

DAY
43
DSA

17 Oct-2021

To understand recursion, one must first understand Recursion
—Stephen Hawking

Recursion, cont....

Agenda for Today

1. Factorial → CW
2. Power linear → CW
3. Power logarithmic → CW
4. Print Zigzag → pending
5. Tower of Hanoi → pending

Ques 1.

Factorial

A number n will be given to us.
We need to find the factorial of n .

The factorial of a number is the product of all the integers from 1 to that number.

For eg \rightarrow factorial of $n = 4$, is $4 * 3 * 2 * 1$

Very
Imp.
Point

You must know that the Factorial of both 0 and 1 is 1.

$$\begin{aligned}L0 &= 1 \\L1 &= 1\end{aligned}$$

Factorial of negative integers do not exist.

High Level Thinking

Expectation

Faith

Expectation meets Faith

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Expectation	Faith	Expectation meets faith
$f(5) = 5 * 4 * 3 * 2 * 1$	$f(4) = 4 * 3 * 2 * 1$	$f(5) = 5 * f(4)$.

1. Self Expectation

We expect that ~~AMR~~ input $5 \in \mathbb{N}$ so $n=5$

Output देगा $\rightarrow f(5) = 5 \times 4 \times 3 \times 2 \times 1$

2. Build Faith :- ऐसे ही faith से ना चलें कि
AMR हमारा code $n=5$ के लिए output देसकता है
तो $n=4$ के लिए यह output जबर देगा।

You just need to believe this.

Don't focus on how this will happen.
वहस मानों की factorial for $n=4$ AMR देगा

$$f(4) = 4 \times 3 \times 2 \times 1$$

3. Expectation meets Faith :-

For printing the desired output for $n=5$.
we could just print "5" before the output
for $n=4$.

$$f(5) = 5 \times f(4)$$

We are done with High Level Thinking

→ Now, we need to do the Low Level Thinking
to find the base case.

Dayrum & Drave Stack

Why is a base case Important?

If we don't have a base condition

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that check if the recursion function should be called or not then the program wouldn't know when to stop.

अगर कोई base condition नहीं होती तो program कभी समाप्त होना नहीं प्रogram infinitely चला रहा।

हमें ये नहीं चाहिए

तो हमें एक base condition या base case, जिससे रुकनी होती ताकि हमारे program को पैर पता हो कि अब ऐसी कोई condition सच होगी या जब मेरे पास ऐसा कोई case होगा तो मुझे इसके जाना है।

तो अब हम low level thinking करते हैं।

तो इसे पता चलता है
हमारा Base Case

Low Level Thinking

जोड़े दी लाइन
a, b, c जोड़े
लाइन डिफरेंस
तो ये function

जितना भी
memory
करें

कुरवे थे,
यो सब वापस
दूर दूर कर दी

bcz now these
variables are no

longer of use JT Java Garbage Collector garbage collect करता।

using the Stacking process

Using the
Stacking
Process

संतत
abc
तो यहाँ
abc रुकता है

Dry run &
draw stack

if ($n == 0$) { return; }
line a ↓
int fnm1 = factorial($n - 1$);
line b ↓
int fn = n * fnm1;
line c ↓
return fn;

m	fnm1	fn
0		
1	1	1
2	1	2
3	2	6
4	6	24
5	24	120

m fnm1 fn

अगर function void ना हो (मतलब अगर return type
void ना हो (क्षेत्र और हो return type
like int/string/char) etc).

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तो हमेशा याद रख कि recursive call लगाते
हुए result को receive करना क्षी न भूल।

unit fn = factorial(n-1);

ऐसे receive
करना भी ज़रूरी नहीं यहाँ ans.print नहीं
होगा।

FACTORIAL <code>

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```
→ import java.util.*;  
→ import java.io.*;  
→ public class Main  
{|  
→ }  
→ public static void main(String[] args)  
{|  
    Scanner scan = new Scanner(System.in);  
    int n = scan.nextInt();  
    int fn = factorial(n);  
    System.out.println(fn);  
}|  
//Expectation → f(5) → 5 * 4 * 3 * 2 * 1  
//Faith → f(4) = 4 * 3 * 2 * 1.  
//Expectation meet faith = 5 * f(4).
```

```
→ public static int factorial(int n)
```

```
|  
{| if (n == 0)  
|}|
```

```
| return 1; |
```

```
|  
| factorial of n minus 1 |
```

```
| int fnm1 = factorial(n-1); // n-1 - 3x2x1  
| int fn = n * fnm1 // n x n-1 x — 3x2x1.
```

```
| return fn; |
```

```
|
```

```
|
```

Powerlinear

→ We are given a number x and a no. n , you are required to calculate x raised to the power n i.e. x multiplied n times

$$\text{eg} \rightarrow x = 5 \\ n = 4$$

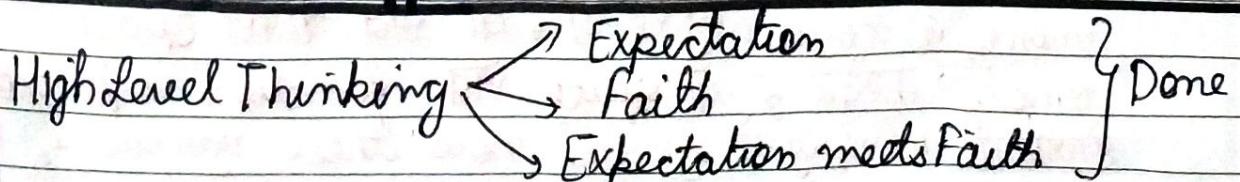
$$\text{Power}(x, n) = 5 \times 5 \times 5 \times 5 = (5)^4$$

↓ Generalizing it

$$\boxed{\text{power}(x, n) = x * x * x - x(n \text{ times}) = x^n}$$

→ High Level Thinking

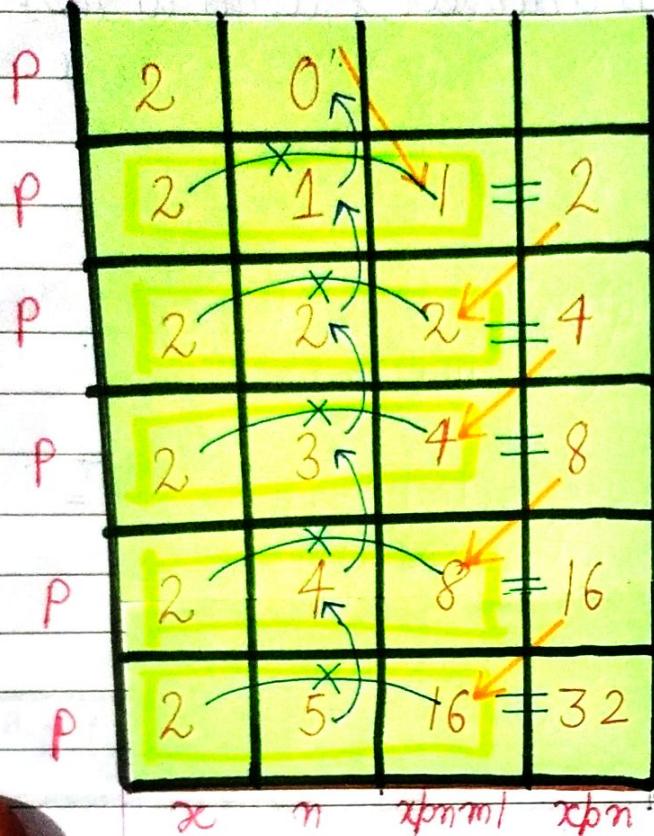
Expectation	Faith	Expectation meets Faith
$\text{Power}(x, n)$. $= x * x * x$ $- x * x (n \text{ times})$ $= x^n$	$\text{Power}(x, n-1)$. $= x * x * x * x$ $- x * x (-1 \text{ times})$ $= x^{n-1}$	$\text{Power}(x, n)$. $= x * (x^{n-1})$ $= x * \text{Power}(x, n-1)$



Low Level Thinking

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It's done to find the base case.



$x=2$
 $n=5$

it becomes linear approach.
इसमें recursive calls
के Tower की height $n+1$
होती है।

~~it becomes recursive calls~~
~~ERET~~
∴ Time complexity
is $O(n+1)$

```
public static int power(int x, int n)
{
    if (n == 0)
        { return 1; }
}
```

Time Complexity

$$O(n)$$

because we are
recursively calling
subproblem with $n-1$
from n , hence $n+1$
recursive call will
be made (+1)

Space Complexity

$O(1)$ → No data structure used: no extra space used

जैसे ही line a, b, c तीनों line एटा func की
तो ये function के सभी variables ने liftoff की
memory cover कर रहे थे, तो सब वापस space
empty कर देंगे, तो सब functions were out
of stack memory से because now these variables are
no longer of use, तो Java Garbage Collector जैसे as
Garbage collect कर देगा।

Power & linear <code>

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```
→ import java.util.*;
→ import java.io.*;
→ public class Main
{
    → public static void main(String[] args)
    {
        Scanner scan = new Scanner(System.in);
        int x = scan.nextInt();
        int n = scan.nextInt();
        xpn → x raised to the power n
        int xpn = power(x, n);
        System.out.println(xpn);
        } // Expectation → p(2,5) = 2·2·2·2·2
        // Faith → p(2,4) = 2·2·2·2
        // Expectation meets faith → p(2,4) = 2
    → public static int power(int x, int n)
    {
        if (n == 0)
            return 1;
        } // x to the power n minus 1
        int xpnm1 = power(x, n-1);
        int xpn = xpnm1 * x;
        } // x to the power n
        return xpn;
    }
```

Power Logarithmic

→ We are given a number x and a number n , you are required to calculate x raised to the power n .

→ If x^n as output provide POT^n
for input x, n .

→ The Input & Output of Power linear & Power Logarithmic is same.

→ Then what is the difference between power linear & power logarithmic

Power linear

For even & odd positive powers

$$\rightarrow x^n = x \cdot x^{n-1}$$

$$\text{eg} \rightarrow x=2$$

$$n=8$$

$$2^8 = 2 \cdot 2^7 \rightarrow 256$$

$$2^7 = 2 \cdot 2^6 \rightarrow 128$$

$$2^6 = 2 \cdot 2^5 \rightarrow 64$$

$$2^5 = 2 \cdot 2^4 \rightarrow 32$$

$$2^4 = 2 \cdot 2^3 \rightarrow 16$$

$$2^3 = 2 \cdot 2^2 \rightarrow 8$$

$$2^2 = 2 \cdot 2^1 \rightarrow 4$$

$$2^1 = 2 \cdot 2^0 \rightarrow 2$$

$$2^0 = 1 \rightarrow 1$$

Power logarithmic

for even +ve powers

$$\rightarrow x^n = x^{n/2} * x^{n/2}$$

for odd +ve powers

$$\rightarrow x^n = x * x^{n/2} * x^{n/2}$$

x extra
multiply x if n is odd

Integer of Case II

$$2^8 = 2^4 \cdot 2^4 \rightarrow 256$$

$$2^4 = 2^2 \cdot 2^2 \rightarrow 16$$

$$2^2 = 2^1 \cdot 2^1 \rightarrow 4$$

$$2^1 = 2^0 \cdot 2^0 \cdot 2^1 \rightarrow 2$$

Pow	2	0	bcz n is odd
Pow	2	1	$\rightarrow (1^2 - 1) = 2$
Pow	2	2	$\rightarrow (2^2 - 4) = 4$
Pow	2	4	$\rightarrow (4^2 - 16) = 16$
Pow	2	8	$\rightarrow (16^2 - 256) = 256$
	x^n	x^{pb2}	x^n

Recursive
call

Power Linear \rightarrow Time Complexity $\rightarrow O(n)$

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→ Since you are recursively calling for a subproblem with $n-1$ from n , hence $n+1$ recursive calls will be made ($+1$ when it hits the base case $n=0$). You can also count the number of recursive calls in the call stack. So, the time complexity turns out to be $O(n)$ (which is independent of the value of x).

Power Logarithmic \rightarrow Time Complexity $\rightarrow O(\log_2 n)$

→ Since you are recursively calling for a subproblem with $n/2$ from n , hence at max $\log_2(n)$ recursive calls will be made. You can also count the number of recursive calls in the call stack. So, the time complexity turns out to be $\log_2(n)$ (which is independent of value of x).
Space Complexity = $O(1)$ (bcz no extra data structure used.)

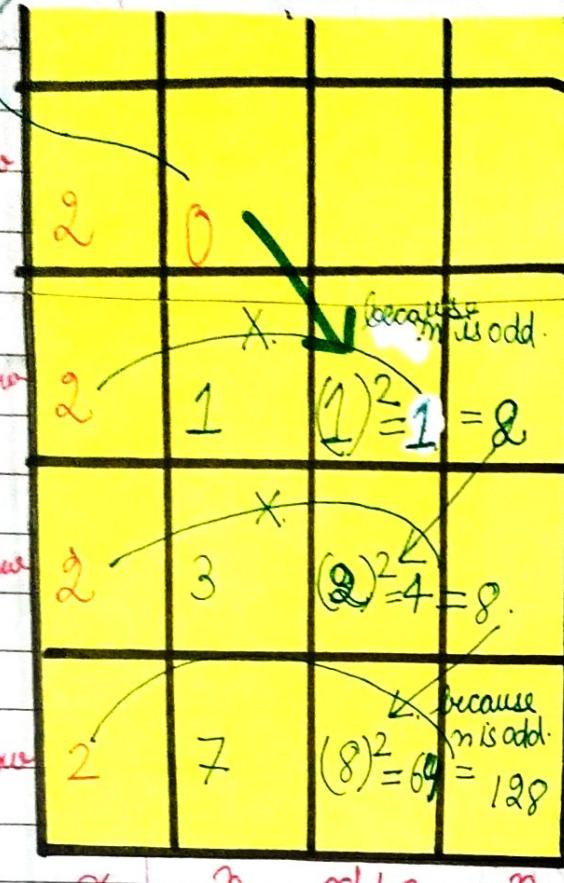
High Level Thinking of Power Logarithmic

Expectation	Faith	Expectation meets faith
<p>If n is even $x^n = x^{n/2} * x^{n/2}$</p> <p>If n is odd $x^n = x * x^{n/2} * x^{n/2}$</p>	<p>$P(x, n) = x * x * x -$ --- $x * x - \frac{x^n}{2}$</p> <p>times 3 $= x^{\frac{n}{2}}$</p>	<p>If n is even $P(x, n) = P(x, \frac{n}{2}) * P(x, \frac{n}{2})$</p> <p>If n is odd $P(x, n) = x * P(x, \frac{n}{2}) * P(x, \frac{n}{2})$</p>

So, we are done with High Level Thinking

Low Level Thinking

→ Base Case



dine a
dine b

Final b

Line C

public static void power
(int x, int n)

{ ref(n == 0) { return 1; } }

unit xpb2 = power(x, n/2);

$$\text{unt } x_{fbn} = x_{fb62} * x_{fb62}$$

$\text{if } (n \% 2 == 1)$

$$\{ \exp n = xpn * x_0^n \}$$

3 4 1

3

abcd

abcd

abcd

→ Height = 4

$$\therefore \log_2 7$$

Do the Square of $x^2 b^2$
& if n is odd then multiply by x .

Pseudocode for Power Logarithmic

- 1) If n is 0, then return 1.
 2. else
 - a.) Get the value of $\text{Power}(x, \frac{n}{2})$ in a variable named " x^{pnb2} ".
 - b.) If n is odd, then return $x * x^{pnb2} * x^{pnb2}$
 - c.) Else (n is even), return $x^{pnb2} * x^{pnb2}$

PowerLogarithmic <code>

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```
→ import java.util.*;  
→ import java.io.*;  
→ public class Main  
S  
S {  
    public static void main (String [] args)  
    S {  
        Scanner sc = new Scanner (System.in);  
        int x = sc.nextInt();  
        int n = sc.nextInt();  
        int xpn = power (x, n);  
        L x to the power  
        raised .
```

System.out.println (xpn);

If n is even If n is odd.
//Expectation $\rightarrow P(x^n) = (x^{n/2}) * P(x^{n/2})$ | $P(x^n) = * P(x^{\frac{n}{2}}) * P(x^{\frac{n}{2}})$
//Faith $\rightarrow P(x^{n/2}) = x * x - \frac{n}{2}$ times $= x^{\frac{n}{2}}$
//Expectation meets faith $= power(x, n) = power(x, \frac{n}{2})$

If n is even. * power (x, n)

If n is odd. $\begin{cases} power(x, n) = x * \\ power(x, \frac{n}{2}) * \\ power(x, \frac{n}{2}). \end{cases}$

public static int power (int x, int n).

{ if (n == 0) { return 1; }
 L x to the power by 2

int xpb2 = power (x, n/2);

int xpn = xpb2 * xpb2;

if (n % 2 == 1) { xpn = xpn * n; }

return xpn;