

CSE2315 — Assignment 2

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Exercise 1

Suppose we have the following language L over the alphabet $\Sigma = \{a, 0, 1\}$:

$$L = \left\{ ua^k v \mid u, v \in \{0, 1\}^* \text{ and } c_0(u) + c_1(v) > k \right\}$$

Where $c_x(w)$ is defined as the number of occurrences of x in w . Use the pumping lemma to show that L is not regular.

Proof. By contradiction.

Assume to the contrary that L is regular. Let p be the pumping length given by the pumping lemma. Choose s to be the string $a^p 1^p 1$. Because s is a member of L and s has length more than p , the pumping lemma guarantees that s can be split into three pieces, $s = xyz$, where for any $i \geq 0$ the string $xy^i z$ is in L . We show this outcome is impossible.

Condition 3 of the pumping lemma says that $|xy| \leq p$. Thus, in the case of our word, y must consist of only as . Let k_{xyz} and k_{xyyz} represent the number of as in xyz and $xyyz$ respectively. $k_{xyz} = p$ and $k_{xyyz} > k_{xyz} = p$, but the number of 1s in $xyyz$ is still $p + 1$. Thus in $xyyz$ we have $c_0(u) + c_1(v) \leq k_{xyyz}$, so $xyyz$ cannot be in L . Therefore, s cannot be pumped and we have reached our desired contradiction. \square

Exercise 2

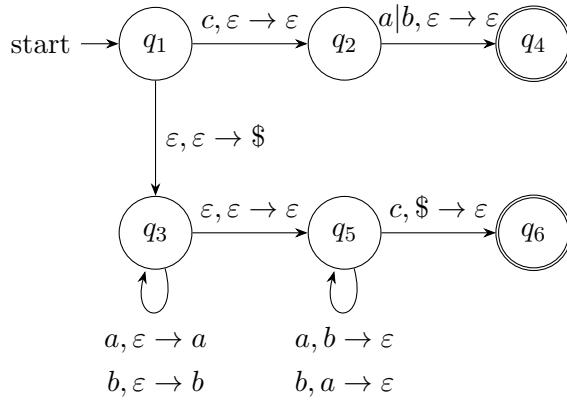
Suppose we have the context free grammar $G = (\{S, T, V\}, \{a, b, c\}, R, S)$, with R containing the following rules:

$$\begin{aligned} S &\rightarrow Tc \mid cV \\ T &\rightarrow aTb \mid bTa \mid \varepsilon \\ V &\rightarrow a \mid b \end{aligned}$$

(a) Describe the set $L(G)$ in your own words.

$L(G)$ includes the words ca and cb as well as antipalindromes formed of as and bs with a c appended to it.

(b) Construct a PDA M such that $L(M) = L(G)$. Use no more than 8 states.



(c) Explain how your PDA works.

So the PDA first nondeterministically chooses if it's in the c first or c last case. In the c first case, it doesn't use the stack, it just reads a c and then an a or a b from the input. In the other case, it starts pushing the characters it sees in the input and then nondeterministically decides where the middle of the word is. After the middle, it pops characters from the stack and makes sure that they are the opposite of what it's reading from the input. At the end, the stack should be empty and the last character should be a c .

Exercise 3

Consider the following rules R of a CFG $G = (\{S, A, B\}, \{a, b\}, R, S)$:

$$\begin{aligned} S &\rightarrow bB \mid abB \\ A &\rightarrow bA \mid AA \\ B &\rightarrow AS \mid SA \mid \epsilon \end{aligned}$$

Convert G into Chomsky normal form.

First, we add a new start variable.

$$\begin{aligned} S_0 &\rightarrow S \\ S &\rightarrow bB \mid abB \\ A &\rightarrow bA \mid AA \\ B &\rightarrow AS \mid SA \mid \epsilon \end{aligned}$$

Then, we remove the ϵ -rules.

$$\begin{aligned} S_0 &\rightarrow S \\ S &\rightarrow bB \mid abB \mid b \mid ab \\ A &\rightarrow bA \mid AA \\ B &\rightarrow AS \mid SA \mid \epsilon \end{aligned}$$

Third, we remove all unit rules.

$$\begin{aligned}
 S_0 &\rightarrow S \mid \mathbf{bB} \mid \mathbf{abB} \mid \mathbf{b} \mid \mathbf{ab} \\
 S &\rightarrow bB \mid abB \mid b \mid ab \\
 A &\rightarrow bA \mid AA \\
 B &\rightarrow AS \mid SA
 \end{aligned}$$

Convert the remaining rules into proper form.

$$\begin{aligned}
 S_0 &\rightarrow VB \mid TB \mid b \mid UV \\
 S &\rightarrow VB \mid TB \mid b \mid UV \\
 A &\rightarrow VA \mid AA \\
 B &\rightarrow AS \mid SA \\
 T &\rightarrow UV \\
 U &\rightarrow a \\
 V &\rightarrow b
 \end{aligned}$$

Exercise 4

Let L be the set of all TUDON words.

- (a) Give a CFG G such that $L(G) = L$, using no more than 6 variables and 9 + 26 rules.

Let $G = (\{S, T, U, V, X\}, \{<, >, :, ;, a, b, \dots, z\}, R, S)$, with R containing the following rules:

$$\begin{aligned}
 S &\rightarrow < T > \\
 T &\rightarrow U : V X \\
 X &\rightarrow ; T \mid \varepsilon \\
 U &\rightarrow \Sigma U \mid \varepsilon \\
 V &\rightarrow S \mid U
 \end{aligned}$$

Where Σ refers to any character a, b, \dots, z

- (b) Explain why your grammar generates L

So basically the start variable makes one of those dictionaries, and then defers to a new variable for the contents of the dict. Then that variable T makes a key value pair, appended with an X. Then the X can be nothing, or a semicolon and then another T. And then the actual key is any number of alphabet characters, and the value is also any number of alphabet characters or another dictionary.

- (c) Is your grammar ambiguous? Give an example or an argument why not.

The grammar is not ambiguous, the only part where it could be ambiguous, is when creating the key-value pairs in the dictionary, you could substitute them and creating them in different orders, but I take care of that by appending them one by one do what is already there. This way, there is no ambiguous word.

Exercise 5

Let the operator @ be defined as $A@ = \{a^{|w|} \mid w \in A\}$. Prove that regularity is closed under @.

Proof. Let A be a regular language. We will prove that $A@$ is also regular.

Since A is regular, there is a DFA $M = (Q, \Sigma, \delta, q_0, F)$ such that $L(M) = A$. Now we will build the NFA $M' = (Q, \{a\}, \Delta, q_0, F)$, where $\Delta(q, a) = \{\delta(q, x) \mid x \in \Sigma\}$. In other words, take M and replace every symbol on transition arrows with a .

Next, we must show that the set of states reachable in M' by reading n a 's from the input is equivalent to the set of states reachable in M by reading any n characters from an input word. This can be shown using induction. It is left as an exercise to the reader. \square

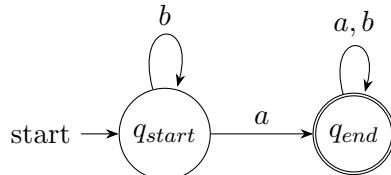
Exercise 6

- (a) Let N be an arbitrary NFA with k states. Explain why if $L(N) \neq \emptyset$, then there is a word of length $\leq k$ in $L(N)$.

If $L(N) \neq \emptyset$, then N must recognize some word. This means that there must be a path from the starting state to some accepting state given by a set of transitions. More formally, take any $w \in L(N)$, where $w = w_1 w_2 \dots w_m$. There must be a sequence of states r_0, r_1, \dots, r_m , such that r_0 is the starting state, r_m is an accepting state, and for all $0 \leq i < m$, $r_{i+1} \in \delta(r_i, w_{i+1})$. Now, if $m \leq k$, we are done. However, if $m > k$, then by the pigeonhole principle some state must be repeated. So, in the sequence $r_0, \dots, r_x, \dots, r_x, \dots, r_m$, we can remove everything between the r_x s, including one of the r_x s themselves, and we can keep doing this procedure until the length of the sequence becomes $\leq k$. Now, this sequence exactly describes an input word of length $\leq k$ that is in $L(N)$.

- (b) Let D be an arbitrary DFA with $k > 1$ states. Give a counterexample to the claim that there is a DFA D' with $k - 1$ states, such that $L(D') = L(D)$.

Let D be the DFA described by the following state diagram (assume the alphabet is $\{a, b\}$):



This DFA only recognizes words that contain at least one a inside. If we were to construct a DFA with $k - 1$ states, it would have to have 1 state. This state could either be accepting, or not. Which means D' would either accept every word, or no words. Neither of which are equivalent to D .

Bonus Exercise