

# Progressive Generation of Canonical Sums of Products Using a SAT Solver

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Austin, Texas

# Sum of Products (SOP) Representation

- ▶ **Sum of products (SOP)** is a two-level representation that can represent a Boolean function as a sum (OR, +) of cubes,  $S = c_1 + \dots + c_m$
- ▶ **Cube**  $c$  is a Boolean product (AND,  $\cdot$ ) of literals,  $c_i = l_1 \cdot \dots \cdot l_k$
- ▶ **Literal**  $l$  is a variable  $v$  or its negation  $\bar{v}$

Truth table				On-set SOP			
$x_1$	$x_2$	$x_3$	$F$	$x_1$	$x_2$	$x_3$	$f$
0	0	0	0	-	1	1	1
0	0	1	0	1	1	-	1
0	1	0	0				
0	1	1	1				
1	0	0	0				
1	0	1	0				
1	1	0	1				
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$f = x_2x_3$

minterms

cubes

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minterms

On-set SOP

$x_1$	$x_2$	$x_3$	$f$
-	1	1	1
1	1	-	1

cubes

$$f = x_2x_3 + x_1x_2$$

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minterms

On-set SOP

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Off-set SOP

$x_1$	$x_2$	$x_3$	$\bar{f}$
0	0	-	0
0	-	0	0
1	0	-	0

$$\bar{f} = \bar{x}_1 \bar{x}_2$$

cubes

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1	0	1	0				
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Off-set SOP			
$x_1$	$x_2$	$x_3$	$\bar{f}$
0	0	-	0
0	-	0	0
1	0	-	0

$f = x_2x_3 + x_1x_2$

$\bar{f} = \bar{x}_1\bar{x}_2 + \bar{x}_1\bar{x}_3$

minterms

cubes

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minterms

On-set SOP

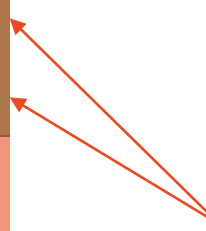
$x_1$	$x_2$	$x_3$	$f$
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1	1	-	1

$$f = x_2x_3 + x_1x_2$$

Off-set SOP

$x_1$	$x_2$	$x_3$	$\bar{f}$
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$$\bar{f} = \bar{x}_1 \bar{x}_2 + \bar{x}_1 \bar{x}_3 + x_1 \bar{x}_2$$



cubes

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minterms

On-set SOP

$x_1$	$x_2$	$x_3$	$f$
-	1	1	1
1	1	-	1

$$f = x_2x_3 + x_1x_2$$

SOP size: 2 cubes, 4 literals

smaller SOP

Off-set SOP

$x_1$	$x_2$	$x_3$	$\bar{f}$
0	0	-	0
0	-	0	0
1	0	-	0

$$\bar{f} = \bar{x}_1 \bar{x}_2 + \bar{x}_1 \bar{x}_3 + x_1 \bar{x}_2$$

SOP size: 3 cubes, 6 literals

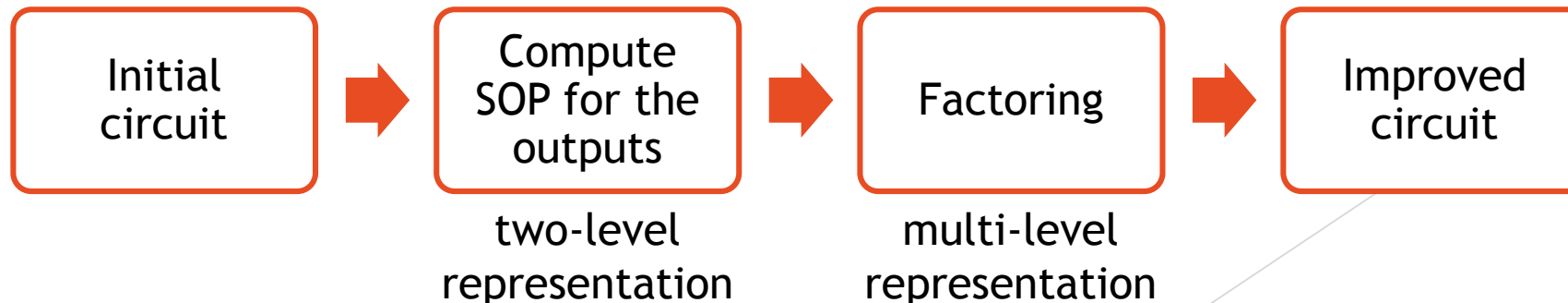
cubes

# Use of SOPs

- ▶ Mapping into Programmable Logic Arrays (PLAs)

Supported in many tools for logic optimization (in academia and industry)

- ▶ Delay optimization in technology independent synthesis and technology mapping
- ▶ Test generation
- ▶ Multi-level logic synthesis for global circuit restructuring (case-study)





# Existing SOP Generation and Minimization

- ▶ From Binary Decision Diagrams (BDDs)
  - **Problem 1:** BDD generation is not scalable for some functions
    - The BDD size is exponential in the number of input variables → BDD memory explosion problem
  - **Problem 2:** Incompatible with incremental applications
    - The complete BDD should be build before converting it to SOP
- ▶ ESPRESSO style minimizers: version enhanced by SAT solving
  - **Problem 1:** Requires as input a computed SOP: do not generate a new SOP from a multi-level representation
  - **Problem 2:** The existing SOP generation is based on enumeration on satisfying assignments

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  - **Problem 1:** Requires as input a computed SOP: do not generate a new SOP from a multi-level representation
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Solution?

# SAT-based SOP Generation and Minimization

- From Binary Decision Diagrams (BDDs)

Feasibility of computation

BDD generation is not scalable for some functions

size is exponential in the number of input variables → BDD memory explosion problem

Progressive generation

Incompatible with incremental applications

Partial SOPs

complete BDD should be build before converting it to SOP

- ESPRESSO style minimizers: version enhanced by SAT solving

Completely based on SAT solving

Requires as input a computed SOP: do not generate a new SOP from a multi-level function

The existing SOP generation is based on enumeration on satisfying assignments

Solution?

# SAT-based SOP Generation and Minimization

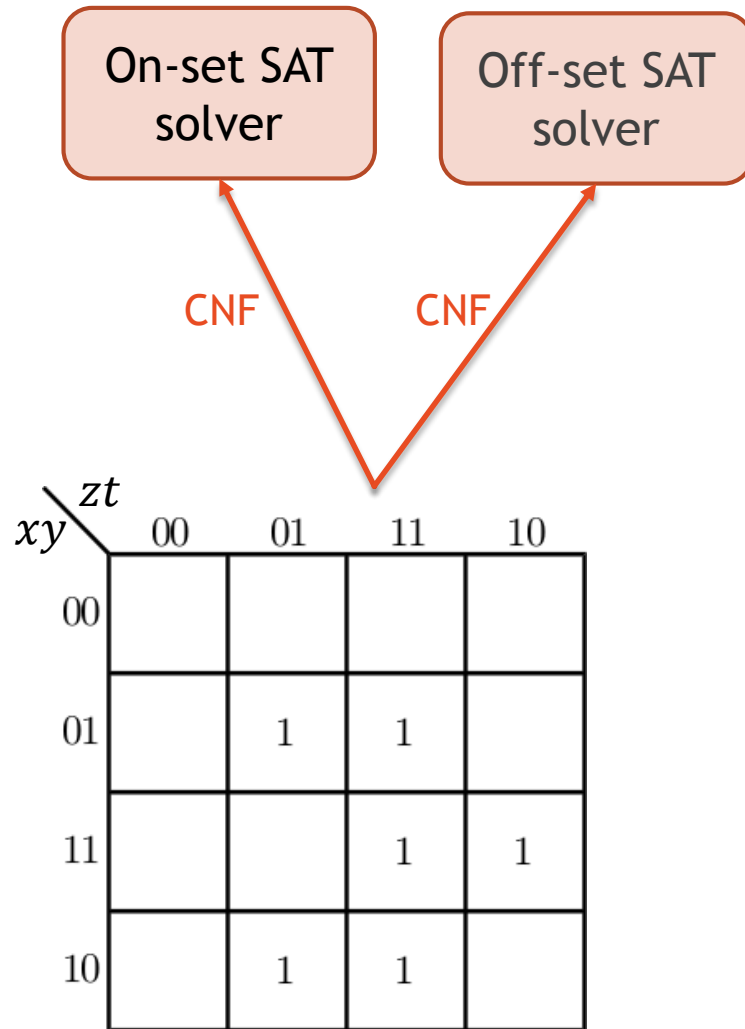
- ▶ Completely based on SAT solvers
- ▶ **Progressive generation:** cube by cube (generated and minimized)
  - Generation of a **partial SOP**
  - Prediction of **feasibility of computation**
- ▶ SAT-based generation of **canonical SOPs**
  - Unique SOP for *a given function* and *variable order*, independent of
    - The initial structure of the circuit
    - The used SAT solver
    - CNF generation algorithm
    - Operating system
- ▶ The generated SOPs are **irredundant**
  - No literal and no cube can be removed without changing the function

# Outline

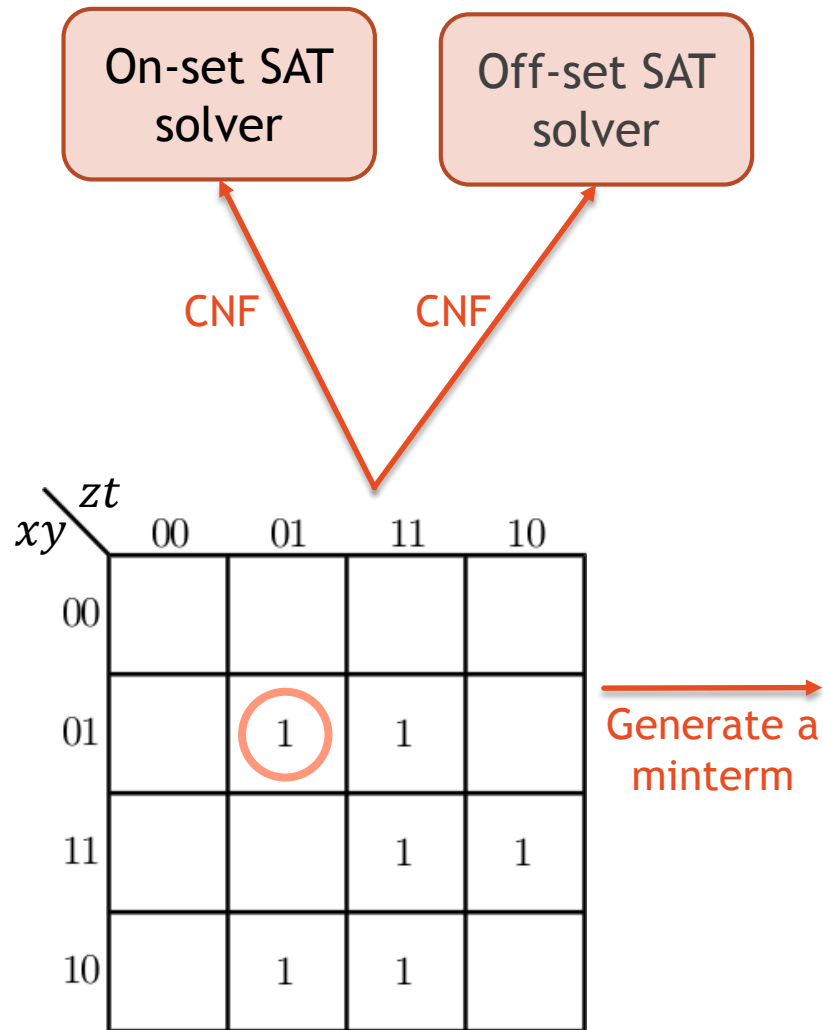
- ▶ SAT-based SOP Generation
- ▶ Methods for Runtime Improvement
- ▶ Experimental Results
- ▶ Conclusion

# Progressive SAT-based SOP Generation: Global view

- Initialize SAT solvers

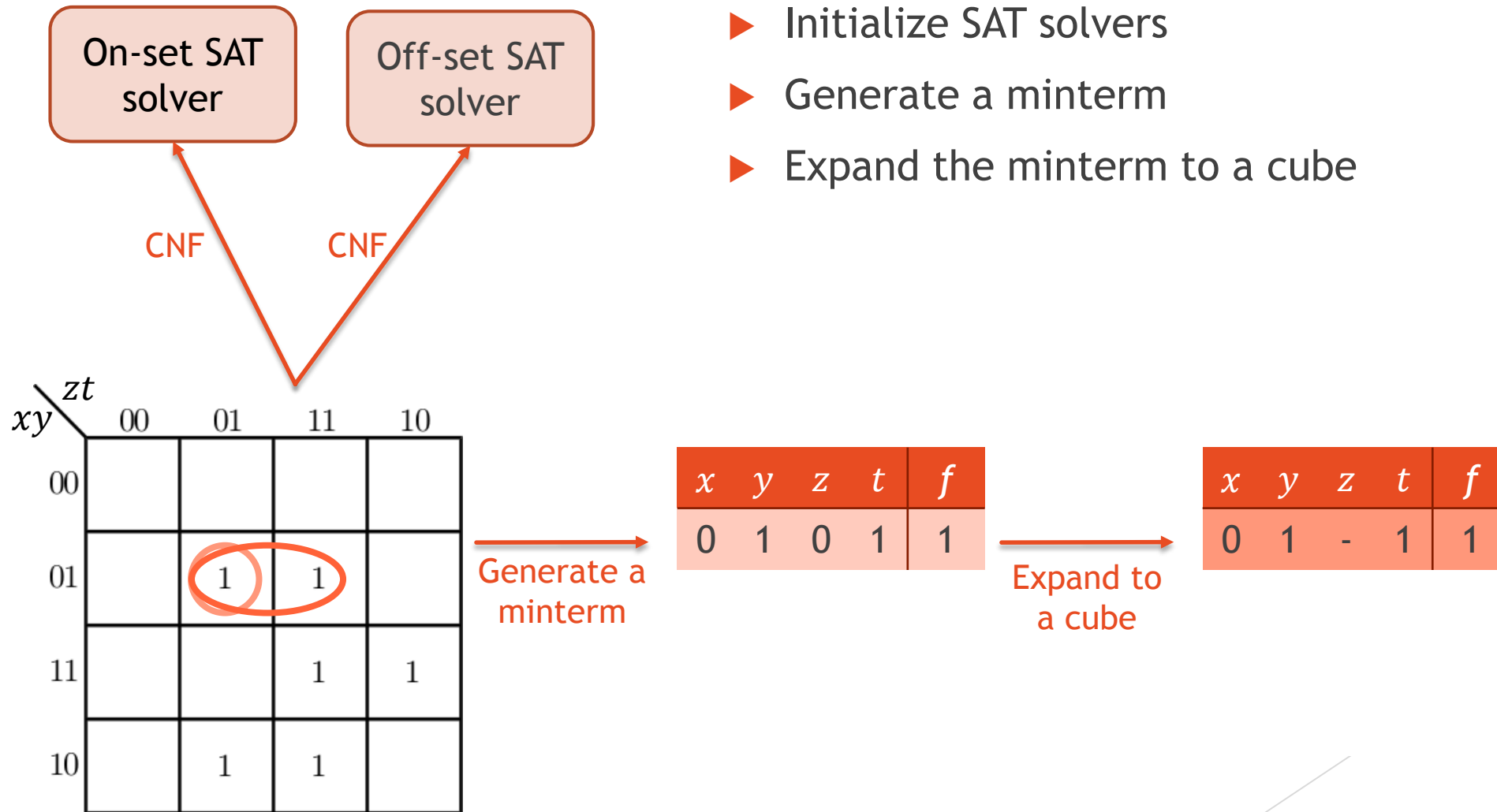


# Progressive SAT-based SOP Generation: Global view



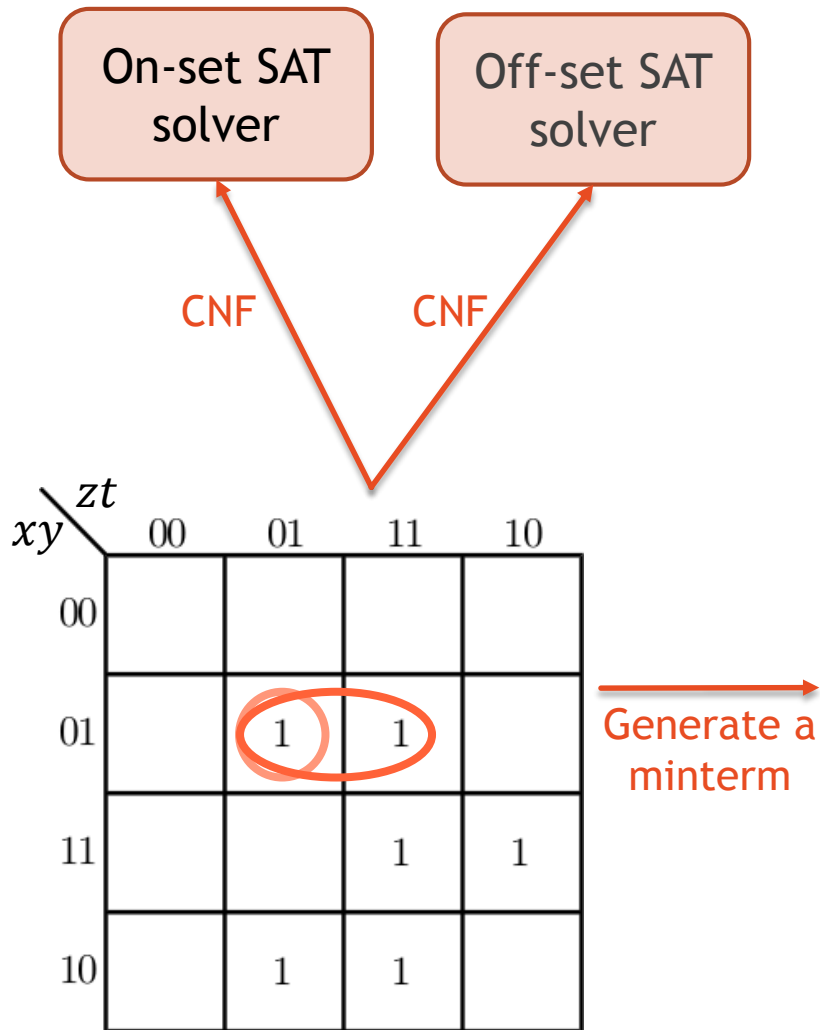
- ▶ Initialize SAT solvers
- ▶ Generate a minterm

# Progressive SAT-based SOP Generation: Global view





# Progressive SAT-based SOP Generation: Global view



<i>x</i>	<i>y</i>	<i>z</i>	<i>t</i>	<i>f</i>
0	1	0	1	1

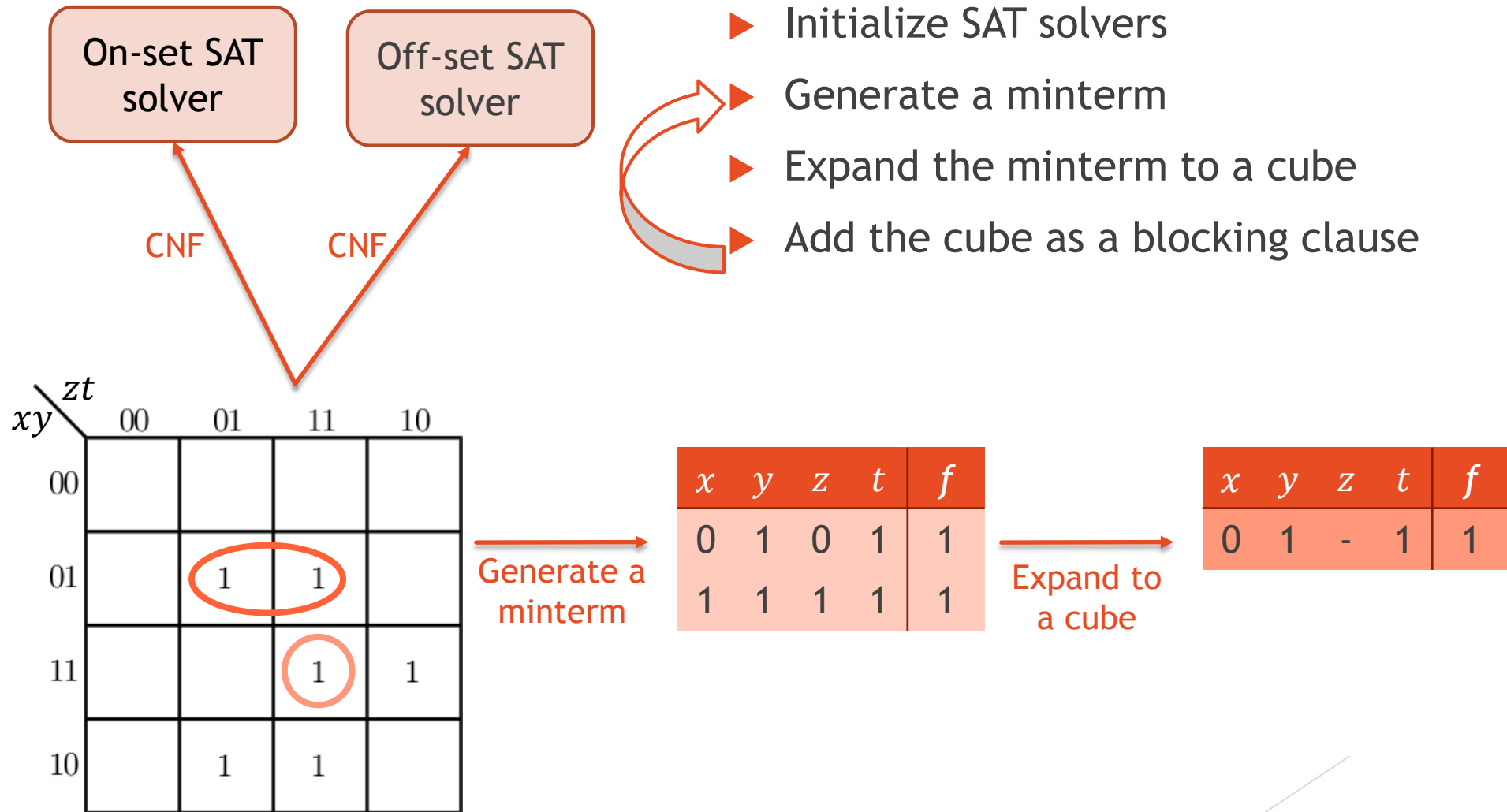
Expand to a cube

<i>x</i>	<i>y</i>	<i>z</i>	<i>t</i>	<i>f</i>
0	1	-	1	1

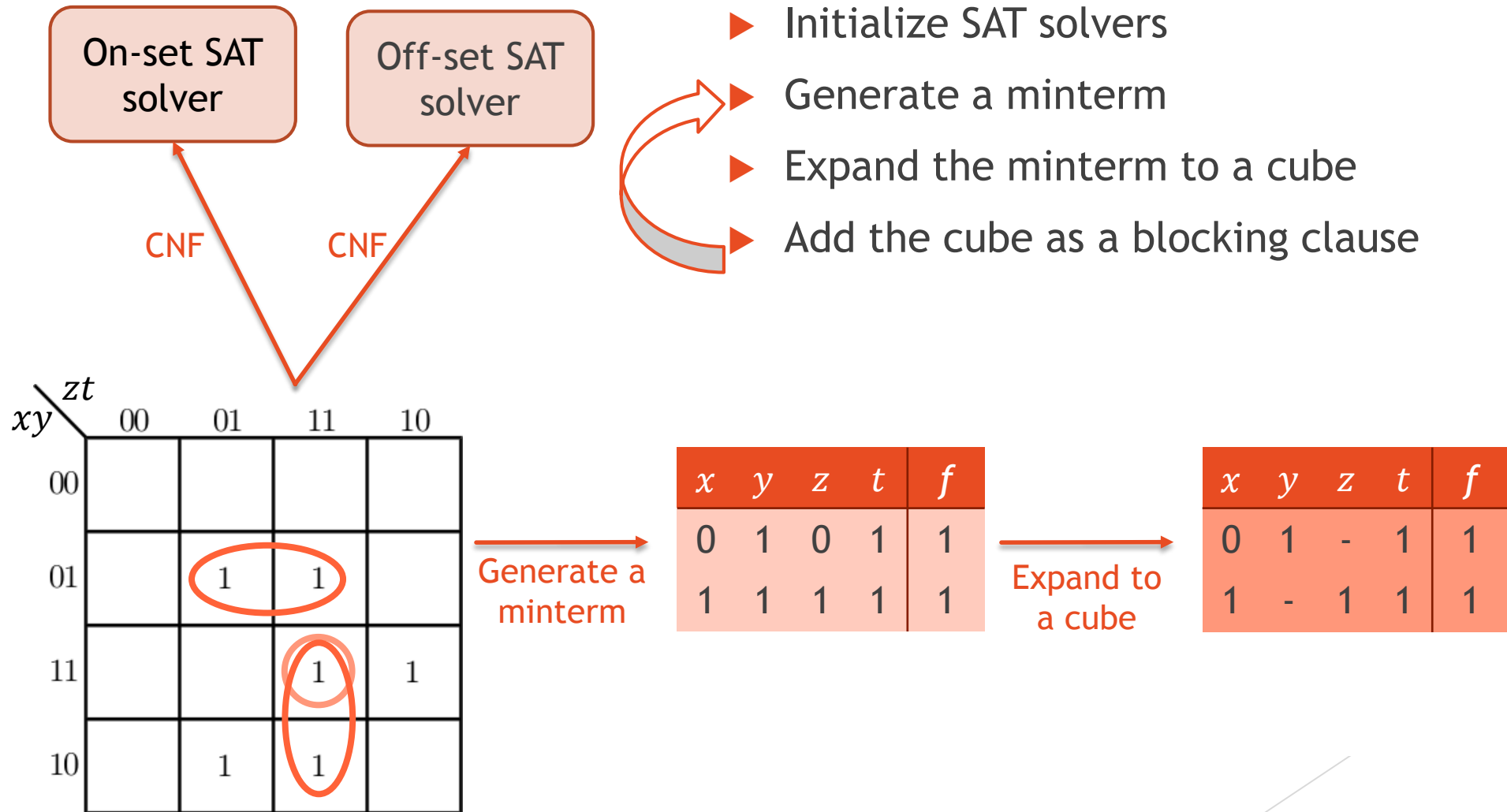
$$\overline{\bar{x}yt} = (x + \bar{y} + \bar{t})$$

- ▶ Initialize SAT solvers
- ▶ Generate a minterm
- ▶ Expand the minterm to a cube
- ▶ Add the cube as a blocking clause

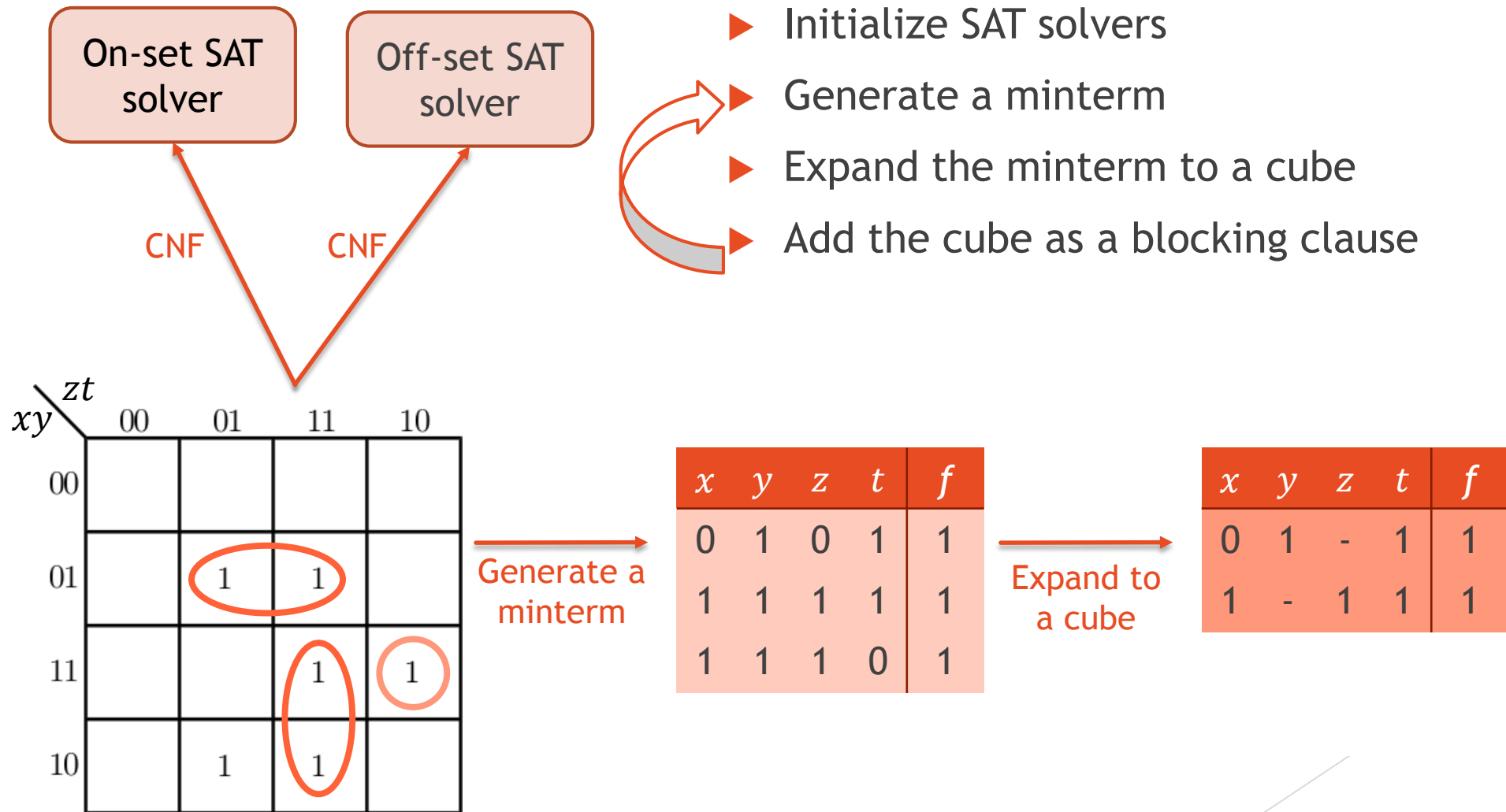
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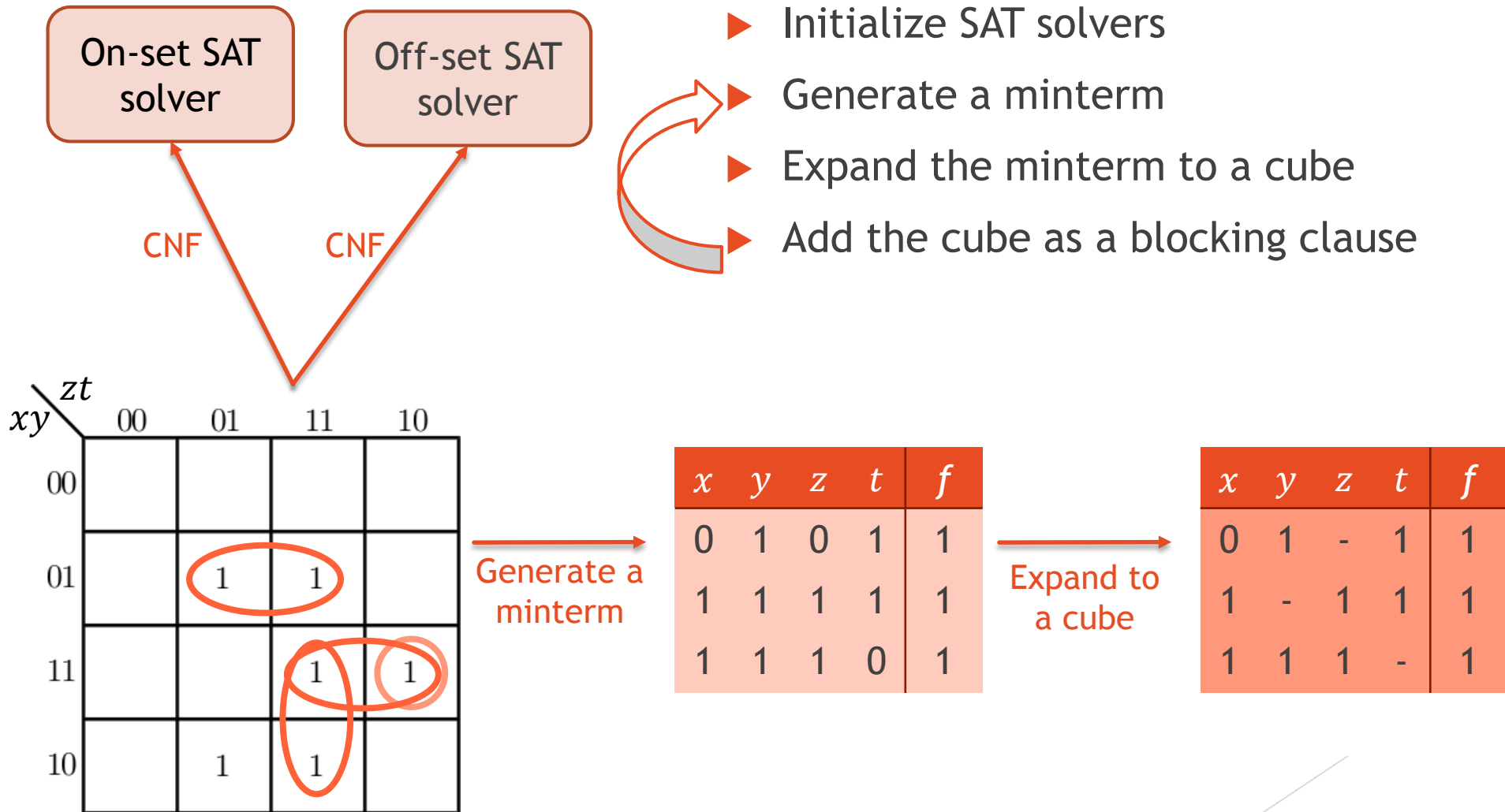
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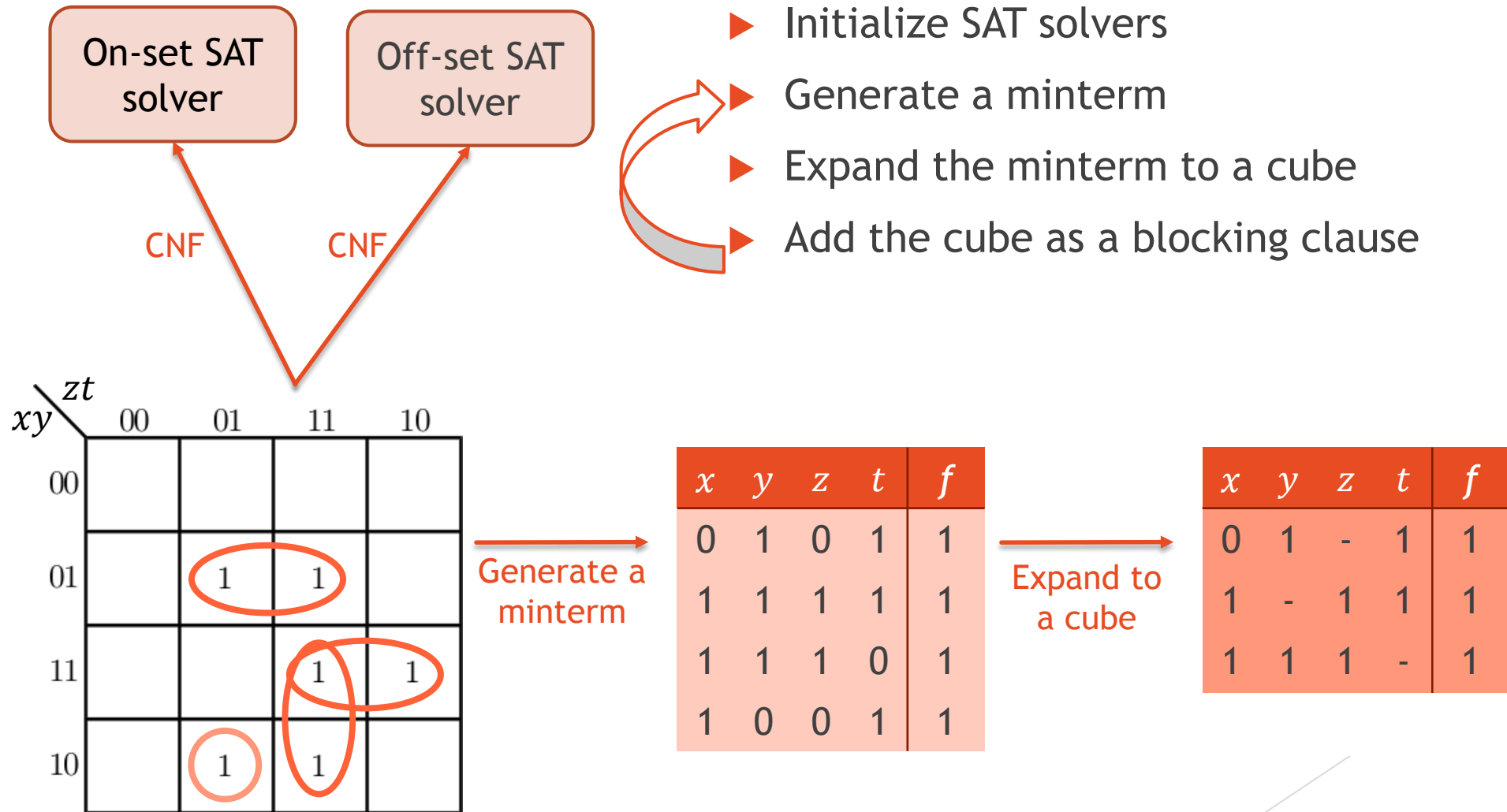
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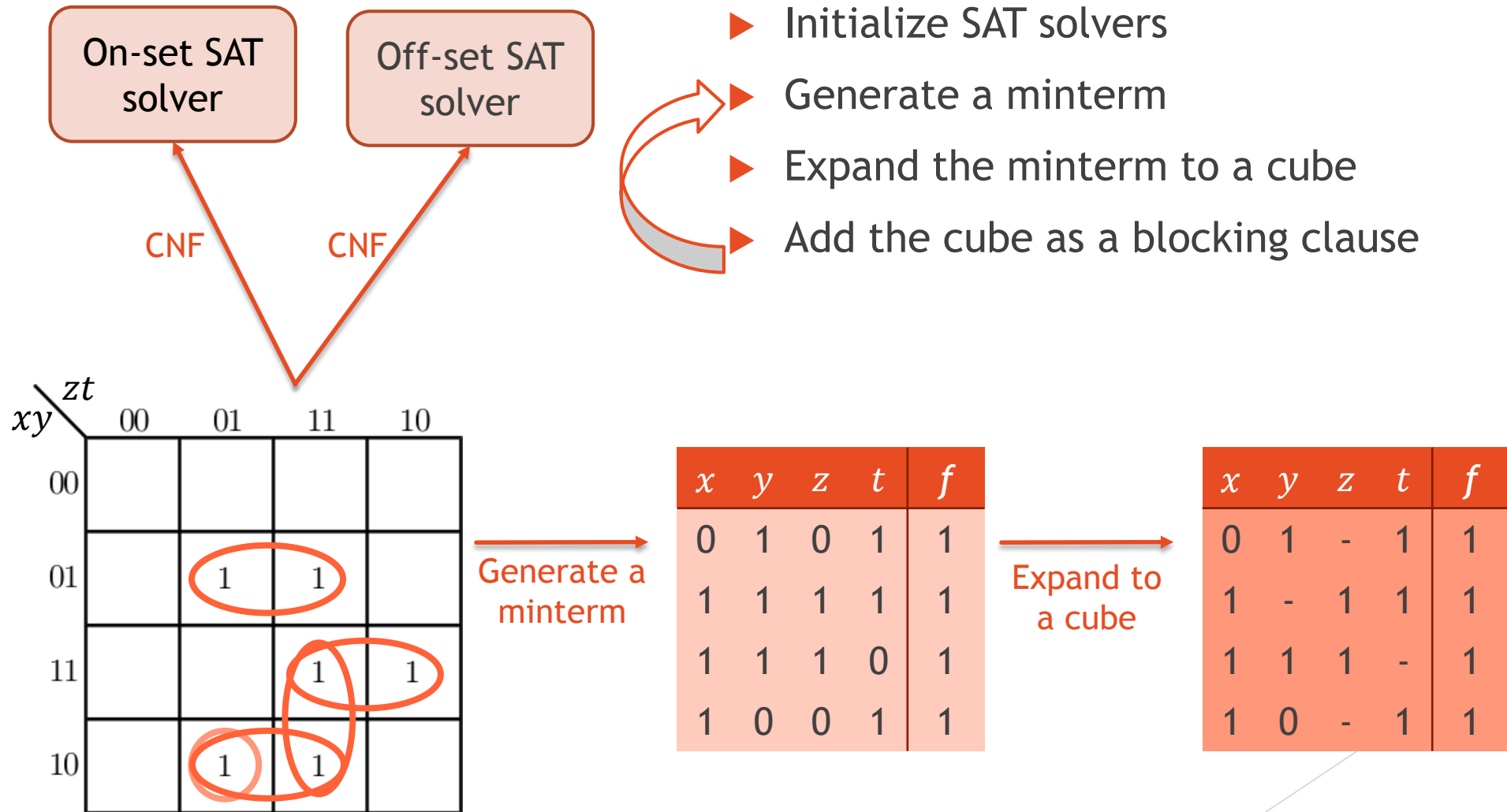
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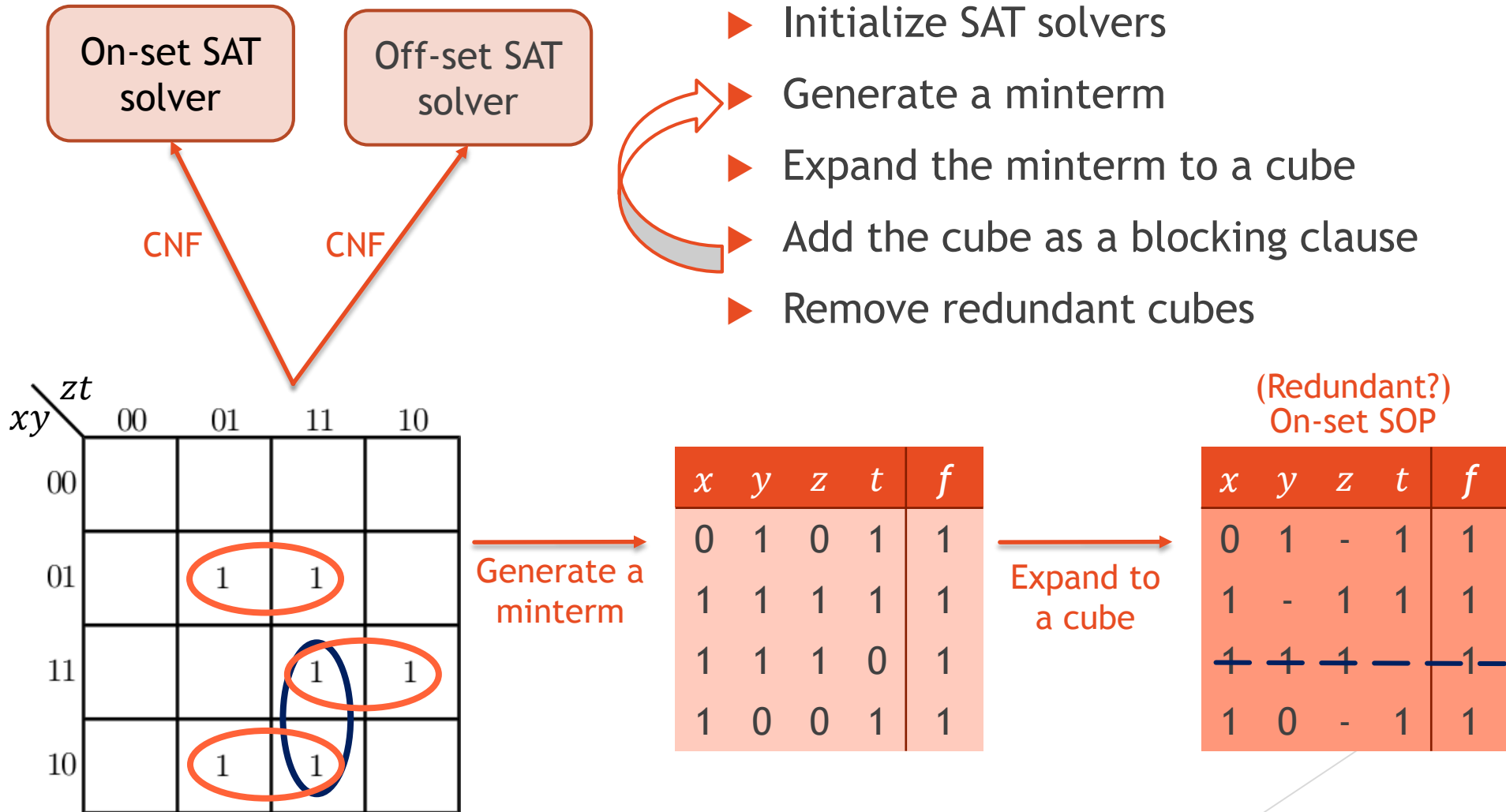
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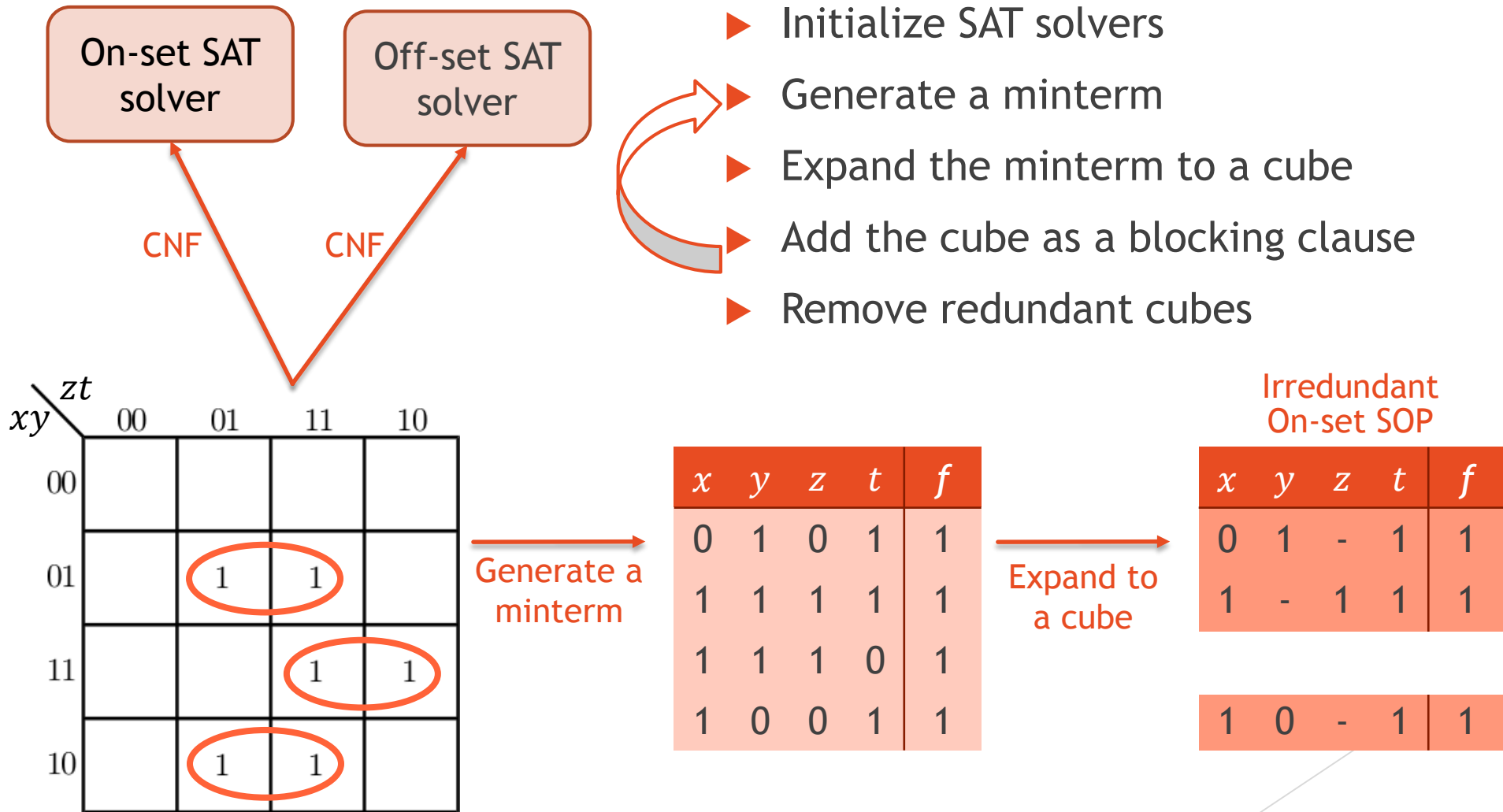


# Progressive SAT-based SOP Generation: Global view





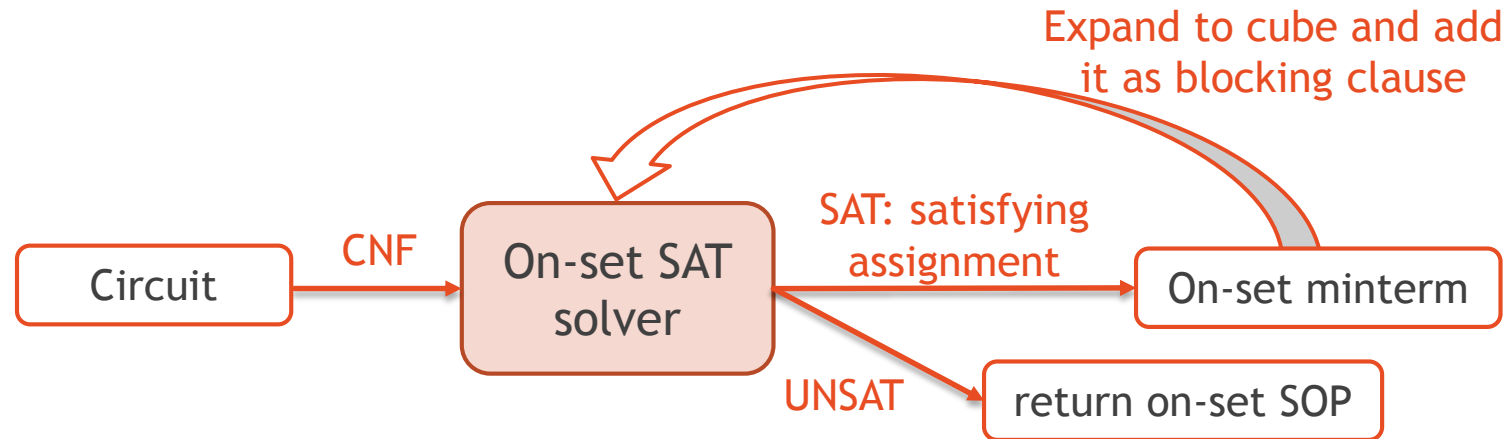
# Progressive SAT-based SOP Generation: Global view



# Outline

- ▶ SAT-based SOP Generation
  - Generation of Minterms
  - Expansion of Minterms to Cubes
  - Removing Redundant Cubes
- ▶ Methods for Runtime Improvement
- ▶ Experimental Results
- ▶ Conclusion

# Generation of Minterms

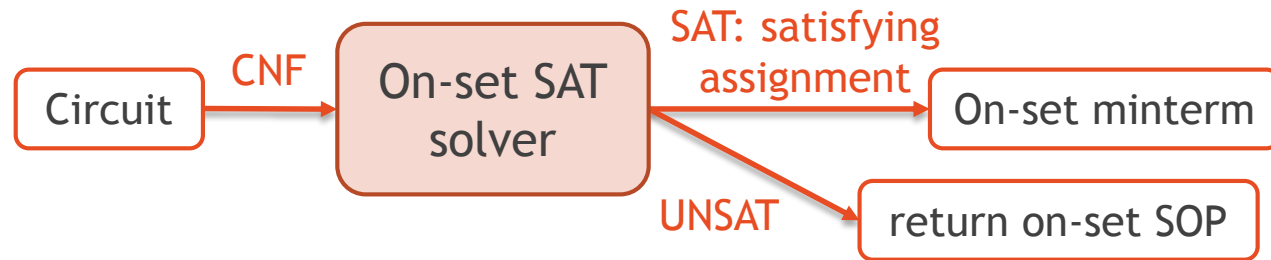


- ▶ *Non-canonical SOPs:* A SAT assignment
- ▶ *Canonical SOPs:* The lexicographically smallest SAT assignment under the variable order
  - Use the LEXSAT algorithm
  - Deterministic algorithm: generate the same minterms in the same order

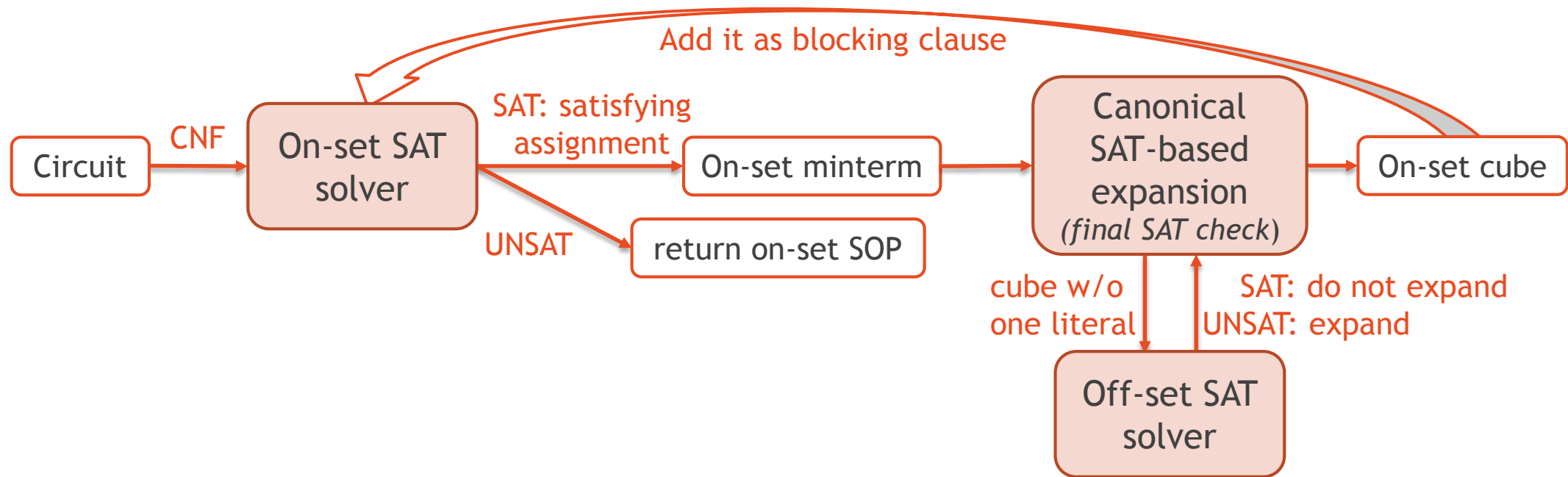
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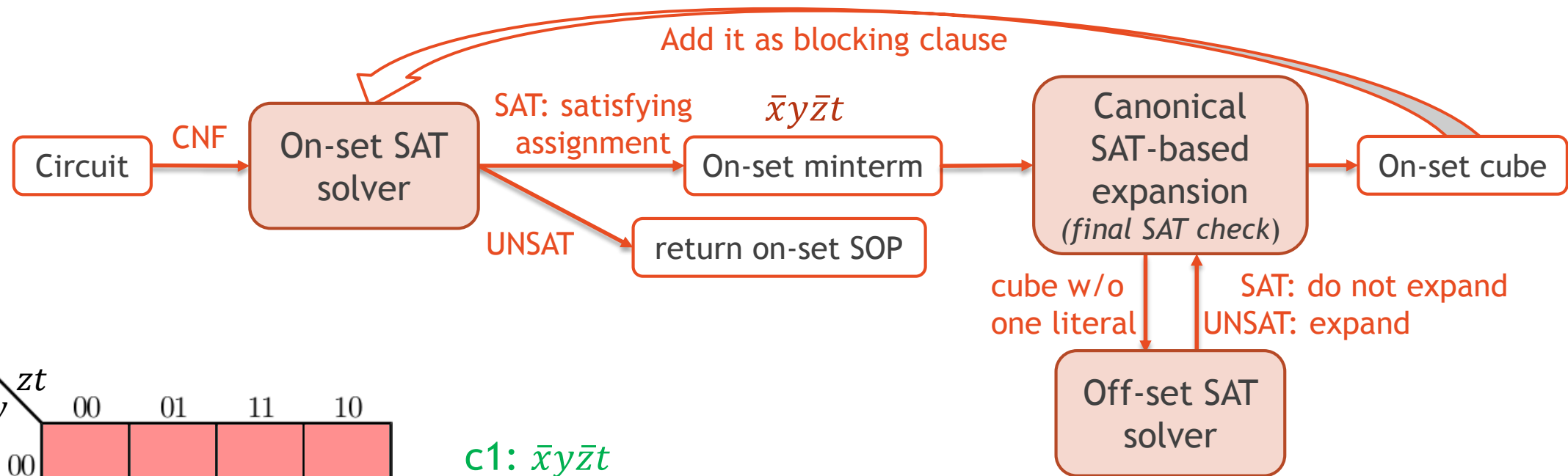
# Expansion of Minterms to Cubes



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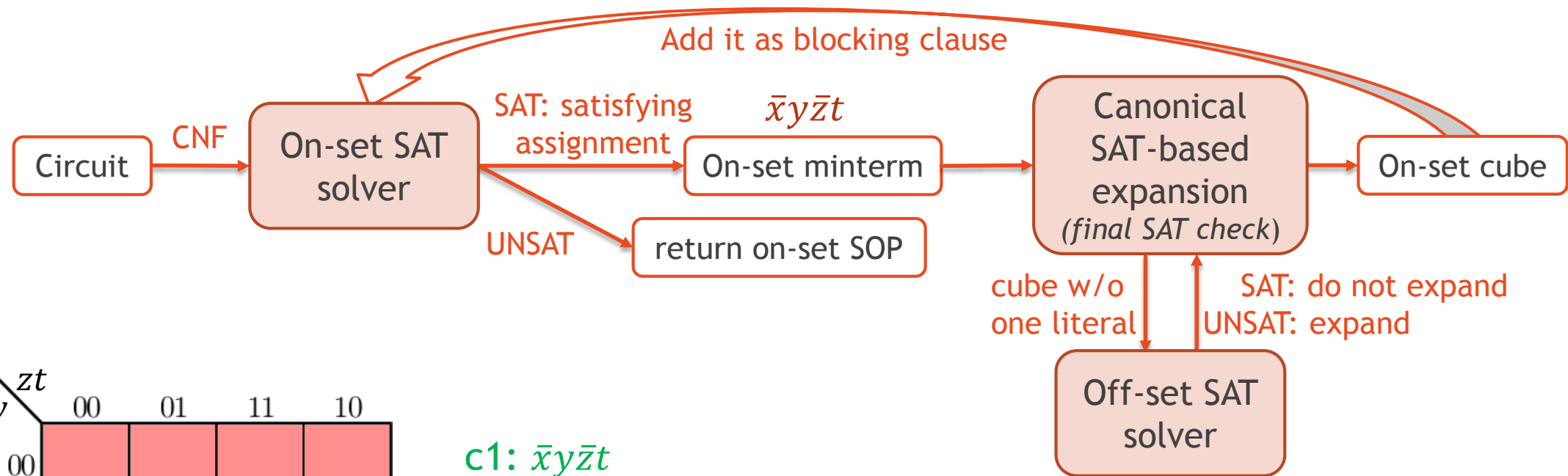
# Expansion of Minterms to Cubes



xy \ zt				
	00	01	11	10
00				
01		1	1	
11			1	1
10		1	1	

c1:  $\bar{x}yzt$

# Expansion of Minterms to Cubes

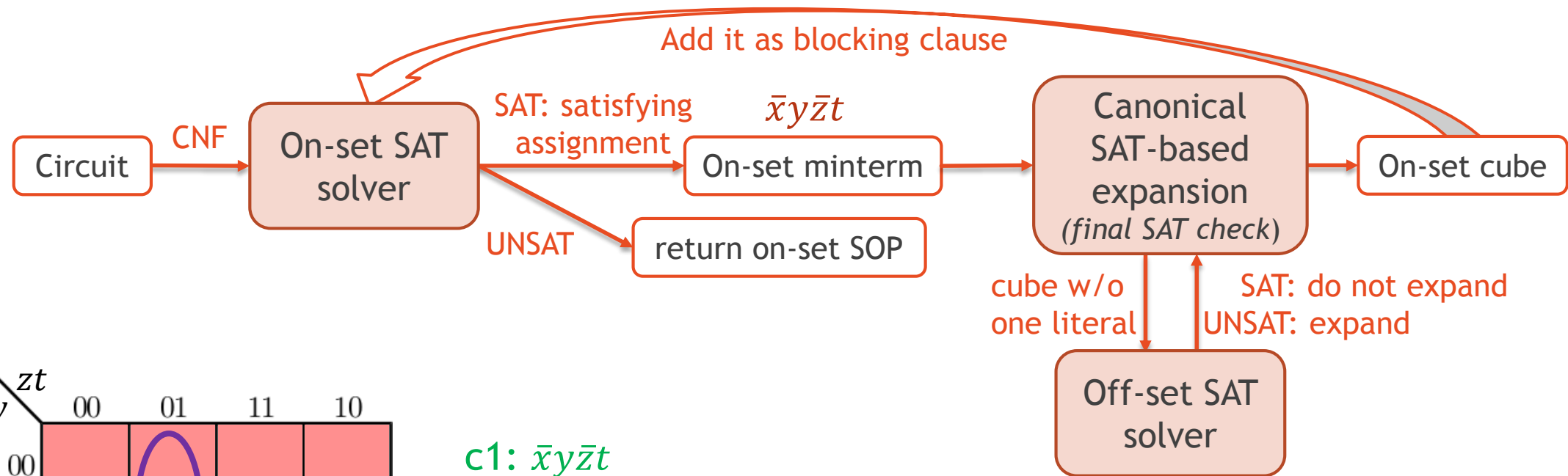


xy \ zt	00	01	11	10
00				
01		1	1	
11			1	1
10		1	1	

c1:  $\bar{x}yzt$   
c2:  $\neg yzt$



# Expansion of Minterms to Cubes



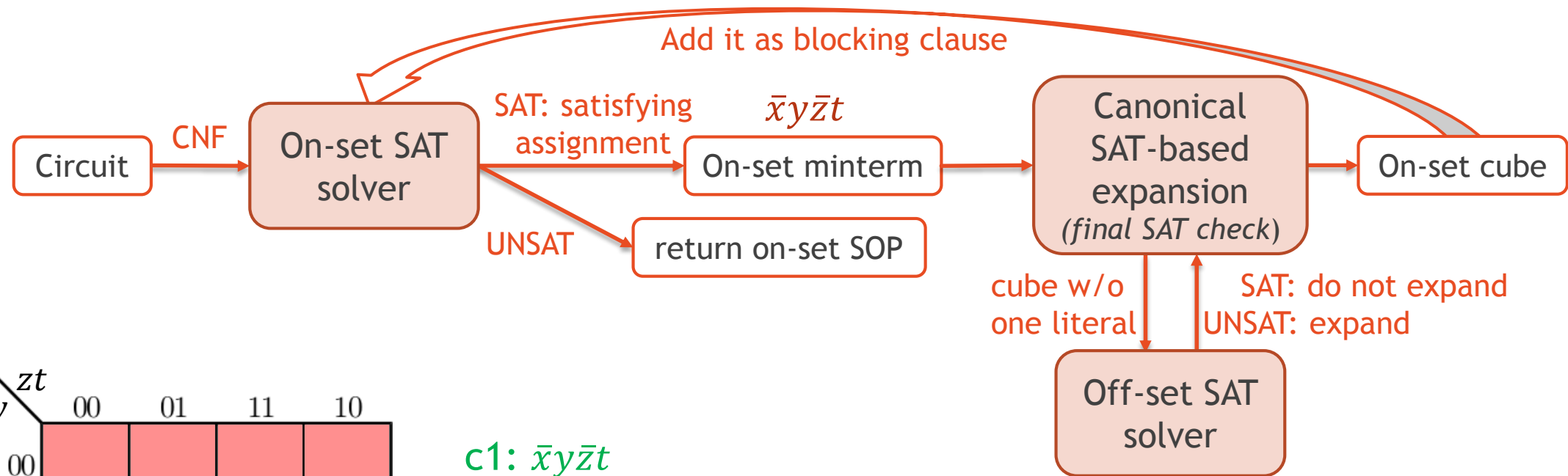
xy \ zt	00	01	11	10
00				
01		1	1	
11			1	1
10		1	1	

c1:  $\bar{x}yzt$

c2:  $\neg yzt$

c3:  $\bar{x}\neg zt$

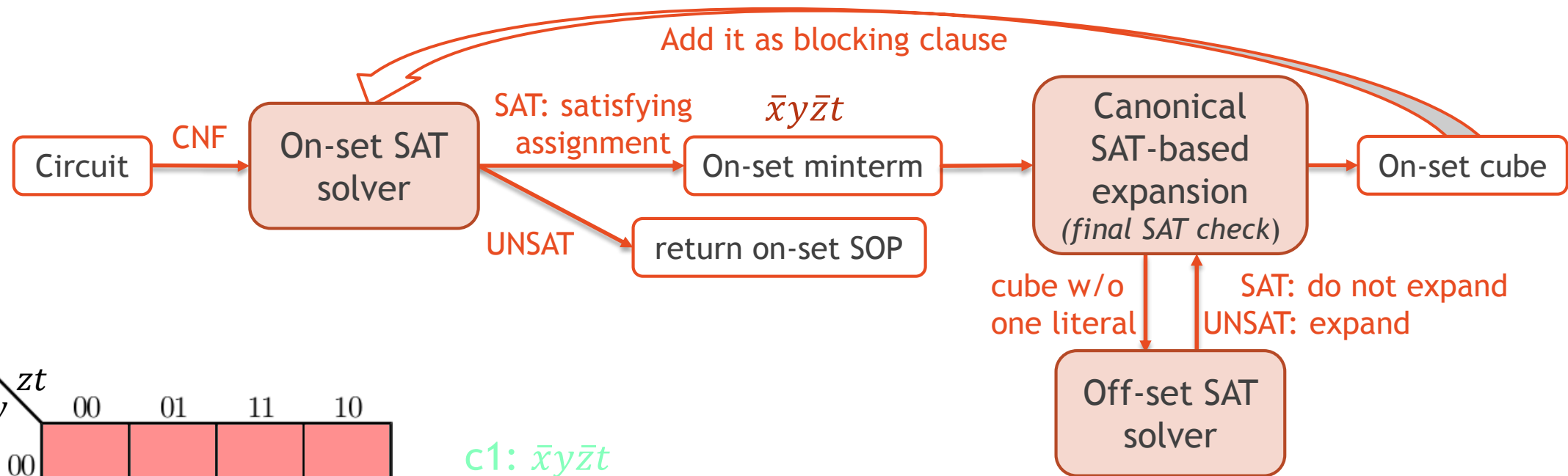
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		zt			
xy		00	01	11	10
00					
01			1	1	
11				1	1
10			1	1	

- c1:  $\bar{x}yzt$
- c2:  $\bar{y}zt$
- c3:  $\bar{x}\bar{z}t$
- c4:  $\bar{x}y\bar{t}$

# Expansion of Minterms to Cubes



		zt			
xy		00	01	11	10
	00				
	01		1	1	
	11			1	1
	10		1	1	

c1:  $\bar{x}yzt$

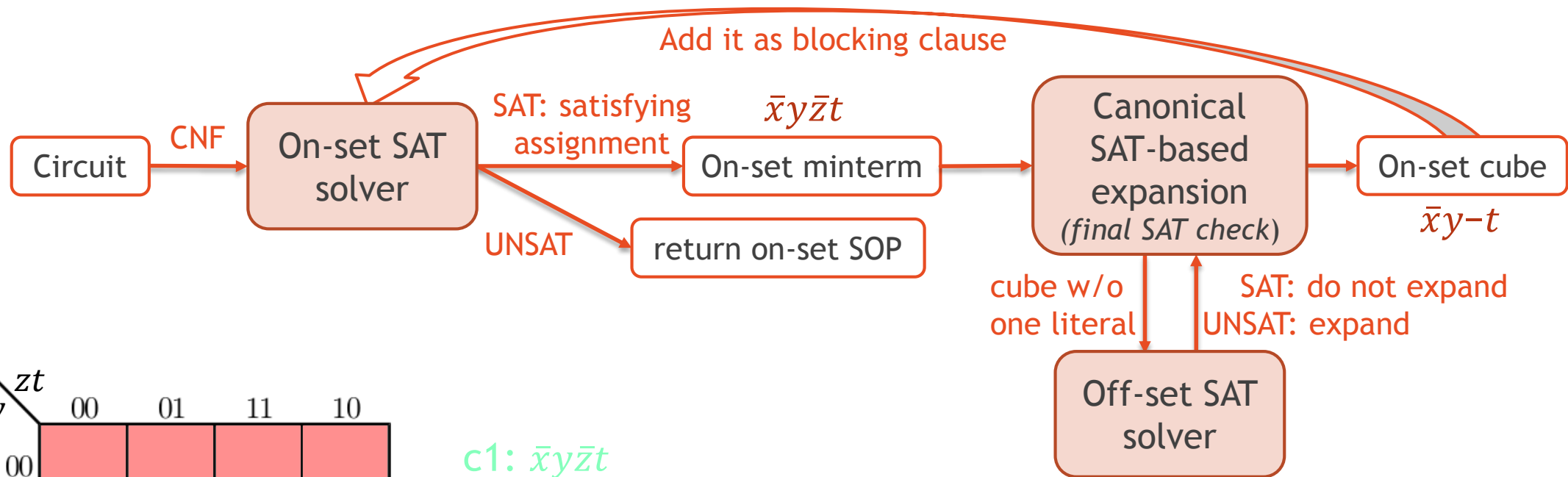
c2:  $\neg yzt$

c3:  $\bar{x}\neg zt$

c4:  $\bar{x}y\neg t$

c5:  $\bar{x}y--$

# Expansion of Minterms to Cubes



		zt			
xy		00	01	11	10
	00				
01			1	1	
11				1	1
10			1	1	

c1:  $\bar{x}yzt$

c2:  $\neg yzt$

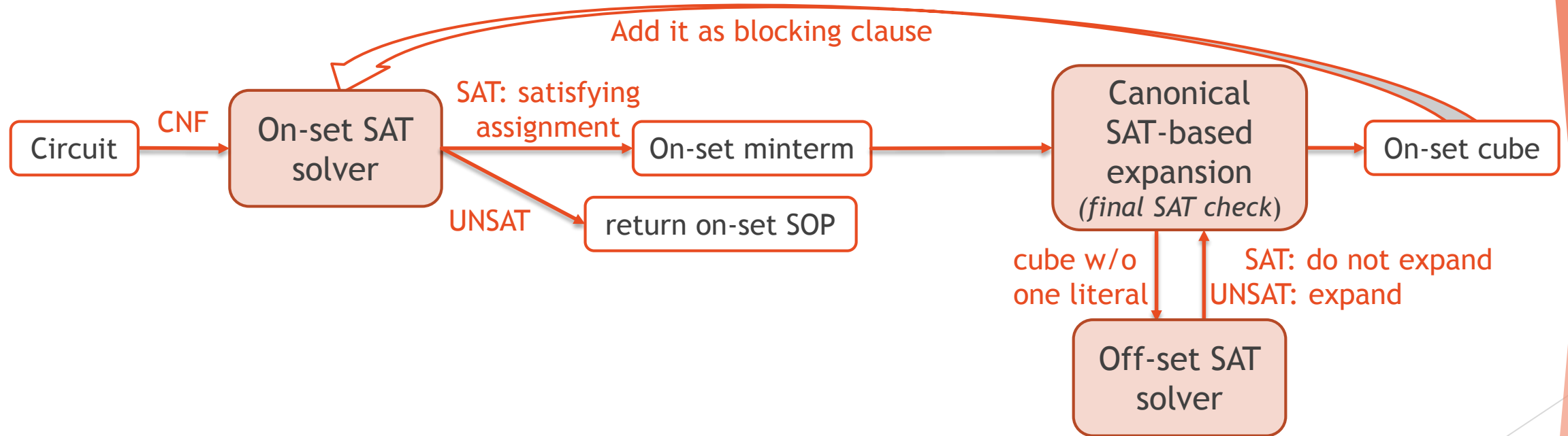
c3:  $\bar{x}\neg zt$

c4:  $\bar{x}y-t$

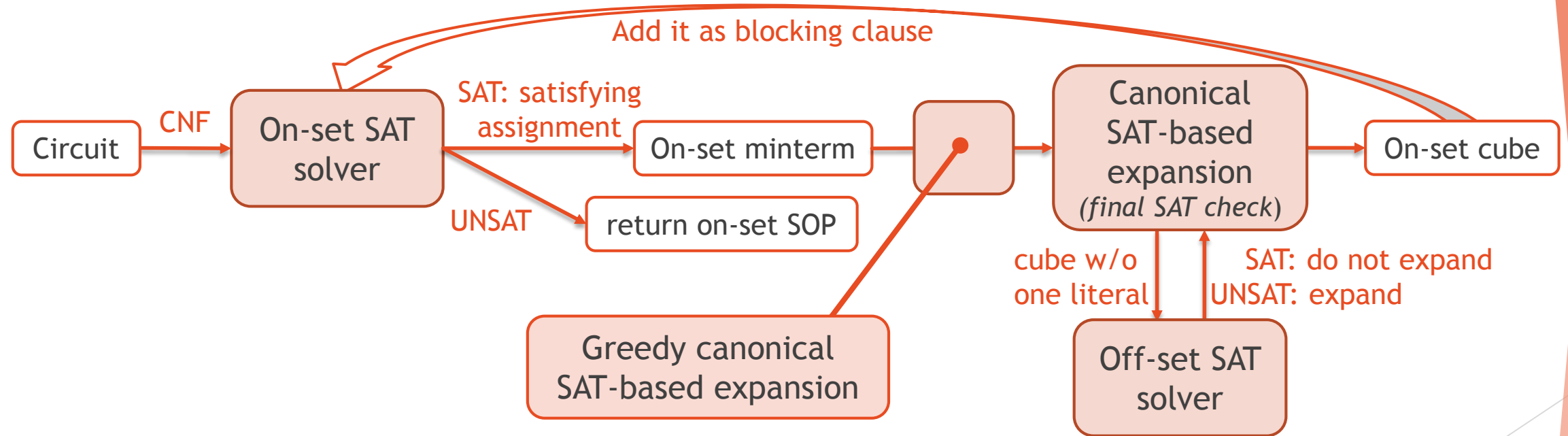
c5:  $\bar{x}y--$

- ▶ Canonical expansion: literals are always removed in the same order
- ▶ Prime cubes: no literal can be removed without changing the value of the function

# Expansion of Minterms to Cubes

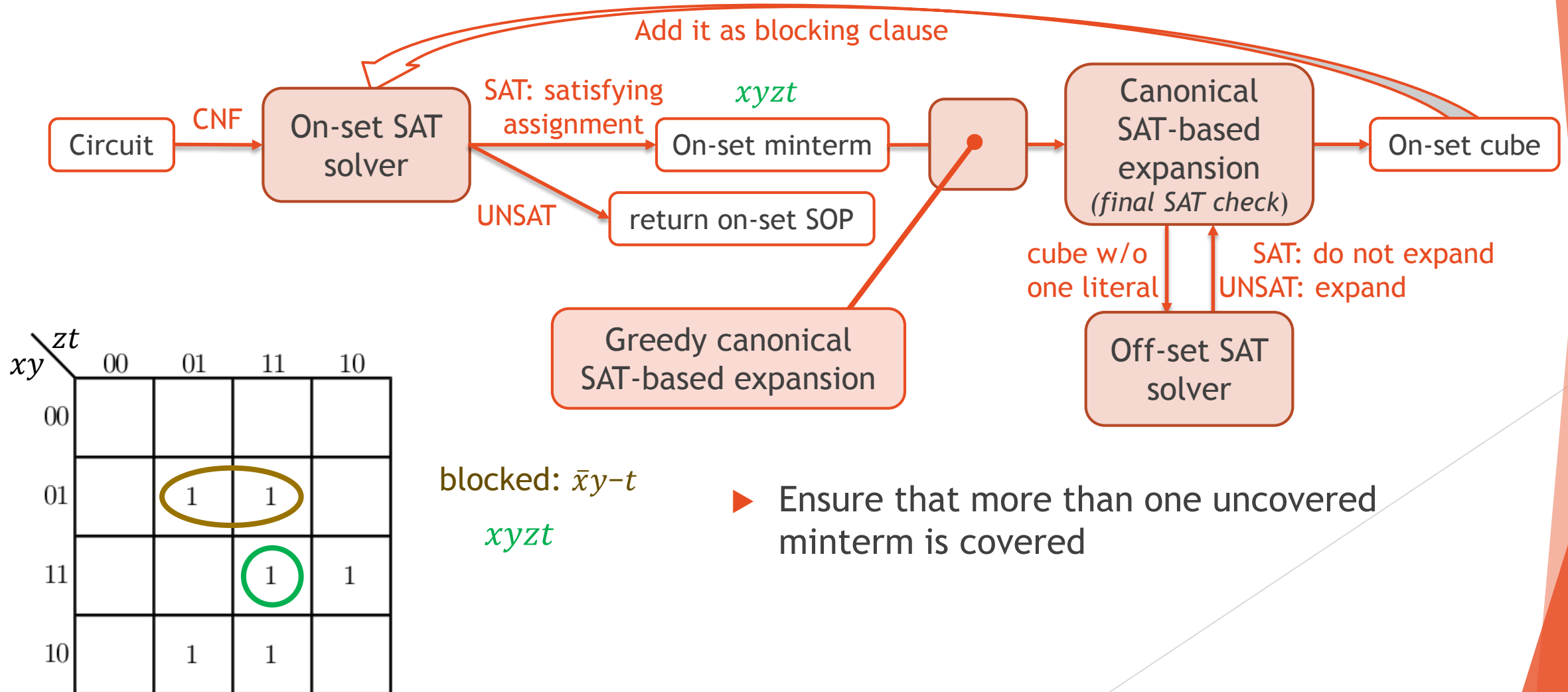


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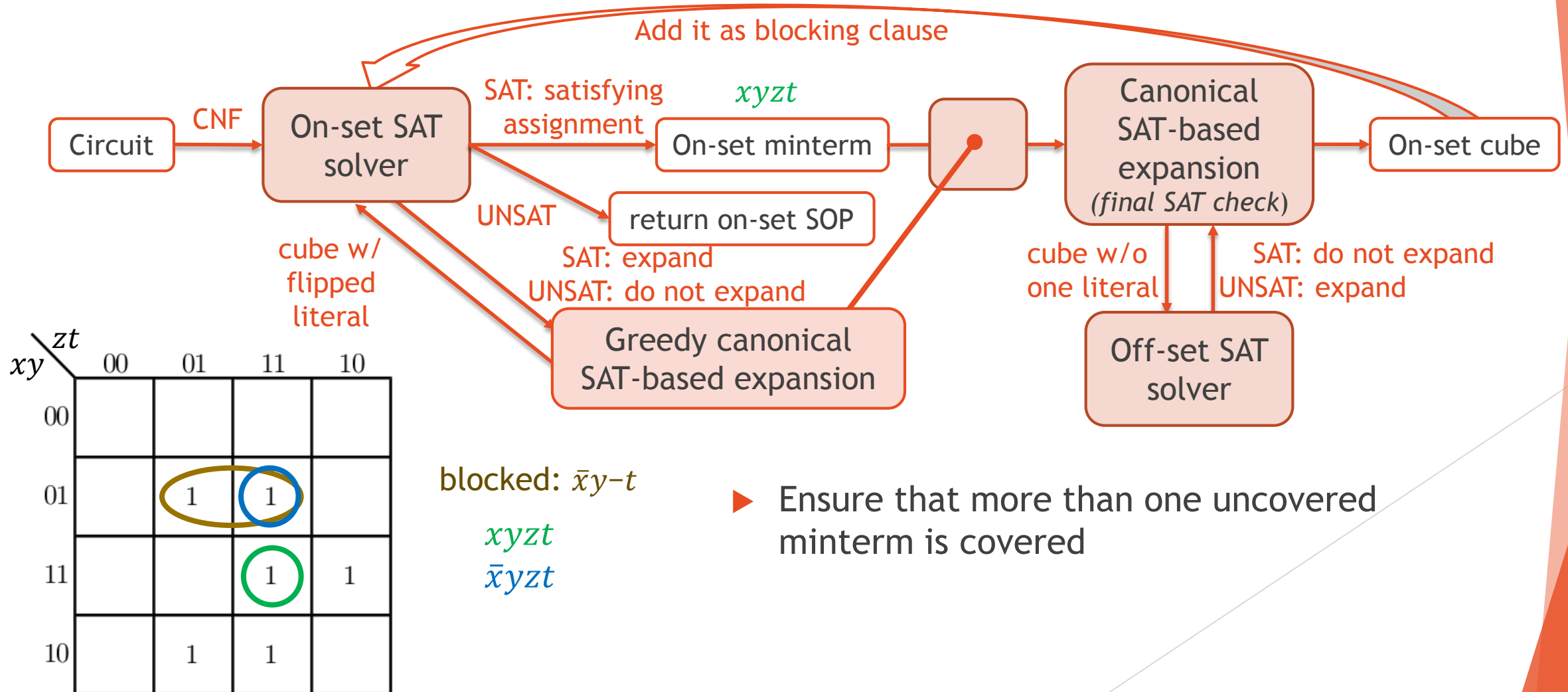


- Ensure that more than one uncovered minterm is covered

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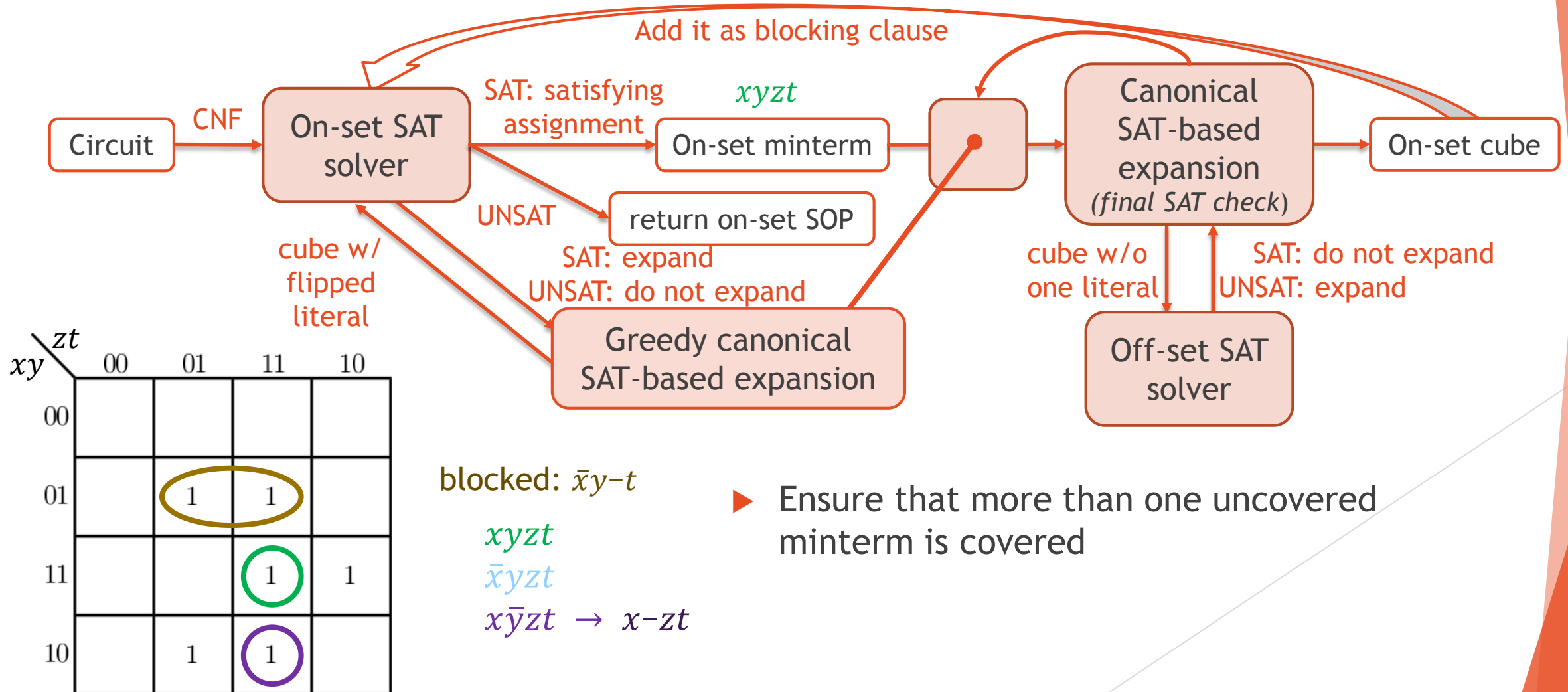


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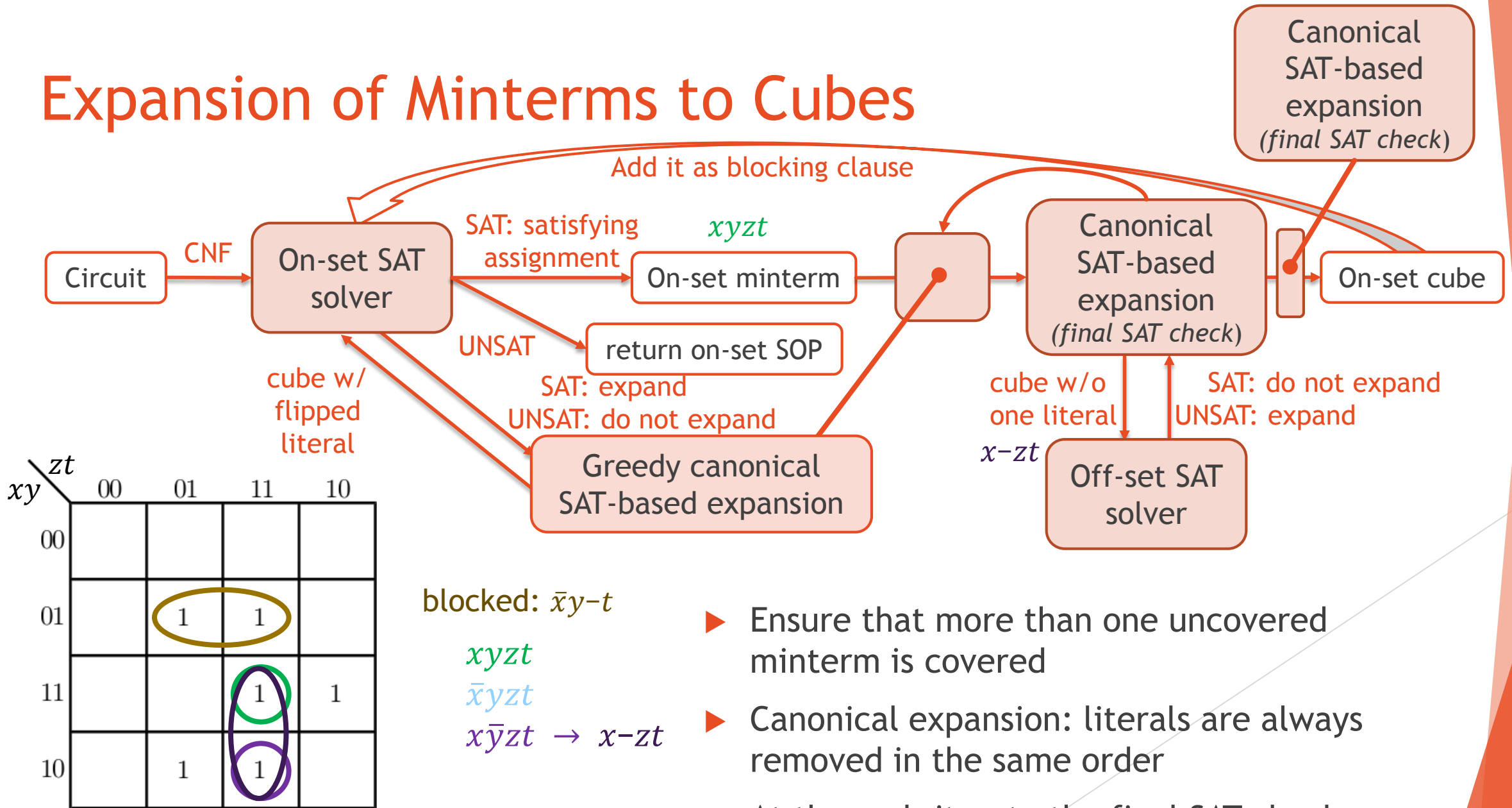


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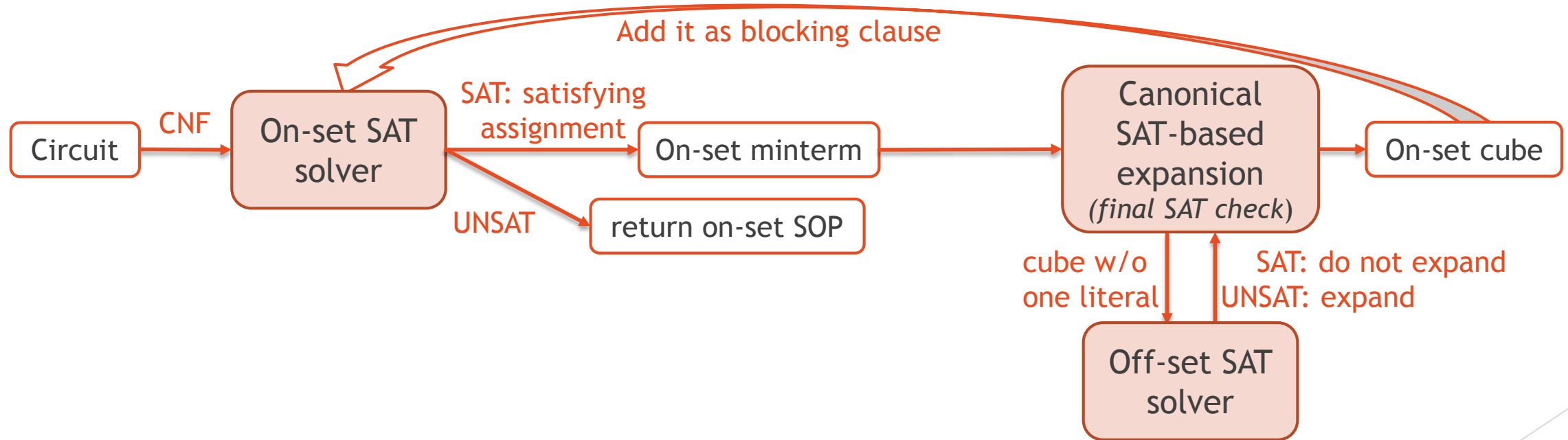
- Ensure that more than one uncovered minterm is covered

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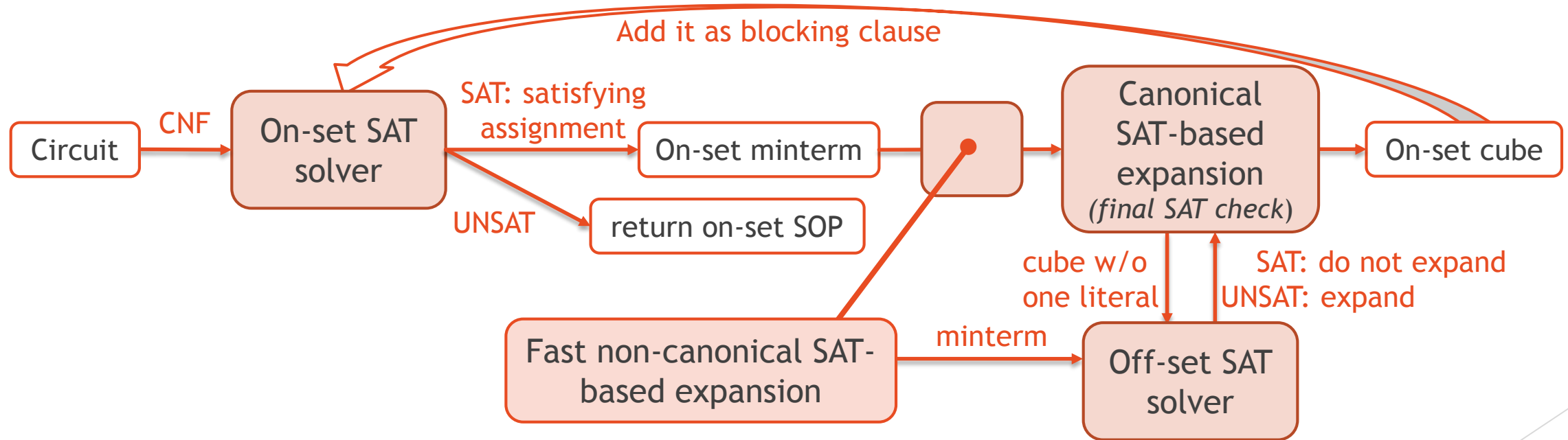


- ▶ Ensure that more than one uncovered minterm is covered
- ▶ Canonical expansion: literals are always removed in the same order
- ▶ At the end: iterate the final SAT check on all literals

# Expansion of Minterms to Cubes

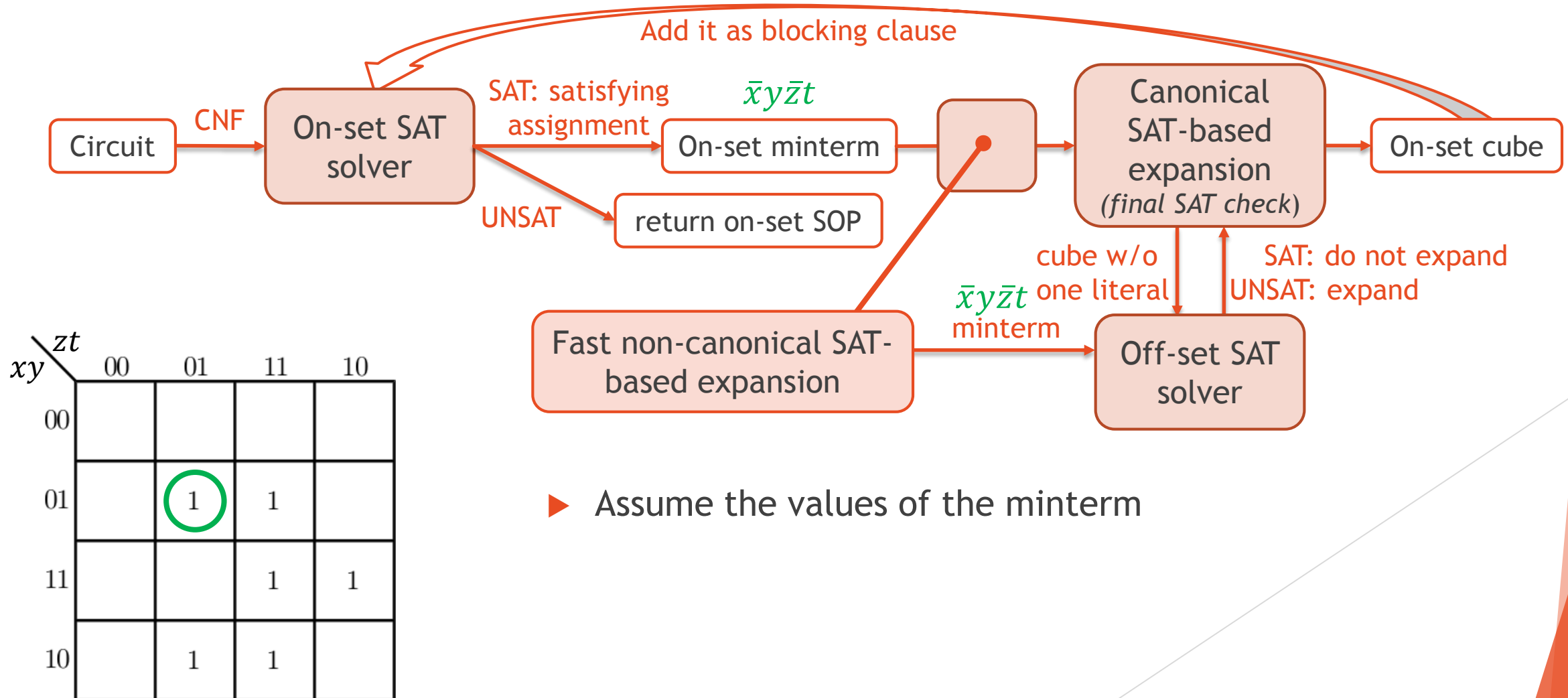


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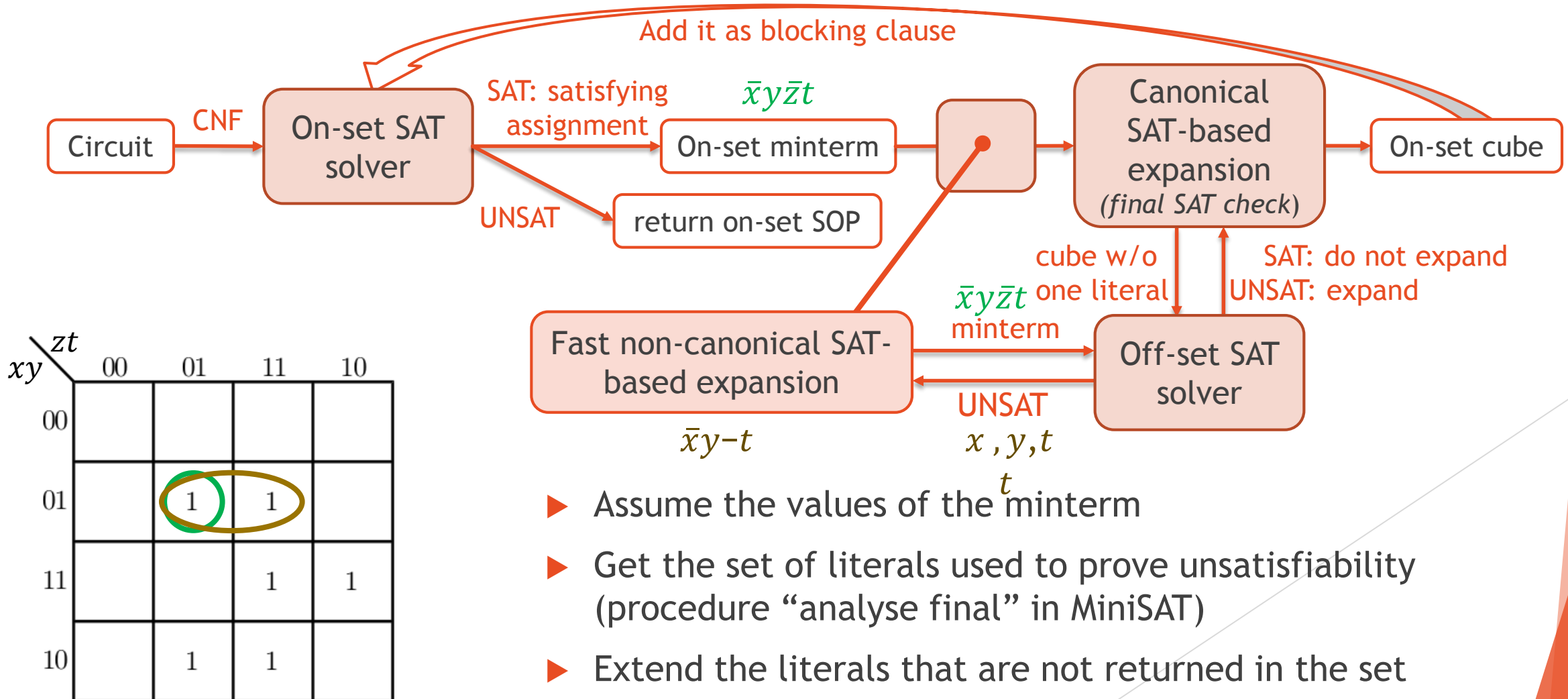
- Assume the values of the minterm

# Expansion of Minterms to Cubes



► Assume the values of the minterm

# Expansion of Minterms to Cubes



- ▶ Assume the values of the minterm
- ▶ Get the set of literals used to prove unsatisfiability (procedure “analyse final” in MiniSAT)
- ▶ Extend the literals that are not returned in the set

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# Removing Redundant Cubes

zt \ xy	00	01	11	10
00				
01		1	1	
11			1	1
10		1	1	

$$\begin{aligned} c1 &= \bar{x}yt \\ c2 &= xzt \\ c3 &= xyz \\ c4 &= x\bar{y}t \end{aligned}$$

$$c1 = 1$$

$$\begin{aligned} x &= 0 \\ y &= 1 \\ t &= 1 \end{aligned}$$

$$\begin{aligned} c2 &= 0 \\ c3 &= 0 \\ c4 &= 0 \end{aligned}$$

$$\begin{aligned} c2 &= 0 \\ c3 &= 0 \\ c4 &= 0 \end{aligned}$$

SAT solver  
with cubes

SAT:  $\bar{x}yzt$   
 $c1$  is irredundant

- ▶ Initialize a SAT solver with all cubes
- ▶ Find an assignment for which  $c_i = 1$  and the other irredundant cubes evaluate to 0
  - SAT:  $c_i$  is irredundant
  - UNSAT:  $c_i$  is redundant, remove it from the SOP
- ▶ Canonical: cubes are removed always in the same order



# Removing Redundant Cubes

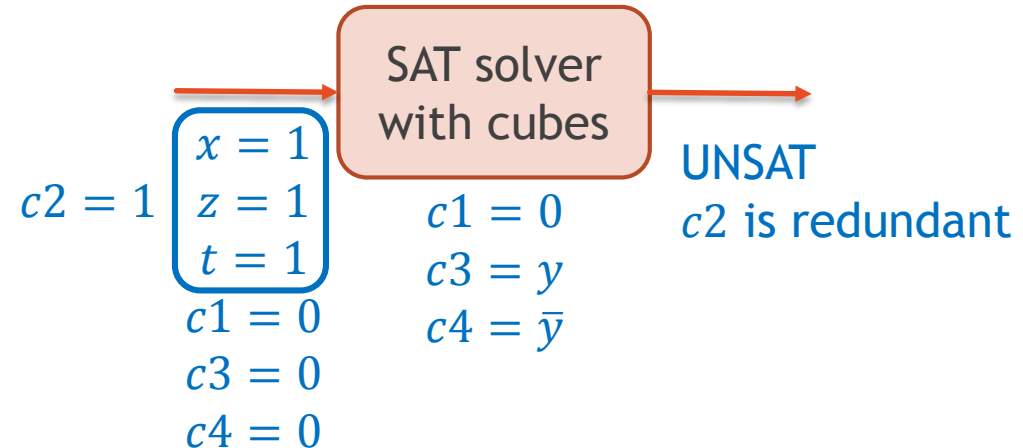
$xy \backslash zt$	00	01	11	10
00				
01		1	1	
11			1	1
10		1	1	

$$c1 = \bar{x}yt$$

$$c2 = xzt$$

$$c3 = xyz$$

$$c4 = x\bar{y}t$$



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- ▶ Find an assignment for which  $c_i = 1$  and the other irredundant cubes evaluate to 0
  - SAT:  $c_i$  is irredundant
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  - Expansion of Minterms to Cubes
  - Removing Redundant Cubes
- ▶ Methods for Runtime Improvement
- ▶ Experimental Results
- ▶ Conclusion

# Methods for Runtime Improvement

- ▶ Generate simultaneously on-set and off-set SOPs

On-set SOP

$x_1$	$x_2$	$x_3$	$f$
-	1	1	1
1	1	-	1

$$f = x_2x_3 + x_1x_2$$

Off-set SOP

$x_1$	$x_2$	$x_3$	$f$
0	0	-	0
0	-	0	0
1	0	-	0

$$\bar{f} = \bar{x}_1\bar{x}_2 + \bar{x}_1\bar{x}_3 + x_1\bar{x}_2$$

Avoid generating first the larger SOP

Future work: generate each set in parallel

- ▶ Prioritize outputs with large SOP

- Sort outputs by size of their input support
- Generate SOPs in decreasing order

Early termination when the large SOP exceed resource limits

Future work: generate SOP of each output in parallel

S<sub>3</sub>: 9 inputs

S<sub>2</sub>: 7 inputs

S<sub>1</sub>: 5 inputs

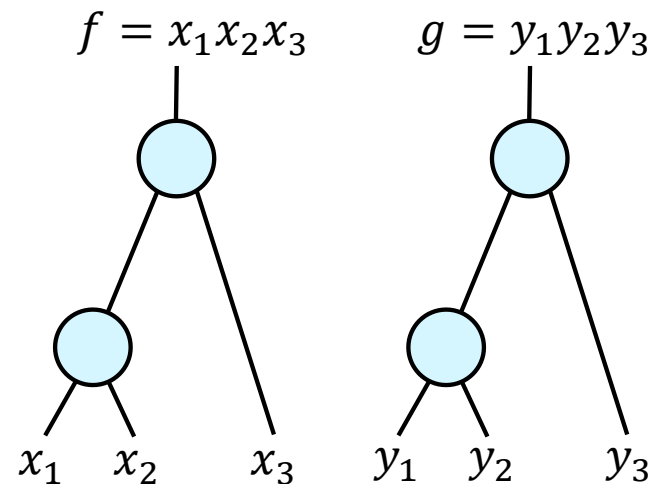
S<sub>0</sub>: 3 inputs

# Methods for Runtime Improvement

## ► Detect isomorphic outputs

Benefit from structure sharing

- Isomorphic functions: implement an identical function using different inputs



### Isomorphic class 1

Func:  $f, g$

Repr:  $f = x_1x_2x_3$

- Generate SOP only for one output per class (the representative)
- Duplicate the generated SOP for the other outputs from the class
- Share one CNF among all outputs

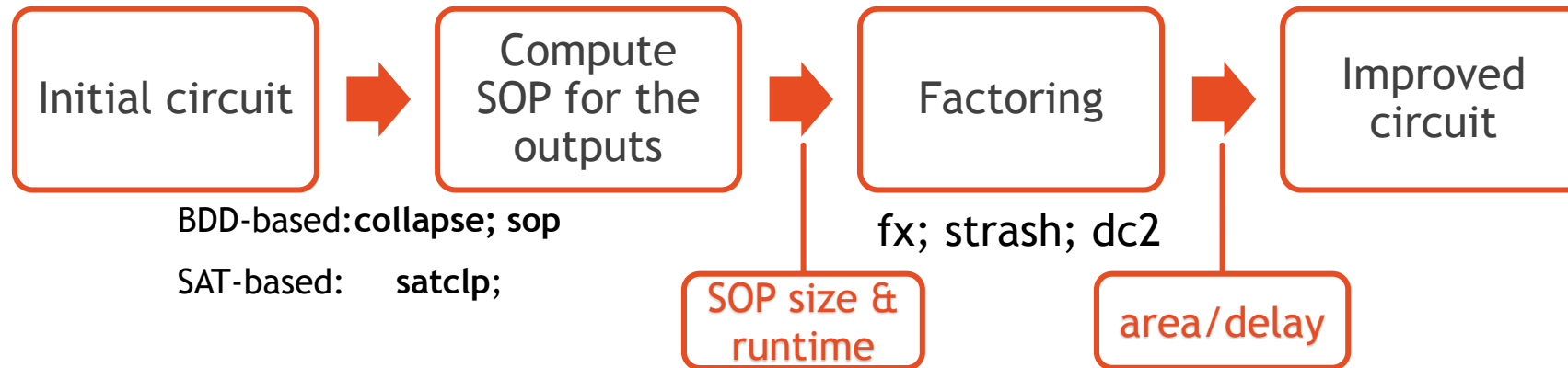
Benefit from logic sharing

# Outline

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# Experimental Setup

- ▶ Implemented the SAT-based SOP generation in ABC
- ▶ Compared flows for SOP generation



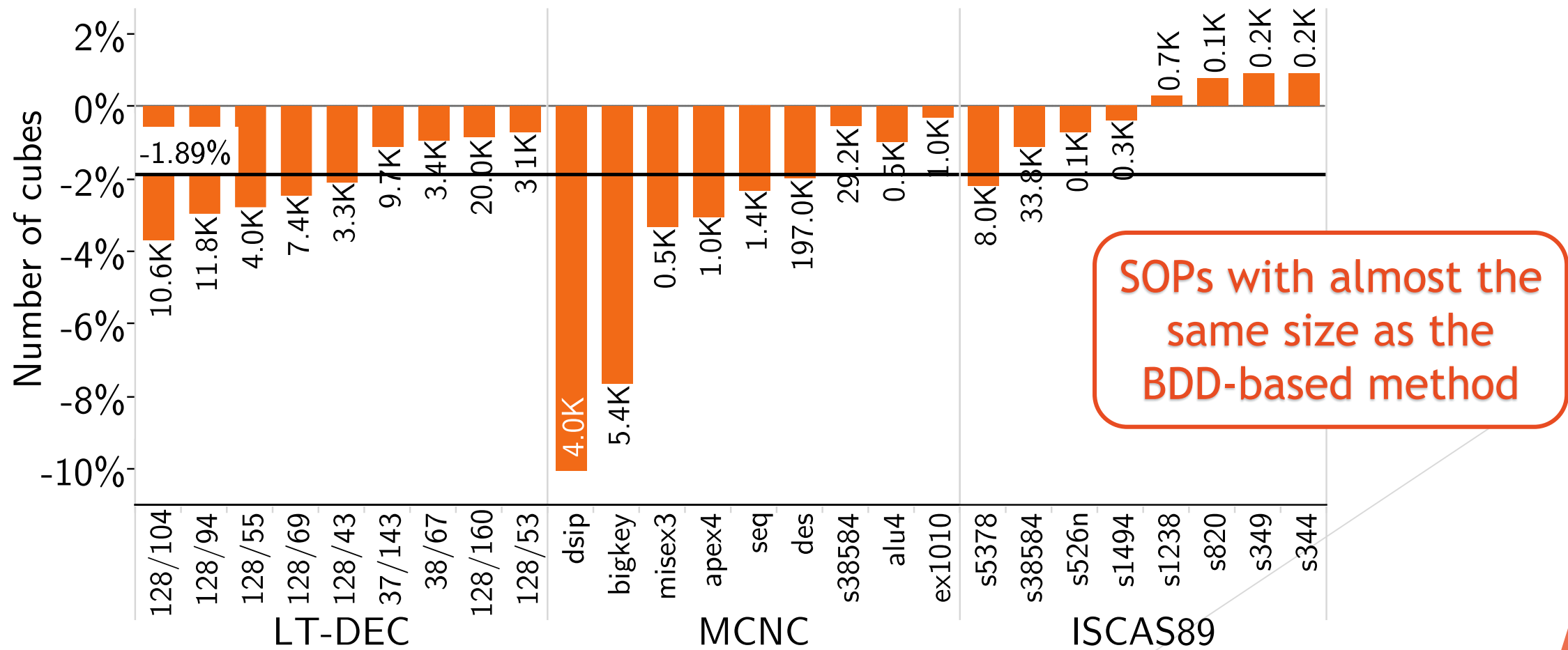
- ▶ For the BDD-based SOPs, generate 5 SOPs using different variable orders
- ▶ For the SAT-based SOPs, generate 12 SOPs using different options
  - Non-canonical or canonical SOPs
  - Order variables by fanout number
  - Reverse variable order
  - Shared CNF by primary outputs

# Experimental Setup

## ► Benchmarks

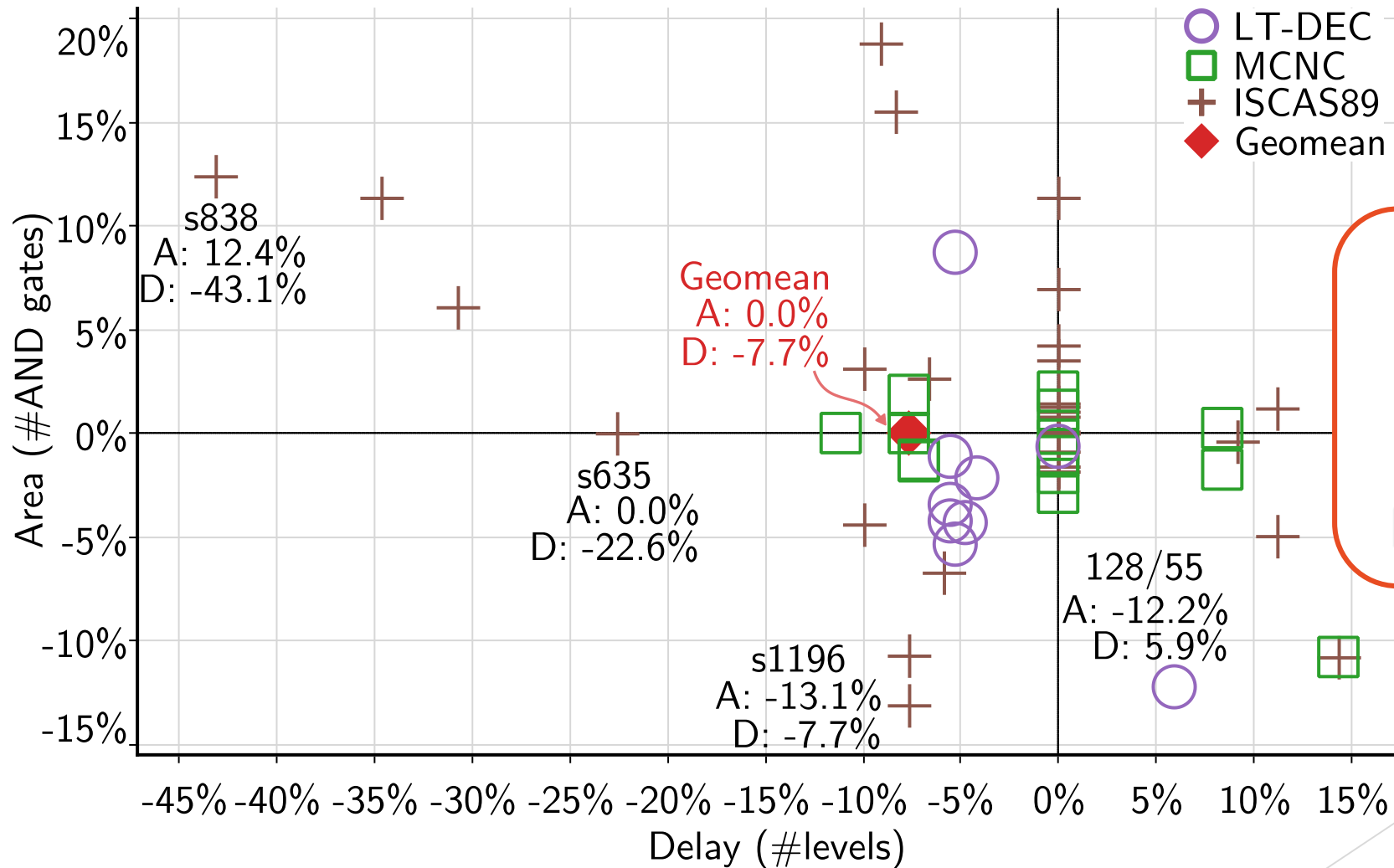
- Set of 20 large MCNC benchmarks
- Set of 32 ISCAS'89 benchmarks
- Set of 9 logic tables from instruction decoder unit (LT-DEC)

# Comparison of SOP Size





# Comparison of Area-Delay Product



Pareto optimal  
multi-level  
implementations  
compared to the  
BDD-based method

# Comparison of Runtime

## ▶ Public benchmarks

- On average, 8x slower than the BDD-based method
- Faster for circuits with many isomorphic outputs
  - S35932 - 10 isomorphic classes for 2048 combinatorial outputs ->  
SAT-based: 0.06 sec  
BDD-based: 1.83 sec

## ▶ Industrial benchmarks

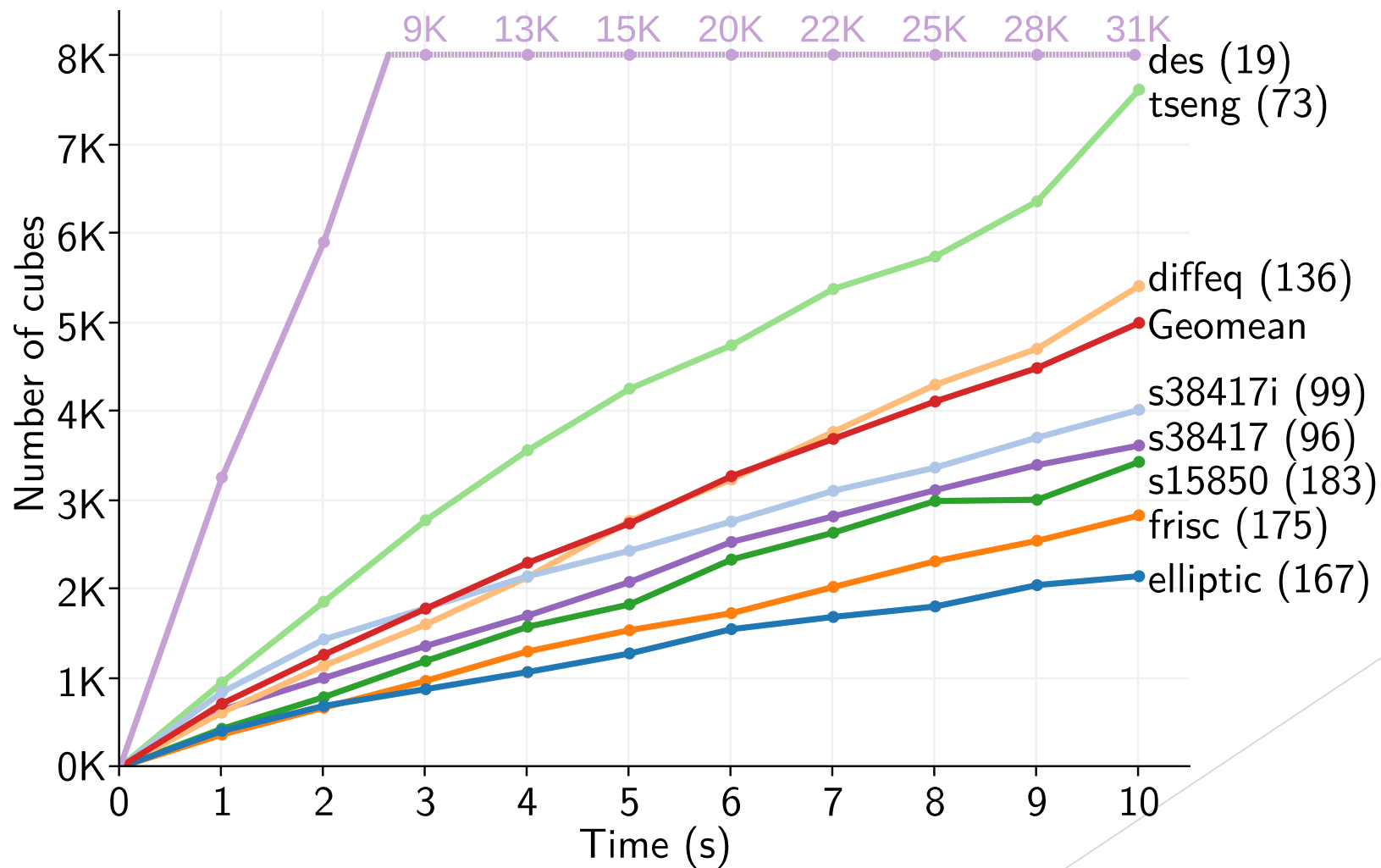
- Faster due to isomorphism: generated SOPs for only ~30% of the outputs
- Hundreds of PIs/POs -> constructing the global BDD in one manager is problematic

# Comparison of Runtime

## ▶ Runtime profile (for the LT-DEC benchmarks)

- Generating minterms 8%
- Expanding minterms to cubes 84%
- Removing redundant cubes 2%
- Other operations 6%

# Partial SOPs



# Outline

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# Conclusion

- ▶ We presented a complete SAT-based algorithm for **progressive generation of irredundant (non-)canonical SOPs**
- ▶ Quality of results
  - Decreased SOP size up to 10% compared to the BDD-based method
  - Area-delay product post factoring can be improved up to 36%
- ▶ Runtime: 8x slower on average, but faster for isomorphic circuits
- ▶ Future work
  - Improving the runtime: using parallelism, faster cube expansion
  - Dedicated SAT-based multi-output SOP computation