

Construction of Longest Monotonically Increasing Subsequence (N log N)

In my previous post, I have explained about longest [monotonically increasing sub-sequence](#) (LIS) problem in detail. However, the post only covered code related to querying size of LIS, but not the construction of LIS. I left it as an exercise. If you have solved, cheers. If not, you are not alone, here is code.

If you have not read my previous post, read [here](#). Note that the below code prints LIS in reverse order. We can modify print order using a stack (explicit or system stack). I am leaving explanation as an exercise (easy).

```
#include <iostream>
#include <string.h>
#include <stdio.h>
using namespace std;

// Binary search
int GetCeilIndex(int A[], int T[], int l, int r, int key) {
    int m;

    while( r - l > 1 ) {
        m = l + (r - l)/2;
        if( A[T[m]] >= key )
            r = m;
        else
            l = m;
    }

    return r;
}

int LongestIncreasingSubsequence(int A[], int size) {
```

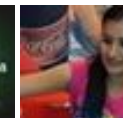
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```
// Add boundary case, when array size is zero
// Depend on smart pointers

int *tailIndices = new int[size];
int *prevIndices = new int[size];
int len;

memset(tailIndices, 0, sizeof(tailIndices[0])*size);
memset(prevIndices, 0xFF, sizeof(prevIndices[0])*size);

tailIndices[0] = 0;
prevIndices[0] = -1;
len = 1; // it will always point to empty location
for( int i = 1; i < size; i++ ) {
    if( A[i] < A[tailIndices[0]] ) {
        // new smallest value
        tailIndices[0] = i;
    } else if( A[i] > A[tailIndices[len-1]] ) {
        // A[i] wants to extend largest subsequence
        prevIndices[i] = tailIndices[len-1];
        tailIndices[len++] = i;
    } else {
        // A[i] wants to be a potential candidate of future subsequence
        // It will replace ceil value in tailIndices
        int pos = GetCeilIndex(A, tailIndices, -1, len-1, A[i]);

        prevIndices[i] = tailIndices[pos-1];
        tailIndices[pos] = i;
    }
}
cout << "LIS of given input" << endl;
for( int i = tailIndices[len-1]; i >= 0; i = prevIndices[i] )
    cout << A[i] << " ";
cout << endl;

delete[] tailIndices;
delete[] prevIndices;

return len;
}

int main() {
    int A[] = { 2, 5, 3, 7, 11, 8, 10, 13, 6 };
    int size = sizeof(A)/sizeof(A[0]);

    printf("LIS size %d\n", LongestIncreasingSubsequence(A, size));
}
```

```
    return 0;  
}
```

Exercises:

1. You know **Kadane's** algorithm to find **maximum sum sub-array**. Modify Kadane's algorithm to trace starting and ending location of maximum sum sub-array.
2. Modify **Kadane's** algorithm to find maximum sum sub-array in a circular array. Refer GFG forum for many comments on the question.
3. Given two integers A and B as input. Find number of Fibonacci numbers existing in between these two numbers (including A and B). For example, A = 3 and B = 18, there are 4 Fibonacci numbers in between {3, 5, 8, 13}. Do it in $O(\log K)$ time, where K is $\max(A, B)$. What is your observation?

— **Venki**. Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.



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- Kth smallest element in a row-wise and column-wise sorted 2D array | Set 1
- Find the number of zeroes
- Find if there is a subarray with 0 sum
- Divide and Conquer | Set 5 (Strassen's Matrix Multiplication)
- Count all possible groups of size 2 or 3 that have sum as multiple of 3



2



0



0

Writing code in comment? Please use ideone.com and share the link here.

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Manisha Barnwal · 11 months ago

/*Construction of Longest Monotonically Increasing Subsequence (N log N).
what about this?*/
#include<iostream>

```
struct store
{
    int index;.

    int len;.
};
using namespace std;.
```

```
//O(N LOG N).
void sort(store *arr, int size)
```

705



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```
void sort(store arr, int size).
```

```
{
```

```
int i, j, k;
```

```
store key;
```

see more

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binary001 • a year ago

```
int search(int in [] ,int a[],int i,int j,int x){
```

```
    while(i<j){
```

```
        int m=i+(j-i)/2;
```

```
        if(in[m]>x)
```

```
            j=m-1;
```

```
        else
```

```
            i=m+1;
```

```
    }
```

```
    return i;
```

```
}
```

```
int lis(int in[],int n){
```

```
    int ans=0;
```

see more

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abhishek08aug • a year ago



Trying out something. Shall update later how it performs.

Given an array of integers: array[n]

1) Create array greatest[n][n]

where,

greatest[i][j]=1, if there is no element in the array between index j
=0, if there is atleast one element in the array between
=-1, if j>n-1-i

2) k=n-1, l=0, last_added_element=Some number lower than all integers

```
while(k>0 && l<=n-1)
    if greatest[k][l]==1
        if k-1>=0 && greatest[k-1][l]!=1
```

[see more](#)

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abhishek08aug → abhishek08aug • a year ago

Here is the implementation of above algo with some fixes:

```
#include<stdio.h>
#include<stdlib.h>

int LongestIncreasingSubsequence(int array[], int n) {
    int ** greatest=(int **)malloc(sizeof(int *)*n);
    int i;
```

```

for(i=0; i<n; i++) {
    *(greatest+i)=(int *)malloc(sizeof(int)*n);
}

int j, largest_on_right;
for(i=0, largest_on_right=-9999; i<n; i++) {
    for(j=n-1; j>=0; j--) {
        if(j>n-1-i) {

```

[see more](#)

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anirudh beria • a year ago

won't the run time be n^2 because of search for min and max elements ?? I th will give $n \log n$ sol (possible if numbers are less than 10^6).

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Venki → anirudh beria • a year ago

@Anirudh, please be more specific. Where are we searching for min a

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ashish • a year ago

#include

```
void longest(int *,int);
```

```
void main()
```

```
{
```

```
int n,i,ary[50];
```

```
printf("enter the no of term:- ");
```

```
scanf("%d",&n);
```

```

for(i=0;i<n;i++)
scanf("%d",&ary[i]);

longest(ary,n);
}

void longest(int ary[],int n)

{
int i,j,l=0,maxlen=0;

```

[see more](#)

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suresh • a year ago

Why we need array prevIndices when tailIndices stores the Indices correspond

```

/* Paste your code here (You may delete these lines if not writing code)

```

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vivek • a year ago

How to do the fibonacci question here ?

```

/* Paste your code here (You may delete these lines if not writing code)

```

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Venki → vivek • a year ago

Okay, here it is. What do we need? Number of fibonacci numbers in between A and B. If we know the position of A and B in the Fibonacci sequence, we can find the number of fibonacci numbers between them.

Fibonacci numbers grow at the rate approximated by golden ratio, i.e. means, to calculate (n+1) Fibonacci number in same time of n-th num 61.8% faster machine.

Here, A and B corresponds to F(n1) and F(n2). We can easily find n1 & difference. Here is my code,

```
int FibonacciInBetween(long unsigned a, long unsigned b) {  
    int x = (int)(4.785 * log10((double)a));  
    int y = (int)(4.785 * log10((double)b));  
  
    return abs(y-x);  
}
```

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Ronny → Venki • 11 months ago

@venki @geeksforgeeks

Should not the returned value be $\text{abs}(y-x)-1$
since the question asks for number of fibonacci numbers BETWEEN

For a fibonacci series starting with 0,

0 1 1 2 3 5 8 13 21 34 55

pos of 5 is 6

pos of 8 is 7

the above method returns difference in position of a and b.

for the above case it will return 1, but there are no fibonacci num

So the return value should be $\text{abs}(y-x)-1$.

Kindly update the comment to avoid confusion

Kindly update the comment to avoid confusion.

^ | v • Reply • Share ›



sachin → Venki • a year ago

How did you get 4.785? I am not getting how you are getting the

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Venki → sachin • a year ago

\log_{10} to the base 1.618 = 4.785. Easy one, I expected it

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Venki → vivek • a year ago

It is easy and can be done in $O(1)$ time. Here is hint, at what rate Fibonacci relation between their growth rate and decimal numbers? Hope you will

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vivek → Venki • a year ago

I am still not able to figure out the soln. I was thinking to find the range of those 2 fib nos. to find other fibs in the range A to B.

How is $O(1)$ achieved? Using some Fibonacci nos. related formula?

```
/* Paste your code here (You may delete these lines if r
```

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vivek • a year ago

What is the solution to problem:

>> Find number of Fibonacci numbers existing in between these two numbers

asked above.

```
/* Paste your code here (You may delete these lines if not writing co
```

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nomind • a year ago

```
memset(tailIndices, 0, sizeof(tailIndices[0])*size);  
memset(prevIndices, 0, sizeof(prevIndices[0])*size);
```

replace these two lines with

```
memset(tailIndices, -2, sizeof(tailIndices[0])*size);  
memset(prevIndices, -2, sizeof(prevIndices[0])*size);
```

then it will work fine :)

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Venki → nomind • a year ago

Thanks.

In my code I have initialized it to 0xFF (effectively -1). Missed here. Upc

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nomind • a year ago

```
input int A[] = { 2, 5, 1, 3, 7, 11, 8, 10, 13, 6 }
```

output

LIS of given input

13 10 8 7 3 1 2

LIS size 6

output is wrong.

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rafi • a year ago

But just try this input :

```
int A[] = { 26, 13, 24, 25, 28, 10, 15, 4, 7, 21, 20, 23, 19, 22, 30,
```

And you have this:

LIS of given input

18 17 16 14 11 8 7 4 26

LIS size 8

So you have to set the sentinel each time you renew tailIndices[0].

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Venki → rafi • a year ago

@Rafi, we need a delimiter to recognize end of trace back. I have used following output on my machine.

Input :

26 13 24 25 28 10 15 4 7 21 20 23 19 22 30 29

Output :

Length of Longest Increasing Subsequence is 8, and it is

4 7 8 11 14 16 17 18

Note that there is possibility of many such LIS of same length.

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Venki · a year ago

An email comment to author: Comment by Rafi:

You forget to reset prevIndices here : if(A[i] < A[tailIndices[0]]) { // new smaller
be a : prevIndices[i] = -1; BTW thank for your 2 posts about LIS they're ji

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Venki → Venki · a year ago

@Rafi, this is not required. If you observe, the location prevIndices[0] a

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