CSE2315 — Assignment 2

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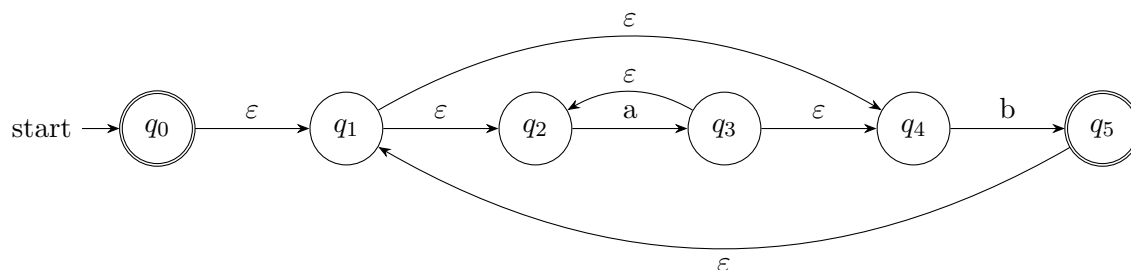
Exercise 1

Consider the language $L = (a^*b)^*$. Construct an NFA N for which $L(N) = L$. To this end, answer the following:

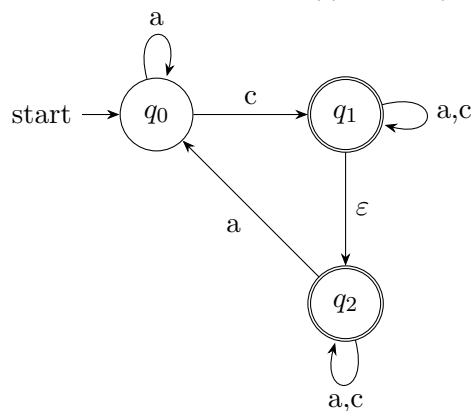### (a) Construct an NFA N_1 for which $L(N_1) = \{x\}$.

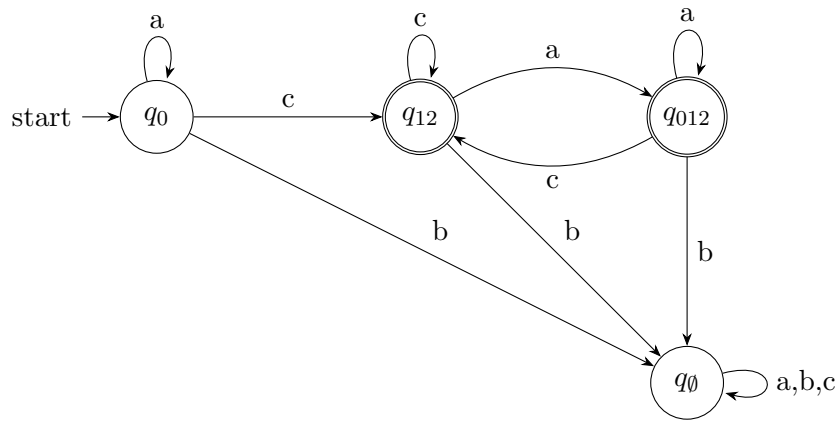### (b) Construct N in a systematic way starting from N_1.

Exercise 2

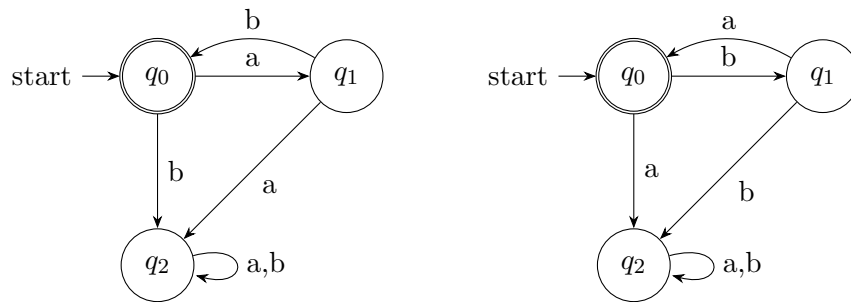Consider the NFA $N = \{\{q_0, q_1, q_2\}, \{a, b, c\}, \delta, q_0, \{q_1, q_2\}\}$ depicted below:

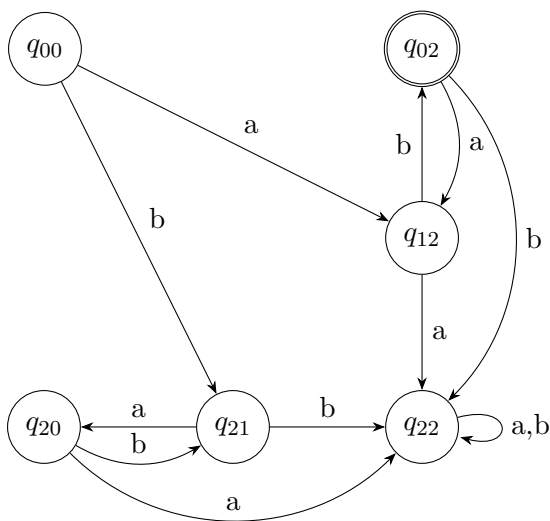Construct a DFA D such that $L(D) = L(N)$:

Exercise 3

Consider the following two DFAs D_1, D_2:

(a)   Describe the languages of D_1 and D_2

$$L(D_1) = (ab)^*, L(D_2) = (ba)^*$$

(b)   Construct a DFA D such that $L(D) = L(D_1) \cap L(D_2)^c$.

Exercise 4

Suppose we have the following language L over the alphabet $\Sigma = \{a, b, c\}$:

$L = \{w \in \Sigma^* | w$ has an a and every a after the first a in w is immediately followed by a $c\}$

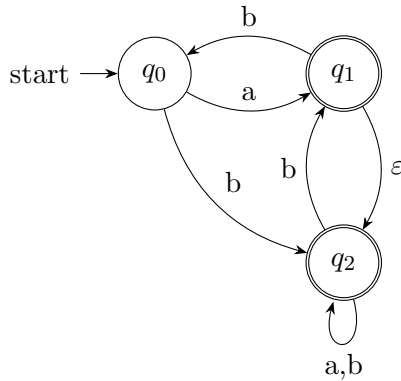### (a) Give a regular expression R such that $L(R) = L$.

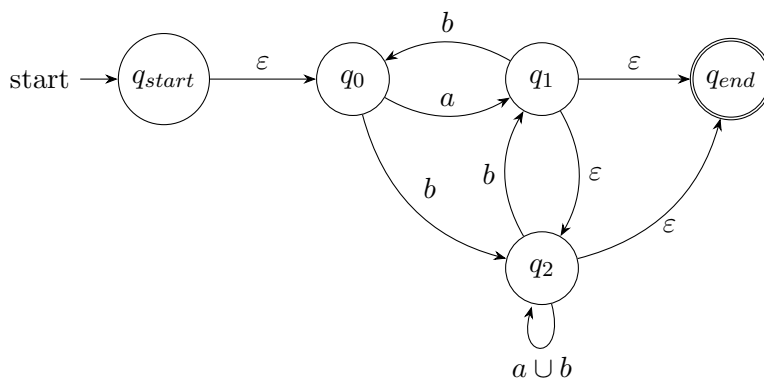$$R = (b \cup c)^* a (b \cup c \cup ac)^*$$

(b) Explain the answer

In order to enforce the one a in the word, we have to go $(b \cup c)^* a (b \cup c)^*$. And then to add ac blocks after the first a, we simply or those ac blocks with the $b \cup c$.
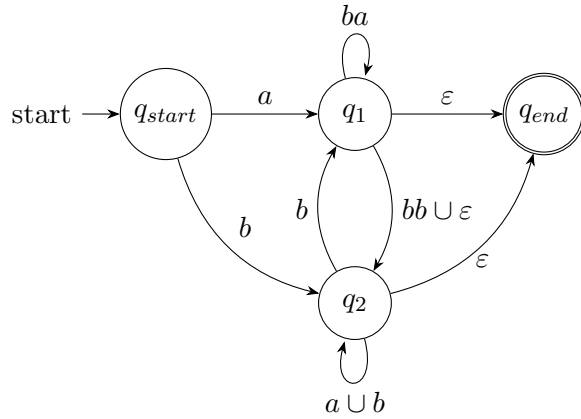
Exercise 5

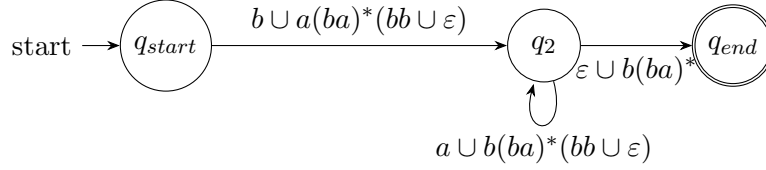Convert the following NFA N to a regular expression R such that $L(R) = L(N)$

We begin by converting it to a GNFA (I omit arrows on which the regular expresion would be $\emptyset$):
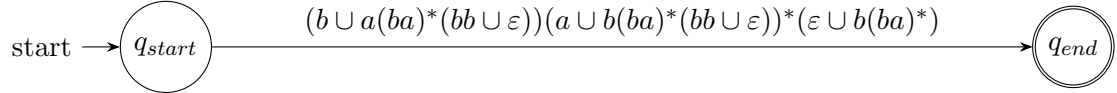

Then, we rip out the state q_0:
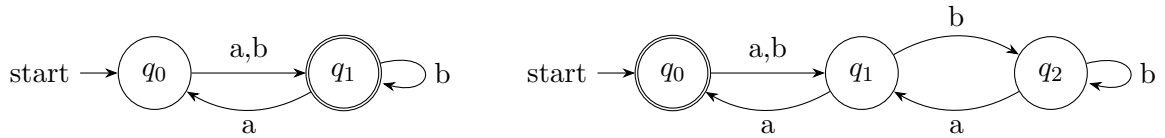
Now, we rip out q_1:

Now, we rip out q_2

Therefore, $R = (b \cup a(ba)^*(bb \cup \varepsilon))(a \cup b(ba)^*(bb \cup \varepsilon))^*(\varepsilon \cup b(ba)^*)$, which simplifies nicely to:
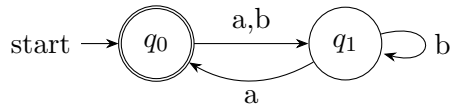
$$\boxed{R = \Sigma^+}$$

Exercise 6

Consider the following two DFAs, D on the left, D' on the right:
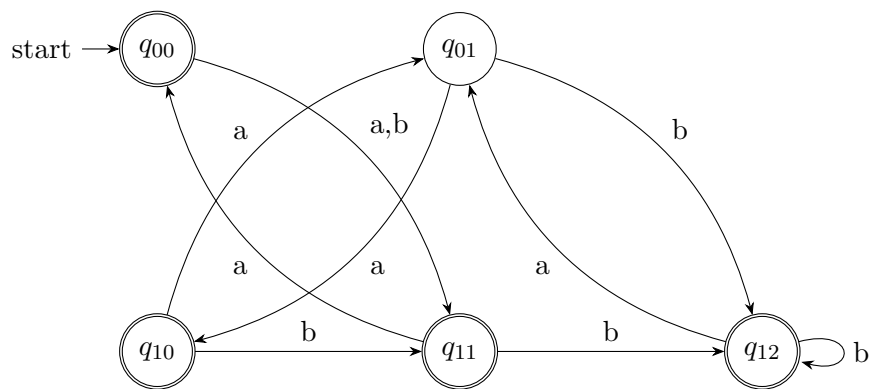

(a)   Construct a new DFA D'' of at most 2 states, such that $L(D'') = \overline{L(D)}$

All you have to do is invert accepting and non-accepting states.

(b) **Construct a DFA** D'''**, such that** $L(D''') = L(D) \cup L(D')$

Bonus Exercise