

## Move all zeroes to end of array

Given an array of random numbers, Push all the zero's of a given array to the end of the array. For example, if the given arrays is {1, 9, 8, 4, 0, 0, 2, 7, 0, 6, 0}, it should be changed to {1, 9, 8, 4, 2, 7, 6, 0, 0, 0, 0}. The order of all other elements should be same. Expected time complexity is  $O(n)$  and extra space is  $O(1)$ .

There can be many ways to solve this problem. Following is a simple and interesting way to solve this problem.

Traverse the given array 'arr' from left to right. While traversing, maintain count of non-zero elements in array. Let the count be 'count'. For every non-zero element arr[i], put the element at 'arr[count]' and increment 'count'. After complete traversal, all non-zero elements have already been shifted to front end and 'count' is set as index of first 0. Now all we need to do is that run a loop which makes all elements zero from 'count' till end of the array.

Below is C++ implementation of the above approach.

```
// A C++ program to move all zeroes at the end of array
#include <iostream>
using namespace std;

// Function which pushes all zeros to end of an array.
void pushZerosToEnd(int arr[], int n)
{
    int count = 0; // Count of non-zero elements

    // Traverse the array. If element encountered is non-zero, then
    // replace the element at index 'count' with this element
    for (int i = 0; i < n; i++)
        if (arr[i] != 0)
            arr[count++] = arr[i]; // here count is incremented
```

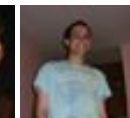
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```
// Now all non-zero elements have been shifted to front and 'count
// set as index of first 0. Make all elements 0 from count to end.
while (count < n)
    arr[count++] = 0;
}
```

```
// Driver program to test above function
int main()
{
    int arr[] = {1, 9, 8, 4, 0, 0, 2, 7, 0, 6, 0, 9};
    int n = sizeof(arr) / sizeof(arr[0]);
    pushZerosToEnd(arr, n);
    cout << "Array after pushing all zeros to end of array :\n";
    for (int i = 0; i < n; i++)
        cout << arr[i] << " ";
    return 0;
}
```

Output:

```
Array after pushing all zeros to end of array :
1 9 8 4 2 7 6 9 0 0 0 0
```

**Time Complexity:**  $O(n)$  where  $n$  is number of elements in input array.

**Auxiliary Space:**  $O(1)$

This article is contributed by **Chandra Prakash**. Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

.....

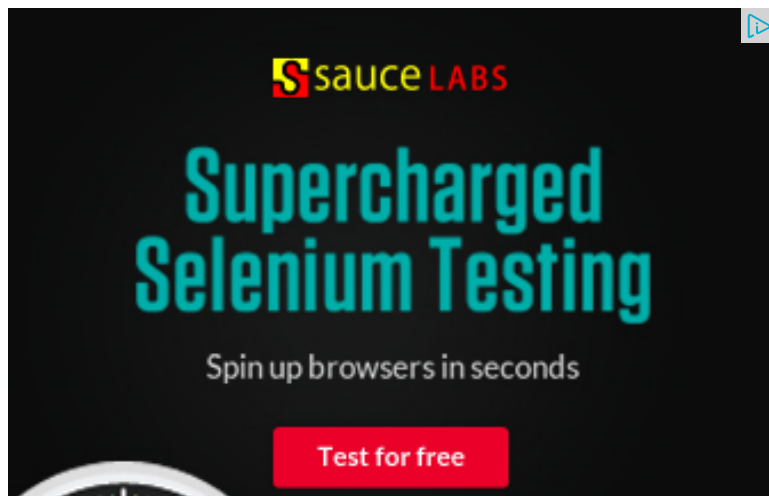
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**Lohith Ravi** · 18 days ago

```
public static void moveZeroToEnd(int array[]) {
```

```
    int NumberPtr = 0;
```

```
    int ZeroPtr = 0;
```

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```
while (array[NumberPtr] == 0 && NumberPtr < array.length - 1 )
```

```
NumberPtr++;
```

```
while (array[ZeroPtr] != 0 && ZeroPtr < array.length - 1)
```

```
ZeroPtr++;
```

```
if(ZeroPtr < NumberPtr)
```

```
{
```

```
swap(array, ZeroPtr, NumberPtr):
```

[see more](#)

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**abc** • a month ago

```
#include<iostream>
```

```
using namespace std;
```

```
#define swap(a,b,c) {c=a;a=b;b=c;}
```

```
int main()
```

```
{
```

```
int a[]={1,0,5,0,6,3,0,4,7,0};
```

```
int len=sizeof(a)/sizeof(a[0]);
```

```
int low=0,high=len-1;
```

705



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```
int i=0,temp=0;
```

```
while(low<high) {="" if(a[high]=="=0)" {="" high--;"="}="" if(a[low]=="=0&amp;&an  
"<<high<<endl;"=" swap(a[low],a[high],temp);"=" low++;"=" high--;"="}="" low+  
cout<<a[i]<<"=" ";"=" system("pause);"="}="">
```

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**sumit gaur** • 2 months ago

Your code's complexity will be  $O(n+cn)$  as upto some extend you are traversir  
here's with  $O(n)$  traverse only once

```
void rearrange(int a[],int size)
```

```
{
```

```
int end=0;
```

```
for(int i=0;i<size;i++) {="" if(a[i])="" {="" int="" tmp="a[end];" a[end]="a[i];" a[i]="
```

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**Guest** • 2 months ago

Your code's complexity will be  $O(n+cn)$  as upto some extend you are traversir

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```

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Guest • 2 months ago

Your code's complexity will be  $O(n+cn)$  as upto some extend you are traversir here's with  $O(n)$  traverse only once

```
void rearrange(int a[],int size)
```

```
{
```

```
int end=0;
```

```
for(int i=0;i<size;i++) {="" if(a[i])="" {="" int="" tmp="a[end];" a[end]="a[i];" a[i]="
```

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Ramesh • 3 months ago

```
private void move(int[] arr) {
```

```
int countZeroes = 0;
```

```
for(int i=0; i < arr.length; i++) {
```

```
if(arr[i] == 0) {
```

```
countZeroes++;
```

```
} else {
```

```
if(countZeroes != 0) {
```

```
arr[i-countZeroes] = arr[i];
```

```
arr[i] = 0;
```

```
}
```

```
}
```

```
}
```

```
}
```

3 ^ | v • Reply • Share ›



**swap** • 3 months ago

```
public class MoveZerosToEnd {  
  
    public static void main(String... args) {  
  
        MoveZerosToEnd m = new MoveZerosToEnd();  
  
        System.out.println(Arrays.toString(m.moveZero(new int[]{0, 1, 2, 3, 0, 9, 4})));  
  
    }  
  
    public int[] moveZero(int[] arr) {  
  
        int k = 1;  
  
        int i = 0;  
  
        while (i < arr.length) {  
  
            if (arr[i] == 0) {  
  
                for (int i = i + 1; i < arr.length; i++) {
```

[see more](#)

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**Nodirbek** • 3 months ago

How about my algorithm:

For each element of array from right to left check if it's zero then swap it with l: of last non-zero element, and shift to left as we swap. Since the problem state



(not actually shifting non-zero numbers to left) this is correct solution. Below c

<http://ideone.com/12oEWe>

2 ^ | v • Reply • Share ›



**rr** • 4 months ago

```
import java.util.Arrays;
```

```
class test {
```

```
    public static void main(String[] args) {
```

```
        int[] x = {0,1,2,3};
```

```
        int count = 0;
```

```
        if(x == null){
```

```
            return;
```

```
        }
```

```
        for(int i = 0; i<x.length; i++){
            if(x[i] != "0"){
                x[count] = x[i];
                count++;
            }
            x[j++] = "0";
            system.out.println(Arrays.toString(x));
        }
```

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**ravi jadhav** • 4 months ago

Here is a simple fact to exploit. Multiplying zero with any element produces a zero element, just keep looking whether the multiplication of successive pairs is zero or not. If both are zero, then the result is zero. You can use subtraction to check if both are zero or not. If both are zero, then the result is zero. You can use subtraction to check if both are zero or not. If both are zero, then the result is zero.

Here is the java code i wrote. I believe when the number of zeroes are lesser or equal to the number of non-zeroes, the complexity is lesser number of comparison. On average  $n/2$ .

<http://ideone.com/IIYYsr>

4 ^ | v • Reply • Share ›



**Rahul Ramesh** → ravi jadhav • 2 months ago

But multiplying causes overhead right????

I am not very sure....

^ | v • Reply • Share ›



**rihansh** • 4 months ago

this can also be done by selecting 0 as pivot element as in case of quick sort I elements

^ | v • Reply • Share ›



**rihansh** • 4 months ago

BUT THIS WILL TAKE TWO PASSES WHEN ALL ELEMENT IN THE ARRAY AN BETTER OPTION

START iterating from right side kept one pointer at end and other one is iterain element we replace it with the pointer position and decrementing hope u get m

1 ^ | v • Reply • Share ›



**rj** → rihansh • 4 months ago

@rihansh but in this case, order of non-zero elements will not retain

1 ^ | v • Reply • Share ›



**bikash** • 5 months ago

I think we can approach this problem by keeping two pointers indicating the cu position, and swapping them. This continues till all zeros are moved to the enc

<http://ideone.com/UMHryG>

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**Jonathan** ➔ bikash · a month ago



000000111

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Proved sol in  $O(n)$ .

```
using namespace std;
```

{

```
for(int j=0;j<n;j++) {="" if(arr[j]!="0") {="" i++;="" arr[i]=arr[j];="" }="" else="" conti
arr[j]="0;" }="" int="" main()="" {="" int="" arr[]={1," 9,"="" 8,"="" 4,"="" 0,"="" 0,"=""
n="sizeof(arr)" sizeof(arr[0]);="" pushzerostoend(arr,="" n);="" cout="" <<="" ",
zeros="" to="" end="" of="" array="" :\n";="" for="" (int="" i="0;" i="" <="" n;="" i-
";="" return="" 0;="" }="">
```



trojansmith • 5 months ago

Hi,,this partitions the array into non zero and zero

```
int[] arr = {1,9,8,4,0,0,2,7,0,6,0};
    int temp=0;
    int i=0,j=0;

    //reach first zero
    while(arr[i]!=0)i++;
    j=i;
    //partition array into 0 and non-zero
    while(i<arr.length &&"" j<arr.length){="" if(arr[i]!="0){" s
```

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trojansmith • 5 months ago

trojansmith...hi this i have given will partition the array into non zero and zeroes:

```
int[] arr = {1,9,8,4,0,0,2,7,0,6,0};
    int temp=0;
    int i=0,j=0;

    //reach first zero
    while(arr[i]!=0)i++;
    j=i;
    //partition array into 0 and non-zero
    while(i<arr.length &&"" j<arr.length){="" if(arr[i]!="0){" s
```

^ | v • Reply • Share ›



xyz • 6 months ago



int i=0,j=0;

while(j!=length)

{

if(array[j]!=0)

{

int temp=array[j];

array[j]=0;

array[i]=temp;

i++;

j++;

}

else

{

j++;

}

}

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**Guest** • 6 months ago

#include<stdio.h>

```
int main(){

int n;

scanf("%d",&n);

int i,arr[n];

for(i=0;i<n;i++) scanf("%d",&arr[i]);="" int="" j="0;" i="0;" while(i<n="" &&="" j<r
i++;="" if(i<n){="" j="i+1;" while(!arr[j]="" &&="" j<n)="" j++;="" if(j="">=n)

break;

int temp = arr[i];

arr[i] = arr[j];

arr[j] = temp;
```

---

[see more](#)

^ | v • Reply • Share ›



**Sheetal Garg** • 6 months ago

```
#include<stdio.h>
```

```
int main()

{

int flag=0,temp,i,j,arr[10]={1,0,9,8,0,0,6,2,0,1};

i=0;j=i+1;

while(j<10)
```

```
if(arr[i]!=0&&flag==0)
```

```
{
```

```
  i++;
```

```
  i++;
```

[see more](#)

1 ^ | v • Reply • Share ›



**Pradeep Patel** • 7 months ago

Performance can be improved if self replacement can be avoided. In case the

```
for (int i = 0; i < n; i++)
    if (arr[i] != 0)
    {
        if(count != i) //avoid self replacement
            arr[count] = arr[i]; // here count is incremented
        ++count;
    }
```

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**krrish** → Pradeep Patel • 6 months ago

You can not bet it to be improvement as perhaps conditional evaluation assignment(depends on processor).Also you are performing n conditic

3 ^ | v • Reply • Share ›



**Harish Kumar** • 7 months ago

"Correct me if i m wrong"

Well we can do it by two pointers and we traverse the whole array only once & writes

Take two pointers  $i=0$  and  $j=0$

do these steps until  $j$  is less than length

--> assign 'i' the first zeroth index from left

--> assign 'j' the first non Zeroth index from left after 'i' (greater than i)

--> Swap the values at those indices

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**Pushkar** → Harish Kumar • 5 months ago

Worst Case your algo might take  $O(n^2)$

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**Guest** • 7 months ago

"Correct me if i m wrong"

Well we can do it by two pointers and we traverse the whole array only once & writes

Take two pointers  $i=0$  and  $j=0$

do these steps until  $j < n$  --> assign 'i' the first zeroth index from left

--> assign 'j' the first non Zeroth index from left after 'i' (greater than i)

--> Swap the values at those indices

The code in c is as follows

```
void shiftZeros(int array[],int n)
```



```

int i=0,j=0;
i = getNextZeroIndex(i,n,array);
if( i == -1)
return;
j = i+1;
while(j<n) {="" if(array[j])="" {="" array[i]="array[j];" array[j]="0;" i++;="" i="getNe
return;="" j="i;" }="" j++;="" }="" printarray(array,n);="" }="" int="" getnextzeroinc
while(i<n="" &&="" array[i]="" !="0") {="" if(="" i="" n)="" return="" -1;="" i++;=""

```

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**Pravin** → Guest • 5 months ago

Please forgive my ignorance.. wats the programming language that is I

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**Himanshu Mishra** • 7 months ago

My in python. Surprisingly same Algorithm :)

```

import time
def main():
    start = time.time()
    a = [1, 9, 8, 4, 0, 0, 2, 7, 0, 6, 0]
    b = []
    count = 0
    for i in a:
        if i != 0:b.append(i)
        else:count += 1
    c = [0] * count
    answer = b + c
    print(answer)
    print("time taken : ",time.time() -start)
    return 0

```

```
main()
```

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**Dinesh Panchananam** → Himanshu Mishra • 5 months ago

# Pyth 2.7 Shorter code.

```
f= [1,9,8,4,0,0,2,7,0,6,0,23,32]
g = filter(lambda x:x!=0 ,f)
print g+(len(f)-len(g))*[0]
```

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**sathish** • 7 months ago

I tried a Simple one.If anything wrong ,please advice.The code is working 100%

no extra space used.

```
void movee(int *a,int n)
```

```
{
```

```
int i,c=n-1;
```

```
for(i=0;i<=c;i++)
```

```
{
```

```
if(*(a+i)==0 && *(a+c)==0)
```

```
{
```

```
c--;
```

```
while(c>i)
```

[see more](#)

1 ^ | v • Reply • Share ›



**shashi** → sathish • 7 months ago

dude, maintain order

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**Ruhi Sharma** • 7 months ago

We can also do this by taking two variables i and j, initialize i as index 0 and j as  
Now, keep incrementing i till we find a 0 and then keep decrementing j till we find a 0  
Continue this until i >= j

C code for the above approach is :

```
front_index=0;
end_index=no_of_elem-1;
while(front_index < end_index)
{
    if(arr[front_index]==0)
    {
        if(arr[end_index]!=0){
            arr[front_index]= arr[end_index];
            arr[end_index]=0;
            front_index++;
            end_index--;
        }
        else{
            end_index--;
        }
    }
    else{
        front_index++;
    }
}
```

```
}  
}
```

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**abhisek** → Ruhi Sharma • 7 months ago

we cannot use dutch flag here as the order has to be maintained

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**guest** → Ruhi Sharma • 7 months ago

The order of non-zero elements are not preserved by this method.

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**atiq** • 7 months ago

@GeeksforGeeks ....Good exploitation but

No need to do this part.....

// Now all non-zero elements have been shifted to front and 'count' is

// set as index of first 0. Make all elements 0 from count to end.

while (count < n)

arr[count++] = 0

just do replace a[i] with 0 at the same moment

for (int i = 0; i < n; i++)

if (arr[i] != 0)

{ arr[count++] = arr[i];

arr[i]=0;

}

That's how in one pass..... ;)

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**ssd** → atiq • 7 months ago

 this will not work. try it on the array with one element [ 1]

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**rk** → ssd • 7 months ago

We can set `arr[i]=0` only when `(i != count-1)` . Is there any other

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**pavansrinivas** • 7 months ago

Code in JAVA (similar to partition in Quicksort)...order is not preserved...

```
void arrange(){
    int[]a = {1, 9, 8, 4, 0, 0, 2, 7, 0, 6, 0};
    int l = 0;
    int r = a.length-1;
    while(l<r){ while(a[l]!="0&&l<r){ l++;=""
                r--;
            }
            a[l] = a[r];
            a[r] = 0;
            l++;
            r--;
        }
    for(int i=0;i<a.length;i++){ system.out.print(a[i]+"="}
```

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**Rashik** • 7 months ago

```
#define N 15
```

```
void main()
```

```
{
```

```
int ar[N]={0,0,20,30,0,0,0,40,0,0,50,0,60,0};
```

```
int j=0,i=0;

while(1)

{

while(ar[i]!=0)

i++;

j=i ;

while(ar[j]==0)
```

[see more](#)

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**Satyendra Kumar Singh** • 7 months ago

```
void ShiftZeroToEnd(int arr[],int n)
{
int cur=-1;
for(int i=0;i<n;i++) {="" if(arr[i]=="0") {="" if(cur=="-1") cur="i;" }="" else="" if(cu
cur++;" }="" }="" }="" the="" above="" program="" will="" push="" all="" zeros
```

1 ^ | v • Reply • Share ›



**sijayaraman** • 7 months ago

```
void move_zero_to_end(int arr[],int len)
{
int j = -1;
for(int i=0;i<len;i++) {="" if(arr[i]!="0" )="" {="" j++;" swap(&arr[i],&arr[j]);="" }=""
*b)="" {="" int="" t="" *a;" *a="" *b;" *b="" t;" }="">
```

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guest • 7 months ago

```
#include<iostream>
using namespace std;
int main()
{
int a[]={1, 9, 8, 4, 0, 0, 2, 7, 0, 6, 0, 9};
int n=sizeof(a)/sizeof(a[0]);
int m=n,k;
for(int i=0;i<n;i++) {="" if(a[i]=="0)" for(int="" j="i+1;j<=n;j++)" if(a[j]!="0)" {="" k
}="" for(int="" i="0;i<=n;i++)" cout<<a[i]<<" ";;="" }="">
```

1 ^ | v • Reply • Share ›



Guest • 7 months ago

```
#include<iostream>
using namespace std;

int main() {
while (1) {
int array[30];
int arraySize;
cin >> arraySize;
for (int i = 0; i < arraySize; i++) {
cin >> array[i];
}
int writePtr = 0;
int temp;
for (int readPtr = 0; readPtr < arraySize; readPtr++) {
temp = array[readPtr];
if (temp != 0) {
array[writePtr] = temp;
writePtr++;
}
```

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**REYAZ** • 7 months ago

another simple method is to maintain two indices  $i$  and  $j$ , traverse the array till  $i$  point to this element and then we increment  $j$  till the end of the array while doing if element pointed to by  $j$  is 0 then  $j++$ , else if element pointed to by  $j$  is other than 0, then swap the elements of position  $i$  and  $j$ .

COMPLEXITY:  $O(n)$ .

this is a modified version of DNF algo, in which while segregating three types, one type of objects with their order preserved.. plz correct me if i am wrong

6 ^ | v • Reply • Share ›



**Vijay Kumar** → REYAZ • 7 months ago

```
int i = 0;
for( i=0; i<n; i++){if(a[i]!=0){break;}} for(int j=i; j<n; j++){if(a[j]==0){swap(a[i], a[j]); i++;}}
```

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**Hector** → REYAZ • 7 months ago

I think this is correct!!!!

1 ^ | v • Reply • Share ›



**its\_dark** → REYAZ • 7 months ago

OR

you can also start  $i$  as index 0 and  $j$  as index ' $n-1$ '.

Now, keep incrementing  $i$  till we find a 0 and then keep decrementing  $j$  them.

continue till  $i > j$ .

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**Jekin** → its\_dark • 7 months ago

Yes. But then it won't preserve original order.

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