



Practical Network Defense

Master's degree in Cybersecurity 2018-19

Network traffic regulation with firewalls

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Today's agenda

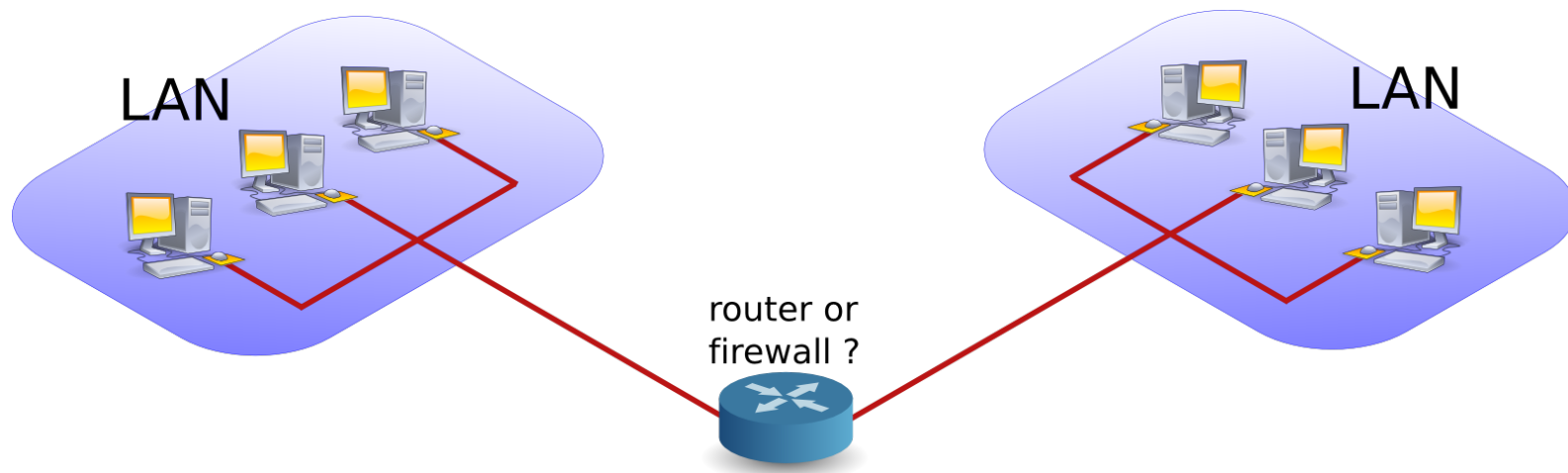
- Traffic regulation
 - Packet filtering
- Filter rules
- Stateful firewall
- Other types of firewall
 - Application-level filtering
 - Circuit-level gateway



Traffic regulation

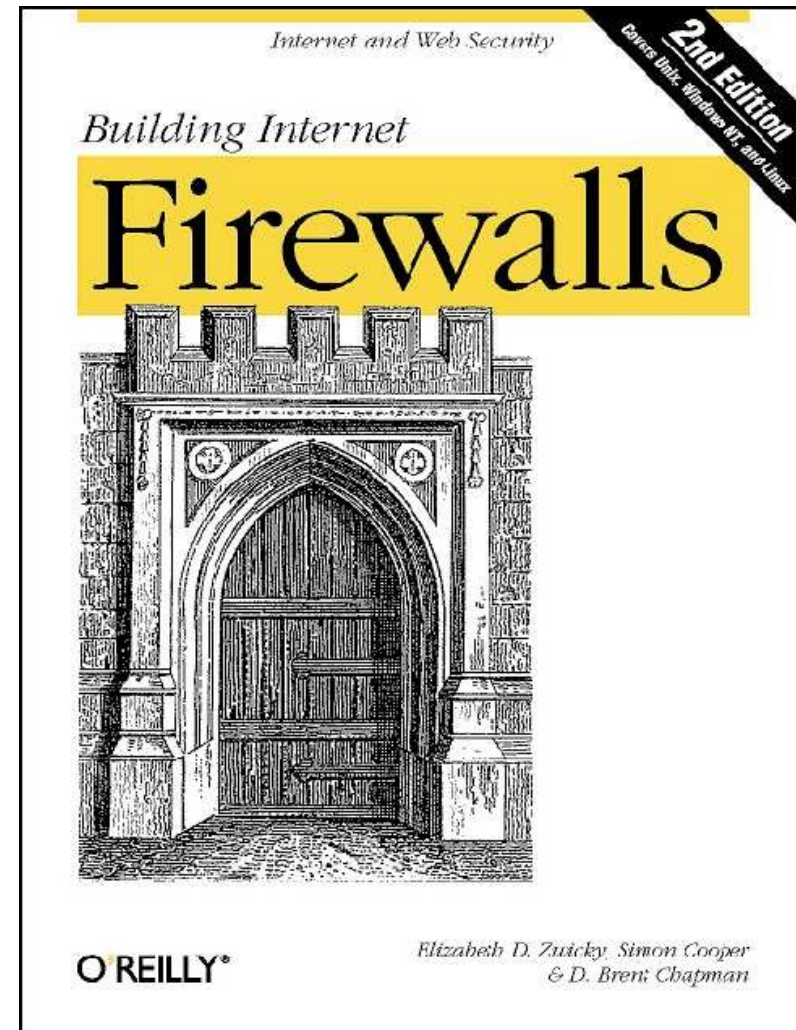
Regulate traffic: routers and firewalls

- A router is a device that connects two networks
- A firewall is a device that besides acting as a router, also contains (and implements) rules to determine whether packets are allowed to travel from one network to another
 - Router also can perform some form of screening (packet filter)



Why firewalls?

- Restricts access from the outside
 - Internet = millions of people together → bad things happen
- Prevents attackers from getting too close
- Restricts people from leaving



Secure traffic regulation

- To attain a certain level of network security, you can:
 - Regulate which traffic is allowed (sources, destinations, services, ...)
 - Protect the traffic by encryption
 - Monitor the traffic for “bad behaviour”
 - Monitor the hosts for “bad behaviour”
- The choice will depend on the security policy to be fulfilled (particularly the CIA targets).



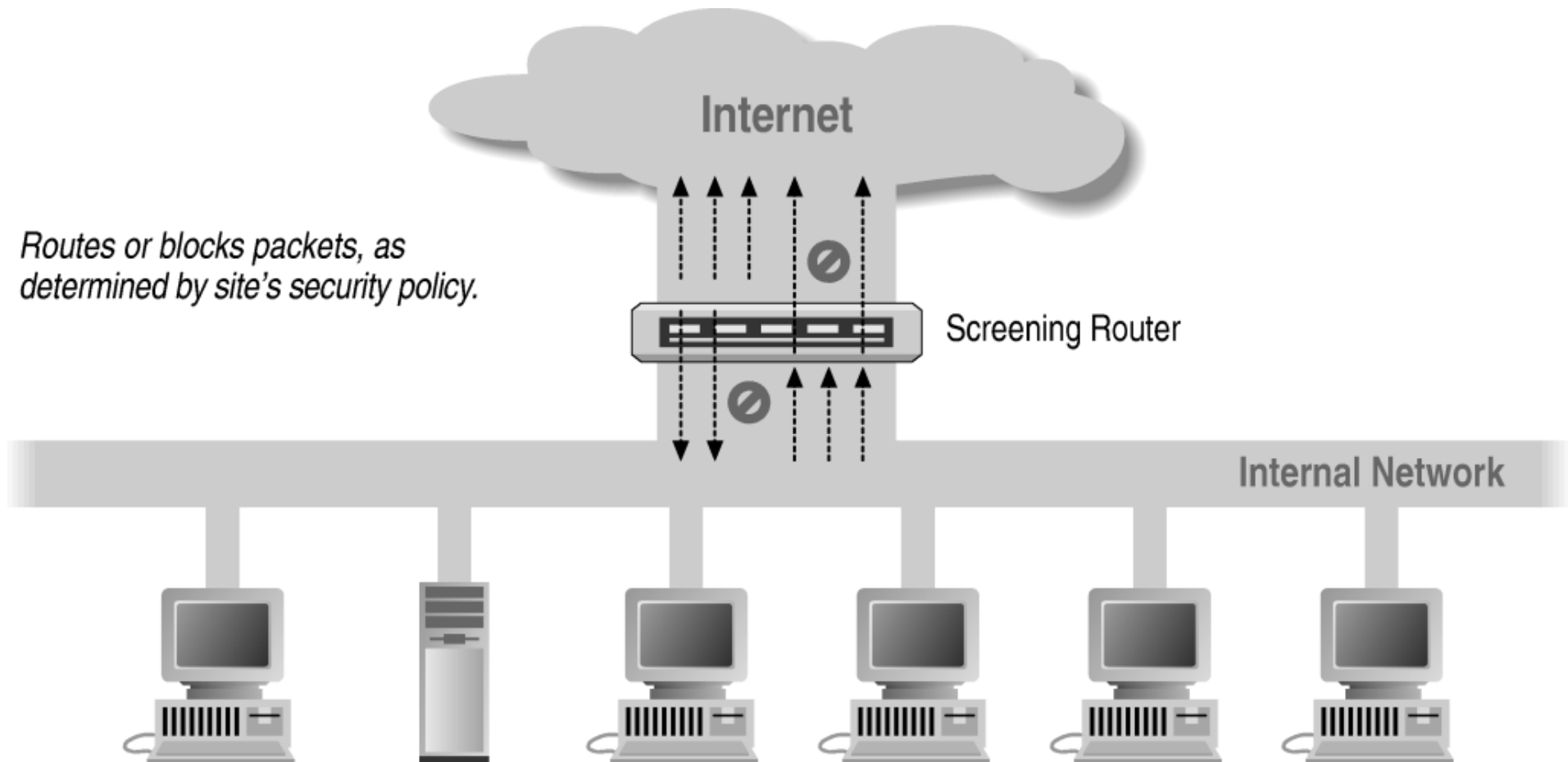
Firewall Design & Architecture Issues

- Least privilege
- Defense in depth
- Choke point
- Weakest links
- Fail-safe stance
- Universal participation
- Diversity of defense
- Simplicity

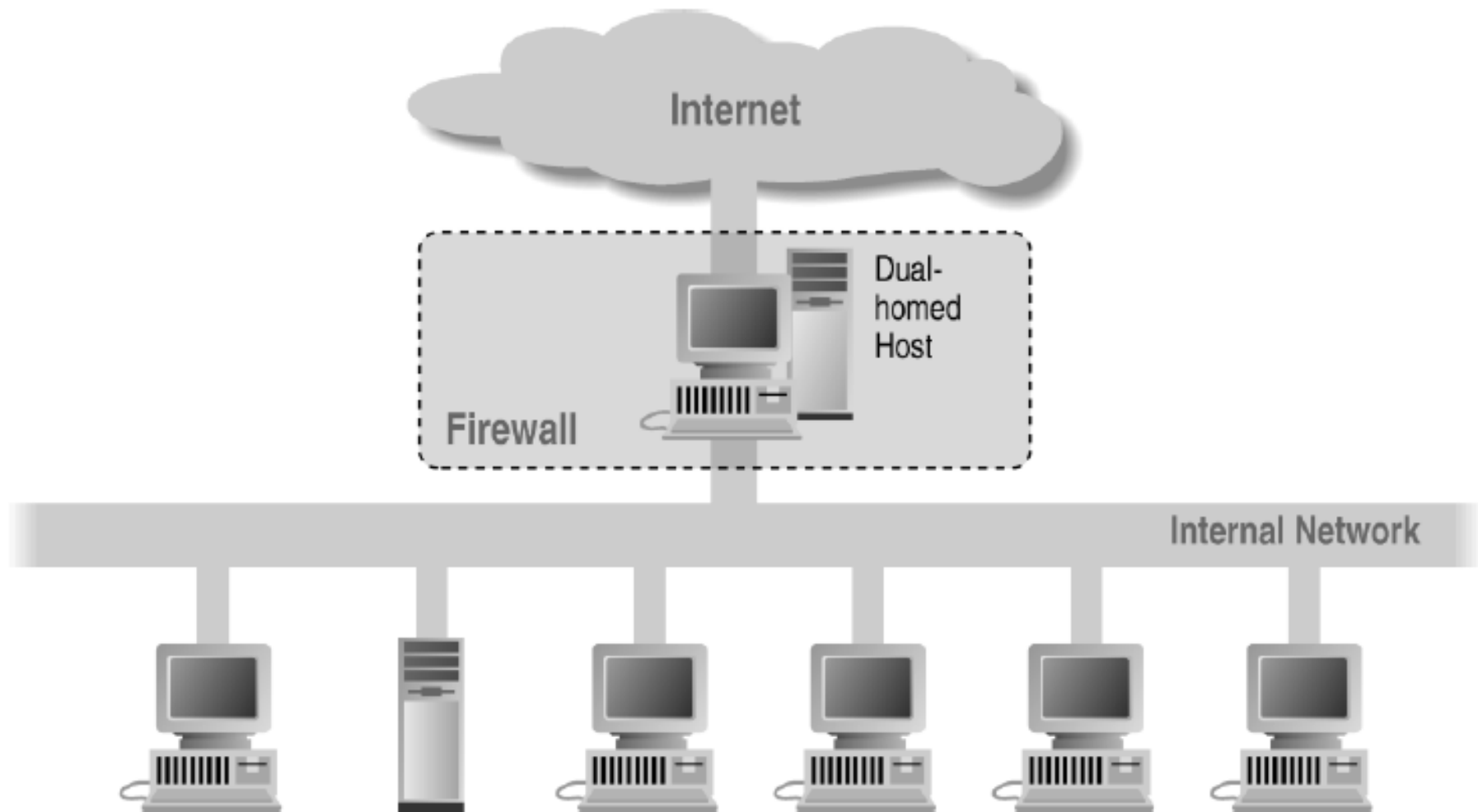
Host based packet filters

- Kind of firewall that disciplines the traffic in/out a single host
- It specifies the packets that can be received and sent
 - Ex: iptables, windows firewall and all the so called “personal firewalls”
- Vendor products generally work per-app: each installed application has a known policy that has to obey

Screening router (ACL-based)



Dual-homed host



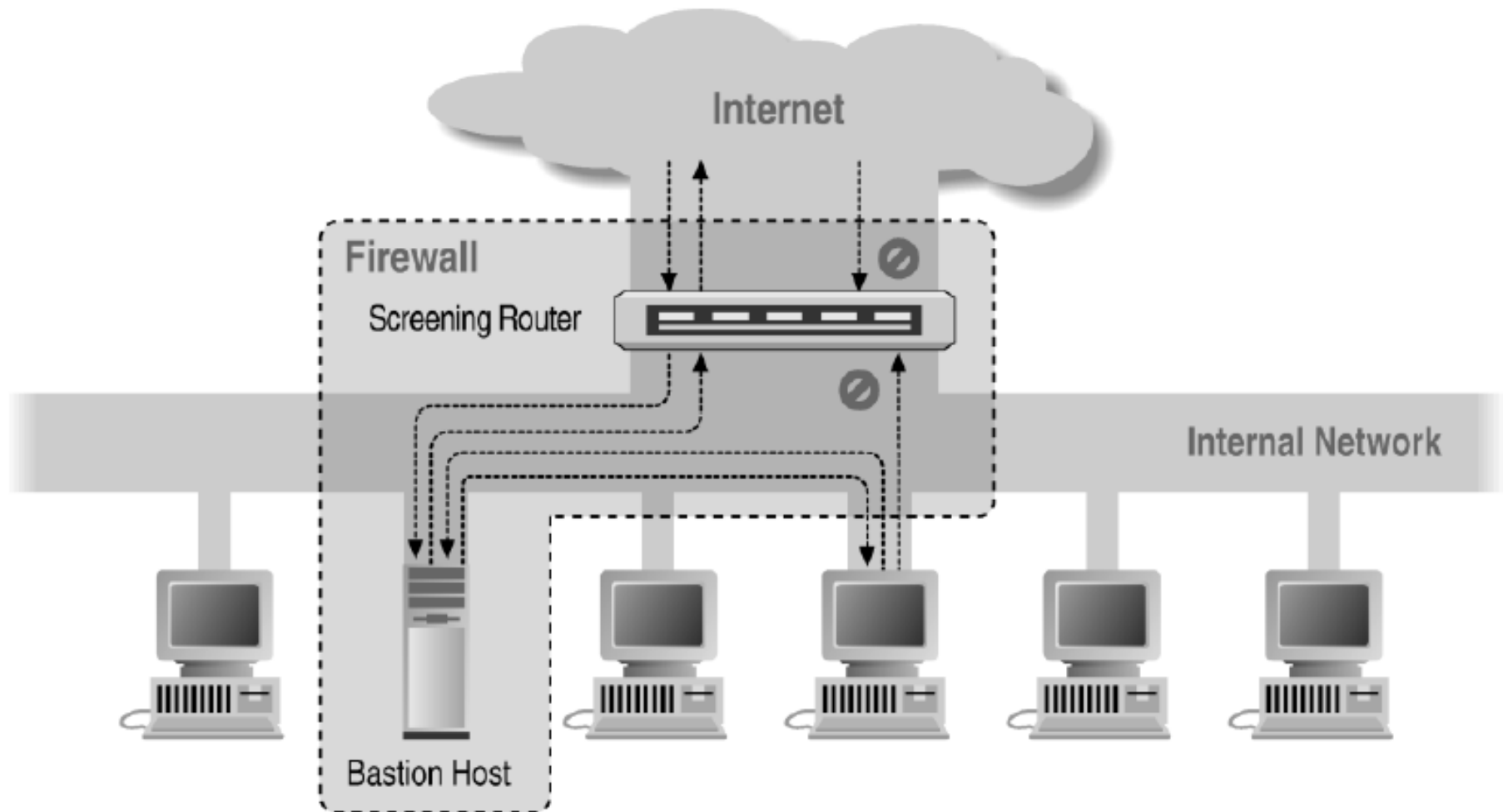
Bastion host

- Hardened computer used to deal with all traffic coming to a protected network from outside
 - Hardening is the task of reducing or removing vulnerabilities in a computer system:
 - Shutting down unused or dangerous services. Strengthening access controls on vital files
 - Removing unnecessary account permissions
 - Using "stricter" configurations for vulnerable components, such as DNS, sendmail, FTP, Apache Tomcat, etc
- Specially suitable for use as Application Proxy Gateways

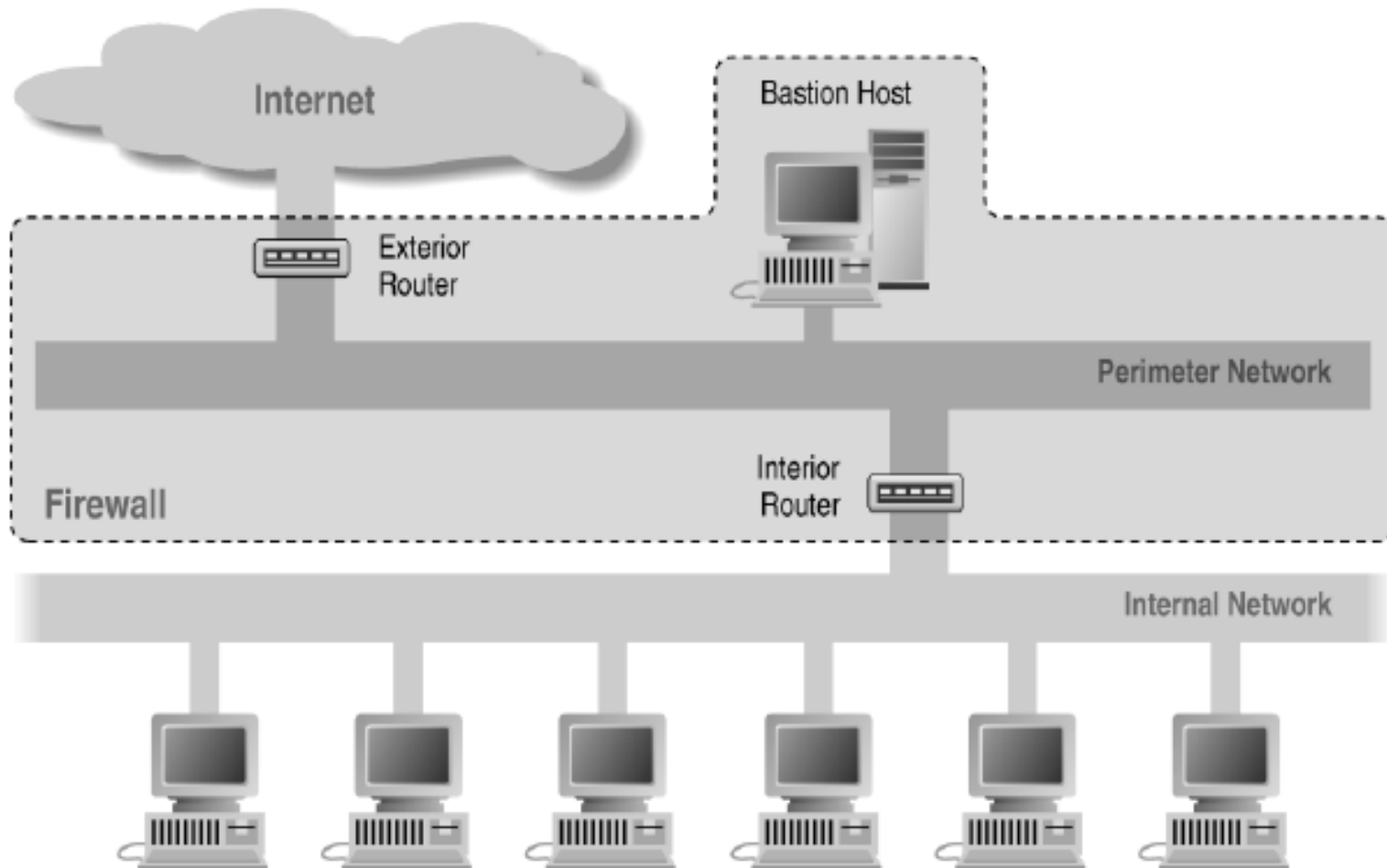
What is a DMZ

- DMZ (demilitarized zone)
 - Computer host or small network inserted as a “neutral zone” between a company’s private network and the outside public network
 - Network construct that provides secure segregation of networks that host services for users, visitors, or partners
- DMZ use has become a necessary method of providing a multilayered, **defense-in-depth** approach to security

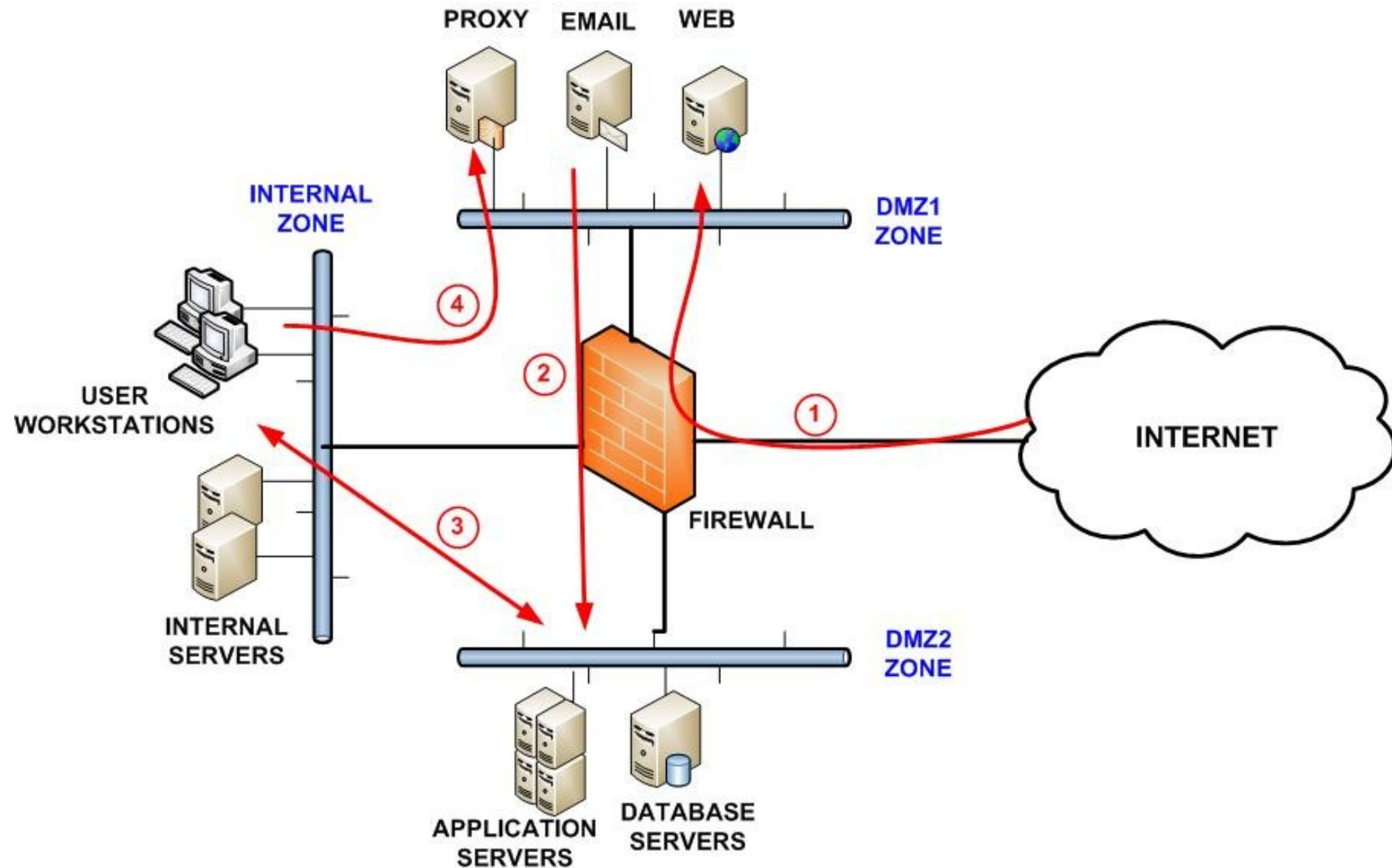
DMZ as a screened Host



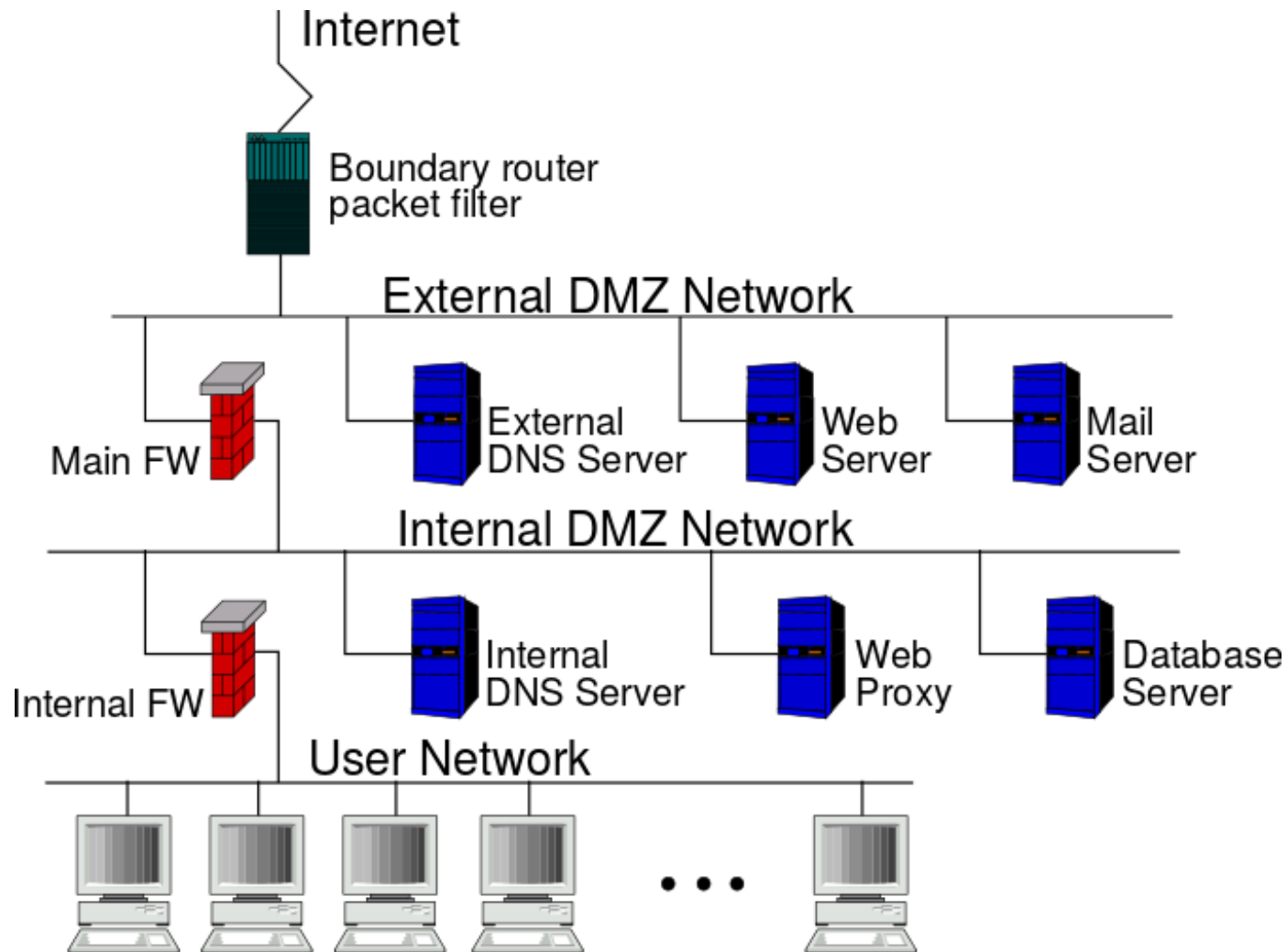
Screened Subnet Using Two Routers/Firewalls



DMZ to segment the network

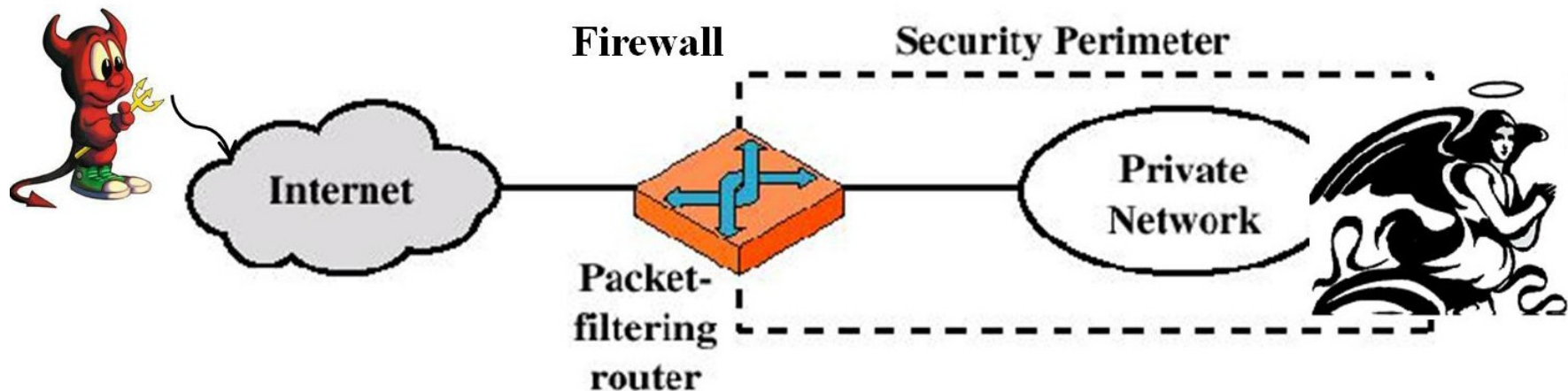


Security in depth: split DMZ



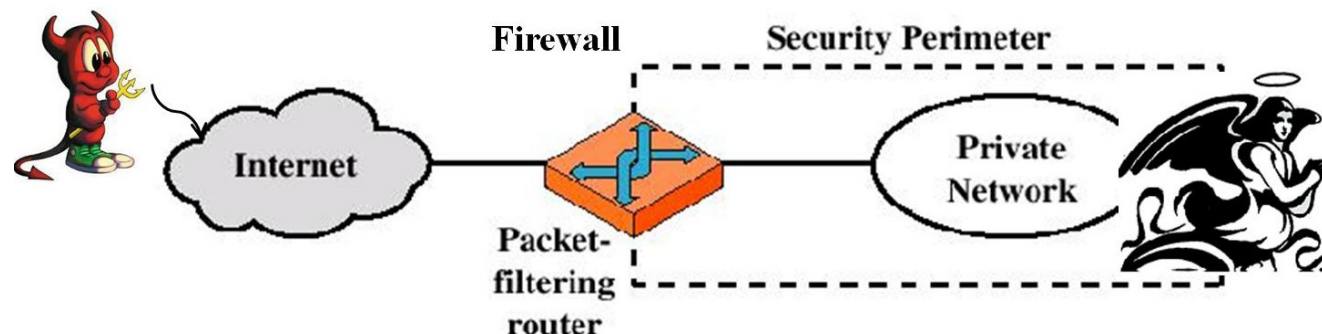
A simple plan for network security

- Use a firewall to filter ingoing and outgoing traffic between “your” network (or individual PC) and the Internet



Assumptions

1. You have security policy stating what is allowed and not allowed.
2. You can identify the “good” and the “bad” traffic by its IP-address, TCP port numbers, etc, ...
3. The firewall itself is immune to penetration.
 - A question of assurance – needs for a trusted system, secure OS etc.

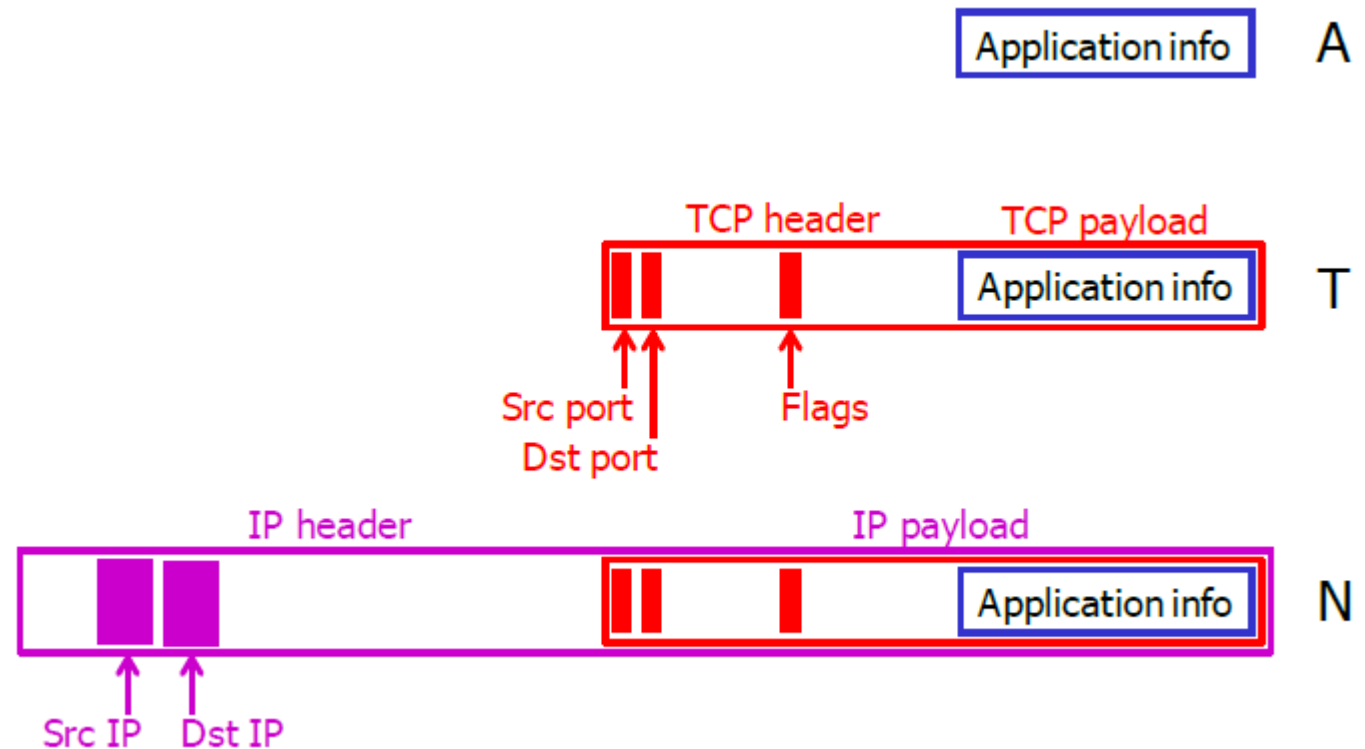
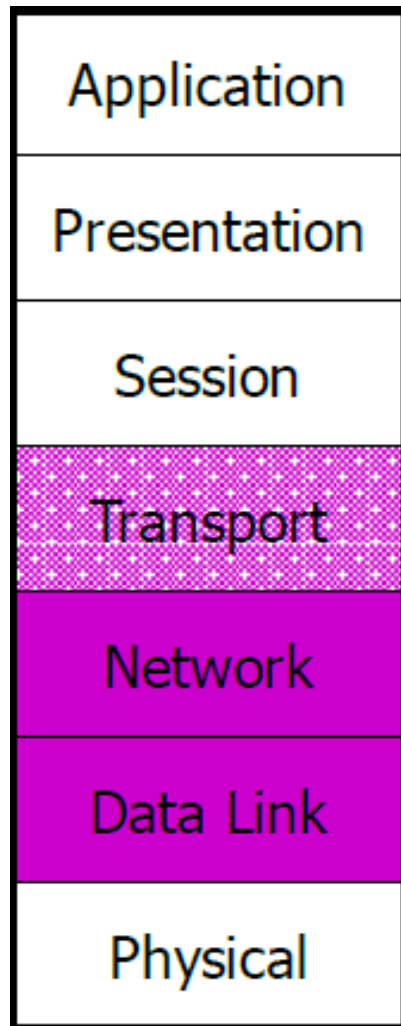


Packet filters (stateless firewall)

- Drop packets based on their source or destination addresses or port numbers or flags
- No context, only contents
- Can operate on
 - incoming interface
 - outgoing interface
 - both
- Check packets with fake IP addresses:
 - from outside (“ingress filtering”)
 - from inside (“egress filtering”)



Packet filters operating layers



Three-step process

1. Know your policy
2. Translate the policy in a formal language
 - E.g.: logical expression on packet fields
3. Rewrite the policy in terms of the firewall syntax

General mechanism:

- Rules are checked from top to bottom
- The first matching rule is applied
- One implicit rule is assumed if no rule matches
 - Block/Allow everything

action	ourhost	port	theirhost	port	comment
block	*	*	*	*	<i>default</i>

Example

- Policy:
 - allow inbound email (SMTP, port 25) only to our-gateway machine: **Mailgw**
 - refuse all traffic from a known spamming site: **demon**
- Possible rules:

action	ourhost	port	theirhost	port	comment
block	*	*	demon	*	<i>don't trust spammers</i>
allow	Mailgw	25	*	*	<i>connection to our SMTP</i>

Example, continued

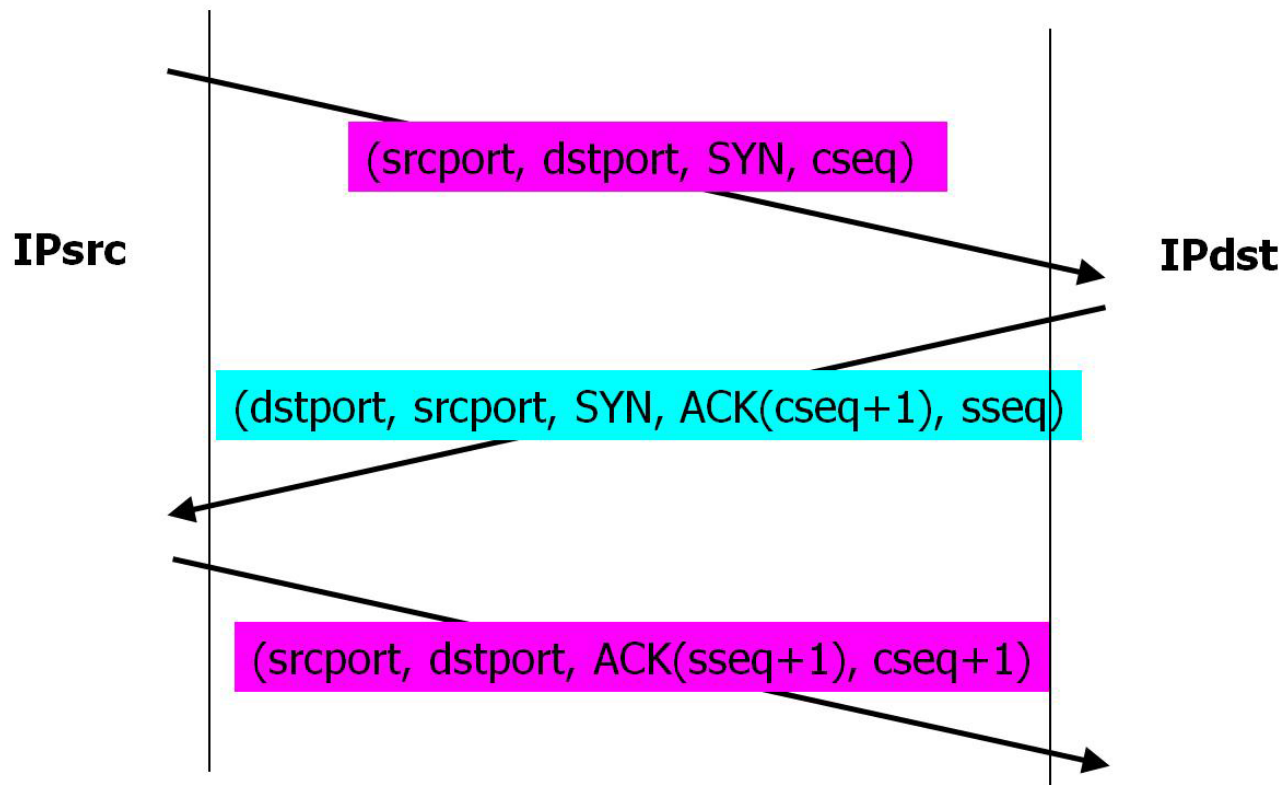
- Add the policy:
 - any inside host can send mail to the outside

action	ourhost	port	theirhost	port	comment
allow	*	*	*	25	<i>connection to their SMTP</i>

- Very bad: we can not control the type of traffic originated from port 25 and coming from the outside
- Then: rules have to specify the direction of the traffic

How to check the direction of TCP?

- Consider the TCP flags



Example with traffic direction

- We distinguish the replies to our SMTP connection considering the ACK flag

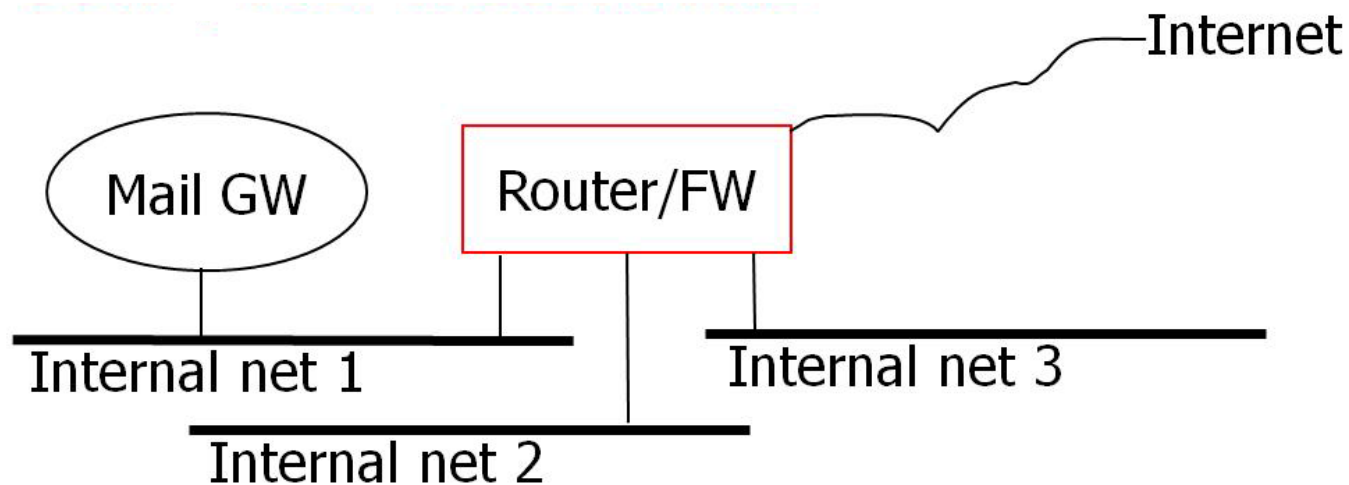
action	src	port	dest	port	flags	comment
allow	{our hosts}	*	*	25		<i>connection to their SMTP</i>
allow	*	25	*	*	ACK	<i>their replies</i>
block	*	*	*	*		<i>default</i>

- Very easy case...



Filter rules for network firewalls

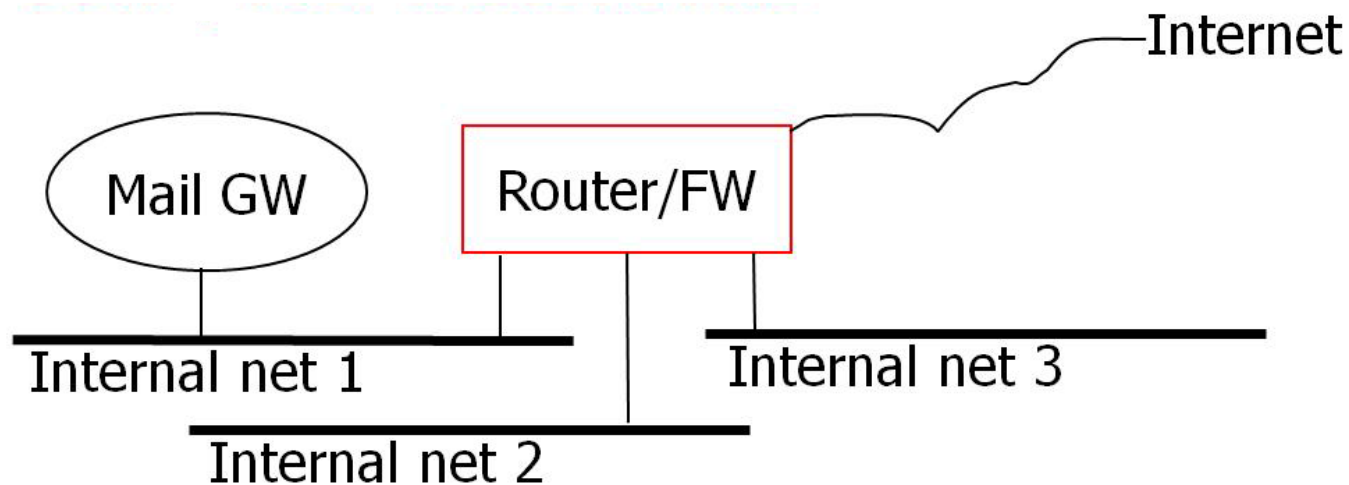
More complex network topology



- Policy:

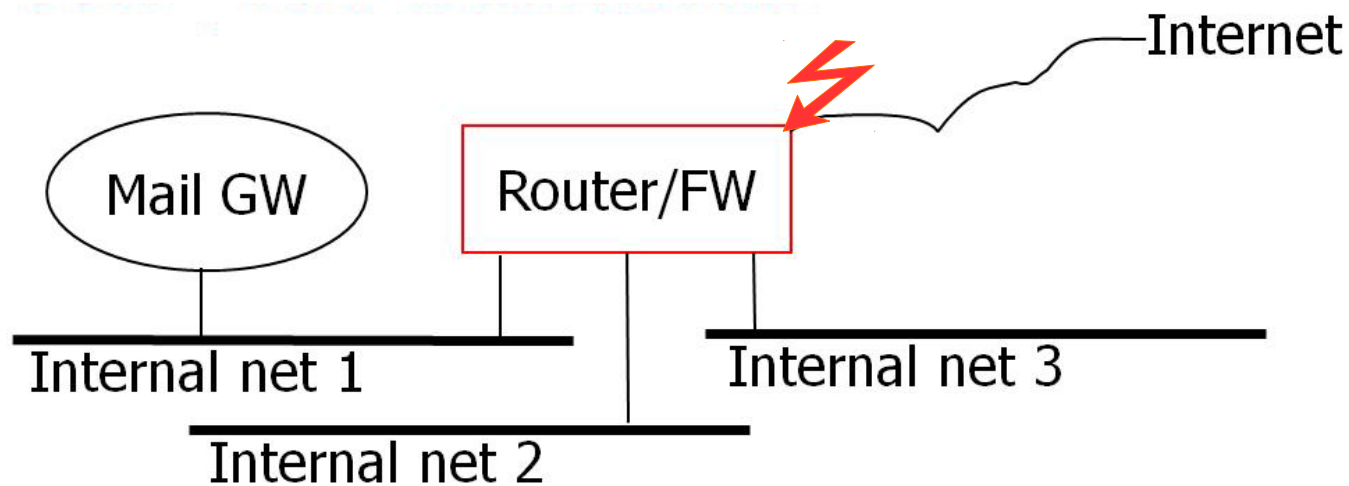
- Internal Net 1 is a DMZ and only hosts Mail GW
- Very limited connections between Mail GW and Internet (only partner servers)
- Limited connections allowed between Mail GW and net 2 and net 3
- Anything can pass between net 2 and net 3
- Outgoing requests only between net 2 or net 3 and the link to the Internet

Requirements



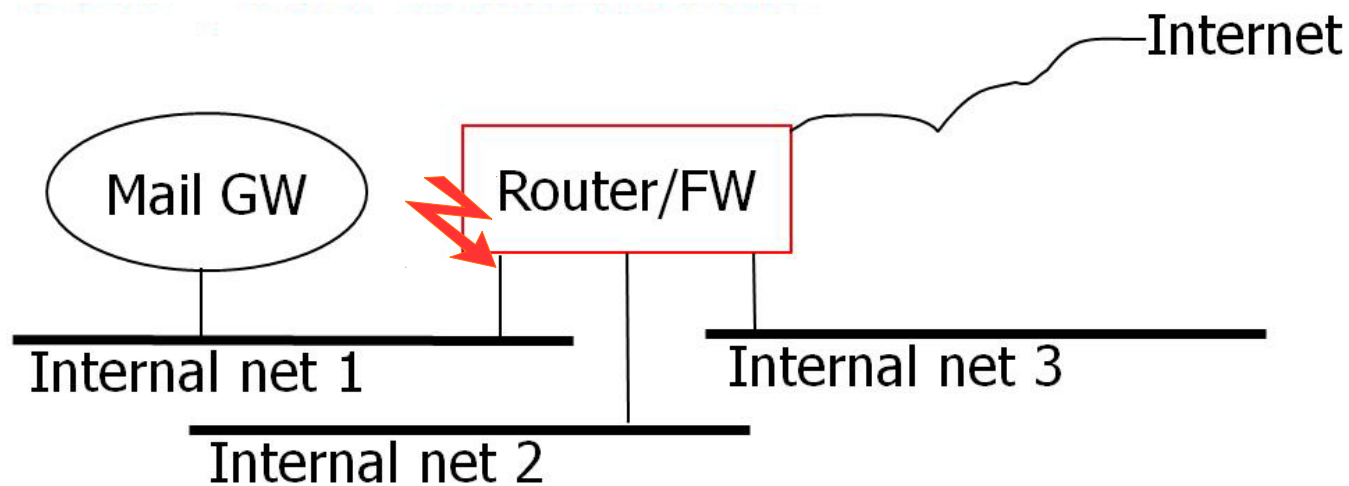
- We cannot only consider where packets have to go (destination → **output filtering**)
 - Open access to net 2 only allowed for traffic with source address in net 3
 - No way to avoid fake source addresses (*address spoofing*) from outside
- We need to define rules based on **from where packets are arriving**, (source → **input filtering**)

Interface towards Internet



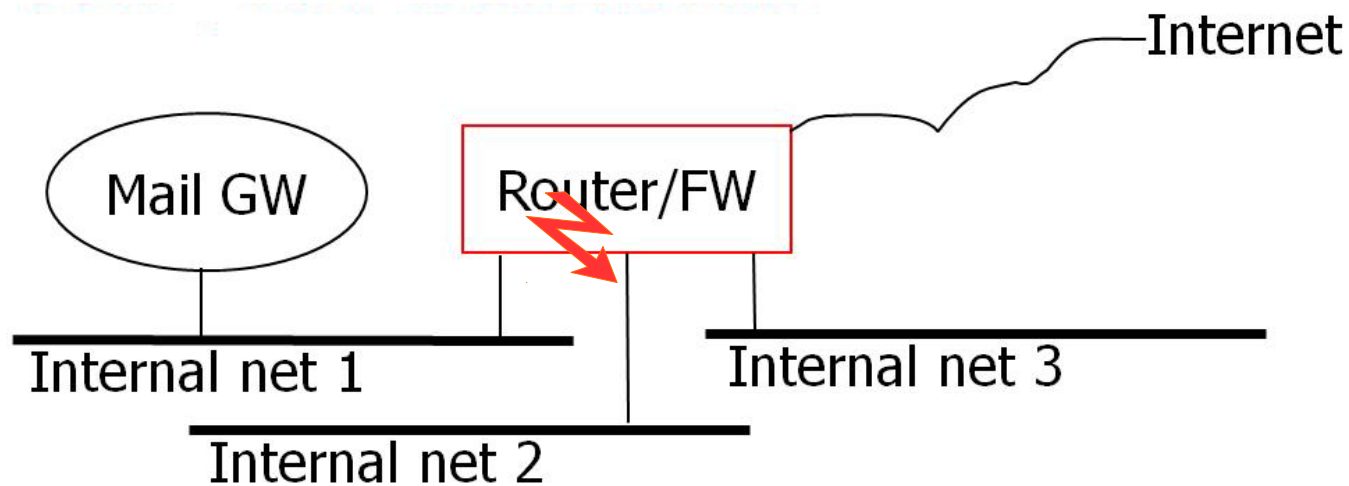
Action	IPsrc	srcport	IPdst	dstport	flags
block	"Net 1"	*	*	*	
block	"Net 2"	*	*	*	
block	"Net 3"	*	*	*	
allow	*	*	GW	25	
allow	*	*	"Net 2"	*	ACK
allow	*	*	"Net 3"	*	ACK

Interface on net 1



Action	IPsrc	srcport	IPdst	dstport	flags
allow	GW	*	"partners"	25	
allow	GW	*	"Net 2"	*	ACK
allow	GW	*	"Net 3"	*	ACK
block	GW	*	"Net 2"	*	
block	GW	*	"Net 3"	*	
allow	GW	*	*	*	

Interface on net 2 (net 3 is similar)



action	IPsrc	srcport	IPdst	dstport	flags
allow	"Net 2"	*	*	*	
block	*	*	*	*	

Problems with Packet Filters

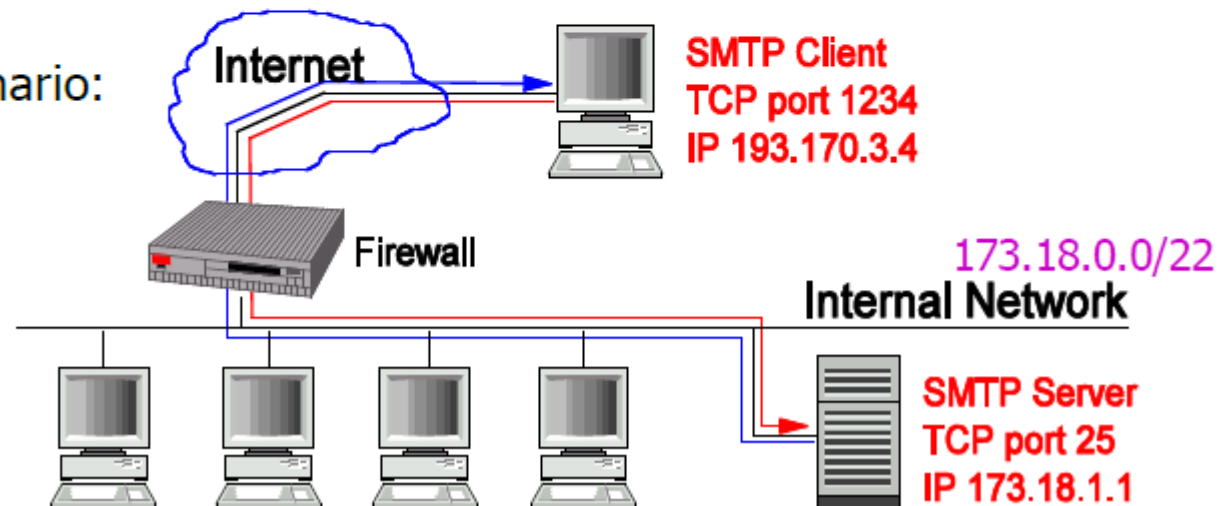
- Only a small number of parameters
 - it is (unfortunately) easy to specify filtering rules which are too specific or too general
- Payload of TCP packet is not inspected
 - No protection against attacks based on upper-layer vulnerabilities
- Limited logging ability (restricted to the few parameters used by the filter)
- No authentication facilities
- Susceptible to attacks based on vulnerabilities in various implementations of TCP and/or IP

Filter rules, 1

- Example: Filtering in- and outgoing SMTP traffic. Try rules:

Rule	In/out	IPsrc	IPdst	Proto	dstport	Action
A	Inward	External	Internal	TCP	25	Allow
B	Outward	Internal	External	TCP	>1023	Allow
C	Outward	Internal	External	TCP	25	Allow
D	Inward	External	Internal	TCP	>1023	Allow
E	*	*	*	*	*	Block

- Scenario:



Filter rules, 2

- Example: Filtering in- and outgoing SMTP traffic. Try rules:

Rule	In/out	IPsrc	IPdst	Proto	dstport	Action
A	Inward	External	Internal	TCP	25	Allow
B	Outward	Internal	External	TCP	>1023	Allow
C	Outward	Internal	External	TCP	25	Allow
D	Inward	External	Internal	TCP	>1023	Allow
E	*	*	*	*	*	Block

- Scenario:

Packet	In/out	IPsrc	IPdst	Proto	dstport	Action
1	Inward	193.170.3.4	173.18.1.1	TCP	25	Allow (A)
2	Outward	173.18.1.1	193.170.3.4	TCP	1234	Allow (B)

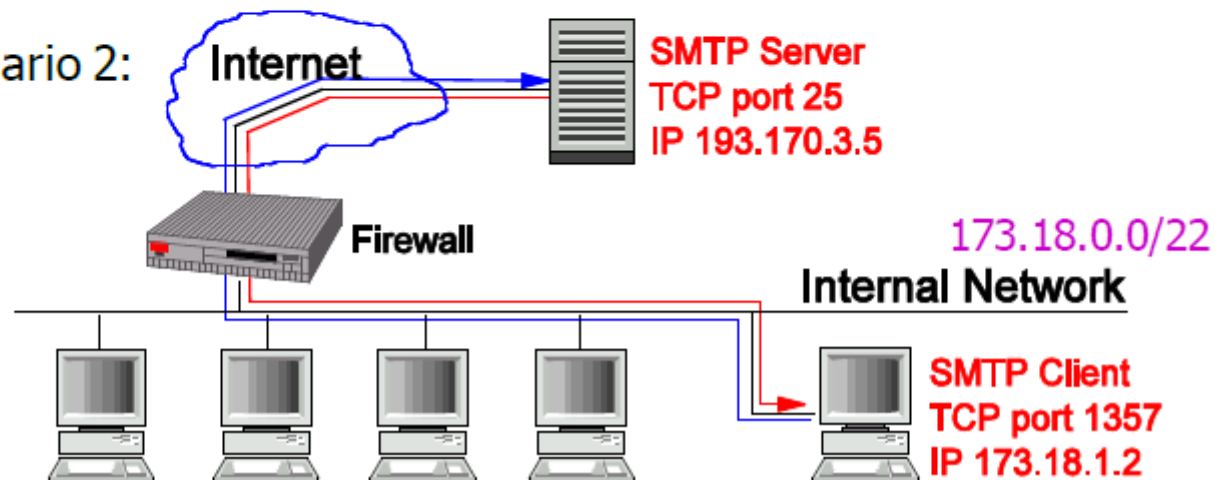
- Conclusion: This looks OK!

Filter rules, 3

- Example: Filtering in- and outgoing SMTP traffic. Try rules:

Rule	In/out	IPsrc	IPdst	Proto	dstport	Action
A	Inward	External	Internal	TCP	25	Allow
B	Outward	Internal	External	TCP	>1023	Allow
C	Outward	Internal	External	TCP	25	Allow
D	Inward	External	Internal	TCP	>1023	Allow
E	*	*	*	*	*	Block

- Scenario 2:



Filter rules, 4

- Example: Filtering in- and outgoing SMTP traffic. Try rules:

Rule	In/out	IPsrc	IPdst	Proto	dstport	Action
A	Inward	External	Internal	TCP	25	Allow
B	Outward	Internal	External	TCP	>1023	Allow
C	Outward	Internal	External	TCP	25	Allow
D	Inward	External	Internal	TCP	>1023	Allow
E	*	*	*	*	*	Block

- Scenario 2:

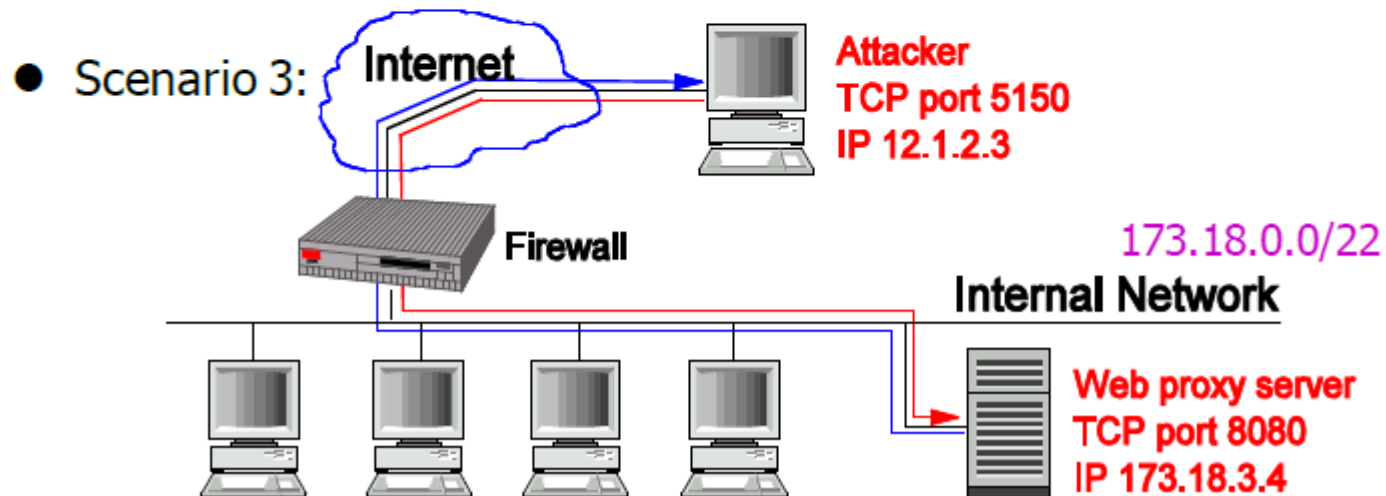
Packet	In/out	IPsrc	IPdst	Proto	dstport	Action
3	Outward	173.18.1.2	193.170.3.5	TCP	25	Allow (C)
4	Inward	193.170.3.5	173.18.1.2	TCP	1357	Allow (D)

- Conclusion: This also looks OK!

Filter rules, 5

- Example: Filtering in- and outgoing SMTP traffic. Try rules:

Rule	In/out	IPsrc	IPdst	Proto	dstport	Action
A	Inward	External	Internal	TCP	25	Allow
B	Outward	Internal	External	TCP	>1023	Allow
C	Outward	Internal	External	TCP	25	Allow
D	Inward	External	Internal	TCP	>1023	Allow
E	*	*	*	*	*	Block



Filter rules, 6

- Example: Filtering in- and outgoing SMTP traffic. Try rules:

Rule	In/out	IPsrc	IPdst	Proto	dstport	Action
A	Inward	External	Internal	TCP	25	Allow
B	Outward	Internal	External	TCP	>1023	Allow
C	Outward	Internal	External	TCP	25	Allow
D	Inward	External	Internal	TCP	>1023	Allow
E	*	*	*	*	*	Block

- Scenario 3:

Packet	In/out	IPsrc	IPdst	Proto	dstport	Action
5	Inward	12.1.2.3	173.18.3.4	TCP	8080	Allow (D)
6	Outward	173.18.3.4	12.1.2.3	TCP	5150	Allow (B)

- Conclusion: Oh, dear! That doesn't look good at all!
- Rules allow all connections where both ends use ports >1023.

Filter rules, 7

- Filtering in- and outgoing SMTP traffic. Include **srcport** in rules:

Rule	In/out	Ipsrc	IPdst	Proto	srcport	dstport	Action
A	Inward	External	Internal	TCP	>1023	25	Allow
B	Outward	Internal	External	TCP	25	>1023	Allow
C	Outward	Internal	External	TCP	>1023	25	Allow
D	Inward	External	Internal	TCP	25	>1023	Allow
E	*	*	*	*	*		Block

- Scenario 3:

Packet	In/out	Ipsrc	IPdst	Proto	srcport	dstport	Action
5	Inw.	12.1.2.3	173.18.3.4	TCP	5150	8080	Deny (E)
6	Outw.	173.18.3.4	12.1.2.3	TCP	8080	5150	Deny (E)

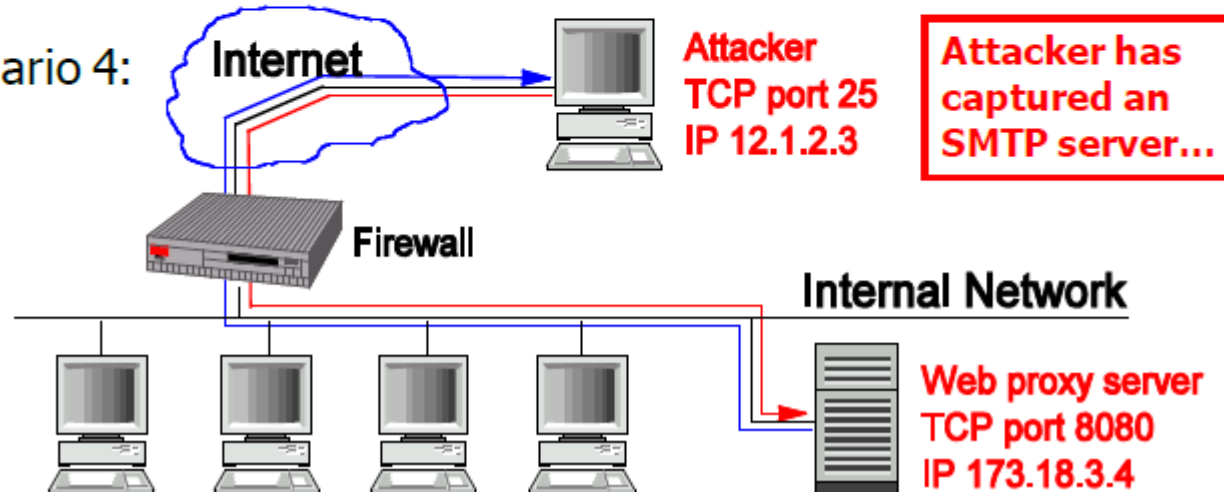
- Conclusion: This looks OK again!
- Check for yourselves that packets 1, 2, 3, 4 are treated OK.

Filter rules, 8

- Filtering in- and outgoing SMTP traffic. Include **srcport** in rules:

Rule	In/out	IPsrc	IPdst	Proto	srcport	dstport	Action
A	In	External	Internal	TCP	>1023	25	Allow
B	Out	Internal	External	TCP	25	>1023	Allow
C	Out	Internal	External	TCP	>1023	25	Allow
D	In	External	Internal	TCP	25	>1023	Allow
E	*	*	*	*	*		Block

- Scenario 4:



Filter rules, 9

- Filtering in- and outgoing SMTP traffic. Include **srcport** in rules:

Rule	In/out	IPsrc	IPdst	Proto	srcport	dstport	Action
A	Inward	External	Internal	TCP	>1023	25	Allow
B	Outward	Internal	External	TCP	25	>1023	Allow
C	Outward	Internal	External	TCP	>1023	25	Allow
D	Inward	External	Internal	TCP	25	>1023	Allow
E	*	*	*	*	*		Block

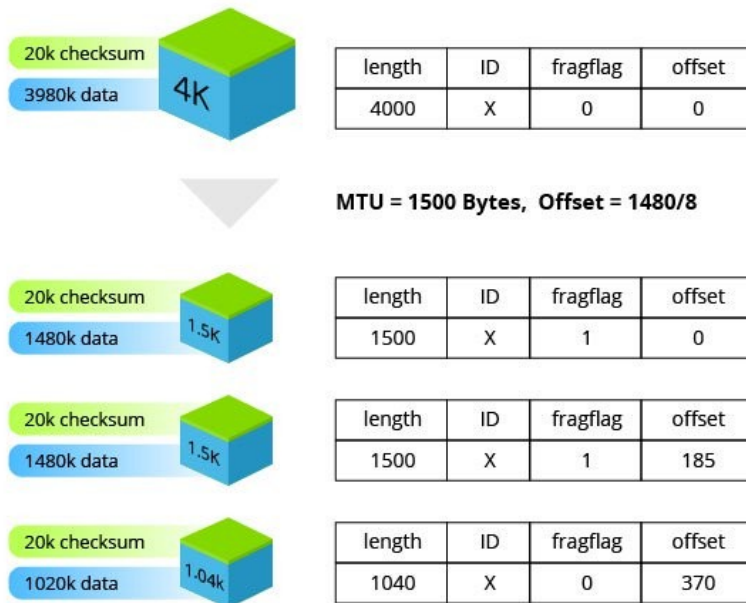
- Scenario 4:

Packet	In/out	IPsrc	IPdst	Proto	srcport	dstport	Action
7	Inw.	12.1.2.3	173.18.3.4	TCP	25	8080	Allow (D)
8	Outw.	173.18.3.4	12.1.2.3	TCP	8080	25	Allow (C)

- Conclusion: This looks bad again!
- Need yet more information (e.g. Flags) to get desired effect: Rules B and D must require ACK flag to be set in order to accept packet.

IP fragmentation

IP Fragmentation and Reassembly (Example)



Length - The size of the fragmented datagram

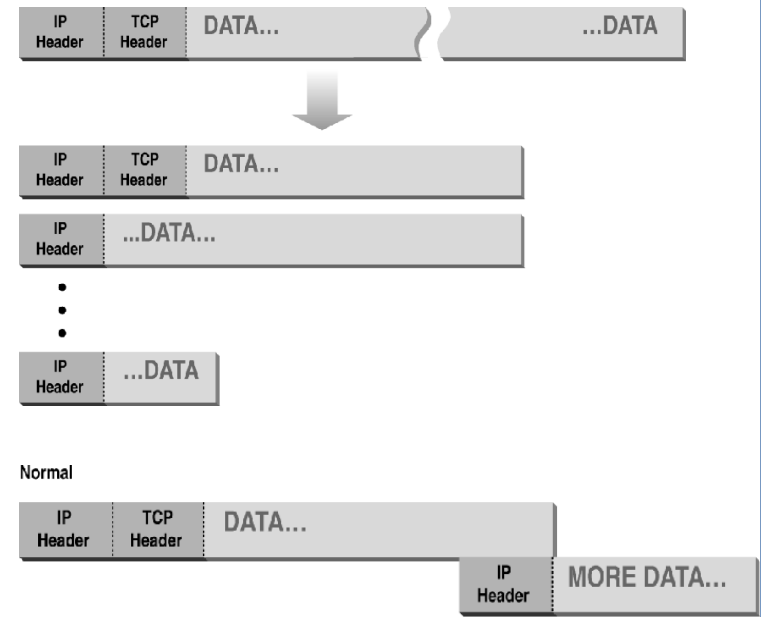
ID - The ID of the datagram being fragmented

Fragflag - Indicates whether there are more incoming fragments

Offset - Details the order the fragments should be placed in during reassembly

<https://www.incapsula.com/ddos/attack-glossary/ip-fragmentation-attack-teardrop.html>

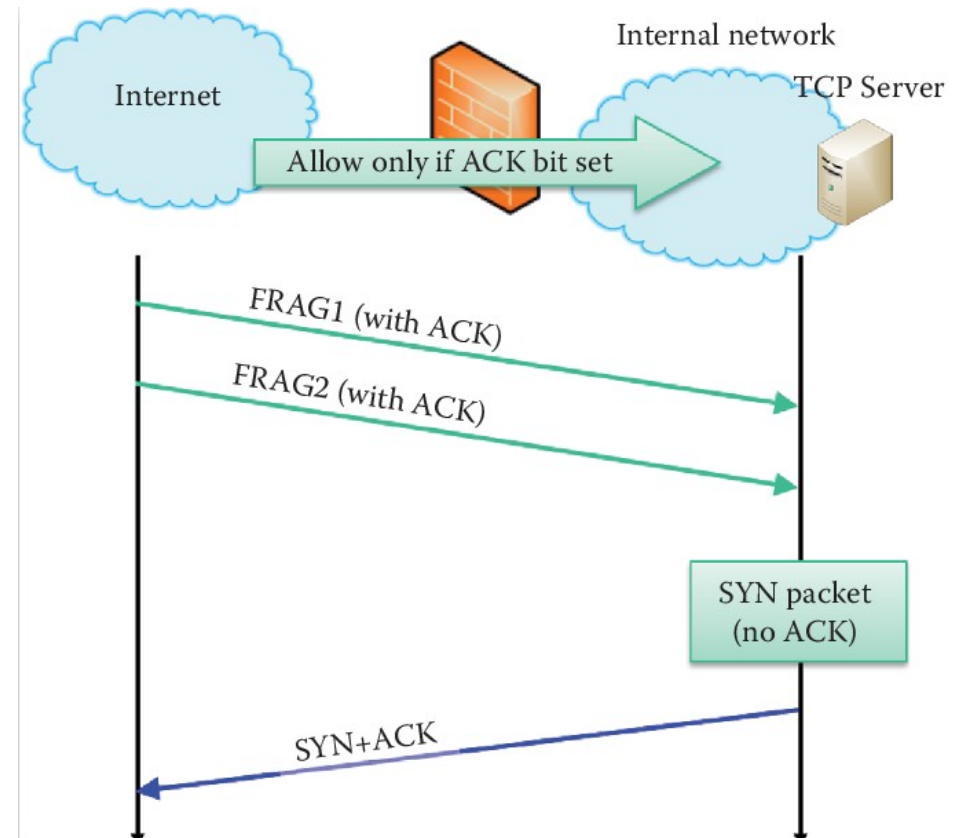
Normal fragmentation



Abnormal fragmentation

Incoming TCP connections with IP frag

- Firewall blocks any incoming TCP connection
- ACK packet is allowed for outgoing packets
- Internal host reassembles a packet with the SYN bit set because two fragment offsets are chosen in order to set the SYN bit
- Attacks
 - SYN scan
 - Create TCP connection
 - SYN flood - DoS





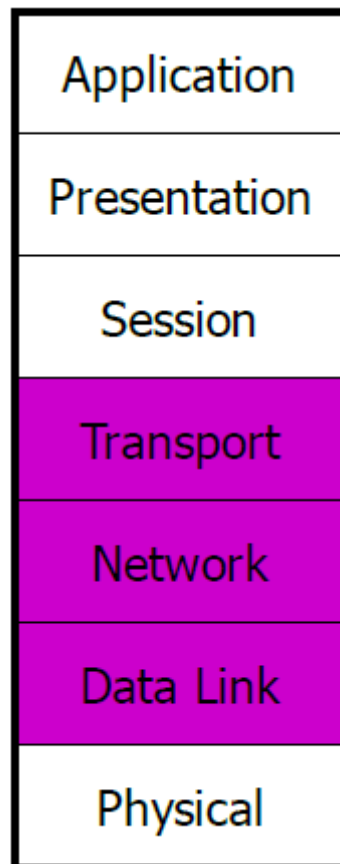
Stateful firewalls

Stateful packet inspection

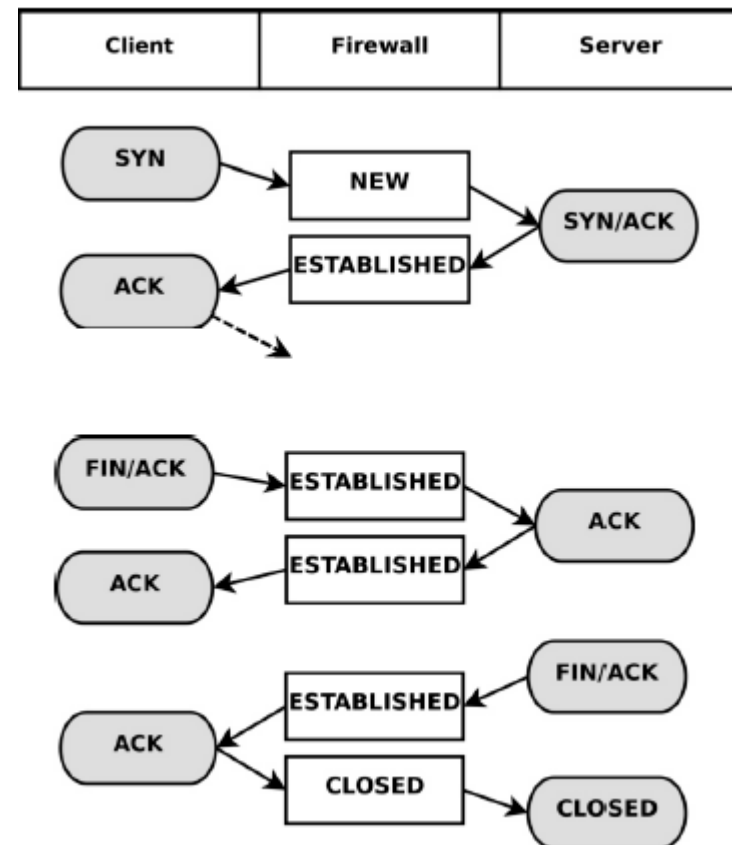
- Stateful Inspection Firewalls (or Dynamic Packet Filters) can keep track of established connections
- Can drop packets based on their source or destination IP addresses, port numbers and possibly TCP flags
 - Solve one major problem of simple packet filters, since they can check that incoming traffic for a high-numbered port is a genuine response to a previous outgoing request to set up a connection

Stateful firewall

- Considered layers



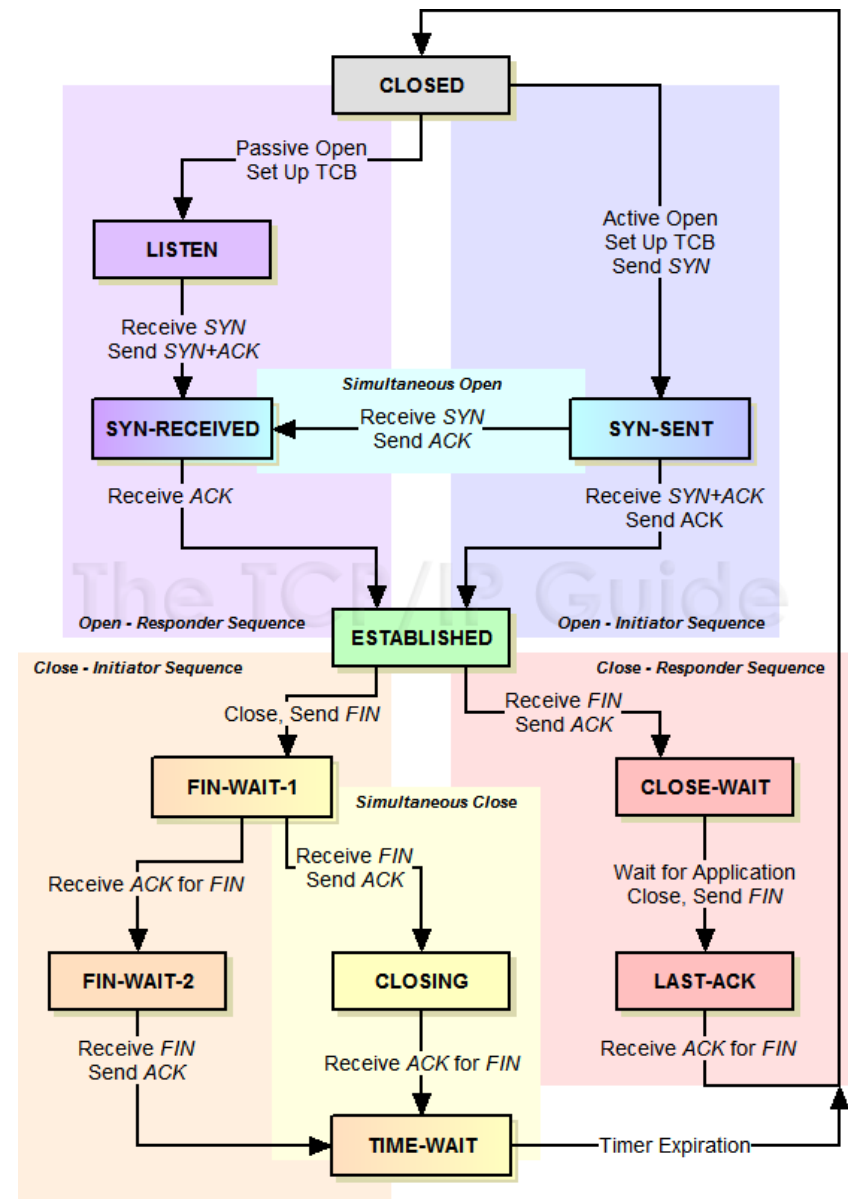
- Connection tracking



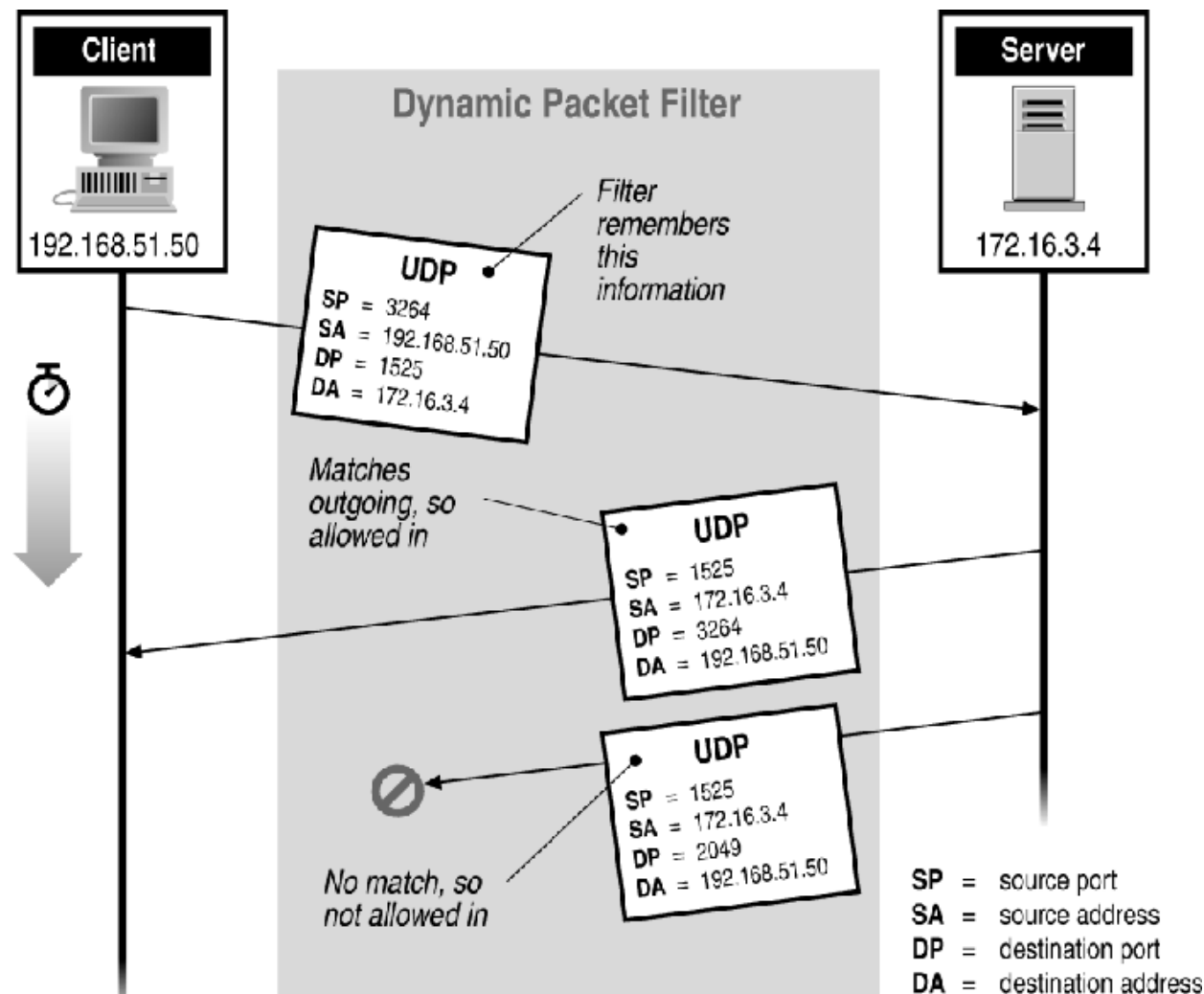
Connection tracking

Considered TCP States

- Setting up connection:
 - client calls from (high-numbered) port to port for application on server
 - server replies to (high-numbered) port on client
 - connection is considered established when the server gives correct SYN/ACK response.
- Closing connection:
 - both parties have to close the connection by sending a TCP packet with FIN flag set before connection is considered closed



Stateful firewall example

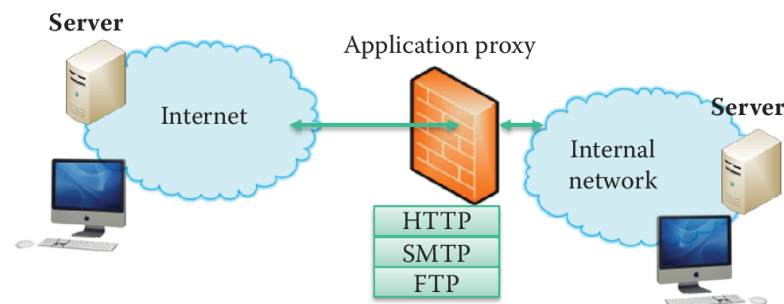




Other types of firewalls

Application-Level filtering (proxy)

- Deal with the details of services
 - Example: mail filters, FTP, HTTP proxy
 - Big overhead, but can log and audit all activity
- Can support user-to-gateway authentication
 - Log into the proxy server with username and password
 - Example: Microsoft ISA, SQUID



Proxy pro and cons

- Logging capacity
- Caching
- Intelligent filtering
- User-level authentication
- Protect for wrong implementations
- Introduce lag
- Application-specific
- Non transparent

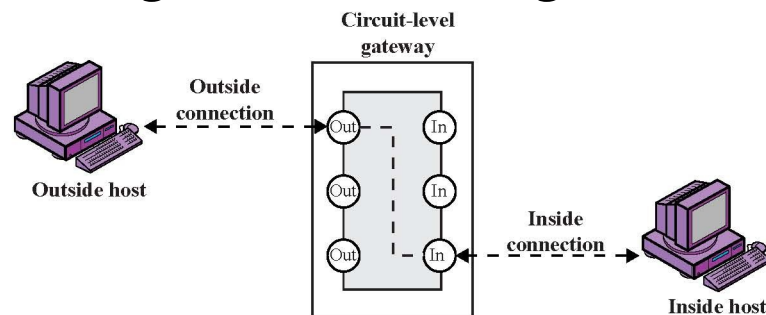


Circuit-level gateways (or generic proxy)

- Also known as a TCP relay
 - Able to deal with several protocols
- SOCKS (v5.0: Internet RFC1928) is the de facto standard
 - It also works with UDP
 - WinSock for Windows
- SOCKS performs at Layer 5 of the OSI model
 - The Session Layer above transport layer
 - TOR (the onion routing): socks-like interface
- The client connect to a proxy that relays its connections in a protocol-independent manner
- Provide user-authentication
- Usually no content filtering

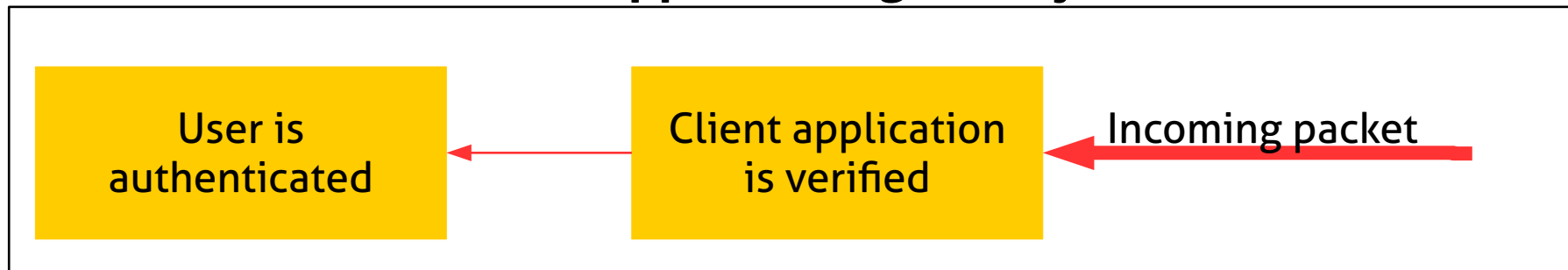
TCP relay

- Splices and relays TCP connections
 - Does not examine the contents of TCP segments
 - Can use ACL (like packet filtering, i.e. dst IP/dst port)
 - Less control than application-level gateway
- Client applications must be adapted for SOCKS
 - “Universal” interface to circuit-level gateways
- Example: `ssh -D 12345 <remote_host>`
 - More on this when talking about tunneling

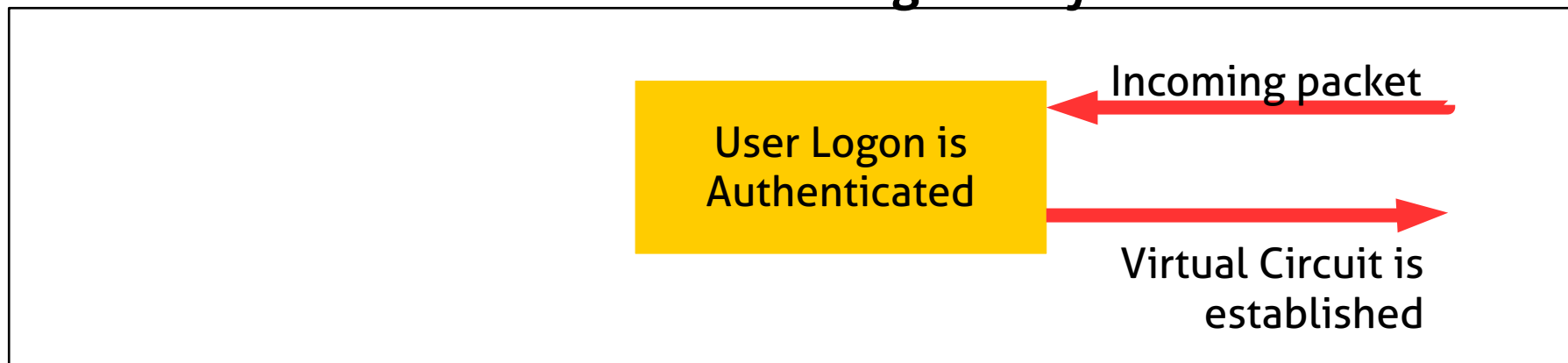


Application vs circuit level gateway

Application gateway



Circuit level gateway



Common firewall weaknesses

- No content inspection causes the problems
 - Software weakness (e.g. buffer overflow, and SQL injection exploits)
 - Protocol weakness (WEP in 802.11)
- No defense against
 - Denial of service
 - Insider attacks
- Firewall failure has to be prevented
 - Firewall cluster for redundancy

NG-Firewalls

- Next Generation firewalls try to include additional features
- Not only traffic filtering, but also:
 - Intrusion Detection System
 - VPN gateway
 - Deep Packet Inspection
 - Traffic shaping



Summary!

Summary

- Traffic regulation: routers and firewall
 - Decide the packets that can pass through the node
- Firewall architectures: where they go in the network?
 - Network segmentation and DMZ
- Types of firewalls:
 - Host firewall, stateless, stateful, application-gateway, circuit-gateway
- Stateless firewall weaknesses
 - No state, IP fragmentation

That's all for today

- **Questions?**
- See you next lecture!
- **Next Thursday we start at 12.40**
- **Resources:**
 - “Building internet firewalls”, Elizabeth D. Zwicky, Simon Cooper, D. Brent Chapman, O'Reilly 2nd ed.
 - https://docstore.mik.ua/oreilly/networking_2ndEd/fire/index.htm (I don't know if it is legal... but it is there...)
 - “Firewalls and Internet security: repelling the wily hacker”, William R. Cheswick, Steven M. Bellovin, Aviel D. Rubin, Addison-Wesley 2nd ed.



Practical Network Defense

Master's degree in Cybersecurity 2018-19

Network traffic regulation with firewalls 2

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Lab activity iptables

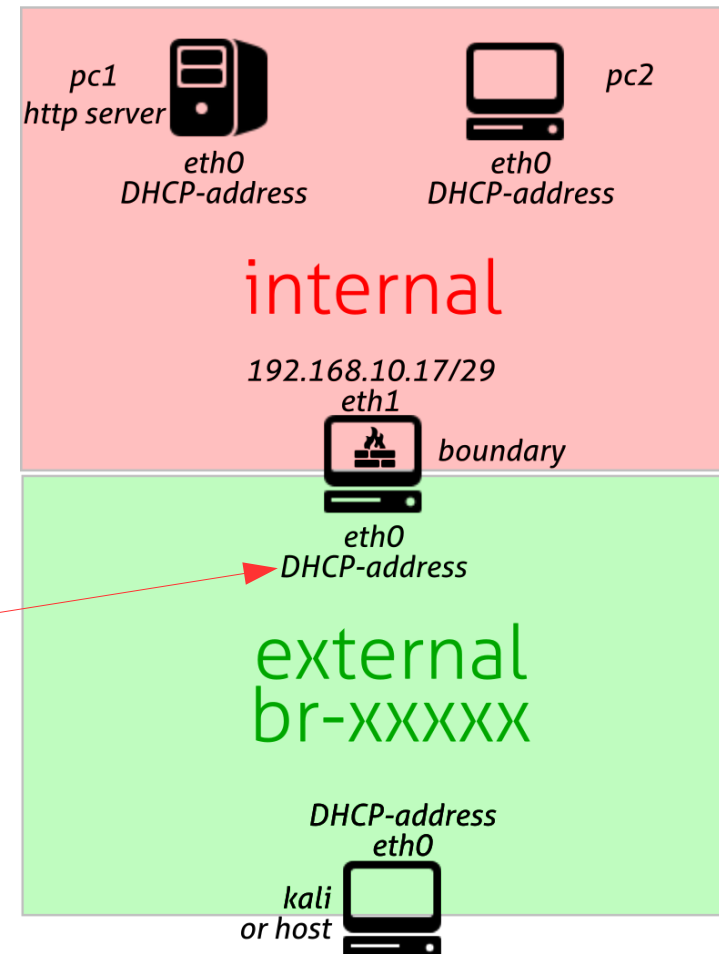
Network setup

- Use pnd-lab2-es1
- In pc1 create an index.html file
- Start a webserver and sshd
- Add a route towards internal via boundary-eth0

```
host$ sudo ip r add 192.168.10.16/29 via  
172.16.0.2 ← change this according to
```

```
pc1# /etc/init.d/ssh start
```

```
pc1# python -m SimpleHTTPServer 80
```





First Demo

Objective: **block any ping to our pc1**

- Start capturing with wireshark
- Firstly, verify we can ping from pc2 and from host
- Then raise our firewall, using iptables

```
iptables -A INPUT -p icmp --icmp-type echo-request -j DROP
```

- Verify we cannot ping pc1 anymore, but we can ping the others
- Check with tcpdump what's going on...
- When done, clean iptables rules

```
iptables -F
```

Second demo

Objective: **exclude any service but HTTP**

- Start capturing with wireshark
- Firstly, verify we can connect from host and pc2 to pc1 ssh and web server (through the different ports)
- Then raise our firewall on pc1, using iptables

```
iptables -A INPUT -p tcp --destination-port 80 -j ACCEPT
iptables -A INPUT -j REJECT
```
- Verify we cannot reach pc1 any more (where?)
- Check with wireshark what's going on...
- When done, clean iptables rules

```
iptables -F
```


Iptables

- It is the implementation of a packet filtering firewall for Linux that runs in kernel space
 - It is the evolution of ipchains and ipfw. Coming successor will be nftables
- iptables tool inserts and deletes rules from the kernel's packet filtering table
- It can also operate at the Transport layer (TCP/UDP)
- Old but still extremely valuable tutorial:

www.frozentux.net/iptables-tutorial/iptables-tutorial.html

Iptables fundamentals

- The rules are grouped in **tables**
 - For now, we focus on the **FILTER** table
- Each table has different **CHAINS** of rules
- Each packet is subject to each rule of a table
- Packet fates depend on the **first matching rule**
- To see chains and rules of the filter table

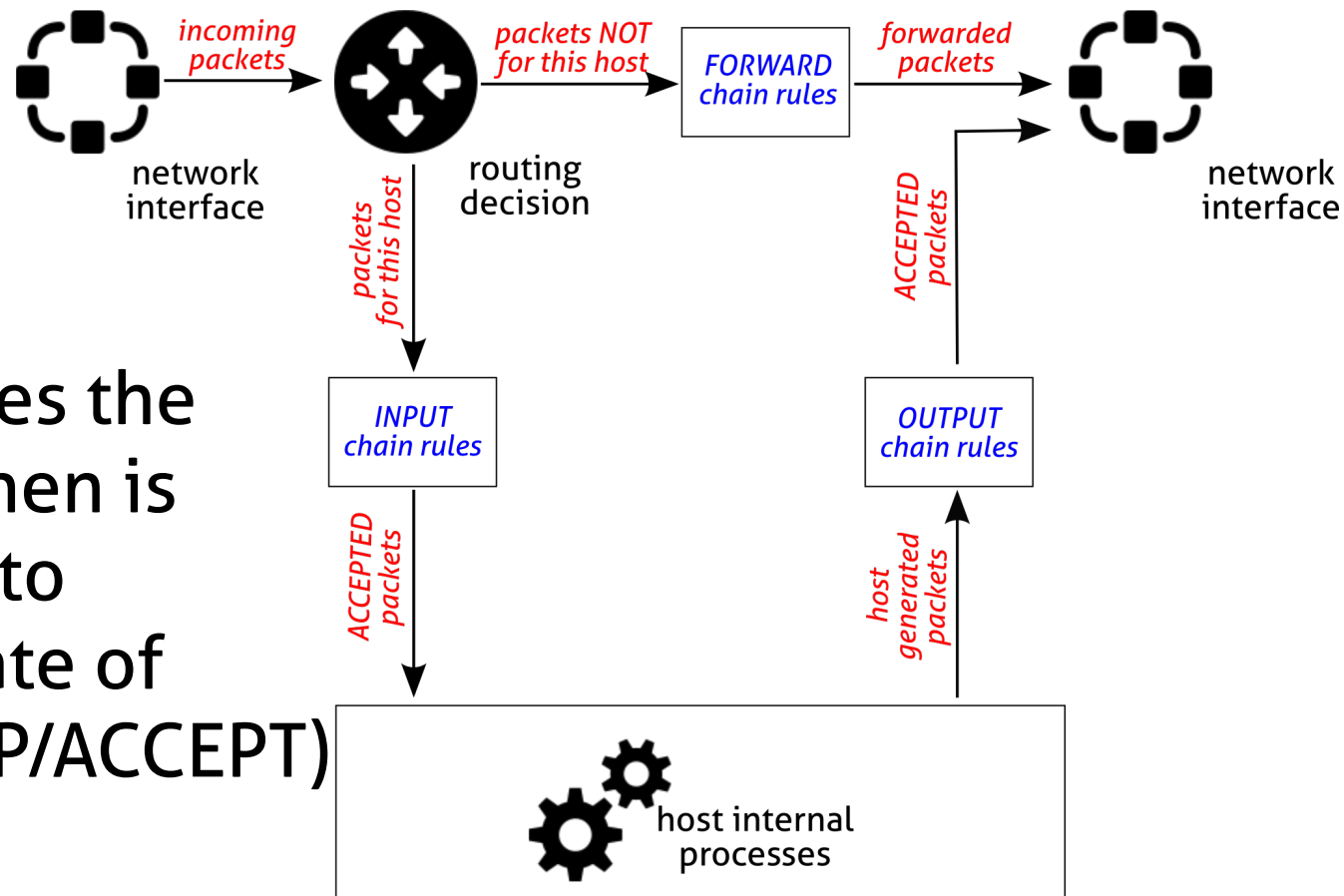
```
iptables -L
```

or (better)

```
iptables -L -n -v --line-numbers
```

Filter table

- Three built-in rule chains:
 - INPUT
 - OUTPUT
 - FORWARD
- If a packet reaches the end of a chain, then is the chain policy to determine the fate of the packet (DROP/ACCEPT)



Create and save a rule set

- You can save in a shell script the sequence of the iptables commands

- Typical structure of iptables_rules.sh

```
#!/bin/bash
```

```
# flush (clean) the filter table  
iptables -t filter -F
```

```
# allow only service XX  
iptables ...
```

- Or you can use the built in commands
 - iptables-save > iptables_rules.bk
 - iptables-restore < iptables_rules.bk

Useful iptables command switches

iptables switches	Description
-t table	Specifies the table (filter if not given)
-j target	Jump to the target (it can be another chain)
-A chain	Append a rule to the specified chain
-F	Flush a chain
-P policy	Change the default policy
-p protocol	Match the protocol type
-s ip-address	Match the source IP address
-d ip-address	Match the destination IP address
-p tcp --sport port	Match the tcp source port (also works for udp)
-p tcp --dport port	Match the tcp destination port (also works for udp)
-i interface-name	Match input interface (from which the packet enters)
-o interface-name	Match output interface (on which the packet exits)

Review the rulesets of demos

```
iptables -A input -p icmp -icmp-type echo-request -j DROP
```

```
iptables -A input -p tcp --destination-port 80 -j ACCEPT
```

```
iptables -A input -j REJECT
```

- We can specify different “targets” (this is a subset):
 - **ACCEPT**: the packet is handed over to the end application or the operating system for processing
 - **DROP**: the packet is blocked.
 - **REJECT**: the packet is blocked, but it also sends an error message to the source host of the blocked packet
 - reject-with <qualifier> <qualifier> is an ICMP message*
 - **LOG**: the packet is sent to the syslog daemon for logging.
 - `iptables` continues processing with the next rule in the table.
 - You can't log and drop at the same time → use two rules (*--log-prefix "reason"*)

Other useful iptables command switches

iptables switches	Description
-p tcp --sport port	Match the tcp source port
-p tcp --dport port	Match the tcp destination port
-p udp --sport port	Match the udp source port
-p udp --dport port	Match the udp destination port
--icmp-type type	Match specific icmp packet types
-m <i>module</i>	Uses an extension module
-m state --state s	Enable connection tracking. Match a packet which is in a specific state: NEW: the packet is the start of a new connection ESTABLISHED: the packet is part of an established connection RELATED: the packet is the starting of a related connection (like FTP data) INVALID: the packet could not be identified
-m multiport ...	Enable specification of several ports with one single rule

Modules examples

- Allow both port 80 and 443 for the webserver on inside:

```
iptables -A FORWARD -s 0/0 -i eth0 -d 192.168.1.58 -o eth1 -p TCP \
--sport 1024:65535 -m multiport --dport 80,443 -j ACCEPT
```

- The return traffic from webserver is allowed, but only if sessions are opened:

```
iptables -A FORWARD -d 0/0 -o eth0 -s 192.168.1.58 -i eth1 -p TCP \
-m state --state ESTABLISHED -j ACCEPT
```

- If sessions are used, you can reduce an attack called half open

Half open is known to consume server all free sockets (tcp stack memory) and is sensed as a denial of service attack, but it is not.

Sessions are usually waiting 3 minutes.

- The diagram is divided into two main sections: TCP and UDP, each showing a sequence of packet exchanges between a Client, a Firewall, and a Server.

TCP Section:

 - Initial State:** A table at the top shows the initial state of the connections: Client (empty), Firewall (empty), and Server (empty).
 - SYN Exchange:**
 - Client sends a **SYN** packet to the Firewall.
 - Firewall forwards the **SYN** packet to the Server.
 - Server responds with a **SYN/ACK** packet to the Firewall.
 - Firewall forwards the **SYN/ACK** packet to the Client.
 - Established Connection:**
 - Client sends an **ACK** packet to the Firewall.
 - Firewall forwards the **ACK** packet to the Server.
 - Server sends a **FIN/ACK** packet to the Firewall.
 - Firewall forwards the **FIN/ACK** packet to the Client.
 - Connection Closure:**
 - Client sends an **ACK** packet to the Firewall.
 - Firewall forwards the **ACK** packet to the Server.
 - Server sends a **CLOSED** packet to the Firewall.
 - Firewall forwards the **CLOSED** packet to the Client.

UDP Section:

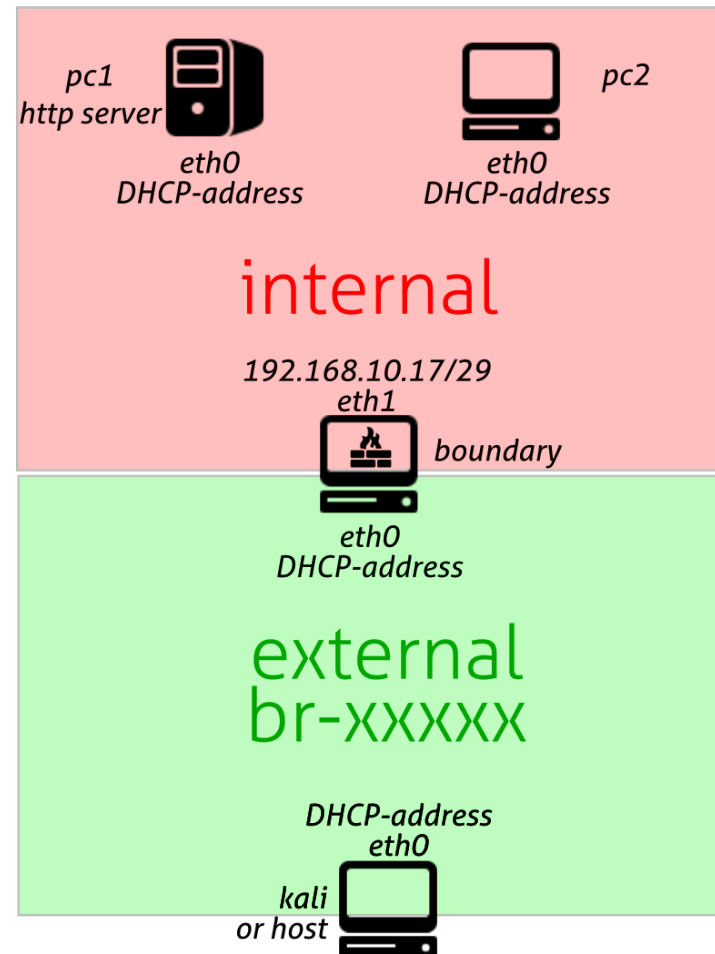
 - Initial State:** A table at the top shows the initial state of the connections: Client (empty), Firewall (empty), and Server (empty).
 - UDP Packet Exchange:**
 - Client sends a **UDP Packet** to the Firewall.
 - Firewall forwards the **UDP Packet** to the Server.
 - Server responds with a **UDP Packet** to the Firewall.
 - Firewall forwards the **UDP Packet** to the Client.
 - Established Connection:**
 - Client sends a **....** packet to the Firewall.
 - Firewall forwards the **....** packet to the Server.
 - Server sends a **....** packet to the Firewall.
 - Firewall forwards the **....** packet to the Client.



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Activity 1

- Start with the previous setting
- Turn on pc1 web server
- Protect the website only from the external network
 - Configure boundary
- Try with other services or ports
 - Ex: telnet, ssh, http on different ports



Activity 2

- Reconsider the ARP spoofing scenario
- Configure the iptables on the pc machines in order to avoid the attack
 - Several options:
 - Use static arp associations
 - Use arptables

That's all for today

- Questions?
- See you tomorrow!
- Resources:
 - “Building internet firewalls”, Elizabeth D. Zwicky, Simon Cooper, D. Brent Chapman, O'Reilly 2nd ed.
 - https://docstore.mik.ua/orelly/networking_2ndEd/fire/index.htm
(I don't know if it is legal... but it is there...)
 - “Firewalls and Internet security: repelling the wily hacker”, William R. Cheswick, Steven M. Bellovin, Aviel D. Rubin, Addison-Wesley 2nd ed.
 - www.frozentux.net/iptables-tutorial/iptables-tutorial.html