

$$13. \quad P(x) = 99 \mid (\text{begin}(\text{set } x \ 3) \ x) = 3$$

$$\frac{x \in \text{dom} \quad P(x) = 99}{\langle \text{VAR}(x), \underline{\underline{99}}, \phi, P \rangle \Downarrow \langle 99, \underline{\underline{99}}, \phi, P \rangle} \quad \langle \text{LITERAL}(3), \underline{\underline{3}}, \phi, P \rangle \Downarrow \langle 3, \underline{\underline{3}}, \phi, P \rangle$$

$$\frac{x \in \text{dom } P \quad P(x) = 3}{\langle \text{VAR}(x), \underline{\underline{3}}, \phi, P \rangle \Downarrow \langle 3, \underline{\underline{3}}, \phi, P \rangle}$$

$$\langle \text{SET}(x, 3), \underline{\underline{3}}, \phi, P \rangle \Downarrow \langle 3, \underline{\underline{3}}, \phi, P \{x \mapsto 3\} \rangle$$

$$\langle \text{BEGIN}(\text{SET}(x, 3), \text{VAR}(x)), \underline{\underline{3}}, \phi, P \rangle \Downarrow \langle 3, \underline{\underline{3}}, \phi, P \rangle$$



14. IF

$$\langle \text{IF}(\text{VAR}(x), \text{VAR}(x), \text{LITERAL}(0)), \xi, \emptyset, P \rangle \Downarrow \langle V_1, \xi', \emptyset, P' \rangle$$

and

$$\langle \text{VAR}(x), \xi, \emptyset, P \rangle \Downarrow \langle V_2, \xi'', \emptyset, P'' \rangle$$

Then Prove

$$V_1 = V_2$$

~~Let~~ Let  $\text{VAR}(x)$  evaluate to true;

Formal Definition of IF

$$\text{TRUE: } \frac{\langle e_1, \xi, \emptyset, P \rangle \Downarrow \langle V_1, \xi', \emptyset, P' \rangle \quad V_1 \neq 0 \quad \langle e_2, \xi', \emptyset, P' \rangle \Downarrow \langle V_2, \xi'', \emptyset, P'' \rangle}{\langle \text{IF}(e_1, e_2, e_3), \xi, \emptyset, P \rangle \Downarrow \langle V_2, \xi'', \emptyset, P'' \rangle}$$

$$\text{FALSE: } \frac{\langle e_1, \xi, \emptyset, P \rangle \Downarrow \langle V_1, \xi', \emptyset, P' \rangle \quad V_1 = 0 \quad \langle e_3, \xi', \emptyset, P' \rangle \Downarrow \langle V_3, \xi'', \emptyset, P'' \rangle}{\langle \text{IF}(e_1, e_2, e_3), \xi, \emptyset, P \rangle \Downarrow \langle V_3, \xi'', \emptyset, P'' \rangle}$$

The formal Definition of IF requires 3 expressions, if the first expression is true, then evaluate the second expression, If the first expression is false then evaluate the third. If we suppose  $e_1 = \text{VAR}(x)$ ,  $e_2 = \text{VAR}(x)$  &  $e_3 = \text{LITERAL}(0)$  and we let  $e_1$  be true, then by definition  $e_2$  must be evaluated. By the Second premise we know  $\text{VAR}(x)$  evaluates to  $V_2$   
 $\therefore$  if  $e_1$  is true,  $\text{VAR}(x) = V_2$  and  $\text{VAR}(x) = V_1$ , so  $V_1 = V_2$

Now Suppose  $e_1$  is false, so  $e_1 = 0$  and we now evaluate the third expression or  $e_3$ .  $e_3$  is the  $\text{LITERAL}(0)$  so  $e_3 = 0$  and so  $V_1 = 0$ .

But since  $e_1$  evaluated to 0 that must mean that  $\text{VAR}(x) = 0$  and by the second premise again  $\text{VAR}(x) = V_2 \rightarrow 0 = V_2$

$\therefore V_1 = V_2$  in all cases  
 Both Q.E.D



21.

a) Awk-like (Create Global)

$$\frac{x \notin \text{dom } P \quad x \notin \text{dom } E \quad \langle e, E, \phi, P \rangle \Downarrow \langle V, E', \phi, P \rangle}{\langle \text{SET}(x, e), E, \phi, P \rangle \Downarrow \langle V, E' \{x \mapsto \emptyset\}, \phi, P' \rangle}$$

b) Icon-like (Create Locally)

$$\frac{x \notin \text{dom } P \quad x \notin \text{dom } E \quad \langle e, E, \phi, P \rangle \Downarrow \langle V, E', \phi, P' \rangle}{\langle \text{SET}(x, e), E, \phi, P \rangle \Downarrow \langle V, E', \phi, P' \{x \mapsto \emptyset\} \rangle}$$

c) Even though Global Variables are frowned upon, I think in this case I would prefer the Awk-like change since if you do not yourself specify the domain of the variable, then the variable should be accessible everywhere by default. ~~If you don't specify the domain~~ If you randomly create a variable in a locally intended domain then you should create it in that domain. Not specifying to me means "this variable already exists outside of this function and I wish to use it here" then if that variable is not found it is created Globally.



20.

~~Begin with~~

$$\langle e_0, \xi, \phi, P \rangle \Downarrow \langle v_1, \xi', \phi, P' \rangle$$

$$\langle e_0, \xi, \phi, P \rangle \Downarrow \langle v_2, \xi'', \phi, P'' \rangle$$

$$\text{if } e_0 = v_1 \quad e_0 = v_2 \quad \therefore v_1 = v_2$$

~~$$\langle e_0, \xi, \phi, P \rangle \Downarrow \langle v_1, \xi', \phi, P' \rangle$$~~

$$\# \quad \text{fv}(\text{VAR}(x)) = \{x\}$$

$$\frac{x \in \text{dom } P}{\langle \text{VAR}(x), \xi, \phi, P \rangle \Downarrow \langle P(x), \xi, \phi, P \rangle}$$

$$\text{Var}(x) = e_0 \quad P(x) = v_1$$

$$\text{Since } x \in \text{dom } P \\ \text{fv}(\text{VAR}(x)) = \{x\} \subseteq \text{dom } P \subseteq (\text{dom } P \cup \text{dom } \xi)$$

$$\text{So } \frac{x \in \text{dom } P}{\langle \underset{e_0}{\text{VAR}(x)}, \xi, \phi, P \rangle \Downarrow \langle v_1, \xi, \phi, P \rangle}$$

$$\text{fv}(e_0) = \{v_1\} \subseteq \text{dom } P \subseteq (\text{dom } P \cup \text{dom } \xi)$$

$$\text{Since } x = v_2 \notin \text{dom } P \\ \therefore P(x) = v_2$$

$$e_0 \rightarrow P(x) = v_1$$

$$e_0 \rightarrow P(x) = v_2$$

$$\therefore v_1 = v_2$$