# Almost-Matching-Exactly for Treatment Effect Estimation under Network Interference

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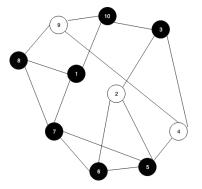
# Setting - Causal Inference

- ⊚ We have i = 1, ..., n experimental units
- ⊚ Treatment  $t_i \in \{0,1\}$  with  $\mathbf{t} \in \{0,1\}^n$  is a binary vector with the treatment level of every unit.
- © Potential outcomes  $Y_i(t_i, \mathbf{t})$  are random variables and depend on both treatment of unit i (1st argument), and treatment of **all other units** (2nd argument).
- ⊚ Observed treatment  $\mathbf{T} \in \{0,1\}^n$  is assigned **uniformly at** random.
- ⊚ Observed outcome:  $Y_i = T_i Y(1, \mathbf{T}) + (1 T_i) Y_i(0, \mathbf{T})$ 
  - Since treatment is randomized:  $\mathbb{E}[Y_i|\mathbf{T}=\mathbf{t},T_i=t]=\mathbb{E}[Y_i(t,\mathbf{T})]$  (Ignorability).
- ⊚ Units are connected in a network, G, in which unit i's treated neighborhood subgraph is  $G_{N_i}^{\mathbf{t}}$ .

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#### The Problem: No SUTVA!

Usually we assume SUTVA: that *units' treatments don't influence other units' outcomes*, but we can't do that here because our units are connected in a network:



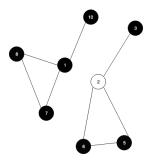
It could be that the treatment assigned to j influences the outcome of i through their connection in the network.

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# Similar Graphs Carry Similar Interference

#### Idea

What if the amount of interference experienced by a unit depended on the *shape* of its treated neighborhood subgraph?



Then, in expectation, two units with the same treated neighborhood graph will respond similarly to the treatment. We can use this idea to do matching to reduce interference.

# Assumptions

- 1. Outcome model:  $Y_i = \alpha + t_i \beta_i + f_i(G_{N_i}^{\mathbf{t}}) + \epsilon_i$ 
  - f is some interference function dependent on  $G_{N_i}^t$ , the **treated neighborhood subgraph** of unit i.
- 2.  $\mathbb{E}[\epsilon_i|T_i] = \mathbb{E}[\epsilon_i] = 0$ 
  - Ignorability
- 3. If  $G_{\mathcal{N}_i}^{\mathbf{t}} \simeq G_{\mathcal{N}_j}^{\mathbf{t}}$  then  $f_i(G_{\mathcal{N}_i}^{\mathbf{t}}) = f_j(G_{\mathcal{N}_j}^{\mathbf{t}}) = f(G_{\mathcal{N}_i}^{\mathbf{t}})$ .
  - Two units with isomorphic neighborhood sugraphs experience the same interference.
  - Together with (1), this assumption encodes a version of SANIA (Airoldi and Sussman, 2018) conditional on unit's neighborhood subgraphs.

# Matching

- © Under these assumptions we can match units to recover the effect we are interested in  $(\tau_i)$
- ⊚ For a treated unit i, find a control unit j, such that  $G_{N_i}^{\mathbf{t}} \simeq G_{N_i}^{\mathbf{t}}$
- ⊚ Subtract  $Y_i$  from  $Y_i$  to recover  $\tau_i$  in expectation

$$\mathbb{E}[Y_i - Y_j] = \mathbb{E}[\tau_i + f(G_{\mathcal{N}_i}^{\mathbf{t}}) - f(G_{\mathcal{N}_i}^{\mathbf{t}}) + \epsilon_i + \epsilon_j] = \tau_i$$

If no unit has exactly  $G_{\mathcal{N}_j}^{\mathbf{t}} \simeq G_{\mathcal{N}_i}^{\mathbf{t}}$  we can find a unit that has a neighborhood graph that is not quite equal to  $G_{\mathcal{N}_i}^{\mathbf{t}}$  but is similar enough.

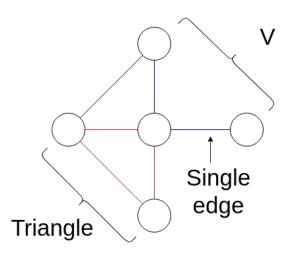
#### Problem

How do we represent similarity between neighborhood graphs?

# Subgraph Counts

- We'll say that neighborhood graphs are similar if they contain similar counts of subgraphs
- In fact, if the counts of subgraphs are the exact same, then the neighborhood graphs must be isomorphic.
- We match together units that have the most similar subgraphs
- If no unit matches exactly to another on all subgraph counts, then we must choose which subgraphs we want to match exactly on and which ones we can afford to ignore.

# Subgraph Counts



This graph has 2 triangles, 6Vs and 4 single edges.

#### ·AME-Networks

#### Problem #2

How do we choose which subgraphs we should use to represent the treated neighborhood graphs of our units?

# We need a matching method that selects which subgraphs to match on

- 1. Enumerate (up to isomorphism) all p subgraphs  $S_1, \ldots, S_p$  seen across all the  $\mathcal{N}_i$ ,  $i = 1, \ldots, n$
- 2. For each unit i, define the p-dimensional vector  $S(G_{N_i^t})$  as the vector whose j'th entry is the number of  $S_j$  in  $N_i$
- 3. These are likely a lot (maximum on the order of  $|\mathcal{N}_i|^2$ ) and it's unlikely that many units will have identical counts
- 4. The FLAME matching algorithm will automatically select the best subgraphs to match on

#### FLAME: An Overview

- FLAME (Fast Large-Scale Almost Matching Exactly) is a method for creating interpretable matches between units with discrete covariates that performs variable selection while matching.
  - 1. Match units exactly on as many covariates as possible
  - Drop a covariate
  - 3. Repeat
- At each step, drop the covariate maximizing match quality:

$$MQ = C \cdot BF - PE$$

- ⊚ BF = prop. controls matched + prop. treated matched
- PE = prediction error achieved by remaining covariates
- Tradeoff between making matches and accurate prediction

### A Small Change to FLAME

We know from our theoretical setup that network statistics should do two things well:

- 1. Predict the **outcomes**
- 2. Predict the **network**

To measure how well the our network statistics (subgraphs) are predicting the network, we model the edges of the network as independent conditional on the observed statistics

$$E_{ij}|x_i, x_j \stackrel{iid}{\sim} \text{Bern}(\text{logit}(\beta_1'x_i + \beta_2'x_j))$$

and consider the AIC of the resulting model, AIC<sub>network</sub>.

#### The Modified PE Function

To ensure FLAME strikes a balance between predicting both the **outcomes** and the **network**, we modify the PE function:

$$PE = \sum_{t=0}^{1} \underset{f \in \mathcal{F}}{\operatorname{arg \, min}} \frac{1}{n} \sum_{i=1}^{n} (Y_i - f(S(G_{\mathcal{N}_i}^{\mathbf{t}})))^2$$
$$- \underbrace{D \cdot AIC_{\text{network}}}_{\text{new component}}$$

# Simulation Setup

- © Recall the outcome model:  $Y_i = \alpha + \beta_i t_i + f(G_{N_i^{\dagger}}) + \epsilon_i$
- The graph G is generated according to Erdos-Renyi or Stochastic Block models
- $\odot$  We consider various forms of f based off different features:
  - $\circ$   $d_i$ : the degree of unit i
  - $\Delta_i$ : the number of triangles in  $\mathcal{N}_i$
  - $t_i^k$ : the number of units in  $\mathcal{N}_i$  with degree  $\geq k$
  - $\bigstar_i^k$ : the number of *k*-stars in  $\mathcal{N}_i$
  - $B_i$ : the vertex betweenness of unit i
  - $C_i$ : the closeness centrality of unit i

#### **Estimators**

- True: nearest neighbor (NN) on true interference
- Naive: naive difference in means
- ⊚ Eigen All: NN on eigenvalues of adjacency matrix *A*
- ⊚ Eigen All: NN on largest eigenvalue of A
- Stratified Naive: Stratified degree estimator
- SANIA: MIVLUE under SANIA
- FLAME: Our approach

#### Simulation Results

Type 1:  $f_i = 10\Delta_i$ 

Type 2:  $f_i = 10d_i + 10\Delta_i$ 

Type 3:  $f_i = -5d_i + 1\Delta_i 1 + 10 \bigstar_i^2 + 1 \bigstar_i^4 + 1 \dagger_i^3 + 10B_i - 5C_i$ 

Type 4:  $f_i = t_i^3$ 

	Interference			
Estimator	Туре 1	Type 2	Туре 3	Type 4
True	0.93	1.56	1.75	0.72
FLAME	2.10	9.22	6.18	0.37
SANIA	4.07	11.49	9.29	0.43
Stratified	4.35	4.04	4.94	4.90
Eigen	2.10	8.68	8.05	0.49
Naive	4.07	11.49	9.29	0.43

Mean ATE Over 50 Simulations

# Social Network in India: Background

- Using data from Banerjee et al. (2013)
- Estimates how centrality of first informed individuals predicts eventual adoption of the microfinance program
- Network information for sample villages
- Collected 6 months before the actual intervention of information diffusion
- Leaders were the source through which information was disseminated (injection points)
- Village level outcome
- Structural model estimated using simulated moments

# Social Network in India: Survey Design

- O Village Survey
  - o No of HHs, facilities, module on village elders
- HH Survey (census)
  - o Dwelling information (wealth index), religion, caste
- Individual survey of a subsample of villagers
  - Respondents randomly selected stratified by religion and geographical sublocation.
  - Included social network data across 12 dimensions such as Visitors, socialize, medical advice, borrow/lend money/material goods, general advice, pray together.
  - Education, election card, SHGs, language, caste, employment etc.

# Social Network in India: Application

- Effect of education on participation in self-help groups (SHGs)
  - Hypothesis: Less educated -> more likely to participate in SHGs
  - Network endogeneity: Less educated rely more on social network and that in turn have an effect on SHG participation
- Effect of education on election participation
  - Election participation measured by ownership of election card.
  - Hypothesis: Less educated -> less likely to hold an election card
  - However, +ve effect is observed using linear regression.
    Biased might be caused by omission of social network
- Effect of gender on participation in SHGs

# Social Network in India: Application

- ⊚ 75 villages
  - Spread across 5 districts, mean distance between villages is 46 km.
  - o Villages can be considered as independent from each other
- Variation across villages and by wealth status of viilages

#### Plan: We Want to Hit the October 8 AISTATS Deadline

- 1. We need to implement the revised PE function
- 2. We need to redo the simulations
- 3. More theory? Statements on how well the subgraph count can "encode" a given graph would be nice
- 4. Find a good application
- 5. Write!