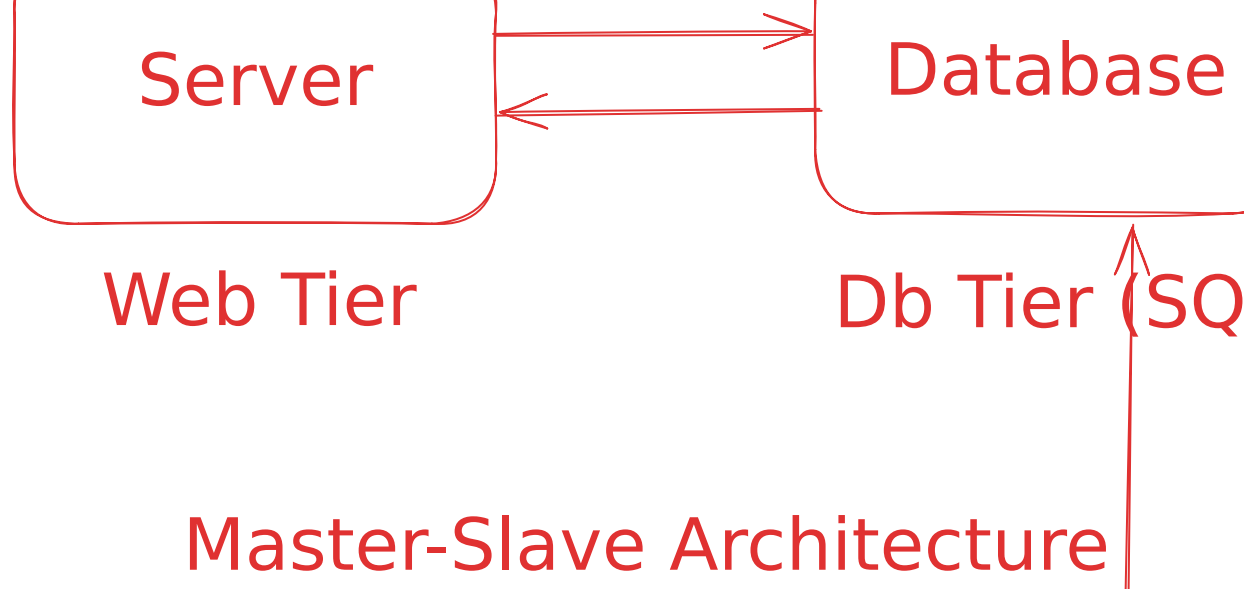


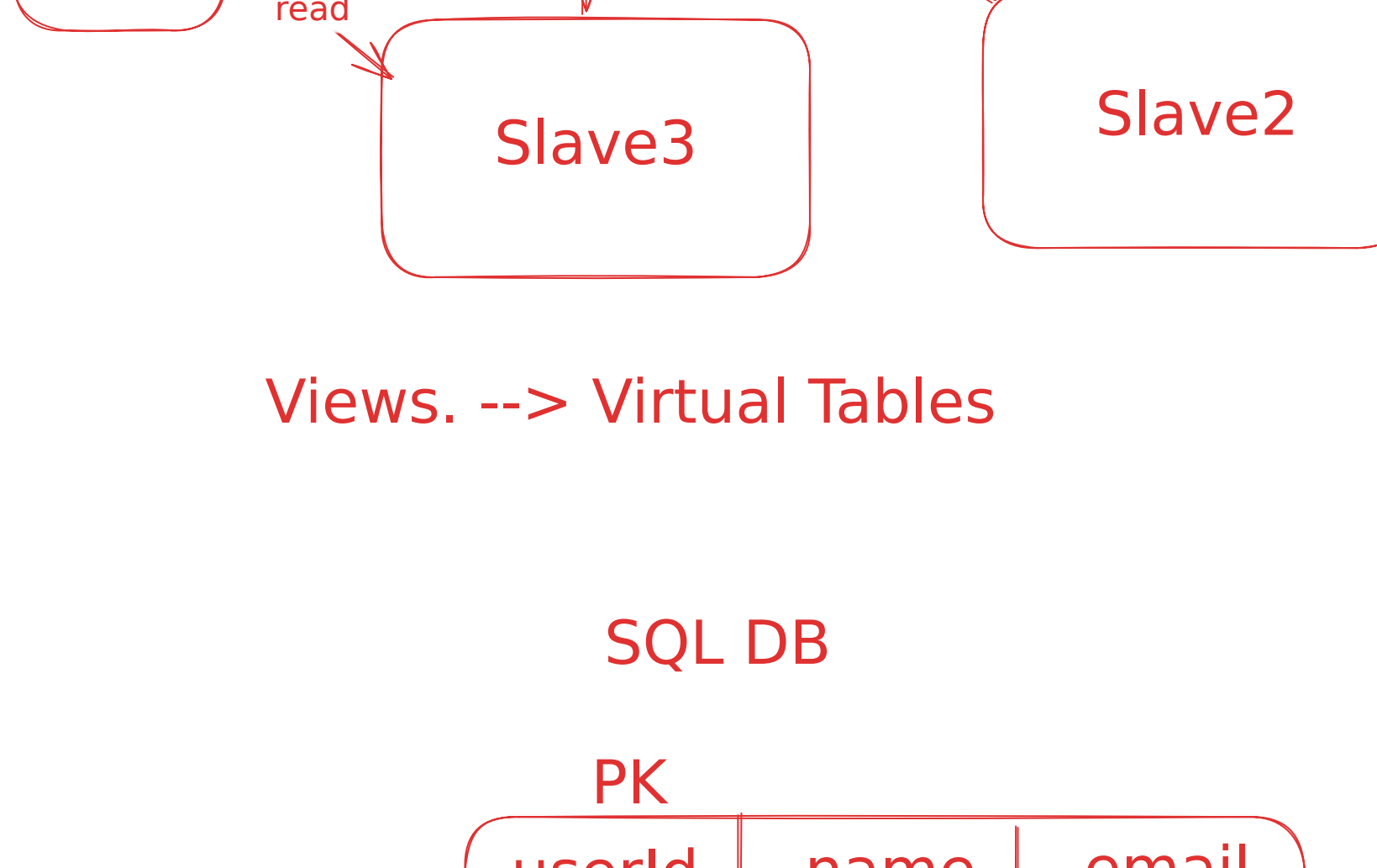
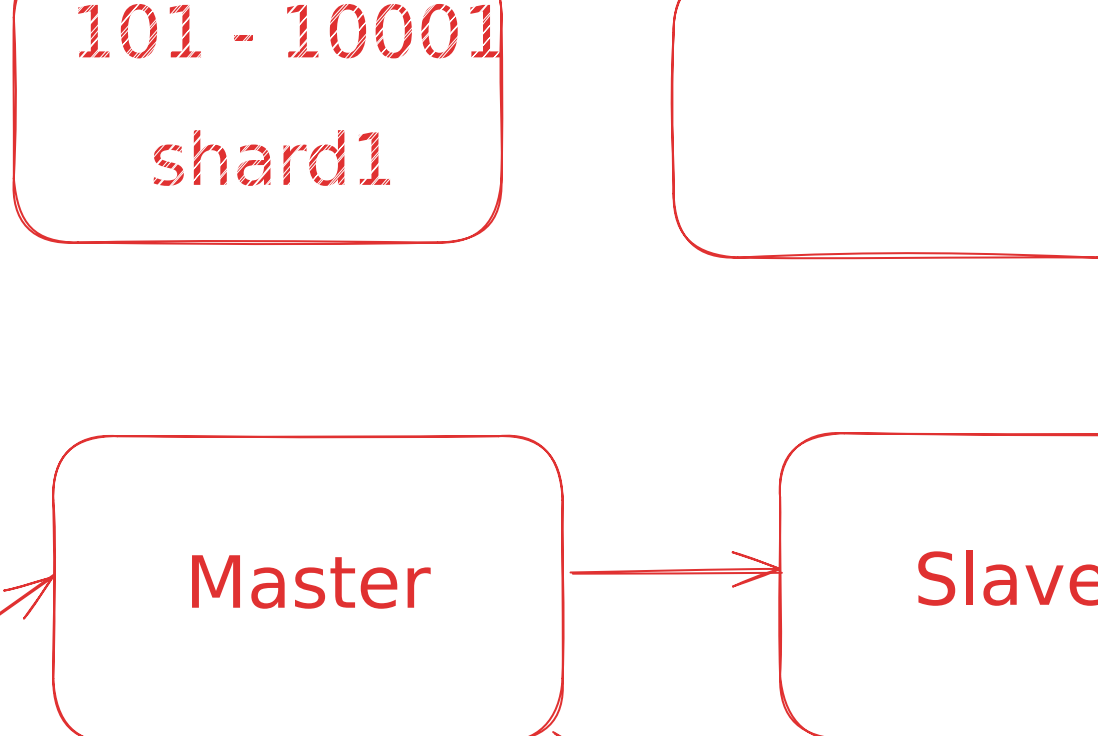
# SQL / Relational Architecture



User

PK	userId	name	email
	101	Aditya	xyz
	102	Rohit	ert
	103	Mukesh	wer

Horizontal Scaling --> Sharding



Views. --> Virtual Tables

## SQL DB

User

PK	userId	name	email
	101	Aditya	xyz
	102	Rohit	ert
	103	Mukesh	wer

Relational DB. --> Rows and Columns

Logical Arrangement

Stored not in Volatile memory

Stored in permanent memory (Disc)



Data Block

Sector

DBMS

Table

PK	userId	name	email
	101	Aditya	xyz
	102	Rohit	ert
	103	Mukesh	wer

Data Page

Headers

96 bytes

Data Records

8 KB = 8192 Bytes

offset

36 bytes

1 Million Rows

row = 32 bytes

8060 bytes

1 Data Page = ~251 rows.

Total page. = 3984 pages



Page Number, checksum

Array of indexed rows

Data pages are managed by DBMS. But practically, All data pages will be stored in memory.

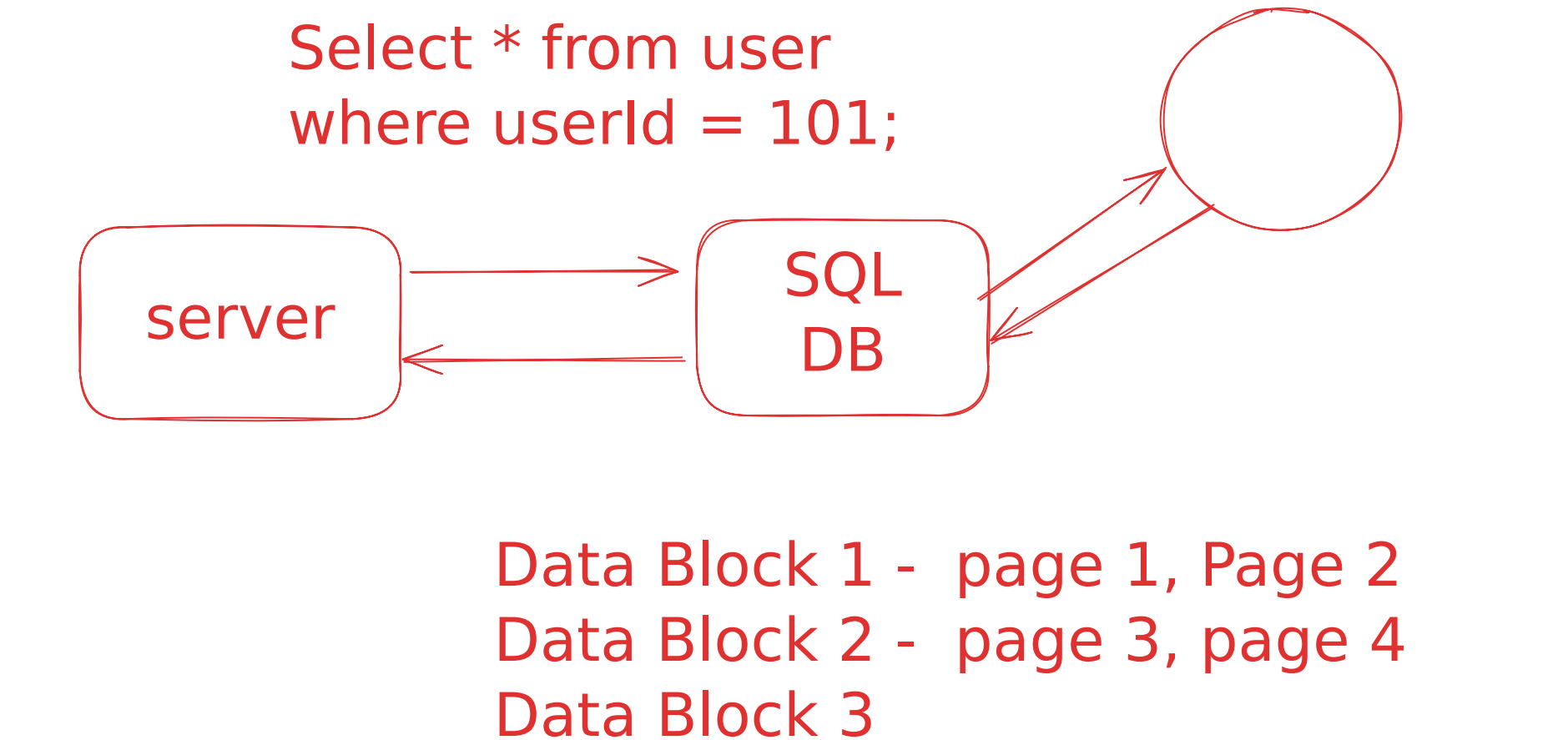
1 Data Block --> 8 KB - 64 KB

If 1 Data block = 8KB

1 Data block = 1 Data Page

Data Block :

It is the smallest possible value we can expect from an IO operation.



Select \* from user where userId = 101;



Data Block 1 - page 1, Page 2  
Data Block 2 - page 3, page 4  
Data Block 3

O(N) time complexity --> 10 million

Data Blocks are managed by Hard disk. DBMS manages Data Pages.

- ✓ DBMS has a mapping table (User)
- ✓ Data page 1 --> Data block 1
- ✓ Data page 2 --> Data block 1
- ✓ Data page 3 --> Data block 2
- ....
- Data page 3984 --> Data block 1992

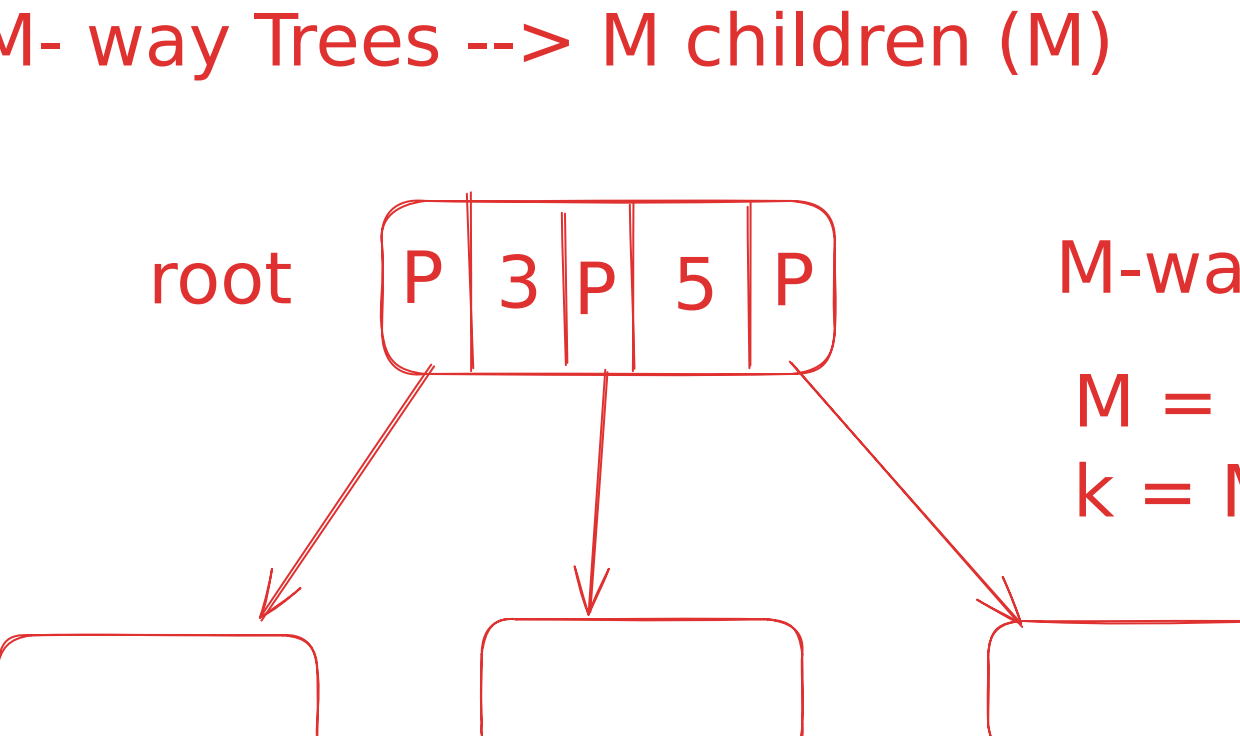
Indexing : used to increase the performance of our Db Queries. --> Making Queries run faster.

O(N) --> Reduce my search result.

DSA --> Binary Search /Binary Search Tree (BST)

B+ Trees are used for DB indexing.

O(N) --> O(logN)



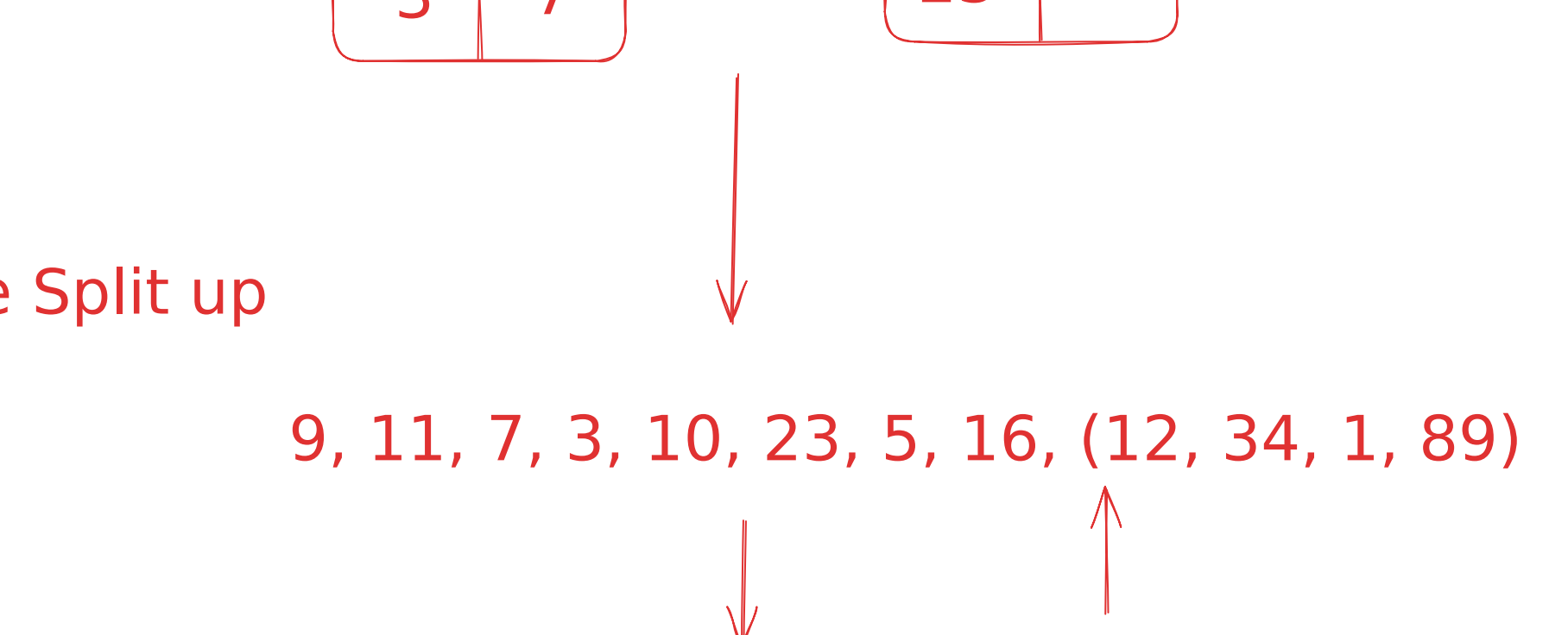
## B Trees

9, 11, 7, 3, 23, 5, 16



BST {  
int data;  
BST p1;  
BST p2;  
}

M-way  
M = 2  
Keys = 1 (M-1)



M-way tree  
M = 3  
k = M-1 (3-1) = 2

{

int data1;

int data2;

Btree p1;

Btree p2;

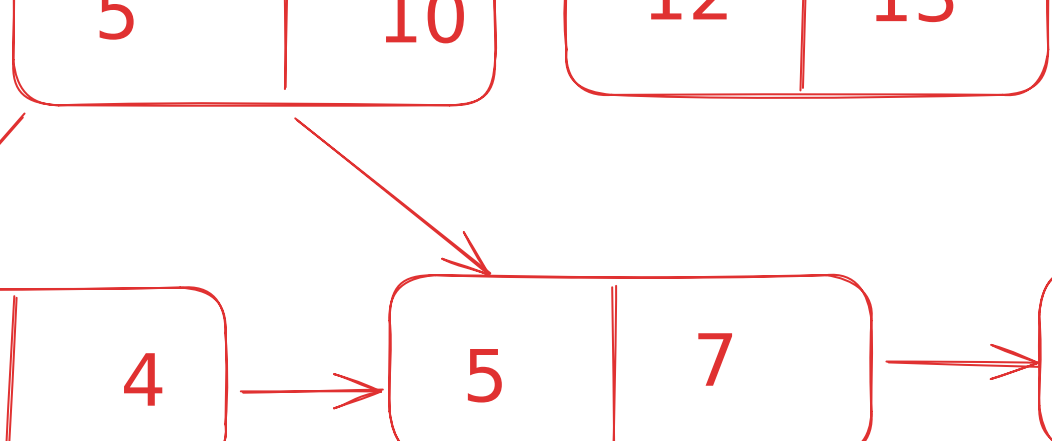
btree p3;

}

## B Trees

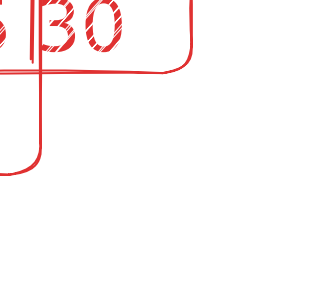
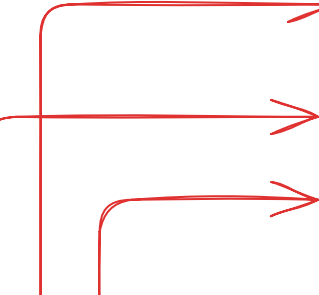
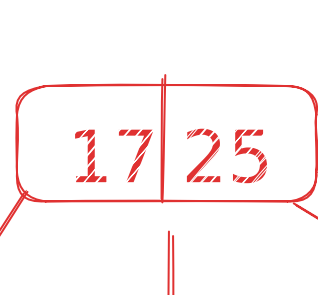
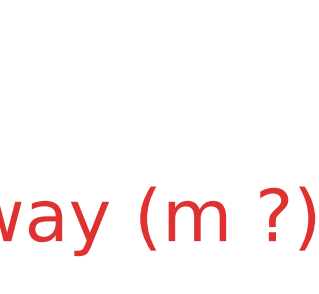
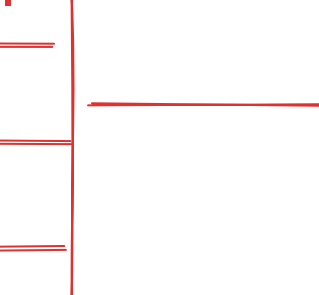


M-Way tree  
M = 3  
keys = M-1 = 2



B-Tree Split up

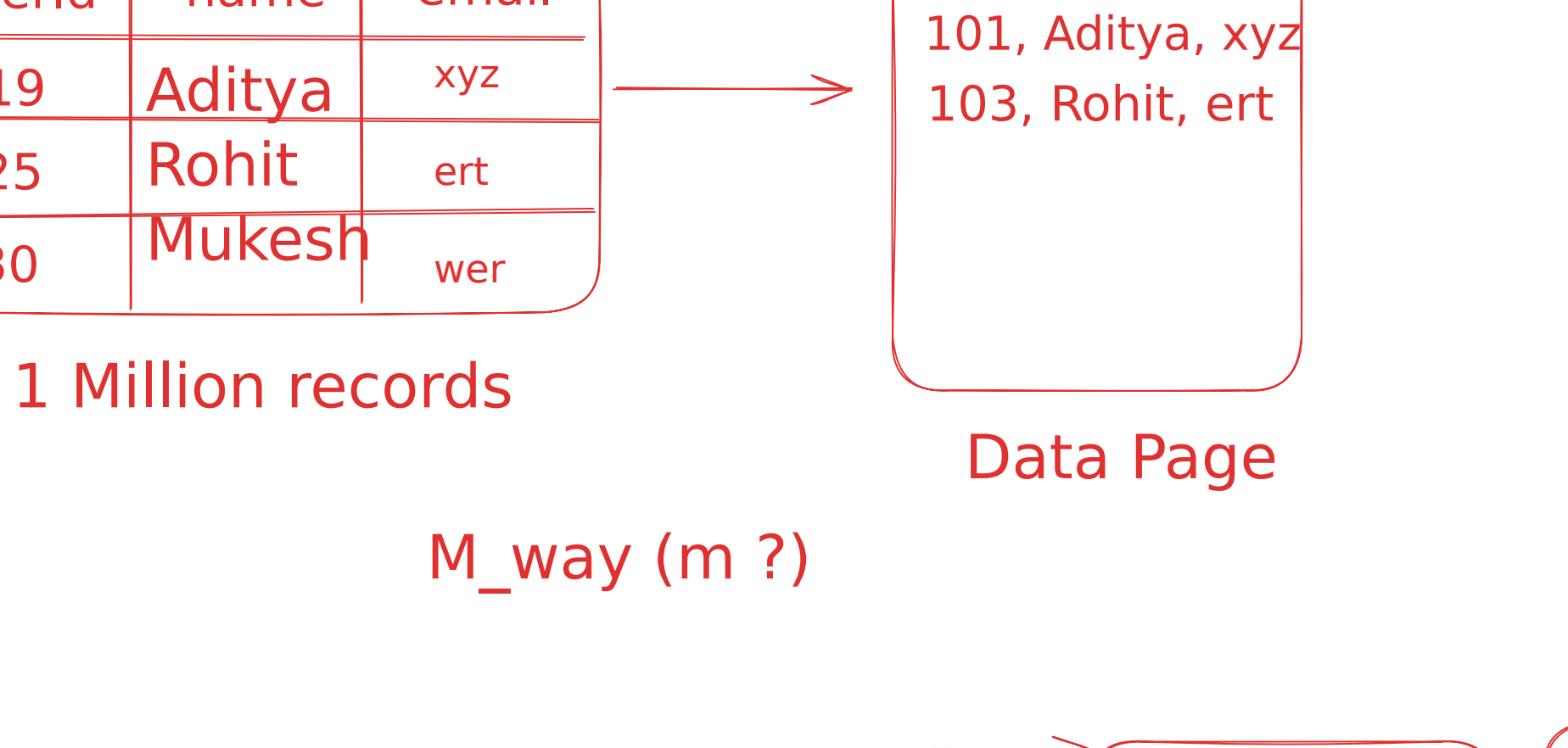
9, 11, 7, 3, 10, 23, 5, 16, (12, 34, 1, 89)



B+ Trees

leaf nodes connected

Intermediate nodes also store parent left value



B-Tree  
B+ Tree  
Code

Indexing

B+ Trees

user

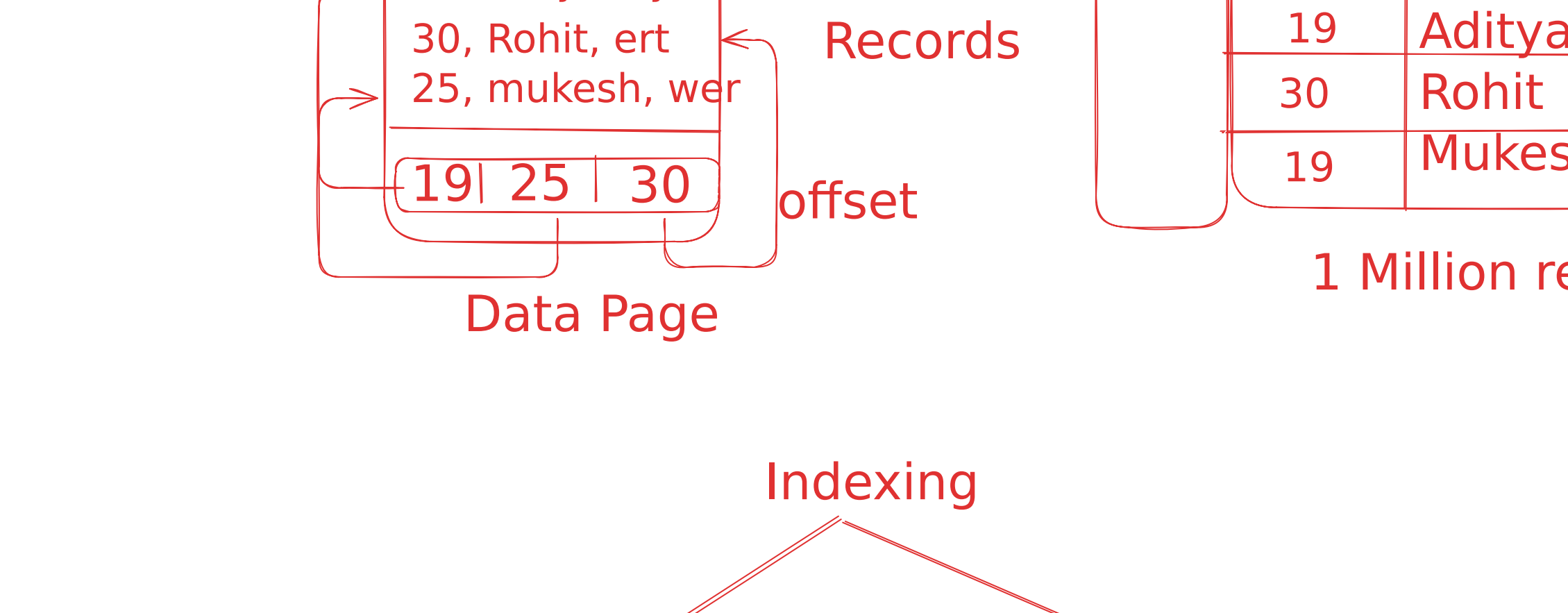
PK	userId	name	email
	19	Aditya	xyz
	25	Rohit	ert
	30	Mukesh	wer

1 Million records

101, Aditya, xyz  
103, Rohit, ert

Data Page

M\_way (m ?)



Data page 1

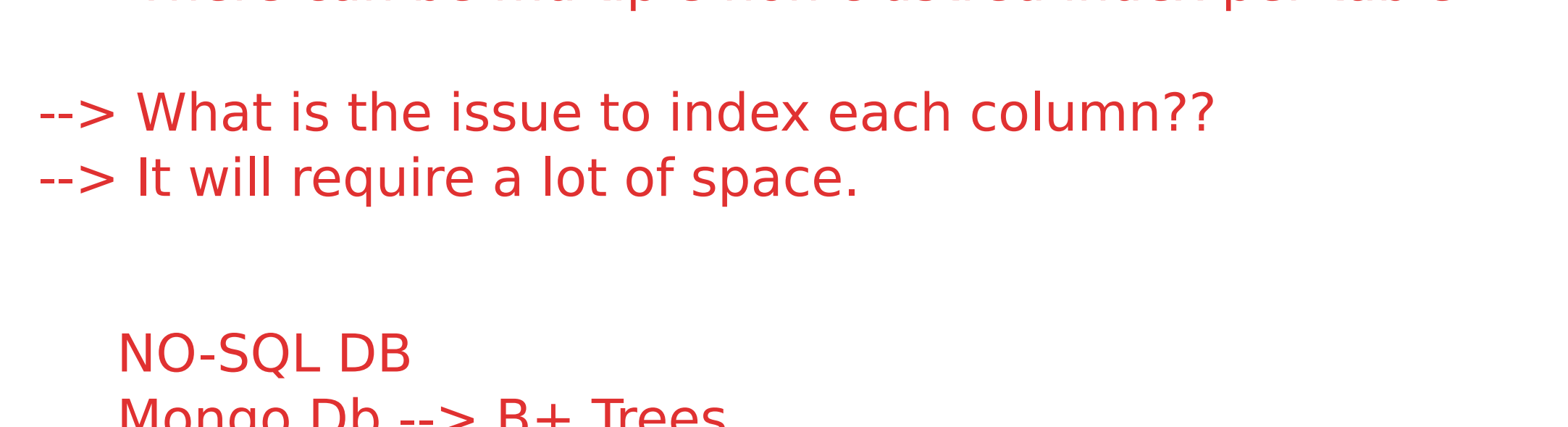
250 values

select \* from user where userId = 19;

DBMS data page --> Data block

## B+ Tree

Range Query --> 101 - 205



1 Million records

Indexing

Clustered

Non Clustered

Leaf nodes of B+ Trees store complete row information

leaf nodes of B+ trees store only id information (page no.)

Indexing :

Clustered vs Non-Clustered

--> There can be only 1 clustered index per Table.

--> There can be multiple non clustered index per table.

--> What is the issue to index each column??

--> It will require a lot of space.

## NO-SQL DB

Mongo Db --> B+ Trees

Cassandra DB --> LSM Tree