

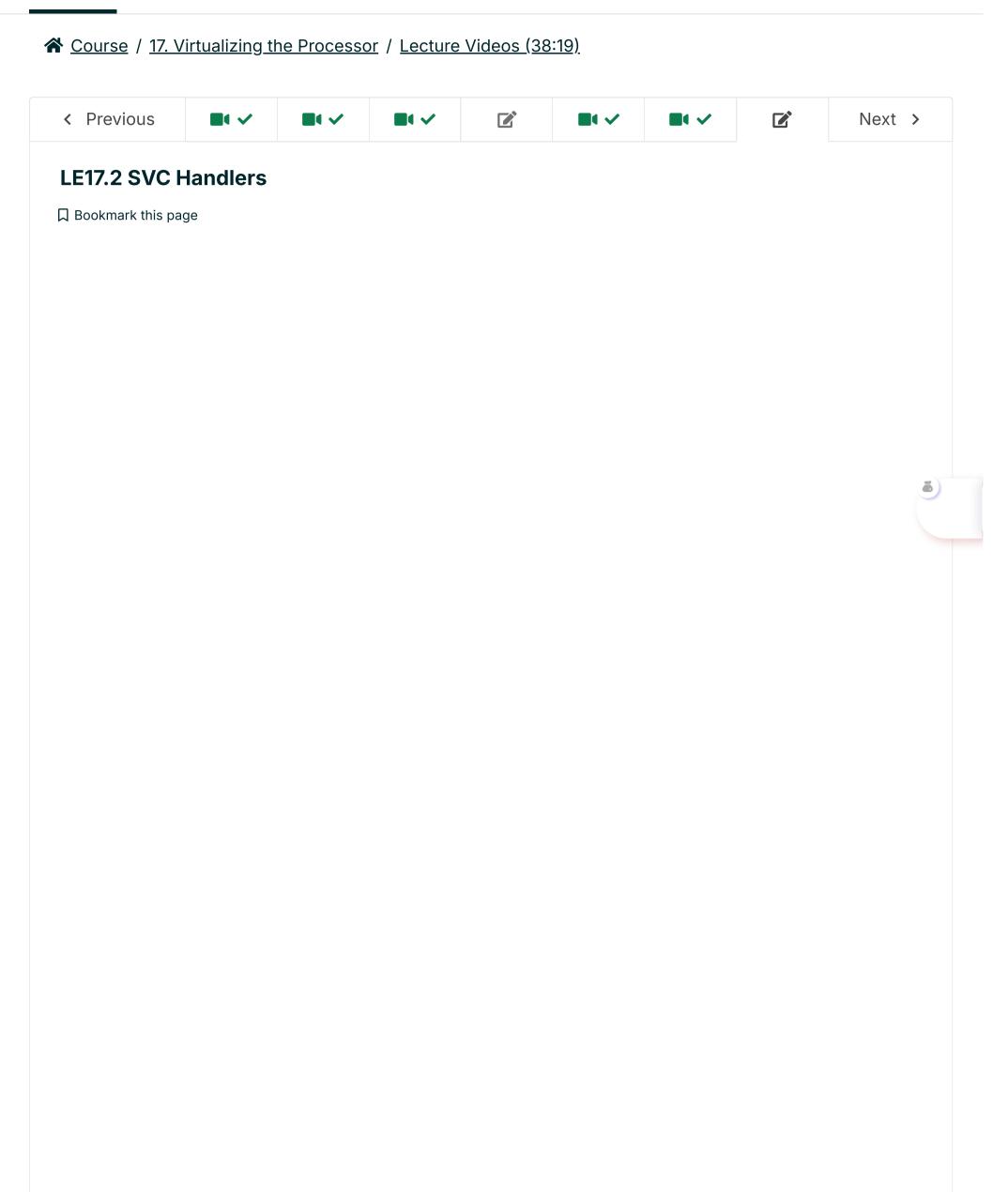
Computation Structures 3: Computer Organization

<u>Help</u>





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LE17.2.1: Wait SVC Handler

0.0/1.0 point (ungraded)

The WAIT() supervisor call expects the address of a user's memory location in R0. The WAIT handler examines the value of that memory location and if it's positive, decrements the value in the memory location by 1 and resumes execution of the user's program at the instruction following the WAIT() SVC. If the memory location is not positive, the handler should arrange to re-execute the WAIT() SVC next time the user's program runs and then yield the remainder of the current time slice.

Your goal is to fill in the required code for Wait_h(), the kernel-mode handler for the WAIT SVC. You can use the kernel procedure GetUserLocn(addr) to get the values of a location in the user's memory and SetUserLocation(addr,v) to set the value of a memory location. The kernel subroutine Scheduler() is also available for your use. Select the correct choice of instructions to include in the if and else clauses of the Wait_h() stub shown below. Make sure to include all the required instructions. Assume that the selected instructions are executed in the order they appear.

```
Wait_h() {
   int locn = UserMState.Regs[0]; // user location to test
   int value = GetUserLocation(locn); // current value
   if (value > 0) {
     // if clause code
   } else {
     // else clause code
if clause code (select all that apply):
      UserMState.Regs[XP] = UserMState.Regs[XP] - 4;
      SetUserLocation(locn, value)
      SetUserLocation(locn, value-1)
      GetUserLocation(locn)
      Scheduler();
else clause code (select all that apply):
      UserMState.Regs[XP] = UserMState.Regs[XP] - 4;
      SetUserLocation(locn, value)
      SetUserLocation(locn, value-1)
      GetUserLocation(locn)
      Scheduler();
```

of the user's program should resume at the instruction following the WAIT() SVC. This is achieved by making a call to SetUserLocation with a value of value - 1 inside the if clause.

If the memory location is not positive, then the handler should re-execute the WAIT() SVC the next time the user's program runs. This is achieved by decrementing the value of the user's XP register by 4 so that it re-executes the supervisor call. This is followed by a call to Scheduler() in order to schedule the next process to run right away rather than waiting for the time slice of the current process to complete before moving on to the next process.

Submit

1 Answers are displayed within the problem

LE17.2.2: YieldN SVC Handler

0.0/1.0 point (ungraded)

The Yield() SVC can be used in user-mode programs on a time-sharing system to give up the remainder of their current time slice. The kernel implementation of the Yield simply calls the kernel Scheduler() routine to choose another process to execute. When the yielding process is next scheduled, user-mode execution resumes with the instruction following the Yield() SVC.

Complete the code for the handler for a new SVC, YieldN(), which expects a numeric value, N, in the user's RO and behaves as if the user program had contained N consecutive Yield() SVCs. When execution resumes following the completion of YieldN(), RO should contain O. Select the correct choice of instructions to include in the if and else clauses of the YieldN_h() stub shown below. Make sure to include all the required instructions. Assume that the selected instructions are executed in the order they appear.

```
YieldN_h() {
  if (UserMState.Regs[0] > 0) {
    // if clause code
  } else {
    // else clause code
  }
}
```

if clause code (select all that apply):

UserMState.Regs[0] = UserMState.Regs[0] - 1;	
UserMState.Regs[0] = UserMState.Regs[0] - N;	
UserMState.Regs[0] = 0;	
<pre>UserMState.Regs[XP] = UserMState.Regs[XP] - 4;</pre>	
UserMState.Regs[LP] = UserMState.Regs[LP] - 4;	

Scheduler();

else clause code (select all that apply):

YieldN();

UserMState.Regs[0] = UserMState.Regs[0] - 1;

⊞ Calculator

nes. To do this, each time it executes the instruction, it needs to decrement the remove by decrementing register R0 by 1. It then needs to re-execute the call to YieldNecrementing the XP register by 4 to point to that instruction rather than the one imprough, it needs to call Scheduler() in order to yield to other processes. This happers > 0. When this is not true, that means that we are done and R0 should be set to 0 init. Submit Answers are displayed within the problem iscussion pic: 17. Virtualizing the Processor / LE17.2 SVC Handlers	maining count by 1. This is I(), this is achieved by mediately after it. Each time as long as the value of R0
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Scheduler();	
YieldN();	
UserMState.Regs[LP] = UserMState.Regs[LP] - 4;	
UserMState.Regs[XP] = UserMState.Regs[XP] - 4;	
✓	
UserMState.Regs[0] = 0;	

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