

Sorting

Today's checklist



- 1. Sorting
- 2. Bubble sort Algorithm
- Time complexity and space complexity
- 4. Bubble sort optimization
- 5. Stable and unstable sort
- **6.** Practice Questions
- 7. Selection sort Algorithm
- 8. Time complexity and space complexity
- 9. Insertion sort Algorithm
- 10. Time complexity and space complexity
- 11. Stability of both
- 12. Programming questions

What is sorting?

Increasing

Ascending

Non-decreasing



$$arr = \{10, 1, 2, 18, 4, 5\}$$

$$\{1, 2, 4, 5, 10, 18\}$$
Sorted Order
$$Any given array is said to be sorted$$

$$ficreating if [arr[i] <= arr[i+1]] for every i$$

$$Ascending where i = 0,1,2...n-2$$



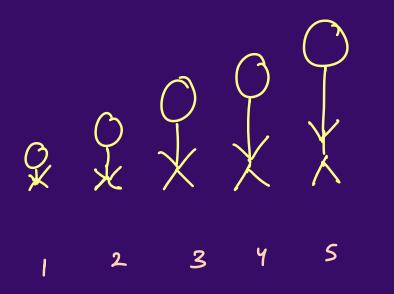
Que Check if given array is sooted array.



arr =
$$\{7, 1, 2, 8, -4\}$$

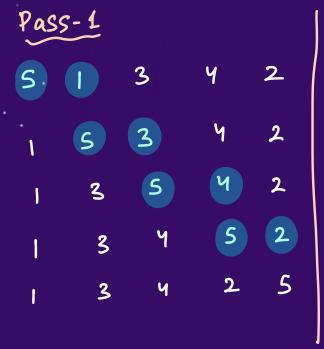
 $\int sort$
 $\{-4, 1, 2, 7, 8\}$

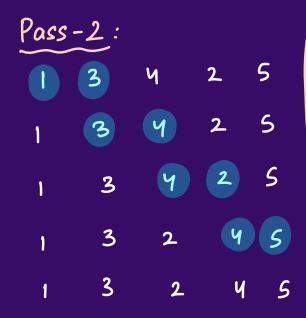




technique:
Swapping every 2
adjacent elements if
arr[i] > arr[i+1]











$$arr = \{10, 4, 1, 0, -2\}$$
 $n=5$

Pacs:1

















Pass: 2



















Pass:3

















Pass: 4

















the maximum posses will always be 4 (n-1)

Example



Time and Space complexity



```
// Bubble Sort - L
for(int x=1;x<=n-1;x++){ // n-1 passes
    for(int i=0;i< n-1;i++){
         if(arr[<u>i</u>]>arr[<u>i</u>+1]){
              int temp = arr[i];
              arr[i] = arr[i+1];
              arr[\underline{i}+1] = temp;
```

no of op! =
$$(n-1)^2 = n^2 - 2n + 1$$

S.C. =
$$O(n)$$
 -> Given array of 'n' size

Auxillary Space -> $O(1)$

T.C. = $O(n^2-2n+1)$
 $\simeq O(n^2)$

Time and Space complexity



```
// Bubble Sort - 🤽
for(int x=0;x<n-1;x++){ // n-1 passes
    for(int i=0;i< n-1-x;i++){
         if(arr[<u>i</u>]>arr[<u>i</u>+1]){
              int temp = arr[i];
              arr[i] = arr[i+1];
              arr[\underline{i}+1] = temp;
```

```
T.n. 0 =
   x=0 - n-1 times
            - n-2 times
   n-1+n-2+n-3...1
    = \frac{n^2}{2} - \frac{n}{2} \left[ \text{T.C.} = O(n^2) \right]
```

$$3 + 2 + 3 \cdot \cdot \cdot n - 1 = \underbrace{\frac{n(n-1)}{2}}_{2} = \frac{n^{2}}{2} - \frac{n}{2} \left[7 \cdot C \cdot = O(n^{2}) \right]$$

Can we optimize it further?



yes. We can

For arrays like 1 arr = {3,1,2,5,43 = This can be sorted in one bass only

1) So, after every bass, we can check if that array is sorted or not

Can we optimize it further?



```
// Bubble Sort Optimised
for(int \underline{x}=0;\underline{x}< n-1;\underline{x}++){ // n-1 passes
     boolean flag = true;
     for(int i=0;i< n-1-x;i++) {
          if (arr[i] > arr[i + 1]) {
               int temp = arr[i];
               arr[\underline{i}] = arr[\underline{i} + 1];
               arr[\underline{i} + 1] = temp;
               flag = false;
     if(flag==true) break;
```

Stable and Unstable sort

SKILLS

Bubble Sort

worst case



Q1: How much maximum swaps are needed to sort array of length 6?

Ex: arr =

5+4+3+2+1 = 15

n-1 + n-2 + n-3 ... 2

 $=\frac{n(n-1)}{2}$ max swaps

4 3 2 1 6

Pass:1 -> 5 swaps



Q2: Push zeroes to end while maintaining the relative order of other elements. $a_{YY} = \begin{cases} 0 & 1 & 0 \\ 3 & 12 \end{cases}$

- 0 1 0 3 12
 - 3 12
 - 0 0 3 12
 - 1 0 3 0 12
 - 0 3 12

Pass- 2

- 1 0 3 12 0
- 1 6 3 12 0
- 1 3 0 12 0
- 3 12 0

Selection Sort Algorithm



In every pass, we find the Kth smallest element k put it in its right place

 $arr = \begin{bmatrix} 3 & 5 & 2 & 1 & 4 \\ 1 & 5 & 2 & 3 & 4 \end{bmatrix} pass 1$ 1 & 2 & 3 & 4 & pass 3 1 & 2 & 3 & 4 & pass 4 1 & 2 & 3 & 4 & 5

Selection Sort Algorithm



'n-1' passes -4 20 1 8 -6 pass - D pass-1 1 & 10 20 ID pass-3 pass-4

Selection sort Code and dry run



```
// Selection Sort
for(int <u>i</u>=0;<u>i</u><n-1;<u>i</u>++){
     int min = Integer.MAX_VALUE;
     int mindx = -1;
     for(int i=i;i<n;i++){</pre>
          if(arr[j]<min){</pre>
               \min = \operatorname{arr}[j];
               mindx = i
     swap(arr, i, mindx);
```

```
public static void swap(int[] arr, int i, int j){
   int temp = arr[i];
   arr[i] = arr[j];
   arr[j] = temp;
}
```

Time and Space complexity



```
// Selection Sort
for(int i=0;i< n-1;i++){
    int min = Integer.MAX_VALUE;
    int mindx = -1;
    for(int j=<u>i;j</u><n;j++){
         if(arr[j]<min){</pre>
             min = arr[j];
             mindx = i;
    swap(arr,i,mindx);
```

```
i = 0 \rightarrow j = 0,1,2,3...n-1 \rightarrow n \text{ finel}
         1 -> j= 1,2,3...n-1 -> n-1 times
        2 -> j=2,3,...n-1 -n-2 times
        n-2
No of operations: n+n-1+n-2+n-3 ... 1
               = \frac{n(n+1)}{2} = \frac{n^2}{2} + \frac{n}{2}
```

$$T.C. = O(n^2)$$

Auxillary Space = 0(1)

Stable and Unstable sort



Time and Space complexity



Selection Sort

Cannot be optimised

Can be optimised

Unstable Sort

No. of swoops are less

No. of swoops

I

n-1 enops

Can be estimised

Stable Sort

No. of swops are more

1

1

1(n-1) max swops

9f cost of swofping is something to consider then selection cost is better.

Best case T.C.

is O(n)is $O(n^2)$

Time and Space complexity



	Time Conflexity	
	<u>B-S</u>	2 -1
Aug. Case	0 (n²)	0(n²)
Worst · Case	o (n²)	0(12)
Best Case	O(n)	0(n ²)

Homework:



Sort a given array in decreasing order using bubble sort

$$arr = \{3, 1, 2, 5, 43 \rightarrow \{5, 4, 3, 2, 1\}$$

Hint: After every pass the smallest element will be at and.

iffarr[i] < arr[i+1]) &was

Homework:



Sort a given Array in increasing order using selection sort, but in each pass, put the kth maximum element at the right position.

Insertion Sort Algorithm

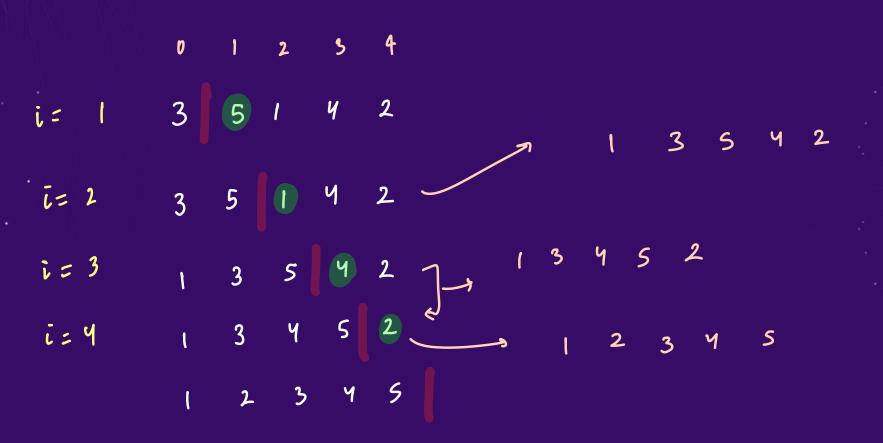


Prantle: you have playing cards numbered from 1 to 10 in random order, you have to sort them.

123 456 78910

Insertion Sort Algorithm





Insertion Sort Algorithm



```
n-1 passes - no of times outer loop will run.
                       i=1; i<n
   i=0; i2n-1
Grover Loop - we stout from
                            if (arr [j] < arr [j-1])
                                Swap (arr [i], arr [i-1])
                     arr[0] < arr[-1]
```

Insertion sort Code and dry run

```
SKILLS
```

```
// Insertion Sort
for(int i=1;i<n;i++){ // n-1 passes
    for(int j=i;j>=1;j--){
        if(arr[j]<arr[j-1])
            swap(arr,j,j-1);
        else break;
    }
}</pre>
```

```
Worst come:

0 1 2 3 1

1 2 3 4 5

j-1 j
```

No. of swaps =
$$1 + 2 + 3 + 4$$

 $1 + 2 + 3 + \cdots + n - 1 = \frac{n(n-1)}{2}$
 $\frac{2}{max}$
Swaps

Time and Space complexity



Avg. Case =
$$b(n^2)$$

Stability of Insertion and Selection & Skills Sort

Insertion Sort is a stable sort.

Bubble Sort

Selection Sort

Instion Sort

• O(n) T·C· in best case

• O(n²) T.C.in best case • O(n) T.C. in best case

· Stable

• Stable

· Unstable

. n (n-1) max swaps

. n(n-1) max swaps

. n-1 swoops

. Always optimised

- you have to obtainable extra variable (boolean)

. Cormot Oftimise



Q4: What will the array look like after the first iteration of selection sort [2,3,1,6,4]

- a) [1,2,3,6,4]
- b) [1,3,2,4,6]
- c) [1,3,2,6,4]
- d) [2,3,1,4,6]





Q5: Which sorting technique is used here?

A player is sorting a deck of cards numbered from 1 to 52. She first picks one card then picks the next card and puts it after the first card if it is bigger or before the first card if it is smaller, then she picks another card and puts it into its proper position.

- a) Bubble sort
- **b)** Insertion sort
- c) Selection sort
- d) None of these



Q6: Which of the following is not a stable sorting algorithm?

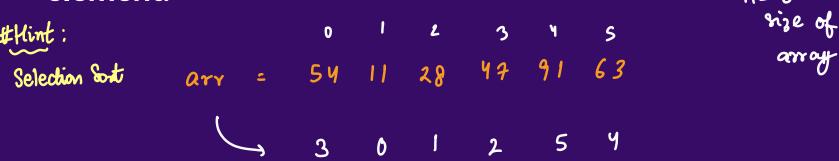
- a) Insertion sort
- め) Selection sort
- c) Bubble sort
- d) None of these



Q7: Majority Element

1 L all are + ve

Q8: Given an array with N distinct elements, convert the given array to a form where all elements are in the range from 0 to N-1. The order of elements is the same, i.e., 0 is placed in the place of the smallest element, 1 is placed for the second smallest element, ... N-1 is placed for the largest element.





THANKYOU