

Find the element that appears once

Given an array where every element occurs three times, except one element which occurs only once. Find the element that occurs once. Expected time complexity is $O(n)$ and $O(1)$ extra space.

Examples:

Input: `arr[] = {12, 1, 12, 3, 12, 1, 1, 2, 3, 3}`

Output: 2

We can use sorting to do it in $O(n\log n)$ time. We can also use hashing, but the worst case time complexity of hashing may be more than $O(n)$ and hashing requires extra space.

The idea is to use bitwise operators for a solution that is $O(n)$ time and uses $O(1)$ extra space. The solution is not easy like other XOR based solutions, because all elements appear odd number of times here. The idea is taken from [here](#).

Run a loop for all elements in array. At the end of every iteration, maintain following two values.

ones: The bits that have appeared 1st time or 4th time or 7th time .. etc.

twos: The bits that have appeared 2nd time or 5th time or 8th time .. etc.

Finally, we return the value of 'ones'

How to maintain the values of 'ones' and 'twos'?

'ones' and 'twos' are initialized as 0. For every new element in array, find out the common set bits in the new element and previous value of 'ones'. These common set bits are actually the bits that should be added to 'twos'. So do bitwise OR of the common set bits with 'twos'. 'twos' also gets some extra bits that appear third time. These extra bits are removed later.

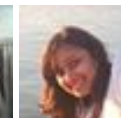
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Update 'ones' by doing XOR of new element with previous value of 'ones'. There may be some bits which appear 3rd time. These extra bits are also removed later.

Both 'ones' and 'twos' contain those extra bits which appear 3rd time. Remove these extra bits by finding out common set bits in 'ones' and 'twos'.

```
#include <stdio.h>
```

```
int getSingle(int arr[], int n)
{
    int ones = 0, twos = 0 ;

    int common_bit_mask;

    // Let us take the example of {3, 3, 2, 3} to understand this
    for( int i=0; i< n; i++ )
    {
        /* The expression "one & arr[i]" gives the bits that are
           there in both 'ones' and new element from arr[]. We
           add these bits to 'twos' using bitwise OR

           Value of 'twos' will be set as 0, 3, 3 and 1 after 1st,
           2nd, 3rd and 4th iterations respectively */
        twos = twos | (ones & arr[i]);

        /* XOR the new bits with previous 'ones' to get all bits
           appearing odd number of times

           Value of 'ones' will be set as 3, 0, 2 and 3 after 1st,
           2nd, 3rd and 4th iterations respectively */
        ones = ones ^ arr[i];

        /* The common bits are those bits which appear third time
           So these bits should not be there in both 'ones' and 'twos'
           common_bit_mask contains all these bits as 0, so that the b
           be removed from 'ones' and 'twos'

           Value of 'common_bit_mask' will be set as 00, 00, 01 and 10
           after 1st, 2nd, 3rd and 4th iterations respectively */
        common_bit_mask = ~(ones & twos);

        /* Remove common bits (the bits that appear third time) from 'ones' and 'twos' */
    }
}
```



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```

        Value of 'ones' will be set as 3, 0, 0 and 2 after 1st,
        2nd, 3rd and 4th iterations respectively */
        ones &= common_bit_mask;

        /* Remove common bits (the bits that appear third time) from 'ones'

        Value of 'twos' will be set as 0, 3, 1 and 0 after 1st,
        2nd, 3rd and 4th iterations respectively */
        twos &= common_bit_mask;

        // uncomment this code to see intermediate values
        //printf (" %d %d \n", ones, twos);
    }

    return ones;
}

int main()
{
    int arr[] = {3, 3, 2, 3};
    int n = sizeof(arr) / sizeof(arr[0]);
    printf("The element with single occurrence is %d ",
           getSingle(arr, n));
    return 0;
}

```

Output:

2

Time Complexity: $O(n)$

Auxiliary Space: $O(1)$

Following is another $O(n)$ time complexity and $O(1)$ extra space method suggested by *aj*. We can sum the bits in same positions for all the numbers and take modulo with 3. The bits for which sum is not multiple of 3, are the bits of number with single occurrence.

Let us consider the example array {5, 5, 5, 8}. The 101, 101, 101, 1000

Sum of first bits%3 = $(1 + 1 + 1 + 0)\%3 = 0$;

Sum of second bits%3 = $(0 + 0 + 0 + 0)\%3 = 0$;

Sum of third bits%3 = $(1 + 1 + 1 + 0)\%3 = 0$;

Sum of fourth bits%3 = $(1)\%3 = 1$.

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Sum of fourth bits is 1000
Hence number which appears once is 1000

```
#include <stdio.h>
#define INT_SIZE 32

int getSingle(int arr[], int n)
{
    // Initialize result
    int result = 0;

    int x, sum;

    // Iterate through every bit
    for (int i = 0; i < INT_SIZE; i++)
    {
        // Find sum of set bits at ith position in all
        // array elements
        sum = 0;
        x = (1 << i);
        for (int j=0; j< n; j++ )
        {
            if (arr[j] & x)
                sum++;
        }

        // The bits with sum not multiple of 3, are the
        // bits of element with single occurrence.
        if (sum % 3)
            result |= x;
    }

    return result;
}

// Driver program to test above function
int main()
{
    int arr[] = {12, 1, 12, 3, 12, 1, 1, 2, 3, 2, 2, 3, 7};
    int n = sizeof(arr) / sizeof(arr[0]);
    printf("The element with single occurrence is %d ",
           getSingle(arr, n));
    return 0;
}
```

7

705



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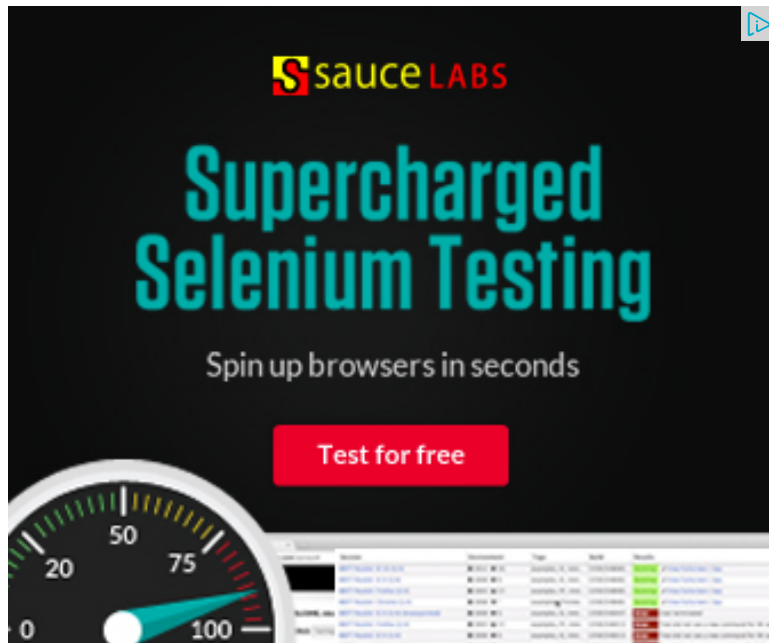
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AN • a month ago

can u explain the logic behind 2nd solution?

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Naveen • 3 months ago

Second implementation is incorrect, it will pass only when non-unique elemen

Failure Case:

1. input[] = { 5, 5, 1}

Sum of first bits = { 1 + 1 + 1}%3 = 0

Sum of second bits = { 0 + 0 + 0}%3 = 0

Sum of third bits = {1 + 1 + 0}%3 = 2

ans = 4;

^ | ▾ • Reply • Share ›



wliao → Naveen • 23 days ago

The problem statement, "Given an array where every element occurs t

^ | ▾ • Reply • Share ›



Iallu • 3 months ago

but the compleity is not O(n) it is O(nlog(MAXelement in the array))

am i right or correct me if i am wrong

^ | ▾ • Reply • Share ›



Guest • 3 months ago

if in first method ,in given example suppose there is no occurence of 4th integ

2 correct me

^ | v • Reply • Share ›



rihansh • 3 months ago

```
int allsum,allXOR;
```

```
for i<-0 to n
```

```
allsum+=arr[i];
```

```
allxor^=arr[i];
```

```
then result =(3*allxor-allsum)/2;
```

hope everybody got it if i am wrong then correct me and if anyone is having an

^ | v • Reply • Share ›



rihansh • 3 months ago

hey kindl look at the following solution this one is intersting one hope every boc

^ | v • Reply • Share ›



rihansh • 3 months ago

this can be done n the following manner in o(n)

take the sum of all numbers and allso XOR of all number the multiply the xor b
will give the answer

3 ^ | v • Reply • Share ›



Rajesh Pal • 4 months ago

/*beauty of map (stl container) ,easily can be solved in o(n) time. :P*/

```
#include <cstdlib>
```

```
#include <stdio.h>
```

```
#include <vector>
```

```
#include <string>
```

```
#include<map>
```

```
#include<iterator>
```

```
#include <iostream>

using namespace std;

int main()
{
    map<int ,int=""> mymap;
    vector<int> myvector;
    int n;
    cin>>n;
```

[see more](#)

^ | v • Reply • Share ›



Castle Age → Rajesh Pal • 3 months ago

O(1) space though.

^ | v • Reply • Share ›



Jing Conan Wang • 4 months ago

There is better way to understand the first method.

Ones and twos are two masks.

(1) If the digit i of ones is 1, the total number of 1s appeared on digit $i \% 3 == 1$

(2) If the digit i of twos is 1, the total number of 1s appeared on digit $i \% 3 == 2$.

ones stores the bits which have appeared once. $arr[i]$ store the bits that will appear once & $arr[i]$ stores the bits that become to satisfy (2).

As a result, the update rule for twos

$twos = twos | (ones \& arr[i]);$

From this analysis, we can conclude that the two methods are equivalent. The efficient way to calculate the mod.

^ | v • Reply • Share ›



groomnestle • 5 months ago

2nd answer is more straightforward and clear.

2 ^ | v • Reply • Share ›



Srinivas • 6 months ago

difference of all numbers in the array gives the answer

4 ^ | v • Reply • Share ›



Guest • 8 months ago

```
#include<stdio.h>
```

```
int partition(int *a, int s, int e);
```

```
void q_sort(int *arr, int s, int e)
```

```
{ int q;
```

```
if(s<e) {="" q="partition(arr," s,="" e);="" q_sort(arr,="" s,="" q-1);="" q_sort(arr  
partition(int="" *a,="" int="" s,="" int="" e)="" {="" int="" x;="" x="a[e];" int="" j="s  
e)="" {="" if(a[j]<x)="" {="" i="i+1;" temp="a[j];" a[j]="a[i];" a[i]="temp;" }="" j="j+1  
a[e]="temp;" return="" i+1;="" }="" int="" find(int="" *a,="" int="" n)="" {="" int=""  
1);="" for(i="0;" i<n;i++)="" printf("\t%d",a[i]);="" printf("\n");="" for(i="1;i<n;i++  
else="" if(count="0)" break;="" else="" count="0;" }="" return="" a[i-1];="" }=""  
arr[]={3," 3,="" 2,="" 3};="" int="" n="sizeof(arr)" sizeof(arr[0]);="" printf("the=""  
occurrence="" is="" %d="" \n",find(arr,="" n));="" return="" 0;="" }="">
```

^ | v • Reply • Share ›



Guest • 8 months ago

```
void q_sort(int *arr, int s, int e)
```

```
{ int q;
```

```

if(s<e) {="" q="partition(arr," s,="" e);="" q_sort(arr,="" s,="" q-1);="" q_sort(arr
partition(int="" *a,="" int="" s,="" int="" e)="" {="" int="" x;="" x="a[e];" int="" j="s
e)="" {="" if(a[j]<x)="" {="" i="i+1;" temp="a[j];" a[j]="a[i];" a[i]="temp;" }="" j="j+1
a[e]="temp;" return="" i+1;="" }="" int="" find(int="" *a,="" int="" n)="" {="" int=""
1);="" for(i="0;" i<n;i++)="" printf("\t%d",a[i]);="" printf("\n");="" for(i="1;i<n;i++
else="" if(count="0)" break;="" else="" count="0;" }="" return="" a[i-1];="" }=""
arr[]={3," 3,="" 2,="" 3};="" int="" n="sizeof(arr)" sizeof(arr[0]);="" printf("the=""
occurrence="" is="" %d="" \n",find(arr,="" n));="" return="" 0;="" }="">

```

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gargsanjay • 10 months ago

the last method is $n \log(\text{no. of bits})$

```

/* Paste your code here (You may delete these lines if not writing c

```

^ | v • Reply • Share ›



Sandeep Jain • a year ago

Could u provide a sample input for which it failed?

^ | v • Reply • Share ›



Let's Make a New India • a year ago

2nd method is not a general method....it fails in many cases.

^ | v • Reply • Share ›



Ajay Kumar • a year ago

2nd method is gd.

^ | v • Reply • Share ›



Hanish • a year ago

Can we also write

$\text{twos} = \text{twos} \wedge (\text{ones} \& \text{arr}[i]);$

instead of |

Will it affect the code in any way ??

Since, it is never possible that a bit is set in both ones and twos

^ | v · Reply · Share ›



Biren Barodia · a year ago

if I am understanding solution1 correctly, shouldn't the algo look like this:

ones: The bits that have appeared 1st time or 3rd time or 5th time.. etc.

twos: The bits that have appeared 2nd time or 4th time or 6th time.. etc.

^ | v · Reply · Share ›



Lokesh Gopu · a year ago

```
int getSingle(int arr[], int n).
```

```
{
```

```
int f = 0, s = 0, t = 0;.
```

```
for(int i=0;i<n;i++).
```

```
{.
```

```
int tempf;.
```

```
tempf = ( t & arr[i]) ;.
```

```
t = t & (t ^ arr[i]);.
```

```
t = t | ( s & arr[i]);.
```

```
s = s & (s ^ arr[i]);.
```

```
s = s | (f & arr[i]);.
```

[see more](#)

1 ^ | v • Reply • Share ›



Lokesh • a year ago

```
/* Paste your code here (You may delete these lines if not writing c  
int getSingle(int arr[], int n)  
{  
    int f = 0, s = 0, t = 0;  
    for(int i=0;i<n;i++)  
    {  
        int tempf;  
  
        tempf = ( t & arr[i]) ;  
        t = t & (t ^ arr[i]);  
  
        t = t | ( s & arr[i]);  
        s = s & (s ^ arr[i]);  
  
        s = s | (f & arr[i]);  
        f = f & (f ^ arr[i]);  
  
        f = tempf | f;  
  
        f = f | (arr[i] & (arr[i] ^ (s | t)));  
    }  
  
    return f;  
}
```

^ | v • Reply • Share ›



GeekstoriGeeks · a year ago

Please take a closer look at the problem statement. The array {15, 15, 4} does

^ | v · Reply · Share ›



Koda Rahul · a year ago

second method does not for for all cases.

eg:{15,15,4}

^ | v · Reply · Share ›



Koda Rahul · a year ago

second method doesn't work for evry input.

eg: {15,15,4}

^ | v · Reply · Share ›



swati · 2 years ago

can't we just sum numbers and take modulo by 3. Is it necessary to consider i

```
/* Paste your code here (You may delete these lines if not writing co
```

^ | v · Reply · Share ›



Parminder → swati · 2 years ago

Would fail if the number we are looking for has 3 as factor

```
/* Paste your code here (You may delete these lines if not wr
```

^ | v · Reply · Share ›



Kartik → swati · 2 years ago

This approach has been discussed earlier and doesn't work. Please s

^ | v · Reply · Share ›



HLS · 2 years ago

The approach doesn't work.

e.g. 1,1,1,3

sum=6

sum%3=0

what next?

```
/* Paste your code here (You may delete these lines if
```

^ | v · Reply · Share ›



Angel → HLS · 8 months ago

If you are referring to the second approach, then that's r

You add the corresponding bits from each number, sum

sum the third bits:

1 1 1 3 result

||||

v v v v v

0+0+0+0=0%3=0

0+0+0+1=1%3=1

1+1+1+1=4%3=1

when we put together the results 011 we get the 3 see?

^ | v · Reply · Share ›



Sandy · 2 years ago

I doubt the values of common_bit_mask during the iterations.

I think it should be -1,-1,-3,-2

Can you explain me a bit?



Sumit → Sandy · 2 years ago

Only the values of last two bits are mentioned in comments. When you get negative value because the sign bit is inverted when ~ operator

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aj · 2 years ago

Infact there is another approach. You can take the sum of bits in same position with 3. The number you get is the number that appears only once.

e.g.

101, 101, 101, 1000

sum of first bits%3 = $(1+1+1+0)\%3 = 0$;

sum of second bits%3 = $(0+0+0+0)\%3 = 0$;

sum of third bits%3 = $(1+1+1+0)\%3 = 0$;

sum of fourth bits%3 = $(1)\%3 = 1$;

hence number which appears once is 1000

2 ^ | v · Reply · Share ›



a2 → aj · 2 years ago

This approach will not work if the numerals have the digits 3 and above

Will it ?

^ | v · Reply · Share ›



Kartik → a2 · 2 years ago

I think it works for any input. If you have counter example, then I will be glad to hear it.

^ | v · Reply · Share ›



a2 → Kartik · 2 years ago

Sry , I was mistaken .. had got confused !!



Kartik → aj · 2 years ago

@aj: I have implemented your approach and it seems to be working fin

```
#include <stdio.h>
#define INT_SIZE 32
int getSingle(int arr[], int n)
{
    int x = 1;
    int result = 0;

    // Let us take the example of {3, 3, 2, 3} to understand th
    for (int i = 0; i < INT_SIZE; i++)
    {
        int count = 0;
        for (int j=0; j< n; j++ )
        {
            x = (1 << i);
            if (arr[j] & x)
                count++;
        }
    }
}
```

[see more](#)

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rohit → Kartik · 2 years ago

gud code but consider array of size = $O(n^2)$

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Kartik → rohit · 2 years ago

@rohit: I could not understand your point, please elabor

^ | v · [Reply](#) · [Share](#) ›



Kartik → Kartik · 2 years ago

Thanks for replying rohit.

When we talk about time complexity, we always talk about
So it is valid to say the time complexity of the algo is $O(n)$

The character for less than operator is considered as a

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rohit → Kartik · 2 years ago

again same thing, consider array of size less than and ϵ

^ | v · Reply · Share ›



rohit → Kartik · 2 years ago

@kartik-site error, only half comment came. i was saying

^ | v · Reply · Share ›



Kartik → aj · 2 years ago

@aj: This solution looks good to be added to the original post. Could you

^ | v · Reply · Share ›



aj → Kartik · 2 years ago

I will try to write down the code today if possible...But I have a

```
/* Paste your code here (You may delete these lines if
```

1 ^ | v · Reply · Share ›



angel · 2 years ago

sum of array of integers will always be in form of $3(a+b+c)+d$

so $\text{sum} \% 3$ always returns the number which is occurring only once, as per my u

^ | v • Reply • Share ›



Kartik → angel • 2 years ago

@angel: Thanks for sharing your thoughts. I think this will give the required number Please correct me if I am wrong.

^ | v • Reply • Share ›



angel → Kartik • 2 years ago

yaa you r right

```
/* Paste your code here (You may delete these lines if
```

^ | v • Reply • Share ›



aj → angel • 2 years ago

what if d is divisible by 3??

^ | v • Reply • Share ›



rohit • 2 years ago

1.add bitwise xor of all elements of the array to a variable

2.do step 1 3 times

3.value of the variable is our answer

correct me if i m wrong :)

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