DYNAMIC SCHEDULING

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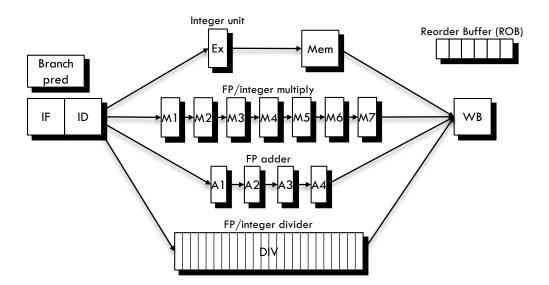
Overview

- Announcement
 - Homework 3 will be uploaded tonight

- This lecture
 - Dynamic scheduling
 - Forming data flow graph on the fly
 - Register renaming
 - Removing false data dependence
 - Architectural vs. physical registers

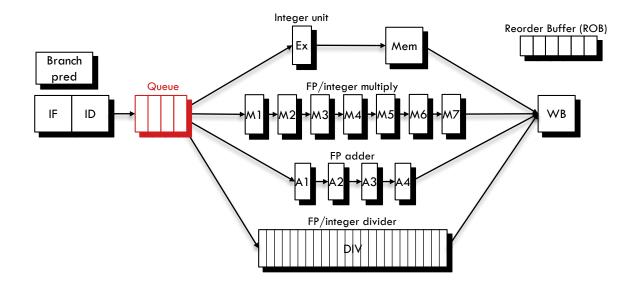
Big Picture

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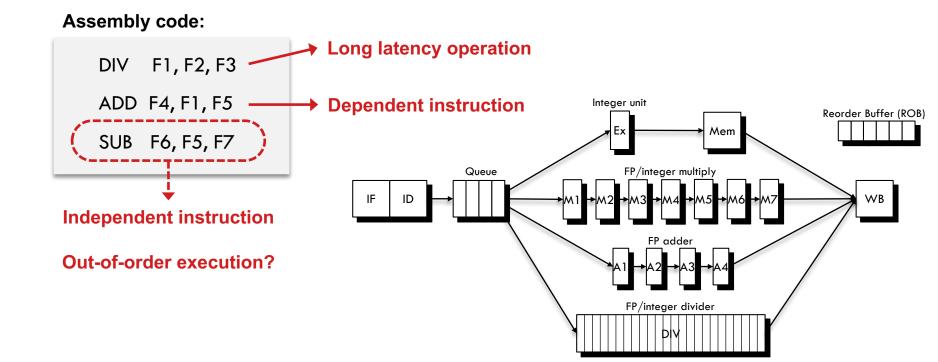


Big Picture

- Goal: exploiting more ILP by avoiding stall cycles
 - Branch prediction can avoid the stall cycles in the frontend
 - More instructions are sent to the pipeline
 - Instruction scheduling can remove unnecessary stall cycles in the execution/memory stage
 - Static scheduling
 - Complex software (compiler)
 - Unable to resolve all data hazards (no access to runtime details)
 - Dynamic scheduling
 - Completely done in hardware

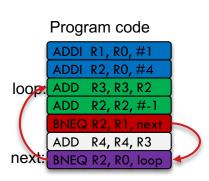
Dynamic Scheduling

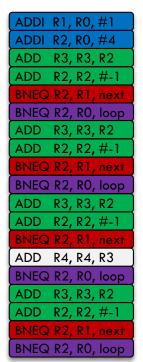
- Key idea: creating an instruction schedule based on runtime information
 - Hardware managed instruction reordering

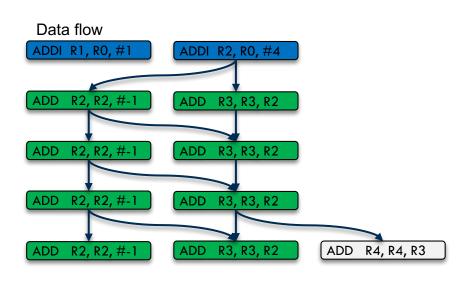


Dynamic Scheduling

- Key idea: creating an instruction schedule based on runtime information
 - Hardware managed instruction reordering
 - Instructions are executed in data flow order

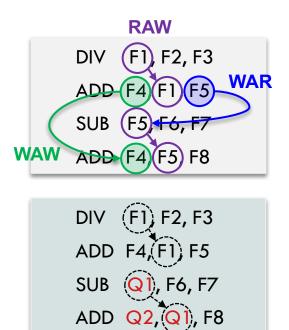




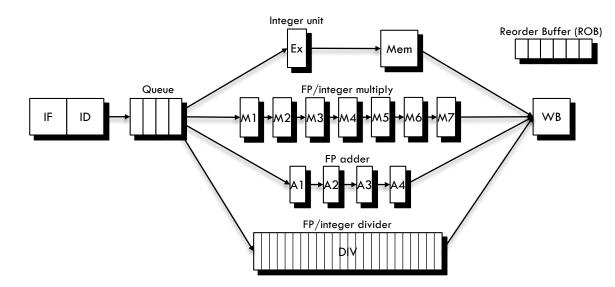


How to form data flow graph on the fly?

- Eliminating WAR and WAW hazards
 - Change the mapping between architectural registers and physical storage locations



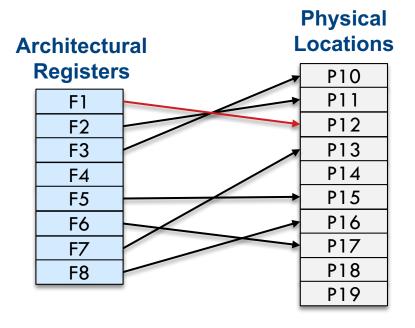
WAR and WAW hazards can be removed using more registers



- Eliminating WAR and WAW hazards
 - 1. allocate a free physical location for the new register
 - 2. find the most recently allocated location for the register

DIV F1, F2, F3
ADD F4, F1, F5
SUB F5, F6, F7
ADD F4, F5, F8

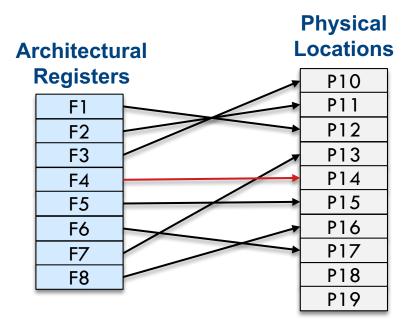
DIV P12, P11, P10



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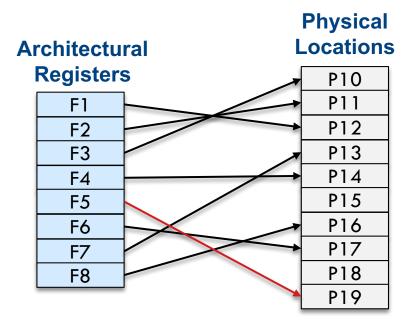
DIV P12, P11, P10 ADD P14, P12, P15



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DIV P12, P11, P10 ADD P14, P12, P15 SUB P19, P17, P13



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DIV P12, P11, P10 ADD P14, P12, P15 SUB P19, P17, P13 ADD P18, P19, P16

