

# Projection of Varieties and Elimination Ideals

Applications: Word-Level Abstraction from Bit-Level Circuits,  
Combinational Verification, Reverse Engineering Functions from Circuits

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Nov 19-24, 2014

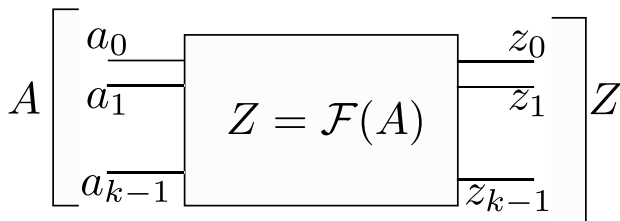
# We will employ everything we have learnt so far....

- Hilbert's Nullstellensatz over  $\mathbb{F}_q$
- Gröbner basis theory
- Efficient term ordering from circuits
- Canonical representations of circuits  $f : \mathbb{B}^k \rightarrow \mathbb{B}^k$  to  $f : \mathbb{F}_{2^k} \rightarrow \mathbb{F}_{2^k}$

And learn a new concept: Elimination ideals

- Apply these techniques to circuit analysis and verification

# Polynomial Interpolation from Circuits



- Circuit:  $f : \mathbb{B}^k \rightarrow \mathbb{B}^k$
- Model it as a polynomial function  $f : \mathbb{F}_{2^k} \rightarrow \mathbb{F}_{2^k}$
- Interpolate a word-level polynomial from the circuit:  $Z = \mathcal{F}(A)$
- Obtain  $Z = \mathcal{F}(A)$  as a **unique, canonical, word-level, polynomial** representation from the *bit-level* circuit
- Why?

# Hierarchical Abstraction and Verification

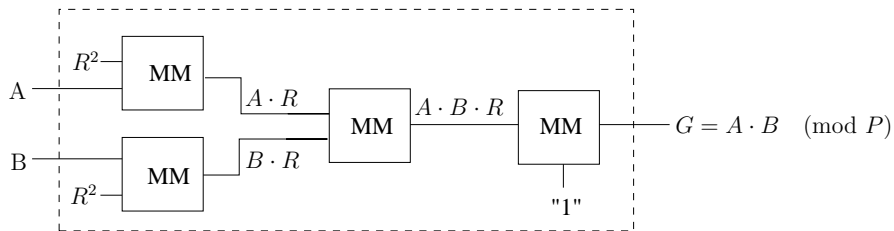
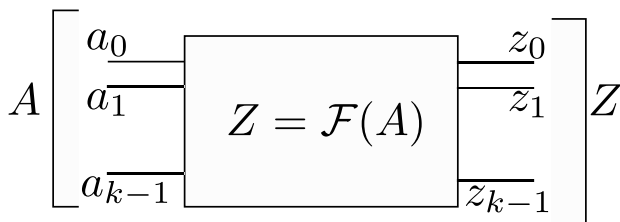


Figure: *Montgomery multiplier over  $\text{GF}(2^k)$*

Montgomery Multiply:  $F = A \cdot B \cdot R^{-1}$ ,  $R = \alpha^k$

# Projection of Variety



- Represent the polynomials of the circuit as ideal  $J$  (or  $J + J_0$ )
- Consider  $V_{\mathbb{F}_q}(J)$
- Let  $x_i$  denote the bit-level variables of the circuit:  $J \subset \mathbb{F}_q[x_i, Z, A]$
- **Project**  $V_{\mathbb{F}_q}(J)$  on  $Z, A$ , denoted by  $V_{\mathbb{F}_q}(J)|_{Z,A}$ 
  - Does this recover the function of the circuit?

## Definition

Given variety  $V = \mathbf{V}(f_1, \dots, f_s) = \mathbf{V}(J) \subset \mathbb{F}_q^n$ . The  $l^{th}$  projection map  $\pi_l : \mathbb{F}_q^n \rightarrow \mathbb{F}_q^{n-l}$ ,  $\pi_l((c_1, \dots, c_n)) = (c_{l+1}, \dots, c_n)$

- We may also denote  $\pi_l$  by  $\text{Proj}[V(J)]_{l+1, \dots, n}$ , or by  $V(J)|_{l+1, \dots, n}$
- In some sense, we have eliminated the first  $l$  variables from the system
- This is related to **elimination ideals** and variable elimination

## Definition (*Elimination Ideal*)

Given  $J = \langle f_1, \dots, f_s \rangle \subset \mathbb{F}_q[x_1, \dots, x_n]$ , the  $l$ th *elimination ideal*  $J_l$  is the ideal of  $\mathbb{F}_q[x_{l+1}, \dots, x_n]$  defined by  $J_l = J \cap \mathbb{F}_q[x_{l+1}, \dots, x_n]$ .

In other words, the  $l$ th elimination ideal does not contain variables  $x_1, \dots, x_l$ , nor do the generators of it.

## Theorem (*Elimination Theorem*)

Let  $J \subset \mathbb{F}_q[x_1, \dots, x_n]$  be an ideal and let  $G$  be a Gröbner basis of  $J$  with respect to a lex ordering where  $x_1 > x_2 > \dots > x_n$ . Then for every  $0 \leq l \leq d$ , the set  $G_l = G \cap \mathbb{F}_q[x_{l+1}, \dots, x_n]$  is a Gröbner basis of the  $l$ th elimination ideal  $J_l$ .

# A Gröbner basis example [From Cox/Little/O'Shea]

Solve the system of equations over  $\mathbb{C}$ :

$$f_1 : x^2 - y - z - 1 = 0$$

$$f_2 : x - y^2 - z - 1 = 0$$

$$f_3 : x - y - z^2 - 1 = 0$$

Gröbner basis  $G$  with lex term order  $x > y > z$

$$g_1 : x - y - z^2 - 1 = 0$$

$$g_2 : y^2 - y - z^2 - z = 0$$

$$g_3 : 2yz^2 - z^4 - z^2 = 0$$

$$g_4 : z^6 - 4z^4 - 4z^3 - z^2 = 0$$

- $G_1 = G \cap \mathbb{C}[y, z] = \{g_2, g_3, g_4\}$
- $G_2 = G \cap \mathbb{C}[z] = \{g_4\}$
- $G$  is *triangular*: solve  $g_4$  for  $z$ , then  $g_2, g_3$  for  $y$ , and then  $g_1$  for  $x$ 
  - Solutions to  $z$  are  $0, 1, -1 + \sqrt{2}, -1 - \sqrt{2}$
  - $V(G) = \{(1, 0, 0), (0, 1, 0), (0, 0, 1), (-1 + \sqrt{2}, -1 + \sqrt{2}, -1 + \sqrt{2}), (-1 - \sqrt{2}, -1 - \sqrt{2}, -1 - \sqrt{2})\}$



# Projection of Variety and Elimination Ideals

- Using elimination, obtain partial solution to  $V(I_I)$ , then extend it to  $V(I)$ , one variable at a time
- However, all partial solutions to  $V(I_I)$  **may not lift** to  $V(I)$

## Example

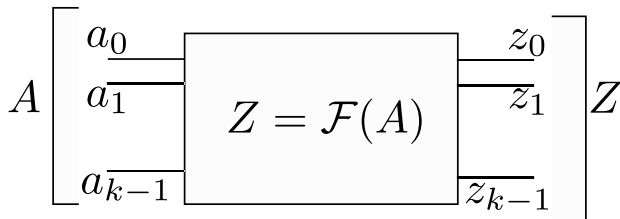
Consider  $f_1 : xy - 1$ ,  $f_2 : xz - 1$ . Eliminate  $x$ , you get  $f_3 : y - z$ . All points  $(a, a)$  are solutions to  $f_3$ . All points  $(1/a, a, a)$  extend to complete solutions, except  $(0, 0)$ .

Given  $J = \langle f_1, \dots, f_s \rangle \subseteq \mathbb{F}[x_1, \dots, x_n]$ ,  $\pi_I(V(J)) \subset V(J_I)$   
In other words,  $\text{Proj}[V(J)]_{x_{I+1}, \dots, x_n} \subset V(J_I)$

## Theorem (Over $\mathbb{F}_q$ Elimination ideals give Projection exactly)

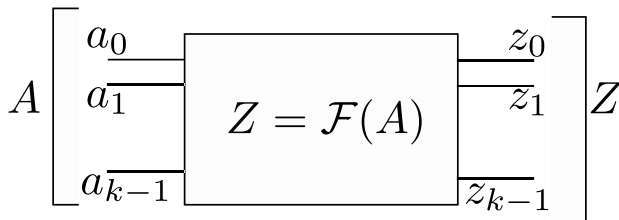
Over Galois fields,  $\mathbb{F}_q$ , let  $J$  be any ideal, and  $J_0$  be the ideal of vanishing polynomials. Let  $I = J + J_0$ . The projection of variety is **equal** to the variety of the elimination ideal. In other words,  $\pi_I(V(I)) = V(I_I)$ .

# Abstraction from Circuits

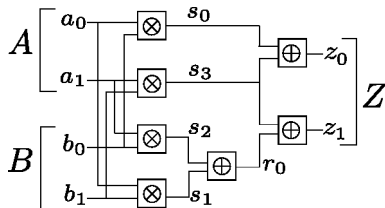


- To obtain,  $Z = \mathcal{F}(A)$ :
- Denote  $x_i$  as bit-level variables,  $A, Z$  as word-level variables
- Obtain  $J + J_0$  from the circuits
- Compute Gröbner basis  $G$  with lex order with  $x_i > Z > A$
- $G_{x_i}$  be the elimination ideal that eliminates  $x_i$
- Projection of variety onto  $Z, A$  is equal to  $V(G_{x_i})$ ,
- This recovers the function of the circuit  $Z = \mathcal{F}(A)$

# Abstraction from Circuits



- $G$  is computed with  $\text{lex } x_i > Z > A$
- There exists a polynomial  $A^q - A$  in  $G$
- There exists a polynomial  $Z = \mathcal{F}(A)$  in  $G$ 
  - Why? Can you prove it?
- The rest is irrelevant for us



$f_1 : z_0 + z_1\alpha + Z$ ;  $f_2 : b_0 + b_1\alpha + B$ ;  $f_3 : a_0 + a_1\alpha + A$ ;  $f_4 :$   
 $s_0 + a_0 \cdot b_0$ ;  $f_5 : s_1 + a_0 \cdot b_1$ ;  $f_6 : s_2 + a_1 \cdot b_0$ ;  $f_7 : s_3 + a_1 \cdot b_1$ ;  $f_8 :$   
 $r_0 + s_1 + s_2$ ;  $f_9 : z_0 + s_0 + s_3$ ;  $f_{10} : z_1 + r_0 + s_3$ . Ideal  $J = \langle f_1, \dots, f_{10} \rangle$ .

Add  $J_0$  and compute  $GB(J + J_0)$  with  $x_i > Z > A > B$ , then  $G :$

$g_1 : z_0 + z_1\alpha + Z$ ;  $g_2 : b_0 + b_1\alpha + B$ ;  $g_3 : a_0 + a_1\alpha + A$ ;  $g_4 :$   
 $s_3 + r_0 + z_1$ ;  $g_5 : s_1 + s_2 + r_0$ ;  $g_6 : s_0 + s_3 + z_0$ ;  $g_7 : Z + AB$ ;  $g_8 :$   
 $a_1b_1 + a_1B + b_1A + z_1$ ;  $g_9 : r_0 + a_1b_1 + z_1$ ;  $g_{10} : s_2 + a_1b_0$

# To Conclude

- Lex orders are elimination orders, but Deglex and DegRevLex are not elimination orders
- Computing GB with Lex orders is hard, gives very large output
- One can use block orders (I will give you a singular file with a block order)
- Projection of varieties can be solved exactly using Elimination ideals over Galois fields, not so over  $\mathbb{R}, \mathbb{Q}, \mathbb{C}$