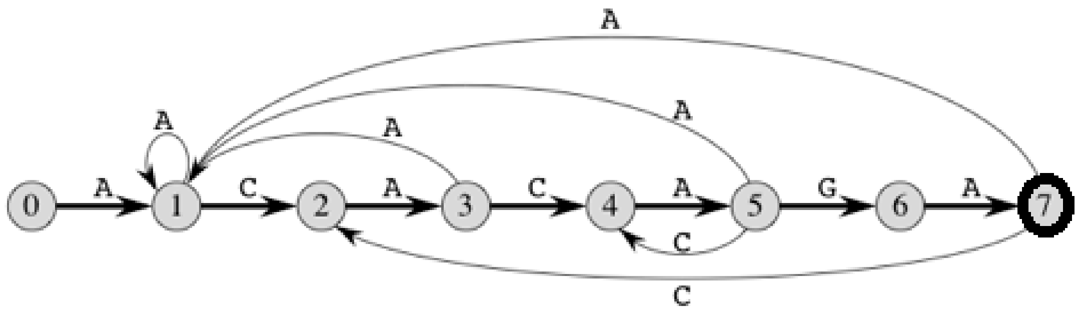
Pattern Searching | Set 6 (Efficient Construction of Finite Automata)

In the previous post, we discussed Finite Automata based pattern searching algorithm. The FA (Finite Automata) construction method discussed in previous post takes O((m^3)*NO_OF_CHARS) time. FA can be constructed in O(m*NO_OF_CHARS) time. In this post, we will discuss the O(m*NO_OF_CHARS) algorithm for FA construction. The idea is similar to lps (longest prefix suffix) array construction discussed in the KMP algorithm. We use previously filled rows to fill a new row.



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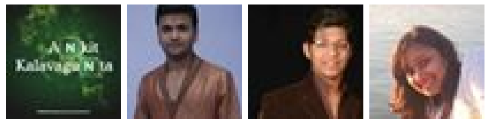

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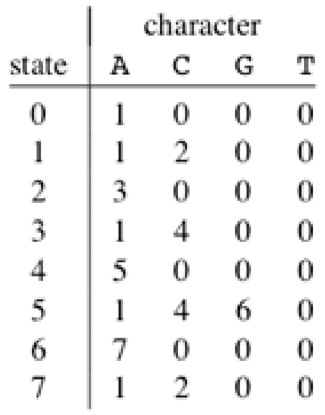

	character			
state	A	C	G	T
0	1	0	0	0
1	1	2	0	0
2	3	0	0	0
3	1	4	0	0
4	5	0	0	0
5	1	4	6	0
6	7	0	0	0
7	1	2	0	0

The abvoe diagrams represent graphical and tabular representations of pattern ACACAGA.

Algorithm:
1) Fill the first row. All entries in first row are always 0 except the entry for pat[0] character. For pat[0] character, we always need to go to state 1.
2) Initialize lps as 0. lps for the first index is always 0.
3) Do following for rows at index i = 1 to M. (M is the length of the pattern)
…..a) Copy the entries from the row at index equal to lps.
…..b) Update the entry for pat[i] character to i+1.
…..c) Update lps "lps = TF[lps][pat[i]]" where TF is the 2D array which is being constructed.

Implementation
Following is C implementation for the above algorithm.

```
#include<stdio.h>
#include<string.h>
#define NO_OF_CHARS 256

/* This function builds the TF table which represents Finite Automata
   given pattern  */
void computeTransFun(char *pat, int M, int TF[][NO_OF_CHARS])
{
    int i, lps = 0, x;

    // Fill entries in first row
    for (x =0; x < NO_OF_CHARS; x++)
       TF[0][x] = 0;
    TF[0][pat[0]] = 1;

    // Fill entries in other rows
```

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```c
    for (i = 1; i<= M; i++)
    {
        // Copy values from row at index lps
        for (x = 0; x < NO_OF_CHARS; x++)
            TF[i][x] = TF[lps][x];

        // Update the entry corresponding to this character
        TF[i][pat[i]] = i + 1;

        // Update lps for next row to be filled
        if (i < M)
          lps = TF[lps][pat[i]];
    }
}

/* Prints all occurrences of pat in txt */
void search(char *pat, char *txt)
{
    int M = strlen(pat);
    int N = strlen(txt);

    int TF[M+1][NO_OF_CHARS];

    computeTransFun(pat, M, TF);

    // process text over FA.
    int i, j=0;
    for (i = 0; i < N; i++)
    {
       j = TF[j][txt[i]];
       if (j == M)
       {
         printf ("\n pattern found at index %d", i-M+1);
       }
    }
}

/* Driver program to test above function */
int main()
{
    char *txt = "GEEKS FOR GEEKS";
    char *pat = "GEEKS";
    search(pat, txt);
    getchar();
    return 0;
}
```

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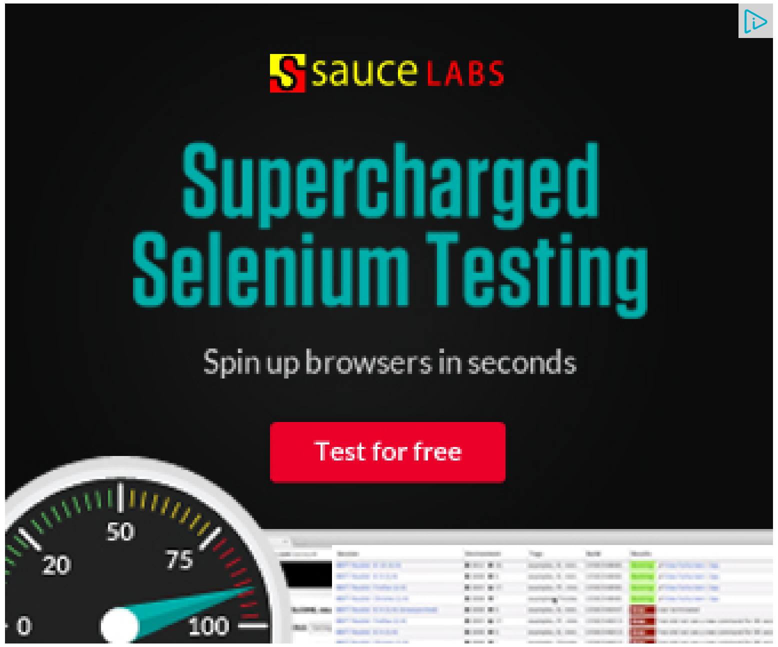
Output:

```
pattern found at index 0
pattern found at index 10
```

Time Complexity for FA construction is O(M*NO_OF_CHARS). The code for search is same as the previous post and time complexity for it is O(n).

Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.



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Guest • 6 months ago

This code runs in O(n) time with constant space ...please correct me if it is inc

int position[26] = {0};

int delta(int state,char input,char *fa,int i){
if(fa[state] == input){
position[input - 97] = state;
return state+1;
}
return position[input-97];
}
int pattern_match(char *str,char *p,int n,int m){
int q = 0;
for(int i = 0 ; i < n; i++){
int old_q = q;
q = delta(q,str[i],p,i);
if(old_q != 0 && old_q - q > 1)
i--;
else if(q == m){
cout<<"pattern found at : "<<i-m+1<<endl; q="0;" }="" }="" }="" int="" main(){='

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char="" pattern[]="geeks" ;="" pattern_match(str,="" pattern,="" sizeof(str)="" s
sizeof(char)="" -1);="" return="" 0;="" }="">

alien · 9 months ago
@GeeksforGeeks: Could you please explain why this algorithm is able to fill TI
at a given time?
1

Dhiren · 10 months ago
Consider this example
Pattern – A C A C
At state-0, we have only "A", so lps = 0
Transition from state-0 to state-1, probable cases may be
Case-1 a new 'A' comes, then we go back to our longest prefix suffix till now w
what if a 'A' comes, in this case it is 1
Case-2 a new 'G' comes, then we go back to our longest prefix suffix till now a
will be 0
Case-3 a new 'C' comes, then also value will be 0
That's why we are first copying the lps row values into the current ith row.
Then we update the state transition for patt[i] in this case for 'C' state will be 2.
Then we calculate the current lps value, that is "AC" but still lps =0 as there is

Calculation of lps can be clear from state transition-2 to 3.
Current lps is 0, now 'A' comes so that new lps is 1 for "ACA" which can be fo

Arvind · 11 months ago
How does this algorithm work ? where is the proof for this ?
1

Ram · 11 months ago

Wher is the proof ????????????? please post the proof

1 ^ | ⌄ •

Akshay khare • a year ago

when i will become equal to M
then TF[i][pat[i]] = i+1. will give segmentation fault
since pat has length M and its index can be upto M-1 only
how will TF[M][pat[M]] will work..pat[M] -> this location not exists..pls explain h
me if i am wrong...

2 ^ | ⌄ •

Shiwakant Bharti ➦ Akshay khare • 9 months ago

Akshay khare Nice findings! I got the same error in Java.

Here is the test case which should break:

```
          char[] txt2 = "AABAACAADAABAAABAA".toCharArray();
          char[] pat2 = "AABA".toCharArray();

//This code fix worked for me. Not sure if this robust enough.
          for (i = 1; i <= M; i++) {
               // Copy values from row at index lps
               // Is this powerful enough to handle case of i = M
               // transition)?
               for (ch = 0; ch < NoOfChars; ch++) {
                    TFDP[i][ch] = TFDP[lps][ch];
               }
               // This is special case where the last halt state
               // for regular processing.
               // Here pat[M] is out of bound and further calcula
               //
```

see more

abhishek08aug • a year ago

Intelligent :D

vinu • 2 years ago

Yes... @GeeksForGeeks can you please me more clear with reasoning of ste

```
/* Paste your code here (You may delete these lines if not writing c
```

spandan • 2 years ago

If someone can please explain how has the transition function been computed

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