LE9.5

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LE9.5.1: Branch Instructions

1/1 point (ungraded)

- Summary of Instruction Formats (PDF)
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Consider the execution of a short program that loops to sum the elements of an array with 4 elements. The first element of the array is stored at location 0×2000.

```
  . = 0                 // first instruction is at location 0
  ADDC(r31,array,r0)    // r0 = pointer to next array element
  ADDC(r31,4,r1)        // r1 = number of array elements remaining
  ADDC(r31,0,r2)        // r2 = accumulated sum
loop:
  LD(r0,0,r3)           // load next value from array
  ADD(r3,r2,r2)         // add to sum
  ADDC(r0,4,r0)         // increment pointer to next word
  SUBC(r1,1,r1)         // decrement counter
  BNE(r1,loop,r31)      // loop if more elements to go
  ST(r2,result,R31)     // write result to memory
  // execution stops here

  . = 0x2000
array:
  LONG(1)               // array[0] = 1
  LONG(2)               // array[1] = 2
  LONG(3)               // array[2] = 3
  LONG(4)               // array[3] = 4
result:
  LONG(0)               // where result will be stored
```

Program execution starts with the first instruction and halts after execution of the ST instruction.

(A) What value does the assembler give the label "loop"? 0×000C ✔ Answer: 0xc

(B) After execution, number of times LD is executed? 4 ✔ Answer: 4

(C) After execution, value left in r0? 0×2010 ✔ Answer: 0×2010

(D) After execution, value left in r1? 0×0 ✔ Answer: 0

(E) After execution, value left in r2? 0×0A ✔ Answer: 10

(F) After execution, value left in r3? 0×04 ✔ Answer: 4

Explanation
"loop:" labels the fourth instruction of the program and the first instruction is at location 0. So the value of the symbol "loop" is 0xC.
At the end of the loop, the loop counter (R1) is decremented and if it's still non-zero, we branch back for another loop iteration. The initial value for R1 was 4, so the LD instruction is executed 4 times.
The array pointer (R0) is incremented by 4 on each loop iteration. The initial value of the pointer was 0×2000 and there were four loop iterations, so the value left in R0 is 0×2010.
The loop only terminates when the value of the loop counter (R1) is 0. So at the end of execution, the value left in R1 is 0.
R2 is the accumulated sum of the array values = 1 + 2 + 3 + 4 = 10.
R3 is used to hold the value of the array element accessed during each iteration. The last array element accessed had the value 4.

The encoding for the OPCODE, RC and RA fields of a branch instruction is just like the encodings for other I

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instructions. Figuring out the value for the 16-bit constant field takes a little more work. The offset value is the number of words between the instruction following the branch (ST in this example) to target instruction (LD in this example). Positive values indicate a forward branch to a subsequent location with a higher address; negative offset values indicate a backward branch to a location with a lower address.

In this example, we'd start counting instructions backwards from the store instruction until we reached the LD instruction. Since it's a backwards branch, we'd encode the count as a negative number in the 16-bit constant field of the BNE instruction.

(G) What is the binary encoding for `BNE(r1,loop,r31)`? 0b01110111111000011111 ✔ Answer: 0×77E1FFFB

Explanation
Here's a diagram showing how the offset of -5 is calculated:

```
offset    instruction
             . = 0                 // first instruction is at location 0
             ADDC(r31,array,r0)    // r0 = pointer to next array element
             ADDC(r31,4,r1)        // r1 = number of array elements remaining
             ADDC(r31,0,r2)        // r2 = accumulated sum
          loop:
-5           LD(r0,0,r3)           // load next value from array
-4           ADD(r3,r2,r2)         // add to sum
-3           ADDC(r0,4,r0)         // increment pointer to next word
-2           SUBC(r1,1,r1)         // decrement counter
-1           BNE(r1,loop,r31)      // loop if more elements to go
 0           ST(r2,result,R31)     // write result to memory
```

The 32-bit binary encoding of the BNE instruction is

```
      opcode  RC    RA    constant (= -5)
      011101 11111 00001 1111 1111 1111 1011
```

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- Value of R31 — 6
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- R31? — 5
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- Reg[Rc] ←Mem[EA] — 2
  "The location in memory specified by EA is read into register Rc" Does this mean: Rc register now has memory address stored or th...
- Value of RC in this branch — 6
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- BNE

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BNE instruction is supposed to work like that NPC = 0x1c+4 = 0x20 if (REG[RT]!=0) pc = NPC + 4*offset if the condition is true, the...

hallo Ich bin mohammed hany elemam .	2
reversing the bits in r0 Hi everybody, In the assembly program provided in the " Introduction to Bsim" section what does it mean to reverse the bits in r0 ?	2
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