vertex v and try different colors */
    for (int c = 1; c <= m; c++)
    {
        /* Check if assignment of color c to v is fine*/
        if (isSafe(v, graph, color, c))
        {
            color[v] = c;

            /* recur to assign colors to rest of the vertices */
            if (graphColoringUtil (graph, m, color, v+1) == true)
                return true;

            /* If assigning color c doesn't lead to a solution
            then remove it */
            color[v] = 0;
        }
    }

    /* If no color can be assigned to this vertex then return false */
    return false;
}

/* This function solves the m Coloring problem using Backtracking.
It mainly uses graphColoringUtil() to solve the problem. It returns
```

stack!

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```

false if the m colors cannot be assigned, otherwise return true and
prints assignments of colors to all vertices. Please note that there
may be more than one solutions, this function prints one of the
feasible solutions.*/
bool graphColoring(bool graph[V][V], int m)
{
    // Initialize all color values as 0. This initialization is needed
    // correct functioning of isSafe()
    int *color = new int[V];
    for (int i = 0; i < V; i++)
        color[i] = 0;

    // Call graphColoringUtil() for vertex 0
    if (graphColoringUtil(graph, m, color, 0) == false)
    {
        printf("Solution does not exist");
        return false;
    }

    // Print the solution
    printSolution(color);
    return true;
}

```

```

/* A utility function to print solution */
void printSolution(int color[])
{
    printf("Solution Exists:"
           " Following are the assigned colors \n");
    for (int i = 0; i < V; i++)
        printf(" %d ", color[i]);
    printf("\n");
}

// driver program to test above function
int main()
{
    /* Create following graph and test whether it is 3 colorable
    (3)---(2)
    |    /  |
    |   /   |
    |  /    |
    (0)---(1)
    */
    bool graph[V][V] = {{0, 1, 1, 1},
                        {1, 0, 1, 0},
                        {1, 1, 0, 1},

```

705



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```

        {1, 0, 1, 0},
    };
    int m = 3; // Number of colors
    graphColoring (graph, m);
    return 0;
}

```

Output:

Solution Exists: Following are the assigned colors

1 2 3 2

References:

http://en.wikipedia.org/wiki/Graph_coloring

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sr7 · 3 months ago

you can further bring down the complexity by avoiding the check for each color and no of vertices is N .

Checking for each color will take $O(n^2)$ time . Instead

I assume the base color (uncolored node) is -1 . Colors start from 0 .

I will check all the adjacent nodes of the given vertex in $O(n)$ time .

If i th node is adjacent to vertex v and $color[i] \geq 0$ (colored node) , then $temp[color[i]]$ (the color is in use by one of its adjacent nodes) , finally I will search the 'temp'

if $temp[i] = 1$, then that is the least color not used by any of the adjacent nodes on the whole, instead of $O(n^2)$.

```
int temp[V];
```

```
//Initialize all temp values to zero
```

```
for (i=0;i<v;i++) temp[i]="0;" for="" (i="0;i<V;i++)" if="" (graph[v][i]="" &&="" c  
colored
```

```
temp[color[i]]=1;
```

```
for (i=0; i<m ;="" i++)="" {="" only="" 0="" to="" m-1="" colors="" are="" allowed  
if="" (graphcoloringutil(graph,m,color,v+1))="" return="" true;="" color[v]="-1;" k
```

^ | v .



pavi.8081 · 10 months ago

Can we modify this algorithm to find the min. number of colors to color the gra
tried though.

If we can do so then the modified algorithm can be used to solve problems like
halls required to organize parties with given starting and finishing times.

Please comment...

^ | v .



soupboy · 2 years ago

Hi, shouldn't the brute force algorithm's complexity be m^V and not V^m as wri

^ | v .



coolguy → soupboy · 2 years ago

You can understand graph coloring using backtracking by watching this
youtu.be/CI3A_9hokjU

^ | v .



Rohit Raj · 2 years ago

```
bool isSafe (int v, bool graph[V][V], int color[], int c)  
{
```

```

for (int i = 0; i < V; i++)
if (graph[v][i] && c == color[i])
return false;
return true;
}

```

as the problem statement says that if $(i == j)$ then also the $graph[i][j] = 1$; so s condition in this function?

i.e if($graph[v][i] \&\& c == color[i] \&\& v != i$).

^ | v .



kartik → Rohit Raj · 2 years ago

@Rohit Raj

We don't need to check for this as we assign a color c only after isSafe we assign the color value as 0 which is not a valid value. When we call than 0 and value of color[v] will be 0. So color[v] will never be equal to c

^ | v .



gautam · 2 years ago

Can be use some dynamic programming to reduce complexity

```

/* Paste your code here (You may delete these lines if not writing co

```

^ | v .



Venki · 2 years ago

There is small bug in the code. The colour vector should be of size V in lieu of tries to access elements past the buffer end (while checking validity of colour :

What is the worst case complexity? $O(V * m V)$ as we are exploring all configurations.

^ | v .



GeeksforGeeks → Venki · 2 years ago

@Venki: Thanks for pointing this out. This has been fixed. Yes, the worst case as we may have to try all configurations in worst case.

^ | v ·



Guddu sharma · 2 years ago

What is the need of following if condition:

```
if (graphColoringUtil (graph, m, color, v+1) == true)
return true;
```

I think when all the vertices have been assigned color, the first condition if(v==\

Please comment if i have understood something wrong..

^ | v ·



Venki → Guddu sharma · 2 years ago

@Guddu sharma, We need to find *one* feasible solution. The condition ends infinite recursion. When we reach the end branch in the state space, we need to backtrack to one level upper and restore the state.

Similarly, if the current assignment is safe, (means the branch generates a feasible state), we recur to next level for next feasible color.

The condition

```
/* recur to assign colors to rest of the vertices */
if (graphColoringUtil (graph, m, color, v+1) == true)
    return true;
```

checks next color. If this next color reaches last color, we are done with

```
return true;
```

Try to draw the state space tree for small values of V and m, you can see the procedure.

^ | v .



kartik → Guddu sharma · 2 years ago

@Guddu sharma: The statement serves two purposes:

- 1) It calls the function recursively for the other vertices.
- 2) It returns true as soon as one of the color assignments lead to a solution checked as soon as a color assignment works.

Also, just putting following line is not sufficient. We need to forward the calls.

```
if (v == V)
    return true;
```

1 ^ | v .

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