



Binary Search

Today's checklist

1. Basics of Binary Search
2. Time complexity Analysis
3. Lower and upper bounds
4. Questions on binary search
5. Binary search on answer questions

What and Why? → It is very very efficient



- ↓ It is used to search a given element in a sorted space(array)
- ↘ It is used to reduce the search space in half after every turn
- Usually we apply B.S when it is given that 'it is a sorted array'

arr = { 10, 15, 21, 34, 81, 105, 180, 500 }

target = 105

Linear Search → T.C. = $O(n)$

What and Why?

Sorted array



arr = { ⁰10, 15, 21, 34, _m81, 105, 180, 500 ⁿ⁻¹ }

target = 81

comparisons
↓
1+1+1

if(arr[mid] < target) go right
else if(arr[mid] > target) go left
else if(arr[mid] == target) done —

Binary search algorithm

arr = { 10, 15, 21, 34, 81, 105, 180, 500 }

target = 110

```

if(arr[mid] < target) go right
else if (arr[mid] > target) go left
else if (arr[mid] == target) done —

```

Binary search algorithm

arr = { 0 1 2 3 4 5 6 7 8
10, 15, 21, 34, 81, 105, 180, 500, 614 }

lo hi
mid

target = 21

```
flag = false;  
while (lo <= hi) {  
    int mid = (lo + hi) / 2;  
    if (arr[mid] < target) lo = mid + 1;  
    else if (arr[mid] > target) hi = mid - 1;  
    else if (arr[mid] == target) flag = true; break;  
}
```

Ques: $100000 \rightarrow 50000 \rightarrow 25000 \rightarrow 12500 \rightarrow 6250 \rightarrow 3125$
 \downarrow
 $1 \leftarrow 3 \leftarrow 6 \leftarrow 12 \leftarrow 25 \leftarrow 50 \leftarrow 100 \leftarrow 200 \leftarrow 400 \leftarrow 800 \leftarrow 1600$

Q1: Binary Search

↓
Reduction in search space by half

Consider that we have 100000 elements present in a sorted array

Linear Search

$$T_{no} = 100000$$

Binary Search

$$T_{no} \approx 17$$

↓
 $T.C. = O(\log n)$

[Leetcode 704]

Time complexity analysis

$$n \rightarrow \frac{n}{2} \rightarrow \frac{n}{4} \rightarrow \frac{n}{8} \dots \dots 2 \rightarrow 1$$

or

$$1 \rightarrow 2 \rightarrow 4 \rightarrow \dots \dots \frac{n}{2} \rightarrow n$$

x terms

$$1 \rightarrow 2^1 \rightarrow 2^2 \rightarrow 2^3 \dots \dots 2^{x-1}$$

$$T.C. = O(x)$$

$$T.C. = O(\log_2 n)$$

$$T.C. = O(\log n)$$

$$2^{x-1} = n$$

$$\log_2 n = x-1 \text{ or}$$

$$x = 1 + \log_2 n$$

Time complexity analysis

$$\text{mid} = \frac{\text{lo} + \text{hi}}{2}$$

lo & hi are 'int'

$$-2^{31} \leq \text{lo}, \text{hi} \leq 2^{31} - 1$$

What if $(\text{lo} + \text{hi}) > 2^{31} - 1$ (Error: Integer overflow Error)

→ But if $\text{lo} = 1 \quad \Delta \quad \text{hi} = 2^{31} - 1$ then $\text{mid} = 2^{30} \leq 2^{31} - 1$

$$\text{mid} = \frac{\text{lo} + \text{hi}}{2} = \frac{(\text{hi} - \text{lo}) + (\text{lo} + \text{lo})}{2} = \frac{\text{hi} - \text{lo}}{2} + \text{lo}$$

*Lower bound \rightarrow index

Q2 : Given a sorted integer array and an integer 'x', find the lower bound of x. *smallest idx such that $\text{arr}[\text{idx}] \geq x$*

arr = $\begin{matrix} & 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ \{ & 10, & 20, & 30, & 30, & 40, & 50, & 60, & 70 \} \end{matrix}$

x = 40
lb = 4

x = 30
lb = 2

x = 35
lb = 4

x = 25
lb = 2

x = 80
lb = 8

x = 5
lb = 0

*Lower bound

$arr[idx] \geq x$
↓
minimise



Q2 : Given a sorted integer array and an integer 'x', find the lower bound of x.

arr = { 0 1 2 3 4 5 6 7
10, 20, 30, 30, 40, 50, 60, 70 }
hi lo
mid

x = 80

```
int lb = 8;  
while(lo <= hi){  
    mid = (lo+hi)/2;  
    if(arr[mid] >= x){  
        lb = min(lb, mid);  
        hi = mid-1;  
    }  
    else lo = mid+1;  
}
```

lb = 8

T.C. = $O(\log n)$

*Upper bound

Q3 : Given a sorted integer array and an integer 'x', find the upper bound of x. *ub is the minimum idx such that $\text{arr}[\text{idx}] > x$*

arr = { 10, 20, 30, 30, 40, 50, 60, 70 }

x = 30

lb = 2

ub = 4

x = 35

lb = 4

ub = 4

x = 5

lb = 0

ub = 0

x = 80

lb = 8

ub = 8

Same code for
lower bound except
 $\text{arr}[\text{mid}] \geq x$ the
statement is $\text{arr}[\text{mid}] > x$

Ques: lower bound ?
 ↑
 upper bound ?

Q4 : Find First and Last Position of Element in Sorted Array

arr = {
 0 1 2 3 4 5 6 7 8 9 10 11
 10, 10, 20, 20, 20, 20, 20, 30, 30, 30, 40, 40 }

x = 20

fp = 2 = lb

lp = 6 = ub - 1

M-I : 1) Check if ele is
 present or not

2) Find lb

3) Find ub

Ques:

Q4 : Find First and Last Position of Element in Sorted Array

m-2:

arr = { 10, 10, 20, 20, 20, 20, 20, 30, 30, 30, 40, 40 }

lo mid hi

target = 20

```
if(arr[mid] == target){
    if(arr[mid] == arr[mid-1]) hi = mid - 1;
    else {fp = mid; break;}
}
else if (arr[mid] < target) lo = mid + 1;
else if (arr[mid] > target) hi = mid - 1;
```

[Leetcode 34]

arr = { 10, 10, 20, 20, 20, 20, 20, 30, 30, 30, 40, 40 }

target = 20

```

if (arr[mid] == target) {
    if (arr[mid] == arr[mid+1]) lo = mid+1;
    else { lp p = mid; break; }
}
else if (arr[mid] < target) lo = mid+1;
else if (arr[mid] > target) hi = mid-1;

```

[Leetcode 34]

Ques:

Q5 : Peak index in a Mountain Array

arr = { 10, 20, 30, 50, 40, 20, 10 }

↑
peak

peak = 3

if (arr[i] > arr[i-1] && arr[i] > arr[i+1]) → peak = i

Binary → array must be sorted ?

search

Not everytime!

[Leetcode 852]

Ques:

Q5 : Peak index in a Mountain Array

arr = { 10, 20, 30, 40, 50, 60, 70, 20, 10 }

 0 1 2 3 4 5 6 7 8

 lo mid hi

```
if (arr[mid] > arr[mid-1] && arr[mid] > arr[mid+1]) return mid;  
else if (arr[m] > arr[m-1] && arr[m] < arr[m+1]) lo = mid + 1;  
else if (arr[m] < arr[m-1] && arr[m] > arr[m+1]) hi = mid - 1;
```

[Leetcode 852]

Ques:

Q6 : Sqrt(x)

for any no. x , $0 \leq \text{sqrt}(x) \leq x$ hint

$x = 16$

0 1 2 3

4

5 6 7

8 9 10

11

12

13

14

15

16

lo
hi
m

if ($m^2 == x$) return m ;

else if ($m^2 > x$) $hi = m - 1$;

else if ($m^2 < x$) $lo = m + 1$;

[Leetcode 69]

Ques:

Q6: Sqrt(x)

$$hi = 2^{25}$$

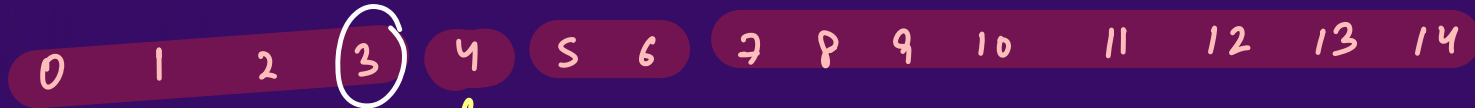
$$lo = 0$$

$$mid = 2^{24}$$

$$mid * mid = 2^{24} \times 2^{24} = 2^{48} > \text{Integer.MAX-VALUE}$$



$$x = 14$$



if ($m^2 == x$) return m;

else if ($m^2 > x$) hi = m - 1;

else if ($m^2 < x$) lo = m + 1;

$$m^2 == x$$

or

$$m == x/m$$

long

[Leetcode 69]

Ques:

$$T.C. = O(\log n)$$

Q7 : Search in Rotated Sorted Array [Very Important]

arr =

0	1	2	3	4	5	6	7
0	1	2	3	4	5	6	7

↪ 3 4 5 6 (7) (0) 1 2

Method-1: 1) Find the pivot idx

↓
 $O(\log n)$

2) Binary Search in
 $(0, p)$ & $(p+1, n-1)$

[Leetcode 33]

Ques: pivot \rightarrow largest

Q7 : Search in Rotated Sorted Array

	0	1	2	3	4	5	6	7	
arr	=	3	4	5	6	7	0	1	2
					lo	m		hi	

int p;

if (arr[m] > arr[m-1] && arr[m] > arr[m+1]) p = mid; break;

else if (arr[m] < arr[m-1] && arr[m] < arr[m+1]) p = mid-1; break;

else if (arr[m] > arr[m-1] && arr[m] < arr[m+1]) {

 if (arr[m] > arr[n-1]) lo = mid + 1;

 else hi = mid - 1;

}

[Leetcode 33]

Ques:

{ 0 1 2 3 4 5 6 7 8 9 }

Q7 : Search in Rotated Sorted Array

target = 8

arr =

0	1	2	3	4	5	6	7	8	9
7	8	9	0	1	2	3	4	5	6

lo m hi

```
else if (arr[mid] <= arr[hi]) { //right
    if (arr[mid] <= target && target <= arr[hi]) lo = mid + 1;
    else hi = mid - 1;
```

```
    }
    else if (arr[mid] > arr[hi]) { // you are in left sorted array
        if (arr[lo] <= target && target <= arr[mid]) hi = mid - 1;
        else lo = mid + 1;
    }
}
```

[Leetcode 33]

Ques:

Q7 : Search in Rotated Sorted Array

```
int n = arr.length;
int lo = 0, hi = n-1;
while(lo<=hi){
    int mid = lo + (hi-lo)/2;
    if(arr[mid]==target) return mid;
    else if(arr[mid]<=arr[hi]){ // i am in right sorted array (mid to high everything is sorted)
        if(target>arr[mid] && target<=arr[hi]) lo = mid + 1;
        else hi = mid - 1;
    }
    else{ // i am in left sorted array (lo to mid everything is sorted)
        if(target>=arr[lo] && target<arr[mid]) hi = mid - 1;
        else lo = mid + 1;
    }
}
return -1;
```

Ques:

Q8 : Find K Closest Elements

arr = { 1, 2, ⁱ3, ^j4, 5 }

lb

x = 3

K=2

arr = { 1, 3, 5, ⁱ8, ^j9, 10 }

lb

x = 6

K=3

↓
lb = 8

arr = { 1, 2, 3, 4, 5 } x = -1 k = 4



[Leetcode 658]

Ques:

Q8 : Find K Closest Elements

arr = { 1, 2, 3, 4, 5 }

0 1 2 3 4 5

x = 7 , k = 3

↓

lb = 5

For lb → arr[lb] ≥ x

↓
minimise

[Leetcode 658]

Ques:

Q8 : Find K Closest Elements

```
List<Integer> ans = new ArrayList<>();
int n = arr.length;
if(x<arr[0]){
    for(int i=0;i<k;i++){
        ans.add(arr[i]);
    }
    return ans;
}
if(x>arr[n-1]){
    for(int i=n-1;i>=n-k;i--){
        ans.add(arr[i]);
    }
    Collections.sort(ans);
    return ans;
}
```

Base Cases

```
int lb = n;
int lo = 0, hi = n-1;
while(lo<=hi){
    int mid = lo + (hi-lo)/2;
    if(arr[mid]>=x){
        lb = mid;
        hi = mid - 1;
    }
    else lo = mid + 1;
}
```

Binary Search to
find lb of x

2 pointer

```
int j = lb, i = lb - 1;
while(k>0 && i>=0 && j<n){
    int di = Math.abs(x-arr[i]);
    int dj = Math.abs(x-arr[j]);
    if(di<=dj){
        ans.add(arr[i]);
        i--;
    }
    else{
        ans.add(arr[j]);
        j++;
    }
    k--;
}
while(i<0 && k>0){
    ans.add(arr[j]);
    j++;
    k--;
}
while(j==n && k>0){
    ans.add(arr[i]);
    i--;
    k--;
}
Collections.sort(ans);
return ans;
```

Ques:

Q9 : Find the **smallest** Divisor given a Threshold

$$sd = 5$$

$$arr = \{5, 1, 9, 2\}$$

$$t = 6$$

max sum of array

$$d = 1 \quad \text{sum} = 17 \quad \alpha$$

$$d = 2 \quad \text{sum} = 3 + 1 + 5 + 1 = 10 \quad \alpha$$

$$d = 3 \quad \text{sum} = 2 + 1 + 3 + 1 = 7 \quad \alpha$$

$$d = 4 \quad \text{sum} = 2 + 1 + 3 + 1 = 7 \quad \alpha$$

$$d = 5$$

$$\text{sum} = 1 + 1 + 2 + 1 = 5 < 6 \quad \checkmark$$

$$d = 9 \quad \text{sum} = \frac{5}{9} + \frac{1}{9} + \frac{9}{9} + \frac{2}{9}$$

$$1 + 1 + 1 + 1 = 4 < 6 \quad (\text{Acceptable})$$

[Leetcode 1283]

Ques:

Q9 : Find the smallest Divisor given a Threshold

```
int d;  
for(d=1;d<=mx;d++){  
    int sum = 0;  
    for(int i=0;i<n;i++){  
        if(arr[i]%d==0) sum += arr[i]/d;  
        else sum += arr[i]/d + 1;  
    }  
    if(sum<=t) return d;  
}  
return d;
```

$$\begin{aligned} \text{arr}[i] &\leq 10^6 \\ mx &\leq 10^6 \\ \log_2 mx &\leq \log_2 10^6 \\ &\downarrow \\ &\leq \boxed{24} \end{aligned} \qquad \begin{aligned} &\downarrow \\ 6 \log_2 10 &\downarrow \\ &\boxed{3-4} \end{aligned}$$

$$\begin{aligned} \text{T.C.} &= O(n^* mx) \\ &\quad \downarrow \\ &\text{largest element of array} \end{aligned}$$

[Leetcode 1283]

Ques:

Q9 : Find the smallest Divisor given a Threshold

$$\text{arr} = \{ 44, 22, 33, 11, 1 \} \quad t = 6$$

lo hi
d = 1 to 44 ✓ not linearly

$$lo = \cancel{1} \cancel{25} \cancel{28} 31$$

$$hi = \cancel{44} 32$$

$$mid = \cancel{22} \cancel{33} \cancel{27} \cancel{30} 31$$

$$d = 22 \quad \text{sum} = 2 + 1 + 2 + 1 + 1 = 7 \alpha$$

$$d = 33 \quad \text{sum} = 2 + 1 + 1 + 1 + 1 = 6 \leq 6 \checkmark$$

$$d = 27 \quad \text{sum} = 2 + 1 + 2 + 1 + 1 = 7 \alpha$$

$$d = 30 \quad \text{sum} = 2 + 1 + 2 + 1 + 1 = 7 \alpha$$

[Leetcode 1283]

Ques:

Q9 : Find the smallest Divisor given a Threshold

```
int d = 1;
int lo = 1, hi = mx;
while(lo<=hi){ → log(mx)
    int mid = lo + (hi-lo)/2;
    if(isLess(mid,arr,t)){ O(n)
        d = mid;
        hi = mid - 1;
    }
    else lo = mid + 1;
}
return d;
```

```
public boolean isLess(int mid, int[] arr, int t){
    int sum = 0;
    for(int i=0;i<arr.length;i++){
        if(arr[i]%mid==0) sum += arr[i]/mid;
        else sum += arr[i]/mid + 1;
    }
    if(sum<=t) return true;
    return false;
}
```

$$T.C. = O(n \cdot \log mx)$$

[Leetcode 1283]

Ques: $hi = sum$, $lo = mx$

Q10 : Capacity to ship packages within D days

$arr = \{ 1, 2, 3, 4, 5, 6, 7, 8, 9, 10 \}$ $d=5$

1 2 3 4 5 6 7

$C = 8 \rightarrow days = 8 \propto$ NEVER

$C = 11 \rightarrow days = 6 \propto$

$C = 55 \rightarrow days = 1 \checkmark$

\downarrow
ship capacity

[Leetcode 1011]

Ques: $lo = mx$, $hi = sum$

Q10 : Capacity to ship packages within D days

$arr = \{ 3, 2, 2, 4, 1, 4 \}$ $days = 3$

$\underline{3_2} \quad \underline{2} \quad \underline{4_1} \quad \underline{4}$

$minC = 10$

$lo = 1$

$hi = 15$

$mid = 10$

$C = 10 \rightarrow d = 2 \checkmark$

$C = 6 \rightarrow d = 3 \checkmark$

$C = 4 \rightarrow d = 5 \times$

$C = 5 \rightarrow d = 4 \times$

[Leetcode 1011]

arr = { 3, 2, 2, 4, 1, 4 } days = 3

$$\begin{array}{r} 322 \\ \hline \end{array} \quad \begin{array}{r} 41 \\ \hline \end{array} \quad \begin{array}{r} 4 \\ \hline \end{array}$$
$$D = 1 \neq 3$$
$$C = 8$$

load = \emptyset β β \neq γ γ
 γ

```
if (load + arr[i] <= c) load += arr[i]
```

```
else {
```

```
load = arr[i];  
D++;
```

[Leetcode 1011]

Ques:

Q11 : Minimized maximum of Products distributed to any store

warehouse
11 units of iPhone
6 units of Samsung
 $arr = \{ 11, 6 \}$
 $n = 6$
 $i \quad s$

11i 6s 0 0 0 0 $max = 11$
Retail Stores

3i 3i 3i 2i 3s 3s $max = 3$
Retail Stores

[Leetcode 2064]

Ques:

Q11 : Minimized maximum of Products distributed to any store

$$\text{arr} = \left\{ \underset{a}{15}, \underset{b}{10}, \underset{c}{10} \right\} \quad n=7$$

$$\underline{s_a} \quad \underline{s_a} \quad \underline{s_a} \quad \underline{s_b} \quad \underline{s_b} \quad \underline{s_c} \quad \underline{s_c}$$

$$lo = 5$$

$$hi = 15$$

$$mid = 8$$

$$\text{maxAns} = 8$$

[Leetcode 2064]

Ques:

Q11 : Minimized maximum of Products distributed to any store

$$\text{arr} = \{ 15, 12, 10 \} \quad n=7$$

$$\text{stores} = 10$$

$$\text{maxQ} = 4$$

$$\begin{array}{cccc|ccc|ccc} 4 & 4 & 4 & 3 & 4 & 4 & 4 & 4 & 4 & 2 \end{array}$$

$$\text{stores} += \text{arr}[i] / \text{maxQ} + \underbrace{1}_{\downarrow}$$

$$\text{arr}[i] \% \text{maxQ} \neq 0$$

[Leetcode 2064]

Ques:

Q11 : Minimized maximum of Products distributed to any store

$arr = \{100, 10, 10\}$ $n = 8$

$maxAns = 10$

$lo = 1$

$hi = mx \text{ of array}$

[Leetcode 2064]

Ques:

Aggressive Cows, Koko Eating Bananas,
min. Time to complete Trips



Products distributed to any

```
public int minimizedMaximum(int n, int[] arr) {  
    int m = arr.length, mx = Integer.MIN_VALUE;  
    for(int i=0;i<m;i++){  
        mx = Math.max(mx,arr[i]);  
    }  
    int lo = 1, hi = mx;  
    int ans = 0;  
    while(lo<=hi){  
        int mid = lo + (hi-lo)/2;  
        if(isPossible(mid,n,arr)){  
            ans = mid;  
            hi = mid - 1;  
        }  
        else lo = mid + 1;  
    }  
    return ans;  
}
```

```
public boolean isPossible(int maxQ, int n, int[] arr){  
    int stores = 0;  
    for(int i=0;i<arr.length;i++){  
        if(arr[i]%maxQ == 0) stores += arr[i]/maxQ;  
        else stores += arr[i]/maxQ + 1;  
    }  
    if(stores>n) return false;  
    return true;  
}
```

[Leetcode 2064]

Ques:

$O(\log n)$
↑

Q12 : Kth Missing Positive Number (Binary Search + Maths)

arr = { 2, 3, 4, 7, 11 } k = 5

0 1 2 3 4

1 1 1 3 6

✓ 1 2 3 4 5 6 7 8 9 10 11

✓
B-S.



Q12 : Kth Missing Positive Number

arr = { 2, 3, 4, 7, 11 } k = 5

$$\text{missed} = \text{arr}[\text{mid}] - (\text{mid} + 1)$$

```
if (missed < K) lo = mid + 1
```

```
else hi = mid - 1
```

$$k^{\text{th}} \text{ missing number} = \text{arr}[hi] + \text{extra}$$
$$\begin{aligned} \text{extra} &= k - (\text{missed no.s till hi}) \\ &= k - (\text{arr[hi]} - (\text{hi} + 1)) \end{aligned}$$

Observations : The k^{th} missing no. is b/w $\text{arr}[hi]$ & $\text{arr}[lo]$ [Leetcode 1539]

Ques:

Q12 : Kth Missing Positive Number

$$\begin{aligned} k^{\text{th}} \text{ missing no} &= \text{arr}[hi] + \text{extra} \\ &= \text{arr}[hi] + k - (\text{arr}[hi] - (hi+1)) \\ &= \cancel{\text{arr}[hi]} + k - \cancel{\text{arr}[hi]} + (hi+1) \\ &= \boxed{k + (hi+1)} \\ &\quad \text{or} \\ &\quad k + lo \end{aligned}$$

[Leetcode 1539]

Ques:

$arr = \begin{matrix} & 0 & 1 & 2 & 3 & 4 \\ \{ & 2, & 3, & 4, & 7, & 11 \} \end{matrix}$ $x = 5$
 $hi \quad lo$

Q12 : Kth Missing Positive Number

K^{th} missing no. is b/w $arr[hi]$ & $arr[lo]$

$$\rightarrow K^{th} \text{ missing no} = arr[lo] - x$$

$$x = \text{missing no.s till } arr[lo] - K - 1$$

$$x = arr[lo] - (lo + 1) - K - 1$$

$$K^{th} \rightarrow arr[lo] - (arr[lo] - (lo + 1) - (K - 1))$$

$$\cancel{arr[lo]} - \cancel{arr[lo]} + \cancel{lo + 1} + (K - 1)$$

[Leetcode 1539]

Ques:

Q12 : Kth Missing Positive Number

```
int lo = 0, hi = arr.length - 1;
while(lo<=hi){
    int mid = lo + (hi-lo)/2;
    int missed = arr[mid] - (mid+1);
    if(missed<k) lo = mid + 1;
    else hi = mid - 1;
}
return k + lo;
```

[Leetcode 1539]

◀ **THANK YOU** ▶