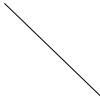


Simple Disjunctive Decomp. On BDDs

- $f = w'x'z' + wx'z + w'yz + wyz'$
- $(w'z')x' + (wz)x' + (w'z)y + (wz')y$
- $f = F(\emptyset(w, z), x, y)$
- $\emptyset = \text{XNOR}(w, z);$
- $\emptyset' = \text{XOR}(w, z)$
- Orthonormal Expansion of f w.r.t. $\emptyset = \emptyset x' + \emptyset' y$
- Decomposition through partition matrix
- Test for existence of a decomp:
- Reduce partition matrix, get column multiplicity
- If col. Mult. ≤ 2 , you WILL FIND A DECOMP!

Decomp. May or maynot exist



0	1	0	1
0	0	1	0
0	0	1	0
0	1	0	1

Significance of Column Multiplicity

0	0
1	0
0	1
1	0

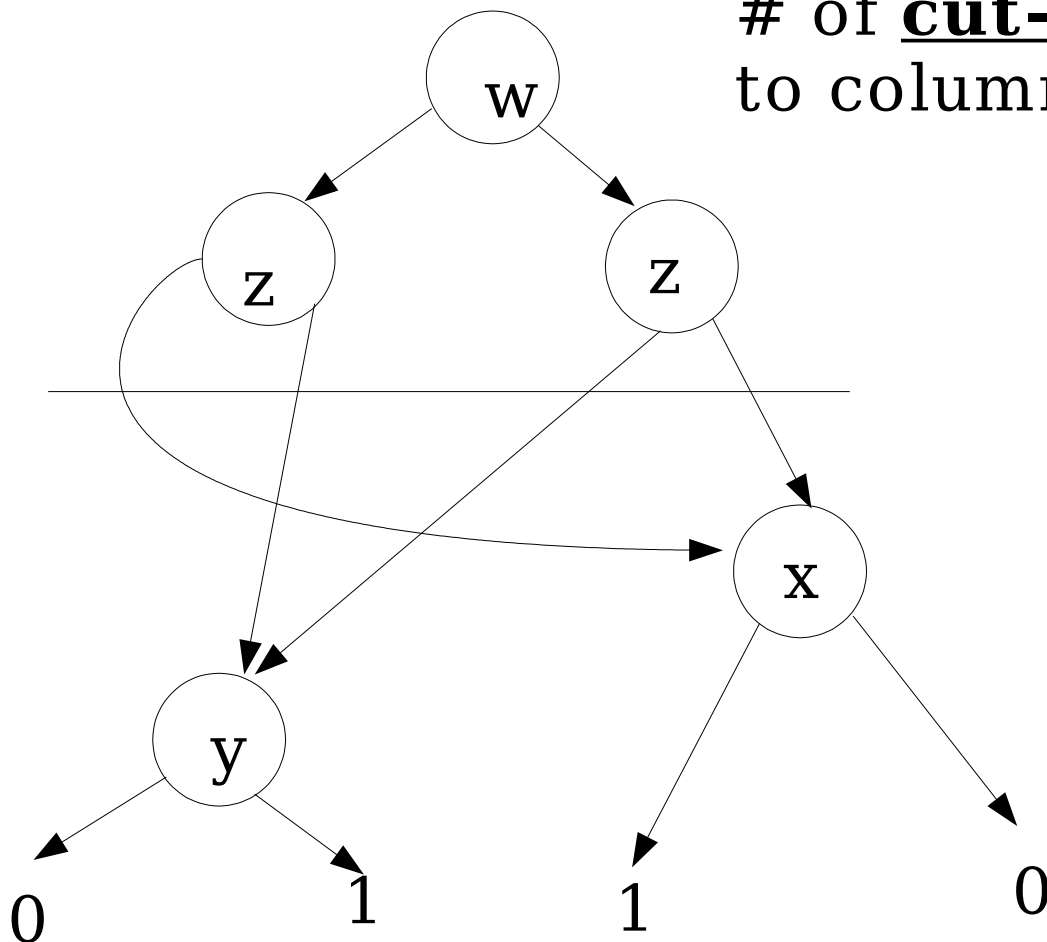
- Column mult = 2:
- Bound set function corresponds to 2 values (0 / 1 or \emptyset/\emptyset')
- One-bit encoding of $\emptyset(w, z)$ exists
- This gives us the simple decomp w.r.t. Orthonormal basis

Simple Decomp on BDDs: Basic Concepts

$$F = w'z'x' + wzx' + w'zy + wz'y$$

$$F = H(G(w, z), x, y)$$

of **cut-set nodes** corresponds
to column multiplicity: BDD MAGIC



Free set term: x'
Associated bound-set
cubes:

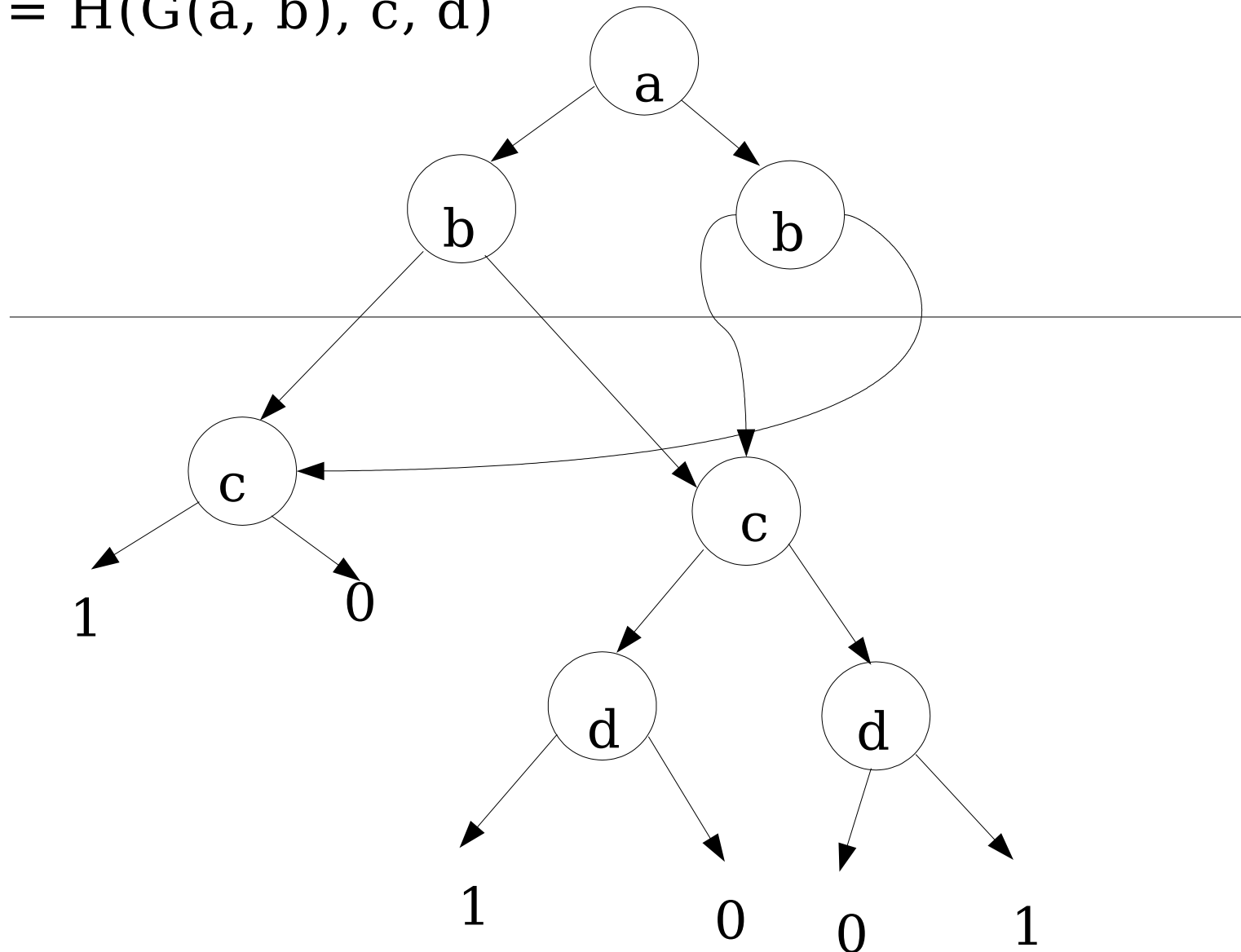
Free-set term: y
Bound-set cubes:

Orthonormality guaranty!

Simple Decomp on BDDs: Basic Concepts

$$F = a'b'c' + a'bc'd' + a'bcd + abc' + ab'cd' + ab'cd$$

$$F = H(G(a, b), c, d)$$



BDD-Based Simple Decomp. Algorithm

- $F = H(G(w, z), x, y)$
- Build BDD for F .
- Bound set vars together on top
- Free set on Bottom
- Perform a “CUT”: Partitioning of vars
- #cut-set nodes ≤ 2 : orthonormal expansion corresponding to simple decomp exists! Problem solved!
- **But what if #cut-set nodes = 3 ?**