

Reduction of the Symmetric Eigenproblem $Ax = \lambda Bx$ and Related Problems to Standard Form^{*}

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1. Theoretical Background

In many fields of work the solution of the eigenproblems $Ax = \lambda Bx$ and $ABx = \lambda x$ (or related problems) is required, where A and B are symmetric and B is positive definite. Each of these problems can be reduced to the standard symmetric eigenproblem by making use of the Cholesky factorization [4] of B .

For if L is defined by

$$LL^T = B \quad (1)$$

then^{**}

$$Ax = \lambda Bx \quad \text{implies} \quad (L^{-1}AL^{-T})(L^Tx) = \lambda(L^Tx) \quad (2)$$

and

$$ABx = \lambda x \quad \text{implies} \quad (L^TAL)(L^Tx) = \lambda(L^Tx). \quad (3)$$

Hence the eigenvalues of $Ax = \lambda Bx$ are those of $Py = \lambda y$, where P is the symmetric matrix $L^{-1}AL^{-T}$, while the eigenvalues of $ABx = \lambda x$ are those of $Qy = \lambda y$, where Q is the symmetric matrix L^TAL . Note that in each case the eigenvector x of the original problem is transformed into L^Tx .

There are a number of closely related problems which are defined by the equations

$$\begin{aligned} y^T AB &= \lambda y^T & \text{or} & & (L^{-1}y)^T(L^TAL) &= \lambda(L^{-1}y)^T \\ BAy &= \lambda y & \text{or} & & (L^TAL)(L^{-1}y) &= \lambda(L^{-1}y) \\ x^T BA &= \lambda x^T & \text{or} & & (L^Tx)^T(L^TAL) &= \lambda(L^Tx)^T. \end{aligned} \quad (4)$$

Here it is assumed throughout that it is B which is required to be positive definite (A may or may not be positive definite). It will be observed that all the problems can be reduced if we have procedures for computing L^TAL and $L^{-1}AL^{-T}$. Further if we denote an eigenvector of the derived problem by z , we need only procedures for solving

$$L^Tx = z \quad \text{and} \quad L^{-1}y = z \quad \text{i.e.} \quad y = Lz. \quad (5)$$

If z is normalized so that $z^Tz = 1$ then x and y determined from Eqs. (5) satisfy

$$x^TBx = 1 \quad \text{and} \quad y^TB^{-1}y = 1, \quad y^Tx = 1 \quad (6)$$

and these are the normalizations that are usually required in practice.

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^{**} Throughout this paper we have used L^{-T} in place of the cumbersome $(L^{-1})^T$ and $(L^T)^{-1}$.

Any decomposition $B = X X^T$ may be used in a similar way, but the Cholesky decomposition is the most economical, particularly when carried out as described in Section 5. If the standard eigenproblem for the matrix B is solved by a method which gives an accurately orthogonal system Y of eigenvectors then we have

$$B = Y \operatorname{diag}(d_i^2) Y^T = Y D^2 Y^T \quad (7)$$

where d_i^2 are the eigenvalues (positive) of B . Hence corresponding to $A x = \lambda B x$ we may solve $(D^{-1} Y^T A Y D^{-1})(D Y^T x) = \lambda (D Y^T x)$ and similarly for the other problems. Generally this is much less efficient than the use of the Cholesky decomposition, but it may have slight advantages in very special cases (see the discussion on pp. 337–338, 344 of ref. [7]).

2. Applicability

reduc 1 may be used to reduce the eigenproblem $A x = \lambda B x$ to the standard symmetric eigenproblem $P z = \lambda z$, where $P = L^{-1} A L^{-T}$ and $B = L L^T$.

reduc 2 may be used to reduce the eigenproblems $y^T A B = \lambda y^T$, $B A y = \lambda y$, $A B x = \lambda x$, $x^T B A = \lambda x^T$ to the standard symmetric eigenproblem $Q z = \lambda z$, where $Q = L^T A L$ and $B = L L^T$.

When A and B are narrow symmetric band matrices of high order these procedures should not be used as they destroy the band form.

The derived standard symmetric eigenproblem may be solved by procedures which are included in this series [1, 3, 5, 6]. The eigenvalues of the derived standard problem are those of the original problem but the vectors are related as indicated by Eq. (2), (3), (4). The eigenvectors of the original problem may then be obtained using the procedures *rebak a* and *rebak b* described below.

rebak a provides eigenvectors x of the problems $A x = \lambda B x$, $A B x = \lambda x$, $x^T B A = \lambda x^T$ from the corresponding eigenvectors $z = L^T x$ of the derived standard symmetric problem.

rebak b provides eigenvectors y of the problems $y^T A B = \lambda y^T$, $B A y = \lambda y$ from the corresponding eigenvectors $z = L^{-1} y$ of the derived standard symmetric problem.

3. Formal Parameter List

3.1. Input to procedure *reduc 1*

- n order of matrices A and B , or negative order if L already exists.
- a elements of the symmetric matrix A given as upper-triangle of an $n \times n$ array (strict lower-triangle can be arbitrary).
- b elements of the symmetric positive definite matrix B given as upper-triangle of an $n \times n$ array (strict lower-triangle can be arbitrary).

Output of procedure *reduc 1*

- a elements of symmetric matrix $P = L^{-1} A L^{-T}$ given as lower-triangle of $n \times n$ array. The strict upper-triangle of A is still preserved but its diagonal is lost.

- b* sub-diagonal elements of matrix L such that $LL^T = B$ stored as strict lower-triangle of an $n \times n$ array.
- dl* the diagonal elements of L stored as an $n \times 1$ array.
- fail* the exit used if B , possibly on account of rounding errors, is not positive definite.

3.2. Input to procedure *reduc 2*

- n* order of matrices A and B or negative order if L already exists.
- a* elements of the symmetric matrix A given as upper-triangle of an $n \times n$ array (strict lower-triangle can be arbitrary).
- b* elements of the symmetric positive definite matrix B given as upper-triangle of an $n \times n$ array (strict lower-triangle can be arbitrary).

Output from procedure *reduc 2*

- a* elements of symmetric matrix $Q = L^T A L$ given as lower-triangle of $n \times n$ array. The strict upper-triangle of A is still preserved but its diagonal is lost.
- b* sub-diagonal elements of matrix L such that $LL^T = B$ stored as strict lower-triangle of an $n \times n$ array.
- dl* the diagonal elements of L stored as an $n \times 1$ array.
- fail* the exist used if B , possibly on account of rounding errors, is not positive definite.

3.3. Input to procedure *rebak a*

- n* order of matrices A and B .
- m1, m2* eigenvectors $m1, \dots, m2$ of the derived standard symmetric eigenproblem have been found.
- b* the sub-diagonal elements of the matrix L such that $LL^T = B$ stored as the strict lower-triangle of an $n \times n$ array (as given by *reduc 1* or *reduc 2*).
- dl* the diagonal elements of L stored as an $n \times 1$ array.
- z* an $n \times (m2 - m1 + 1)$ array containing eigenvectors $m1, \dots, m2$ of the derived symmetric eigenproblem.

Output from procedure *rebak a*

- z* an $n \times (m2 - m1 + 1)$ array containing eigenvectors of the original problem; output $z = L^{-T}(\text{input } z)$.

3.4. Input to procedure *rebak b*

- n* order of matrices A and B .
- m1, m2* eigenvectors $m1, \dots, m2$ of the derived standard symmetric eigenproblem have been found.

- b* the sub-diagonal elements of the matrix L such that $LL^T = B$ stored as the strict lower-triangle of an $n \times n$ array (as given by *reduc 1* or *reduc 2*).
- dl* the diagonal elements of L stored as an $n \times 1$ array.
- z* an $n \times (m2 - m1 + 1)$ array containing eigenvectors $m1, \dots, m2$ of the derived symmetric eigenproblem.

Output from procedure *rebak b*

- z* an $n \times (m2 - m1 + 1)$ array containing eigenvectors of the original problem. output $z = L(\text{input } z)$.

4. ALGOL Programs

procedure *reduc1* (*n*) *trans*: (*a*, *b*) *result*: (*dl*) *exit*: (*fail*);

value *n*; **integer** *n*; **array** *a*, *b*, *dl*; **label** *fail*;

comment Reduction of the general symmetric eigenvalue problem

$$A \times x = \text{lambda} \times B \times x,$$

with symmetric matrix A and symmetric positive definite matrix B , to the equivalent standard problem $P \times z = \text{lambda} \times z$.

The upper triangle, including diagonal elements, of A and B are given in the arrays $a[1:n, 1:n]$ and $b[1:n, 1:n]$.

L ($B = L \times L^T$) is formed in the remaining strictly lower triangle of the array b with its diagonal elements in the array $dl[1:n]$, and the lower triangle of the symmetric matrix P ($P = \text{inv}(L) \times A \times \text{inv}(L^T)$) is formed in the lower triangle of the array a , including the diagonal elements. Hence the diagonal elements of A are lost.

If $n < 0$, it is assumed that L has already been formed.

The procedure will fail if B , perhaps on account of rounding errors, is not positive definite;

```
begin
  integer i, j, k;
  real x, y;
  if  $n < 0$  then  $n := -n$ 
  else
    for  $i := 1$  step 1 until  $n$  do
      for  $j := i$  step 1 until  $n$  do
        begin  $x := b[i, j]$ ;
          for  $k := i - 1$  step  $-1$  until 1 do
             $x := x - b[i, k] \times b[j, k]$ ;
          if  $i = j$  then
            begin if  $x \leq 0$  then go to fail;
               $y := dl[i] := \text{sqrt}(x)$ 
            end
          else  $b[j, i] := x/y$ 
        end  $j$ ;
    comment  $L$  has been formed in the array b;
    for  $i := 1$  step 1 until  $n$  do
      begin  $y := dl[i]$ ;
        for  $j := i$  step 1 until  $n$  do
```

```

    begin  $x := a[i, j]$ ;
      for  $k := i - 1$  step  $-1$  until  $1$  do
         $x := x - b[i, k] \times a[j, k]$ ;
         $a[j, i] := x/y$ 
      end  $j$ 
    end  $i$ ;
  comment The transpose of the upper triangle of  $inv(L) \times A$  has been
    formed in the lower triangle of the array  $a$ ;
  for  $j := 1$  step  $1$  until  $n$  do
    for  $i := j$  step  $1$  until  $n$  do
      begin  $x := a[i, j]$ ;
        for  $k := i - 1$  step  $-1$  until  $j$  do
           $x := x - a[k, j] \times b[i, k]$ ;
          for  $k := j - 1$  step  $-1$  until  $1$  do
             $x := x - a[j, k] \times b[i, k]$ ;
             $a[i, j] := x/dl[i]$ 
          end  $k$ 
        end  $k$ 
      end  $ij$ 
    end  $i$ 
  end  $reduc1$ ;

```

procedure *reduc2* (n) *trans*: (a, b) *result*: (dl) *exit*: (*fail*);

value n ; **integer** n ; **array** a, b, dl ; **label** *fail*;

comment Reduction of the general symmetric eigenvalue problems

$$A \times B \times x = \text{lambda} \times x, \quad y^T \times A \times B = y^T \times \text{lambda},$$

$$B \times A \times y = \text{lambda} \times y, \quad x^T \times B \times A = x^T \times \text{lambda},$$

with symmetric matrix A and symmetric positive definite matrix B , to the equivalent standard problem $Q \times z = \text{lambda} \times z$.

The upper triangle, including diagonal elements, of A and B are given in the arrays $a[1:n, 1:n]$ and $b[1:n, 1:n]$.

L ($B = L \times L^T$) is formed in the remaining strictly lower triangle of the array b with its diagonal elements in the array $dl[1:n]$, and the lower triangle of the symmetric matrix Q ($Q = L^T \times A \times L$) is formed in the lower triangle of the array a , including the diagonal elements. Hence the diagonal elements of A are lost.

If $n < 0$, it is assumed that L has already been formed.

The procedure will fail if B , perhaps on account of rounding errors, is not positive definite;

```

begin   integer  $i, j, k$ ;
        real  $x, y$ ;
        if  $n < 0$  then  $n := -n$ 
        else
          for  $i := 1$  step  $1$  until  $n$  do
            for  $j := i$  step  $1$  until  $n$  do
              begin  $x := b[i, j]$ ;
                for  $k := i - 1$  step  $-1$  until  $1$  do
                   $x := x - b[i, k] \times b[j, k]$ ;
                  if  $i = j$  then

```

```

    begin if  $x < 0$  then go to fail;
         $y := dl[i] := \text{sqrt}(x)$ 
    end
    else  $b[j, i] := x/y$ 
end  $ji$ ;
comment  $L$  has been formed in the array  $b$ ;
for  $i := 1$  step 1 until  $n$  do
for  $j := 1$  step 1 until  $i$  do
begin  $x := a[j, i] \times dl[j]$ ;
    for  $k := j + 1$  step 1 until  $i$  do
         $x := x + a[k, i] \times b[k, j]$ ;
    for  $k := i + 1$  step 1 until  $n$  do
         $x := x + a[i, k] \times b[k, j]$ ;
     $a[i, j] := x$ 
end  $ji$ ;
comment The lower triangle of  $A \times L$  has been formed in the lower
        triangle of the array  $a$ ;
for  $i := 1$  step 1 until  $n$  do
begin  $y := dl[i]$ ;
    for  $j := 1$  step 1 until  $i$  do
        begin  $x := y \times a[i, j]$ ;
            for  $k := i + 1$  step 1 until  $n$  do
                 $x := x + a[k, j] \times b[k, i]$ ;
             $a[i, j] := x$ 
        end  $j$ 
    end  $i$ 
end reduc2;
```

procedure rebaka (n) data: ($m1, m2, b, dl$) trans: (z);

value $n, m1, m2$; **integer** $n, m1, m2$; **array** b, dl, z ;

comment This procedure performs, on the matrix of eigenvectors, Z , stored in the array $z[1:n, m1:m2]$, a backward substitution $LT \times X = Z$, overwriting X on Z .

The diagonal elements of L must be stored in the array $dl[1:n]$, and the remaining triangle in the strictly lower triangle of the array $b[1:n, 1:n]$. The procedures *reduc1* and *reduc2* leave L in this desired form.

If x denotes any column of the resultant matrix X , then x satisfies $xT \times B \times x = zT \times z$, where $B = L \times LT$;

```

begin    integer  $i, j, k$ ;
        real  $x$ ;
        for  $j := m1$  step 1 until  $m2$  do
        for  $i := n$  step  $-1$  until 1 do
```

```

begin  $x := z[i, j]$ ;
      for  $k := i + 1$  step 1 until  $n$  do
         $x := x - b[k, i] \times z[k, j]$ ;
         $z[i, j] := x/dl[i]$ 
      end  $ij$ 
end rebaka;

```

```

procedure rebakb ( $n$ ) data: ( $m1, m2, b, dl$ ) trans: ( $z$ );
value  $n, m1, m2$ ; integer  $n, m1, m2$ ; array  $b, dl, z$ ;

```

comment This procedure performs, on the matrix of eigenvectors, Z , stored in the array $z[1:n, m1:m2]$, a forward substitution $Y = L \times Z$, overwriting Y and Z .

The diagonal elements of L must be stored in the array $dl[1:n]$, and the remaining triangle in the strictly lower triangle of the array $b[1:n, 1:n]$. The procedures *reduc1* and *reduc2* leave L in this desired form.

If y denotes any column of the resultant matrix Y , then y satisfies $y^T \times inv(B) \times y = z^T \times z$, where $B = L \times LT$;

```

begin   integer  $i, j, k$ ;
        real  $x$ ;
        for  $j := m1$  step 1 until  $m2$  do
          for  $i := n$  step  $-1$  until 1 do
            begin  $x := dl[i] \times z[i, j]$ ;
                  for  $k := i - 1$  step  $-1$  until 1 do
                     $x := x + b[i, k] \times z[k, j]$ ;
                     $z[i, j] := x$ 
                  end  $ij$ 
            end rebakb;

```

5. Organisational and Notational Details

In *reduc 1* and *reduc 2* the symmetric matrices A and B are stored in $n \times n$ arrays though only the upper-triangles need be stored. The procedures *choldet 1* and *cholsol 1* [4] could have been used but since we wish to take full advantage of symmetry in computing the derived symmetric matrices special purpose procedures are preferable.

The matrix L such that $LL^T = B$ is determined in either case. Its strict lower-triangle is stored in the strict lower-triangle of B and its diagonal elements are stored in the linear array dl . Hence full information on B is retained. Since the factorization of B may already be known (if, for example several matrices A are associated with a fixed matrix B) provision is made for omitting the Cholesky factorization by inputting $-n$ instead of n .

In *reduc 1* the matrix $P = L^{-1}AL^{-T}$ is formed in the two steps

$$LX = A, \quad PL^T = X. \quad (8)$$

The matrix X is not symmetric but only half of it is required for the computation of P . The transpose of the upper-triangle of X is computed and written on the lower-triangle of A . The diagonal of A is therefore destroyed and must be copied

before using *reduc 1* if the residuals or Rayleigh quotients are required. The lower triangle of P is then determined and overwritten on X (see pp. 337–340 of ref. [7]).

Similarly in *reduc 2* the matrix $Q = L^T A L$ is computed in the steps

$$Y = A L, \quad Q = L^T Y. \quad (9)$$

Here only the lower-triangle of Y is required (although it is not symmetric). It is overwritten on the lower-triangle of A . The lower-triangle of Q is then formed and overwritten on Y . Again the diagonal of A is lost and a copy must be made independently if the residuals or Rayleigh quotients are required.

The derived standard symmetric eigenproblem may be solved in a variety of ways and the algorithms are given elsewhere in this series [1, 3, 5, 6]. Assuming eigenvectors of the derived problem have been found by some method and are normalized according to the Euclidean norm, then the corresponding eigenvectors of the original problem may be obtained using *rebak a* or *rebak b*, as appropriate (see 2).

6. Discussion of Numerical Properties

In general, the accuracy of the procedures described here is very high. The computed \bar{L} satisfies

$$\bar{L} \bar{L}^T = B + E \quad (10)$$

and we have discussed the bounds for E in [4]. The errors made in computing $\bar{L} A \bar{L}^T$ or $\bar{L}^{-1} A \bar{L}^{-T}$ are comparatively unimportant. In general the errors in the eigenvalues are such as could be produced by very small perturbations in A and B . However, with the eigenproblem $A x = \lambda B x$ small perturbations in A and B can correspond to quite large perturbations in the eigenvalues if B is ill-conditioned with respect to inversion and this can lead to the inaccurate determination of eigenvalues.

It might be felt that in this case such a loss of accuracy is inherent in the data but this is not always true. Consider, for example, the case

$$A = \begin{bmatrix} 2 & 2 \\ 2 & 1 \end{bmatrix}, \quad B = \begin{bmatrix} 1 & 2 \\ 2 & 4 + \varepsilon \end{bmatrix} \quad (11)$$

for which the eigenvalues are the roots of

$$-2 - \lambda(1 + 2\varepsilon) + \varepsilon \lambda^2 = 0$$

i.e.

$$\lambda = [+(1 + 2\varepsilon) \pm (1 + 12\varepsilon + 4\varepsilon^2)^{\frac{1}{2}}]/2\varepsilon. \quad (12)$$

As $\varepsilon \rightarrow 0$ one root tends to -2 and the other to $+\infty$. It may be readily verified that the root -2 is not at all sensitive to small random changes in the elements of A and B .

Rather surprisingly, well-conditioned eigenvalues such as the -2 of the previous example are often given quite accurately, in spite of the ill-condition of B . In the above example for instance the true matrix $P = L^{-1} A L^{-T}$ is

$$\begin{bmatrix} 2 & -2/\varepsilon^{\frac{1}{2}} \\ -2/\varepsilon^{\frac{1}{2}} & 1/\varepsilon \end{bmatrix}, \quad (13)$$

and for small ε it has some very large elements. When rounding errors are involved one might easily obtain a matrix \bar{P} of the form

$$\bar{P} = \begin{bmatrix} 2/(1 + \varepsilon_1) & -2(1 + \varepsilon_2)/(\varepsilon + \varepsilon_3)^{\frac{1}{2}} \\ -2(1 + \varepsilon_2)/(\varepsilon + \varepsilon_3)^{\frac{1}{2}} & (1 + \varepsilon_4)/(\varepsilon + \varepsilon_3) \end{bmatrix} \quad (14)$$

where the ε_i are of the order of machine precision. If ε itself is of this order the computed \bar{P} will differ substantially from the true matrix and the large elements may well have a high relative error. However \bar{P} has an eigenvalue close to -2 for all relevant values of the ε_i .

As regards the eigenvectors, there are procedures which give an eigensystem of the derived problem which is accurately orthogonal. However, when transforming back to the eigensystem of the original eigenproblem, the computed set of vectors will be ill-conditioned when B (and hence L) is ill-conditioned. This deterioration would appear to be inherent in the data.

In our experience the considerations have not been important in practice since the matrices B which arise are, in general, extremely well-conditioned.

7. Test Results

To give a formal test of the procedures described here the matrices

$$F = \begin{bmatrix} 10 & 2 & 3 & 1 & 1 \\ 2 & 12 & 1 & 2 & 1 \\ 3 & 1 & 11 & 1 & -1 \\ 1 & 2 & 1 & 9 & 1 \\ 1 & 1 & -1 & 1 & 15 \end{bmatrix}, \quad G = \begin{bmatrix} 12 & 1 & -1 & 2 & 1 \\ 1 & 14 & 1 & -1 & 1 \\ -1 & 1 & 16 & -1 & 1 \\ 2 & -1 & -1 & 12 & -1 \\ 1 & 1 & 1 & -1 & 11 \end{bmatrix}$$

were used. Since both F and G are positive definite either matrix can be used as A or B in either of the main procedures. When used in conjunction with procedures for solving the standard eigenproblems this gives quite extensive checks on the overall performance. Since both A and B are well-conditioned with respect to inversion good agreement should be achieved between computed matrices $L^{-1}AL^{-T}$, L^TAL , eigenvalues and eigenvectors obtained on different computers.

A selection from the results obtained on KDF 9, a computer using a 39 binary digit mantissa are given in Tables 1–4. The eigensystems of the derived standard eigenproblems were found using the procedures *tridi* [3], *bisect* [4], *tridi inverse iteration* [6] and *rebak* [3]. Any accurate procedures for dealing with the standard eigenproblem could be used including the procedure *jacobi* [5] already published.

Table 1 gives the output matrices a obtained from *reduc* 1 and *reduc* 2 taking $A=F$, $B=G$.

Table 2 gives the output matrices b and vectors dl obtained from *reduc* 1 and *reduc* 2 taking $A=F$, $B=G$. Since the computations involved in the two procedures are identical in every detail the output matrices are identical.

Table 3 gives the eigenvalues of $F - \lambda G$ and those of $G - \lambda F$ obtained using *reduc* 1; these involve the Cholesky factorizations of G and F respectively. The product of corresponding eigenvalues is equal to unity to about eleven decimal places. It also gives the eigenvalues of $FG - \lambda I$ and $GF - \lambda I$ calculated using

Table 1

Column 1	Column 2	Column 3	Column 4	Column 5
<i>Output matrix a from reduc 1 with A = F, B = G</i>				
+ 8.3333 3333 334 ₁₀ - 1;	+ 2.0000 0000 000 ₁₀ + 0;	+ 3.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
+ 9.0279 3771 290 ₁₀ - 2;	+ 8.4331 3373 256 ₁₀ - 1;	+ 1.0000 0000 000 ₁₀ + 0;	+ 2.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
+ 2.7151 9129 955 ₁₀ - 1;	- 4.4710 9421 464 ₁₀ - 3;	+ 7.2690 8512 779 ₁₀ - 1;	+ 1.0000 0000 000 ₁₀ + 0;	- 1.0000 0000 000 ₁₀ + 0;
- 3.3434 1703 880 ₁₀ - 2;	+ 2.1334 5628 682 ₁₀ - 1;	+ 6.2497 8723 923 ₁₀ - 2;	+ 8.1234 1947 476 ₁₀ - 1;	+ 1.0000 0000 000 ₁₀ + 0;
- 1.6323 2422 021 ₁₀ - 2;	+ 2.4171 4784 237 ₁₀ - 2;	- 1.5020 8557 974 ₁₀ - 1;	+ 1.3606 1578 747 ₁₀ - 1;	+ 1.4260 4956 979 ₁₀ + 0;
<i>Output matrix a from reduc 2 with A = F, B = G</i>				
+ 1.3100 0000 000 ₁₀ + 2;	+ 2.0000 0000 000 ₁₀ + 0;	+ 3.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
+ 4.4649 6005 171 ₁₀ + 1;	+ 1.6856 8862 275 ₁₀ + 2;	+ 1.0000 0000 000 ₁₀ + 0;	+ 2.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
+ 3.1565 2935 415 ₁₀ + 1;	+ 2.6097 6424 615 ₁₀ + 1;	+ 1.7183 5449 373 ₁₀ + 2;	+ 1.0000 0000 000 ₁₀ + 0;	- 1.0000 0000 000 ₁₀ + 0;
+ 2.8625 9509 072 ₁₀ + 1;	+ 1.5509 5804 016 ₁₀ + 1;	+ 8.7809 0481 507 ₁₀ + 0;	+ 1.0313 4644 279 ₁₀ + 2;	+ 1.0000 0000 000 ₁₀ + 0;
+ 2.9269 1592 858 ₁₀ + 1;	+ 2.2283 8464 931 ₁₀ + 1;	- 1.1468 4872 835 ₁₀ + 0;	- 3.9498 1287 698 ₁₀ + 0;	+ 1.6046 1044 072 ₁₀ + 2;

Table 2

<i>Output matrix b from reduc 1 and reduc 2 with A = F, B = G</i>		
+ 1.2000 0000 000 ₁₀ + 1;	+ 1.0000 0000 000 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
+ 2.8867 5134 595 ₁₀ - 1;	+ 1.4000 0000 000 ₁₀ + 1;	+ 1.0000 0000 000 ₁₀ + 0;
- 2.8867 5134 595 ₁₀ - 1;	+ 2.9039 8583 546 ₁₀ - 1;	+ 1.0000 0000 000 ₁₀ + 0;
+ 5.7735 0269 190 ₁₀ - 1;	- 3.1273 6936 127 ₁₀ - 1;	+ 1.0000 0000 000 ₁₀ + 0;
+ 2.8867 5134 595 ₁₀ - 1;	+ 2.4572 1878 385 ₁₀ - 1;	+ 1.0000 0000 000 ₁₀ + 0;
<i>Output vector d1 from reduc 1 and reduc 2 with A = F, B = G</i>		
	+ 3.4641 0161 514 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
	+ 3.7305 0488 093 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
	+ 3.9789 8672 143 ₁₀ + 0;	+ 1.0000 0000 000 ₁₀ + 0;
	+ 3.3961 8010 923 ₁₀ + 0;	- 1.0000 0000 000 ₁₀ + 0;
	+ 3.2706 8845 017 ₁₀ + 0;	+ 1.1000 0000 000 ₁₀ + 1;

Table 3. Eigenvalues

$A - \lambda B, A = F, B = G$	$A - \lambda B, A = G, B = F$	$AB - \lambda I, A = F, B = G$	$AB - \lambda I, A = G, B = F$
+ 4.3278 7211 020 ₁₀ - 1;	+ 6.7008 2644 107 ₁₀ - 1;	+ 7.7697 1911 953 ₁₀ + 1;	+ 7.7697 1911 963 ₁₀ + 1;
+ 6.6366 2748 402 ₁₀ - 1;	+ 9.0148 1958 801 ₁₀ - 1;	+ 1.1215 4193 247 ₁₀ + 2;	+ 1.1215 4193 246 ₁₀ + 2;
+ 9.4385 9004 670 ₁₀ - 1;	+ 1.0594 8027 732 ₁₀ - 0;	+ 1.3468 6463 320 ₁₀ + 2;	+ 1.3468 6463 320 ₁₀ + 2;
+ 1.1092 8454 002 ₁₀ + 0;	+ 1.5067 8940 837 ₁₀ + 0;	+ 1.6748 4878 917 ₁₀ + 2;	+ 1.6748 4878 915 ₁₀ + 2;
+ 1.4923 5323 254 ₁₀ + 0;	+ 2.3106 0432 137 ₁₀ + 0;	+ 2.4297 7273 320 ₁₀ + 2;	+ 2.4297 7273 319 ₁₀ + 2;

Table 4. Eigenvectors

<i>1st eigenvector of $F - \lambda G$</i>	<i>last eigenvector of $G - \lambda F$</i>	<i>1st eigenvector of $FG - \lambda I$ via $G = LL^T$</i>	<i>1st eigenvector of $FG - \lambda I$ via $F = LL^T$</i>
+ 1.3459 0573 962 ₁₀ - 1;	- 2.0458 6718 183 ₁₀ - 1;	+ 2.3491 1413 526 ₁₀ - 1;	+ 2.0706 5038 597 ₁₀ + 0;
- 6.1294 7224 718 ₁₀ - 2;	+ 9.3172 0977 419 ₁₀ - 2;	- 4.1091 5167 469 ₁₀ - 2;	- 3.6220 5325 515 ₁₀ - 1;
- 1.5790 2562 211 ₁₀ - 1;	+ 2.4002 2507 111 ₁₀ - 1;	- 3.8307 5945 797 ₁₀ - 2;	- 3.3766 6162 397 ₁₀ - 1;
+ 1.0946 5787 725 ₁₀ - 1;	- 1.6639 5354 480 ₁₀ - 1;	- 2.0590 0367 490 ₁₀ - 1;	- 1.8149 2958 995 ₁₀ + 0;
- 4.1473 0117 966 ₁₀ - 2;	+ 6.3041 7653 099 ₁₀ - 2;	- 7.3470 7965 853 ₁₀ - 2;	- 6.4761 5758 762 ₁₀ - 1;

reduc 2 and again using the Cholesky factorization of G and F respectively. Corresponding eigenvalues have a maximum disagreement of 1 in the eleventh decimal.

Table 4 gives the first eigenvector of $F - \lambda G$ and the last eigenvector of $G - \lambda F$. These should be the same apart from the normalization. The first is normalized so that $x^T G x = 1$ and the second so that $x^T F x = 1$. The angle between the vectors is smaller than 10^{-11} radians. It also gives the first eigenvector of $FG - \lambda I$ obtained via the factorization of G and via the factorization of F . Again the angle between the two vectors is smaller than 10^{-11} radians. These results are typical of those obtained for all five eigenvectors in each case.

The complete sets of results were subjected to extensive cross checking and proved generally to be "best possible" for a 39 digit binary computer. A second test was carried out by C. REINSCH at the Technische Hochschule, München, and confirmed the above results.

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