

# Ques:

Q : Longest Common Substring (Tabulation)

	ϕ	f	a	b	g	d	k
ϕ	0	0	0	0	0	0	0
a	0	0	1	0	0	0	0
b	0	0	0	2	0	0	0
c	0	0	0	0	0	0	0
d	0	0	0	0	0	1	0
e	0	0	0	0	0	0	0
sl							

(m+1) \* (n+1)

ans is not  
 $dp[m][n]$

# Ques:

*length of*  
**Q : Longest Common Substring (Tabulation)**

```
int m = a.length(), n = b.length();
int[][] dp = new int[m+1][n+1];
int max = 0;
for(int i=1;i<=m;i++){
    for(int j=1;j<=n;j++){
        if(a.charAt(i-1)==b.charAt(j-1))
            dp[i][j] = 1 + dp[i-1][j-1];
        else
            dp[i][j] = 0;
        max = Math.max(max,dp[i][j]);
    }
}
return max;
```

**Ques:**

↑ print LCS.

**Q : Shortest Common Supersequence** → string (length)

$s1 = "abac"$      $s2 = "cab"$   
└──────────┘ → LCS = ab

SCS = "cabac"

Find the length of shortest common supersequence of  $s1$  &  $s2$

• Think about LCS of  $s1, s2$

$$\begin{aligned} & \downarrow \\ & \text{LCS} + s1 - \text{LCS} \\ & + s2 - \text{LCS} \end{aligned}$$

[Leetcode 1092]

# Ques:

Q : Shortest Common Supersequence (print)

s1 = a d b e c f

s2 = g a h b i c

LCS = abc  
↑  
K

scs =  
g a d h b e i c f  
g a h d b i e c f  
g a d h b e e c f  
g a h d b i i c f

[Leetcode 1092]

# Ques:

Q : Shortest Common Supersequence

$s1 = a d b e f c$   
 $s2 = z g a h b i c m n$

$i$   
 $\uparrow$

$LCS = abc$

$\uparrow$   
 $k$

$j$

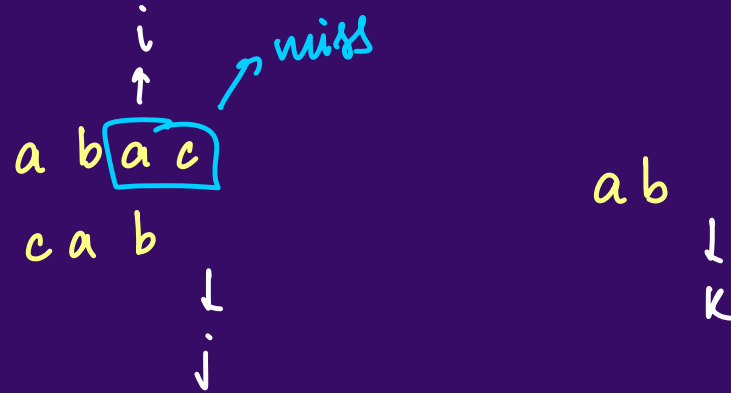
$while (s1[i] \neq lcs[k]) i++$   
 $while (s2[j] \neq lcs[k]) j++$

$scs = z g a d h b e f i c$

[Leetcode 1092]

# Ques:

Q : Shortest Common Supersequence



SCS = c a b ac

[Leetcode 1092]

# Ques:

## Q : Shortest Common Supersequence

Diagram illustrating the Shortest Common Supersequence (SCS) problem. Two strings are shown: "b b b a a a b a" (top) and "b b a b a b b b" (bottom). The characters are indexed from 1 to 10. The Longest Common Subsequence (LCS) is highlighted in green: "b b a a b". The indices for the LCS are marked as  $i$  (top) and  $j$  (bottom).

LCS = b b a a b

SCS = b b b a b a a b b b a  
b b b a a a b a b b b

[Leetcode 1092]

# Ques:

## Q : Shortest Common Supersequence

S1 = a d b e f c

S2 = g a h b i c

a d b e f c

g a h b i c

SCS = g a d h b i e f c

store LCS wali  
string



# Ques:

## Q : Shortest Common Supersequence

```
public String shortestCommonSupersequence(String a, String b) {  
    String lcs = LCS(a,b);  
    int i = 0, j = 0, k = 0; // i->a, j->b, k->lcs  
    String scs = "";  
    while(k<lcs.length()){  
        while(a.charAt(i)!=lcs.charAt(k)){  
            scs += a.charAt(i++);  
        }  
        while(b.charAt(j)!=lcs.charAt(k)){  
            scs += b.charAt(j++);  
        }  
        scs += lcs.charAt(k);  
        i++; j++; k++;  
    }  
    while(j<b.length()){  
        scs += b.charAt(j++);  
    }  
    while(i<a.length()){  
        scs += a.charAt(i++);  
    }  
    return scs;  
}
```

[Leetcode 1092]

**Ques:** (Tabulation)

**Q : Palindromic Substrings**

$s = g a b c b a b$

$g, a, b, c, b, a, b$

$bcb, bab$

$abcba$

Method-1 : Brute Force

↓

generate all substrings & check each substring

[Leetcode 647]

**Ques:**

$$\sum k = \frac{k(k+1)}{2} \quad \sum k^2 =$$



**Q : Palindromic Substrings**

For ex:  $s = a b c b a$   $n=5$

$a, ab, abc, abcb, abcba$

$$1+2+\dots+n + 1+2+\dots+n-1 + 1+2+\dots+n-2$$

$$\sum_{k=1}^n 1+2+3+\dots+k = \sum_{k=1}^n \frac{k(k+1)}{2} \approx \sum k^2 + k$$

$$\sum_1^n k^2 = \frac{n(n+1)(2n+1)}{6} = \boxed{n^3}$$

$$T.C. = O(n^3)$$

**[Leetcode 647]**

## Ques:

Q : Palindromic Substrings

$s =$ 

0	1	2	3	4	5	6
g	a	b	c	b	a	b

$s =$ 

	$i$					$j$
	a	b	c	b		a
	↑					↑

$s[i] == s[j] \Rightarrow (i+1, j-1)$  substring is palindrome

[Leetcode 647]

# Ques:

## Q: Palindromic Substrings

$s =$ <sup>0 1 2 3 4 5 6</sup>  
g a b c b a b



2-length substring  
palindrome  
kab hoti hai?

A.S. =  $O(n^2)$   
T.C. =  $O(n^2)$

$i, j \rightarrow i+1, j+1$

	0	1	2	3	4	5	6
0	1	0	0	0	0	0	0
1	X	1	0	0	0	1	0
2	X	X	1	0	1	0	0
3	X	X	X	1	0	0	0
4	X	X	X	X	1	0	1
5	X	X	X	X	X	1	0
6	X	X	X	X	X	X	1

Every cell stores  
true// if substring  
from  $i$  to  $j$   
is palindrome

[Leetcode 647]

# Homework:



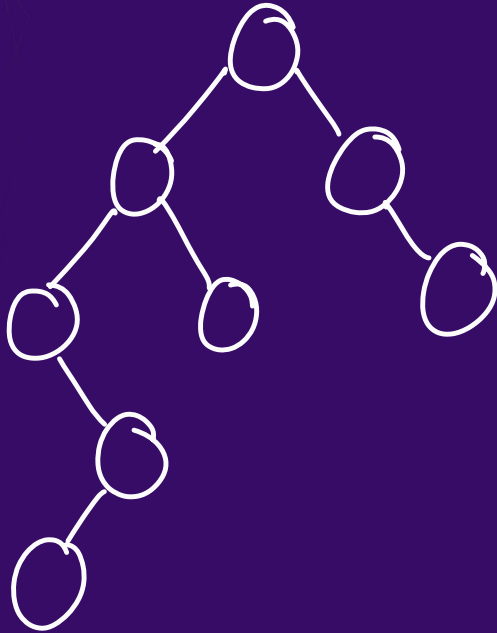
**Q : Longest Palindromic Substring**

**[Leetcode 5]**

**Ques:**

→ longest **path**

**Q : Diameter of a Binary Tree**



maximum = diameter  
↑  
path length =  $n - 1$   
↓  
no. of nodes  
in the path

$$n = 1 + \text{levels}(\text{left}) + \text{levels}(\text{right})$$

for a particular peak (node)  $\leftarrow \text{dia} = \text{levels}(\text{left}) + \text{levels}(\text{right})$

[Leetcode 543]

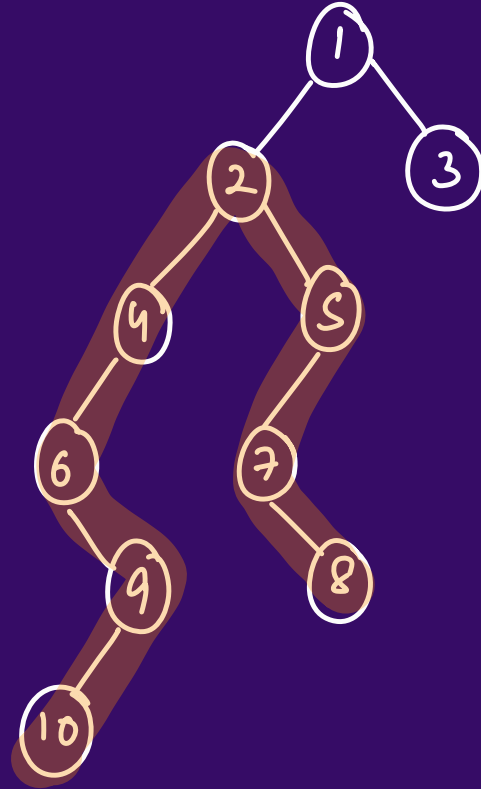
## Ques:

$$dia = levels(left) + levels(right)$$



## Q : Diameter of a Binary Tree

does not  
passes through  
the root



$O(n^2)$  T.C.

bcz for every node  
we are calling levels  $\rightarrow O(n)$

we can memoize using Hashmap

```
Map < TreeNode, Integer> dp;
```

## [Leetcode 543]



# Ques:

(Recursion + Memoization)

## Q : Diameter of a Binary Tree

```
public int levels(TreeNode root, Map<TreeNode,Integer> dp) {  
    if(root==null) return 0;  
    if(dp.containsKey(root)) return dp.get(root);  
    int leftLevels = levels(root.left,dp);  
    int rightLevels = levels(root.right,dp);  
    dp.put(root,1 + Math.max(leftLevels,rightLevels));  
    return dp.get(root);  
}  
  
public int diameter(TreeNode root, Map<TreeNode,Integer> dp) {  
    if(root==null) return 0;  
    int myDia = levels(root.left,dp) + levels(root.right,dp);  
    int leftDia = diameter(root.left,dp);  
    int rightDia = diameter(root.right,dp);  
    return Math.max(myDia,Math.max(leftDia,rightDia));  
}  
  
public int diameterOfBinaryTree(TreeNode root) {  
    Map<TreeNode,Integer> dp = new HashMap<>();  
    return diameter(root,dp);  
}
```

T.C. =  $O(n)$

A.S. =  $O(n)$

[Leetcode 543]

# Ques:

## Q : Diameter of a Binary Tree

```
public int levels(TreeNode root, int[] dia) {  
    if(root==null) return 0;  
    int leftLevels = levels(root.left,dia);  
    int rightLevels = levels(root.right,dia);  
    int path = leftLevels + rightLevels; // extra  
    dia[0] = Math.max(dia[0],path); // extra  
    return 1 + Math.max(leftLevels,rightLevels);  
}  
public int diameterOfBinaryTree(TreeNode root) {  
    int[] dia = {0};  
    levels(root,dia);  
    return dia[0];  
}
```

$$T.C. = O(n)$$

$$A.S. = O(h)$$

↓  
recursive  
stack

[Leetcode 543]

# Ques:

**Q : Diameter of a Binary Tree** (*without memoization*) DP

[Leetcode 543]

# Homework:



**Q : Balanced Binary Tree**

**[Leetcode 110]**

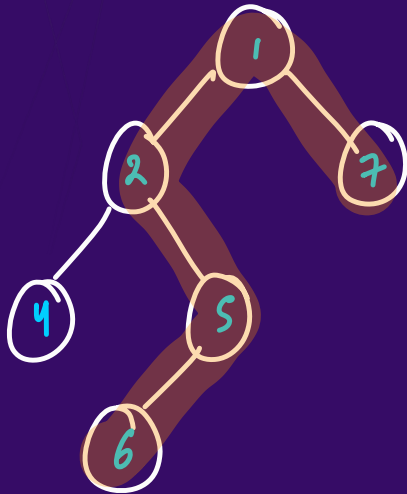
diameter/balance  $\rightarrow$  levels

**Ques:** max path sum  $\rightarrow$  ?

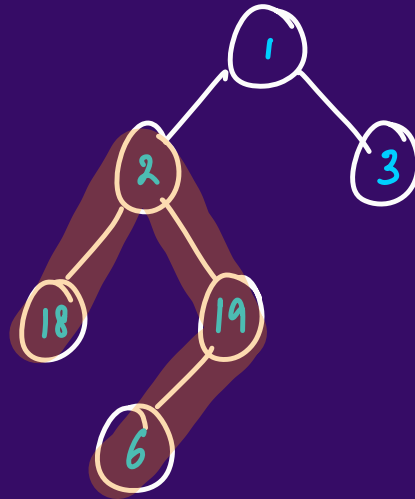
**Q:** Binary Tree Maximum **path** sum

PW SKILLS

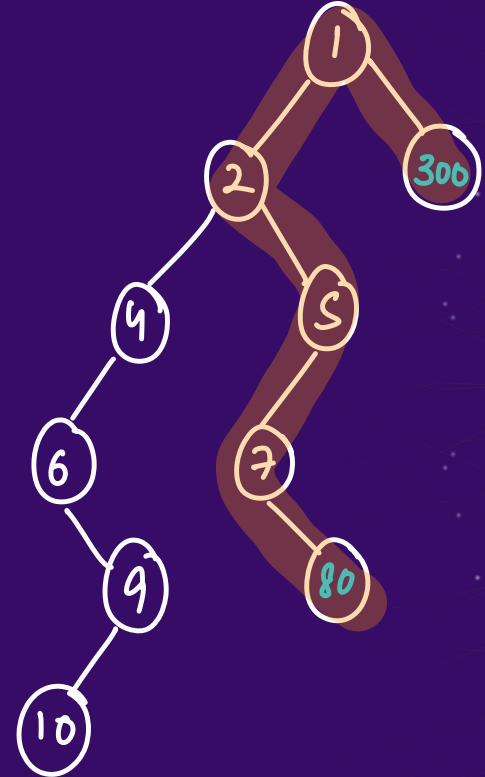
• Diameter of B.T.



$$6 + 5 + 2 + 1 + 7 = 21$$



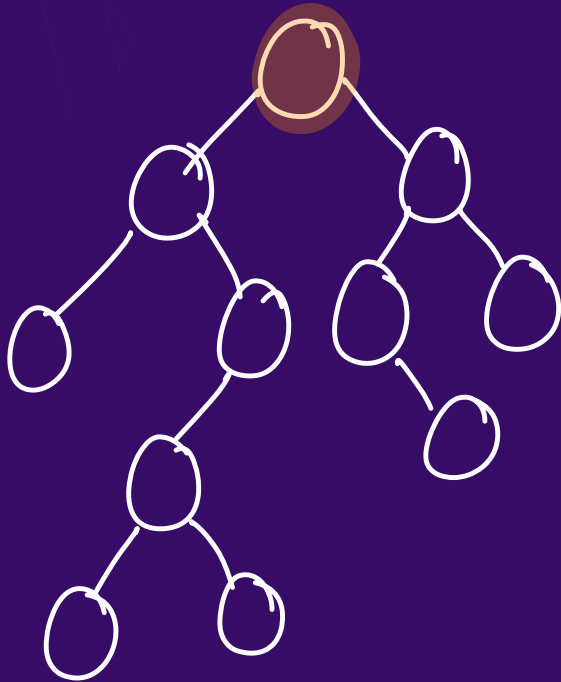
$$18 + 2 + 19 + 6$$



[Leetcode 124]

**Ques:** As of now I am considering leaf to leaf

**Q:** Binary Tree Maximum path sum



↑ root → leaf

$$\text{pathSum} = \text{root.val} + \text{lineSum}(\text{left}) + \text{lineSum}(\text{right})$$

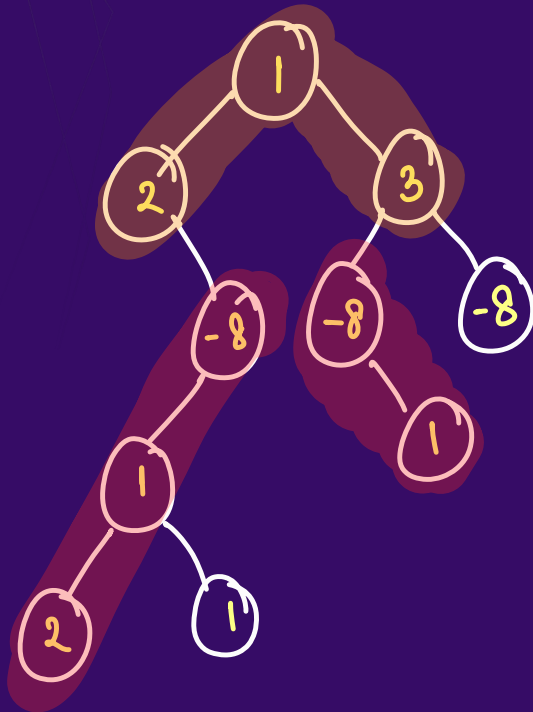
↓  
line sum ko 'maximise'

$$\text{maxSum} = \max(\text{sum}, \text{maxSum});$$

$\text{lineSum}(\text{root}) = \text{root.val} + \max(\text{lineSum}(\text{left}), \text{lineSum}(\text{right}))$  [Leetcode 124]

# Ques:

Q : Binary Tree Maximum path sum



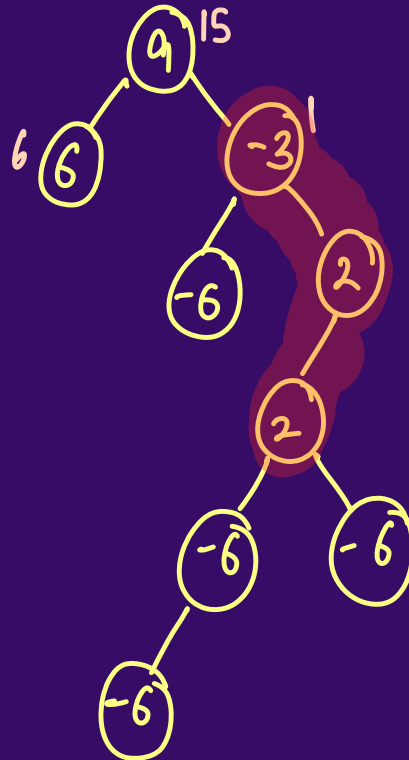
$$1 + 2 + 3 = \text{maxSum}$$

$$-3 \rightarrow \text{ans} = -3$$

[Leetcode 124]

# Ques:

Q : Binary Tree Maximum path sum



Out  $\rightarrow 15$

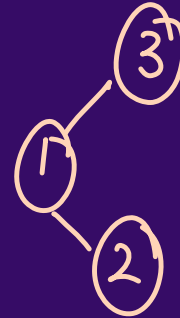
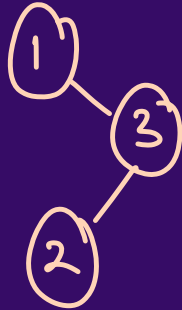
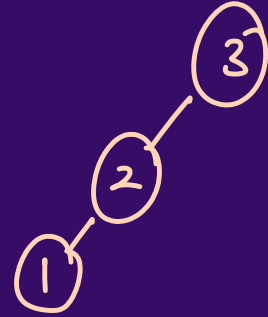
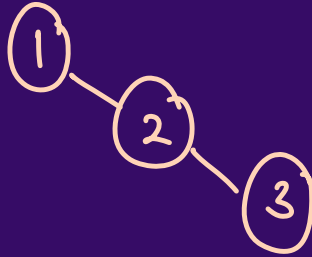
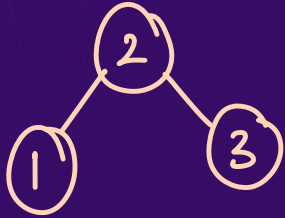
Exp  $\rightarrow 16$

[Leetcode 124]



**Ques:** dp??

**Q : Unique Binary Search Trees**  $n=3$



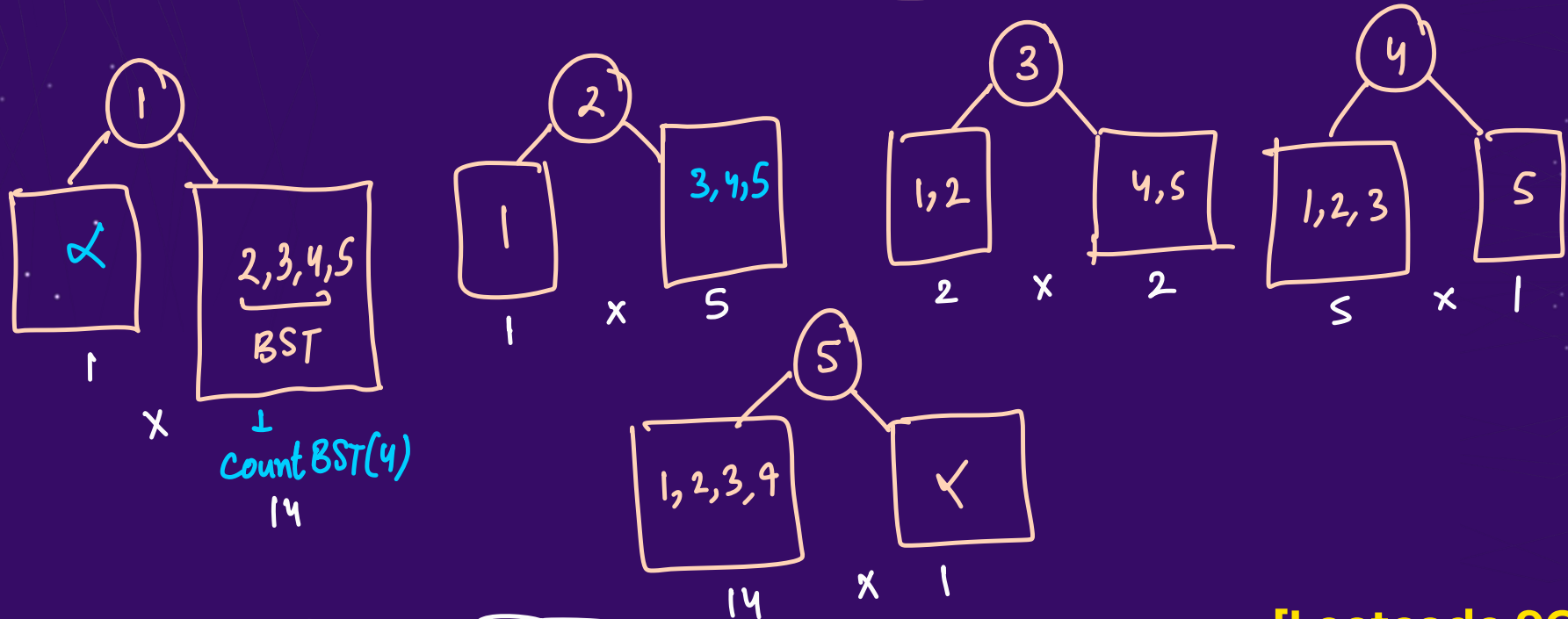
**[Leetcode 96]**

# Ques:

Q: Unique Binary Search Trees

$n=5$

Count BST(n)



[Leetcode 96]

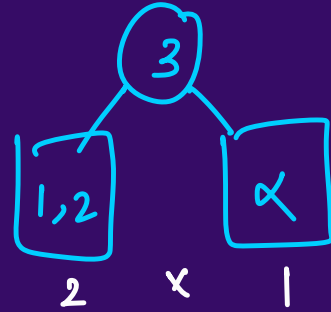
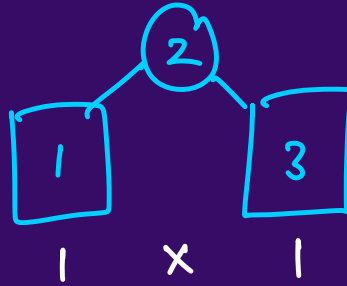
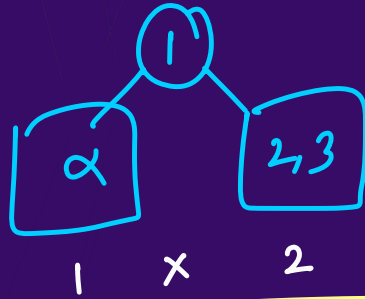
**Ques:**

dp

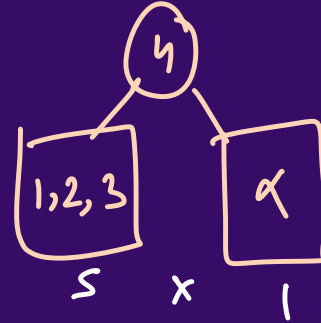
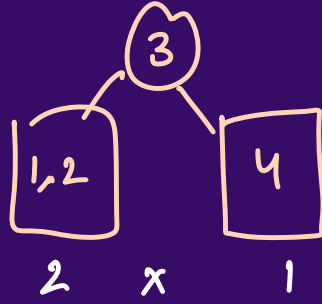
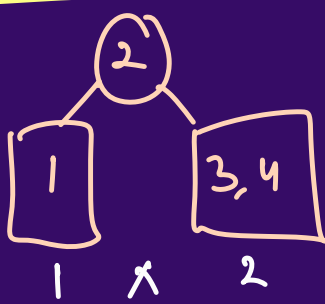
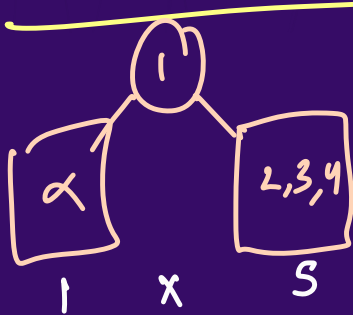
0	1	2	3	4	5
1	1	2	5	14	42



**Q:** Unique Binary Search Trees (1D DP) (nested loop)



= 5



[Leetcode 96]

# Ques:

## Q : Unique Binary Search Trees

```
public int numTrees(int n) {  
    if(n<=1) return 1;  
    int[] dp = new int[n+1];  
    dp[0] = 1; dp[1] = 1; dp[2] = 2;  
    for(int i=3;i<=n;i++){  
        for(int j=1;j<=i;j++){  
            dp[i] += (dp[j-1]*dp[i-j]);  
        }  
    }  
    return dp[n];  
}
```

$$T.C. = O(n^2)$$

$$A.S. = O(n)$$

[Leetcode 96]

# Ques:

Q : Longest Increasing Subsequence

arr = { 10, 9, 2, 5, 3, 7, 101, 18 }

1 1 1 2 2

↓  
2 2  
5 3  
7 7

---

arr = { 10, 9, 2, 5, 3, 7, 101, 18, 6 }

dp = { 1, 1, 1, 2, 2, 3, 4, 4, 3 }

[Leetcode 300]

# Ques:

## Q : Longest Increasing Subsequence

```
public int lengthOfLIS(int[] nums) {  
    int n = nums.length;  
    int[] dp = new int[n];  
    int max = 0;  
    for(int i=0;i<n;i++){  
        for(int j=0;j<=i-1;j++){  
            if(nums[j]<nums[i])  
                dp[i] = Math.max(dp[i],dp[j]);  
        }  
        dp[i] += 1;  
        max = Math.max(max,dp[i]);  
    }  
    return max;  
}
```

$$T.C. = O(n^2)$$

$$A.S. = O(n)$$

**[Leetcode 300]**

**Ques:**

Largest possible length of mountain array after  
↑  
removals



**Q : Minimum Number of Removals to make Mountain Array**

(2) 1 (1) 5 6 2 (3) 1 → ans = 3

1 5 6 2 1

2 5 6 3 1

**[Leetcode 1671]**

# Ques:

Q : Minimum Number of Removals to make Mountain Array

arr	2	1	1	5	6	2	3	1	
dp1	1	1	1	2	3	2	3	1	Normal LIS
dp2	2	1	1	3	3	2	2	1	Reverse LIS
	2	1	1	4	5	3	4	1	
	x	x	x					x	

[Leetcode 1671]



# Ques:

**Q : Minimum Number of Removals to make Mountain Array**

9 8 1 7 6 5 4 3 2 1

dp1

1 1

dp2

9 8

9 8

[Leetcode 1671]

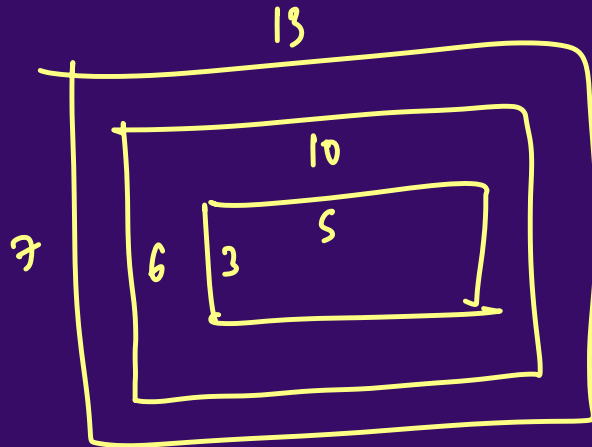
**Ques:** ↑ <sup>toy</sup>



**Q : Russian Doll Envelopes**

arr = { {3, 5}, {7, 13}, {6, 10}, {8, 4}, {6, 7}, {7, 11} }

ans = 3



(a, b) (c, d)

(a, b) can get inside (c, d)

if & only if

$a < c$  &  $b < d$

[Leetcode 354]

# Ques:

## Q : Russian Doll Envelopes

Hint  $\rightarrow$  1 Sorting (Custom) (Comparable)

Hint  $\rightarrow$  2 LIS

arr = { {3,5}, {7,13}, {6,10}, {8,4}, {6,7}, {7,11} }

sort(row) = {3,5} {6,7} {6,10} {7,11} {7,13} {8,4}

sort(col) = { {8,4}, {3,5}, {6,7}, {6,10}, {7,11}, {7,13} }

# Ques:

## Q : Russian Doll Envelopes

arr = { {3, 5}, {7, 13}, {6, 10}, {8, 4}, {6, 7}, {7, 11} }

Sort(w) = { 3, 5 } { 6, 7 } { 6, 10 } { 7, 11 } { 7, 13 } { 8, 4 }

Sort(w: asc)  
h: des = { 3, 5 } { 6, 10 } { 6, 7 } { 7, 13 }, { 7, 11 }, { 8, 4 }

◀ **THANK YOU** ▶