

Lecture with Computer Exercises: Modelling and Simulating Social Systems with MATLAB

Project Report

Stable Marriage Problem

Valentin Junet & Samuel Imfeld

Zurich December 2014

Agreement for free-download

We hereby agree to make our source code for this project freely available for download from the web pages of the SOMS chair. Furthermore, we assure that all source code is written by ourselves and is not violating any copyright restrictions.

Valentin Junet

Samuel Imfeld

5. Durcred



Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich

Declaration of originality

The signed declaration of originality is a component of every semester paper, Bachelor's thesis, Master's thesis and any other degree paper undertaken during the course of studies, including the respective electronic versions.

Lecturers may also require a declaration of originality for other written papers compiled for their courses.

I hereby confirm that I am the sole author of the written work here enclosed and that I have compiled it in my own words. Parts excepted are corrections of form and content by the supervisor.

Title of work (in block letters):

Stable Marriage Problem

Authored by (in block letters):

For papers written by groups the names of all authors are required.

Name(s):	First name(s):	
Junet	Valentin	
Imfeld	Samuel	

With my signature I confirm that

- I have committed none of the forms of plagiarism described in the '<u>Citation etiquette</u>' information sheet
- I have documented all methods, data and processes truthfully.
- I have not manipulated any data.
- I have mentioned all persons who were significant facilitators of the work.

I am aware that the work may be screened electronically for plagiarism.

Place, date Zunch, 40.12.2014	Signature(s)

For papers written by groups the names of all authors are required. Their signatures collectively guarantee the entire content of the written paper.

Contents

1	Abs	stract		1				
2	Ind	Individual contributions						
3	Intr	oducti	ion and Motivations	1				
	3.1	Resear	rch Question	2				
4	Des	Description of the Model						
	4.1	Gale a	and Shapley's algorithm	2				
	4.2	Modif	ications	3				
5	Imp	olemen	tation	4				
	5.1	Conve	entions	4				
	5.2	Gener	rating the Preference Lists	5				
		5.2.1	Random	5				
		5.2.2	The Plane Algorithm	5				
	5.3	Makin	ng Matches	5				
	5.4	Check	ing the Engagements	7				
	5.5	Simula	ation	8				
6	Sim	ulatio	n Results and Discussion	9				
	6.1	Instab	oilities	9				
		6.1.1	Parameters	9				
		6.1.2	Effects of preference changes	9				
		6.1.3	Effects of visibility radius	10				
		6.1.4	Dependence on who changes preferences	10				
	6.2	Single	s	12				
	6.3	6.3 Optimality Index						
		6.3.1	Parameters	13				
		6.3.2	Effects of preference changes	14				
		6.3.3	Effects of visibility radius					
		6.3.4	Input size dependency	14				
7	Sun	nmary	and Outlook	16				

8	References	17
9	Appendix A: MATLAB Codes	17
10	Appendix B: Plots	38
	10.1 Optimality Index	- 38

1 Abstract

Online dating or match-making websites are flourishing these days. More and more people rely on their algorithms when searching for their Mr. Right, or Mrs. Right, respectively. Algorithms for match-making are therefore of quite some interest.

An important result in this concern is the so-called Gale Shapley algorithm proposed by Gale and Shapley [1962]. In 2012, Shapley even received the Nobel Prize in Economics for his work. The goal of this paper is to discuss the original model described by Gale and Shapley [1962] and advance the model in some sense. Namely we are going to introduce two changes to the model: The incompleteness of information and the possibility of preference changes. The modifications will be discussed further in section 4.2.

It is however not our claim that these changes applied to the model will make it an exact description of reality. Our goal is to study the repercussions on the stability and other significant indicators that show up when applying the modifications.

2 Individual contributions

Code was written by both. In the report, Valentin focused on the chapters 3, 4, 6.1 and 6.2 whereas Samuel concentrated on chapters 1, 5 and 6.3.

3 Introduction and Motivations

The stable marriage problem is quite well known; it is originally an interesting mathematical problem, but it also has something catchy: Everyone has his own idea of what stability is, how we think it is optimal, and how we would like to be matched, or how we think it should be done. Then there comes this emotionless algorithm that pretends – with the support of some definitions and mathematical proofs – to find all of this for us and to do it as least as good as we would. Of course this is a little exaggerated. But when we think about those internet sites that match people according to some criteria, is't it what we're doing? Letting some numbers find an objective answer to those questions, to our subjective questions?

In order to prevent misunderstandings one must very precisely answer the following questions: What is stability? How can a stable situation be optimal? Of course answering those questions would result in a rather philosophical discussion which is not the goal of our paper. Nevertheless stability and other related indicators (which were going to introduce soon enough) will be the center of our attention. But in order to reduce complexity will only be considering some very abstract and theoretical nodes which, for some practical reasons, will be called men or women.

The key points of our work presented here are to generate random links (or "friendship") between men and women (which in the following discussions will be referred to by nodes), to classify those links in a random way in order to get a preference list and to apply the algorithm of Gale and Shapley which matches the nodes according to their preferences. Then we modify some parameters and interpret the results. For a more precise description of the procedure, you may refer to the next section.

3.1 Research Question

We want to analyze the influence that the modifications we are going to apply (see section 4.2) has on the simulation result, i. e. the indicators for stability and optimality (see sections 4.2 and 5.5 for detailed explanation). In particular the questions are:

For stability

- How is the stability influenced when the visibility radius becomes smaller (and therefore also the preference lists)?
- Are there more instabilities if the preference changes are more frequent?
- Does it matter whether the men or the women change preferences?

We expect that there are more instabilities for a higher preference change rate. We also think it will make a difference when women change preferences instead of men because the algorithm is not symmetric.

For optimality

- How is the optimality influenced when the visibility radius becomes smaller?
- Is the output less optimal if the preference changes are more frequent?
- Does the optimality depend on the input size?

Here we expect to see the following results: The result is less optimal for a smaller visibility radius because then it is more likely that nodes have no partner (which is considered as suboptimal, see section 4.2). It is also less optimal for a higher preference change rate.

4 Description of the Model

4.1 Gale and Shapley's algorithm

Let's first describe the algorithm of Gale and Shapley. The setting is of a very abstract nature. There is a certain number of nodes, half of them are men and the other half are

women. The first step is that every node makes a preference list, i. e. a ranking of all the nodes of oppposite sex. Then the algorithm proceeds as follows:

- 1. single men propose to their favorite women
- 2. women accept their favorite suitor or their only suitor
 - if a woman has a fiance and a suitor, she chooses the one she likes better
 - if a woman rejects a man, then the next woman on the rejected man's list is his new favorite
- 3. if the single men still have women on their list of preferences, go to step 1

So the algorithm terminates when all nodes have a partner. This is obviously always achieved if the length of all the preference lists corresponds to the number of nodes of opposite sex (see Gale and Shapley [1962, p. 14, Theorem 1]). Gale and Shapley also proved that this algorithm always terminates with an optimal stability. For this purpose they gave precise definitions of stability and optimality:

Definition 1: An assignment of applicants to colleges will be called unstable if there are two applicants α and β who are assigned to colleges A and B, respectively, although β prefers A to B and A prefers β to α .¹

Definition 2: A stable assignment is called optimal if every applicant is at least as well off under it as under any other assignment.²

These definitions are in terms of colleges and applicants because in their paper they first start with this setting and then go on to the special case of stable marriages. Translated to the words of match-making the definitions would be something like: "A set of marriages is called unstable if under it there are a man and a woman who are not married to each other but prefer each other to their actual mates" and "A stable assignment is called optimal if every man is at least as well off under it as under any other assignment".

4.2 Modifications

We will now present in more detail the modifications we made to this algorithm. Our purpose was to modify and adapt it so that it represents more accurately the reality. Of course it is fundamentally impossible to make a perfect representation of reality through this: If for example a man goes into a bar, he doesn't particularly make a list of the

¹Gale and Shapley [1962, p. 10]

²Gale and Shapley [1962, p. 10]

³Gale and Shapley [1962, p. 11]

women he prefers and propose to each woman according to his list (or at least this is not systematically observed). However, there are some modifications we can make to make it more realistic. One of these is that a man doesn't necessarily know every women (and vice versa). To improve this point, we modified the list of preferences so that men and women only know each other in a restricted area. For example, a man who lives in China doesn't know a woman who lives in Switzerland, so he won't propose to her (we didn't consider the case where the man in China could know the Swiss woman though internet or by traveling, he only knows the women "around" him). We also considered the case in which a man knows a woman in his area but she doesn't know him and then this man can propose to her and she might accept or not. The other modification we made is that people, obviously, can change their minds. We modified it so that during the process either only the men, only the women or both can change their mind. These changes of preferences are made each "round" (each time a single man proposes to a woman) with a certain probability which is fixed.

We also introduced some new concepts, namely the number of 'singles" and "'dumps" and the "optimality index". The number of singles tells us how many nodes didn't find a partner. Since we chose to do the simulation with the same number of men and women, there will be exactly the same number of single men and women. The number of "dumps" is the number of times a woman rejected her fiance for another man. This can be seen has an indication of time, the more people "get dumped", the longer it takes to reach a final assignment. The optimality index is defined as the sum over all ranks of the actual partners in the preference list of their respective partners, divided by two times the number of men/women squared (resulting in a normalization, e. g. it only takes values between zero and one). If a node has no partner, it counts as the maximal rank because having no partner is certainly not optimal. The optimality index should somehow give information about the optimality of the outcome, although it does not correspond to definition 2 in section 4.1. A low optimality index corresponds to a rather optimal assignment whereas a high optimality index indicates a not so optimal assignment.

5 Implementation

5.1 Conventions

The nodes (men resp. women) are referenced by integers from one to n where n is the input size (number of men/women).

Accordingly, a preference list is a permutation of the first n integers: The element $\sigma(1)$ is the most preferred node, $\sigma(2)$ the second most preferred and so on. In the case of non-complete information, there are m < n distinct integers followed by zeros to complete the sequence.

5.2 Generating the Preference Lists

5.2.1 Random

For testing purposes we first coded a generator for random preference matrices which basically calls randperm(n) n times, resulting in a matrix whose rows represent the preference list of a single node. The generated matrix can be used to simulate the classic case with complete information.

5.2.2 The Plane Algorithm

The question that now arises is the following: How can I generate a preference matrix with non-complete information that represents reality in an appropriate manner? In particular there should be a mechanism for controlling the number of non-zero entries in the preference lists (i. e. controlling the degree of completeness of information). The number of non-zero elements should, however, not be the same for each node (this would mean that each node knows the same fixed number of people, which is clearly not a good representation of real situations). Additionally, situations like 'man x knows woman y but woman y does not know man x' should be considered. To achieve these characteristics we developed the 'generatePlane' algorithm: It is based on the idea that a node only knows his neighbours (e. g. the people in his town or community). Therefore his preference list should consist only of nodes that are closer than a certain distance, which we will call the visibility radius in the further discussions. We realized this by assigning to each node a random position in a two dimensional plane. For simplicity they are distributed in $[0,1] \times [0,1]$.

Definition 3: The visibility radius defines the neighbourhood of a node in $[0,1] \times [0,1]$. A node can only know other nodes that have a distance smaller than the visibility radius.

To avoid border effects, the edges are connected (one could view it as a torus). For the actual generation of the preference lists one just has to iterate through all nodes and determine all nodes of opposite sex that are in the disc of visibility radius. We then chose to use a random permutation of these neighbourhood nodes to keep it simple.

A constant visibility radius implies in particular that if node x knows node y, then also node y knows node x. But as we don't want this to be the case all the time we also made an option for a random visibility radius that is updated every step.

5.3 Making Matches

This is the part where the algorithm proposed by Gale and Shapley [1962] actually comes into play. An implementation in MATLAB is shown in listing 1 (The basic ideas for the implementation are taken from RosettaCode [2014] and adapted to MATLAB).

```
1 function [ engaged, stable ] = makeMatch( m, f )
   %makeMatch finds engagements for preferences according to Gale-Shapley ...
       algorithm
3
4 freemen = [(1:n)', ones(n,1)];
   engaged = zeros(n, 2);
5
   while ~isempty(find(freemen(:,2)==1,1))
6
       theman = find(freemen(:,2) == 1,1);
       thegirl = m(theman,1);
8
       index = find(engaged(:,2) == thegirl,1);
9
       if(isempty(index))
10
11
            engaged(theman, 1) = theman;
12
            engaged(theman, 2) = thegirl;
            freemen(theman, 2) = 0;
13
       else
14
            fiance = engaged(index,1);
15
            girlprefers = f(thegirl,:);
16
            if(find(girlprefers==theman,1)<find(girlprefers==fiance,1))</pre>
17
                engaged(theman,1) = theman;
18
                engaged(theman, 2) = thegirl;
19
                engaged(fiance, 1) = 0;
20
                engaged(fiance, 2) = 0;
21
                freemen (theman, 2) = 0;
22
23
                freemen(fiance, 2) = 1;
24
            else
25
                m(theman,:) = [m(theman,2:n) 0];
26
            end
27
       end
   end
28
29
   stable = checkEngagements(engaged, m, f);
31
  end
```

Listing 1: makeMatch code skeleton

The main structures are the arrays freemen and engaged. The array freemen contains all free men (first column: indices of men, second column: 1 for free, 0 for not free) and engaged contains the engagements produced by the algorithm (first column: indices of men, second column: indices of women). The main loop starts in line 6 and continues until all men are in an engagement. Then in the loop the first free man is picked (line 7) and his most preferred girl is determined (line 8). Now one has to distinguish between two cases:

- The girl has no engagement: Everything is fine, the man and the girl are now engaged (lines 11 and 12). Also the man is not free anymore (line 13).
- The girl alread has a fiance: In this case one has to do another distinction:
 - The girl prefers the new man to her fiance: A new engagement is made and the old one is cancelled (lines 18-21). One also has to update freemen.

- The engagements remain unchanged, but the girl is removed from the man's preference list because she is not attainable for him.

In the end the engagements are checked with the checkEngagements algorithm described in section 5.4. However, when applying the modifications described in section 4.2 and still using this algorithm one runs into problems (as expected). Therefore we had to adapt the makeMatch algorithm to be able to handle the following situations:

- An unknown man proposes to a woman who is not engaged: We decided that in this case the proposing man should have a chance to succeed with his proposal, but this should not always be the case. Therefore we implemented a random decision with a certain probability for accepting (typically around 0.25). This can be seen as a simulation of the real-life situation 'the woman gets to know the unknown man and gets to like him (or not)'.
- An unknown node proposes to a woman who is engaged: We decided to apply the same procedure as above. One could argue that the probability for accepting should be lower because she already has a fiance, but we left that out for the sake of simplicity.
- The preference list of a man is empty but he is not engaged: In this case the man is just left with no partner.

The implementation of the above points will not be discussed here any further. The final algorithm is in the appendix for reference.

Finally we added the preference change functionality: In each iteration step a random decision is made between changing preferences or not, using the probability given in the additional parameter changerate. If the answer is positive, the preference list of one randomly chosen node is perturbed a little: A random node in the preference list is chosen and switched with the node one rank lower (for example [4,2,3,1] becomes [4,3,2,1] after switching nodes 2 and 3). There is also a parameter to determine whether only men, only women or both should change their preferences.

5.4 Checking the Engagements

An important indicator for the later discussion is the stability of the engagements. It can be checked using this algorithm. The main loop is shown in listing 2 (again the basic structure is inspired by RosettaCode [2014])

```
while he<=n
she = engaged(he,2);</pre>
```

```
% make sure that the man has a partner i. e. she is not 0
       % otherwise continue because nothing to check
4
5
       hisindex = find(f(she,:) == he, 1);
6
       herindex = find(m(he,:) == she, 1);
       helikesbetter = m(he,1:herindex);
8
9
       shelikesbetter = f(she,1:hisindex);
       % check for her
10
       for i=1:size(shelikesbetter)
11
           guy = shelikesbetter(i);
12
13
           guysgirl = engaged(guy,2);
           guylikes = m(guy,:);
15
           if (find(guylikes==she,1)<find(guylikes==guysgirl,1))</pre>
                stable = false;
16
           end
17
       end
18
       % check for him
19
20
21
       he=he+1;
22 end
```

Listing 2: checkEngagements code fragment

The main loop is an iteration over all men (line 1), but only those who are engaged because having no partner is not considered as an instability (lines 3-5). After having retrieved all the indices and preference lists, one iterates over all nodes that appear before the actual partner in their respective preference lists and checks whether there is an instability (see definition 1 in section 4.1). This results in two for loops, one for the man and one for the woman he is engaged to.

5.5 Simulation

The actual simulation makes iterated calls to generatePlane and makeMatch to simulate the modified Gale-Shapley Algorithm for different parameter settings. The input parameters are:

- n: the input size $n \in \mathbb{N}$
- mode, radius: mode determines the choice of the visibility radius (either random or constant), and if the radius is constant then it can be set with the parameter radius $r \in [0, 0.5]$
- changerate: the rate at which preference changes are performed $c \in [0,1]$
- p: determines whose preferences are changed $p \in \{0, 0.5, 1\}$

The resulting output variables are:

- no. of instabilities: number of instabilities according to definition 1 in section 4.1
- no. of dumps: number of fiances that have been dumped during match-making, see section 4.2
- no. of singles: number of single nodes that remain after match-making
- optimality index: as defined above, see section 4.2

6 Simulation Results and Discussion

Let's first make a general remark for the interpretation of the results. We labelled our nodes – like in the original algorithm – with man and woman, but a more accurate notation would be "the ones who choose" for the men and "the ones who are chosen" for the women. For the sake of interpretation, this indeed makes more sense since: In real life, when someone asks a person if he or she wants to go out with him/her, the one asking chooses someone that he/she prefers. The person who is asked makes his/her decision based on his/her preferences (at least according to this model). So we could do exactly the same by switching or even mixing both categories and find through our interpretation an even more realistic information about the difference in the results between choosing and being chosen.

It is also important to notice that every kind of interpretation with respect to reality is purely hypothetical. This is not a realistic model; it only shows us how this specific way of choosing and being chosen responds to the input parameters.

6.1 Instabilities

6.1.1 Parameters

For the considerations about the instability and the number of singles we did the following simulation: Generate a matrix of 2^n persons with n in 1:11 and either a random radius or a constant radius r in $\{0.1, 0.25, 0.4\}$. To make the matches, we chose a probability p – for the change of preferences – in 0:0.25:1. We did this with the change of preferences only for men, only for women and for both. We iterated this simulation 11 times and then took the algebraic mean value as data.

6.1.2 Effects of preference changes

We will start by some considerations about the number of instable (as defined by Shapley and Gale) couples counted at the end of the process. We first note, that if the probability of changing preferences is 0, then there is like in the original algorithm no instability for

any kind of matrices. This would mean that the algorithm still holds with a restricted knowledge of the others. If we think of the plane as small or big cities, then the persons who follow this algorithm without changing their mind will be able to find a stable love regardless the size of their respective city.

We can also quickly mention that the bigger the probability that a change occurs is, the bigger is the number of instabilities (for both, men and women changing their mind), so when people are indecisive and change their mind, the marriages are less stable.

6.1.3 Effects of visibility radius

We found out an interesting relation between the visibility radius and the number of instabilities. In general (meaning for every kind and probability in the change of preferences and every size of matrices) we found out that for the "big" radius 0.25 and 0.4, there is significantly less instability than for the 0.1 radius. Actually, for the two big radii, the mean number of instability is always less than two. For example for 2¹¹ persons, the mean number of instability (w.r.t. all the probabilities) is 12.236 for r=0.1, 0.491 for r=0.25 and 0.218 for r=0.4. This would mean that the people with restricted possibilities who change their mind are more likely to end up in an unstable union than the people who have more choice.

6.1.4 Dependence on who changes preferences

Let's now consider the influence of whether only the men or only the women change their mind on the instability of the marriages. We noticed that for constants radius there is always more instability when the men change their preferences but, as mentioned in last paragraph, the number of instability for r=0.25 and 0.4 is so low that this difference, although present, is negligible. Except randomness, we cannot find an explanation for this in the way we made the changes in the algorithm since the change of preferences are made at the beginning of each main loop regardless if the change is made by a man or a woman. So this means that, up to randomness, the explanation should lay in the way of choosing and making matches. With this algorithm, once a man is matched, there nothing we can do to change it, so if he changes his mind and the woman he now likes better likes him better too, theres nothing to do about it, it will be unstable, whereas if a woman changes her mind then the man she likes better might propose to her and even if she is engaged she can be with him. This means then that a person should make sure to be able to leave his partner if he changes his mind...not very surprising. However, what is very curious is that if we consider this relation not with a constant radius but rather with a random radius, then – from a certain amount of people (2^7) on – the opposite is true: there is more instability when the women change their preferences! With our previous considerations it is clear that, up to randomness, the explanation should be find in the construction of the

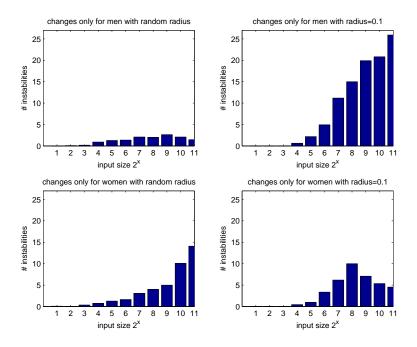


Figure 1: Number of instabilities in function of the number of persons, for a probability in the change of preferences fixed at 100% and for different radius.

plane. The fact that the radius are random influences the men and the women similarly, so the true difference (and the reason we introduced it) between constant and random radius is that, with the random one, a man can know someone who doesn't know him and propose to her and respectively a woman can be proposed to by someone she doesn't know (in this case, we arbitrarily decided that she accepts 25% of the time). As we can see in the graphs, in this case the undecided women don't produce particularly more instabilities, but the undecided however produce less instability than with a constant radius. This might either be caused by the fact that a random radius can be bigger than 0.1 and as we've seen in this case there is less instabilities, or because if they like better a woman who doesnt know him, then this wont be unstable. We can then suppose that with any radius the most influent factor of instability is the indecision of the men, the one who choses and can't break a union.

We also noticed that in the case in which both, the men and the women change their minds we observe that the results are approximately the mean between the ones in which only the men and the ones in which only the women change their mind.

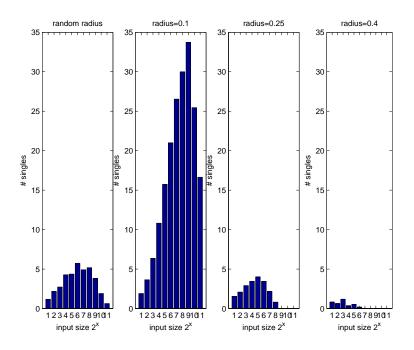


Figure 2: Number of singles in function of the number of person, for different radii.

6.2 Singles

We now want to consider how this variances of the algorithm influence the number of singles person. However let's first start by mentioning what does not have any effects on this number: The change of preferences, more precisely its rate and whether only men, women or both change them. Indeed, for any fixed radius and fixed number of persons in our simulation, we observe the same results by varying the rate of change and the group of people changing their minds.

The influent factors are then the number of people and the radius. Looking at figure 2 we can easily notice, that with the number of people increasing, the number of singles has a normal distribution. We performed a Kolmogorow-Smirnow-Test for all data on the number of singles in function of the number of persons. The test was always positive. We can observe that for the small radius 0.1, there is a lot more of singles men and women than for the others, actually for radius 0.4, it is even unusual to find a single pair. For radius 0.4, resp. radius 0.25 from 2⁷ to 2⁹ persons there isn't any single person left! So the more people know each other, the less singles there are. The random radius gives something intermediate, there is though more singles than for the radius 0.4 and 0.25 since the radius

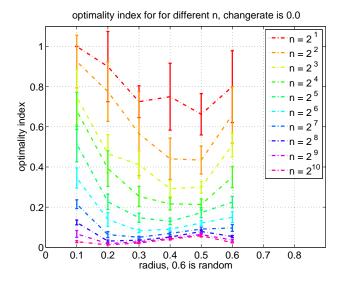


Figure 3: Errorbar plot for the optimality index in dependence of the visibility radius. The different lines correspond to different input sizes n. The preference change rate was set to 0.

0.1 has way more influence on the number of singles, resp. persons there isnt any single So the more people know each other, the less single there is. The random radius gives something intermediate, although there are more single persons than for the radius 0.4 and 0.25 since the radius 0.1 has way more influence on the number of singles.

6.3 Optimality Index

6.3.1 Parameters

The input parameters were chosen as follows: The input size is a power of 2. Due to the quadratic dependence of the generatePlane algorithm, the highest power evaluated is $2^{10} = 1024$. Originally we wanted to work with basis 10 but we had to change to basis 2 because of performance reasons as well. For the visibility radius we took the discrete values 0.1:0.1:0.5, and for the preference change rates values in 0:0.5:1. For each triple of parameters the simulation was performed 10 times.

6.3.2 Effects of preference changes

Surprisingly the expected effect of the preference changes could not be observed at all. The optimality index is about the same for the extreme values 0 and 1 for the change rate, and also for 0.5. For small input sizes there is a small variation of the optimality index (that seems to be caused more by random effects than an actual correlation). For higher n the plot lines are at constant height (see also section 10.1 in the Appendix for plots).

6.3.3 Effects of visibility radius

The results of the visibility radius variation do also not look very promising. In figure 3 one can see that for low input sizes (n = 2 to 32) there seems to be a dependence on the visibility radius, namely the optimality index decreases for higher visibility radii (as we expected). But for higher n the effect disappears, and the optimality index even gets higher again for the visibility radius close to 0.5. Note that the points labeled with 0.6 on the x axis correspond to random radius and not visibility radius 0.6.

One way to explain this is that there might be several 'hidden' effects that cause this weird looking behavior. One of them could be that the visibility radius is chosen to be constant for all input sizes. It is now not the same when thousand instead of only ten nodes are placed in the plane, because for more nodes the number of nodes that happen to be in the neighborhood defined by a constant visibility radius is much higher.

6.3.4 Input size dependency

Looking at figure 4, one can clearly observe that the optimality index gets smaller for increasing input size, meaning that the result is more optimal for higher n. This is a bit counterintuitive at first thought. Why should the result be less optimal for smaller n? One could think that this effect is somehow produced by the modifications that were made to the Gale-Shapley algorithm. But it turns out that the effect can also be observed for randomly generated preference lists using generateRandom, see figure 5. For this type of preference lists none of the 'special treatments' in the modified algorithm (e. g. empty preference lists) apply, but the effect can still be observed. Onecan also observe that the higher the visibility radius gets, the closer is the curve to the one with randomly generated preferences.

So one might think that there is some power law relating the input size to the optimality index. Conducting an analysis (using loglog plot and polyfit for linear regression) yields

$$opt = c \cdot n^{-0.36}$$

But how can this relation be explained? We suspect that it has to do with the following: The optimality index measures the rank of the assigned partners. If now the number of

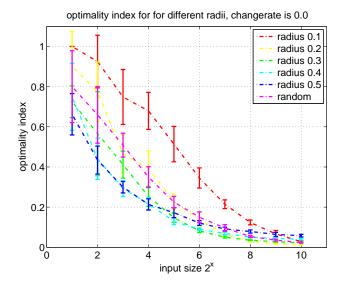


Figure 4: Optimality index errorbar plot in dependence of input size n. Preferences were generated by the generatePlane algorithm. Different lines represent different visibility radii.

nodes is high, the chance that for example two women choose the same man as their number one preference is higher than it would be for lower input sizes. But as only one of those women gets assigned to this man, the other women gets another man who is now at most second best on her list.

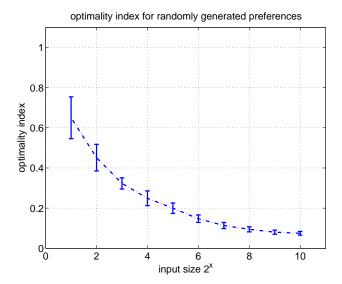


Figure 5: Optimality index errorbar plot in dependence of input size n. Preferences were generated randomly.

7 Summary and Outlook

The important points related to the stability of the modified model are the following:

- Without change of preferences stability still holds.
- For the "big" radii 0.25 and 0.4 there is practically no instability.
- The men who change their mind induce more instabilities than the women.

The important points related to the number of singles of the modified model are the following:

- The change of preferences has no influences on the number of single persons.
- The small radius 0.1 increases the number of single persons.
- When the input size is big enough, there are no single persons.

And for the optimality index:

• Neither change of preferences nor non-complete information through the visibility radius has an effect on stability

• The stability is related to the input size by a power law

8 References

D. Gale and L. S. Shapley. College admissions and the stability of marriage. *The American Mathematical Monthly*, 69(1):pp. 9–15, 1962. ISSN 00029890. URL http://www.jstor.org/stable/2312726.

Youtube Numberphile. Stable marriage problem, 2014a. URL https://www.youtube.com/watch?v=Qcv1IqHWAzg.

Youtube Numberphile. Stable marriage problem, 2014b. URL https://www.youtube.com/watch?v=LtTV6rIxhdo.

RosettaCode. Stable marriage problem, 2014. URL http://rosettacode.org/wiki/Stable_marriage_problem#Python.

9 Appendix A: MATLAB Codes

generateRandom.m

```
function [ m, f ] = generateRandom( n )
generateRandom generates random preference matrices
m = zeros(n,n);
f = zeros(n,n);
for i=1:n
m(i,:) = randperm(n,n);
f(i,:) = randperm(n,n);
end
```

generatePlane.m

```
1 function [ mpref,fpref ] = generatePlane( n ,mode, radius)
2 %GENERATEPLANE generates preference lists for men and women
3 % based on a plane where women and men are represented by points
```

```
4 %
       they have a limited visibility radius
       n: number of men and women
   00
5
   양
       mode: visibility radius mode, optional argument
6
7
   응
         1 --> const, one constant radius for all nodes
   응
         2 --> random, a new random radius is generated in each iteration
               value is between 0.1 and 0.5
10
   9
       default mode is const
       mpref: mens preferences in nxn matrix
   응
11
12 %
       fpref: womens preferences in nxn matrix
13
14 global verbosity
16 if (nargin >= 2 && mode == 1)
       assert (nargin==3);
17
       r = radius;
18
19 end
_{20} if(nargin < 2)
       mode = 1;
22
       r = 0.2; %default value
23 end
24
25 % generate random coordinates
26 % and extend to torus
27 men = zeros(3,9*n);
28 rnd = rand(2,n);
29 men(:, (0*n)+1:1*n) = [(1:n); rnd];
30 men(:,(1*n)+1:2*n)=men(:,(0*n)+1:1*n)+[zeros(1,n);ones(1,n);zeros(1,n)];
31 men(:,(2*n)+1:3*n)=men(:,(0*n)+1:1*n)+[zeros(1,n);ones(1,n);ones(1,n)];
32 \text{ men}(:, (3*n)+1:4*n) = \text{men}(:, (0*n)+1:1*n) + [zeros(1,n); zeros(1,n); ones(1,n)];
33 men(:, (4*n)+1:5*n) = men(:, (0*n)+1:1*n) + [zeros(1,n); -ones(1,n); ones(1,n)];
34 men(:, (5*n)+1:6*n) = men(:, (0*n)+1:1*n) + [zeros(1,n); -ones(1,n); zeros(1,n)];
35 men(:, (6*n)+1:7*n) = men(:, (0*n)+1:1*n) + [zeros(1,n); -ones(1,n); -ones(1,n)];
36 men(:, (7*n)+1:8*n) = men(:, (0*n)+1:1*n) + [zeros(1,n); zeros(1,n); -ones(1,n)];
37 men(:,(8*n)+1:9*n) = men(:,(0*n)+1:1*n)+[zeros(1,n);ones(1,n);-ones(1,n)];
38
39 women = zeros(3,9*n);
40 rnd = rand(2, n);
41 women(:, (0*n)+1:1*n) = [(1:n); rnd];
42 women(:, (1*n)+1:2*n) =women(:, (0*n)+1:1*n) +[zeros(1,n);ones(1,n);zeros(1,n)];
43 women (:, (2*n)+1:3*n) =women (:, (0*n)+1:1*n)+[zeros(1,n);ones(1,n);ones(1,n)];
44 women(:, (3*n)+1:4*n) =women(:, (0*n)+1:1*n) +[zeros(1,n);zeros(1,n);ones(1,n)];
45 women(:, (4*n)+1:5*n) =women(:, (0*n)+1:1*n) +[zeros(1,n);-ones(1,n); ones(1,n)];
46 women(:, (5*n)+1:6*n)=women(:, (0*n)+1:1*n)+[zeros(1,n);-ones(1,n);zeros(1,n)];
47 women(:, (6*n)+1:7*n) = women(:, (0*n)+1:1*n) + [zeros(1,n); -ones(1,n); -ones(1,n)];
   women (:, (7*n)+1:8*n) = women (:, (0*n)+1:1*n) + [zeros(1,n); zeros(1,n); -ones(1,n)];
   women (:, (8*n)+1:9*n) =women (:, (0*n)+1:1*n)+[zeros(1,n);ones(1,n);-ones(1,n)];
49
50
51 %plotting
52 % if verbosity~=0
         plot (men(2,1:n), men(3,1:n), 'o', women(2,1:n), women(3,1:n), 'o');
```

```
54 %
          label1 = cellstr( num2str(women(1,1:n)') );
          label2 = cellstr( num2str(men(1,1:n)') );
55 %
          text (women (2,1:n), women (3,1:n), label1);
56
          text (men(2,1:n), men(3,1:n), label2);
57 %
          title('nodes in plane');
          legend('men','women');
59
60 % end
61
62 d = zeros(2, 9*n);
63 mpref = zeros(n,n);
64 fpref = zeros(n,n);
66
67 for i=1:n
        man = men(:,i);
68
        for j=1:9*n
69
            woman = women(:,j);
70
71
            d(:,j) = [woman(1,1); norm(man(2:3)-woman(2:3),2)];
72
73
        if mode==2
            r = rand*0.4+0.1;
74
75
        end
        index = find(d(2,:) < r);
76
77
        available = women(:,index);
        sz = size(available, 2);
79
        if sz>n
            available = available(:,1:n);
80
            sz = n;
81
        end
82
83
        perm = randperm(sz);
        mpref(i,1:sz) = available(1,perm);
85 end
   for i=1:n
86
        woman = women(:,i);
87
        for j=1:9*n
88
89
            man = men(:,j);
            d(:,j) = [man(1,1); norm(man(2:3)-woman(2:3),2)];
90
91
        end
92
        if mode==2
            r = rand*0.4+0.1;
93
        end
94
        index = find(d(2,:)<r);
95
        available = men(:,index);
96
        sz = size(available, 2);
97
98
            available = available(:,1:n);
99
            sz = n;
100
        end
101
        perm = randperm(sz);
102
        fpref(i,1:sz) = available(1,perm);
103
```

```
104 end
105 end
```

vprintf.m

```
function vprintf(varargin)
function vprintf
```

makeMatch.m

```
1 function [ engaged, output ] = makeMatch( m, f, changerate, p )
\mathbf{2} %makeMatch finds engagements for preferences according to Gale—Shapley ...
      algorithm
      men an women encoded as integers from 1 to n
4 %
  양
      input:
      m: preference matrix of the men. Each row corresponds to a man and
      the elements are the women listed according to his preferences.
      f: preference matrix of the women. Each row corresponds to a woman and
10
      the elements are the men listed according to her preferences.
      changerate: rate at which preference changes are performed, e. g. if
11
      changerate=0.2 then only in 20% of iterations preferences are changed
12 %
13 %
      p: change preferences for men (p=1) / women (p=0.5) / both (p=0.5)
14 %
15 %
      dimensions must be correct, m=nxn, f=nxn.
16 %
17 %
      returns:
18 %
      engaged: nx2 Matrix containing matches
      output: output data -->
19 %
      output(1,1): number of instabilities
20 %
      output(1,2): number of singles
21 %
22 %
      output(1,3): number of dumps
23
      output(1,4): optimality index
24
25 % optional test prints
26 global verbosity
27 vprintf('mens preferences:\n');
28 if verbosity~=0 disp(m); end
29 vprintf('womens preferences:\n')
30 if verbosity~=0 disp(f); end
31 % assign local variables
```

```
32 initialm = m;
33 initialf = f;
34 n = size(m, 1);
35 n2 = size(f, 1);
36 % make sure dimensions agree
37 assert(n == size(m, 2));
38 assert (n==n2);
39 if nargin > 2
       assert (nargin==4);
40
       assert (changerate<=1);</pre>
41
42
       assert (changerate>=0);
       assert(^{\sim}isempty(find(p==[0,1,0.5],1)));
45 % more local variables
46 freemen = [(1:n)', ones(n,1)]; % column 1= men; column 2= 1 \rightarrow man is free, ...
       0 -> man isn't free
47 engaged = zeros(n,2);% column 1= men; column 2= women
48 dumped=0; % no of dumps
49 acceptrate = 0.75; % rate at which unknown nodes are accepted
50 % main loop
51 while ~isempty(find(freemen(:,2)==1,1)) % iterate as long as there are free men
       % preference changes
52
       if nargin > 2 % only if changerate and p are given
53
           if rand < changerate % change prefs?</pre>
54
               node = randi(n); % node whose prefs to change
56
               if rand 
                    pref = nonzeros(m(node,:))';
57
                    len = size(pref, 2);
58
                    if len>1
59
                        k = randi([2, len]); % where in pref to change
60
                        girl1 = pref(k); % the girl to swap
61
62
                        i1 = find(initialm(node,:) == girl1,1); % index of girl1 ...
                           in initialm
                        girl2 = m (node, k-1); % girl to be swapped with
63
                        i2 = find(initialm(node,:) == girl2,1); % index of girl2 ...
64
                            in initialm
                        initialm(node, i2) = girl1;
65
                        initialm(node, i1) = girl2;
                        m(node, i1) = girl2;
                        m (node, i1-1) = qirl1;
68
                    end
69
               else
70
                    pref = nonzeros(f(node,:))';
71
                    len = size(pref, 2);
72
                    if len>1
73
                        k = randi([2,len]); % where in pref to change
74
                        man1 = pref(k); % the man to swap
75
                        i1 = find(initialf(node,:) == man1,1); % index of man1 in ...
76
                            initialf
                        man2 = f(node, k-1); % man to be swapped with
77
```

```
i2 = find(initialf(node,:) == man2,1); % index of man2 in ...
78
                             initialf
                         initialf(node, i2) = man1;
79
                         initialf(node, i1) = man2;
80
                         f(node, i1) = man2;
                         f(node, i1-1) = man1;
82
                    end
83
                end %if_2
84
                vprintf('preferences changed\n');
85
            end %if_1
86
87
        end
88
        응 +++
89
        theman = find(freemen(:,2)==1,1); % the first man free on the list
        thegirl = m(theman,1); % his first choice
90
            if thegirl==0; % theman doesn't know any free girls who want him, ...
91
                he'll be alone : (
                freemen (theman, 2) = 0;
92
                engaged (theman,:) = 0;
            else
                index = find(engaged(:,2) == the girl, 1); % index of possible ...
95
                    fiance of his first choice
                if(isempty(index) ) % thegirl is free -> theman will be engaged ...
96
                    to thegirl
97
                    if isempty(find(f(thegirl,:)==theman,1))
                         vprintf('man %d proposed to women %d, she does not know ...
                             him\n', theman, thegirl);
                         if rand>acceptrate % man accepts with a certain rate
99
                             engaged(theman, 1) = theman; % make new engagement
100
                             engaged(theman, 2) = thegirl;
101
                             vprintf('she accepts\nman %d is engaged to girl ...
102
                                 d^n, theman, thegirl);
103
                             freemen(theman, 2) = 0; % man is not free anymore
                             f(thegirl,:) = [theman, f(thegirl,1:n-1)]; % update ...
104
                                 preferences
                             initialf(thegirl,:) = [theman, ...
105
                                 initialf(thegirl,1:n-1)]; % also in initial ...
                                 matrix (will be used for checking)
106
                         else
107
                             vprintf('she declines\n');
                             m(theman,:) = [m(theman,2:n) 0]; % make pref list ...
108
                                 of theman smaller
                         end % if_4
109
                    else
110
                         engaged (theman, 1) = theman; % make new engagement
111
                         engaged(theman,2) = thegirl;
112
                         vprintf('man %d is engaged to girl %d\n', theman, thegirl);
113
                         freemen(theman,2) = 0; % man is not free anymore
114
                    end % if_3
115
                else % thegirl is already engaged \rightarrow check if thegirl prefers ...
116
                    theman to her fiance
```

```
fiance = engaged(index,1); % her fiance
117
                    girlprefers = f(thegirl,:); % pref list of thegirl
118
                    howgirllikestheman=find(girlprefers==theman,1); % themans ...
119
                        index on thegirls preferences list
                    howgirllikesfiance=find(girlprefers==fiance,1); % fiances ...
120
                        index on thegirls preferences list
                    if(isempty(howgirllikestheman)) % thegirl doesn't know ...
121
                        theman -> thegirl accepts with a certain rate
                        if rand > 0.75
122
                             % thegirl prefers theman -> update pref list
123
                             f(thegirl,:) = [f(thegirl,1:howgirllikesfiance), ...
124
                                 theman, f(thegirl,howgirllikesfiance+1:n-1)];
                             initialf(thegirl,:) = ...
125
                                 [initialf(thegirl,1:howgirllikesfiance), ...
                                 theman, ...
                                 initialf(thegirl,howgirllikesfiance+1:n-1)]; % ...
                                 also initial
126
                        end % if_4
127
                    end % if_3
                    if(find(girlprefers==theman,1)<find(girlprefers==fiance,1)) ...</pre>
128
                        % thegirl prefers theman ->change engagement
                        engaged(theman,1) = theman; % change fiance of the girl
129
                        engaged(theman,2) = thegirl;
130
                        engaged(fiance,1) = 0; % fiance is free again
131
132
                        engaged(fiance, 2) = 0;
133
                        vprintf('girl %d dumped man %d for man %d\n', thegirl, ...
                            fiance, theman);
                        dumped=dumped+1;
134
                         freemen (theman, 2) = 0;
135
136
                        freemen(fiance, 2) = 1;
137
                    else
138
                        m(theman,:) = [m(theman,2:n) 0]; % thegirl prefers her ...
                            fiance -> take thegirl out of themans preference list
                    end % if_3
139
                end % if_2
140
            end % if_1
141
142 end % while
143 % result printing (suppressed if verbositiy set to 0)
144
        vprintf('\n%d man has been dumped for another\n\n', dumped);
145
146 else
       vprintf('\n%d men have been dumped for others\n\n', dumped);
147
148 end % if
   single = size(find(engaged(:,2)==0),1); % number of single nodes
   if single==1
150
       vprintf('There is %d single man/woman\n\n', single);
151
152 else
       vprintf('There are %d single men/women\n', single);
153
154 end % if
```

```
155 [stable, counter] = checkEngagements(engaged,initialm,initialf); % check ...
       the engagements
156 if (stable)
        vprintf('marriages are stable\n');
157
158 else
        vprintf('marriages are unstable\n');
159
        if counter==1
160
            vprintf('there is %d unstable mariage\n', counter);
161
        else
162
           vprintf('there are %d unstable mariages\n', counter);
163
        end % if_2
165 end % if
166 % calculate optimality index
167 \text{ opt} = 0;
168 for i = 1:n
        he = i;
169
        she = engaged(he, 2);
170
171
        if she^=0
172
            hisindex = find(initialf(she,:) == he, 1);
173
            herindex = find(initialm(he,:) == she, 1);
174
        else
            hisindex = n;
175
            herindex = n;
176
177
        end
        opt = opt + hisindex + herindex;
179 end
180 opt = opt/(2*n*n);
181 vprintf('optimality index is %1.2f\n',opt);
182 % set output
183 output = zeros(1,4);
184 output (1,1) = counter;
185 output (1,2) = single;
186 output (1,3) = dumped;
187 output (1,4) = \text{opt};
188 end
```

checkEngagements.m

```
1 function [ stable,counter ] = checkEngagements( engaged, m, f )
2 %checkEngagements checks whether a set of engagements is stable
3 %
4 % men an women encoded as integers from 1 to n
5 %
6 % input:
7 % engaged: engagement matrix
8 % m,f: preference matrices
9 %
10 % dimensions must be correct, m=nxn, f=nxn, engaged=nx2
11 %
```

```
returns:
       stable: true for stable engagements, false otherwise
14 %
       counter: the number of unstable mariages
15
16 n = size(m, 1); % input size
17 % reverse the engaged matrix such that the new matrix has the index of the
18 % women on the column one and those of their respective husbands in row two
invengaged=zeros(n,2);
20 copy = engaged(:,[2,1]);
21 i=1;
22 while i~=n+1
23
       index=copy(i,1);
       while index == 0 && i~=n % find first index that is nonzero
           i=i+1;
25
           index=copy(i,1);
26
       end % while
27
       if index==0 && i==n
28
           break;
29
       end % if
31
       invengaged(index,:) = copy(i,:);
       i=i+1;
32
33 end % while
34 % assign local variables
35 stable=true;
36 he=1;
37 counter=0;
38 \text{ inst} = [0,0];
39 % main loop
40 while he<=n
       she = engaged(he,2); % she is engaged to he
41
       while (she==0 && he~=n) % he is not engaged, so there is no instability ...
           -> check the next man
           he = he+1;
43
           she = engaged (he, 2);
44
       end % while
45
       if she==0 % -> he=n is not engaged, nothing to check.
46
47
           break;
48
       end %if
       % get indexes in pref lists
       hisindex = find(f(she,:)==he,1);
50
       herindex = find(m(he,:) == she, 1);
51
       helikesbetter = m(he,1:herindex);
52
       shelikesbetter = f(she,1:hisindex);
53
       % check for her
54
       if "isempty(shelikesbetter) % there is no one on earth she likes better
           for i=1:size(shelikesbetter) % loop to check if there is ...
               unstability for the girl
               guy = shelikesbetter(i); % all the guys she likes better
57
               guysgirl = engaged(guy,2); % the guy she is engaged to
58
```

```
if quysqirl == 0 \&\& ~isempty(find(m(quy,:) == she,1)) % if this ...
59
                   guy isn't engaged, then she could be with him -> unstable, ...
                   unless he doesn't know her.
                   stable = false;
60
                   vprintf('man %d and woman %d like each other better\n', guy, ...
                      she);
                  inst = [guy, she; inst];
62
63
               else
64
                    guylikes = m(guy,:); % the ordered preferences of guy
65
                    if (find(guylikes==she,1)<find(guylikes==guysgirl,1)) % if ...</pre>
                        guy also likes she better than his wife -> unstable
                        stable = false;
67
                        vprintf('man %d and woman %d like each other better\n', ...
68
                            guy, she);
                        inst = [guy, she; inst];
69
                    end % if 3
70
               end % if_2
71
           end % for
73
       end % if_1
       % now the other way round, check for him
74
       if ~isempty(helikesbetter) % there is no one on earth he likes better
75
           for i=1:size(helikesbetter) % loop to check if there is unstability ...
76
               for the man
77
               girl = helikesbetter(i); % all the girls he likes better
78
               girlsquy = invengaged(girl,2); % the girl he is engaged to
               if girlsguy == 0 && ~isempty(find(f(girl,:) == he,1))% if this ...
79
                   girl isn't engaged, then she could be with her -> unstable
                    stable = false;
80
                    vprintf('man %d and woman %d like each other better\n', he, ...
81
                       girl);
82
                    inst = [he,girl;inst];
               else
83
                    girllikes = f(girl,:);% the ordered preferences of girl
84
                    if (find(girllikes==he,1)<find(girllikes==girlsquy,1)) % if ...</pre>
85
                       guy also likes she better than his wife \rightarrow unstable
                        stable = false;
86
                        vprintf('man %d and woman %d like each other better\n', ...
                            he, girl);
                        inst = [he,girl;inst];
88
                    end % if_3
89
               end % if_2
90
           end % for
91
       end % if_1
92
       he=he+1; % go to the next man
95 end % while
96 % delete duplicate instabilities
97 inst = unique(inst, 'rows');
98 counter = size(inst, 1) -1;
```

simulation.m

```
1 function [ data ] = simulation( saveit )
2 %simulation perform simulation
3 %
       input:
4 ^{\circ}_{\circ}
       saveit: if 1 data is saved, if 0 not
 5 %
 6 %
7 %
     returns:
 8 % data: simulation data
10
11 %simulation
13 % simulate match making
14 % n is 2et, t from 1 to 6
15 % radius is either constant or random
16 % when constant, in 0.1:0.05:0.5
17 % frequency
19 global verbosity
20 verbosity = 0;
21
22 assert(~isempty(find(saveit==[0,1],1)));
23 \text{ tmax} = 6;
24 t = 2.^(1:tmax);
25 r = 0.1:0.05:0.5;
26 data = zeros(tmax, 10, 4);
27 seed = rng;
28 if saveit==1
29
       dirname = sprintf('data/%s',datestr(now,'yyyy_mm_dd_HH_MM_SS'));
30
       mkdir(dirname);
31 end
32
33 % radius random
34 for i=1:tmax
     n = t(i);
      [a,b] = generatePlane(n,2);
      [x,y] = makeMatch(a,b);
38
      data(i, 10, :) = y;
39 end
40
41 % radius const
42 for i=1:tmax
43
      for j=1:9
          n = t(i);
           radius = r(j);
```

```
[a,b] = generatePlane(n,1,radius);
46
           [x,y] = makeMatch(a,b);
47
           data(i,j,:) = y;
48
49
      end
50 end
51 % plot optimality index for each radius
52 hold on
53 handle = figure(1);
54 \text{ col} = \text{hsv}(10);
55 %set(groot, 'defaultAxesLineStyleOrder', {'-*', ':', 'o'});
56 for i=1:10
      plot(1:tmax,data(:,i,4),'color', col(i,:), 'marker', '*','linestyle','--');
      title('optimality index for for different radiuses');
58
59
60 end
62 xlabel('input size 2^x');
63 ylabel('optimality index');
64 legend([num2str(r', 'radius %1.3f');arr]);
65 if saveit==1
       saveas(handle,sprintf('%s/figure_1.pdf', dirname));
66
67 end
68 hold off
70 % plot no of dumps for each radius
71 handle = figure(2);
72 for i=1:10
      subplot(3,4,i);
73
      bar(1:tmax, data(:,i,3));
74
      xlabel('input size 2^x');
75
      ylabel('number of dumps');
76
77
      ylim([0,100]);
      if i~=10
78
           title(sprintf('plotting #dumps for radius %1.3f',r(i)));
79
80
      else
          title('plotting #dumps for radius random');
81
82
      end
83
85 if saveit==1
      saveas(handle,sprintf('%s/figure_2.pdf', dirname));
86
87 end
88 % saving
89 if (saveit==1)
       save(sprintf('%s/data.mat',dirname),'data','seed');
91 end
92
93 end
```

simulation2.m

```
1 function [ data ] = simulation2( saveit )
2 %simulation perform simulation
3 %
  90
       input:
4
5
   응
       saveit: if 1 data is saved, if 0 not
      returns:
      data: simulation data
8
10
11 %simulation
12
13 % simulate match making
14 % n is 2et, t from 1 to 6
15 % radius is either constant or random
16 % when constant, r=0.4
17 %makeMatch with preference changes for men, women, men/women and
18 %probability of changing in p=0.1:0.1:0.9
19
20 global verbosity
21 verbosity = 0;
22
23 assert(~isempty(find(saveit==[0,1],1)));
24 \text{ tmax} = 11;
25 t = 2.^(1:tmax);
p = 0.0:0.25:1.0;
pmax=size(p,2);
28 m=12; %number of different makeMatch called in simulation2
29 data = zeros(tmax, m*pmax, 4);
30 r = 0.1:0.15:0.4;
31 seed = rng;
32 if saveit==1
       dirname = sprintf('data/%s',datestr(now,'yyyy_mm_dd_HH_MM_SS'));
       mkdir(dirname);
34
35 end
36
37 % radius random
38 for i=1:tmax
       n = t(i);
       [a,b] = generatePlane(n,2);
40
       fprintf('Plane with rand radius, size %d is generated\n', n);
41
       for j=1:pmax
42
          [x1,y1] = makeMatch(a,b, p(j),1);% x: engagement matrix; y(1): ...
43
              #unstable mariage, y(2): #single men/women, y(3): #dumps, y(4): ...
              optimality index
          [x2,y2] = makeMatch(a,b, p(j),0);
44
          [x3,y3] = makeMatch(a,b, p(j),0.5);
45
          data(i,j,:) = y1;
46
          data(i,pmax+j,:)=y2;
47
          data(i,2*pmax+j,:)=y3;
```

```
end
49
  end
50
51
  %constant radius in r=[0.1, 0.25, 0.4]
   for i=1:tmax
       n = t(i);
       [a,b] = generatePlane(n,1, r(1));
55
       fprintf('Plane with radius r=0.1, size %d is generated\n', n);
56
       for j=1:pmax
57
          [v1,w1] = makeMatch(a,b, p(j),1);% x: engagement matrix; y(1): ...
              #unstable mariage, y(2): #single men/women, y(3): #dumps, y(4): ...
              optimality index
59
          [v2, w2] = makeMatch(a, b, p(j), 0);
          [v3, w3] = makeMatch(a, b, p(j), 0.5);
60
          data(i,3*pmax+j,:) = w1;
61
          data(i, 4*pmax+j, :) = w2;
62
          data(i,5*pmax+j,:) = w3;
       end
65
  end
66
  for i=1:tmax
67
       n = t(i);
68
       [a,b] = generatePlane(n,1, r(2));
69
70
       fprintf('Plane with radius r=0.25, size %d is generated\n', n);
71
       for j=1:pmax
          [v1, w1] = makeMatch(a, b, p(j), 1); % x: engagement matrix; y(1): ...
              #unstable mariage, y(2): #single men/women, y(3): #dumps, y(4): ...
              optimality index
          [v2, w2] = makeMatch(a,b, p(j),0);
73
          [v3, w3] = makeMatch(a, b, p(j), 0.5);
74
          data(i,6*pmax+j,:) = w1;
75
76
          data(i,7*pmax+j,:) = w2;
77
          data(i, 8*pmax+j, :) = w3;
       end
78
   end
79
80
  for i=1:tmax
       n = t(i);
83
       [a,b] = generatePlane(n,1, r(3));
       fprintf('Plane with radius r=0.4, size %d is generated\n', n);
84
85
       for j=1:pmax
          [v1,w1] = makeMatch(a,b, p(j),1);% x: engagement matrix; y(1): ...
86
              #unstable mariage, y(2): #single men/women, y(3): #dumps, y(4): ...
              optimality index
           [v2, w2] = makeMatch(a, b, p(j), 0);
87
          [v3, w3] = makeMatch(a,b, p(j), 0.5);
88
          data(i, 9*pmax+j, :) = w1;
89
          data(i,10*pmax+j,:) = w2;
90
          data(i,11*pmax+j,:) = w3;
91
       end
```

simulation3.m

```
1 function [ data ] = simulation3( saveit , tmax )
 2 %SIMULATION3 perform simulation
4
   9
       simfeld
 5 %
 6 %
       input:
       saveit: if 1 data is saved, if 0 not
 7 %
 8 %
       tmax: maximal exponent form input size (base 2)
10 %
     returns:
11 % data: simulation data
12
13 % simulate match making
14 % n is 2^t, t from 1 to tmax
15 % radius is either constant or random
16 % when constant, in 0.1:rstep:0.5
17 % frequency changerate 0:fstep:1
18 % in makeMatch argument list p=0.5 \rightarrow change pref for both men and women
19
20 global verbosity
21 verbosity = 0;
23 assert(~isempty(find(saveit==[0,1],1)));
24 if nargin < 2
25
       tmax = 6;
26 end
27 t = 2.^(1:tmax);
28 \text{ rstep} = 0.1;
r = 0.1:rstep:0.5;
30 fstep = 0.5;
31 f = 0:fstep:1;
32 \text{ sizef} = \text{size}(f,2);
33 sizer = size(r,2)+1;
34 m = 10; % number of iterations
35 data = zeros(tmax, sizer, sizef, m, 4);
```

```
36 % data dimensions:
37 % 1 input size n, tmax
38 % 2 radius, sizer
39 % 3 frequency, sizef
40 % 4 iterations, m
41 % 5 output values, 4
42 seed = rng;
43 disp(seed);
44 if saveit==1
       dirname = sprintf('data/%s',datestr(now,'yyyy_mm_dd_HH_MM_SS'));
       mkdir(dirname);
47 end
48 tic
49 fprintf('simulating for radius random\n');
50 % radius random
51 for i=1:tmax
       n = t(i);
52
53
       for k=1:sizef
54
           freq = f(k);
           for l=1:m
55
                %[a,b] = generatePlane(n,2);
56
                [a,b] = generateRandom(n);
57
                fprintf('.');
58
59
                [x,y] = makeMatch(a,b,freq,0.5);
                data(i, sizer, k, l, :) = y;
61
           end
           fprintf('\n');
62
63
       fprintf('n = %4d complete after %5.1f(n', n, toc);
64
65 end
66 if 1==0
67 fprintf('simuation for radius const\n')
68 % radius const
  for i=1:tmax
69
       n = t(i);
70
       for j=1:sizer-1
71
           radius = r(j);
72
73
           for k=1:sizef
74
                freq = f(k);
                for l=1:m
75
                    %[a,b] = generatePlane(n,1,radius);
76
                    [a,b] = generateRandom(n);
77
                    fprintf('.');
78
                    [x,y] = makeMatch(a,b,freq,0.5);
79
                    data(i,j,k,l,:) = y;
80
                end
81
                fprintf('\n');
82
           end
83
           fprintf('radius %1.1f complete\n', radius);
84
85
       end
```

```
fprintf('n = %4d complete after %5.1f\n', n, toc);
end
se end
se end
plot3(data);
ff (saveit==1)
save(sprintf('%s/data.mat',dirname),'data','seed');
end
end
end
```

iteratesimulation.m

```
1 function [data, dataplot] = iteratesimulation(saveit, n)
       global verbosity
2
3
       verbosity=0;
4
       %!!!tmax, m, p, t, pmax must correspond to simulation2!!!
5
       assert(~isempty(find(saveit==[0,1],1)));
6
7
       seed = rng;
       tmax=11;
       m=12; %number of makeMatch called in simulation2
9
       p = 0.0:0.25:1.0;
10
       pmax=size(p, 2);
11
       t = 2.^(1:tmax);
12
       data = zeros(tmax, m*pmax, 4, n);
13
       dataplot=zeros(tmax, m*pmax, 4);
       if saveit==1
15
16
           dirname = sprintf('data/%s',datestr(now,'yyyy_mm_dd_HH_MM_SS'));
17
           mkdir(dirname);
       end
18
19
       for i=1:n
20
           data(:,:,:,i) = simulation2(0);
21
           dataplot=dataplot+data(:,:,:,i);
22
           fprintf('iteration %d is done\n', i);
23
       end
       dataplot=dataplot/(n*1.);
24
25
26
27
28 % saving
29 if (saveit==1)
30
       save(sprintf('%s/data.mat',dirname),'data','dataplot','seed');
31 end
32
33 end
```

plot3.m

```
1 function [ ] = plot3( data )
2 %PLOT3 plotting for optimality index analysis
4 %
      simfeld
5 if 1==0
6 dirname = sprintf('data/%s',datestr(now,'yyyy_mm_dd_HH_MM_SS'));
7 mkdir(dirname);
8 % define arrays
9 sizer = size(data,2);
10 r=0.1:0.1:0.5;
11 f=0:0.5:1;
12 sizef=size(f,2);
13 tmax = size(data, 1);
14 t=1:tmax;
16 % x-axis n, diff radius, plots freq
17 col = hsv(sizer);
18 for j=1:sizef
       freq = f(j);
19
20
       handle = figure(j);
21
       set(gca, 'FontSize', 16);
       hold on
       for i=1:sizer
23
24
           mm = squeeze(mean(data(:,i,j,:,4),4));
           st = squeeze(std(data(:,i,j,:,4),0,4));
25
           %plot(1:tmax,data(:,i,1,1,4),'color', col(i,:), 'marker', ...
               '*','linestyle','--');
27
           errorbar(1:tmax,mm,st,'color', col(i,:), 'marker', ...
               '.', 'linestyle', '-.', 'LineWidth', 2);
       end
28
       hold off
29
       box on
30
31
       grid on
       title(sprintf('optimality index for different radii, changerate is ...
32
           %1.1f', freq));
       arr = ['r', 'a', 'n', 'd', 'o', 'm', ' ', ' ', ' ', ' '];
33
       xlabel('input size 2^x');
34
       ylabel('optimality index');
35
       ylim([0,1.1]);
36
37
       xlim([0,11]);
       legend([num2str(r', 'radius %1.1f');arr]);
39
       saveas(handle, sprintf('%s/figure_%d.pdf', dirname, j));
40 end
41
42 % x-axis radius, diff n, plots freq
43 col = hsv(tmax);
44 for j=1:sizef
```

```
freq = f(j);
45
       handle = figure(j+3);
46
       set(gca, 'FontSize', 16);
47
       hold on
48
       for i=1:tmax
           mm = squeeze(mean(data(i,:,j,:,4),4));
50
           st = squeeze(std(data(i,:,j,:,4),0,4));
51
           %plot(1:tmax,data(:,i,1,1,4),'color', col(i,:), 'marker', ...
52
                '*','linestyle','--');
           errorbar(0.1:0.1:0.6,mm,st,'color', col(i,:), 'marker', ...
53
               '.', 'linestyle', '-.', 'LineWidth', 2);
       end
       hold off
55
       box on
56
       grid on
57
       title(sprintf('optimality index for for different n, changerate is ...
58
           %1.1f', freq));
       %arr = ['r','a','n','d','o','m',' ',' ',' ',' ',' ',' ',' '];
59
60
       xlabel('radius, 0.6 is random');
       ylabel('optimality index');
61
       ylim([0,1.1]);
62
63
       xlim([0,0.9])
       legend(num2str(t','n = 2^{\frac{3}{2}}));
64
65
       saveas(handle,sprintf('%s/figure_%d.pdf', dirname, j+3));
66
  end
67
  % x-axis freq, diff n, plots radius
68
  col = hsv(tmax);
69
  for j=1:sizer
70
       if j~= sizer radius = r(j); end
71
       handle = figure(j+6);
72
73
       set(gca, 'FontSize', 16);
       hold on
74
       for i=1:tmax
75
           mm = squeeze(mean(data(i,j,:,:,4),4));
76
           st = squeeze(std(data(i,j,:,:,4),0,4));
77
           %plot(1:tmax,data(:,i,1,1,4),'color', col(i,:), 'marker', ...
78
               '*','linestyle','--');
           errorbar(f,mm,st,'color', col(i,:), 'marker', '.','linestyle','-.', ...
               'LineWidth', 2);
       end
80
       hold off
81
       box on
82
       grid on
83
       if j~= sizer
           title(sprintf('optimality index for for different n, radius ...
85
               %1.1f', radius));
       else
86
           title(sprintf('optimality index for for different n, radius random'));
87
88
       end
```

```
89
        xlabel('changerate');
90
        ylabel('optimality index');
91
92
        ylim([0,1.1]);
        xlim([-0.1, 1.4]);
        legend(num2str(t','n = 2^{*2d}'));
94
        saveas(handle, sprintf('%s/figure_%d.pdf', dirname, j+6));
95
96 end
97
98 % surfs
99 handle = figure(13);
100 [a,b]=meshgrid(1:10,0.1:0.1:0.5);
102 mm = squeeze(mean(data(:,1:5,3,:,4),4));
103 view(0,90);
104 box on
105 grid on
106 surf(a,b,mm','EdgeColor','none');
107 colorbar;
108 set(gca, 'FontSize', 16);
109 xlim([1,10]);
110 xlabel('input size2^x');
111 ylabel('radius');
saveas(handle,sprintf('%s/figure_13.pdf', dirname));
113 end
114 %generaterandom
115 handle = figure(14);
set(gca, 'FontSize', 16);
117 box on
118 dd=zeros(10,20,4);
119 for i=1:10
120
        n = 2^i;
        for j=1:20
121
            [a,b] = generateRandom(n);
122
            [x,y] = makeMatch(a,b);
123
124
            dd(i,j,:) = y;
            fprintf('.');
125
126
        end
127
        fprintf('\n');
128 end
129 mm = squeeze(mean(dd(:,:,4),2));
130 st = squeeze(std(dd(:,:,4),0,2));
131 errorbar(1:10,mm,st, 'marker', '.','linestyle','-.', 'LineWidth', 2);
132 %plot(1:10,log2(mm), 'marker', '.','linestyle','-.', 'LineWidth', 2);
133 p = polyfit(1:10, log2(mm'), 1);
134 % line = polyval(p,1:10);
135 hold on;
136 x=1:10;
137 plot(x, (2^p(2)) * (x.^(p(1))));
```

```
139 xlabel('input size 2^x');
140 ylabel('optimality index');
141 title('optimality index for randomly generated preferences');
142 %ylim([0,1.1]);
143 xlim([0,11]);
144 box on
145 grid on
146 %saveas(handle,sprintf('%s/figure_14.pdf', dirname));
147 end
```

10 Appendix B: Plots

Additional Plots

10.1 Optimality Index

