

## Merge Overlapping Intervals

Given a set of time intervals in any order, merge all overlapping intervals into one and output the result which should have only mutually exclusive intervals. Let the intervals be represented as pairs of integers for simplicity.

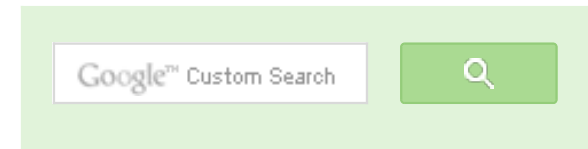
For example, let the given set of intervals be  $\{\{1,3\}, \{2,4\}, \{5,7\}, \{6,8\}\}$ . The intervals  $\{1,3\}$  and  $\{2,4\}$  overlap with each other, so they should be merged and become  $\{1, 4\}$ . Similarly  $\{5, 7\}$  and  $\{6, 8\}$  should be merged and become  $\{5, 8\}$

Write a function which produces the set of merged intervals for the given set of intervals.

A **simple approach** is to start from the first interval and compare it with all other intervals for overlapping, if it overlaps with any other interval, then remove the other interval from list and merge the other into the first interval. Repeat the same steps for remaining intervals after first. This approach cannot be implemented in better than  $O(n^2)$  time.

An **efficient approach** is to first sort the intervals according to starting time. Once we have the sorted intervals, we can combine all intervals in a linear traversal. The idea is, in sorted array of intervals, if interval[i] doesn't overlap with interval[i-1], then interval[i+1] cannot overlap with interval[i-1] because starting time of interval[i+1] must be greater than or equal to interval[i]. Following is the detailed step by step algorithm.

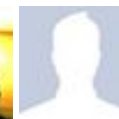
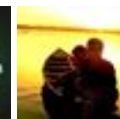
1. Sort the intervals based on increasing order of starting time.
2. Push the first interval on to a stack.
3. For each interval do the following
  - a. If the current interval does not overlap with the stack top, push it.
  - b. If the current interval overlaps with stack top and ending time of current interval is more than that of stack top, update stack top with the ending time of current interval.



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4. At the end stack contains the merged intervals.

Below is a C++ implementation of the above approach.

```
// A C++ program for merging overlapping intervals
#include <iostream>
#include <vector>
#include <algorithm>
#include <stack>
using namespace std;

// An interval has start time and end time
struct Interval
{
    int start;
    int end;
};

// Compares two intervals according to their starting time.
// This is needed for sorting the intervals using library
// function std::sort(). See http://goo.gl/iGspV
bool compareInterval(Interval i1, Interval i2)
{
    return (i1.start < i2.start)? true: false;
}

// The main function that takes a set of intervals, merges
// overlapping intervals and prints the result
void mergeIntervals(vector<Interval>& intervals)
{
    // Test if the given set has at least one interval
    if (intervals.size() <= 0)
        return;

    // Create an empty stack of intervals
    stack<Interval> s;

    // sort the intervals based on start time
    sort(intervals.begin(), intervals.end(), compareInterval);

    // push the first interval to stack
    s.push(intervals[0]);

    // Start from the next interval and merge if necessary
    for (int i = 1 ; i < intervals.size(); i++)
    {
        // get interval from stack top
        Interval top = s.top();
```



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```

// if current interval is not overlapping with stack top,
// push it to the stack
if (top.end < intervals[i].start)
{
    s.push( intervals[i] );
}
// Otherwise update the ending time of top if ending of current
// interval is more
else if (top.end < intervals[i].end)
{
    top.end = intervals[i].end;
    s.pop();
    s.push(top);
}
}

// Print contents of stack
cout << "\n The Merged Intervals are: ";
while (!s.empty())
{
    Interval t = s.top();
    cout << "[" << t.start << "," << t.end << "]" << " ";
    s.pop();
}

return;
}

```

// Functions to run test cases

```

void TestCase1()
{
    // Create a set of intervals
    Interval intvls[] = { {6,8}, {1,9}, {2,4}, {4,7} };
    vector<Interval> intervals(intvls, intvls+4);

    // Merge overlapping intervals and print result
    mergeIntervals(intervals);
}

void TestCase2()
{
    // Create a set of intervals
    Interval intvls[] = { {6,8}, {1,3}, {2,4}, {4,7} };
    vector<Interval> intervals(intvls, intvls+4);

    // Merge overlapping intervals and print result
    mergeIntervals(intervals);
}

```



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**newCoder3006** Code without using while

```
void TestCase3()
{
    // Create a set of intervals
    Interval intvls[] = { {1,3},{7,9},{4,6},{10,13} };
    vector<Interval> intervals(intvls, intvls+4);

    // Merge overlapping intervals and print result
    mergeIntervals(intervals);
}

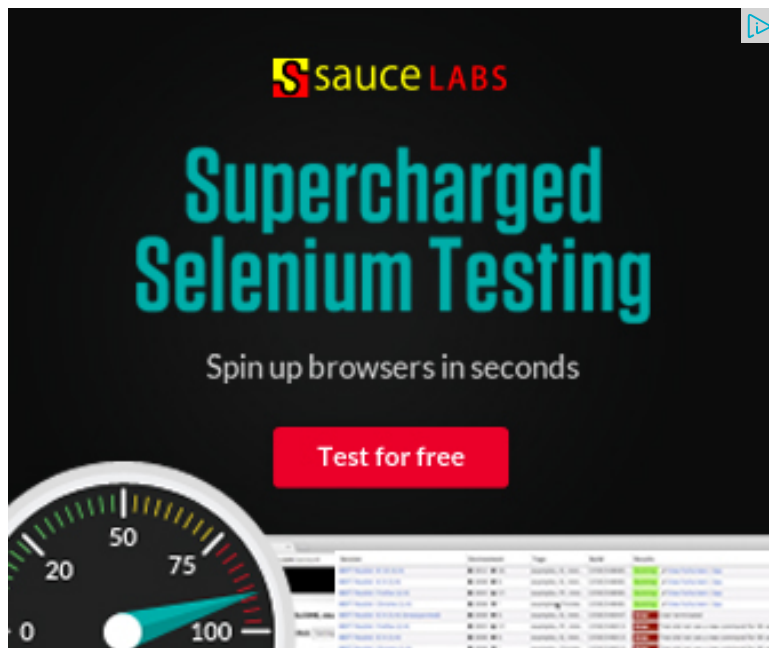
// Driver program to test above functions
int main()
{
    TestCase1();
    TestCase2();
    TestCase3();
    return 0;
}
```

Output:

```
The Merged Intervals are: [1,9]
The Merged Intervals are: [1,8]
The Merged Intervals are: [10,13] [7,9] [4,6] [1,3]
```

Time complexity of the method is  $O(n \log n)$  which is for sorting. Once the array of intervals is sorted, merging takes linear time.

This article is compiled by Ravi Chandra Enaganti. Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.



loop. We can do it...

Find subarray with given sum · 1 hour ago

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with the algorithm...



**Nilesch Kumar** · 3 days ago

Hi Ravi,

Instead of using vector and stack, you can use a `std::list`. Internally, this contains delete in between from a list rather than a vector. So, no stack is required, if you

If interested, you can view the same here : <http://ideone.com/m11IDb>

Regards,

Nilesch

^ | v · Reply · Share ›



**Sreek** · a month ago

It has to be

if  $(\text{top.end} + 1 < \text{intervals}[i].\text{start})$

to support the use case [1,2] [3,4] to give the output [1,4].

^ | v · Reply · Share ›



**anonym2** · 5 months ago

Another solution: A min heap of the pairs with respect to the start element can give  $O(n \log n)$  solution.

2 ^ | v · Reply · Share ›



**Deepak Shrivastava** · 7 months ago

I wrote a java implementation of the problem. Please find it below...

```
import java.util.Arrays;
```

```
import java.util.Comparator;
```

```
public class Solution {
```

```
    public static void main(String[] args) {
```

```
        Integer arr[][] = { { 1, 3 }, { 6, 8 }, { 2, 4 }, { 5, 7 } };
```

```
        Arrays.sort(arr, new Comparator<integer[]>() {
```

```
            @Override
```

```
            public int compare(Integer[] o1, Integer[] o2) {
```

```
                if (o1[0] > o2[0])
```

```
                    return 1;
```

[see more](#)

^ | v • Reply • Share ›



**Sriharsha g.r.v** • 7 months ago

can some one

explain me how is the complexity  $O(n \log n)$ ..what is the worst case complexity i.e when neither of intervals overlap? is the complexity  $n \log n$  occurring because of sorting the intervals??

^ | v • Reply • Share ›



**zzet** → Sriharsha g.r.v • 17 days ago

yeah, the  $O(N \log N)$  is based on sorting intervals.

^ | v • Reply • Share ›



**Guest** • 7 months ago

can some explain me how is the complexity  $O(n \log n)$ ..what is the worst case overlap? is the complexity  $n \log n$  occurring because of sorting the intervals??

^ | v • Reply • Share ›



**Guest** • 7 months ago

can some explain me how is the complexity  $O(n \log n)$ ..what is the worst case overlap?

^ | v • Reply • Share ›



**Ankit Chaudhary** • 7 months ago

Variation of above problem: suppose intervals are coming one by one and u have time when new interval comes.

Means at every instant u have to show merged intervals.

This can be implemented using doubly linked list.

1 ^ | v • Reply • Share ›



**DRAGONWARRIOR** • 8 months ago

IN BODY OF ELSE IF

SINCE END IS UPDATED FOR TOP ELEMENT(top()) function return ref to top  
WHY WE NEED POP STATEMENT

^ | v • Reply • Share ›



**rakitic** • 8 months ago

```
#include<iostream>
```

```
#include<conio.h>
```

```
#include<stack>
```

```
#include<vector>
```

```
#include<algorithm>
```



```
using namespace std;
```

```
struct interval{
```

```
    int start;
```

```
    int end;
```

[see more](#)

^ | v • Reply • Share ›



**rakitic** • 8 months ago

```
#include<iostream>
```

```
#include<conio.h>
```

```
#include<stack>
```

```
#include<vector>
```

```
#include<algorithm>
```

```
using namespace std;
```

```
struct interval{
```

```
int start;
```

```
int end;
```

```
};
```

```
stack<struct interval="">S;
```

[see more](#)

^ | v • Reply • Share ›



**Guest** • 8 months ago

this can be done in  $O(n \lg n)$  by building a segment tree.

^ | v • Reply • Share ›



**Neha Modi** • 9 months ago

how about this....

```
#include<iostream.h>
```

```
#include<stdio.h>
```

```
#include<conio.h>
```

```
void main()
```

```
{ int arr[4][2]={6,8},{1,9},{2,4},{4,7}}, i,j, temp, temp1;.
```

```
for(i=0;i<4;i++).
```

```
{ for(j=i+1;j<4;j++).
```

```
{ if(arr[j][0]<arr[i][0]).
```

```
{ temp=arr[j][0];.
```

```
temp1=arr[j][1];.
```

```
arr[j][0]=arr[i][0];.
```

```
arr[j][1]=arr[i][1];.
```

```
arr[i][0]=temp;.
```

```
arr[i][1]=temp1;.
```

```
}.  
_____
```

[see more](#)

^ | v • Reply • Share ›



**Cracking The Code** • a year ago

how about this.

start loop

if (top.end >= intervals[i].start) // 1,9, 2,4.

{.

top.end = top.end > intervals[i].end ? top.end : intervals[i].end;

s.pop();.

s.push(top);.

}.  
}

else if (top.end < intervals[i].start).

s.push(intervals[i]);.

[see more](#)

^ | v • Reply • Share ›



nikhil • a year ago

// assuming i have a sort function to sort the interval array

```
struct Interval{
    int first;
    int second;
};

void merge(struct Interval a[], int n){
    int i=0, j;
    sort(a,n);
    while(i<n){
        j=i;
```

```

while( a[j].second > a[j+1].first )
j++;
cout<<"["<<a[i].first<<" "<<a[j].second<<"]"<<endl;
i=j+1;
}
}

```

Time complexity will be  $O(n \log n(\text{sort}) + n(\text{merge}))$ .

Space complexity  $O(1)$ .

^ | v • Reply • Share ›



nikhil • a year ago

//assuming i have a sort function to sort the interval array

```

struct Interval{
int first;
int second;
};
void merge(struct Interval a[],int n){
int i=0,j;
sort(a,n);
while(i < a[j+1].first)
j++;
cout<<"["<<a[i].first<<" "<<a[j].second<<"]"<<endl;
i=j+1;
}
}

```

Time complexity will be  $O(n \log n(\text{sort}) + n(\text{merge}))$ .

Space complexity  $O(1)$ .

^ | v • Reply • Share ›



**Mani** • a year ago

I tried to implement in java.

```
import java.util.LinkedHashMap;
import java.util.Map;
import java.util.Set;
import java.util.TreeMap;

public class MergeIntervals {

    /**
     * @param args
     */
    public static void main(String[] args) {
        // TODO Auto-generated method stub
        int a[][] =new int[][]{ {2,3}, {7,9}, {4,6}, {1,13}};
        //int a[][] =new int[][]{{6,8}, {1,9}, {2,4}, {4,7}};
        //int a[][] =new int[][]{{1,3},{4,6},{5,10}};
        //int a[][] =new int[][]{{1,4},{2,4},{5,5},{6,8}};
        int b[][] =new int[4][2]:
```

[see more](#)

^ | v • Reply • Share ›



**Bamftubb** • a year ago

Or another shorter one:

```
var sorted = ranges.OrderBy(r => r[0]).ToArray();
var result = new LinkedList<int[]>();
result.AddFirst(sorted[0]);
foreach (var range in sorted) {
```

```
var top = result.Last.Value,
if (range[0] <= top[1])
    top[1] = top[1] < range[1] ? range[1] : top[1];
else
    result.AddLast(range);
}
return result.ToArray();
```

^ | v • Reply • Share ›



**Bamftubb** • a year ago

Here C# variant without using stack. It uses List for collecting ranges.

[sourcecode language="C#"]

```
class Program
{
    static void Main(string[] args)
    {
        var given = new int[][] {
            new[] { 0, 5 },
            new[] { -1, 5 },
            new[] { 6, 10 },
            new[] { 3, 4 },
            new[] { 15, 20 },
            new[] { -10, 6 }
        };
        Console.WriteLine("Given: [{0}]", string.Join(", ", given.Select(r => string.For
        Console.WriteLine("Result: [{0}]", string.Join(", ", Merge(given).Select(r => st
    }
```

[see more](#)

^ | v • Reply • Share ›



**Daddy** · a year ago

Here is implementation in C :

```
#include<iostream>
```

```
#include<stdio.h>
```

```
#include<conio.h>
```

```
#include<stdlib.h>
```

```
using namespace std;
```

```
struct interval{
```

```
int start;
```

```
int end;
```

```
};
```

```
struct stack{
```

```
struct interval* arr;
```

```
int top;
```

```
};
```

```
struct stack* createstack(int n)
```

```
{
```

[see more](#)

^ | v · Reply · Share ›



**Sambasiva** · a year ago

<http://effprog.wordpress.com/2...>

^ | v · Reply · Share ›



**Anon** · a year ago

Here is another implementation in C++.

```
/* Paste your code here (You may delete these lines if not writing c
```

```
#include <iostream>
```

```

#include <algorithm>
#include <vector>
#include <stack>
using namespace std ;

void print ( vector <pair<int,int> > v )
{
    vector <pair<int,int> > :: iterator it ;
    for ( it = v.begin() ; it != v.end() ; ++it )
        cout << (*it).first << " " << (*it).second << endl ;
    return ;
}

int main()

```

[see more](#)

^ | v • Reply • Share ›



**Gupt** • a year ago

An  $O(n)$  in time and  $O(n)$  in memory approach could be this, where  $n$  is the high

```

    create an array of size n, say sets
    fill it with -1

    current setnb = 1;
    for (each interval)
    {
        if (set[interval.start == -1])
            currentmarker = setnb++;
        else
            currentmarker = set[interval.start]
    }

```



```
    for (i = interval.start to interval.end)
    {
        sets[i] = currentmarker;
    }
}
```

now the sets [array](#) contains different interval, we dont need any sort:  
just extract those intervals by traversing in [array](#).

^ | v • Reply • Share ›



**vishal** • a year ago

Please correct me if i'm wrong but, what will be the output for {{6,8},{1,3},{2,4}{1,9}}.....but as per current code it will be {1,8}.

^ | v • Reply • Share ›



**GeeksforGeeks** → vishal • a year ago

@vishal: Thanks for pointing this out. The post is updated to handle thi:

^ | v • Reply • Share ›



**aasshishh** • a year ago

The above code has a bug.

I changed the inputs.

Test Case 1: {6,8}, {1,9}, {2,4}, {4,7}

Test Case 2: {2,3}, {7,9}, {4,6}, {1,13}

Output were:

Test case 1 : [1,8]

Outputs should be :

Test case 1 : [1,9]

Test case 2 : [1,13]

The code should be as follows

```
if (top.end < intervals[i].start)
{
    s.push( intervals[i] );
}
```

[see more](#)

^ | v • Reply • Share ›



**GeeksforGeeks** → aasishh • a year ago

@aasishh: Thanks for pointing this out and suggesting the resolution

^ | v • Reply • Share ›



**aasishh** • a year ago

There is a bug in the program.

I changed the inputs.

Test Case 1: {6,8}, {1,9}, {2,4}, {4,7}

Test Case 2: {2,3}, {7,9}, {4,6}, {1,13}

Output were:

Test case 1 : [1,8]

Test case 2 : [7,9] [4,6] [1,3]

Outputs should be :

Test case 1 : [1,9]

Test case 2 : [1,13]

The following code has bug.

```
if (top.end < intervals[i].start)
{
    s.push( intervals[i] );
}
```

[see more](#)

^ | v • Reply • Share ›



**Ashwini Chaudhary** • a year ago

Implementation in python:

```
[sourcecode language="Python"]
from stack import Stack
from operator import itemgetter
lis=[(1,3),(2,4),(5,7),(6,8)]
lis.sort(key=itemgetter(0))
s=Stack()
s.push(lis[0])
for x,y in lis[1:]:
    topx,topy=s.top()
    if topx<x<=topy:
        s.pop()
        s.push((topx,y))
    else:
        s.push((x,y))
print s.stack
```

^ | v • Reply • Share ›



**Ashwini Chaudhary** → Ashwini Chaudhary • a year ago

Edit: the push should be

```
s.push((min(x,topx),max(y,topy)))
```

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**Somu** • a year ago

Good work Ravi Chandra!

^ | v • Reply • Share ›



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