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Tutorial : FSM

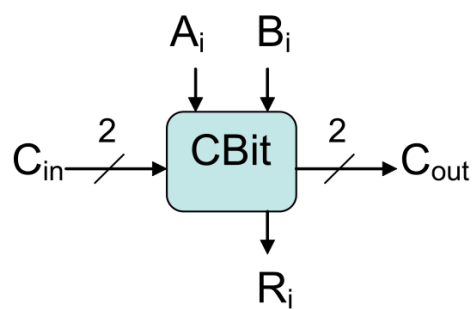
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 Calculator

Sequential Logic Timing of FSM

4/4 points (ungraded)

MaxOut is a Cambridge startup whose products are binary comparators which determine the largest of several unsigned binary integers. A building block common to all MaxOut products is the combinational CBit module depicted below.



Each CBit module takes corresponding bits of two unsigned binary numbers, A and B, along with two C_{in} bits from higher-order CBit modules. Its output bit, R, is the appropriate bit of the larger among A and B, as determined from these inputs; it passes two C_{out} bits to lower-order CBit modules.

The propagation delay of each CBit module is 10ns. The two C_{out} bits indicate, respectively, if $A > B$ or $B > A$ in the bits considered thus far.

MaxOut's latest product is the MAXFSM. The MAXFSM is to be a clocked finite state machine that takes two N-bit binary numbers, $A_{N-1:0}$ and $B_{N-1:0}$ in bit-serial form, most significant bit first, and outputs the larger of these numbers as $R_{N-1:0}$ also in bit-serial form, but delayed by one clock cycle. The MAXFSM is a Moore machine; recall that this means its output is strictly a function of its current state (like those FSMs shown in lecture).

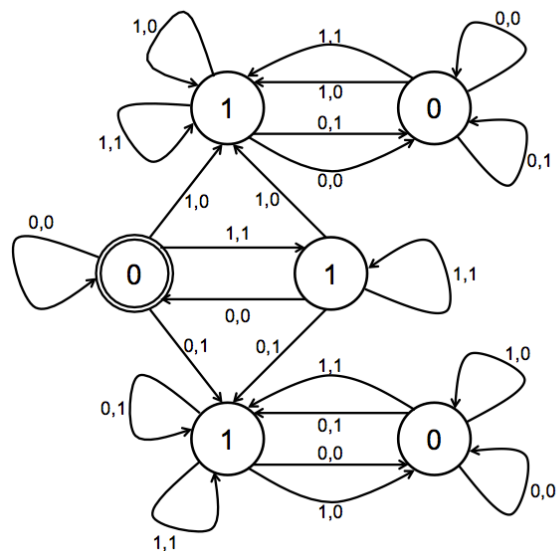


Before each active clock edge, the i^{th} bit of the A and B inputs are applied; during the **next** clock cycle, the i^{th} bit of the larger of the two input numbers appears at the R output. Note that the serial output bits are delayed with respect to the input bits by one clock cycle in order to allow each i^{th} output bit to be influenced by the i^{th} input bits.

(A) What is the minimum number of states necessary to implement a Moore machine obeying the above specifications? ✔ Answer: 6

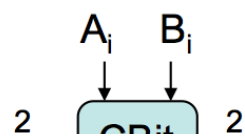
Explanation

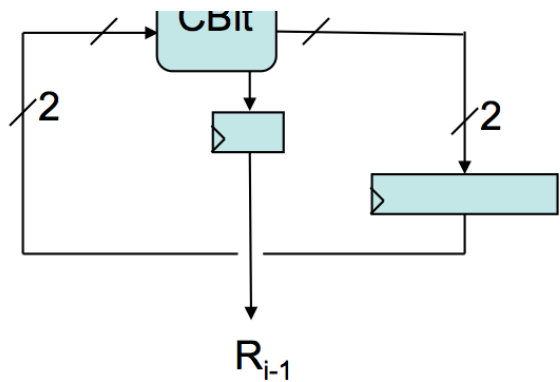
The state machine below implements the specification in 6 states.



The middle two states are the output when A and B have been the same so far. The upper two states are when A is greater than B, so the output follows A. The bottom two states are when B is greater than A, so the output follows B.

Ignoring your answer, MaxOut decides to build the MAXFSM using a CBit module and several flipflops, as shown in the diagram below:





The registers have the following specifications:

t_{pd}	5ns
t_{cd}	2ns
t_s	4ns
t_h	2ns

Recall that the CBit has a 10ns propagation delay; assume that its contamination delay is 0ns.

(B) What is the shortest clock period, in ns, for which this circuit will operate reliably?

Clock period (ns) \geq ✓ Answer: 19

Explanation

The clock period needs to be long enough for the next state to be stored in the register. There are three parts to that duration: the propagation delay of the register (5ns), the propagation delay of the CBit (10ns), and the setup time of the register (4ns). So the minimum clock period is $5 + 10 + 4 = 19$.

(C) What setup and hold time requirements should be specified for this FSM?

FSM Setup Time (ns): ✓ Answer: 14
FSM Hold Time (ns): ✓ Answer: 2

Explanation

The setup time for the FSM is how long the input needs to be stable and valid before the rising clock edge. An input needs to propagate through the CBit and arrive at the register with enough time for the register to setup. So the setup time of the FSM is the propagation delay of the CBit (10ns) + the setup time of the register (4ns), which equals 14ns.

The hold time for the FSM is how long after a rising clock edge the input needs to be stable and valid in order for the register's hold-time specification to be met. The contamination delay of the CBit is zero, so the input to the register becomes invalid as soon as the input to the FSM does. In order to satisfy the register's hold time, the hold time for the FSM must be the same as the hold time of the register, which is 2ns.

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i Answers are displayed within the problem

FSM

1/1 point (ungraded)

A "Froboz" is a clocked device built out of 3 interconnected components, each of which is known to be a 4-state FSM. What numeric upper bound, if any, can you put on the number of states of a Froboz?

✓

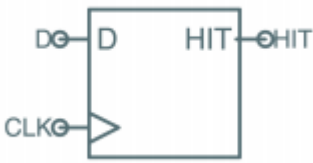
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✔ Correct (1/1 point)

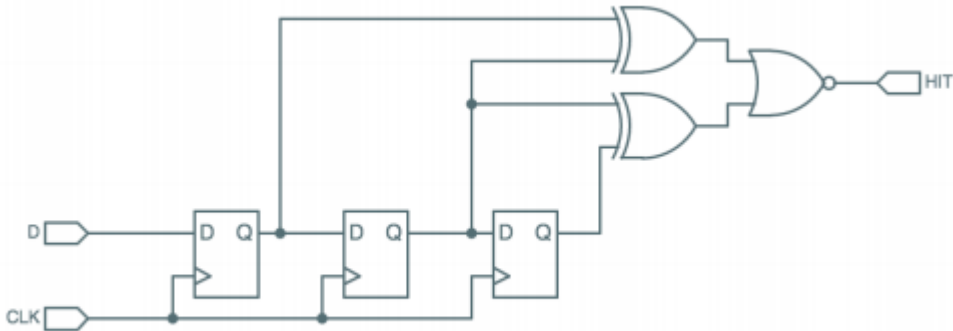
FSM

10/10 points (ungraded)

The NSA has an urgent need for a simple clocked FSM that monitors a sequence of binary data bits, appearing one bit per clock cycle, and signals a HIT whenever the last three data bits have been the same: in other words, whenever the sequence 000 or 111 appears in the input data. The circuit symbol for NSA's desired FSM is shown to the right: the D input senses one bit per clock cycle, and the HIT output must be 1 whenever three consecutive bits have had the same value. Of course, the timing of D input changes must obey the module's setup and hold time specifications. NSA's Chief Engineer, Philip Philop (a name you might recognize from his college celebrity as Captain of Harvard's Curling team) has proposed the following



circuit:



A) How many states does Philip's FSM have? ✔ Answer: 8

Explanation
The FSM keeps track of the last three input bits. These three bits can represent 8 possible combinations, so the FSM has 8 possible states.

Philip arranged to pre-determine the states of the three flip flops in the circuit on power-up, and had them all start in the 0 state. Unfortunately, this made the output logic (which correctly determines if all three flip flops are in the same state) report a HIT when no data had been entered. Philip, defending his circuit, said "Oh, I just picked the wrong starting state".

B) Is there ANY initial state for which Philip's design will work properly? If so, give the 3-bit initial state; otherwise, write NONE.

✔ Answer: NONE

Explanation
There is no starting state for which the circuit will work properly as the initial state will always affect the output for the first couple of data entries. For example if you started with 101, then entering two 1's (not three) would produce a hit.

While emptying the trash, a sharp-eyed NSA janitor notices a partially-completed table showing state transitions for a 7-state FSM, apparently discarded by a recent MIT intern who was hired after acing 6.004 and worked on this project. The table is shown below. On seeing the table, Philip shouts "EUREKA" and immediately orders his FSM to be replaced by the newly-discovered version.

Unfortunately, he needs your help...

C) Fill in the missing entries to the state transition table.

Current State	Input (D)	Next State	HIT
Sx	0	S0	0
Sx	1	S1	0
S0	0	S00	0
S0	1	<input type="text" value="S1"/> ✔ Answer: S1	0
S00	0	S000	0

Sx	0	Sxxx	0
S00	1	<div>S1</div> ✓ Answer: S1	0
S000	0	<div>S000</div> ✓ Answer: S000	<div>1</div> ✓ Answer: 1
S000	1	S1	<div>1</div> ✓ Answer: 1
S1	0	S0	0
S1	1	S11	0
S11	0	S0	<div>0</div> ✓ Answer: 0
S11	1	S111	<div>0</div> ✓ Answer: 0
S111	0	S0	1
S111	1	<div>S111</div> ✓ Answer: S111	1

Explanation

The initial state is S_x which indicates that no inputs have yet been received. The remaining states keep track of the maximum number of 0's or 1's that have recently been seen. Thus, from state S_x , a 0 input takes you to S_0 and a 1 input takes you to S_1 . Continuing this logic, if you are in S_0 and receive another 0 you move to state S_{00} which indicates that you have already seen 2 0's. Similarly, from state S_1 , receiving another 1 input moves you to state S_{11} . The same logic follows from S_{00} on a 0 input you go to S_{000} , and from S_{11} on a 1 input you go to S_{111} . States S_{000} and S_{111} indicate that you have received 3 0's in a row or 3 1's in a row, so when you are in those states, the HIT output is 1. For all other states the HIT output is 0. If you have seen a sequence of 2 0's for example, and then see a 1, that basically means that you are starting over because your sequence of 0's was interrupted. However, you don't go back to state S_x , instead you go to state S_1 to indicate that you just saw your first 1. The same logic explains why if you are in state S_{000} and you receive another 0, you go back to state S_{000} because the last 3 inputs you received were all 0's. You don't need to start counting 0's from scratch as if you hadn't previously gotten any.

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i Answers are displayed within the problem

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