DON BOSCO INSTITUTE OF TECHNOLOGY

Department of Information Science & Engineering

DB

NETWORKS LABORATORY

[SUBJECT CODE: 15CSL57]

LAB MANUAL

Prepared by

Prof. SUMATHI M
Prof. NAYANA D V

Laboratory Technical Support

Mr. RANGASWAMY H
Mr. PAVAN R

Don Bosco Institute of Technology

(Approved by AICTE & Affiliated to VTU, Karnataka)
Mysore Road, Kumbalgodu, Bengalore-560074
Phone: 91-080-28437028/31
Fax: 91-80-26710881,
E-mail: principal@dbit.co.in
Visit us at: www.dbit.co.in

Note : Student is required to solve one problem from PART-A and one problem from PART-B. Both the parts have equal weightage.

PART A – Simulation Exercises

The following experiments shall be conducted using either NS2/NS3 simulators.

1. Implement three nodes point-to-point network with duplex links between them. Set the queue size, vary the bandwidth and find the number of packets dropped.

2. Implement transmission of ping messages/trace route over a network topology consisting of 6 nodes and find the number of packets dropped due to congestion.

3. Implement an Ethernet LAN using n nodes and set multiple traffic nodes and plot congestion window for different source / destination.

4. Implement simple ESS and with transmitting nodes in wire-less LAN by simulation and determine the performance with respect to transmission of packets.

5. Implement and study the performance of GSM on NS2/NS3 (Using MAC layer) or equivalent environment.

6. Implement and study the performance of CDMA on NS2/NS3 (Using stack called Call net) or equivalent environment.

PART B

The following experiments are conducted using Java

7. Write a program for error detecting code using CRC-CCITT (16-bits).

8. Write a program to find the shortest path between vertices using bellman-ford algorithm.

9. Using TCP/IP sockets, write a client-server program to make client sending the file name and the server to send back the contents of the requested file if present. Implement the above program using as message queues or FIFOs as IPC channels.

10. Write a program on datagram socket for client/server to display the messages on client side, typed the server side.

11. Write a program for simple RSA algorithm to encrypt and decrypt the data.

12. Write a program for congestion control using Leaky bucket algorithm.

Introduction to NS

- Widely known as NS2, is simply an event driven simulation tool.
- Useful in studying the dynamic nature of communication networks.
- Simulation of wired as well as wireless network functions and protocols (e.g., routing algorithms, TCP, UDP) can be done using NS2.
- In general, NS2 provides users with a way of specifying such network protocols and simulating their corresponding behaviours.

Basic Architecture of NS2
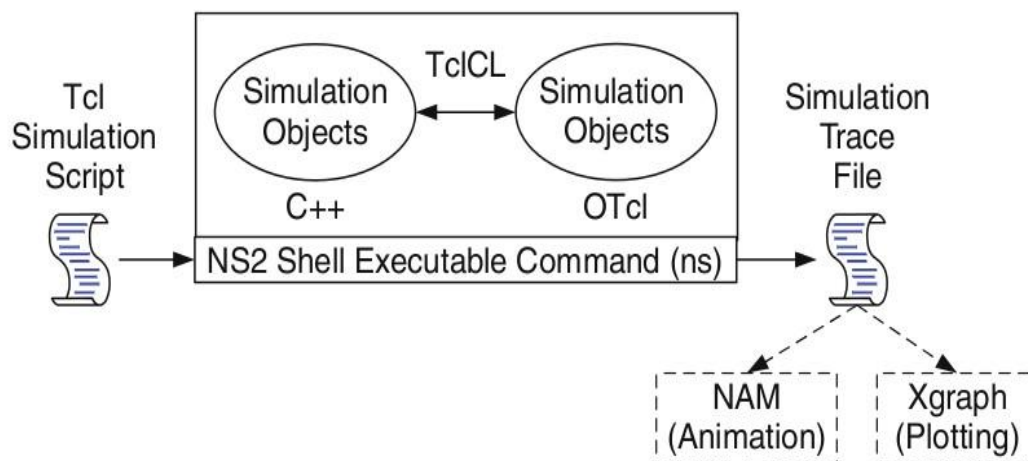
Tcl scripting

- Tcl is a general purpose scripting language. [Interpreter]
- Tcl runs on most of the platforms such as Unix, Windows, and Mac.
- The strength of Tcl is its simplicity.
- It is not necessary to declare a data type for variable prior to the usage.

Basics of TCL

Syntax: command   arg1   arg2   arg3

- **Hello World!**

      puts stdout{Hello,
      World!} Hello,
      World!

- **Variables**   Command Substitution

      set a 5   set len [string length foobar]
      set b $a   set len [expr [string length foobar] + 9]

- **Simple Arithmetic**
    expr 7.2 / 4
- **Procedures**

proc Diag {a b} {
set c [expr sqrt($a * $a +
$b * $b)] return $c }
        puts "Diagonal of a 3, 4 right triangle is
    [Diag 3 4]" Output: Diagonal of a 3, 4 right
    triangle is 5.0

- **Loops**

while{$i < $n} {   for {set i 0} {$i < $n} {incr i} {
. . .   . . .
}

Wired TCL Script Components

- Create the event scheduler
- Open new files & turn on the tracing
- Create the nodes
- Setup the links
- Configure the traffic type (e.g., TCP, UDP, etc)
- Set the time of traffic generation (e.g., CBR, FTP)
- Terminate the simulation

NS Simulator Preliminaries

1. Initialization and termination aspects of the ns simulator.
2. Definition of network nodes, links, queues and topology.
3. Definition of agents and of applications.
4. The nam visualization tool.
5. Tracing and random variables.

Initialization and Termination of TCL Script in NS-2

An ns simulation starts with the command

```
set ns [new Simulator]
```

Which is thus the first line in the tcl script? This line declares a new variable as using the
set command, you can call this variable as you wish, In general people declares it as ns
because it is an instance of the Simulator class, so an object the code [new Simulator] is
indeed the installation of the class Simulator using the reserved word new.
In order to have output files with data on the simulation (trace files) or files used
for visualization (nam files), we need to create the files using "open" command:

#Open the Trace file

```
set tracefile1 [open out.tr w]

$ns trace-all $tracefile1
```

#Open the NAM trace file

```
set namfile [open out.nam w]

$ns namtrace-all $namfile
```

The above creates a data trace file called "out.tr" and a nam visualization trace file called "out.nam". Within the tcl script, these files are not called explicitly by their names, but instead by pointers that are declared above and called "tracefile1" and "namfile" respectively. Remark that they begins with a # symbol. The second line open the file "out.tr" to be used for writing, declared with the letter "w". The third line uses a simulator method called trace-all that have as parameter the name of the file where the traces will go.

The last line tells the simulator to record all simulation traces in NAM input format. It also gives the file name that the trace will be written to later by the command $ns flush-trace. In our case, this will be the file pointed at by the pointer "$namfile",i.e the file "out.tr".

The termination of the program is done using a "finish" procedure.

#Define a 'finish' procedure

```
Proc finish { } {

global ns tracefile1 namfile

$ns flush-trace Close

$tracefile1 Close $namfile

Exec nam out.nam &

Exit 0

}
```

The word proc declares a procedure in this case called finish and without arguments. The word global is used to tell that we are using variables declared outside the procedure. The simulator method "flush-trace" will dump the traces on the respective files. The tcl command "close" closes the trace files defined before and exec executes the nam program for visualization. The command exit will ends the application and return the number 0 as status to the system. Zero is the default for a clean exit. Other values can be used to say that is a exit because something fails.

At the end of ns program we should call the procedure "finish" and specify at what time the termination should occur. For example,

```
$ns at 125.0 "finish"
```

will be used to call "finish" at time 125sec.Indeed, the at method of the simulator allows us to schedule events explicitly.

The simulation can then begin using the command

```
$ns run
```

Definition of a network of links and nodes

The way to define a node is

```
set n0 [$ns node]
```

We created a node that is printed by the variable n0. When we shall refer to that node in the script we shall thus write $n0.

Once we define several nodes, we can define the links that connect them. An example of a definition of a link is:

```
$ns duplex-link $n0 $n2 10Mb 10ms DropTail
```

Which means that $n0 and $n2 are connected using a bi-directional link that has 10ms of propagation delay and a capacity of 10Mb per sec for each direction.

To define a directional link instead of a bi-directional one, we should replace "duplex- link" by "simplex-link".

In NS, an output queue of a node is implemented as a part of each link whose input is that node. The definition of the link then includes the way to handle overflow at that queue. In our case, if the buffer capacity of the output queue is exceeded then the last packet to arrive is dropped. Many alternative options exist, such as the RED (Random Early Discard) mechanism, the FQ (Fair Queuing), the DRR (Deficit Round Robin), the stochastic Fair Queuing (SFQ) and the CBQ (which including a priority and a round-robin scheduler).

In ns, an output queue of a node is implemented as a part of each link whose input is that node. We should also define the buffer capacity of the queue related to each link. An example would be:

```
#set Queue Size of link (n0-n2) to 20

$ns queue-limit $n0 $n2 20
```

Agents and Applications

We need to define routing (sources, destinations) the agents (protocols) the application that use them.

FTP over TCP

TCP is a dynamic reliable congestion control protocol. It uses Acknowledgements created by the destination to know whether packets are well received.

There are number variants of the TCP protocol, such as Tahoe, Reno, NewReno, Vegas.

The type of agent appears in the first line:

```
set tcp [new Agent/TCP]
```

The command $ns attach-agent $n0 $tcp defines the source node of the tcp connection.

The command

```
set sink [new Agent /TCPSink]
```

Defines the behaviour of the destination node of TCP and assigns to it a pointer called sink.

#Setup a UDP connection

```
set udp [new Agent/UDP]

$ns attach-agent $n1 $udp set null

[new Agent/Null]

$ns attach-agent $n5 $null

$ns connect $udp $null

$udp set fid_2
```

#setup a CBR over UDP connection

```
set cbr [new Application/Traffic/CBR]

$cbr attach-agent $udp

$cbr set packetsize_ 100

$cbr set rate_ 0.01Mb

$cbr set random_ false
```

Above shows the definition of a CBR application using a UDP agent
The command $ns attach-agent $n4 $sink defines the destination node. The command

$ns connect $tcp $sink finally makes the TCP connection between the source and destination nodes.

TCP has many parameters with initial fixed defaults values that can be changed if mentioned explicitly. For example, the default TCP packet size has a size of 1000bytes.This can be changed to another value, say 552bytes, using the command $tcp set packetSize_ 552.
When we have several flows, we may wish to distinguish them so that we can identify them with different colors in the visualization part. This is done by the command $tcp set fid_ 1 that assigns to the TCP connection a flow identification of "1".We shall later give the flow identification of "2" to the UDP connection.

CBR over UDP

A UDP source and destination is defined in a similar way as in the case of TCP.
Instead of defining the rate in the command $cbr set rate_ 0.01Mb, one can define the time interval between transmission of packets using the command.

```
$cbr set interval_ 0.005
```

The packet size can be set to some value using

```
$cbr set packetSize_ <packet size>
```

Scheduling Events

NS is a discrete event based simulation. The tcp script defines when event should occur. The initializing command set ns [new Simulator] creates an event scheduler, and events are then scheduled using the format:

```
$ns at <time> <event>
```

The scheduler is started when running ns that is through the command $ns run.
The beginning and end of the FTP and CBR application can be done through the following command

```
$ns at 0.1 "$cbr start"

$ns at 1.0 " $ftp start"

$ns at 124.0 "$ftp stop"
```

Structure of Trace Files

When tracing into an output ASCII file, the trace is organized in 12 fields as follows in fig shown below, The meaning of the fields are:

Event	Time	From Node	To Node	PKT Type	PKT Size	Flags	Fid	Src Addr	Dest Addr	Seq Num	Pkt id

1. The first field is the event type. It is given by one of four possible symbols r, +, -, d which correspond respectively to receive (at the output of the link), enqueued, dequeued and dropped.

2. The second field gives the time at which the event occurs.
3. Gives the input node of the link at which the event occurs.
4. Gives the output node of the link at which the event occurs.
5. Gives the packet type (eg CBR or TCP)
6. Gives the packet size
7. Some flags
8. This is the flow id (fid) of IPv6 that a user can set for each flow at the input OTcl script one can further use this field for analysis purposes; it is also used when specifying stream color for the NAM display.
9. This is the source address given in the form of "node.port".
10. This is the destination address, given in the same form.
11. This is the network layer protocol's packet sequence number. Even though UDP implementations in a real network do not use sequence number, ns keeps track of UDP packet sequence number for analysis purposes
12. The last field shows the unique id of the packet.

XGRAPH

The xgraph program draws a graph on an x-display given data read from either data file or from standard input if no files are specified. It can display upto 64 independent data sets using different colors and line styles for each set. It annotates the graph with a title, axis labels, grid lines or tick marks, grid labels and a legend.

Syntax:

> Xgraph [options] file-name

Options are listed here

/-bd <color> (Border)

This specifies the border color of the xgraph window.
/-bg <color> (Background)

This specifies the background color of the xgraph window.
/-fg<color> (Foreground)

This specifies the foreground color of the xgraph window.
/-lf <fontname> (LabelFont)

All axis labels and grid labels are drawn using this font.

<u>/-t<string> (Title Text)</u>

This string is centered at the top of the graph.

<u>/-x <unit name> (XunitText)</u>

This is the unit name for the x-axis. Its default is "X".

<u>/-y <unit name> (YunitText)</u>

This is the unit name for the y-axis. Its default is "Y".

<u>Awk- An Advanced</u>

awk is a programmable, pattern-matching, and processing tool available in UNIX. It works equally well with text and numbers.

awk is not just a command, but a programming language too. In other words, awk utility is a pattern scanning and processing language. It searches one or more files to see if they contain lines that match specified patterns and then perform associated actions, such as writing the line to the standard output or incrementing a counter each time it finds a match.

Syntax:

```
awk option 'selection_criteria {action}' file(s)
```

Here, selection_criteria filters input and select lines for the action component to act upon. The selection_criteria is enclosed within single quotes and the action within the curly braces. Both the selection_criteria and action forms an awk program.

Example: $ awk '/manager/ {print}'

<u>emp.lst Variables</u>

Awk allows the user to use variables of there choice. You can now print a serial number, using the variable kount, and apply it those directors drawing a salary exceeding 6700:

$ awk –F"|" '$3 == "director" && $6 >_6700 { count =count+1

printf " %3f %20s %-12s %d\n", count,$2,$3,$6 }' empn.lst

THE –f OPTION: STORING awk PROGRAMS INA FILE

You should holds large awk programs in separate file and provide them with the awk extension for easier identification. Let's first store the previous program in the file empawk.awk:

$ cat empawk.awk

Observe that this time we haven't used quotes to enclose the awk program. You can now use awk with the –f *filename* option to obtain the same output:

```
Awk –F"|" –f empawk.awk empn.lst
```

The Begin And End Sections

Awk statements are usually applied to all lines selected by the address, and if there are no addresses, then they are applied to every line of input. But, if you have to print something before processing the first line, for example, a heading, then the BEGIN

section can be used gainfully. Similarly, the end section useful in printing some totals after processing is over.

The BEGIN and END sections are optional and take the form

BEGIN

{action}

END

{action}

These two sections, when present, are delimited by the body of the awk program. You can use them to print a suitable heading at the beginning and the average salary at the end.

Built-In Variables

Awk has several built-in variables. They are all assigned automatically, though it is also possible for a user to reassign some of them. You have already used NR, which signifies the record number of the current line. We'll now have a brief look at some of the other variable. *The FS Variable:* as stated elsewhere, awk uses a contiguous string of spaces as the default field delimiter. FS redefines this field separator, which in the sample database happens to be the |. When used at all, it must occur in the BEGIN section so that the body of the program knows its value before it starts processing:

BEGIN {FS="|"}

This is an alternative to the –F option which does the same thing.

The OFS Variable: when you used the print statement with comma-separated arguments, each argument was separated from the other by a space. This is awk's default output field separator, and can reassigned using the variable OFS in the BEGIN section:

BEGIN { OFS="~" }

When you reassign this variable with a ~ (tilde), awk will use this character for delimiting the print arguments. This is a useful variable for creating lines with delimited fields.

The NF variable: NF comes in quite handy for cleaning up a database of lines that don't contain the right number of fields. By using it on a file, say emp.lst, you can locate those lines not having 6 fields, and which have crept in due to faulty data entry:

$awk 'BEGIN {FS

= "|"} NF! =6 {

Print "Record No ", NR, "has", "fields"}' empx.lst

The FILENAME Variable: **FILENAME** stores the name of the current file being processed. Like grep and sed, awk can also handle multiple filenames in the command line. By default, awk doesn't print the filename, but you can instruct it to do so:

'$6<4000 {print FILENAME, $0 }'

With FILENAME, you can device logic that does different things depending on the file that is processed.

NS2 Installation

- NS2 is a free simulation tool.
- It runs on various platforms including UNIX (or Linux), Windows, and Mac systems.
- NS2 source codes are distributed in two forms: the all-in-one suite and the component- wise.
- 'all-in-one' package provides an "install" script which configures the NS2 environment and creates NS2 executable file using the "make" utility.

NS-2 installation steps in Linux

- ➢ Go to **Computer · File System ·** now paste the zip file **"ns-allinone-2.34.tar.gz"** into opt folder.
- ➢ Now **unzip** the file by typing the following **command**
  [root@localhost opt] # **tar -xzvf ns-allinone-2.34.tar.gz**

- ➢ After the files get extracted, we get ns-allinone-2.34 folder as well as zip file ns-allinone- 2.34.tar.gz
  [root@localhost opt] # **ns-allinone-2.34   ns-allinone-2.34.tar.gz**

- ➢ Now go to ns-allinone-2.33 folder and
  install it [root@localhost opt] # **cd ns-allinone-2.34** [root@localhost ns-allinone-2.33] # **./install**
- ➢ Once the installation is completed successfully we get certain pathnames in that terminal which must be pasted in **".bash_profile"** file.
- ➢ First **minimize the terminal** where installation is done and **open a new terminal** and open the file **".bash_profile"**
  [root@localhost ~] # **vi .bash_profile**

- ➢ When we open this file, we get a line in that file which is shown below

PATH=$PATH:$HOME/bin

To this line we must paste the path which is present in the previous terminal where **ns was installed**. First put **":"** then paste the path in-front of bin. That path is shown below.
":/opt/ns-allinone-2.33/bin:/opt/ns-allinone-2.33/tcl8.4.18/unix:/opt/ns-allinone-2.33/tk8.4.18/unix".

- ➢ In the next line type **"LD_LIBRARY_PATH=$LD_LIBRARY_PATH:"** and paste the **two paths** separated by **":"** which are present in the previous terminal i.e **Important notices section (1)**

  "/opt/ns-allinone-2.33/otcl-1.13:/opt/ns-allinone-2.33/lib"
- ➢ In the next line type **"TCL_LIBRARY=$TCL_LIBRARY:"** and paste the path which is present in previous terminal i.e **Important Notices section (2)**
  "/opt/ns-allinone-2.33/tcl8.4.18/library"
- ➢ In the next line type **"export LD_LIBRARY_PATH"**

➢ In the next line type **"export TCL_LIBRARY"**
➢ The next two lines are already present the file **"export PATH"** and **"unset USERNAME"**
➢ **Save the program (ESC + shift : wq and press enter)**
➢ Now in the terminal where we have opened **.bash_profile** file, type the following command to **check if path is updated correctly or not**

     [root@localhost ~] # **vi .bash_profile**

     [root@localhost ~] # **source .bash_profile**

➢ If **path is updated properly**, then we will **get the prompt** as shown
     below [root@localhost ~] #
➢ Now open the previous terminal where you have installed **ns**
     [root@localhost ns-allinone-2.33] #
➢ Here we need to configure three packages **"ns-2.33", "nam-1.13" and "xgraph-12.1"**
➢ **First**, configure **"ns-2.33"** package as shown
     below [root@localhost ns-allinone-2.33] #
     cd ns-2.33 [root@localhost ns-2.33] #
     ./configure [root@localhost ns-2.33] #
     make clean [root@localhost ns-2.33] #
     make [root@localhost ns-2.33] # **make**
     install [root@localhost ns-2.33] # **ns**
                      %
➢ If we get **"%"** symbol it indicates that **ns-2.33 configuration** was **successful**.
➢ **Second,** configure **"nam-1.13"** package as shown
     below [root@localhost ns-2.33] # **cd ..**
     [root@localhost ns-allinone-2.33] # **cd nam-**
     1.13 [root@localhost nam-1.13] # **./configure**
     [root@localhost nam-1.13] # **make clean**
     [root@localhost nam-1.13] # **make**
     [root@localhost nam-1.13] # **make install**
     [root@localhost nam-1.13] # **ns**

                      %
➢ If we get **"%"** symbol it indicates that **nam-1.13 configuration** was **successful**.
➢ **Third,** configure **"xgraph-12.1"** package as shown
     below [root@localhost nam-1.13] # **cd ..**
     [root@localhost ns-allinone-2.33] # **cd xgraph-**
     12.1 [root@localhost xgraph-12.1] # **./configure**
     [root@localhost xgraph-12.1] # **make clean**
     [root@localhost xgraph-12.1] # **make**
     [root@localhost xgraph-12.1] # **make install**
     [root@localhost xgraph-12.1] # **ns**

                      %

     This completes the installation process of "NS-2" simulator.

PART-A

1. Implement three nodes point-to-point network with duplex links between them. Set the queue size, vary the bandwidth and find the number of packets dropped.

Source Code:

1.tcl

```
set ns [new Simulator]
set nf [open lab1.nam w]
$ns namtrace-all $nf
set tf [open lab1.tr w]
$ns trace-all $tf
proc finish { } {
global ns nf tf
$ns flush-trace
close $nf
close $tf
exec nam lab1.nam &
exit 0
}
set n0 [$ns node]
set n1 [$ns node]
set n2 [$ns node]

$ns duplex-link $n0 $n1 2Mb 100ms DropTail
$ns duplex-link $n1 $n2 200Kb 100ms DropTail

$ns queue-limit $n0 $n1 50
$ns queue-limit $n1 $n2 10
set udp0 [new Agent/UDP]
$ns attach-agent $n0 $udp0
set cbr0 [new Application/Traffic/CBR]
$cbr0 set packetSize_ 500
$cbr0 set interval_ 0.005
$cbr0 attach-agent $udp0



set udp1 [new Agent/UDP]
$ns attach-agent $n1 $udp1
set cbr1 [new Application/Traffic/CBR]
$cbr1 attach-agent $udp1
set null0 [new Agent/Null]
```

```
$ns attach-agent $n2 $null0
$ns connect $udp0 $null0

$ns at 0.1 "$cbr0 start"

$ns at 3.0 "finish"
$ns run
```

AWK file (Open a new editor using "vi command" and write awk file and save with ".awk" extension)

/*immediately after BEGIN should open braces '{'

1.awk

```
BEGIN{
drop=0;
}
{
if($1=="d" )
{
drop++;
printf("%s\t%s\n",$5,$11);
}
}
END{
printf("Total number of %s packets dropped due to congestion
=%d\n",$5,drop);
}
```

Steps for execution

1) Open vi editor and type program. Program name should have the extension " **.tcl** "
   [root@localhost ~]# vi lab1.tcl

2) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

3) Open vi editor and type **awk** program. Program name should have the extension ".**awk** "
   [root@localhost ~]# vi lab1.awk

4) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

5) Run the simulation program
   [root@localhost~]# ns lab1.tcl

   i) Here **"ns"** indicates network simulator. We get the topology shown in the snapshot.
   ii) Now press the play button in the simulation window and the simulation will begins.

6) After simulation is completed run **awk file** to see the output ,
   [root@localhost~]# awk –f lab1.awk lab1.tr

7) To see the trace file contents open the file as ,

[root@localhost~]# vi lab1.tr

Trace file contains 12 columns:-

Event type, Event time, From Node, Source Node, Packet Type, Packet Size, Flags (indicated by --------), Flow ID, Source address, Destination address, Sequence ID, Packet ID

Topology
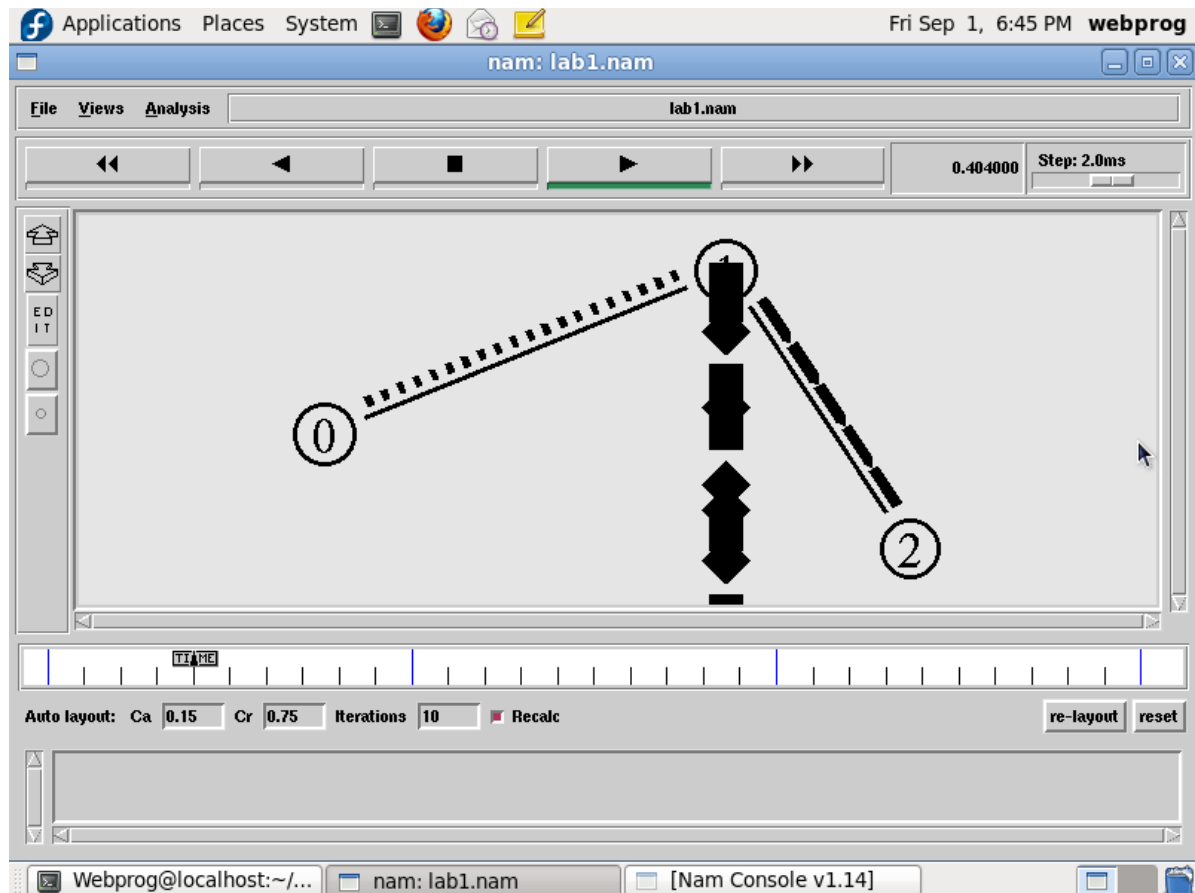
Output

#awk –f lab1.awk lab1.tr

```
cbr     524
cbr     526
cbr     527
cbr     528
cbr     530
cbr     531
cbr     532
cbr     534
cbr     535
cbr     536
cbr     538
cbr     539
cbr     540
cbr     542
cbr     543
cbr     544
cbr     546
cbr     547
cbr     548
cbr     550
cbr     551
cbr     552
cbr     554
cbr     555
cbr     556
cbr     558
cbr     559
Total number of cbr packets dropped due to congestion =411
[Webprog@localhost first]$
```

Note:

1. Set the queue size fixed from n0 to n2 as 10, n1-n2 to 10 and from n2-n3
   as 5. Syntax: To set the queue size

   $ns set queue-limit <from> <to> <size> Eg:
   $ns set queue-limit $n0 $n2 10

2. Go on varying the bandwidth from 10, 20 30 . . and find the number of
   packets dropped at the node 2

2. Implement transmission of ping messages/trace route over a network topology consisting of 6 nodes and find the number of packets dropped due to congestion.

Source Code:

lab2.tcl

```
set ns [ new Simulator ]
set nf [ open lab2.nam w ]
$ns namtrace-all $nf
set tf [ open lab2.tr w ]
$ns trace-all $tf
set n0 [$ns node]
set n1 [$ns node]
set n2 [$ns node]
set n3 [$ns node]
set n4 [$ns node]
```

```
set n5 [$ns node]
$n4 shape box
$ns duplex-link $n0 $n4 1005Mb 1ms DropTail
$ns duplex-link $n1 $n4 50Mb 1ms DropTail
$ns duplex-link $n2 $n4 2000Mb 1ms DropTail
$ns duplex-link $n3 $n4 200Mb 1ms DropTail
$ns duplex-link $n4 $n5 1Mb 1ms DropTail
set p1 [new Agent/Ping]
$ns attach-agent $n0 $p1
$p1 set packetSize_ 50000
$p1 set interval_ 0.0001
set p2 [new Agent/Ping]
$ns attach-agent $n1 $p2
set p3 [new Agent/Ping]
$ns attach-agent $n2 $p3
$p3 set packetSize_ 30000
$p3 set interval_ 0.00001
set p4 [new Agent/Ping]
$ns attach-agent $n3 $p4
set p5 [new Agent/Ping]
$ns attach-agent $n5 $p5
$ns queue-limit $n0 $n4 5
$ns queue-limit $n2 $n4 3
$ns queue-limit $n4 $n5 2
Agent/Ping instproc recv {from rtt} {
$self instvar node_
puts "node [$node_ id] received answer from $from with round trip
time $rtt msec"
}
# please provide space between $node_ and id. No space between $
and from. No
#space between and $ and rtt */
$ns connect $p1 $p5
$ns connect $p3 $p4
proc finish { } {
global ns nf tf
$ns flush-trace
close $nf
close $tf
exec nam lab2.nam &
exit 0
}
$ns at 0.1 "$p1 send"
$ns at 0.2 "$p1 send"
$ns at 0.3 "$p1 send"
$ns at 0.4 "$p1 send"
$ns at 0.5 "$p1 send"
$ns at 0.6 "$p1 send"
$ns at 0.7 "$p1 send"
$ns at 0.8 "$p1 send"
$ns at 0.9 "$p1 send"
$ns at 1.0 "$p1 send"
$ns at 1.1 "$p1 send"
$ns at 1.2 "$p1 send"
$ns at 1.3 "$p1 send"
$ns at 1.4 "$p1 send"
```

```
$ns at 1.5 "$p1 send"
$ns at 1.6 "$p1 send"
$ns at 1.7 "$p1 send"
$ns at 1.8 "$p1 send"
$ns at 1.9 "$p1 send"
$ns at 2.0 "$p1 send"
$ns at 2.1 "$p1 send"
$ns at 2.2 "$p1 send"
$ns at 2.3 "$p1 send"
$ns at 2.4 "$p1 send"
$ns at 2.5 "$p1 send"
$ns at 2.6 "$p1 send"
$ns at 2.7 "$p1 send"
$ns at 2.8 "$p1 send"
$ns at 2.9 "$p1 send"
$ns at 0.1 "$p3 send"
$ns at 0.2 "$p3 send"
$ns at 0.3 "$p3 send"
$ns at 0.4 "$p3 send"
$ns at 0.5 "$p3 send"
$ns at 0.6 "$p3 send"
$ns at 0.7 "$p3 send"
$ns at 0.8 "$p3 send"
$ns at 0.9 "$p3 send"
$ns at 1.0 "$p3 send"
$ns at 1.1 "$p3 send"
$ns at 1.2 "$p3 send"
$ns at 1.3 "$p3 send"
$ns at 1.4 "$p3 send"
$ns at 1.5 "$p3 send"
$ns at 1.6 "$p3 send"
$ns at 1.7 "$p3 send"
$ns at 1.8 "$p3 send"
$ns at 1.9 "$p3 send"
$ns at 2.0 "$p3 send"
$ns at 2.1 "$p3 send"
$ns at 2.2 "$p3 send"
$ns at 2.3 "$p3 send"
$ns at 2.4 "$p3 send"
$ns at 2.5 "$p3 send"
$ns at 2.6 "$p3 send"
$ns at 2.7 "$p3 send"
$ns at 2.8 "$p3 send"
$ns at 2.9 "$p3 send"
$ns at 3.0 "finish"
$ns run
```

AWK file (Open a new editor using "vi command" and write awk file and save with ".awk" extension)

lab2.awk

```
BEGIN {
c=0;
}
```

```
{
if($1=="d")
{
c++;
}
}
END{
printf("The number of packets dropped =%d\n",c);
}
```

Steps for execution

1) Open vi editor and type program. Program name should have the extension " **.tcl** "
   [root@localhost ~]# vi lab2.tcl

2) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

3) Open vi editor and type **awk** program. Program name should have the extension "**.awk** "

   [root@localhost ~]# vi lab2.awk

4) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

5) Run the simulation program

   [root@localhost~]# ns lab2.tcl

   i) Here **"ns"** indicates network simulator. We get the topology shown in the snapshot.

   ii) Now press the play button in the simulation window and the simulation will begins.

6) After simulation is completed run **awk file** to see the output ,

   [root@localhost~]# awk –f lab2.awk lab2.tr

7) To see the trace file contents open the file as ,

   [root@localhost~]# vi lab2.tr

Topology

Output

#awk –f lab2.awk lab2.tr
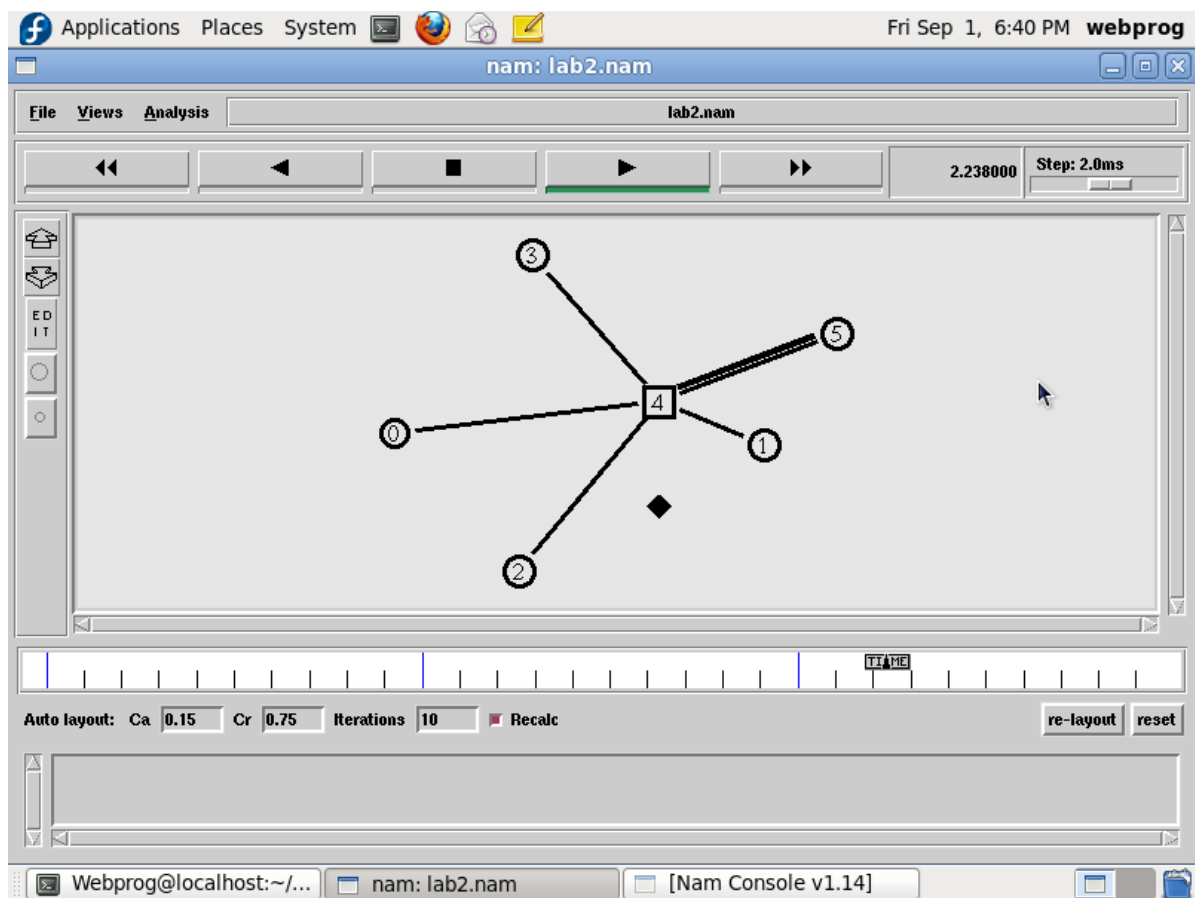
```
Applications  Places  System                          Fri Sep 1, 6:42 PM  webprog

                    Webprog@localhost:~/cn/second

File  Edit  View  Terminal  Help

node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 0 received answer from 5 with round trip time 804.9 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 0 received answer from 5 with round trip time 804.9 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 0 received answer from 5 with round trip time 804.9 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 0 received answer from 5 with round trip time 804.9 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 2 received answer from 3 with round trip time 5.3 msec
node 0 received answer from 5 with round trip time 804.9 msec
node 2 received answer from 3 with round trip time 5.3 msec
[Webprog@localhost second]$ awk -f lab2.awk lab2.tr
The number of packets dropped =20
[Webprog@localhost second]$

   Webprog@localhost:~/...
```

3. Implement an Ethernet LAN using n nodes and set multiple traffic nodes and plot congestion window for different source / destination.

Source Code:

lab3.tcl

```
set ns [new Simulator]
set tf [open lab3.tr w]
$ns trace-all $tf
set nf [open lab3.nam w]
$ns namtrace-all $nf
set n0 [$ns node]
$n0 color "magenta"
$n0 label "src1"
set n1 [$ns node]
$n1 color "red"
set n2 [$ns node]
$n2 color "magenta"
$n2 label "src2"
set n3 [$ns node]
```

```
$n3 color "blue"
$n3 label "dest2"
set n4 [$ns node]
$n4 shape square
set n5 [$ns node]
$n5 color "blue"
$n5 label "dest1"

$ns make-lan "$n0 $n1 $n2 $n3 $n4" 50Mb 100ms LL Queue/DropTail
Mac/802_3
$ns duplex-link $n4 $n5 1Mb 1ms DropTail
$ns duplex-link-op $n4 $n5 orient right
set tcp0 [new Agent/TCP]
$ns attach-agent $n0 $tcp0
set ftp0 [new Application/FTP]
$ftp0 attach-agent $tcp0
$ftp0 set packetSize_ 500
$ftp0 set interval_ 0.0001
set sink0 [new Agent/TCPSink]
$ns attach-agent $n5 $sink0
$ns connect $tcp0 $sink0

set tcp1 [new Agent/TCP]
$ns attach-agent $n2 $tcp1
set ftp1 [new Application/FTP]
$ftp1 attach-agent $tcp1
$ftp1 set packetSize_ 600
$ftp1 set interval_ 0.001
set sink1 [new Agent/TCPSink]
$ns attach-agent $n3 $sink1
$ns connect $tcp1 $sink1

set file1 [open file1.tr w]
$tcp0 attach $file1
set file2 [open file2.tr w]
$tcp1 attach $file2
$tcp0 trace cwnd_
$tcp1 trace cwnd_
proc finish { } {
global ns nf tf
$ns flush-trace
close $tf
close $nf
exec nam lab3.nam &
exit 0
}
$ns at 0.1 "$ftp0 start"
$ns at 5 "$ftp0 stop"
$ns at 7 "$ftp0 start"
$ns at 0.2 "$ftp1 start"
$ns at 8 "$ftp1 stop"
$ns at 14 "$ftp0 stop"
$ns at 10 "$ftp1 start"
$ns at 15 "$ftp1 stop"
$ns at 16 "finish"
$ns run
```

AWK file (Open a new editor using "vi command" and write awk file and save with ".awk" extension)

cwnd:- means congestion window

lab3.awk
```
BEGIN {
}
{
if($6= ="cwnd_") /* don't leave space after writing cwnd_ */
printf("%f\t%f\t\n",$1,$7); /* you must put \n in printf */
}
END {
}
```

Steps for execution

1) Open vi editor and type program. Program name should have the extension " **.tcl** "

   [root@localhost ~]# vi lab3.tcl

2) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

3) Open vi editor and type **awk** program. Program name should have the extension "**.awk** "

   [root@localhost ~]# vi lab3.awk

4) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

5) Run the simulation program

   [root@localhost~]# ns lab3.tcl

6) After simulation is completed run **awk file** to see the output ,

   i. **[root@localhost~]# awk –f lab3.awk file1.tr > a1**
   ii. **[root@localhost~]# awk –f lab3.awk file2.tr > a2**
   iii. **[root@localhost~]# xgraph a1 a2**

7) Here we are using the congestion window trace files i.e. **file1.tr** and **file2.tr** and we are redirecting the contents of those files to new files say **a1** and **a2** using **output redirection operator (>)**.

8) To see the trace file contents open the file as ,
   [root@localhost~]# vi lab3.tr

Topology
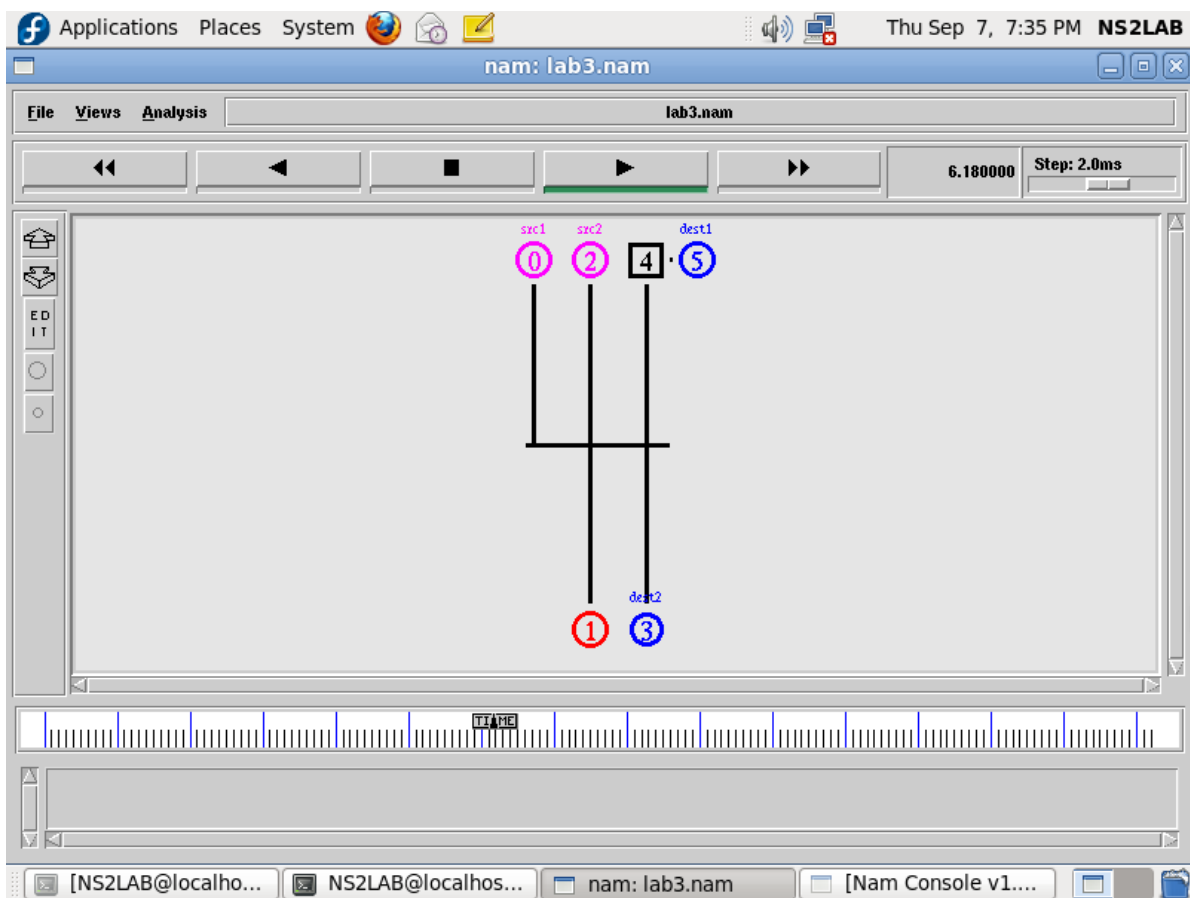
Output

    i.   **[root@localhost~]# awk –f lab3.awk file1.tr > file1.txt**
    ii.   **[root@localhost~]# awk –f lab3.awk file2.tr > file2.txt**
    iii.   **[root@localhost~]# xgraph file1.txt file2.txt**

4. Implement simple ESS and with transmitting nodes in wire-less LAN by simulation and determine the performance with respect to transmission of packets.

Source code

lab4.tcl

```
set ns [new Simulator]
set tf [open lab4.tr w]
$ns trace-all $tf
set topo [new Topography]
$topo load_flatgrid 1000 1000
set nf [open lab4.nam w]
$ns namtrace-all-wireless $nf 1000 1000
$ns node-config -adhocRouting DSDV \
-llType LL \
-macType Mac/802_11 \
-ifqType Queue/DropTail \
-ifqLen 50 \
-phyType Phy/WirelessPhy \
-channelType Channel/WirelessChannel \
-propType Propagation/TwoRayGround \
-antType Antenna/OmniAntenna \
-topoInstance $topo \
-agentTrace ON \
-routerTrace ON
create-god 3
```

```
set n0 [$ns node]
set n1 [$ns node]
set n2 [$ns node]
$n0 label "tcp0"
$n1 label "sink1/tcp1"
$n2 label "sink2"
$n0 set X_ 50
$n0 set Y_ 50
$n0 set Z_ 0
$n1 set X_ 100
$n1 set Y_ 100
$n1 set Z_ 0
$n2 set X_ 600
$n2 set Y_ 600
$n2 set Z_ 0
$ns at 0.1 "$n0 setdest 50 50 15"
$ns at 0.1 "$n1 setdest 100 100 25"
$ns at 0.1 "$n2 setdest 600 600 25"
set tcp0 [new Agent/TCP]
$ns attach-agent $n0 $tcp0
set ftp0 [new Application/FTP]
$ftp0 attach-agent $tcp0
set sink1 [new Agent/TCPSink]
$ns attach-agent $n1 $sink1
$ns connect $tcp0 $sink1
set tcp1 [new Agent/TCP]
$ns attach-agent $n1 $tcp1
set ftp1 [new Application/FTP]
$ftp1 attach-agent $tcp1
set sink2 [new Agent/TCPSink]
$ns attach-agent $n2 $sink2
$ns connect $tcp1 $sink2
$ns at 5 "$ftp0 start"
$ns at 5 "$ftp1 start"
$ns at 100 "$n1 setdest 550 550 15"
$ns at 190 "$n1 setdest 70 70 15"
proc finish { } {
global ns nf tf
$ns flush-trace
exec nam lab4.nam &
close $tf
exit 0
}
$ns at 250 "finish"
$ns run
```

AWK file (Open a new editor using "vi command" and write awk file and save with ".awk" extension)

lab4.awk

```
BEGIN{
count1=0
count2=0
```

```
 pack1=0
 pack2=0
time1=0
time2=0
 }
 {
        if($1= ="r"&& $3= ="_1_" && $4= ="AGT")
        {
                count1++
                pack1=pack1
                +$8
                time1=$2
        }
        if($1= ="r" && $3= ="_2_" && $4= ="AGT")
        {
                count2++
                pack2=pack2
                +$8

                time2=$2
        }
}

END{
printf("The Throughput from n0 to n1: %f Mbps \n",
((count1*pack1*8)/(time1*1000000))); printf("The Throughput from n1 to n2: %f
Mbps", ((count2*pack2*8)/(time2*1000000)));
}
```

Steps for execution

1) Open vi editor and type program. Program name should have the extension " **.tcl** "

   [root@localhost ~]# vi lab4.tcl

2) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

3) Open vi editor and type **awk** program. Program name should have the extension "**.awk** "

   [root@localhost ~]# vi lab4.awk

4) Save the program by pressing **"ESC key"** first, followed by **"Shift and :"** keys simultaneously and type **"wq"** and press **Enter key**.

5) Run the simulation program

   [root@localhost~]# ns lab4.tcl

   iii) Here **"ns"** indicates network simulator. We get the topology shown in the snapshot.

   iv) Now press the play button in the simulation window and the simulation will begins.

6) After simulation is completed run **awk file** to see the output ,

   [root@localhost~]# awk –f lab4.awk lab4.tr

7) To see the trace file contents open the file as ,

[root@localhost~]# vi lab4.tr

Topology
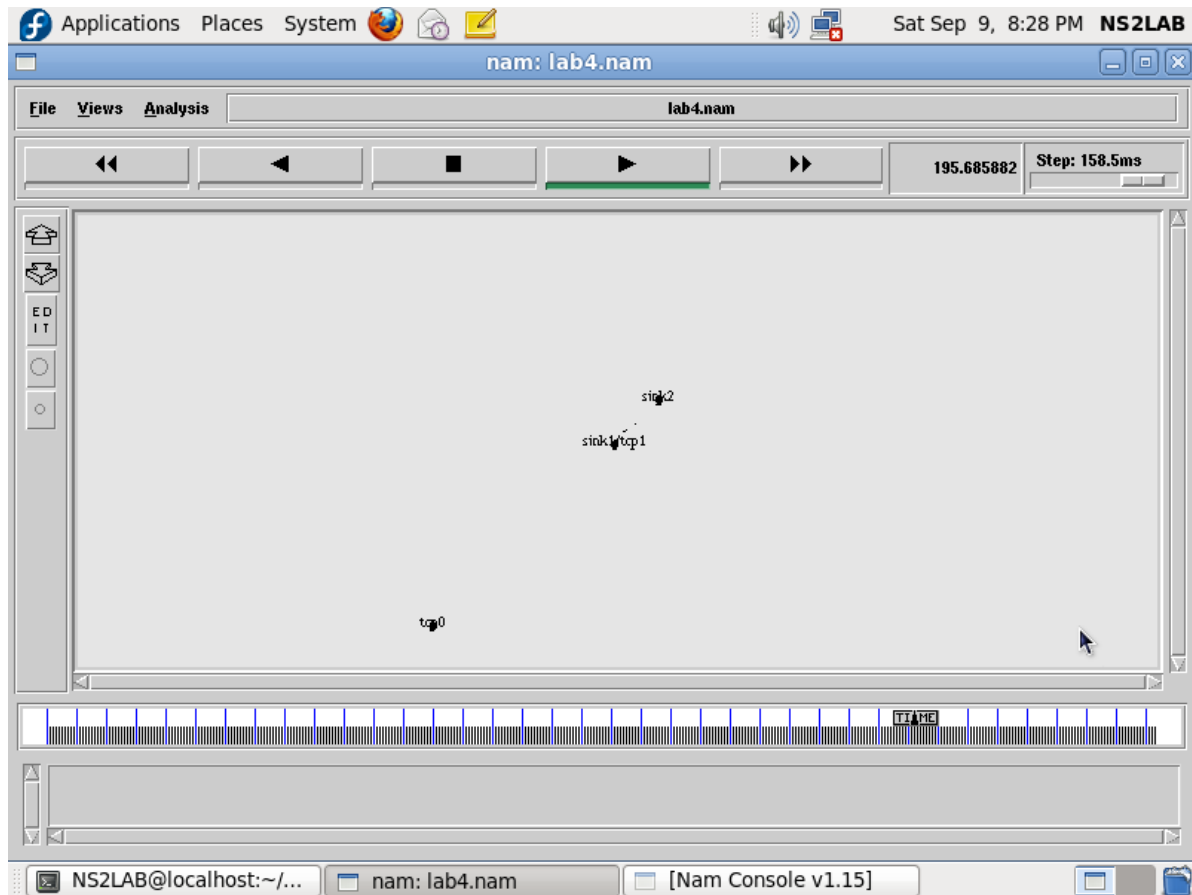
Node 1 and 2 are communicating

Trace File

```
[v]                        root@localhost:~                          [_][=][x]
File  Edit  View  Terminal  Tabs  Help
s 0.036400876 _0_ RTR   --- 0 message 32 [0 0 0 0] ------- [0:255 -1:255 32 0]   [▲]
r 0.037421112 _1_ RTR   --- 0 message 32 [0 ffffffff 0 800] ------- [0:255 -1:255
 32 0]
M 0.10000 0 (50.00, 50.00, 0.00), (50.00, 50.00), 15.00
M 0.10000 1 (100.00, 100.00, 0.00), (100.00, 100.00), 25.00
M 0.10000 2 (600.00, 600.00, 0.00), (600.00, 600.00), 25.00
s 0.182633994 _1_ RTR   --- 1 message 32 [0 0 0 0] ------- [1:255 -1:255 32 0]
r 0.183694230 _0_ RTR   --- 1 message 32 [0 ffffffff 1 800] ------- [1:255 -1:255
 32 0]
s 0.882774710 _2_ RTR   --- 2 message 32 [0 0 0 0] ------- [2:255 -1:255 32 0]
s 5.000000000 _0_ AGT   --- 3 tcp 40 [0 0 0 0] ------- [0:0 1:0 32 0] [0 0] 0 0
r 5.000000000 _0_ RTR   --- 3 tcp 40 [0 0 0 0] ------- [0:0 1:0 32 0] [0 0] 0 0
s 5.000000000 _0_ RTR   --- 3 tcp 60 [0 0 0 0] ------- [0:0 1:0 32 1] [0 0] 0 0
s 5.000000000 _1_ AGT   --- 4 tcp 40 [0 0 0 0] ------- [1:1 2:0 32 0] [0 0] 0 0
r 5.000000000 _1_ RTR   --- 4 tcp 40 [0 0 0 0] ------- [1:1 2:0 32 0] [0 0] 0 0
r 5.004812650 _1_ AGT   --- 3 tcp 60 [13a 1 0 800] ------- [0:0 1:0 32 1] [0 0] 1
 0
s 5.004812650 _1_ AGT   --- 5 ack 40 [0 0 0 0] ------- [1:0 0:0 32 0] [0 0] 0 0
r 5.004812650 _1_ RTR   --- 5 ack 40 [0 0 0 0] ------- [1:0 0:0 32 0] [0 0] 0 0
s 5.004812650 _1_ RTR   --- 5 ack 60 [0 0 0 0] ------- [1:0 0:0 32 0] [0 0] 0 0
r 5.006977357 _0_ AGT   --- 5 ack 60 [13a 0 1 800] ------- [1:0 0:0 32 0] [0 0] 1
 0
s 5.006977357 _0_ AGT   --- 6 tcp 1040 [0 0 0 0] ------- [0:0 1:0 32 0] [1 0] 0 0
"lab8.tr" 128664L, 11456314C                              1,1            Top  [▼]
```

Here "M" indicates mobile nodes, "AGT" indicates Agent Trace, "RTR" indicates Router Trace

Output

5. Implement and study the performance of GSM on NS2/NS3 (Using MAC layer) or equivalent environment.

Second Generation (2G) technology is based on the technology known as global system for mobile communication (GSM). This technology enabled various networks to provide services like text messages, picture messages and MMS. The technologies used in 2G are either TDMA (Time Division Multiple Access) which divides signal into different time slots or CDMA (Code Division Multiple Access) which allocates a special code to each user so as to communicate over a multiplex physical channel.

GSM uses a variation of time division multiple access (TDMA). 2G networks developed as a replacement for first generation (1G) analog cellular networks, and the GSM standard originally described as a digital, circuit-switched network optimized for full duplex voice telephony. This expanded over time to include data communications, first by circuit-switched transport, then by packet data transport via GPRS (General Packet Radio Services).

GSM can be implemented on all the versions of NS2 (Since year 2004: ns-2.27, and later versions of NS2)

Design:
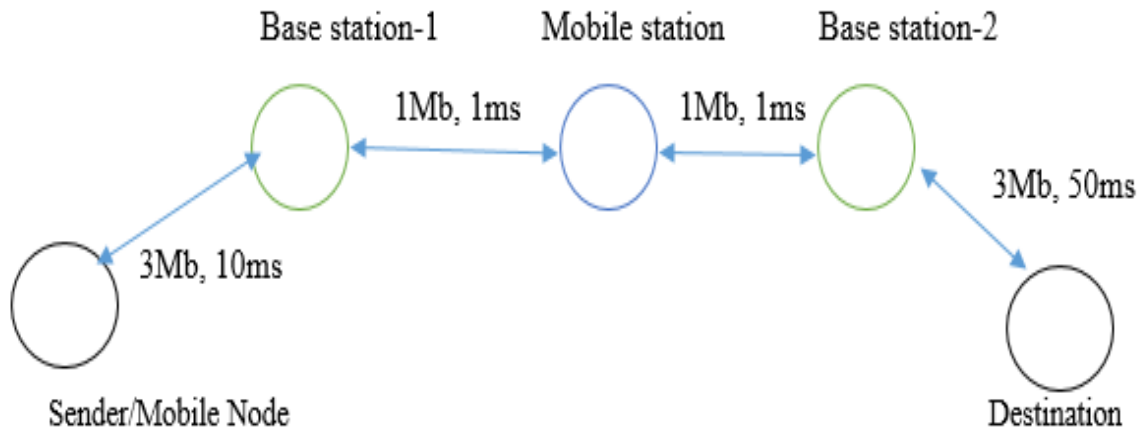
Source Code:

 # General Parameters
set opt(title) zero   ;
set opt(stop) 100   ;# Stop time.
set opt(ecn) 0   ;
 # Topology
set opt(type) gsm   ;#type of link:
set opt(secondDelay) 55   ;# average delay of access links in ms
 # AQM parameters
set opt(minth) 30   ;
set opt(maxth) 0   ;
set opt(adaptive) 1   ;# 1 for Adaptive RED, 0 for plain RED
 # Traffic generation.
set opt(flows) 0   ;# number of long-lived TCP flows
set opt(window) 30   ;# window for long-lived traffic
set opt(web) 2   ;# number of web sessions
 # Plotting statistics.
set opt(quiet)   0   ;# popup anything
set opt(wrap)   100   ;# wrap plots
set opt(srcTrace) is   ;# where to plot traffic
set opt(dstTrace) bs2   ;# where to plot traffic
set opt(gsmbuf) 10   ; # buffer size for gsm

#default downlink bandwidth in bps
set bwDL(gsm) 9600
#default uplink bandwidth in bps
set bwUL(gsm) 9600
#default downlink propagation delay in seconds
set propDL(gsm) .500
#default uplink propagation delay in seconds
set propUL(gsm) .500
#default buffer size in packets
set buf(gsm)   10

```
set ns [new Simulator]
set tf [open out.tr w]
$ns trace-all $tf

set nodes(is) [$ns node]
set nodes(ms) [$ns node]
set nodes(bs1) [$ns node]
set nodes(bs2) [$ns node]
set nodes(lp) [$ns node]

 proc cell_topo {} {
 global ns nodes
 $ns duplex-link $nodes(lp) $nodes(bs1) 3Mbps 10ms DropTail
 $ns duplex-link $nodes(bs1) $nodes(ms) 1 1 RED
 $ns duplex-link $nodes(ms) $nodes(bs2) 1 1 RED
 $ns duplex-link $nodes(bs2) $nodes(is) 3Mbps 50ms DropTail
 puts "Cell Topology"
 }
 proc set_link_params {t} {
 global ns nodes bwUL bwDL propUL propDL buf
 $ns bandwidth $nodes(bs1) $nodes(ms) $bwDL($t) simplex
 $ns bandwidth $nodes(ms) $nodes(bs1) $bwUL($t) simplex
 $ns bandwidth $nodes(bs2) $nodes(ms) $bwDL($t) simplex
 $ns bandwidth $nodes(ms) $nodes(bs2) $bwUL($t) simplex
 $ns delay $nodes(bs1) $nodes(ms) $propDL($t) simplex
 $ns delay $nodes(ms) $nodes(bs1) $propDL($t) simplex
 $ns delay $nodes(bs2) $nodes(ms) $propDL($t) simplex
 $ns delay $nodes(ms) $nodes(bs2) $propDL($t) simplex
 $ns queue-limit $nodes(bs1) $nodes(ms) $buf($t)
 $ns queue-limit $nodes(ms) $nodes(bs1) $buf($t)
 $ns queue-limit $nodes(bs2) $nodes(ms) $buf($t)
 $ns queue-limit $nodes(ms) $nodes(bs2) $buf($t)
}
# RED and TCP parameters
Queue/RED set summarystats_ true
Queue/DropTail set summarystats_ true
Queue/RED set adaptive_ $opt(adaptive)
Queue/RED set q_weight_ 0.0
Queue/RED set thresh_ $opt(minth)
Queue/RED set maxthresh_ $opt(maxth)
Queue/DropTail set shrink_drops_ true
Agent/TCP set ecn_ $opt(ecn)
Agent/TCP set window_ $opt(window)
DelayLink set avoidReordering_ true

source web.tcl

#Create topology
switch $opt(type) {
```

```
 gsm -
 gprs -
 umts {cell_topo}
}
 set_link_params $opt(type)
 $ns insert-delayer $nodes(ms) $nodes(bs1) [new Delayer]
 $ns insert-delayer $nodes(bs1) $nodes(ms) [new Delayer]
 $ns insert-delayer $nodes(ms) $nodes(bs2) [new Delayer]
 $ns insert-delayer $nodes(bs2) $nodes(ms) [new Delayer]
```

 # Set up forward TCP connection
```
if {$opt(flows) == 0} {
        set tcp1 [$ns create-connection TCP/Sack1 $nodes(is) TCPSink/Sack1 $nodes(lp) 0]
        set ftp1 [[set tcp1] attach-app FTP]
        $ns at 0.8 "[set ftp1] start"
}
if {$opt(flows) > 0} {
   set tcp1 [$ns create-connection TCP/Sack1 $nodes(is) TCPSink/Sack1 $nodes(lp) 0]
   set ftp1 [[set tcp1] attach-app FTP]
   $tcp1 set window_ 100
   $ns at 0.0  "[set ftp1] start"
   $ns at 3.5  "[set ftp1] stop"
   set tcp2 [$ns create-connection TCP/Sack1 $nodes(is) TCPSink/Sack1 $nodes(lp) 0]
   set ftp2 [[set tcp2] attach-app FTP]
   $tcp2 set window_ 3
   $ns at 1.0  "[set ftp2] start"
   $ns at 8.0  "[set ftp2] stop"
}

proc stop {} {
        global nodes opt nf
        set wrap $opt(wrap)
        set sid [$nodes($opt(srcTrace)) id]
        set did [$nodes($opt(dstTrace)) id]
        if {$opt(srcTrace) == "is"} {
                set a "-a out.tr"
        } else {
                set a "out.tr"
        }
        set GETRC "../../../bin/getrc"
     set RAW2XG "../../../bin/raw2xg"
     exec $GETRC -s $sid -d $did -f 0 out.tr | \
       $RAW2XG -s 0.01  -m $wrap -r > plot.xgr
     exec $GETRC -s $did -d $sid -f 0 out.tr | \
       $RAW2XG -a -s 0.01 -m $wrap  >> plot.xgr
     exec $GETRC -s $sid -d $did -f 1 out.tr | \
       $RAW2XG -s 0.01 -m $wrap -r >> plot.xgr
     exec $GETRC -s $did -d $sid -f 1 out.tr  | \
       $RAW2XG -s 0.01 -m $wrap -a  >> plot.xgr
        exec ./xg2gp.awk plot.xgr
```

```
    if {!$opt(quiet)} {
        exec xgraph -bb -tk -nl -m -x time -y packets plot.xgr &
    }
        exit 0
}
$ns at $opt(stop) "stop"
$ns run
```

**Output:**

GSM Trace file:

6. Implement and study the performance of CDMA on NS2/NS3 (Using stack called Call net) or equivalent environment.

        3G networks developed as a replacement for second generation (2G) GSM standard network with full duplex voice telephony. CDMA is used as the access method in many mobile phone standards. IS-95, also called cdmaOne, and its 3G evolution CDMA2000, are often simply referred to as CDMA, but UMTS(The Universal Mobile Telecommunications System is a third generation mobile cellular system for networks based on the GSM standard.), the 3G standard used by GSM carriers, also uses wideband CDMA. Long-Term Evolution (LTE) is a standard for high-speed wireless communication which uses CDMA network technology.

        3G technology generally refers to the standard of accessibility and speed of mobile devices. The standards of the technology were set by the International Telecommunication Union (ITU). This technology enables use of various services like GPS (Global Positioning System), mobile television and video conferencing. It not only enables them to be used worldwide, but also provides with better bandwidth and increased speed. The main aim of this technology is to allow much better coverage and growth with minimum investment.

CDMA can be implemented on all the versions of NS2 (Since year 2004: ns-2.27, and later versions of NS2)

Design:
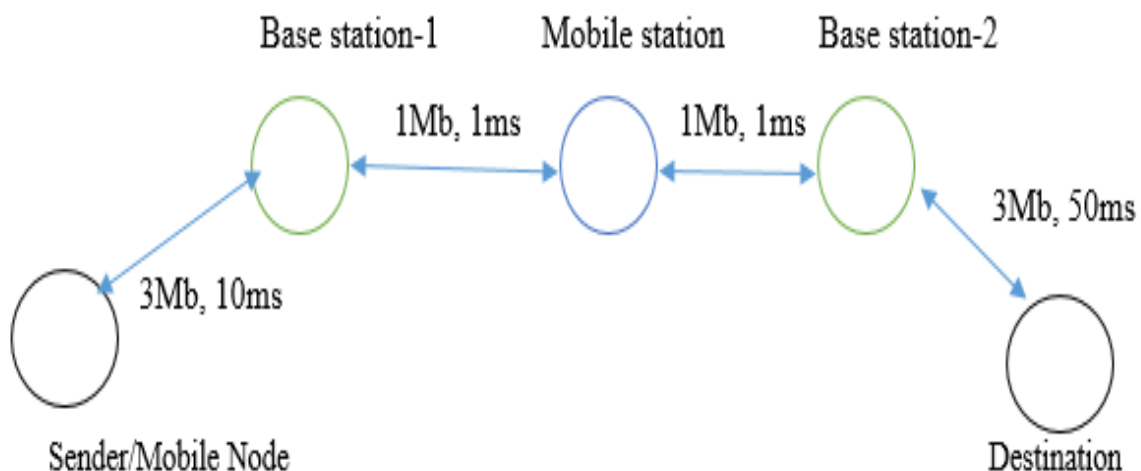
Source Code:

 # General Parameters
```
set opt(title) zero      ;
set opt(stop) 100        ;# Stop time.
set opt(ecn) 0           ;
```
 # Topology
```
set opt(type) umts       ;#type of link:
set opt(secondDelay) 55        ;# average delay of access links in ms
```
 # AQM parameters
```
set opt(minth) 30        ;
```

```
set opt(maxth) 0        ;
set opt(adaptive) 1     ;# 1 for Adaptive RED, 0 for plain RED
# Traffic generation.
set opt(flows) 0        ;# number of long-lived TCP flows
set opt(window) 30      ;# window for long-lived traffic
set opt(web) 2          ;# number of web sessions
# Plotting statistics.
set opt(quiet)   0      ;# popup anything
set opt(wrap)    100    ;# wrap plots
set opt(srcTrace) is    ;# where to plot traffic
set opt(dstTrace) bs2   ;# where to plot traffic
set opt(umtsbuf) 10              ; # buffer size for umts

#default downlink bandwidth in bps
set bwDL(umts) 384000
#default uplink bandwidth in bps
set bwUL(umts) 64000
#default downlink propagation delay in seconds
set propDL(umts) .150
#default uplink propagation delay in seconds
set propUL(umts) .150
#default buffer size in packets
set buf(umts)   20

set ns [new Simulator]
set tf [open out.tr w]
$ns trace-all $tf

set nodes(is) [$ns node]
set nodes(ms) [$ns node]
set nodes(bs1) [$ns node]
set nodes(bs2) [$ns node]
set nodes(lp) [$ns node]

proc cell_topo {} {
  global ns nodes
  $ns duplex-link $nodes(lp) $nodes(bs1) 3Mbps 10ms DropTail
  $ns duplex-link $nodes(bs1) $nodes(ms) 1 1 RED
  $ns duplex-link $nodes(ms) $nodes(bs2) 1 1 RED
  $ns duplex-link $nodes(bs2) $nodes(is) 3Mbps 50ms DropTail
  puts "Cell Topology"
}
proc set_link_params {t} {
  global ns nodes bwUL bwDL propUL propDL buf
  $ns bandwidth $nodes(bs1) $nodes(ms) $bwDL($t) simplex
  $ns bandwidth $nodes(ms) $nodes(bs1) $bwUL($t) simplex
  $ns delay $nodes(bs1) $nodes(ms) $propDL($t) simplex
  $ns delay $nodes(ms) $nodes(bs1) $propDL($t) simplex
  $ns queue-limit $nodes(bs1) $nodes(ms) $buf($t)
  $ns queue-limit $nodes(ms) $nodes(bs1) $buf($t)
```

```
    $ns bandwidth $nodes(bs2) $nodes(ms) $bwDL($t) simplex
    $ns bandwidth $nodes(ms) $nodes(bs2) $bwUL($t) simplex
    $ns delay $nodes(bs2) $nodes(ms) $propDL($t) simplex
    $ns delay $nodes(ms) $nodes(bs2) $propDL($t) simplex
    $ns queue-limit $nodes(bs2) $nodes(ms) $buf($t)
    $ns queue-limit $nodes(ms) $nodes(bs2) $buf($t)
}
```

 # RED and TCP parameters
```
Queue/RED set summarystats_ true
Queue/DropTail set summarystats_ true
Queue/RED set adaptive_ $opt(adaptive)
Queue/RED set q_weight_ 0.0
Queue/RED set thresh_ $opt(minth)
Queue/RED set maxthresh_ $opt(maxth)
Queue/DropTail set shrink_drops_ true
Agent/TCP set ecn_ $opt(ecn)
Agent/TCP set window_ $opt(window)
DelayLink set avoidReordering_ true
```

source web.tcl

#Create topology
```
switch $opt(type) {
 umts {cell_topo}
}
  set_link_params $opt(type)
  $ns insert-delayer $nodes(ms) $nodes(bs1) [new Delayer]
  $ns insert-delayer $nodes(bs1) $nodes(ms) [new Delayer]
  $ns insert-delayer $nodes(ms) $nodes(bs2) [new Delayer]
  $ns insert-delayer $nodes(bs2) $nodes(ms) [new Delayer]
```

 # Set up forward TCP connection
```
if {$opt(flows) == 0} {
        set tcp1 [$ns create-connection TCP/Sack1 $nodes(is) TCPSink/Sack1 $nodes(lp) 0]
        set ftp1 [[set tcp1] attach-app FTP]
        $ns at 0.8 "[set ftp1] start"
}
if {$opt(flows) > 0} {
   set tcp1 [$ns create-connection TCP/Sack1 $nodes(is) TCPSink/Sack1 $nodes(lp) 0]
   set ftp1 [[set tcp1] attach-app FTP]
   $tcp1 set window_ 100
   $ns at 0.0  "[set ftp1] start"
   $ns at 3.5  "[set ftp1] stop"
   set tcp2 [$ns create-connection TCP/Sack1 $nodes(is) TCPSink/Sack1 $nodes(lp) 0]
   set ftp2 [[set tcp2] attach-app FTP]
   $tcp2 set window_ 3
   $ns at 1.0  "[set ftp2] start"
   $ns at 8.0  "[set ftp2] stop"
}
```

```
proc stop { } {
        global nodes opt nf
        set wrap $opt(wrap)
        set sid [$nodes($opt(srcTrace)) id]
        set did [$nodes($opt(dstTrace)) id]
        if {$opt(srcTrace) == "is"} {
                set a "-a out.tr"
        } else {
                set a "out.tr"
        }
        set GETRC "../../../bin/getrc"
    set RAW2XG "../../../bin/raw2xg"

    exec $GETRC -s $sid -d $did -f 0 out.tr | \
      $RAW2XG -s 0.01  -m $wrap -r > plot.xgr
    exec $GETRC -s $did -d $sid -f 0 out.tr | \
      $RAW2XG -a -s 0.01 -m $wrap  >> plot.xgr

    exec $GETRC -s $sid -d $did -f 1 out.tr | \
      $RAW2XG -s 0.01 -m $wrap -r >> plot.xgr
    exec $GETRC -s $did -d $sid -f 1 out.tr  | \
      $RAW2XG -s 0.01 -m $wrap -a  >> plot.xgr

      exec ./xg2gp.awk plot.xgr
    if {!$opt(quiet)} {
        exec xgraph -bb -tk -nl -m -x time -y packets plot.xgr &
    }
      exit 0
}
$ns at $opt(stop) "stop"
$ns run
```

CDMA Trace File

NS2LAB@localhost:~/ns-allinone-2.35/ns-2.35/tcl/ex/wireless-scripts

File Edit View Terminal Help

```
+ 0.8 0 3 tcp 40 ------- 0 0.0 4.0 0 0
- 0.8 0 3 tcp 40 ------- 0 0.0 4.0 0 0
r 0.850107 0 3 tcp 40 ------- 0 0.0 4.0 0 0
+ 0.850107 3 1 tcp 40 ------- 0 0.0 4.0 0 0
- 0.850107 3 1 tcp 40 ------- 0 0.0 4.0 0 0
r 1.00094 3 1 tcp 40 ------- 0 0.0 4.0 0 0
+ 1.00094 1 2 tcp 40 ------- 0 0.0 4.0 0 0
- 1.00094 1 2 tcp 40 ------- 0 0.0 4.0 0 0
r 1.15594 1 2 tcp 40 ------- 0 0.0 4.0 0 0
+ 1.15594 2 4 tcp 40 ------- 0 0.0 4.0 0 0
- 1.15594 2 4 tcp 40 ------- 0 0.0 4.0 0 0
r 1.166047 2 4 tcp 40 ------- 0 0.0 4.0 0 0
+ 1.166047 4 2 ack 40 ------- 0 4.0 0.0 0 1
- 1.166047 4 2 ack 40 ------- 0 4.0 0.0 0 1
r 1.176153 4 2 ack 40 ------- 0 4.0 0.0 0 1
+ 1.176153 2 1 ack 40 ------- 0 4.0 0.0 0 1
- 1.176153 2 1 ack 40 ------- 0 4.0 0.0 0 1
r 1.326987 2 1 ack 40 ------- 0 4.0 0.0 0 1
+ 1.326987 1 3 ack 40 ------- 0 4.0 0.0 0 1
- 1.326987 1 3 ack 40 ------- 0 4.0 0.0 0 1
r 1.481987 1 3 ack 40 ------- 0 4.0 0.0 0 1
+ 1.481987 3 0 ack 40 ------- 0 4.0 0.0 0 1
- 1.481987 3 0 ack 40 ------- 0 4.0 0.0 0 1
r 1.532093 3 0 ack 40 ------- 0 4.0 0.0 0 1
+ 1.532093 0 3 tcp 1040 ------- 0 0.0 4.0 1 2
- 1.532093 0 3 tcp 1040 ------- 0 0.0 4.0 1 2
+ 1.532093 0 3 tcp 1040 ------- 0 0.0 4.0 2 3
- 1.534867 0 3 tcp 1040 ------- 0 0.0 4.0 2 3
"out.tr" 18006L, 901063C                                    2,1        Top
```

NS2LAB@localhost:~/...   [xgraph]   [xgraph]

PART-B

Java is a general-purpose computer programming language that is simple, concurrent, class-based, object-oriented language. The compiled Java code can run on all platforms that support Java without the need for recompilation hence Java is called as "write once, run anywhere" (WORA).The Java compiled intermediate output called "byte-code" that can run on any Java virtual machine (JVM) regardless of computer architecture. The language derives much of its syntax from C and C++, but it has fewer low-level facilities than either of them.

In Linux operating system Java libraries are preinstalled. It's very easy and convenient to compile and run Java programs in Linux environment. To compile and run Java Program is a two-step process:

1. Compile Java Program from Command Prompt
   [root@host ~]# javac Filename.java

   The Java compiler (Javac) compiles java program and generates a byte-code with the same file name and .class extension.

2. Run Java program from Command Prompt
   [root@host ~]# java Filename

   The java interpreter (Java) runs the byte-code and gives the respective output. It is important to note that in above command we have omitted the .class suffix of the byte- code (Filename.class).

7. Write a program for error detecting code using CRC-CCITT (16-bits).

Whenever digital data is stored or interfaced, data corruption might occur. Since the beginning of computer science, developers have been thinking of ways to deal with this type of problem. For serial data they came up with the solution to attach a parity bit to each sent byte. This simple detection mechanism works if an odd number of bits in a byte changes, but an even number of false bits in one byte will not be detected by the parity check. To overcome this problem developers have searched for mathematical sound mechanisms to detect multiple false bits. The CRC calculation or *cyclic redundancy check* was the result of this. Nowadays CRC calculations are used in all types of communications. All packets sent over a network connection are checked with a CRC. Also each data block on your hard disk has a CRC value attached to it. Modern computer

world cannot do without these CRC calculations. So let's see why they are so widely used. The answer is simple; they are powerful, detect many types of errors and are extremely fast to calculate especially when dedicated hardware chips are used.

The idea behind CRC calculation is to look at the data as one large binary number. This number is divided by a certain value and the remainder of the calculation is called the CRC. Dividing in the CRC calculation at first looks to cost a lot of computing power, but it can be performed very quickly if we use a method similar to the one learned at school. We will as an example calculate the remainder for the character 'm'—which is 1101101 in binary notation— by dividing it by 19 or 10011. Please note that 19 is an odd number. This is necessary as we will see further on. Please refer to your schoolbooks as the binary calculation method here is not very different from the decimal method you learned when you were young. It might only look a little bit strange. Also notations differ between countries, but the method is similar.

```
                    1 0 1 = 5
                  -------------
    1 0 0 1 1 / 1 1 0 1 1 0 1
                  1 0 0 1 1 | |
                  --------- | |
                    1 0 0 0 0 |
                    0 0 0 0 0 |
                  --------- |
                    1 0 0 0 0 1
                      1 0 0 1 1
                    ---------
                    1 1 1 0 = 14 = remainder
```

With decimal calculations you can quickly check that 109 divided by 19 gives a quotient of 5 with 14 as the remainder. But what we also see in the scheme is that every bit extra to check only costs one binary comparison and in 50% of the cases one binary subtraction. You can easily increase the number of bits of the test data string—for example to 56 bits if we use our example value "*Lammert*"—and the result can be calculated with 56 binary comparisons and an average of 28 binary subtractions. This can be implemented in hardware directly with only very few transistors involved. Also software algorithms can be very efficient.

All of the CRC formulas you will encounter are simply checksum algorithms based on modulo-2 binary division where we ignore carry bits and in effect the subtraction will be equal to an *exclusive or* operation. Though some differences exist in the specifics across different CRC formulas, the basic mathematical process is always the same:

- The message bits are appended with *c* zero bits; this *augmented message* is the dividend
- A predetermined *c+1*-bit binary sequence, called the *generator polynomial*, is the divisor

- The checksum is the *c*-bit remainder that results from the division operation

Table 1 lists some of the most commonly used generator polynomials for 16- and 32-bit CRCs. Remember that the width of the divisor is always one bit wider than the remainder. So, for example, you'd use a 17-bit generator polynomial whenever a 16-bit checksum is required.

	CRC-CCITT	**CRC-16**	**CRC-32**
Checksum Width	16 bits	16 bits	32 bits
Generator Polynomial	10001000000100001	11000000000000101	100000100110000010001110110110111

International Standard CRC Polynomials

Source Code:

```java
import java.io.*;
class Crc
{
    public static void main(String args[]) throws IOException
    {
        BufferedReader br=new BufferedReader(new
InputStreamReader(System.in)); int[ ] data;
        int[  ]div; int[ ]divisor; int[ ]rem;
        int[ ] crc;
        int data_bits, divisor_bits, tot_length;
        System.out.println("Enter number of data bits : ");
data_bits=Integer.parseInt(br.readLine()); data=new int[data_bits];
        System.out.println("Enter data bits : "); for(int i=0;
i<data_bits; i++)
        data[i]=Integer.parseInt(br.readLine());
System.out.println("Enter number of bits in divisor : ");
            divisor_bits=Integer.parseInt(br.readLine());
divisor=new int[divisor_bits];
        System.out.println("Enter Divisor bits : "); for(int
i=0; i<divisor_bits; i++)
        divisor[i]=Integer.parseInt(br.readLine());
        tot_length=data_bits+divisor_bits-1; div=new
int[tot_length];
        rem=new int[tot_length]; crc=new int[tot_length];
        /*----------------- CRC GENERATION--------------------
--*/
        for(int i=0;i<data.length;i++) div[i]=data[i];
            System.out.print("Dividend (after appending 0's)
are : ");
        for(int i=0; i< div.length; i++)
            System.out.print(div[i]);
        System.out.println();
        for(int j=0; j<div.length; j++)
        {
            rem[j] = div[j];
        }
        rem=divide(div, divisor, rem);
```
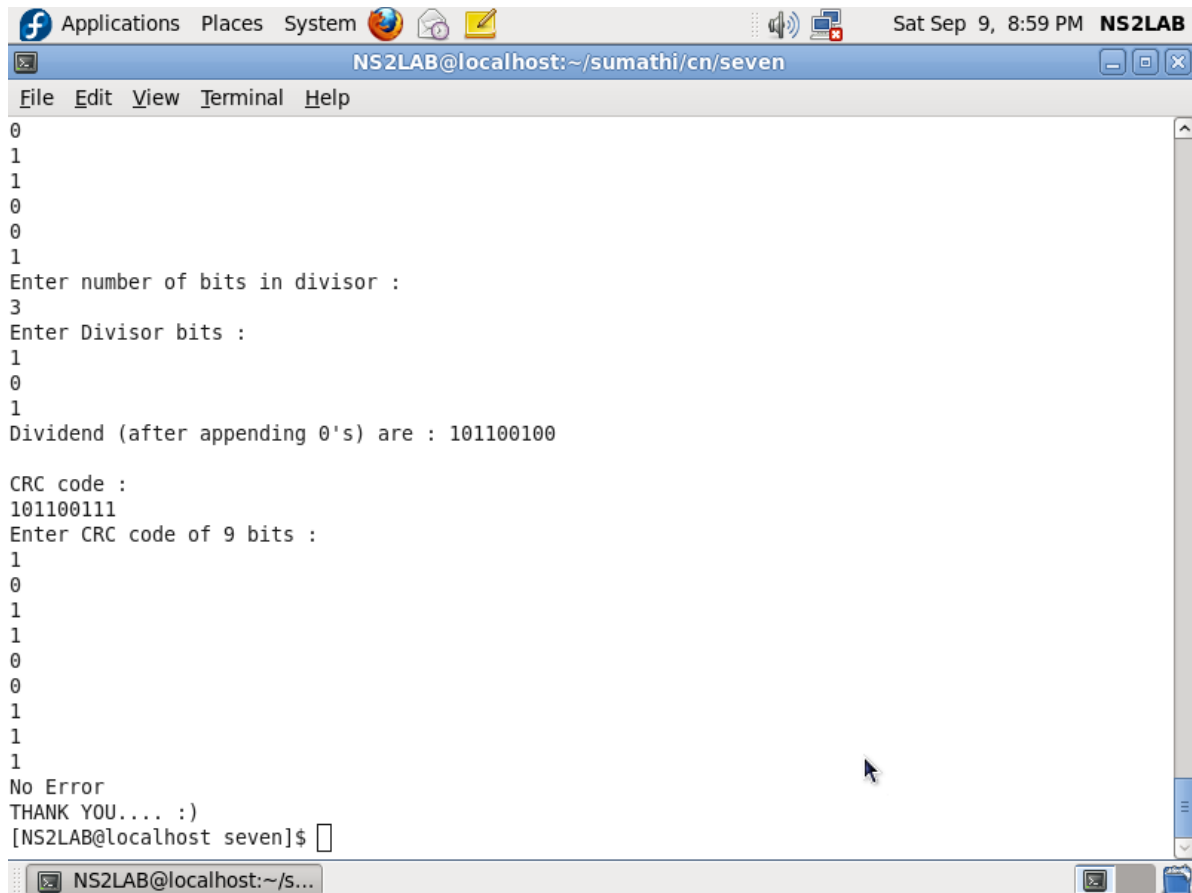
```
        for(int i=0;i<div.length;i++)     //append dividend and
ramainder
        {
            crc[i]=(div[i]^rem[i]);
        }
        System.out.println(); System.out.println("CRC code : ");
        for(int i=0;i<crc.length;i++)
        System.out.print(crc[i]);
        /*------------------ERROR DETECTION-------------------
-*/
        System.out.println();
        System.out.println("Enter CRC code of "+tot_length+"
bits : "); for(int i=0; i<crc.length; i++)
        crc[i]=Integer.parseInt(br.readLine());
        for(int j=0; j<crc.length; j++)
        {
            rem[j] = crc[j];
        }
        rem=divide(crc, divisor, rem); for(int i=0; i<
rem.length; i++)
        {
            if(rem[i]!=0)
            {
                System.out.println("Error"); break;
            }
            if(i==rem.length-1) System.out.println("No Error");
        }
        System.out.println("THANK YOU.... :)");
    }
    static int[] divide(int div[],int divisor[], int rem[])
    {
        int cur=0; while(true)
        {
            for(int i=0;i<divisor.length;i++)
rem[cur+i]=(rem[cur+i]^divisor[i]);
            while(rem[cur]==0 && cur!=rem.length-1) cur++;
            if((rem.length-cur)<divisor.length) break;
        }
        return rem;
    }
}
```

Output:
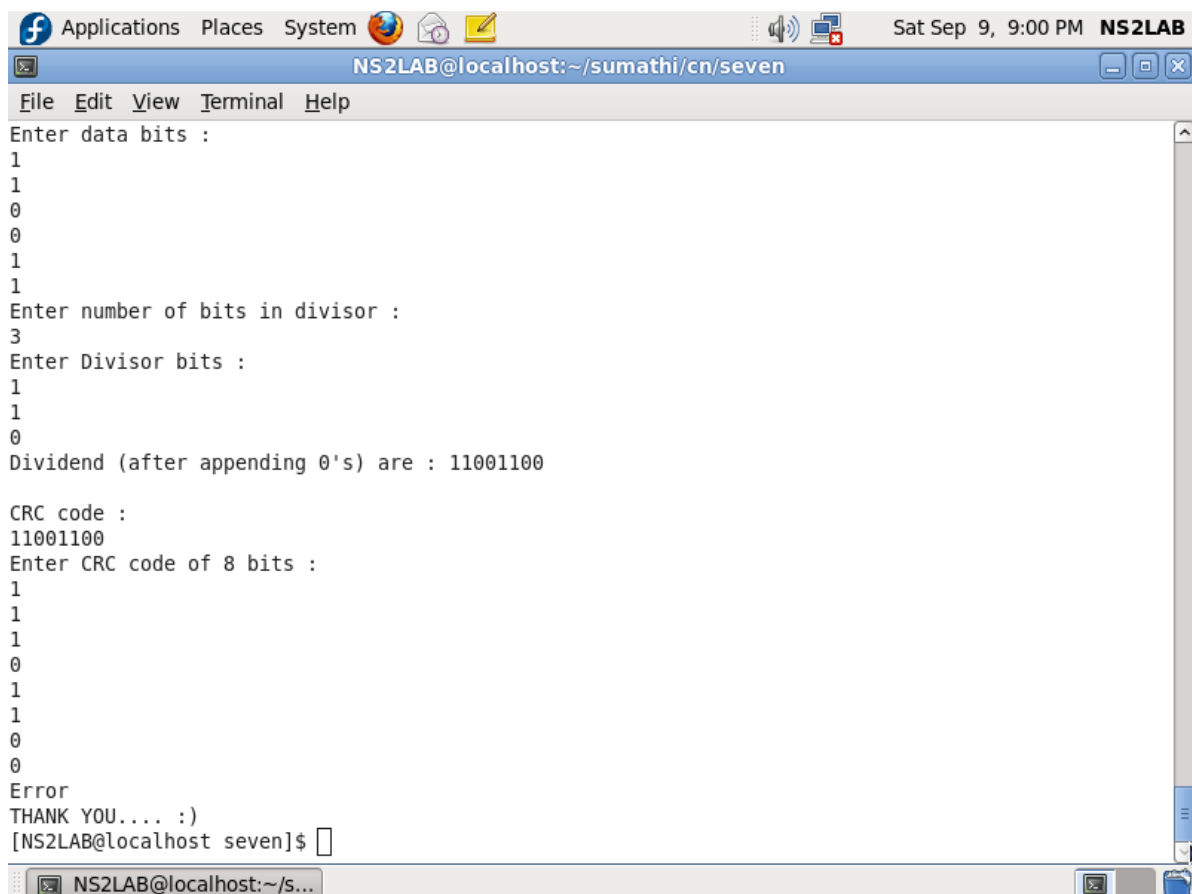
 [root@localhost ~]# vi Crc.java

```
0
1
1
0
0
1
Enter number of bits in divisor :
3
Enter Divisor bits :
1
0
1
Dividend (after appending 0's) are : 101100100

CRC code :
101100111
Enter CRC code of 9 bits :
1
0
1
1
0
0
1
1
1
No Error
THANK YOU.... :)
[NS2LAB@localhost seven]$
```

```
Enter data bits :
1
1
0
0
1
1
Enter number of bits in divisor :
3
Enter Divisor bits :
1
1
0
Dividend (after appending 0's) are : 11001100

CRC code :
11001100
Enter CRC code of 8 bits :
1
1
1
0
1
1
0
0
Error
THANK YOU.... :)
[NS2LAB@localhost seven]$
```

8. Write a program to find the shortest path between vertices using bellman-ford algorithm.

Distance Vector Algorithm is a decentralized routing algorithm that requires that each router simply inform its neighbors of its routing table. For each network path, the receiving routers pick the neighbor advertising the lowest cost, then add this entry into its routing table for re-advertisement. To find the shortest path, Distance Vector Algorithm is based on one of two basic algorithms: the Bellman-Ford and the Dijkstra algorithms.

Routers that use this algorithm have to maintain the distance tables (which is a one- dimension array -- "a vector"), which tell the distances and shortest path to sending packets to each node in the network. The information in the distance table is always up date by exchanging information with the neighboring nodes. The number of data in the table equals to that of all nodes in networks (excluded itself). The columns of table represent the directly attached neighbors whereas the rows represent all destinations in the network. Each data contains the path for sending packets to each destination in the network and distance/or time to transmit on that path (we call this as "cost"). The measurements in this algorithm are the number of hops, latency, the number of outgoing packets, etc.

The Bellman–Ford algorithm is an algorithm that computes shortest paths from a single source vertex to all of the other vertices in a weighted digraph. It is slower than Dijkstra's algorithm for the same problem, but more versatile, as it is capable of handling graphs in which some of the edge weights are negative numbers. Negative edge weights are found in various applications of graphs, hence the usefulness of this algorithm. If a graph contains a "negative cycle" (i.e. a cycle whose edges sum to a negative value) that is reachable from the source, then there is no cheapest path: any path that has a point on the negative cycle can be made cheaper by one more walk around the negative cycle. In such a case, the Bellman–Ford algorithm can detect negative cycles and report their existence

Source code:

```
import java.util.Scanner;

public class BellmanFord
{
        private int D[]; private int num_ver;
        public static final int MAX_VALUE = 999;
        public BellmanFord(int num_ver)
        {
                this.num_ver = num_ver; D = new int[num_ver + 1];
        }

        public void BellmanFordEvaluation(int source, int A[][])
        {
                for (int node = 1; node <= num_ver; node++)
                {
                        D[node] = MAX_VALUE;
                }
                D[source] = 0;
                for (int node = 1; node <= num_ver - 1; node++)
                {
                        for (int sn = 1; sn <= num_ver; sn++)
                        {
```
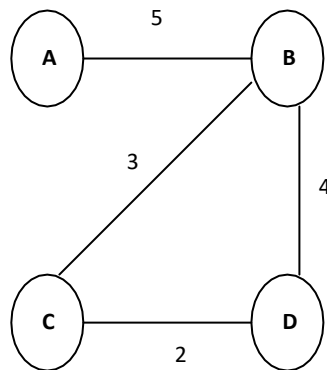
```java
                        for (int dn = 1; dn <= num_ver; dn++)
                        {
                                if (A[sn][dn] != MAX_VALUE)
                                {
                                        if (D[dn] > D[sn]+ A[sn][dn])
                                                D[dn] = D[sn] + A[sn][dn];
                                }
                        }
                }
        }
        for (int sn = 1; sn <= num_ver; sn++)
        {
                for (int dn = 1; dn <= num_ver; dn++)
                {
                        if (A[sn][dn] != MAX_VALUE)
                        {
                                if (D[dn] > D[sn]+ A[sn][dn])
                                        System.out.println("The Graph
contains negative egde cycle");
                        }
                }
        }
        for (int vertex = 1; vertex <= num_ver; vertex++)
        {
                System.out.println("distance of source " + source +
" to "+ vertex + " is " + D[vertex]);
        }
    }
    public static void main(String[ ] args)
    {
        int num_ver = 0; int source;
        Scanner scanner = new Scanner(System.in);
System.out.println("Enter the number of vertices");
        num_ver = scanner.nextInt();
        int A[][] = new int[num_ver + 1][num_ver + 1];
System.out.println("Enter the adjacency matrix");
        for (int sn = 1; sn <=num_ver; sn++)
        {
                for (int dn = 1; dn <= num_ver; dn++)
                {
                        A[sn][dn] = scanner.nextInt(); if (sn == dn)
                        {
                                A[sn][dn] = 0; continue;
                        }
                        if (A[sn][dn] == 0)
                        {
                                A[sn][dn] = MAX_VALUE;
                        }
                }
        }
        System.out.println("Enter the source vertex"); source =
scanner.nextInt();
        BellmanFord b = new BellmanFord (num_ver);
b.BellmanFordEvaluation(source, A); scanner.close();
    }
}
```

Input graph:

Output:

[root@localhost ~]# vi BellmanFord.java

9. Using TCP/IP sockets, write a client-server program to make client sending the file name and the server to send back the contents of the requested file if present. Implement the above program using as message queues or FIFOs as IPC channels.

Socket is an interface which enables the client and the server to communicate and pass on information from one another. Sockets provide the communication mechanism between two computers using TCP. A client program creates a socket on its end of the communication and attempts to connect that socket to a server. When the connection is

made, the server creates a socket object on its end of the communication. The client and the server can now communicate by writing to and reading from the socket.

Source Code:

TCP Client

```java
import java.io.*;
import java.net.*;
import java.util.Scanner;

public class TCPClient
{
      public static void main(String[] args) throws
IOException,InterruptedException
      {
            DataOutputStream out;
            DataInputStream in;
            Scanner scanner = new Scanner(System.in);

            Socket socket = new Socket("127.0.0.1", 6000); //server
IP and Port num

            System.out.println("Client Connected to Server");
            System.out.print("\nEnter the filename to request\n");
            String filename = scanner.nextLine();

            in = new DataInputStream(socket.getInputStream());
            out = new DataOutputStream(socket.getOutputStream());

            out.writeUTF(filename);

            String fileContent = in.readUTF();

            if (fileContent.length() > 0)
                  System.out.println(fileContent);
            else
                  System.out.println("FILE IS EMPTY");
      }
}
```

TCP Server

```java
import java.io.*;
import java.net.*;
import java.nio.file.Path;
import java.nio.file.Paths;
import java.nio.*;
import java.nio.file.Files;

public class TCPServer
{

      public static void main(String[] args) throws IOException
      {
```

```
        ServerSocket server;
        DataOutputStream out = null;
        DataInputStream in;

        try
        {
            server = new ServerSocket(6000, 1); //port number
and num of connections

            System.out.println("Server Waiting for client");
            Socket socket = server.accept();

            System.out.println("Client connected ");

            in = new DataInputStream(socket.getInputStream());
            out = new
DataOutputStream(socket.getOutputStream());

            String fileName = in.readUTF();
            System.out.println("File Requested is : " +
fileName);

                  //  Path p=Paths.get("fileName");
                   //byte[] filedata=File.readAllBytes(p);
            byte[] filedata =
Files.readAllBytes(Paths.get(fileName));
            String fileContent = new String(filedata);

            out.writeUTF(fileContent.toString());
            System.out.println("FILE SENT SUCCESSFULLY");
        }
        catch (Exception e)
        {
            System.out.println(e.getMessage());
            out.writeUTF("FILE DOESN'T EXISTS");
        }
    }
}
```
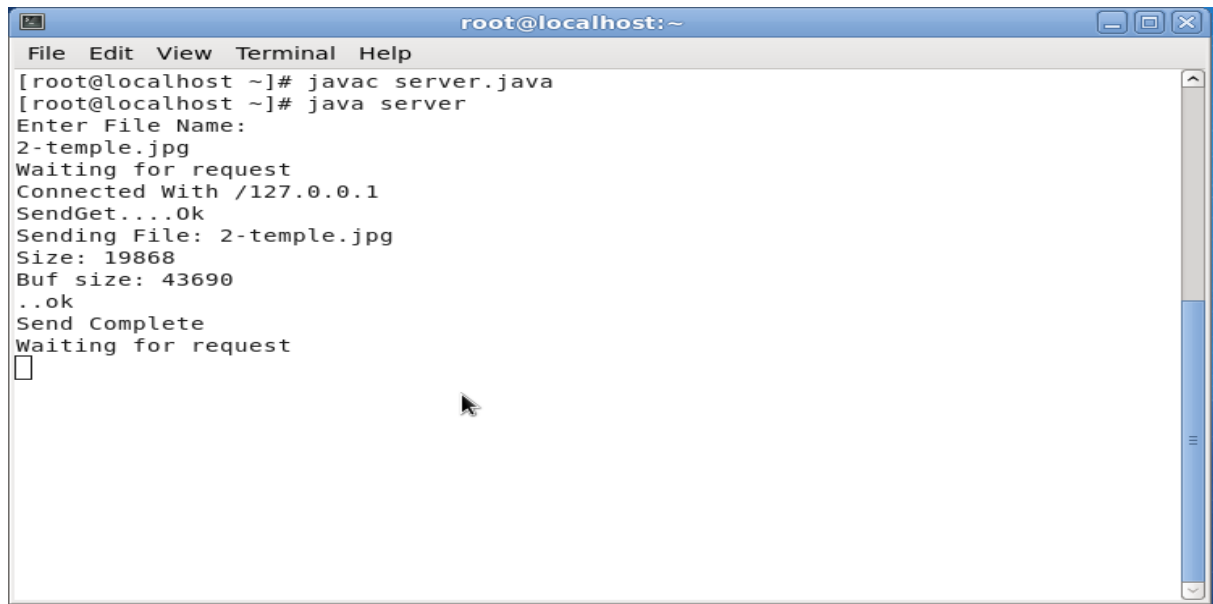
Note: Create two different files Client.java and Server.java. Follow the steps given:
1. Open a terminal run the server program and provide the filename to send
2. Open one more terminal run the client program and provide the IP address of the server. We can give localhost address "127.0.0.1" as it is running on same machine or give the IP address of the machine.
3. Send any start bit to start sending file.
4. Refer https://www.tutorialspoint.com/java/java_networking.htm for all the parameters, methods description in socket communication.
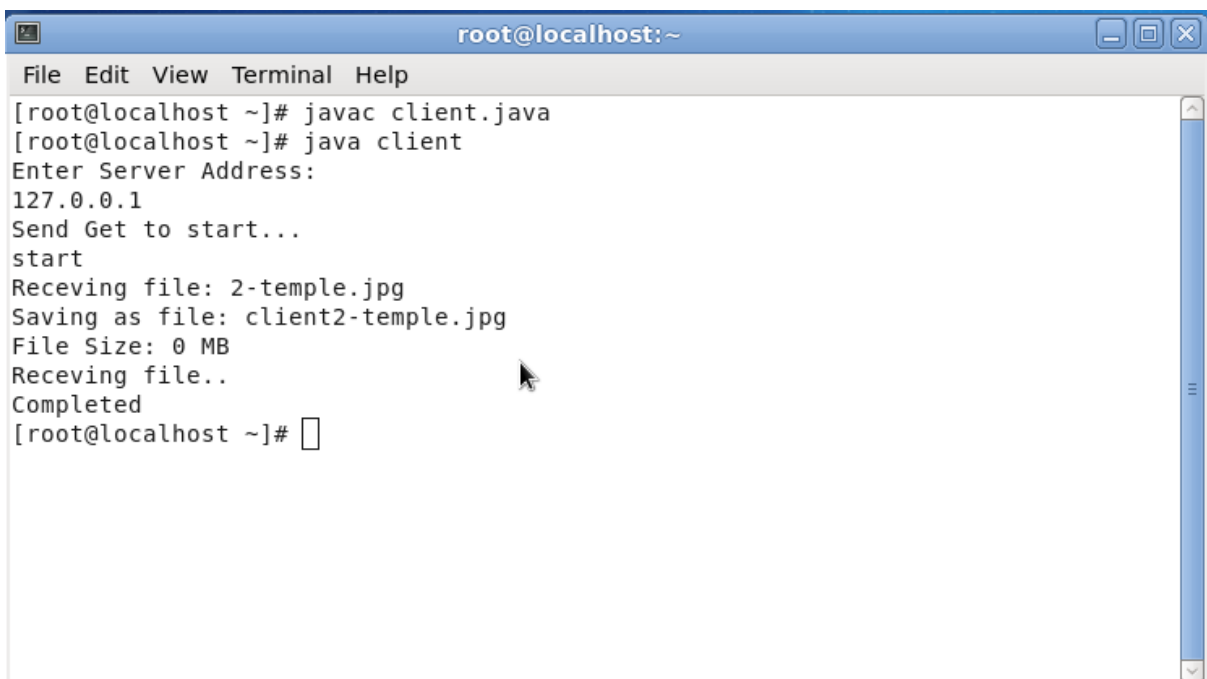

Output:

At server side:

At client side:

10. Write a program on datagram socket for client/server to display the messages on client side, typed the server side.

A datagram socket is the one for sending or receiving point for a packet delivery service. Each packet sent or received on a datagram socket is individually addressed and routed. Multiple packets sent from one machine to another may be routed differently, and may arrive in any order.

Source Code:

UDP Client

```java
import java.net.*;
import java.util.Scanner;
public class DSender
{
     public static void main(String[] args) throws Exception
     {
          System.out.println("Sender");
          DatagramSocket ds = new DatagramSocket();
          Scanner scanner = new Scanner(System.in);

          System.out.println("\nEnter the Message : ");

          while(true)
          {
               String msg = scanner.nextLine();
               InetAddress ip =
InetAddress.getByName("127.0.0.1");

               DatagramPacket dp = new
DatagramPacket(msg.getBytes(), msg.length(), ip, 3000);
               ds.send(dp);
          }
     }
}
```

UDP Server

```java
import java.net.*;
public class DReciever
{
     public static void main(String[] args) throws Exception
     {
          byte[] buf = new byte[1024];
          System.out.println("Receiver");
          DatagramSocket ds = new DatagramSocket(3000);

          while(true)
          {
               DatagramPacket dp = new DatagramPacket(buf, 1024);
               ds.receive(dp);
               String msg = new String(dp.getData(), 0,
dp.getLength());
               System.out.println(msg);
          }
     }
}
```
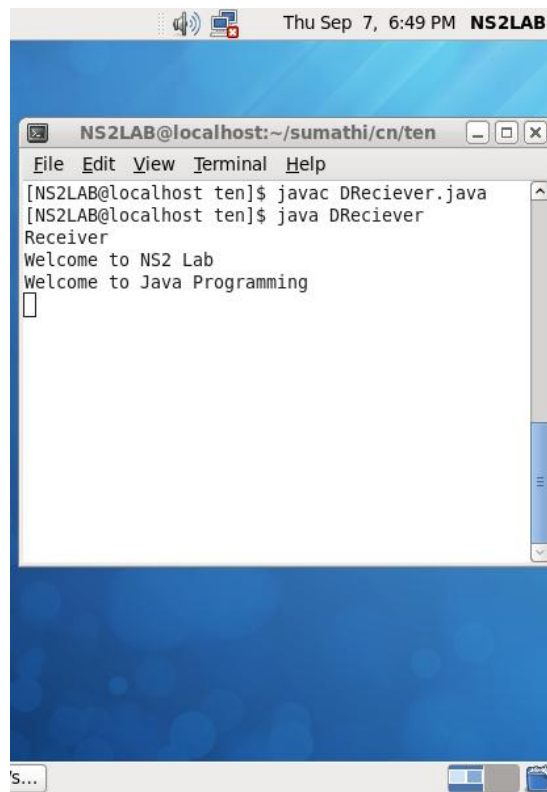
Note: Create two different files UDPC.java and UDPS.java. Follow the following steps:
1. Open a terminal run the server program.

2. Open one more terminal run the client program, the sent message will be received.
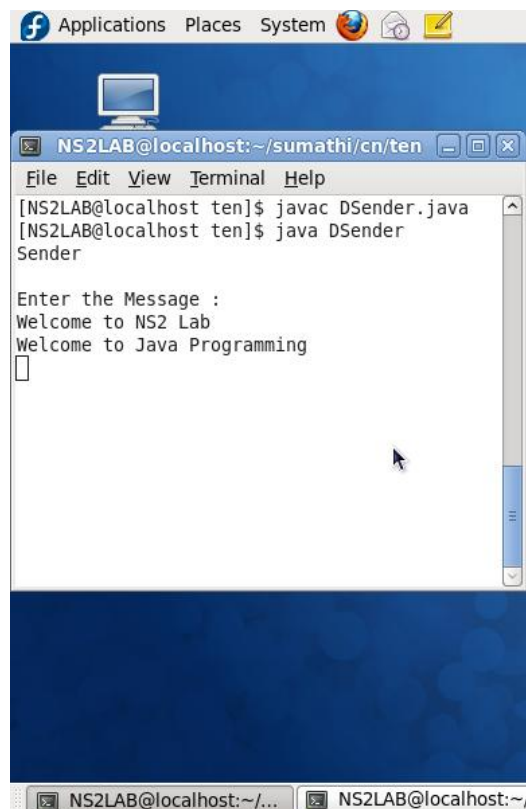
At Server side

At Client side

11. Write a program for simple RSA algorithm to encrypt and decrypt the data.

RSA is an example of public key cryptography. It was developed by Rivest, Shamir and Adelman. The RSA algorithm can be used for both public key encryption and digital signatures. Its security is based on the difficulty of factoring large integers.

The RSA algorithm's efficiency requires a fast method for performing the modular exponentiation operation. A less efficient, conventional method includes raising a number (the input) to a power (the secret or public key of the algorithm, denoted *e* and *d*, respectively) and taking the remainder of the division with *N*. A straight-forward implementation performs these two steps of the operation sequentially: first, raise it to the power and second, apply modulo. The RSA algorithm comprises of three steps, which are depicted below:

Key Generation Algorithm

1. Generate two large random primes, p and q, of approximately equal size such that their product n = p*q
2. Compute n = p*q and Euler's totient function (φ) phi(n) = (p-1)(q-1).
3. Choose an integer e, 1 < e < phi, such that gcd(e, phi) = 1.
4. Compute the secret exponent d, 1 < d < phi,
   such that e*d ≡ 1 (mod phi).
5. The public key is (e, n) and the private key is (d, n). The values of p, q, and phi should also be kept secret.

Encryption

Sender A does the following:-

1. Using the public key (e,n)
2. Represents the plaintext message as a positive integer M
3. Computes the cipher text C = M^e mod n.
4. Sends the cipher text C to B (Receiver).

Decryption

Recipient B does the following:-

1. Uses his private key (d, n) to compute M = C^d mod n.
2. Extracts the plaintext from the integer representative m.

Source Code:

```java
import java.util.*;
import java.math.*;
import java.io.*;
    public class RSAnew {
    public BigInteger p,q,n,phi,e,d;
        public Random r;
```

```java
public  RSAnew()
{
r = new Random();
        p=BigInteger.probablePrime(1024, r);
                q=BigInteger.probablePrime(1024, r);
                n = p.multiply(q);
phi = p.subtract(BigInteger.ONE) .multiply (q.subtract(BigInteger.ONE));
e= BigInteger.probablePrime(512,r);

while(phi.gcd(e).compareTo(BigInteger.ONE) > 0 && e.compareTo(phi) < 0)
                {
                                e.add(BigInteger.ONE);
                }
        d=e.modInverse(phi);
}
public static void main(String[] args) throws IOException {
        RSAnew rs=new RSAnew();
        DataInputStream in= new DataInputStream(System.in);
        String s;
        System.out.println("Enter the string:");
        s = in.readLine();
        System.out.println("String ecrypting is :"+s);
System.out.println("String on byte is "+stringToBytes(s.getBytes()));
        byte[] encd = rs.enc(s.getBytes());
        byte[] decd = rs.dec(encd);
        System.out.println("Decrypted bytes :"+ stringToBytes(decd));
        System.out.println("Decrypted string is:"+new String (decd));
        in.close();
        }
private static String stringToBytes(byte [] m)
{
        String t="";
        for ( byte b:m )
        {
                t=t+ Byte.toString(b);
        }
        return t;
}
```

```java
    public byte[] enc( byte [] m)
    {
    return( new BigInteger (m).modPow(e, n).toByteArray());
    }
    public byte[] dec( byte [] m)
    {
    return( new BigInteger (m).modPow(d, n).toByteArray());
    }
}
```

OUTPUT:

```
Enter the string:
HELLO 123 @#$%^&
String ecrypting is :HELLO 123 @#$%^&
String on byte is 72697676793249505132643536379543832
Decrypted bytes :72697676793249505132643536379543832
Decrypted string is:HELLO 123 @#$%^&
```

12. Write a program for congestion control using Leaky bucket algorithm.

The main concept of the leaky bucket algorithm is that the output data flow remains constant despite the variant input traffic, such as the water flow in a bucket with a small hole at the bottom. In case the bucket contains water (or packets) then the output flow follows a constant rate, while if the bucket is full any additional load will be lost because of spillover. In a similar way if the bucket is empty the output will be zero. From network perspective, leaky bucket consists of a finite queue (bucket) where all the incoming packets are stored in case there is space in the queue, otherwise the packets are discarded. In order to regulate the output flow, leaky bucket transmits one packet from the queue in a fixed time (e.g. at every clock tick). In the following figure we can notice the main rationale of leaky bucket algorithm, for both the two approaches (e.g. leaky bucket with water

(a) and with packets (b)).
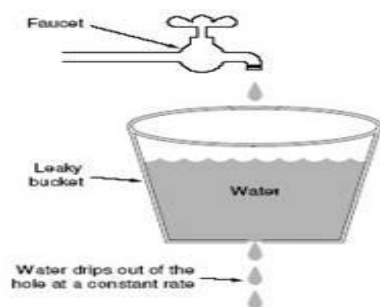
(a) A leaky bucket with water
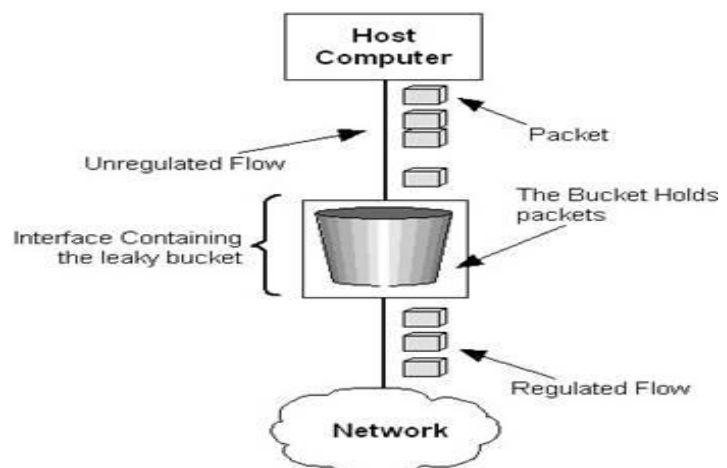
(b) A leaky bucket with packets

While leaky bucket eliminates completely bursty traffic by regulating the incoming data flow its main drawback is that it drops packets if the bucket is full. Also, it doesn't take into account the idle process of the sender which means that if the host doesn't transmit data for some time the bucket becomes empty without permitting the transmission of any packet.

Source Code:

```java
import java.util.*;
public class leaky4
{
    public static void main(String[] args)
    {
        Scanner my = new Scanner(System.in);
        int no_groups,bucket_size;
        System.out.print("\n Enter the bucket size : \t");
        bucket_size = my.nextInt();
        System.out.print("\n Enter the no of groups : \t");
        no_groups = my.nextInt();
        int no_packets[] = new int[no_groups];
        int in_bw[] = new int[no_groups];
        int out_bw,reqd_bw=0,tot_packets=0;
        for(int i=0;i<no_groups;i++)
        {
            System.out.print("\n Enter the no of packets for group " + (i+1) +
"\t");
            no_packets[i] = my.nextInt();
            System.out.print("\n Enter the input bandwidth for the group " +
(i+1) + "\t");
            in_bw[i] = my.nextInt();
            if((tot_packets+no_packets[i])<=bucket_size)
            {
                tot_packets += no_packets[i];
            }
            else
            {
                do
                {
                System.out.println(" Bucket Overflow ");
                System.out.println(" Enter value less than " + (bucket_size-
tot_packets));
                no_packets[i] = my.nextInt();
                }while((tot_packets+no_packets[i])>bucket_size);
                tot_packets += no_packets[i];
            }
            reqd_bw += (no_packets[i]*in_bw[i]);
        }
        System.out.println("\nThe total required bandwidth is " + reqd_bw);
        System.out.println("Enter the output bandwidth ");
        out_bw = my.nextInt();
        int temp=reqd_bw;
        int rem_pkts = tot_packets;
        while((out_bw<=temp)&&(rem_pkts>0))
        {
            System.out.println("Data Sent \n" + (--rem_pkts) + " packets
remaining");
            System.out.println("Remaining Bandwidth " + (temp -= out_bw));
            if((out_bw>temp)&&(rem_pkts>0))
```

```
            System.out.println(rem_pkts + " packet(s) discarded due to
insufficient bandwidth");
        }
    }
}
```

Output

```
Enter the bucket size :    5
 Enter the no of groups : 2
 Enter the no of packets for group 1    2
 Enter the input bandwidth for the group 1    40
 Enter the no of packets for group 2    2
 Enter the input bandwidth for the group 2    20
The total required bandwidth is 120
Enter the output bandwidth
30
Data Sent
3 packets remaining
Remaining Bandwidth 90
Data Sent
2 packets remaining
Remaining Bandwidth 60
Data Sent
1 packets remaining
Remaining Bandwidth 30
Data Sent
0 packets remaining
Remaining Bandwidth 0
```

REFERENCES

1. https://www.isi.edu/nsnam/ns/
2. http://aknetworks.webs.com/e-books
3. **Communication Networks: Fundamental Concepts and Key Architectures -** Alberto Leon, Garcia and Indra Widjaja, 4th Edition, Tata McGraw- Hill, reprint-2012.
4. **Data and Computer Communication,** William Stallings, 8th Edition, Pearson Education, 2009.
5. **Computer Networks: A Systems Approach -** Larry L. Peterson and Bruce S. David, 4th Edition, Elsevier, 2009.
6. **Introduction to Data Communications and Networking** – Wayne Tomasi, Pearson Education, 2009.
7. **Communication Networks – Fundamental Concepts and Key architectures –** Alberto Leon- Garcia and Indra Widjaja:, 2rd Edition, Tata McGraw-Hill, 2009
8. **Computer and Communication Networks –** Nader F. Mir:, Pearson Education, 2012

Viva Questions

1.What do you mean by data communication?
2.What is simplex?
3.What is half-duplex?
4.What is full duplex?
5.What is a network?
6.What is distributed processing?
7.What is point to point connection?
8.What is multipoint connection?
9.What is a topology?
10.Define LAN, MAN and WAN.
11.Define internet?
12.What is a protocol?
13.What is TCP/IP protocol model?
14.Describe the functions of five layers?
15.What is ISO-OSI model?
16. What is multiplexing? 16.What is switching?
17.How data is transmitted over a medium?
18. Compare analog and digital signals?
19.Define bandwidth?
20.What are the factors on which data rate depends?
21.Define bit rate and bit interval?
22.What is Nyquist bit rate formula?
23.Define Shannon capacity?
24.What is sampling?
25.Define pulse amplitude modulation?
26.Define pulse code modulation?
27.What is Nyquist Theorem?
28.What are the modes of data transmission?
29.What is Asynchronous mode of data transmission?
30.What is Synchronous mode of data transmission?
31.What arethe different types of multiplexing?
32.What is FDM and TDM ?
33.What are the different transmission media?
34.What are the different Guided Media?
35.Describe about the different Guided Medias.
36.What do you mean by wireless communication?
37.What do you mean by switching?
38.What are the switching methods?
39.What are the duties of data link layer?
40.What are the types of errors?
41.What do you mean by redundancy?
42.Define parity check.
43. Define cyclic redundancy check (CRC).
44. What is hamming code?
45.What do you mean by flow control?
46.What do you mean by error control?
47.Define stop and wait ARQ.
48.Define Go-Back-N ARQ?
49.Define Selective Repeat ARQ?

50. What do you mean by pipelining, is there any pipelining in error control?
51. What is HDLC?
52. What do you mean by point to point protocol?
53. What do you mean by point to point protocol stack?
54. What do you mean by line control protocol?
55. What do you mean by Authentication protocol?
56. What do you mean by network control protocol?
57. What do you mean by CSMA?
58. What do you mean by Bluetooth?
59. What is IP address?
60. What do you mean by subnetting?
61. What are the advantages of fiber optics cable ?
62. What are the disadvantages of fiber optics cable?
63. What are the propagation type of radio wave ?
64. What do you mean by Geosynchronous Satellites ?
65. What are the factors for evaluating the suitability of the media ?
66. What do you mean by medium access control(MAC) sublayer.
67. What do you mean by ALOHA ?
68. What is pure ALOHA?
69. What is slotted ALOHA?
70. What do you mean by persistent CSMA(carrier sense multiple access) ?
71. What do you mean by non persistent CSMA(carrier sense multiple access) ?
72. What do you mean by p persistent CSMA(carrier sense multiple access) ?
73. What is FDDI?
74. What is Firewalls?
75. What is Repeaters ?
76. What is Bridges?
77. What is Routers ?
78. What is Gateway ?
79. What do you mean by Data Terminal Equipment(DTE) ?.
80. What do you mean by Data Terminating Equipment (DCE) ?
81. What do you mean by protocol stack ?
82. What do you mean by peer ?
83. What do you mean by broadcasting ?
84. What are the advantages of broadcast network.
85. What do you mean by point to point network?
86. What are the design issue of layers ?
87. What are the protocols in application layer ?
88. What are the protocols in transport layer ?
89. Define TCP ?
90. Define UDP ?
91. Define IP ?
92. What do you mean by client server model ?
93. What are the information that a computer attached to a TCP/IP internet must possesses ?
94. What is domain name system(DNS)?
95. What is TELNET ?
96. What do you mean by local login and remote login ?
97. What is Network Virtual Terminal ?
98. What do you mean by Simple Mail Transfer Protocol ?
99. What is Hypertext Transfer Protocol(HTTP) ?

100.What is URL ?
101. What is World Wide Web ?
102.What is HTML ?