



Merge Sort & Inversion Count

Today's checklist

1. Merge 2 Sorted Arrays → *Already Done*
2. Merge Sort Algorithm
3. Merge Sort Time and Space Complexity
4. Stability
5. Applications of ^{Merge}~~Quick~~ Sort
6. Count Inversion Problem

Merge 2 Sorted Arrays (VV Important)

a

0	1	2	3	4
10	30	50	60	80

i c

0	1	2	3	4	5	6	7	8
10	20	30	40	50	60	70	75	80

b

0	1	2	3
20	40	70	75

j

k

$$T.C. = O(m+n)$$

Merge Sort Algorithm → Using Magic

0	1	2	3	4	5	6	7
80	10	70	30	60	40	50	20

0	1	2	3
80	10	70	30

↓ Magic

10	30	70	80
----	----	----	----

0	1	2	3
60	40	50	20

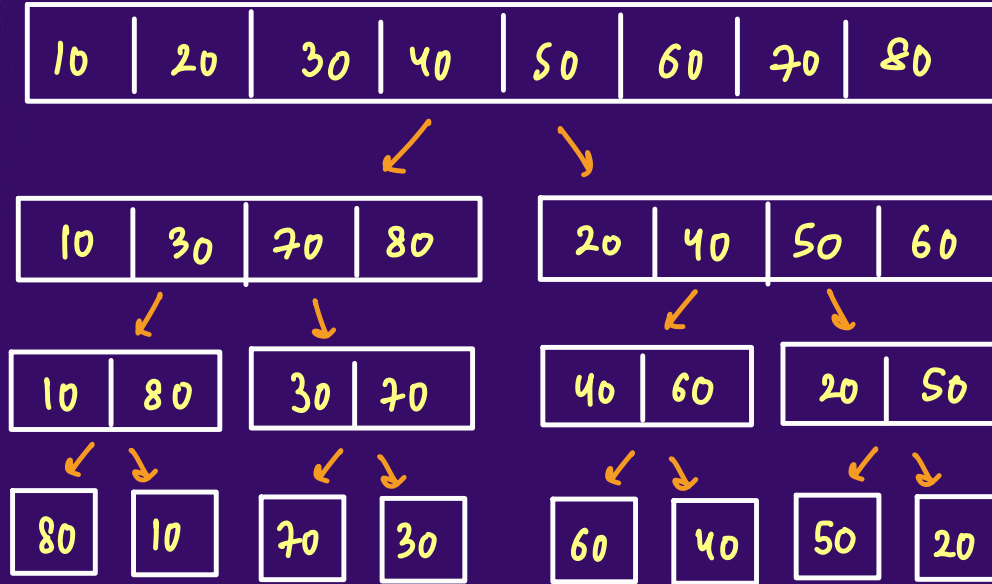
↓ Magic

20	40	50	60
----	----	----	----

↖ merge ↗
↓

10	20	30	40	50	60	70	80
----	----	----	----	----	----	----	----

Merge Sort Algorithm



Divide & Conquer

T.n 0 = *copy pasting*

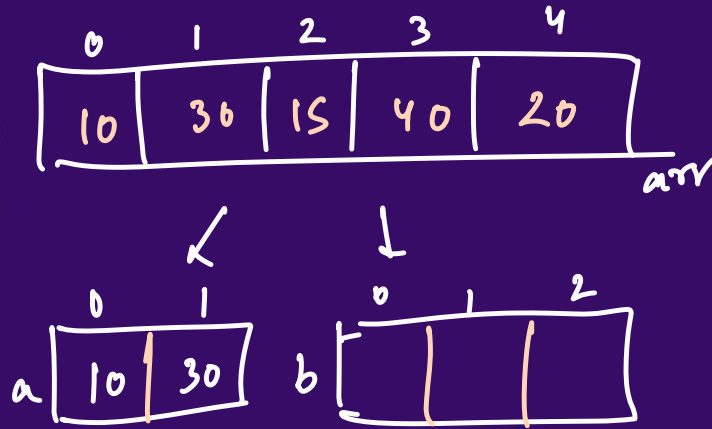
$$\left[\begin{aligned} &n + \left(\frac{n}{2} + \frac{n}{2}\right) \\ &+ \left(\frac{n}{4} + \frac{n}{4} + \frac{n}{4} + \frac{n}{4}\right) \\ &+ \left(\frac{n}{8} + \frac{n}{8} + \frac{n}{8} + \frac{n}{8} + \frac{n}{8} + \frac{n}{8}\right) \\ &+ \left(\frac{n}{16} + \frac{n}{16}\right) \\ &+ n \end{aligned} \right]$$

merging

$$\Rightarrow 2 [n + n + n]$$

$$\Rightarrow 2(Q-1) \cdot n$$

Merge Sort Algorithm



$$b[i] = arr[i + n/2]$$

```

public static void mergesort(int[] arr){
    int n = arr.length;
    if(n==1) return; // base case
    // create two arrays of n/2 size each
    int[] a = new int[n/2];
    int[] b = new int[n-n/2];
    // copy pasting
    for(int i=0;i<n/2;i++){
        a[i] = arr[i];
    }
    for(int i=0;i<n-n/2;i++){
        b[i] = arr[i+n/2];
    }
    // magic
    mergesort(a);
    mergesort(b);
    // merge these 'a' and 'b'
    merge(a,b,arr);
}

```

hm

```

public static void merge(int[] a, int[] b, int[] c){
    int i = 0, j = 0, k = 0;
    while(i<a.length && j<b.length){
        if(a[i]<=b[j]) c[k++] = a[i++];
        else c[k++] = b[j++];
    }
    while(j<b.length) c[k++] = b[j++];
    while(i<a.length) c[k++] = a[i++];
}

```

→ S.C. = $O(n \log n)$

Time and Space complexity

\downarrow
 n
 $n/2$
 $n/8$
 $n/16$
 \vdots
 1

$n \quad \frac{n}{2} \quad \frac{n}{4} \quad \frac{n}{8} \quad \dots \quad 1$

or

$1 \quad 2 \quad 4 \quad \dots \quad \frac{n}{2} \quad n$

$\underbrace{\hspace{15em}}$
 $(\log_2 n + 1) \text{ terms}$
 \downarrow
levels

$$\begin{aligned} T.N.O &= 2 \cdot (l-1) \cdot n = 2 \cdot (\log_2 n + 1 - 1) n \\ &= 2 n \log n \end{aligned}$$

$$T.C. = O(n \log n)$$

Best Case : $O(n \log n)$

Avg Case : $O(n \log n)$

Worst Case : $O(n \log n)$

Time and Space complexity

In this basic code, the amount of space used = $(l-1) \cdot n$
↓
S.C. = $O(n \log n)$

$$= \boxed{n \log_2 n}$$

↓
We can improve it by deletion of a & b after merging into array

Best Case: $O(n)$

Avg Case: $O(n)$

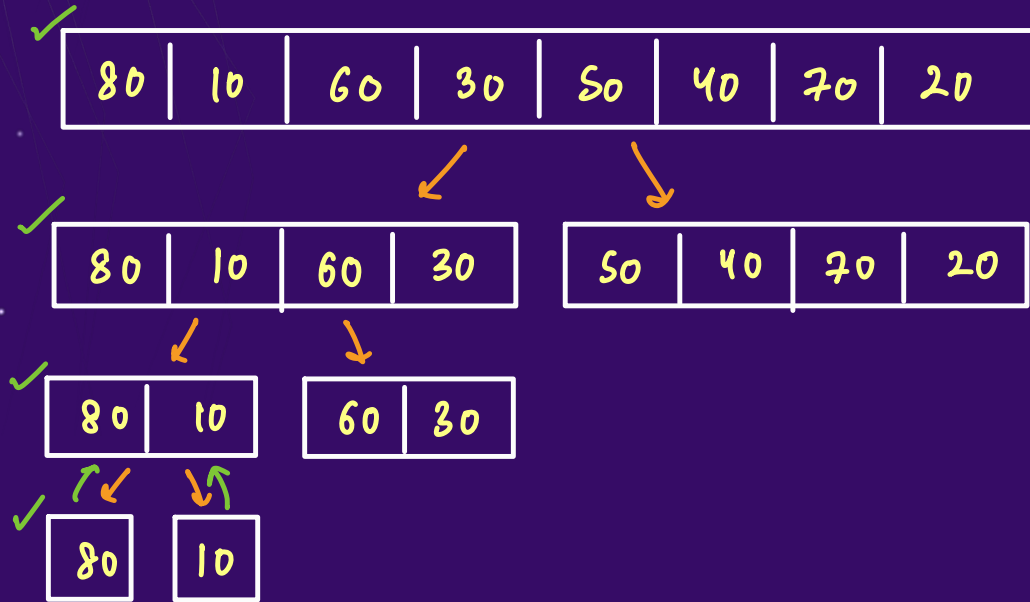
Worst Case: $O(n)$

Time and Space complexity

Maths

$$1 + \frac{1}{2} + \frac{1}{4} + \dots < 2$$

$$n(1 + \frac{1}{2} + \frac{1}{4} + \dots) < 2 * n$$



Extra Space Used: $n + \frac{n}{2} + \frac{n}{4} + \dots = n(1 + \frac{1}{2} + \frac{1}{4} + \dots \infty)$ if $< 2n$

S.C. = $O(n)$

Stability of Merge Sort

↓
yes. Merge Sort is stable

arr = 5 1 2 5* 3
 ↓ sorting
 1 2 3 5 5*

Make sure to code like this

```
if (a[i] <= b[j])  
    c[k++] = a[i++];
```

if I use ' $<$ ' instead of ' $<=$ ' then it won't be stable

Applications of Merge Sort

1. $O(n \log n)$ worst case time complexity
2. Custom Sorting
3. Sorting Linked Lists
4. Inversion Count & Related Problems

Inversions Count Problem

arr = 8 2 5 3 1 4

(8,2) (2,1) (5,3) (3,1)
(8,5) (5,1)
(8,3) (5,4)
(8,1)
(8,4)

ans = 10

Q, You have to count the
no. of inversions

$i < j$

$arr[i] > arr[j]$

Inversions Count Problem

Basic Approach (Brute Force)

```
int count = 0;
for(int i = 0; i < n-1; i++){
    for(int j = i+1; j < n; j++){
        if(arr[i] > arr[j]) count++;
    }
}
```

T.C. = $O(n^2)$ S.C. = $O(1)$

Definitely not the best

Inversions Count Problem

arr = 8 2 5 3 1 4

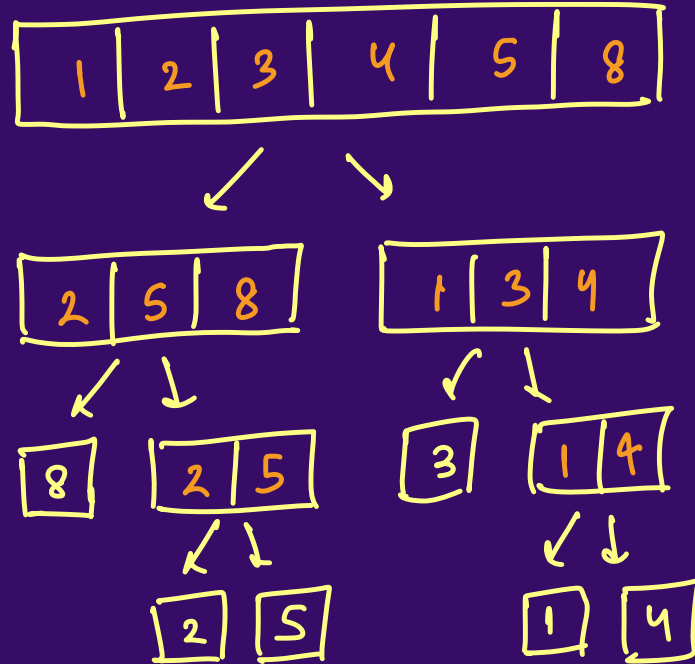
a = 2 5 8
i

b = 1 3 4
j

$arr[i] > arr[j]$
↓ * ↘
ele from a ele from b

Count = 0 3 8 7

Inversions Count Problem



Count

0
1
2
3
6
8
10

Inversions Count Problem

Create two arrays a & b of $n/2$ size each & copy paste.

Sort(a)

Sort(b)

Count inversions in a & b

Merge(a, b, arr)

Inversions Count Problem

```
public static void inversion(int[] a, int[] b){
    int i = 0, j = 0;
    while(i<a.length && j<b.length){
        if(a[i]>b[j]){
            count += (a.length-i);
            j++;
        }
        else i++;
    }
}
```

$$T.C. = O(n^2 \log n)$$

```
public static void mergesort(int[] arr){
    int n = arr.length;
    if(n==1) return; // base case
    // create two arrays of n/2 size each
    int[] a = new int[n/2];
    int[] b = new int[n-n/2];
    // copy pasting
    for(int i=0;i<n/2;i++) a[i] = arr[i];
    for(int i=0;i<n-n/2;i++) b[i] = arr[i+n/2];
    // magic
    mergesort(a);
    mergesort(b);
    inversion(a,b); // Extra
    // merge these 'a' and 'b'
    merge(a,b,arr);
    // delete a and b
    a = null; b = null;
}
```

Homework :

Q : Reverse Pairs



[Leetcode 493]

◀ **THANK YOU** ▶