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Questions

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Deadlock

4/4 points (graded)

In the code below three threads 1, 2, and 3 are trying to acquire locks on objects `alpha`, `beta`, and `gamma`.

Thread 1

```
synchronized (alpha) {  
    // using  $\alpha$   
    // ...  
}  
  
synchronized (gamma) {  
    synchronized (beta) {  
        // using  $\beta$  &  $\gamma$   
        // ...  
    }  
}  
// finished
```

Thread 2

```
synchronized (gamma) {  
    synchronized (alpha) {  
        synchronized (beta) {  
            // using  $\alpha$ ,  $\beta$  &  $\gamma$   
            // ...  
        }  
    }  
}  
// finished
```

Thread 3

```
synchronized (gamma) {  
    synchronized (alpha) {  
        // using  $\alpha$  &  $\gamma$   
        // ...  
    }  
}  
  
synchronized (beta) {  
    synchronized (gamma) {  
        // using  $\beta$  &  $\gamma$   
        // ...  
    }  
}  
// finished
```

This system is susceptible to deadlock.

For each of the scenarios below, determine whether the system is in deadlock if the threads are currently on the indicated lines of code.

Scenario A

Thread 1 inside `using alpha`Thread 2 blocked on `synchronized (alpha)`

Thread 3 finished

☐ deadlock☒ not deadlock**Explanation**

Thread 1 will exit the top synchronized block, release the lock on `alpha`, and the system will continue.

Scenario B

Thread 1 finished

Thread 2 blocked on `synchronized (beta)`

Thread 3 blocked on 2nd `synchronized (gamma)`

☒ deadlock☐ not deadlock**Explanation**

Thread 2 has acquired the lock on `gamma` and is awaiting `beta`. Thread 3 has `beta` and wants `gamma`. Deadlock.

Scenario C

Thread 1 running `synchronized (beta)`

Thread 2 blocked on `synchronized (gamma)`

Thread 3 blocked on 1st `synchronized (gamma)`

☐ deadlock☒ not deadlock**Explanation**

Thread 1 can successfully acquire the lock on `beta`, then exit the synchronized block, and one of the other threads will be able to acquire the lock on `gamma`. (As we saw in scenario B, they could deadlock later!)

Scenario D

Thread 1 blocked on `synchronized (beta)`

Thread 2 finished

Thread 3 blocked on 2nd `synchronized (gamma)`

☒ deadlock☐ not deadlock

Explanation

Thread 1 has acquired the lock on `gamma` and is awaiting `beta`. Thread 3 has `beta` and wants `gamma`. Deadlock.

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1/1 point (graded)

Examine the code again.

In the previous problem, we saw deadlocks involving `beta` and `gamma`.

What about `alpha`?

☐ there is a possible deadlock where thread 1 owns the lock on `alpha`☐ there is a possible deadlock where thread 2 owns the lock on `alpha`☐ there is a possible deadlock where thread 3 owns the lock on `alpha`☒ there are no deadlocks involving `alpha`**Explanation**

We can reason about it this way: in order to encounter deadlock, threads must try to acquire locks in different orders, creating a cycle in the graph of who-is-waiting-for-who.

So we look at `alpha` vs. `beta`: are there two threads that try to acquire these locks in the opposite order? No. Only thread 2 acquires them both at the same time.

Next we look at `alpha` vs. `gamma`: are there two threads that try to acquire these locks in the opposite order? No. Both thread 2 and thread 3 acquire both locks, but both of them acquire `gamma` first, then `alpha`.

Use the same analysis to demonstrate why `beta` and `gamma` are susceptible to deadlock: what order are these locks acquired in by the different threads?

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