Transformacija gramatike

- 1. Statements → Statements; Statement | Statement
- 2. Statement → Assignment | IfStatement
- 3. If $Statement \rightarrow if$ (RelExpression): { Statements }
- 4. *RelExpression* → *Expression* **eq** *Expression*
- 5. Assignment \rightarrow **ID** = Expression
- 6. Expression \rightarrow **ID** | **CONST**

- 1. Statements → Statement Statements'
- 2. Statements' → ; Statement Statements'
- 3. Statements' $\rightarrow \epsilon$
- 4. *Statement* → *Assignment*
- 5. $Statement \rightarrow IfStatement$
- 6. *IfStatement* → **if** (*RelExpression*) : { *Statements* }
- 7. $RelExpression \rightarrow Expression$ **eq** Expression
- 8. Assignment \rightarrow ID = Expression
- 9. Expression \rightarrow ID
- 10. Expression \rightarrow CONST

Primenom postupka za eliminaciju levo-rekurzivnih smena elimininisana je direktno leva rekurzija u skupu smena:

Statements → *Statement* | *Statement* |

Dobijen je sledeći skup smena:

- Statements → Statement Statements'
- Statements' → ; Statement Statements'
- Statements $\rightarrow \epsilon$

Odredjivanje first i follow skupova

	Smena	FIRST(SMENA)						
1.	Statements → Statement Statements'	FIRST(Statement Statements') = { ID , if }						
2.	Statements' → ; Statement Statements'	<pre>FIRST(; Statement Statements') = { ; }</pre>						
3.	Statements' $\rightarrow \epsilon$	$FIRST(\epsilon) = \{\epsilon\}$						
4.	Statement → Assignment	FIRST(Assignment) = { ID }						
5.	Statement → IfStatement	FIRST(IfStatement) = { if }						
6.	IfStatement → if (RelExpression) : { Statements }	FIRST(if (RelExpression) : { Statements }) = { if }						
7.	RelExpression → Expression eq Expression	FIRST(Expression eq Expression) = { ID, CONST }						
8.	Assignment \rightarrow ID = Expression	$FIRST(ID = Expression) = \{ID\}$						
9.	Expression → ID	FIRST(ID) = { ID }						
10.	Expression → CONST	FIRST(CONST) = { CONST }						

1. Statements

} ∈ FOLLOW(*Statements*) # ∈ FOLLOW(*Statements*)

2. Statements'

FOLLOW(*Statements*) ⊂ FOLLOW(*Statements* ')

3. Statement

FIRST(Statements') \subset FOLLOW(Statement) FOLLOW(Statements) \subset FOLLOW(Statements) FOLLOW(Statements') \subset FOLLOW(Statement)

4. IfStatement

FOLLOW(*Statement*) ⊂ FOLLOW(*IfStatement*)

5. RelExpression

 $) \in FOLLOW(RelExpression)$

6. Assignment

FOLLOW(*Statement*) ⊂ FOLLOW(*Assignment*)

7. Expression

 $eq \in FOLLOW(Expression)$

FOLLOW(*RelExpression*) ⊂ FOLLOW(*Expression*)

FOLLOW(*Assignment*) ⊂ FOLLOW(*Expression*)

	Simbol	FOLLOW(Simbol)					
1.	Statements	FOLLOW(Statements) = { }, # }					
2.	Statements'	FOLLOW(Statements') = { }, # }					
3.	Statement	FOLLOW(<i>Statement</i>) = { ;, }, # }					
4.	IfStatement	FOLLOW(<i>IfStatement</i>) = { ;, }, # }					
5.	RelExpression	FOLLOW(RelExpression) = {) }					
6.	Assignment	FOLLOW(<i>Assignment</i>) = { ;, }, # }					
7.	Expression	FOLLOW(<i>Expression</i>) = { eq ,) , ; , } , # }					

Analiza skupova smena

Skupovi smena { *Statements*' \rightarrow ; *Statement Statements*', *Statements*' \rightarrow ϵ }, { *Expression* \rightarrow **CONST** }, { *Statement* \rightarrow *Assignment*, *Statement* \rightarrow *IfStatement*} imaju na levoj strani smene iste terminalne simbole. Za prvi skup važi:

- 1) FIRST(; Statement Statements') \cap FIRST(ϵ) = \emptyset
- 2) Samo jedna od navedenih smena se slika u prazan niz
- 3) FIRST(; Statement Statements') \cap FOLLOW(Statements') = \emptyset

Za drugi skup važi:

- 1) FIRST(**ID**) \cap FIRST(**CONST**) = \emptyset
- 2) Nijedana od smena se ne slika u prazan niz

Za treci skup važi:

- 1) FIRST(Assignment) \cap FIRST(IfStatement) = \emptyset
- 2) Nijedna od smena se ne slika u prazan niz

Za sve skupove smena važe uslovi formalne definicje LL(1) gramatike.

Sintaksna tabela

- 1. $SS \rightarrow SSS$
- 2. $SS' \rightarrow SS'$
- 3. $SS' \rightarrow \epsilon$
- 4. $S \rightarrow A$
- 5. $S \rightarrow IfS$
- 6. If $S \rightarrow if$ (RelE): {SS}
- 7. $RelE \rightarrow E eq E$
- 8. $A \to ID = E$
- 9. $E \rightarrow \mathbf{ID}$
- 10. E → **CONST**

Vrh	Ulazni simboli											
magacina	;	if	()	:	{	}	eq	ID	=	CONST	#
SS		(S SS', 1)							(S SS', 1)			
SS'	(; S SS', 2)						(e, 3)					(e,, 3)
S		(IfS, 5)							(A, 4)			
IfS		(if (RelE): { SS }, 6)										
RelE									$(E \mathbf{eq} E, 7)$		(E eq E, 7)	
A									(ID = E, 8)			
E									(ID , 9)		(CONST, 10)	
;	pop											
if		рор										
(pop									
)				pop								
:					pop							
{						pop						
}							pop					
eq								pop				
ID									pop			
=										pop		
CONST											pop	
#												acc