

(A.L.L.) Demonstrați Teorema L.L.L.

Fie sistemul de inecuații liniare:

$$(S_L) \sum_{j=1}^m a_{ij} x_j \leq b_i, \text{ pt } i \in \overline{1, m}$$

Fie  $x_k$  o variabilă pt eliminare.

$$K_0 := \{i \mid a_{ik} = 0\}, \quad K_+ := \{i \mid a_{ik} > 0\}, \quad K_- := \{i \mid a_{ik} < 0\}$$

Fie sistemul  $(S_L)$ :  $\sum_{j=1}^m a_{ij} x_j \leq b_i \quad i \in K_0$ .

$$a_{ik} \left( \sum_{j=1}^m a_{ij} x_j \right) - a_{ek} \left( \sum_{j=1}^m a_{ej} x_j \right) \leq a_{ik} b_i - a_{ek} b_e$$

$$\text{pt } \forall (i, e) \in K_+ \times K_-$$

$$\bullet \text{ Fie } P := \{x \in \mathbb{R}^n \mid x \text{ sol pt } S_L\}$$

$$P_{\perp} := \{x' \in \mathbb{R}^{n-1} \mid x' \text{ sol pt } S_L\}$$

Avem  $P^k \Rightarrow$  proiectez  $\text{lu } P$  pe direcția  $x_k$ .

$$P^k = \{ (x_1, \dots, x_{k-1}, x_{k+1}, \dots, x_n) \mid (x_1, \dots, x_n) \in P \}$$

Teorema  $P^k \supseteq P_{\perp}$

Demonstratie: Vom auto mbei or  $P^k \subseteq P_1$

He  $x \in P^k$ .  $x = (x_1 \dots x_{k-1}, x_{k+1} \dots x_n)$

$\exists x_k \in \mathbb{R}$  at  $x = (x_1 \dots x_{k-1}, x_k, x_{k+1} \dots x_n) \in P$ .

• Vom auto or i)  $\sum_{j=1}^n a_{ij} x_j \leq b_i$  ;  $i \in K_0$  este verificat de  $x$ .

$$\sum_{j=1}^{k-1} a_{ij} x_j + \underbrace{a_{ik}}_0 x_k + \sum_{j=k+1}^n a_{ij} x_j \leq b_i \text{ deoarece } x \in P.$$

Deu  $x$  verifică meg i) din sistemul S2.

• Vom auto or ii)  $x$  verifică meg.

$$a_{ik} \left( \sum_{j=1}^n a_{ij} x_j \right) - a_{ik} \left( \sum_{j=1}^n a_{ij} x_j \right) \leq a_{ik} b_i - a_{ik} b_i$$

$$(i, k) \in K_+ \times K_-$$

Deoarece  $x \in P$  se verifică:

$$(1) \sum_{j=1}^n a_{ij} x_j \leq b_i \quad | \cdot (-a_{ik})$$

Cum  $a_{ik} \leq 0$ ,  $-a_{ik} \geq 0$

$$\Rightarrow -a_{ik} \sum_{j=1}^n a_{ij} x_j \leq -a_{ik} b_i$$

$$(2) \sum_{j=1}^n a_{ij} x_j \leq b_i \quad | \cdot a_{ik}, a_{ik} > 0$$

$$a_{ik} \sum_{j=1}^n a_{ij} x_j \leq a_{ik} b_i$$

Adunând (1) și (2)  $\Rightarrow x^*$  satisface și această inegalitate.

$x^*$  este max optimă

Ann demonstret das  $p^k \in P_L$ .

Vom demonstre  $P_L \in P^k$

Pre  $x \in P_L$ . Vom aufba  $\exists x^k \in \mathbb{R}^{n^k}$  art-

$x = (x_1, \dots, x_{k-1}, x_{k+1}, \dots, x_n)$ ,  $x \in P_L$ .

$x' = (x_1, \dots, x_{k-1}, x_k, x_{k+1}, \dots, x_n)$

Per  $x' \in P$ .

$x' \in P$  das

$$\sum_{j=1}^n a_{ij} x_j \leq b_i, \quad i=1, \dots, m$$

Poden neue inequalities aufbauen auf:

$$a_{ik} x_k \leq b_i - \underbrace{\sum_{j=1}^{k-1} a_{ij} x_j^0 - \sum_{j=k+1}^n a_{ij} x_j^0}_{c_i}$$

Poden neue...  $a_{ik} x_k \leq c_i$

$\exists$  obere  $\exists$  pt  $i \in K_0$   $x$  verfasst inequalities aufbauen

$$\text{pt } i \in K_+ \Rightarrow x_k \leq \frac{c_i}{a_{ik}} \Rightarrow 0 \leq \frac{c_i}{a_{ik}} - \frac{c_l}{a_{lk}} \mid (-a_{lk} \cdot a_{ik})$$

$$\text{pt } i \in K_- \Rightarrow x_k \leq \frac{-c_l}{a_{lk}} \quad \Downarrow$$

$$0 \leq -c_i a_{lk} + a_{ik} c_l.$$

~~Kann~~  $\Rightarrow$  eine equivalent an addere ~~alle~~  $\text{ineq}$  der  $S_i$ .

Denn  $x \in P_L \Rightarrow \exists x^k$  art  $x' \in P$  das  $p^k \in P_L$ .

Ann  $P^k \subseteq P_L$  u  $P_L \subseteq P^k \Rightarrow P^k = P_L$ . qed.

(H 2.2) Se  $A \in \mathbb{R}^{m \times n}$  u  $b \in \mathbb{R}^m$ . Systemul  $Ax \leq b$  are solutibile dacă  $\exists$  numerele  $y^t b \geq 0$ . u  $y \in \mathbb{R}^m$  cu  $y \geq 0$  u  $y^t A = 0^t$

Dem: " $\Rightarrow$ " Se  $x_0 \in \mathbb{R}^n$  o solutie a sistemului  $Ax_0 \leq b$ .

Se  $y \geq 0$ ,  $y^t A = 0^t$ .

Atunci  $Ax_0 \leq b \Rightarrow y^t A x_0 \leq y^t b$

$$\Rightarrow 0^t x_0 \leq y^t b$$

$$\Rightarrow 0 \leq y^t b \Rightarrow y^t b \geq 0 \quad \checkmark$$

" $\Leftarrow$ " Se  $y \geq 0$  u  $y^t A = 0$ . Conform

Conform ipotezei  $y^t b \geq 0$ .

Pp ca  $\forall x \in \mathbb{R}^n$   $Ax > b \Rightarrow y^t Ax > y^t b$

$$\Rightarrow 0^t x > y^t b$$

$$\Rightarrow 0 > y^t b \geq 0 \quad \underline{\underline{\text{de}}}$$

Atunci  $\exists x_0 \in \mathbb{R}^n$  cu  $Ax_0 \leq b$ . qed.



(A1.3) Un poliedru este o mulțime convexă.

Un poliedru este intersecția unui număr finit de semispafii.

Un semispafiu este definit astfel:

$$H = \{ x \in \mathbb{R}^n \mid a^T \cdot x \leq \beta \mid a \in \mathbb{R}^n, \beta \in \mathbb{R} \}$$

O mulțime convexă  $C \subseteq \mathbb{R}^n$  dacă:

$$\forall x_1, x_2 \in C, \lambda_1, \lambda_2 \geq 0, \lambda_1 + \lambda_2 = 1 \\ \lambda_1 x_1 + \lambda_2 x_2 \in C$$

Dem. Pe  $x_1, x_2 \in P$  poliedru.

Cum  $P$  este ~~un~~ afecțat ca intersecția unui număr finit de semispafii de  $H$  un semispafiu ales arbitrar.

$$\text{Cum } x_1, x_2 \in P \Rightarrow x_1, x_2 \in H \Rightarrow \begin{aligned} a^T \cdot x_1 &\leq \beta \\ a^T x_2 &\leq \beta. \end{aligned}$$

Pentru  $\lambda_1, \lambda_2 \geq 0$  cu  $\lambda_1 + \lambda_2 = 1$ .

$$a^T \cdot x_1 \leq \beta \mid \cdot \lambda_1 \Rightarrow a^T \lambda_1 x_1 \leq \lambda_1 \beta$$

$$a^T x_2 \leq \beta \mid \cdot \lambda_2 \Rightarrow a^T \lambda_2 x_2 \leq \lambda_2 \beta$$

$$a^T (\lambda_1 x_1 + \lambda_2 x_2) \leq \underbrace{(\lambda_1 + \lambda_2)}_1 \beta$$

Deci  $\lambda_1 x_1 + \lambda_2 x_2 \in P$ , deci  $P$  este convexă qed.

(H.L.U) Lem Theorem 1.3.8.

Se  $x$  o soluto feasible a eu LP, ~~yes~~ ou  $y$  o soluto feasible a dual LP. Assume  $x, y$  int optima duals snt complementa pt (P) n (D).

obs  $x, y$  snt complementa duals  $y_i (b_i - a_i x) = 0, \forall i=1, \dots, m$

Dem  $\leftarrow$  Se  $x, y$  snt  $y_i (b_i - a_i x) = 0, \forall i=1, \dots, m$

$$y^T A x = \sum_{i=1}^m y_i (a_i x) = \sum_{i=1}^m y_i \cdot b_i = y^T \cdot b$$

Se  $y^T A x = c^T x$  deoar  $y^T A = c^T$  ( $y$  e feasible a dual)

Se  $c^T x = y^T b \xrightarrow[\text{deoar}]{\text{Dualite}} x, y$  snt optima.

$\Rightarrow$  Se  $x, y$  optima.

Seoar (P) n (D) snt feasible  $\xRightarrow[\text{Seoar}]{\text{Dualite}} y^T b = c^T x = y^T A x \leq y^T b$

$$\Rightarrow y^T (b - A x) = 0$$

$$\Rightarrow \sum_{i=1}^m y_i (b_i - a_i x) = 0 \Rightarrow \underbrace{\sum_{i=1}^m y_i}_{\geq 0} \cdot \underbrace{(b_i - a_i x)}_{\geq 0 (Ax \leq b)} = 0 \Rightarrow y \geq 0$$

$$y_i (b_i - a_i x) = 0 \quad \forall i=1, \dots, m$$

ged.

(A1.5) Dem is duala dualis lvs LP st echivalent cu primala lvs LP.

Avem problema primala LP:

$$\max \{c^t x \mid Ax \leq b\}.$$

$$\text{Avem duala } \min \{y^t b \mid y \geq 0, y^t A = c^t\}$$

Vom construi duala dualis.

$$\text{Avem } y^t A = c^t.$$

$$\forall x \text{ cu } Ax \leq b. \quad \Rightarrow \quad y^t Ax = c^t x.$$

$$\text{Dar } \cancel{y^t Ax} \quad c^t x = y^t Ax \leq y^t b \quad (y \geq 0)$$

$$\text{Av } c^t x \leq y^t b, \quad Ax \leq b.$$

$$\text{Deci duala dualis st } \max \{c^t x \mid Ax \leq b\}$$

care coincide cu primala gata