

# CSE2315 — Assignment 2

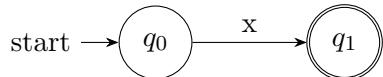
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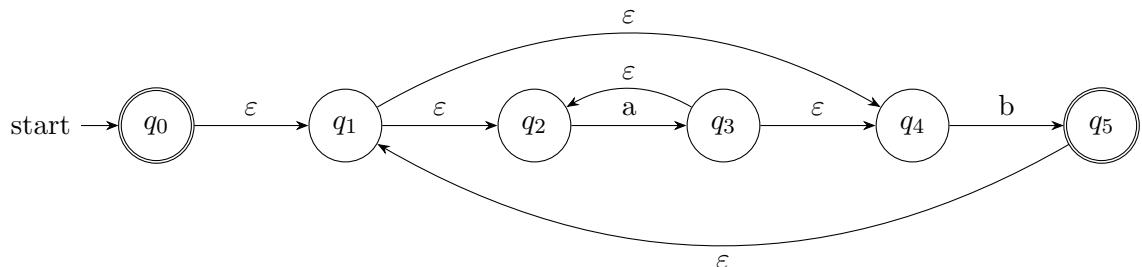
## Exercise 1

Consider the language  $L = (a^*b)^*$ . Construct an NFA  $N$  for which  $L(N) = L$ . To this end, answer the following:

- (a) Construct an NFA  $N_1$  for which  $L(N_1) = \{x\}$ .

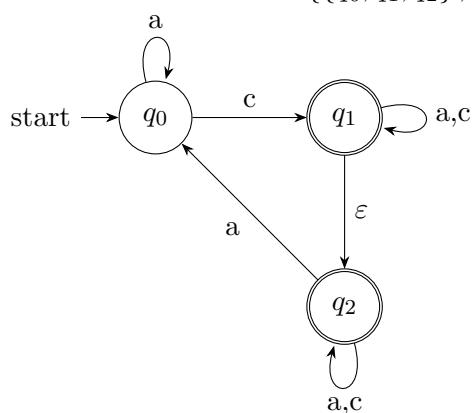


- (b) Construct  $N$  in a systematic way starting from  $N_1$ .

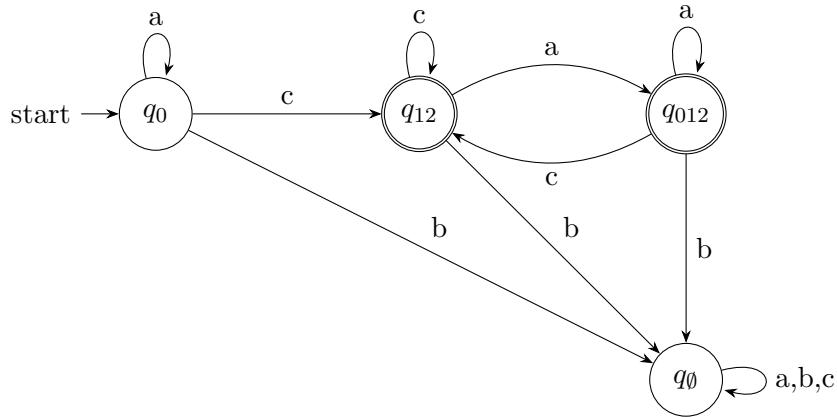


## Exercise 2

Consider the NFA  $N = \{\{q_0, q_1, q_2\}, \{a, b, c\}, \delta, q_0, \{q_1, q_2\}\}$  depicted below:

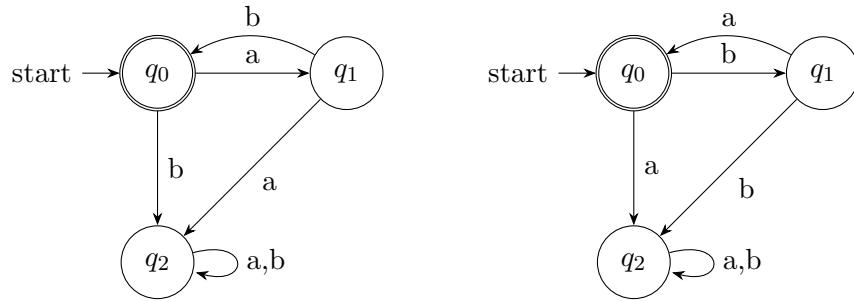


Construct a DFA  $D$  such that  $L(D) = L(N)$ :



### Exercise 3

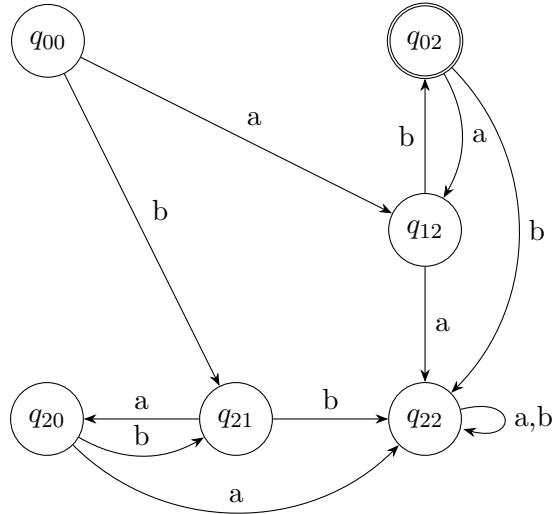
Consider the following two DFAs  $D_1, D_2$ :



(a) Describe the languages of  $D_1$  and  $D_2$

$$L(D_1) = (ab)^*, L(D_2) = (ba)^*$$

(b) Construct a DFA  $D$  such that  $L(D) = L(D_1) \cap L(D_2)^c$ .



## Exercise 4

Suppose we have the following language  $L$  over the alphabet  $\Sigma = \{a, b, c\}$ :

$$L = \{w \in \Sigma^* \mid w \text{ has an } a \text{ and every } a \text{ after the first } a \text{ in } w \text{ is immediately followed by a } c\}$$

- (a) Give a regular expression  $R$  such that  $L(R) = L$ .

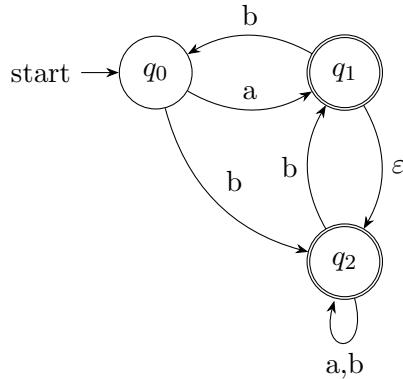
$$R = (b \cup c)^* a (b \cup c \cup ac)^*$$

- (b) Explain the answer

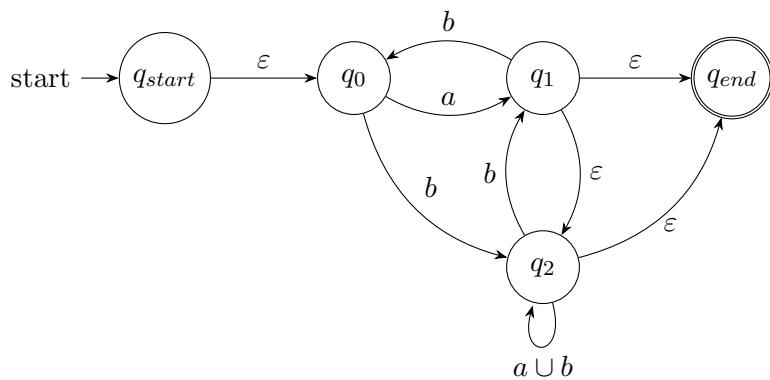
In order to enforce the one  $a$  in the word, we have to go  $(b \cup c)^* a (b \cup c)^*$ . And then to add  $ac$  blocks after the first  $a$ , we simply or those  $ac$  blocks with the  $b \cup c$ .

## Exercise 5

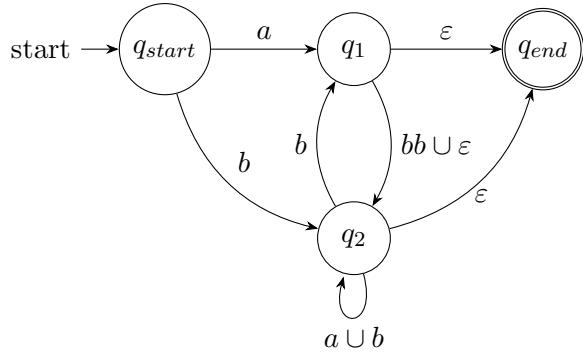
Convert the following NFA  $N$  to a regular expression  $R$  such that  $L(R) = L(N)$



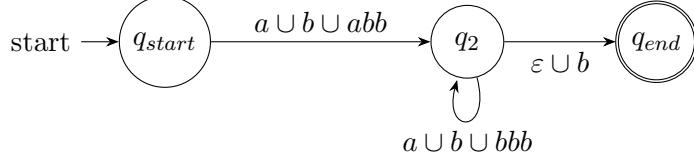
We begin by converting it to a GNFA (I omit arrows on which the regular expression would be  $\emptyset$ ):



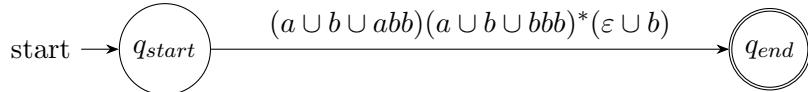
Then, we rip out the state  $q_0$ :



Now, we rip out  $q_1$ :



Now, we rip out  $q_2$

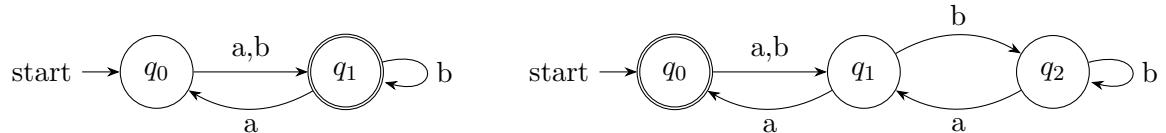


Therefore,  $R = (a \cup b \cup abb)(a \cup b \cup bbb)^*(\varepsilon \cup b)$ , which simplifies nicely to:

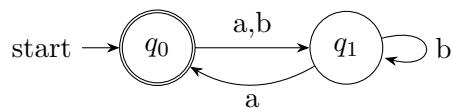
$$R = \Sigma^+$$

## Exercise 6

Consider the following two DFAs,  $D$  on the left,  $D'$  on the right:

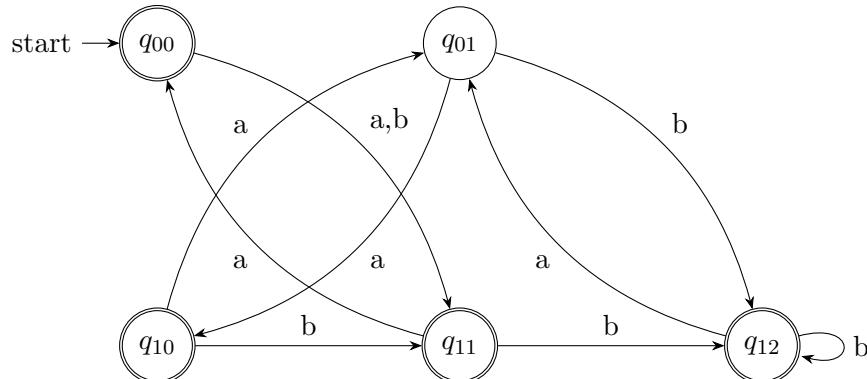


(a) Construct a new DFA  $D''$  of at most 2 states, such that  $L(D'') = \overline{L(D)}$



All you have to do is invert accepting and non-accepting states.

(b) Construct a DFA  $D'''$ , such that  $L(D''') = L(D) \cup L(D')$



## Bonus Exercise