# First Project

### **Summary**

Merge Sort is a timeless sorting algorithm that epitomizes the divide and conquer strategy, efficiently organizing data. This process entails dividing a large dataset into smaller sub-arrays, sorting each, and then merging them into a final sorted array. To visualize, imagine dividing a deck of cards into smaller stacks, sorting each, and then combining them.

Recursion is pivotal in Merge Sort, where a function calls itself to resolve smaller problem instances. This continues until reaching the base case, typically when the dataset reduces to a single element or a small, easily sortable subset. For your visual information, think about Russian nesting dolls, with each smaller doll emerging from within.

Merge Sort's linear-logarithmic time complexity (O(n log n)) ensures time-efficient data processing, particularly beneficial for large datasets. It maintains the original order of equal elements, unlike algorithms like Quick Sort, which has variable performance, or Bubble Sort, less efficient for larger datasets.

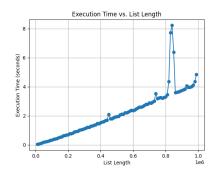
Historically, Merge Sort's development marked a significant advancement in algorithmic efficiency, widely applied in database management for its stable sorting. This stability ensures that records with equal keys retain their original sequence, critical in maintaining data integrity.

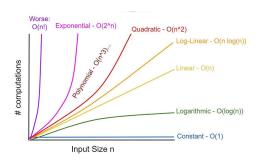
In terms of space complexity, Merge Sort requires additional space proportional to the dataset, a trade-off for its efficiency.

An example with an array [3, 1, 4, 1, 5, 9, 2, 6, 5, 3] demonstrates Merge Sort's process. The array is divided, sorted in sub-arrays, and merged into [1, 1, 2, 3, 3, 4, 5, 5, 6, 9].

To enhance understanding, imagine a flowchart illustrating these steps: splitting the array, sorting subparts, and the final merging. This visual aid complements the textual description, catering to visual learners.

In conclusion, Merge Sort is an algorithm of choice for various applications, from computer science to real-world data sorting. Its predictable performance and efficiency make it a staple in sorting algorithm discussions.





The chart on the left side is the actual relation between the size of the data and the execution time which I get by executing the algorithm thousands of times, the chart on the right side is the graph of n log n.

Through this effort, the time complexity of Merge Sort has experimentally been proven.



#### 이름 You

Could you explain me about Merge sort.



#### ChatGPT

Certainly! Merge sort is a popular and efficient comparison-based sorting algorithm. It follows the divide-and-conquer paradigm to sort a list of elements. The basic idea behind merge sort is to divide the unsorted list into smaller sublists, sort each of these sublists, and then merge them back together to produce a sorted list.

Here are the steps involved in the merge sort algorithm:

- Divide: The unsorted list is divided into two or more sublists. This is typically done by finding the
  middle point of the list. Divide the list into two halves, and then recursively divide each of these
  halves into smaller halves until you have sublists that contain only one element. These oneelement sublists are considered sorted by definition.
- Conquer: Once you have divided the list into one-element sublists, you start merging them back together in a sorted order. Merging involves comparing elements from two sublists and merging them into a new sorted sublist. This process is repeated until you have a single sorted list that contains all the elements.
- 3. Merge: Merging is the critical step in the merge sort algorithm. You take two sorted sublists and merge them into a single sorted sublist. The merge operation involves comparing elements from both sublists and selecting the smaller (or larger, depending on the sorting order) element to be placed in the new merged list. This process continues until both sublists are empty, and all elements are merged into the final sorted list.

Here's a high-level overview of the merge sort process:

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Here's a high-level overview of the merge sort process:

- · Divide the unsorted list into two halves.
- · Recursively sort each of the halves.
- Merge the sorted halves back together to create a single sorted list.

The key to the efficiency of merge sort is that it ensures that each element is compared only a constant number of times during the merge step. The time complexity of merge sort is O(n log n), where "n" is the number of elements in the list. This makes merge sort one of the most efficient sorting algorithms for large datasets.

Additionally, merge sort is a stable sort, meaning that it preserves the relative order of equal elements. It is also suitable for sorting linked lists, which is not the case with some other sorting algorithms.

In summary, merge sort is a reliable and efficient sorting algorithm that uses the divide-and-conquer strategy to sort a list of elements by dividing it into smaller sublists, sorting those sublists, and then merging them back together into a sorted list. Its time complexity of O(n log n) makes it a good choice for sorting large datasets.

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#### ChatGPT

Of course! Let's break down the merge sort algorithm in simpler terms:

Imagine you have a stack of unsorted cards, and you want to arrange them in ascending order. Here's how merge sort works:

- Divide: First, you split the stack into two roughly equal halves. You continue doing this until you
  have smaller stacks, each with just one card. These small stacks are considered sorted because a
  single card is always sorted by itself.
- 2. **Conquer (Sort)**: Now, you start merging the small stacks back together in a sorted manner. To do this, you pick the smallest card from the top of the two stacks you're merging and place it in a new stack. You repeat this process until you've combined all the cards into one sorted stack.
- 3. **Merge**: Merging is the heart of the algorithm. You compare the top cards from two stacks and pick the smaller one to place in the new stack. You keep doing this until both stacks are empty. This way, you're always putting cards in the new stack in the correct order.
- 4. **Repeat**: You continue this process, dividing and merging stacks until you've combined all the cards into a single sorted stack.

The result is a completely sorted stack of cards, and this is exactly what the merge sort algorithm does with a list of elements. It repeatedly divides the list into smaller parts, sorts those parts, and then merges them back together into a fully sorted list.

The key takeaways are that merge sort breaks the problem into smaller, more manageable pieces, and it ensures that the merging step always places elements in the correct order. This approach guarantees that the final result is a sorted list.

Merge sort is efficient and reliable, making it a great choice for sorting data, especially when dealing with large datasets.

I adjusted the level of conversation to suit my level. As a result, I was able to understand merge sort in terms of the image of organizing cards.

o⊪ You

Please explain me about the term "recursively'

ChatGPT

Certainly! The term "recursively" is related to the concept of recursion in computer science and mathematics. In simple terms, recursion is a technique where a function or a process calls itself to solve a smaller or simpler instance of the same problem. This process continues until the problem becomes trivial, and the solutions to the simpler instances are combined to solve the original problem.

Here are some key points to understand about recursion and the term "recursively":

- Recursive Function: A recursive function is a function that calls itself within its own definition.
   When the function is called with a particular input, it breaks the problem down into smaller, similar subproblems and calls itself to solve those subproblems.
- Base Case: In a recursive function, there is typically a base case or termination condition. This is a
  condition that specifies when the recursion should stop. When the base case is met, the function
  stops calling itself and starts returning values back up the call stack.
- 3. Call Stack: Recursive calls are managed by a data structure called the call stack. Each time a function is called, a new stack frame is created to store information about that specific call, including its parameters and the location where it should return to after it's done. When a function returns, its stack frame is removed.
- 4. **Example**: A classic example of recursion is computing the factorial of a number. The factorial of a non-negative integer "n" (denoted as "n!") is the product of all positive integers from 1 to "n." It can be defined recursively as follows:

```
java

n! = 1 if n = 0

n! = n * (n-1)! for n > 0
```

ChatGPT mentioned that a recursive function is defined by calling itself within its definition. It is said that a recursive function calls itself when dealing with smaller problems after division. Additionally, a key concept necessary for recursive processing is the base case, which occurs when the data is in its simplest form and cannot be simplified any further for processing.



To apply a recursive function in the script, I first understood that it involves the function calling itself. Now, I'm asking about the code for determining the base case.



#### You

I want to know more about the time complexity of merge sort



#### ChatGPT

The time complexity of the merge sort algorithm is O(n log n), where "n" is the number of elements in the list being sorted. This makes merge sort one of the most efficient general-purpose sorting algorithms, especially for large datasets.

Here's a breakdown of why merge sort has a time complexity of O(n log n):

- 1. Divide Phase: In the divide phase, the input list is repeatedly divided into two equal halves until you have sublists with only one element each. This process takes O(log n) time because you can divide the list into halves approximately log<sub>2</sub>(n) times.
- 2. Conquer Phase (Merging): In the conquer phase, you merge the sorted sublists. The merging process involves comparing and combining elements from the sublists, and it takes O(n) time in the worst case because you may need to compare and move every element once.

Since the divide phase takes O(log n) time, and the conquer phase takes O(n) time, the overall time complexity of merge sort is O(n log n). This makes it an efficient sorting algorithm, and it's not significantly affected by the initial order of elements, making it a reliable choice for sorting various types of data.

It's important to note that the O(n log n) time complexity of merge sort is guaranteed for both the average and worst-case scenarios, making it a stable and predictable sorting algorithm. This efficiency makes it a popular choice for sorting large datasets in practice, and it is often used as the sorting algorithm in various programming libraries and languages.

I revisited the topic of time complexity that was mentioned earlier. I inquired about the time complexity of merge sort, noting that it allocates log n in the division process and n in the sorting process for each division layer. This is described as a linear-logarithmic function, which increases less sharply compared to other algorithms. As a result, even when processing large data sets, the time complexity does not significantly increase, which is considered an advantage.

### o⊪ You

```
def merge_sort(input_list):
    if len(input_list) // 2 == 0:
        return input_list

    data_a, data_b = divide(input_list)

    data_a = merge_sort(data_a)
    data_b = merge_sort(data_b)

    output_list = conquer_and_merge(data_a, data_b)

    return output_list
```

I created the code about merge sort. Is there a better option to determine the base case?

### **\$**

#### ChatGPT

Your code for merge sort looks almost correct, but there is a small issue with the base case condition. The condition 'len(input\_list) // 2 == 0 'checks if the length of 'input\_list' is exactly 0, which is not the correct base case for merge sort. The correct base case is when the length of the list is 1 or less because, in merge sort, you want to continue dividing the list until you have single-element sublists, which are considered sorted.

So, you should change the base case condition to `len(input\_list) <= 1`. Here's the corrected code:

I wrote code based on the information provided at the end and asked about a better option for the part of the code that determines the base case. ChatGPT recommended using equality over division, so I decided to reflect this in my script.

## **Script**

```
if len(input list) <= 1:</pre>
     return input list
  data a, data b = divide(input list)
  data a = merge sort(data a)
  data b = merge sort(data b)
  return output list
def divide(data):
  split_point = len(data) // 2
  data_a = data[:split_point]
  data_b = data[split_point:]
```

```
def conquer and merge(data a, data b):
         result.append(data a[a index])
          result.append(data_b[b_index])
input list = list(map(int, input().split(',')))
sorted_list = merge_sort(input_list)
```

print(sorted_list)		