

Time Complexity of building a heap

Consider the following algorithm for building a Heap of an input array A.

```
BUILD-HEAP(A)
    heapsize := size(A);
    for i := floor(heapsize/2) downto 1
        do HEAPIFY(A, i);
    end for
END
```

What is the worst case time complexity of the above algo?

Although the worst case complexity looks like $O(n \log n)$, upper bound of time complexity is $O(n)$. See following links for the proof of time complexity.

<http://www.cse.iitk.ac.in/users/sbaswana/Courses/ESO211/heap.pdf/>

http://www.cs.sfu.ca/CourseCentral/307/petra/2009/SLN_2.pdf

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groomnestle · 5 months ago

It is easier to write heapsort with first heap index as 1 instead of 0, in this case with a simple formula:

parent = k, left child = 2k, right child = 2k+1.

^ | v .



stack!

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Venki · 3 years ago

The complexity $O(n \log n)$ is an upper bound. The upper bound is calculated as $\log n$. But in reality it is not the case. Height of node varies in the heap. By definition, path length from that node to leaf node. This way the root is at maximum height.

The build-heap time complexity is a function of node's height $f(h)$. How can we estimate it? Having N elements, we can observe that there will be a maximum of $f(h) = \lceil N/2^h \rceil$. Given $N = 7$ i.e. a full binary tree of height 2, the root is at height 2 and $f(2) = 1$, satisfying our assumption. In fact, it is an approximation and as the height of the tree approaches the exact limit.

Now, it is easy to check the tighter bound on building the heap. To build heap, we call $O(h)$ and we call it on nodes from $n/2$ to 1. These nodes form internal nodes. For computational cost we need to sum the cost of heapifying each node. Yet we need to find the cost of *a node* at each level and integrate (sum) them to find the total cost.

Mathematically $T(n) = \text{summation of } [h \times f(h)]$ in the interval 0 to $\log N$ (i.e. height of tree). We will arrive at $2N$ which is linear in N .

Hope it helps in better understanding the heap build analysis. The tighter bound is $O(n)$.

[see more](#)

3 ^ | v ·



Sandeep → Venki · 3 years ago

@Venki: Both $O(n)$ and $O(n \log n)$ are upper bounds for build heap. $O(n)$ is a tighter bound.

1 ^ | v ·



tk · 3 years ago

Interesting fact. Good to know.

^ | v ·



vinay · 3 years ago

one conceptual question from Cormen: why the loop goes down from heapsize
heapsize/2?

^ | v ·



Amit → Vinay · 3 years ago

If you index $i=1$ to $n/2$ you can observe that you will need to modify you
this and you will be able to see the difference).... so in Cormen to avoid th

Both indexing will work but the efficient way will be the latter case....

^ | v ·



kartik → Vinay · 3 years ago

I think it's because the way Heapify process works. When we Heapify a
the subtrees of i are heapified. Also, the Heapify process must go in up
maximum (or minimum) element is at the top.

1 ^ | v ·

705



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