

Maximum sum such that no two elements are adjacent

Question: Given an array of positive numbers, find the maximum sum of a subsequence with the constraint that no 2 numbers in the sequence should be adjacent in the array. So 3 2 7 10 should return 13 (sum of 3 and 10) or 3 2 5 10 7 should return 15 (sum of 3, 5 and 7). Answer the question in most efficient way.

Algorithm:

Loop for all elements in arr[] and maintain two sums incl and excl where incl = Max sum including the previous element and excl = Max sum excluding the previous element.

Max sum excluding the current element will be max(incl, excl) and max sum including the current element will be excl + current element (Note that only excl is considered because elements cannot be adjacent).

At the end of the loop return max of incl and excl.

Example:

```
arr[] = {5, 5, 10, 40, 50, 35}
```

```
inc = 5
```

```
exc = 0
```

```
For i = 1 (current element is 5)
```

```
incl = (excl + arr[i]) = 5
```

```
excl = max(5, 0) = 5
```

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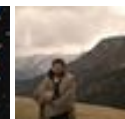
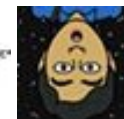
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```
For i = 2 (current element is 10)
incl = (excl + arr[i]) = 15
excl = max(5, 5) = 5
```

```
For i = 3 (current element is 40)
incl = (excl + arr[i]) = 45
excl = max(5, 15) = 15
```

```
For i = 4 (current element is 50)
incl = (excl + arr[i]) = 65
excl = max(45, 15) = 45
```

```
For i = 5 (current element is 35)
incl = (excl + arr[i]) = 80
excl = max(5, 15) = 65
```

And 35 is the last element. So, answer is $\max(\text{incl}, \text{excl}) = 80$

Thanks to [Debanjan](#) for providing code.

Implementation:

```
#include<stdio.h>

/*Function to return max sum such that no two elements
are adjacent */
int FindMaxSum(int arr[], int n)
{
    int incl = arr[0];
    int excl = 0;
    int excl_new;
    int i;

    for (i = 1; i < n; i++)
    {
        /* current max excluding i */
        excl_new = (incl > excl)? incl: excl;

        /* current max including i */
```



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```

        incl = excl + arr[i];
        excl = excl_new;
    }

    /* return max of incl and excl */
    return ((incl > excl)? incl : excl);
}

/* Driver program to test above function */
int main()
{
    int arr[] = {5, 5, 10, 100, 10, 5};
    printf("%d \n", FindMaxSum(arr, 6));
    getchar();
    return 0;
}

```

Time Complexity: $O(n)$

Now try the same problem for array with negative numbers also.

Please write comments if you find any bug in the above program/algorithm or other ways to solve the same problem.

4th Grade Common Core



Easy to Use; Fun for Kids; Math & ELA; Practice & Instruction



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11

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3

Writing code in comment? Please use ideone.com and share the link here.

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Venu Gopal · a month ago

Recursive version as in other DP problems: just the function is changed
<http://ideone.com/rdbd9A>

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rainhacker → Venu Gopal · 8 days ago

What is the $O()$ complexity of your solution ?

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705



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rainhacker → rainhacker · 8 days ago

It is definitely more than $O(n)$

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Mohamed Abdul Rahim · a month ago

How about this?

```
public static int getMax(int[] A) {
    int[] mArray = new int[A.length];
    int maxSoFar = 0;
    for (int i = 0; i < A.length; i++) {
        int m1 = (i - 2 > 0) ? A[i] + mArray[i - 2] : A[i];
        int m2 = (i - 3 > 0) ? A[i] + mArray[i - 3] : A[i];
        mArray[i] = max(m1, m2);
        if (mArray[i] > maxSoFar) {
            maxSoFar = mArray[i];
        }
    }
    return maxSoFar;
}
```

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Aditya Murgai · a month ago

// a recursive approach

```
#include<stdio.h>
```

```
/*Function to return max sum such that no two elements
are adjacent */
```

```
int mymax(int a,int b)
```

```
{
```

```
return a>b?a:b;
```

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```
}
```

```
int FindMaxSum(int arr[], int pos, int sum,int n)
```

```
{
```

```
if(pos>=n)
```

```
return sum;
```

```
else
```

```
return mymax(FindMaxSum(arr,pos+1,sum,n),FindMaxSum(arr,pos+2,sum+arr[pos],n));
```

```
.....
```

[see more](#)

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Venu Gopal → Aditya Murgai • a month ago

saw your code now only, just before I was going to post this same app
my code: <http://ideone.com/rdbd9A>

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Vijay Apurva • 2 months ago

for -ve number we can use the same approach

first we replace all the -ve numbers with 0 . after this we can apply this apprao

^ | v • Reply • Share ›



newCoder • 2 months ago

Here is the code which works on all cases positive and negative or mix:

```
private static int maxNonAdjacentSum(int a[]) {
```

```
if (a.length == 1)
```

```
return a[0];
```

```
if (a.length == 2)
```

```
return max(a[0], a[1]);
```

```

int secondLast = a[0];
int last = max(secondLast, a[1]);
int current = last;

for (int i = 2; i < a.length; i++) {
    current = max(a[i], max(secondLast + a[i], last));
    secondLast = last;
    last = current;
}
return current;
}

```

6 ^ | v • Reply • Share ›



alien • 3 months ago

awesome algo dude

^ | v • Reply • Share ›



xiveman • 3 months ago

Can you explain why we need two different arrays? Why not use only one array such sum with a[i] included:

```

public static int maxSum(int[] a){
    if(a == null || a.length == 0) return 0;
    if(a.length == 1) return a[0];
    if(a.length == 2) return (a[0] > a[1] ? a[0] : a[1]);

    int[] M = new int[a.length];
    M[0] = a[0]; M[1] = a[1];
    int max = 0;
    for(int i = 2; i < a.length; i++){
        M[i] = M[i-2];
        if(i-3 >= 0 && M[i-3] > M[i-2]) M[i] = M[i-3];
    }
}

```

```

        M[i] += a[i];
        max = (max > M[i] ? max : M[i]);
    }
    return max;
}

```

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HRISHIKESH • 3 months ago

//recursive code of above problem

```

#include <iostream>
using namespace std;
int getmaxsum(int a[],int size)
{
    if (size>=2) {
        int temp=getmaxsum(a,size-1);
        int temp2=getmaxsum(a,size-2);
        return temp2 +a[size]>temp?temp2+a[size]:temp;
    }
    else if (size==1)
        return a[0]>a[1]?a[0]:a[1];
    else return a[0];
}
int main () {
    int array[] = {3,8,12,6,2,34,4,19,7,9,11};
    cout<<getmaxsum(array,10); return="" 0;="" }="">

```

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skmahawar • 3 months ago

@orchidmajumder some modification for case of -ve numbers. please comn

[Language : Java]


```
import java.io.*;
public class Program{

    public static void main(String[] args) throws IOException{
        int input[] = {-3,-2,-1,-10};
        int sumUpto[] = new int[4];
        sumUpto[0] = input[0];
        sumUpto[1] = Math.max(input[0],sumUpto[0]);
        for(int i = 2 ; i<4 ; i++){
            sumUpto[i] = Math.max(input[i],Math.max(input[i]+sumUpto[i-2],sumUpto[i-1]))

        }
        System.out.println(sumUpto[3]);

    }

}
```

1 ^ | v • Reply • Share ›



guest • 6 months ago

No need to use DP...A very simple approach..

Just a little modification in above code...

It wont work if all are -ve..we can have one pre check...please let me know if it

```
#include<stdio.h>
#include<conio.h>
/*Function to return max sum such that no two elements
are adjacent */
int max(int a,int b)
{
    if(a>b)
        return a;
    else
```

```
return b;
}
int FindMaxSum(int arr[], int n)
{
int incl = arr[0];
```

[see more](#)

^ | v • Reply • Share ›



Abhay • 6 months ago

//work for negative number also

```
int main()
{
int i,j,sum1=0,sum2=0;
int arr[]={5,5,10,40,50,-35};
int n=sizeof(arr)/sizeof(arr[0]);
for(i=0,j=1;j<=n;i+=2,j+=2)
{
sum1=sum1+arr[i];
sum2=sum2+arr[j];
}
if(sum1<sum2) printf("%d",sum2);="" else="" printf("%d",sum1);="" }="">
```

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zorro • 6 months ago

Very poorly written article.... with complex and probably incorrect solution a..th
have better solutions !!!

1 ^ | v • Reply • Share ›



Garrick → zorro • 5 months ago

Agree. Which solutions below do you feel are better?

Algorithm (2 paragraphs): Contradict each other. Are we excluding the

Example: Is very poor, starting off with duplicate values. eg. Which 5 is

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zorro → Garrick • 5 months ago

I feel the DP solution provided by shek8034 is the best solution.

^ | v • Reply • Share ›



Amit • 6 months ago

Works for -ve values too:

```
#include<stdio.h>

int max(int a,int b)
{
    if(a>=b)
        return a;
    return b;
}

int main()
{
    int a[]={-3 , -2 , -1 , -10};
    int n=4,i,m;
    int f[10]={0};
    f[0]=a[0];
    f[1]=max(a[1],a[0]);
    for(i=2;i<n;++i) {="" if(max(f[i-2],a[i])="">f[i-2]+a[i])
        m = max(f[i-2],a[i]);
    else
        m = f[i-2]+a[i];
    f[i]=m;
}
```

```

        ... = f[i-1] * f[i+1],
        f[i]=max(m, f[i-1]);
    }
    printf("%d\n", f[n-1]);
    return 0;
}

```

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HSIRIHS • 7 months ago

Better way : I don't get the above solution but it's very simple if take maximum elements at even positions in the array - alternatively. No need to remember a

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Shreyans → HSIRIHS • 3 months ago

It won't give correct answer when negative numbers are also included.

^ | v • Reply • Share ›



Gunni → HSIRIHS • 6 months ago

Then solve this: list = { 1, 0, 0, 1 }

3 ^ | v • Reply • Share ›



Pooja → Gunni • 2 months ago

why hsirihs approach is wrong?? plz explain me

^ | v • Reply • Share ›



draganwarrior • 8 months ago

Does this algo handle -ve value also

^ | v • Reply • Share ›



magician.trilok • 8 months ago

```
#include <stdio.h>
```

```

int FindMaxSum(int arr[], int n)
{
    int i,a,b,c;

    a=b=c=0;

    for(i=0;i<n;++i)
    {
        c=arr[i];
        c=( (a+c) > b ) ? (a+c) : b;
        a=b;    b=c;
    }
    return c;
}

int main()

```

see more

^ | v • Reply • Share ›



Anish Singhania • 9 months ago

```
#include<stdio.h>
```

```
#include<iostream>
```

```
using namespace std;
```

```
int maxSumNonAdjacent( int a[ ], int size )
```

```
{
```

```
int excl = 0, excl1 = 0;
```

```
int excl_new, excl_new1;
```

```
int incl = a[ 0 ];
```

```
int incl1 = a[ 1 ];
```

```
for( int i = 2; i < size - 1; i++ )
```

```
{
excl_new = ( incl > excl ) ? incl : excl;
excl_new1 = ( incl1 > excl1 ) ? incl1 : excl1;
incl = excl + a[ i ];
incl1 = excl1 + a[ i + 1 ];
excl = excl_new;
excl1 = excl_new1;
```

[see more](#)

^ | v • Reply • Share ›



Ankur • 10 months ago

```
#include
#include
#define SIZE 6
int check(int *a,int size,int i,int sum)
{
if(sizesum2 ?sum1:sum2);
}
main()
{
int a[SIZE]={3, 2, 5, 10, 7};
int sum,sum1;
sum=check(a,SIZE-1,0,0);
sum1=check(a,SIZE -1,1,0);
printf("maximum sum is %d",sum>sum1?sum:sum1);
getch();
}
```

^ | v • Reply • Share ›



Ankur → Ankur • 10 months ago

```
#include
```

```

#include
#define SIZE 6
int check(int *a,int size,int i,int sum)
{
    if(sizesum2 ?sum1:sum2);
}
main()
{
    int a[SIZE]={5, 5, 10, 40, 50, 35};
    int sum,sum1;
    sum=check(a,SIZE-1,0,0);
    sum1=check(a,SIZE -1,1,0);
    printf("maximum sum is %d",sum>sum1?sum:sum1);
    getch();

}

```

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Mukut • 10 months ago

```

#include<stdio.h>
#define no 20
int n;
int A[no];
int Sum(int i, int s, bool sel)
{
    int a = 0,b = 0;
    if(i == n)
    {
        if(!sel)
            return s + A[i];
        else

```

```
        return s;
    }
    if(!sel)
        a = Sum(i+1,s + A[i], true);
    b = Sum(i+1,s, false);
```

[see more](#)

^ | v • Reply • Share ›



shek8034 • 11 months ago

A Very Simple DP Solution.

Time : O(n).

Space: O(1).

Please go through this algorithm.

Let sum[i] represent maximum non-consecutive subsequence sum till ith element

$sum[i] = \max(sum[i-1], input[i] + input[i-2])$

which says that new sum would either be obtained by not including ith element with last to previous sum i.e input[i-2]. The new sum would be maximum of the

Since space complexity is O(1), instead of using sum[] array, we only need 3 variables, second last values of sum.

I'm using 3 variables here.

a -> for (i-2)th index.

b -> for (i-1)th index.

c -> for ith index. (This stores our answer).

This is 100% working code for all cases (negatives values also).

Check it out.

[see more](#)

^ | v • Reply • Share ›



smith → shek8034 • 5 months ago



smith · shek8034 · 5 months ago

input[i-2] must be sum[i-2]

^ | v · Reply · Share ›



smith → shek8034 · 5 months ago

good one

^ | v · Reply · Share ›



Vijay Daultani → shek8034 · 7 months ago

Why in 2nd test case output is -1 it could have been just 4

^ | v · Reply · Share ›



Ankit Chaudhary → shek8034 · 7 months ago

There are two flaws in ur code.

1. variable c is not initialised. In case array size of 2, function return

garbage value. So before for loop write statement

```
c=b;
```

2. Your code will not work if all elements in array are negative, otherwise

Modification in dp :

```
sum[i]=max(arr[i],sum[i-1],sum[i-2]+arr[i]);
```

Below is modified code . This will work even if all elements are negative

Correct me if I am wrong.

Why my code is not posted in readable form ?

I have tried many times but unable to post it in correct format.

code: modification in for loop :

```
c=b;
```

```
for(int i=2;i<n;i++) c="max(input[i],input[i]+a,b);" a="b;" b="c;" return=""
```

^ | v • Reply • Share ›



khurshid → shek8034 • 11 months ago

@shek8034 :

i think the dp should be

$sum[i] = \max(sum[i-1], input[i] + sum[i-2])$

@GeeksforGeeks: Please verify it .

2 ^ | v • Reply • Share ›



sajal jain → khurshid • 3 months ago

@khurshid : your code is correct..

^ | v • Reply • Share ›



shek8034 → khurshid • 11 months ago

@khurshid : I think my DP is correct and its working for all the c
if you find some difficulty.

Also, the problem statement says that no two element are adja
the two alternate elements or previous stored sum. Correct me

^ | v • Reply • Share ›



gourav pathak → shek8034 • 3 months ago

No, I think @khurshid is right

^ | v • Reply • Share ›



shek8034 • 11 months ago

A Very Simple DP Solution.

Time : $O(n)$.

Space: $O(1)$.

Please go through this algorithm.

Let $sum[i]$ represent maximum non-consecutive subsequence sum till i th element.
 $sum[i] = \max(sum[i-1], input[i] + input[i-2])$

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Since space complexity is $O(1)$, instead of using $sum[]$ array, we only need 3 variables to store the last two values of sum .

I'm using 3 variables here.

$a \rightarrow$ for $(i-2)$ th index.

$b \rightarrow$ for $(i-1)$ th index.

$c \rightarrow$ for i th index. (This stores our answer).

This is 100% working code for all cases (negatives values also).

Check it out.

[see more](#)

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coder! → shek8034 • 10 months ago

your algo is same as above

^ | v • Reply • Share ›



joker • 11 months ago

[sourcecode language="C++"]

```
int sum(vector<int> a)
{
    vector<int> dp(100);
    int i;
    dp[0]=a[0],dp[1]=a[1];
    for(i=2;i<a.size();i++)
        dp[i]=max(dp[i-2],dp[i-2]+a[i]);
    return max(dp[i-1],dp[i-2]);
}
```

```

main()
{
int n,k,x,i;
vector<int> a,ans;

scanf("%d",&n);
for(i=0;i<n;i++)
scanf("%d",&x), a.push_back(x);
printf("sum is: %d\n",sum(a));
}

```

1 ^ | v • Reply • Share ›



orchidmajumder • a year ago

Dynamic programming approach..

```

#include<stdio.h>
int max(int a,int b)
{
    if(a>=b)
        return a;
    return b;
}
int main()
{
    int a[]={3 ,2 ,7 ,10};
    int n=4,i;
    int f[10]={0};
    f[0]=a[0];
    f[1]=max(a[1],a[0]);
    for(i=2;i<n;++i)
        f[i]=max(f[i-2]+a[i],f[i-1]);
}

```

```
    printf("%d\n", r[n-1]);  
    return 0;  
}
```

3 ^ | v • Reply • Share ›



Tuhin Chakrabarty → orchidmajumder • 3 months ago

esob abar kobe :D

^ | v • Reply • Share ›



orchidmajumder → Tuhin Chakrabarty • 3 months ago

bochor khanek aage hobe :P

^ | v • Reply • Share ›



Amit → orchidmajumder • 6 months ago

Please check for this case:

int a[]={-3,-2,-1,-10};

o/p: -2

should be: -1

^ | v • Reply • Share ›



DraganWarrior → Amit • 6 months ago

Plz read Question carefully

This is only for array with +ve value

1 ^ | v • Reply • Share ›



Gaurav pruthi → orchidmajumder • 7 months ago

good one :)

^ | v • Reply • Share ›



nikhil • a year ago



A DP solution...

<http://algorithmsforever.blogspot...>

^ | v • Reply • Share ›



Nikhil Lohia • a year ago

what about a case when we say that "no 3 elements are adjacent"..
how can we modify the code to achieve this.

^ | v • Reply • Share ›



kT • a year ago

Hi,

I think this needs to be corrected :

>> excl = max(5, 15) = 65

instead should be excl = max(65, 45) = 65

Please correct me otherwise.

Thanks.

^ | v • Reply • Share ›



joker • a year ago

```
{{{  
int solve(vector a)  
{ int dp[10000];  
CLR(dp);  
dp[0]=a[0],dp[1]=max(a[0],a[1]);  
for(int i=2;i<a.size();i++) dp[i]=max(dp[i-2]+a[i],dp[i-1]);  
return dp[a.size()-1];  
}  
main()  
{  
int t;
```

```
...  
int b[]={5, 5, 10, 40, 50, 35};  
vector a(b,b+sizeof(b)/sizeof(int));  
printf("%d\n",solve(a));  
system("pause");  
return 0;  
}  
}}}
```

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