

# Hawaiian Shirt: No!

Told Byron: No.

Date: Tuesday, February 12, 2019

## Announcements

- HW #2 due Thursday @ midnight
- Read Ch. 3 of Probabilistic Programming for Hackers by Thursday
- Almost done with grading... will have key on Thursday

[ Slide 19 - Markov Chain Simulation ]

$$\mathbf{P} = \begin{pmatrix} & \text{to} \\ \text{from} & \left| \begin{array}{c} \\ P_{ij} \end{array} \right| \end{pmatrix}$$

$P_{ij} = P$ (moving from state i at time t to state j at time t+1).

Example 1:

state space

Question 2:  $S = \{1, 2\}$

where 1 = no rain  
2 = rain

$$\mathbf{P} = \begin{pmatrix} & 1 \\ 1 & \begin{pmatrix} 1-\beta & \beta \\ 1-\alpha & \alpha \end{pmatrix} \\ 2 & \end{pmatrix} \quad \left. \right\} \text{Add rows to 1 So yes, this is MC!}$$

## Example 2:

$$S = \{1, 2, 3\}$$

where  
 1 = Cheerful  
 2 = So-so  
 3 = Glum

$$P = \begin{array}{c|ccc} & 1 & 2 & 3 \\ \hline 1 & 0.5 & 0.4 & 0.1 \\ 2 & 0.3 & 0.4 & 0.3 \\ 3 & 0.2 & 0.3 & 0.5 \end{array} \quad \left. \begin{array}{l} \text{Rows add to 1,} \\ \text{so yes, } \{x_t, t \geq 0\} \text{ is a MC.} \end{array} \right.$$

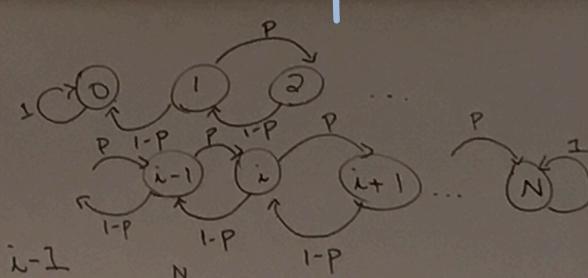
## Example 3:

$$S = \{0, 1, \dots, N\}$$

$$P_{ij} = \begin{cases} 1 & \text{if } i=j=0 \quad (\text{Goes broke}) \\ 1 & \text{if } i=j=N \quad (\text{Attains fortune of } \$N) \\ 1-p & \text{if } i \neq 0, N; j=i-1 \quad (\text{Loses}) \\ p & \text{if } i \neq 0, N; j=i+1 \quad (\text{Wins}) \\ 0 & \text{otherwise} \end{cases}$$

State transition diagram

$$\text{Ex. 3)} \quad S = \{0, 1, \dots, N\}$$

$$P_{ij} = \begin{cases} 1, & i=j=0 \\ 1, & i=j=N \\ 1-p, & i \neq 0, N; j=i-1 \\ p, & i \neq 0, N; j=i+1 \\ 0, & \text{o.w.} \end{cases}$$


$$\sum_{j=0}^N P_{ij} = p + 1-p = 1 \quad \text{for } i \neq 0, N$$

$$\sum_{j=0}^N P_{i,j} = 1 \quad \text{for } i=0, N$$

# Stationary Distribution of a DTMC

$$\pi_j := P(X_t = j)$$

= Long run probability that your stochastic process visits state  $j$ .

It can be calculated using the equation:

$$\pi_j = \sum_{i \in S} \pi_i p_{ij}$$

This means we can solve  $\pi = \pi Q$  to get the stationary distribution  $\pi$ .

[Markov Chain Monte Carlo]

## MCMC Basics

- \* We will start w/ some proposal  $g_i(\theta)$
- \* Subsequently update at each step  $t$  (usage  $g_t(\theta)$ )
  - ↳ The way to do this is to simulate from a transition probability distribution  $g_t(\theta) = T_t(\theta^t | \theta^{t-1})$  so that the standard distribution of the Markov Chain with transition probability matrix  $T_t$  is exactly  $P(\theta|y)$ .
- \* Key take away: If you sample long enough, eventually we'll sample from the Posterior! :)
- \* An (unfortunate) take away is that the first few (maybe thousand) samples could be complete garbage.
- \* This is taken into account with a BurnIn parameter, which specifies how many samples should be tossed out.

- \* Now our goal is to come up with smart  $\pi_t$ 's so that we get  $\pi(\theta|y)$  as our stationary distribution.

Three main algorithms to do this:

### 1) Gibbs Sampler

- !⇒
- most efficient
  - requires knowing the conditional distribution of  $\theta_t | \theta^{t-1}$  which is rarely ever known ;;
  - will likely never use

### 2) Metropolis Algorithm

- 2nd most efficient
- requires a symmetric proposal distribution, which we can always come up with ;;
- will probably use the most often

### 3) Metropolis-Hastings Algorithm

- slowest
- requires no information
- generalizes the Metropolis algorithm
- It can readily be used, but typically stick w/ Metropolis.

$$\begin{pmatrix} \theta_1 \\ \vdots \\ \theta_d \end{pmatrix} \rightarrow \pi(\theta_1, \dots, \theta_d | y)$$

↑ there is dependence among parameters

A valid Symmetric Proposal:

$J(\theta^t | \theta^{t-1}) = \text{Density of } N(\theta^{t-1}, 1)$  → Any Normal distribution w/ mean  $\theta^{t-1}$  works! (x)

$$f(\theta_a | \theta_b) = \frac{1}{\sqrt{2\pi}} e^{-\frac{(\theta_a - \theta_b)^2}{2}} \stackrel{?}{=} f(\theta_b | \theta_a)$$