

Q1]

a)
Ans →

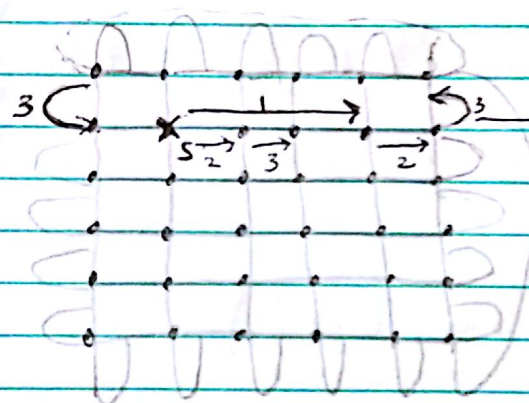
The formula for calculating the communication time for one to all broadcast in a ring, using cut through routing is →

$$T_c = 2 T_s \log P + T_w m \log P + T_h (P-1)$$

∴

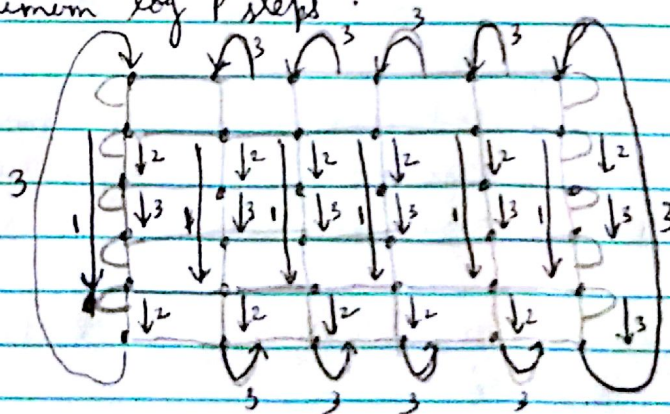
The steps involved in operation ~~see~~ of one to all broadcast in a 2-D torus, using cut through routing are →

- 1) Broadcast the message in the row (ring) in which the source element is present.



communication (broadcast) will be completed in the entire ring in exactly 3 steps (maximum $\log P$ steps).
* The no. above arrow represent the step.

- 2) Now each element in the 2nd ring will broadcast the message to the elements in their respective columns, again in maximum $\log P$ steps:



\therefore the cost of communication in 2-D torus \rightarrow

$$T_c = 2 \left(T_s \log P + t_w \cdot m \cdot \log P + T_h (P-1) \right)$$

here,

$$T_s = 15 \text{ micro seconds.}$$

$$T_h = 3 \text{ micro seconds}$$

$$t_w = 0.1 \text{ micro seconds.}$$

$$m = 1000 \text{ bytes.}$$

$$P = 6$$

\therefore here $P =$ no. of processors
in one row or column.

$$\therefore T_c = 2 \left(15 \times \log 6 + 0.1 \times 1000 \times \log 6 + 3 \cdot (6-1) \right)$$

$$= 2 \left(15 \times 2.58 + 100 \times 2.58 + 15 \right)$$

$$= 2 \left(38.7 + 258 + 15 \right)$$

$$= 2 \times 311.7$$

$$= \boxed{623.4 \text{ micro seconds}}$$

b)

Ans \rightarrow All to All scatter can be performed in the following steps \rightarrow

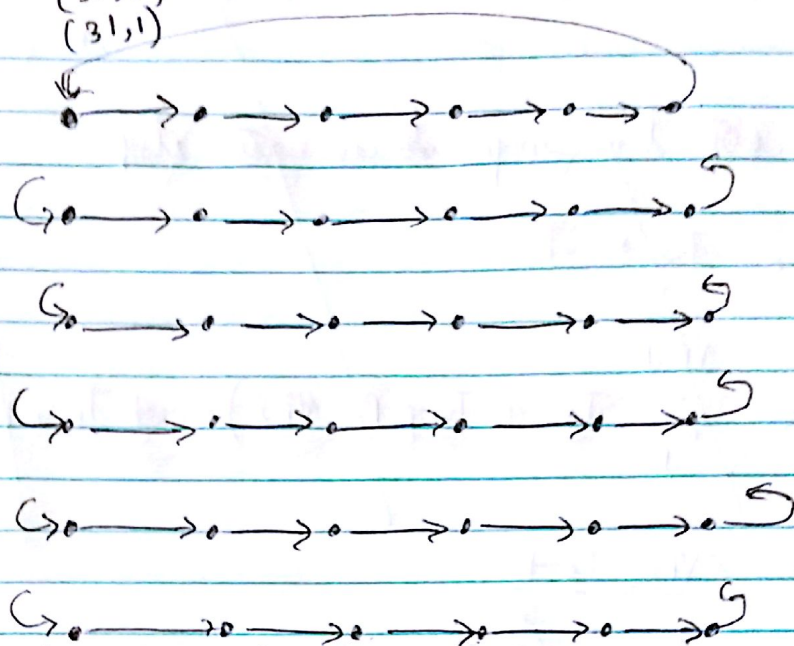
1) ~~In a row.~~

2) In all the rows, each node keeps the messages belonging to the nodes in its column, and transfers rest of all messages to other processors in nodes.

$$\text{no. of steps in distributing message across row} = (P-1)$$

$$\text{message size} = \sum_{i=1}^{P-1} (P-i)$$

$(31, 31)$
 $(31, 25)$
 $(31, 19)$
 $(31, 13)$
 $(31, 7)$
 $(31, 1)$



1st step.

$(1, 1)$	$(2, 2)$
$(1, 7)$	$(2, 8)$
$(1, 13)$	$(2, 14)$
$(1, 19)$	$(2, 20)$
$(1, 25)$	$(2, 26)$
$(1, 31)$	$(2, 32)$

(row, dest).

as we can see, all the processors keep the messages sent for others in some column.

The above process is repeated $(P-1)$ times.

at the end of it, each processor has the processor number for it, and the processor messages sent for others in some column.

2) In the second step, the same operation is carried out along the columns. The message size is reduced in following way $\rightarrow (\sqrt{P}-1), (\sqrt{P}-2) \dots 1$.

$$\text{Total cost} = (2T_s + T_w \cdot mP + T_R) (NP - 1)$$

$$= (2 \times 15 + 0.1 \times 1000 \times 36 + 3) (6 - 1)$$

$$= (30 + 3600 + 3) 5$$

$$= (3633) 5$$

$$= \boxed{18165 \text{ micro seconds.}}$$