



Data Structure and Algorithms

Session-25

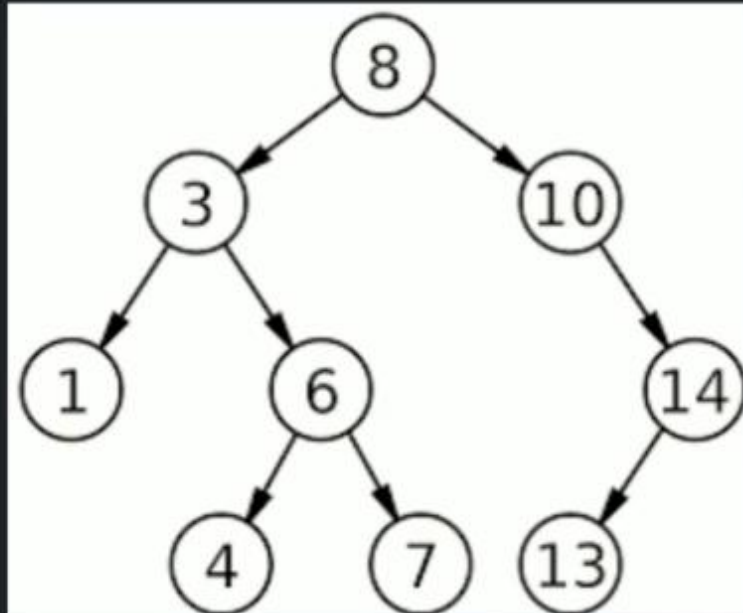
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What is BST ?

Binary Search Tree (BST) is a Binary Tree in which all the nodes follows the below-mentioned properties:

- ✓ *The left sub-tree of a node has a key less than or equal to its parent node's key.*
- ✓ *The right sub-tree of a node has a key greater than to its parent node's key*

Sample BST:




Why should we learn BST ?

Operation	Array	Linked List	Tree
Creation	$O(1)$	$O(1)$	Can we improve ? Let's see...
Insertion	$O(n)$	$O(n)$	
Deletion	$O(n)$	$O(n)$	
Searching	$O(n)$	$O(n)$	
Traversing	$O(n)$	$O(n)$	
Deleting entire Array/LinkedList/Tree	$O(1)$	$O(1)$	
Space Efficient ?	No	Yes	



Common operations of BST:

- ✓ *Creation of BST*
 - ✓ *Search for a value*
 - ✓ *Traverse all nodes*
 - ✓ *Insertion of a node*
 - ✓ *Deletion of a node*
 - ✓ *Deletion of BST*
- 

Algorithm - Creation of blank BST:

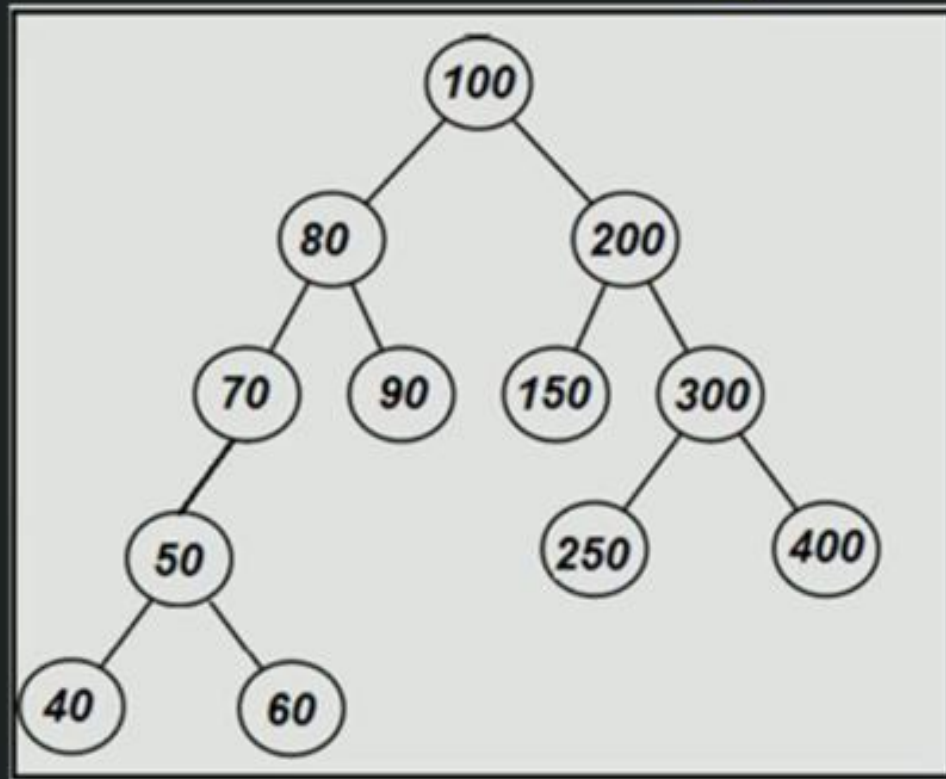
`createBST()`

Initialize Root with null

Time Complexity – $O(1)$

Space Complexity – $O(1)$

Searching a node in BST:



Algorithm - Searching a node in BST:

BST_Search (root, value):

if (root is null)

return null

else if (root == value)

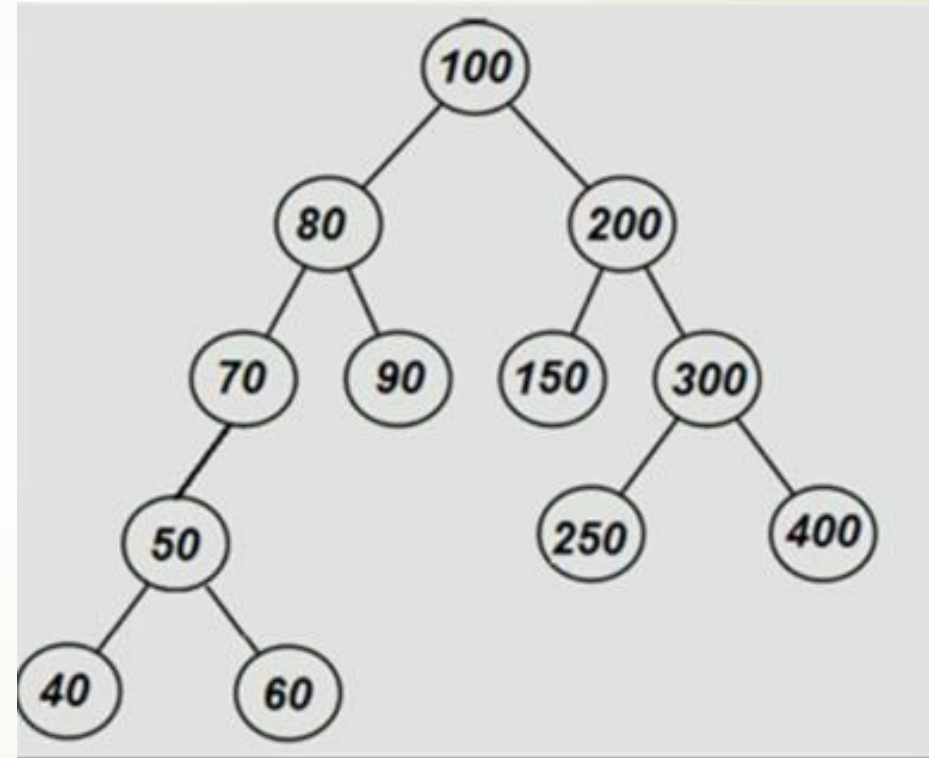
return root

else if (value < root)

BST_Search (root.left, value)

else if (value > root)

BST_Search (root.right, value)



Time & Space Complexity - Searching a node in BST:

```
BST_Search (root, value) -----  $T(n)$   
    if (root is null) -----  $O(1)$   
        return null -----  $O(1)$   
    else if (root == value) -----  $O(1)$   
        return root -----  $O(1)$   
    else if (value < root) -----  $O(1)$   
        BST_Search (root.left, value) -----  $T(n/2)$   
    else if (value > root) -----  $O(1)$   
        BST_Search (root.right, value) -----  $T(n/2)$ 
```

Time Complexity – $O(\log n)$

Space Complexity – $O(\log n)$ (because of recursive call)

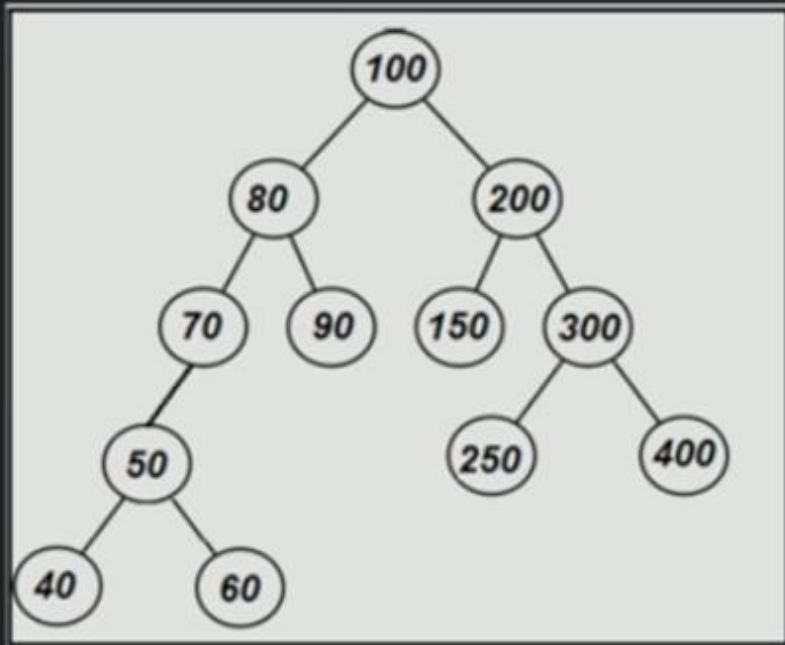
Traversal of BST:

✓ *Depth First Search:*

- ✓ *PreOrder Traversal*
- ✓ *InOrder Traversal*
- ✓ *PostOrder Traversal*


✓ *Breadth First Search:*

- ✓ *LevelOrder Traversal*

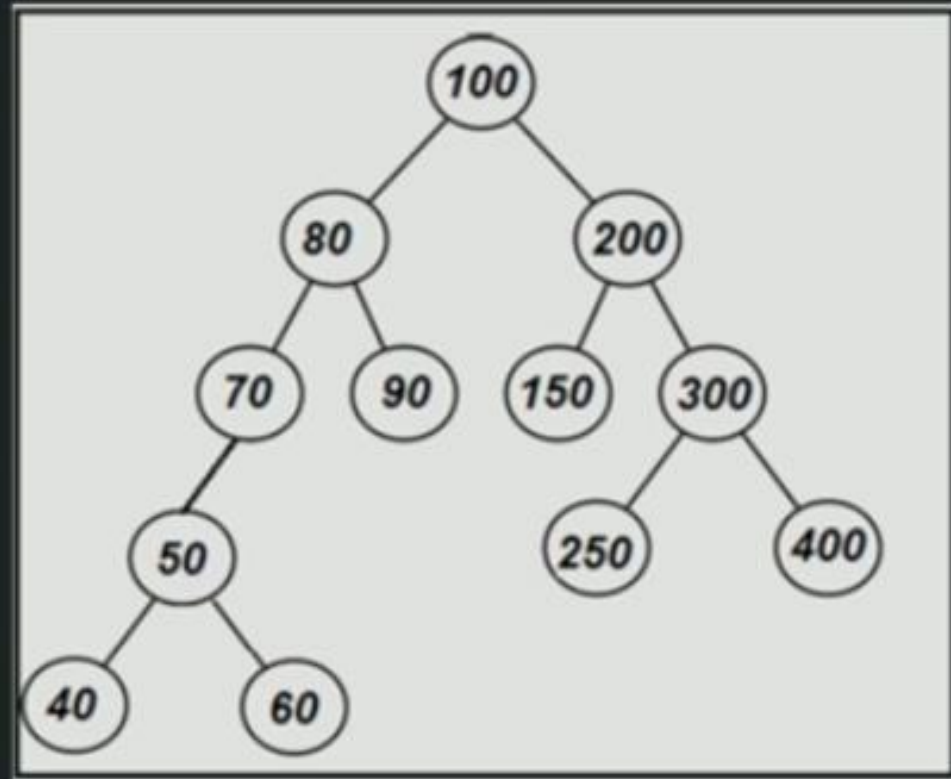


Algorithm - Pre-Order Traversal of BST:

```
preorderTraversal(root)

if (root equals null) 
    return error message

else
    print root
    preorderTraversal (root.left)
    preorderTraversal(root.right)
```



Algorithm - 'In-Order Traversal' of BST:

inOrderTraversal (root)

if (root equals null)

return

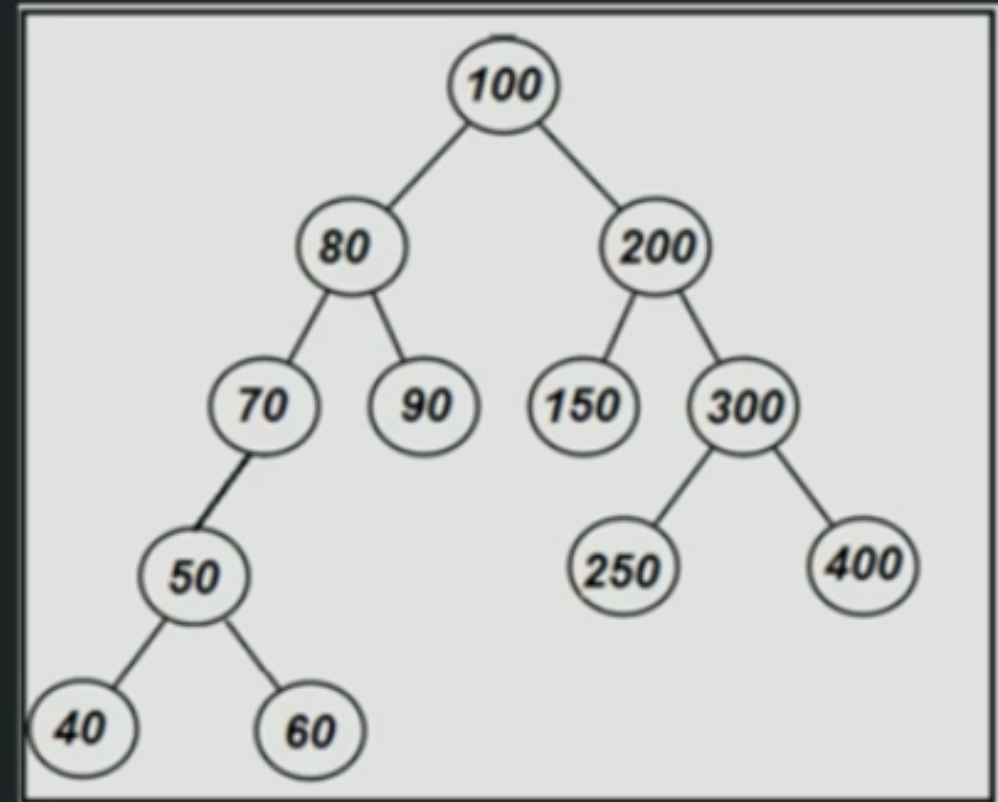


else

inOrderTraversal(root.left)

print root

inOrderTraversal(root.right)



Algorithm - 'Post-Order Traversal' of BST:

postOrderTraversal(root)

if (root equals null)

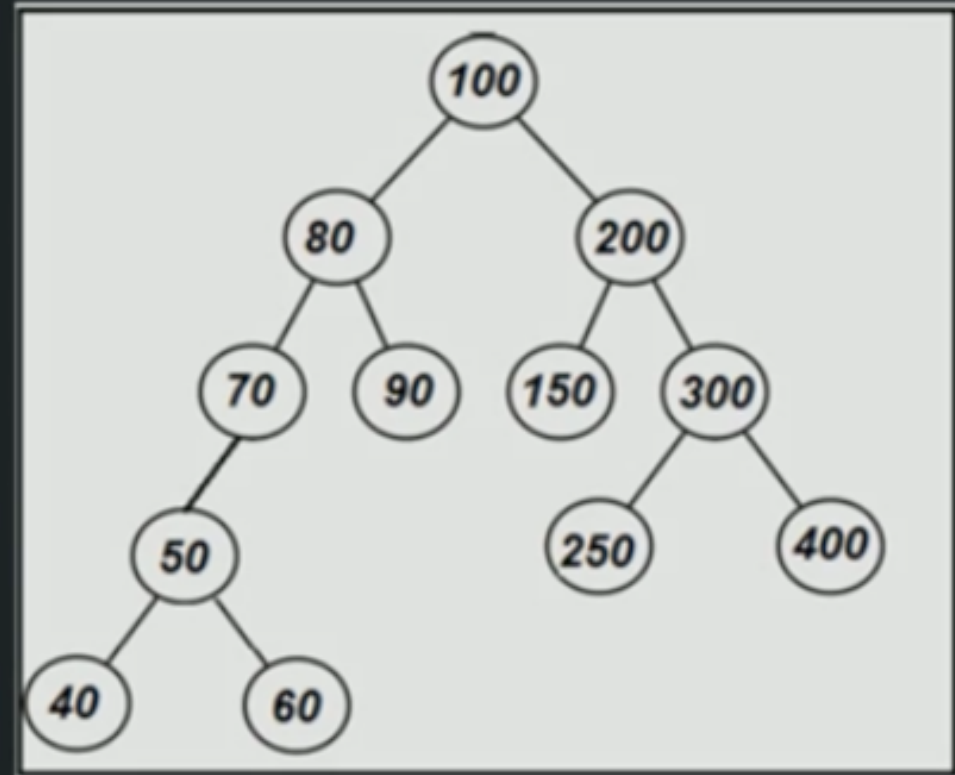
return

else

postOrderTraversal(root.left)

postOrderTraversal(root.right)

print root



Algorithm - 'Level Order Traversal' of BST:

levelOrderTraversal(root)

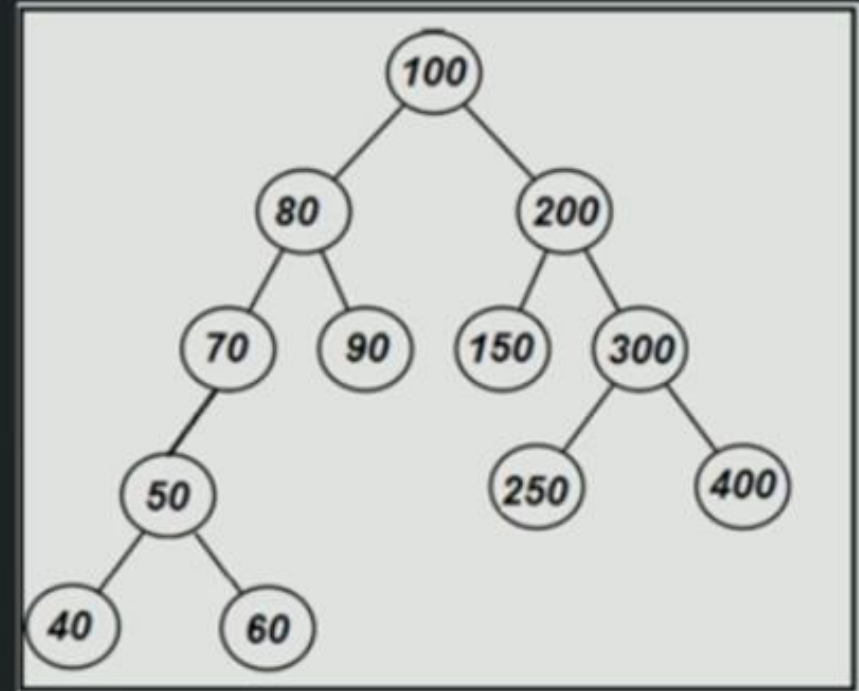
Create a Queue(Q)

enqueue(root)

While (Queue is not empty)

 dequeue() and print

 enqueue() the child of dequeued element



Algorithm - Inserting a node in BST:

Cases:

1. BST is blank
2. BST is non-blank



```
BST_Insert (currentNode, valueToInsert)
```

```
  if (currentNode is null)
```

```
    create a node, insert 'valueToInsert' in it
```

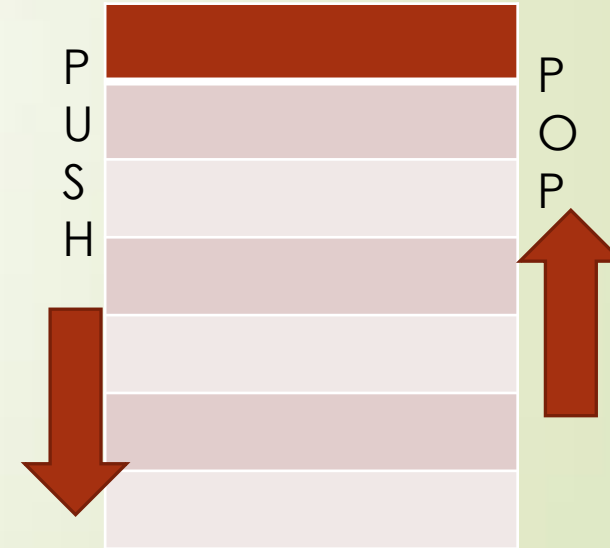
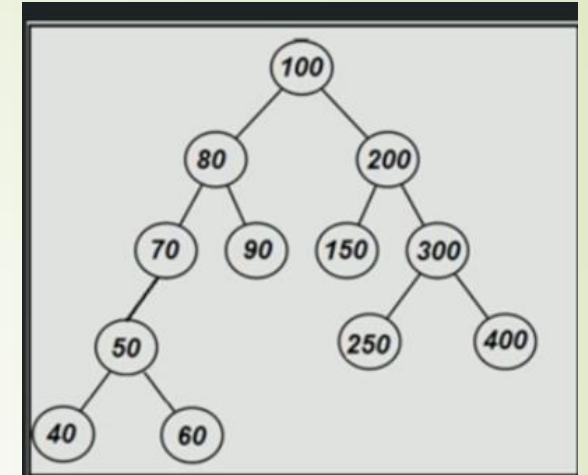
```
  else if (valueToInsert <= currentNode 's value)
```

```
    currentNode.left = BST_Insert (currentNode.left, valueToInsert)
```

```
  else
```

```
    currentNode.right = BST_Insert (currentNode.right, valueToInsert)
```

```
  return currentNode
```



Algorithm - Inserting a node in BST:

Cases:

1. BST is blank

2. BST is non-blank



```
BST_Insert (currentNode, valueToInsert)
```

```
  if (currentNode is null)
```

```
    create a node, insert 'valueToInsert' in it
```

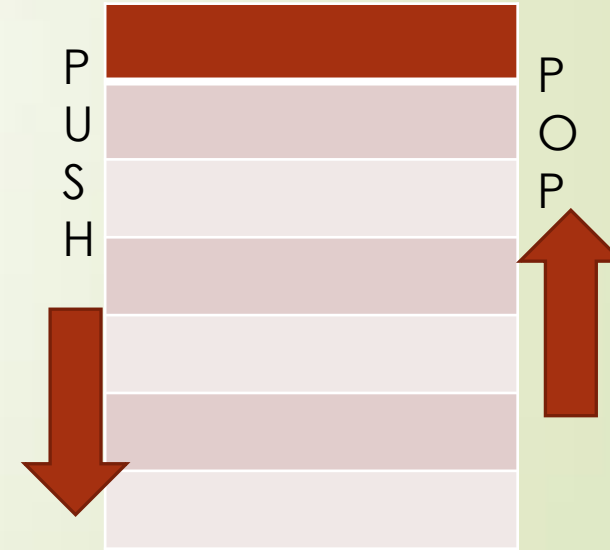
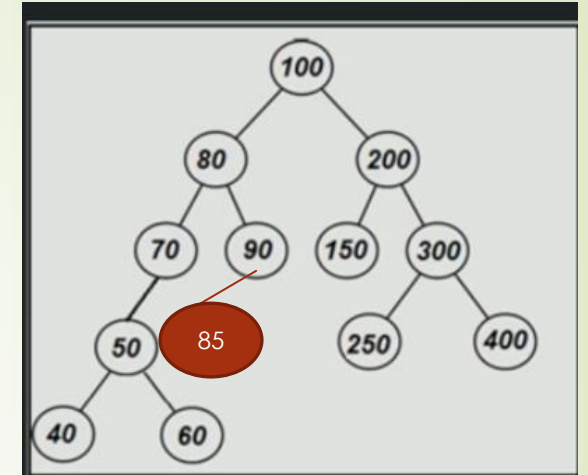
```
  else if (valueToInsert <= currentNode 's value)
```

```
    currentNode.left = BST_Insert (currentNode.left, valueToInsert)
```

```
  else
```


```
    currentNode.right = BST_Insert (currentNode.right, valueToInsert)
```

```
  return currentNode
```



Time & Space Complexity - Inserting a node in B

BST_Insert (currentNode, valueToInsert) ----- $T(n)$

if (currentNode is null) ----- $O(1)$ 

create a node, insert valueToInsert in it ----- $O(1)$

else if (valueToInsert <= currentNode 's value) ----- $O(1)$

currentNode.left = BST_Insert (currentNode.left, valueToInsert) ----- $T(n/2)$

else ----- $O(1)$

currentNode.right = BST_Insert (currentNode.right, valueToInsert) ----- $T(n/2)$

return currentNode ----- $O(1)$

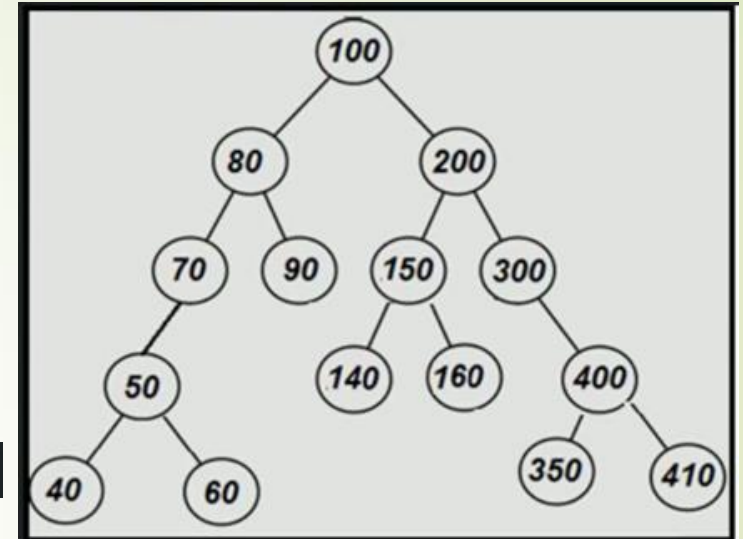
Time Complexity – $O(\log n)$

Space Complexity – $O(\log n)$

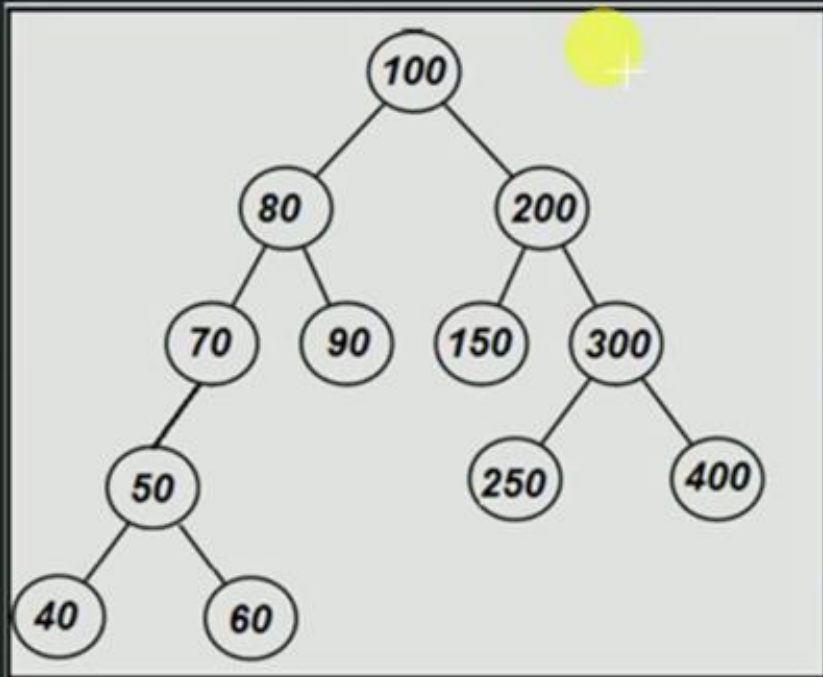
Deletion of a node from BST:

- ✓ Node to be deleted is leaf node
- ✓ Node to be deleted is having 1 child
- ✓ Node to be deleted has 2 children

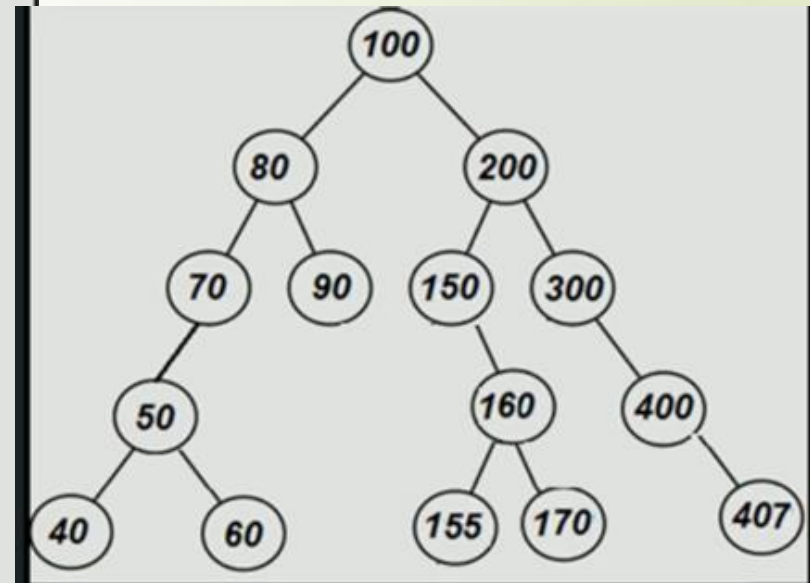
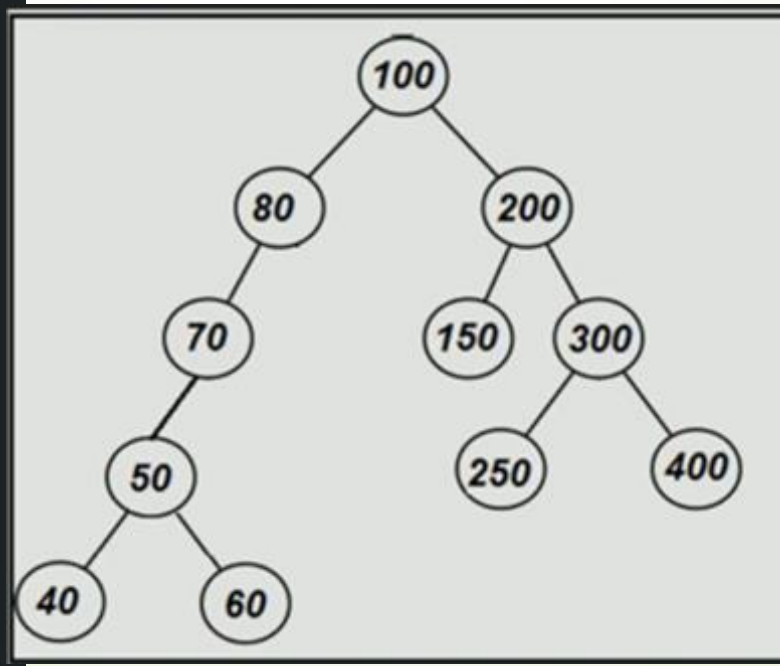
✓ Case#3 – Node to be deleted has 2 children




✓ Case#1 - Node to be deleted is leaf node



✓ Case#2 – Node to be deleted is having 1 child



Algorithm - Deletion of a node from BST:

 deleteNodeOfBST (root, valueToBeDeleted):

if (root == null) return null;

if (valueToBeDeleted < root.Value)

then root.left = deleteNodeOfBST (root.left, valueToBeDeleted)

else if (valueToBeDeleted > root.value)

then root.right = deleteNodeOfBST (root.right, valueToBeDeleted)

else // If currentNode is the node to be deleted

if root have both children, then find minimum element from right subtree (Case#3)

replace current node with minimum node from right subtree and delete minimum node from right

else if nodeToBeDeleted has only left child (Case#2)

then root = root.Left()

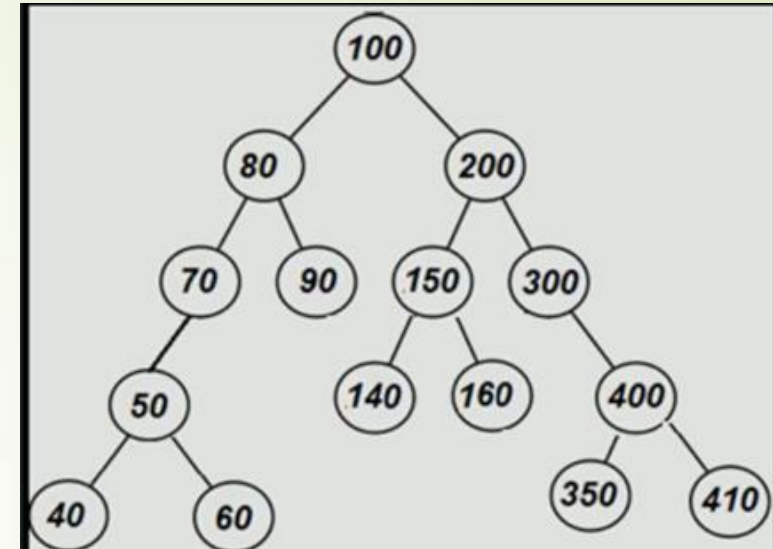
else if nodeToBeDeleted has only right child (Case#2)

then root = root.Right();

else // if nodeToBeDeleted do not have child (Case#1)

root = null;

return root;



Time & Space Complexity

```
deleteNodeOfBST (root, valueToBeDeleted) -----  $T(n)$ 

if (root == null) return null; -----  $O(1)$ 

if (valueToBeDeleted < root.Value) -----  $O(1)$ 

    then root.left = deleteNodeOfBST (root.left, valueToBeDeleted) -----  $T(n/2)$ 

else if (valueToBeDeleted > root.value) -----  $O(1)$ 

    then root.right = deleteNodeOfBST(root.right, valueToBeDeleted) -----  $T(n/2)$ 

else // If currentNode is the node to be deleted -----  $O(1)$ 

    if root have both children, then find minimum element from right subtree -----  $O(\log n)$ 

        replace current node with minimum node from right subtree and delete minimum node from right -----  $O(1)$ 

    else if nodeToBeDeleted has only left child -----  $O(1)$ 

        then root = root.Left() -----  $O(1)$ 

    else if nodeToBeDeleted has only right child -----  $O(1)$ 

        then root = root.Right(); -----  $O(1)$ 

    else // if nodeToBeDeleted do not have child (Leaf node) -----  $O(1)$ 

        root = null; -----  $O(1)$ 

return root -----  $O(1)$ 
```

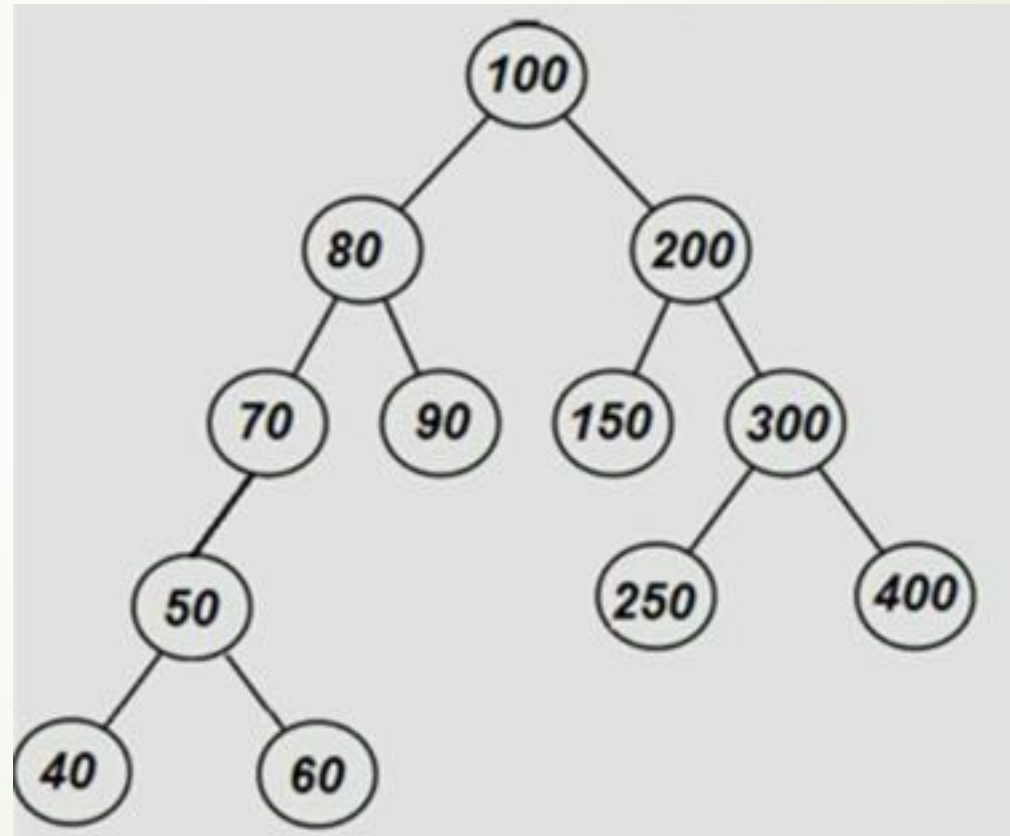
Time Complexity – $O(\log n)$

Space Complexity – $O(\log n)$

Algorithm - Deletion of entire BST:

DeleteBST()

root = null



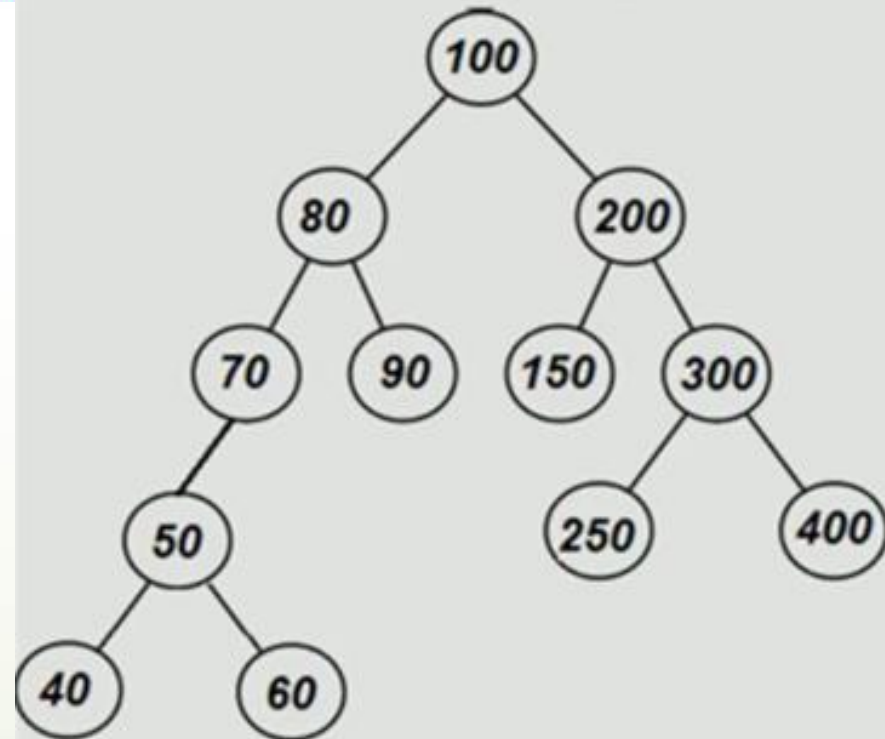
Find Minimum and Maximum of a BST

TREE-MINIMUM (x)

```
while left[x]  $\neq$  NIL do  
  x  $\leftarrow$  left [x]  
return x
```

TREE-MAXIMUM (x)

```
while right[x]  $\neq$  NIL do  
  x  $\leftarrow$  right [x]  
return x
```





Thank
you