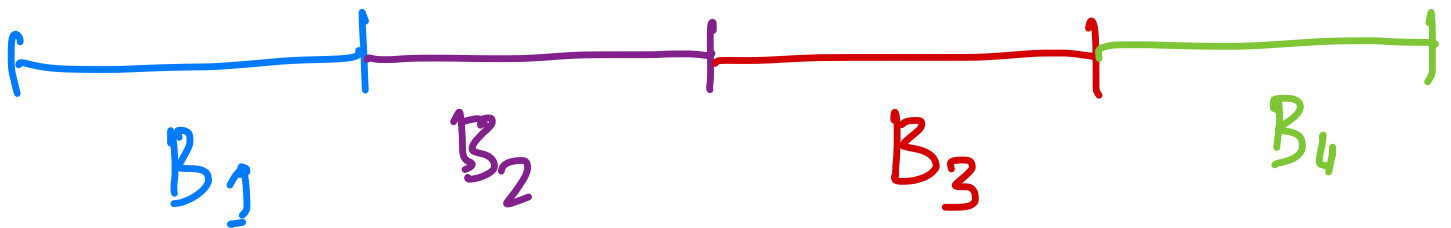
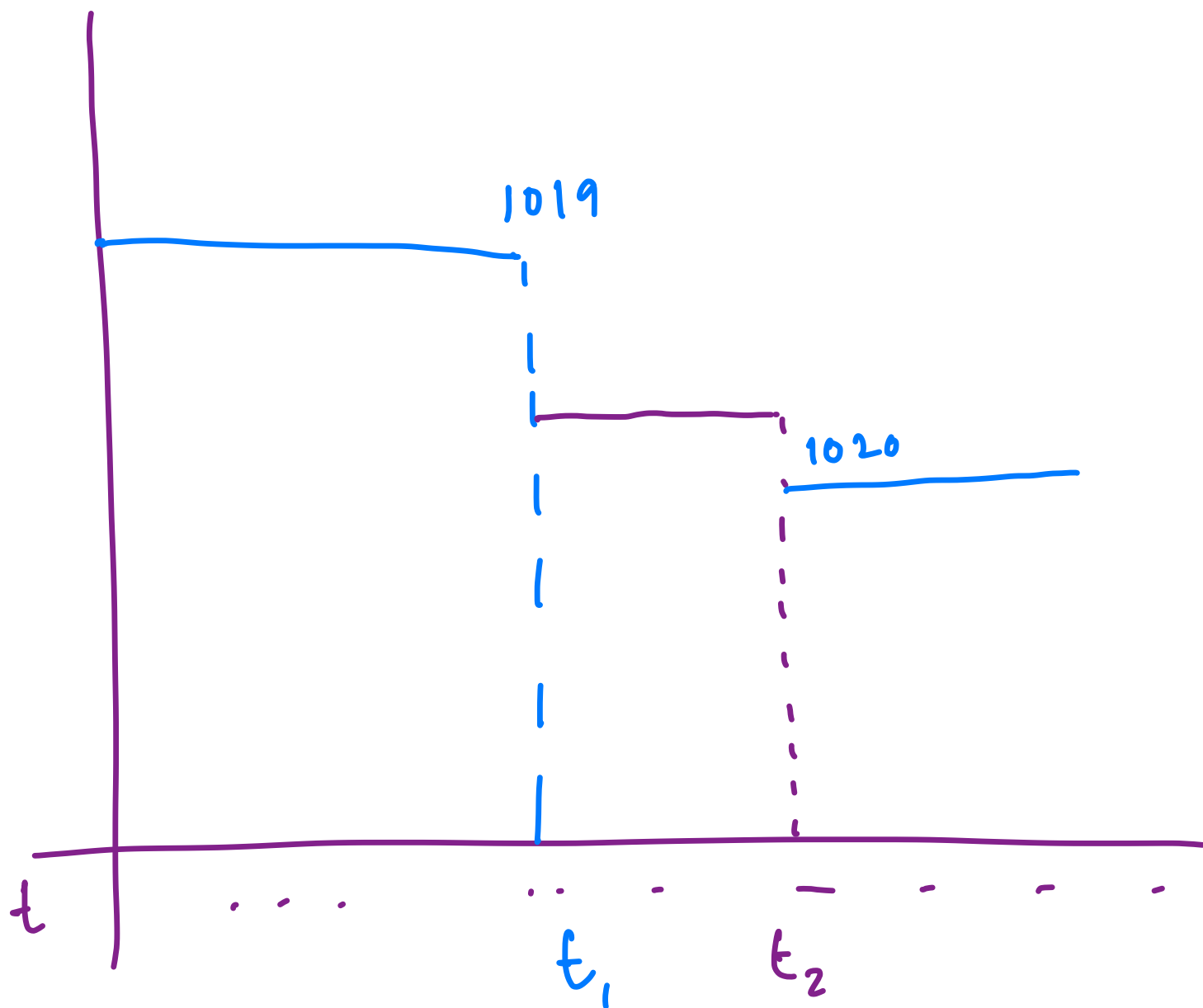


Batch processing



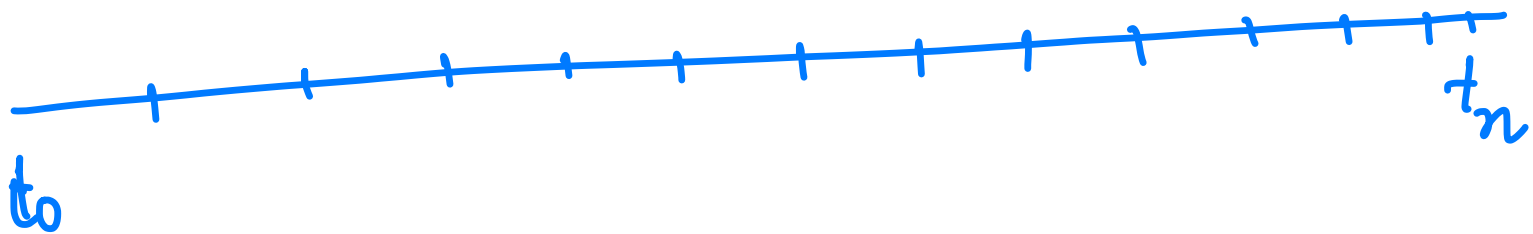
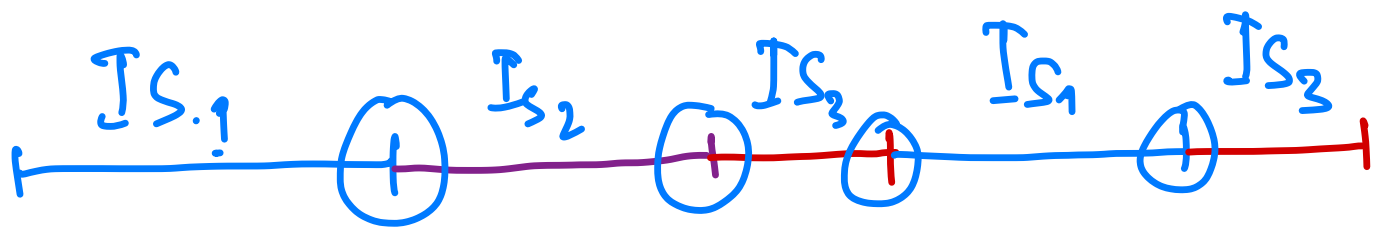
Sequential





O.S.

Coordinator



Instruction Set \rightarrow
processes

Change-over

Context-switching

OS \rightarrow to efficiently manage the CPU time between different processes

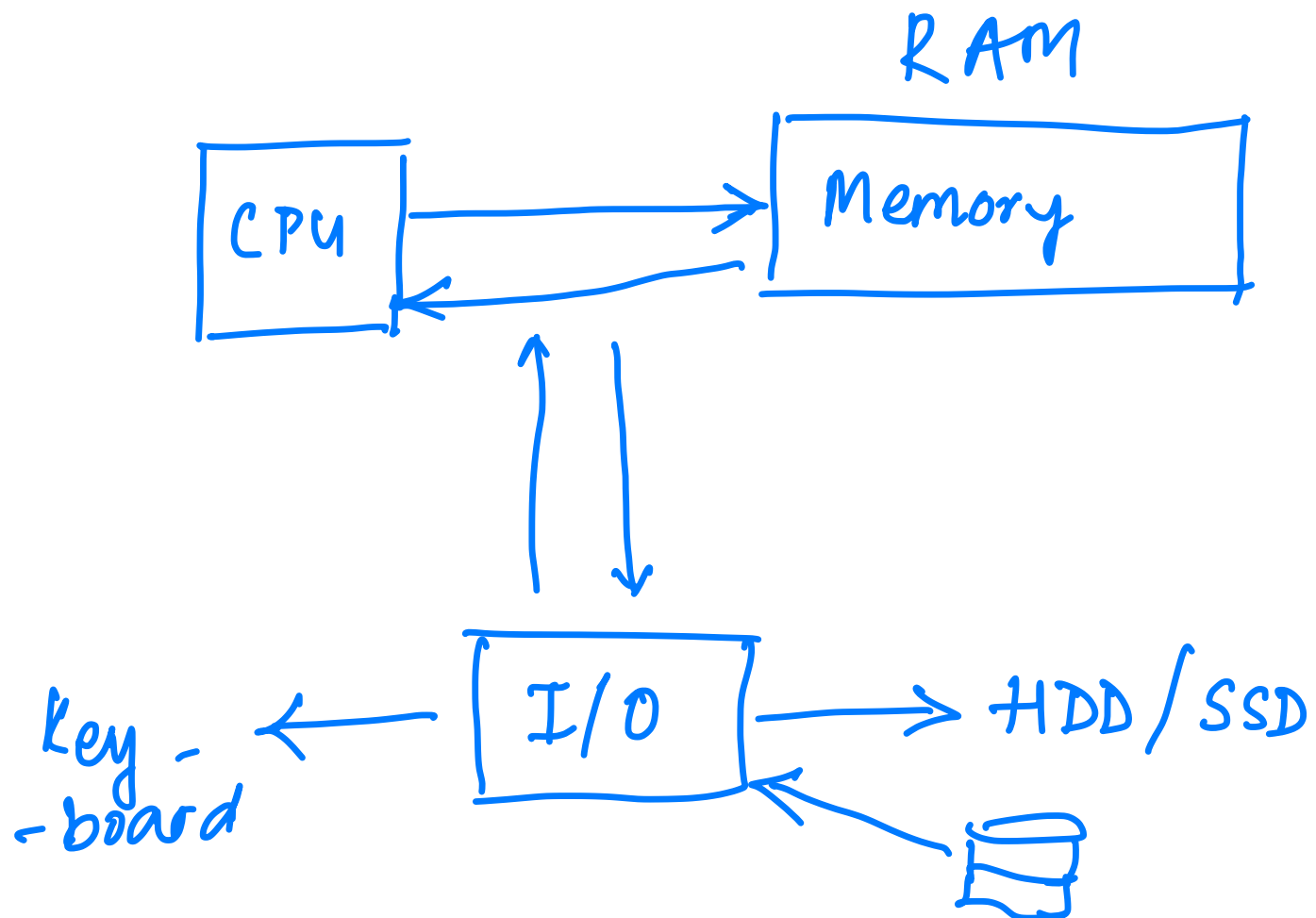
Efficient management of CPU time

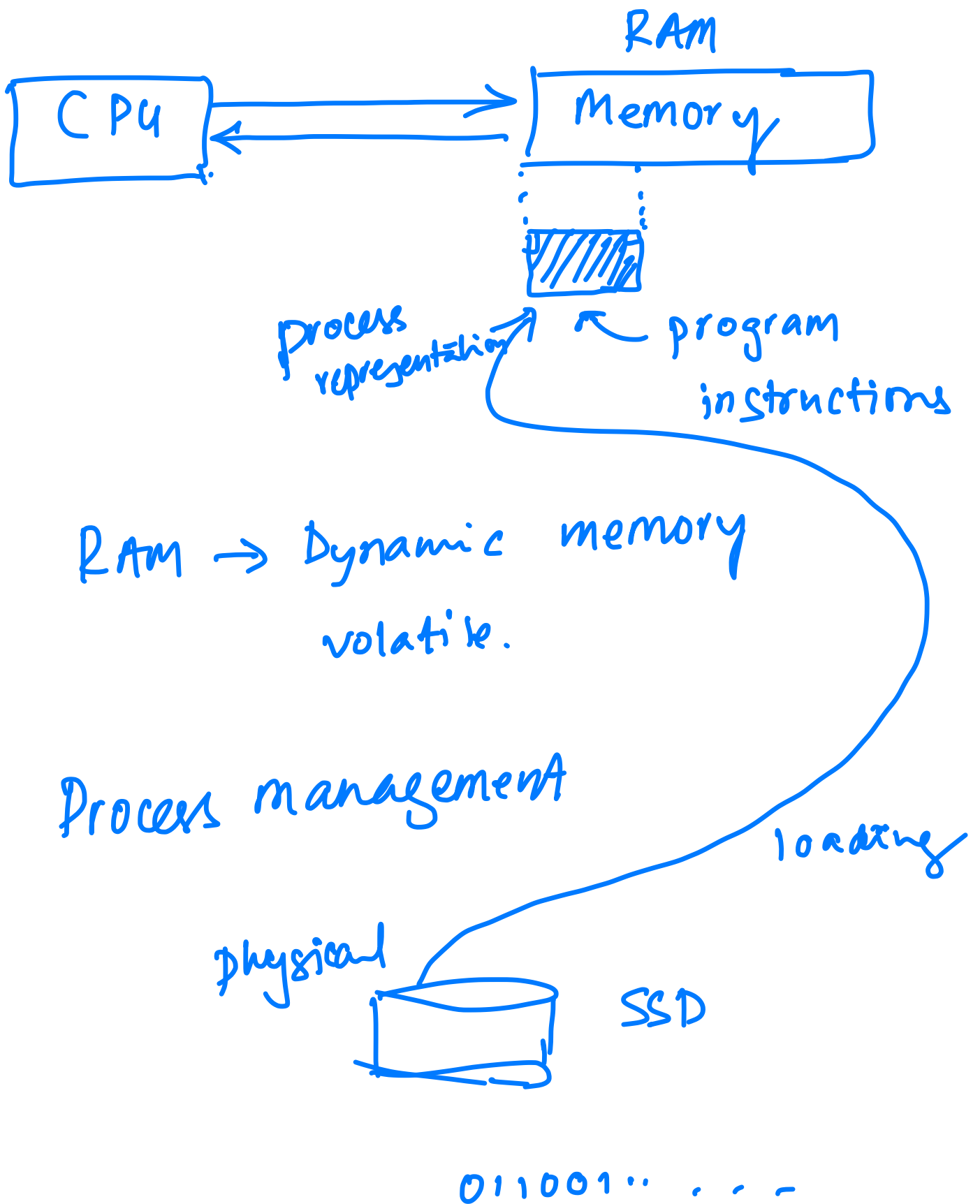
- 1) Equal importance to all processes
(-fairness)
- 2) provide the isolation for the processes and yet provide the way for communication.
- 3) Make sure a process doesn't hog the CPU beyond the limits

process → set of instructions for the CPU.

Usually store these instructions in binary in a file on the file system.

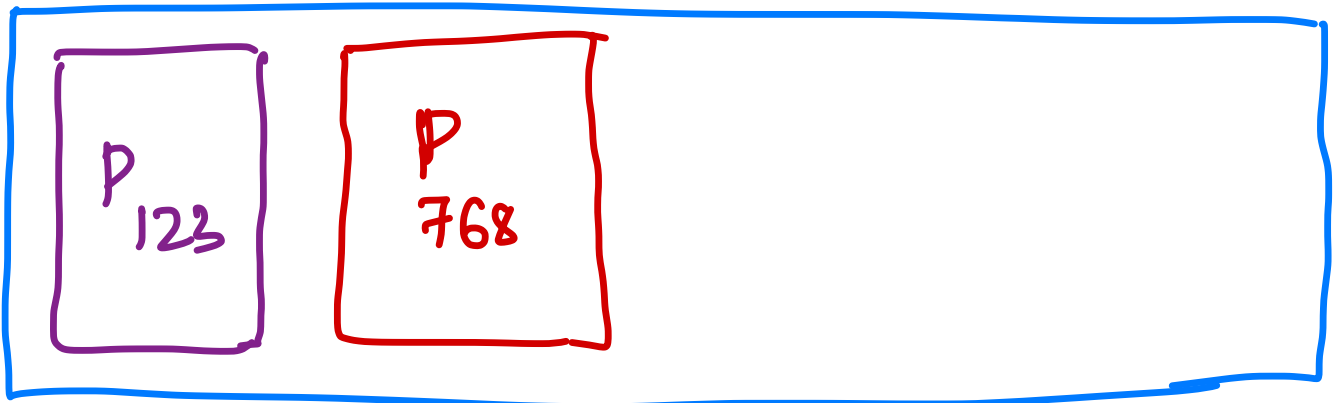
von Neumann Architecture.





Loading happens manually (by user)
or automatically.

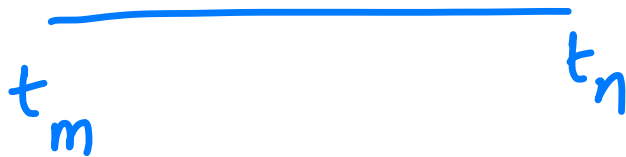
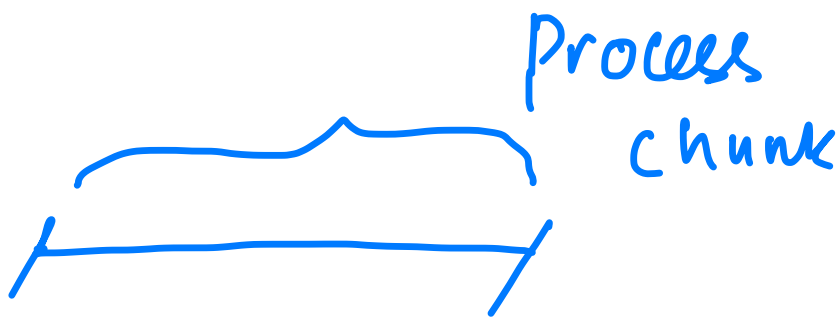
When a process is loaded, the OS provides an abstraction of the slice of memory to this process. This slice is an isolated box. (Sandboxing)



When the CPU is executing P_{123} it is only allowed to access the addresses within the silo of P_{123}

OS, Time slice, Process.

Thread. \rightarrow execution path.



Within a process, we may group the instructions so that they constitute an execution path. This happens usually in the code.

Technically, a process may have more than one such execution paths.

Execution path is called a Thread

There is always at least one thread per process.

A process is basically instructions executing. These instructions are themselves (or part of) functions, or methods. The execution, then, is the chain of method calls, where every method has other instructions.

A group of instructions, then, is the group of such function call chain. Thread is the method call chain.

In a method, we create variables (local variables), we may get parameters passed, we may call other methods.

The programming abstraction of this is called a Stack

In a Stack, we have the ability to declare local variables, parameters, and call chain handling.

A thread gets its own Stack

