

LE8.1.1 Asymptotic Latency and Throughput

1/1 point (ungraded)

If we account for fan in limitations but ignore wire delays, what is the asymptotic latency of the fastest combinational N-input AND circuit we can build?

Asymptotic latency of N-input AND:

$\log(N)$

✓ Answer: $\log(N)$

You entered:

$\log(N)$

Explanation

Accounting for the fan in limitations means that if we have a gate with N inputs, we can't just assume that its propagation delay is the same as a 2-input gate. The way we model this is by turning an N-input gate into a tree of 2-input gates. The equivalent tree of 2-input AND gates would have $\log(N)$ levels in the tree in order to arrive at enough 2-input AND gates to allow for N inputs. Each level of the tree must complete its propagation delay before the next level can begin its computation, therefore the latency of the tree is $\log(N)$.

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LE8.1.2 Asymptotic Latency and Throughput

1 point possible (ungraded)

Is $\Theta(\log_2 N)$ the same as $\Theta(\log_{10} N)$

☐ Yes

☐ No

☐ Only for some N

LE8.1.3 Asymptotic Latency and Throughput

2 points possible (ungraded)

A combinational multiplier is pipelined for maximum throughput. If the multiplier accepts two N-bit operands, what is the appropriate “order of” notation for its throughput and latency?

Throughput $\Theta(\dots)$:

Latency $\Theta(\dots)$:

Discussion

Topic: 8. Design Tradeoffs / LE8.1

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Accounting for the fan in limitations

4

"... if we have a gate with N inputs, we can't just assume that its propagation delay is the same as a..."