

# 5G PHY Layer Processing – receive processing (2)

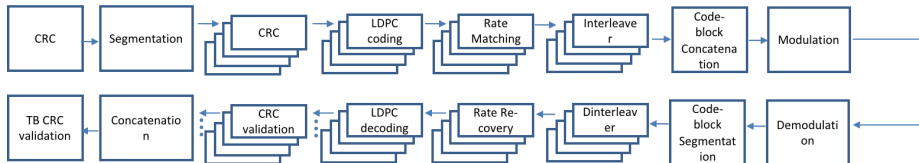
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Simulation-Based Design of 5G Wireless Standard (EE698H)

# Agenda for today

- Discuss receiver processing for the transmit chain discussed
- Discuss scrambling
  - References mentioned later in the slides

# Receiver processing – rate recovery principles



- **Rate matching**

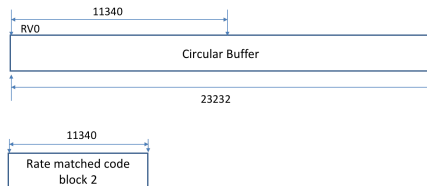
- Puncturing/repeating bits to match the allocated resources

- **Rate recovery**

- LDPC decoder works only for code rate of  $1/3$ . Code rate should be reverted back to  $1/3$
- Zeros are inserted in place of punctured bits
- Repeated bits are combined

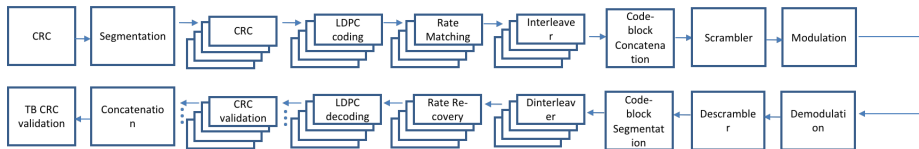
# 5G receiver Processing - rate recovery

- Rate recovery of code blocks e.g., second code block



- Write bits starting from RV0. Write neutral information for bits punctured at transmitter
- Feed rate recovered data to LDPC decoder. Validate CRC of each code block

# 5G transmit and receive chain upto scrambler



# Scrambling principles

- Consider a concatenated block of  $G$  interleaved bits  $b(0), \dots, b(G-1)$
- Scrambled bits  $\tilde{b}(0), \dots, \tilde{b}(G-1)$  are given as

$$\tilde{b}(i) = b(i) \oplus c(i) \quad i = 0, \dots, (G-1)$$

- $c(i)$  is pseudo random sequence
- Example: 8-bit coded sequence  $b = [0\ 0\ 0\ 0\ 0\ 0\ 0\ 0]$  and  $c = [0\ 1\ 1\ 0\ 1\ 0\ 0\ 1]$

$$\tilde{b}(i) = [0\ 1\ 1\ 0\ 1\ 0\ 0\ 1]$$

- Scrambling is done to randomize the output of interleaver
  - both inner and outer signal points in the 16/64/256 QAM constellation to be used

# De-scrambling principles

- For block of bits  $b(0), \dots, b(G-1)$ , where  $G$  is the number of bits in code word
- Received scrambled bits  $\tilde{b}(0), \dots, \tilde{b}(G-1)$  were calculated as

$$\tilde{b}(i) = b(i) \oplus c(i) \quad i = 0, \dots, (G-1)$$

- Received descrambled bits  $\tilde{b}(0), \dots, \tilde{b}(G-1)$  can be recovered as

$$b(i) = \tilde{b}(i) \oplus c(i) \quad i = 0, \dots, (G-1)$$

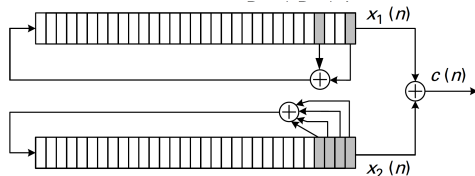
- Example: 8-bit coded sequence  $b = [0\ 0\ 0\ 0\ 0\ 0\ 0\ 0]$  and  $c = [0\ 1\ 1\ 0\ 1\ 0\ 0\ 1]$

$$\tilde{b}(i) = [0\ 1\ 1\ 0\ 1\ 0\ 0\ 1]$$

$$b(i) = \tilde{b}(i) \oplus c(i) = [0\ 0\ 0\ 0\ 0\ 0\ 0\ 0]$$

# Pseudo random sequence generation in 5G<sup>1</sup> (1)

- Pseudo random sequences in 5G are defined by a length-31 Gold sequence



- Output sequence  $c(n)$  of length  $G$  where  $n = 0, 1, \dots, G - 1$

$$c(n) = x_1(n) \oplus x_2(n)$$

$$x_1(n + 31) = x_1(n + 3) \oplus x_1(n)$$

$$x_2(n + 31) = x_2(n + 3) \oplus x_2(n + 2) \oplus x_2(n + 1) \oplus x_2(n)$$

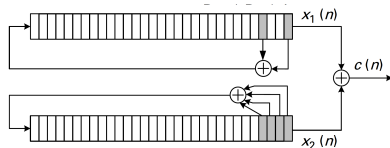
- Standard rejects first 1600 samples and uses  $c'(n) = c(n + 1600)$  instead

<sup>1</sup>Section 5.2 of 38.211



# Pseudo random sequence generation in 5G (2)

- Pseudo random sequences in 5G are defined by a length-31 Gold sequence

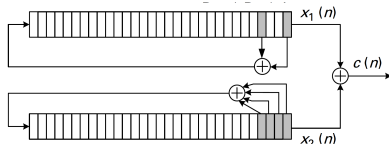


- Output sequence  $c(n)$  of length  $G$  where  $n = 0, 1, \dots, G - 1$
- First sequence  $x_1(n)$  is initialized as

$$\begin{aligned} x_1(n) &= 1 & n &= 0 \\ &= 0 & 0 < n \leq 30 \end{aligned}$$

# Pseudo random sequence generation in 5G (3)

- Pseudo random sequences in 5G are defined by a length-31 Gold sequence



- Output sequence  $c(n)$  of length  $G$  where  $n = 0, 1, \dots, G - 1$
- Second sequence  $x_2(n)$  is initialized by writing a constant  $c_{\text{init}}$  in binary form
- $c_{\text{init}}$  is determined based on a cell ID and RNTI e.g, consider  $c_{\text{init}} = 255$ .  
Second sequence  $x_2(n)$

$$\begin{aligned} x_2(n) &= 1 & n \leq 7 \\ &= 0 & 7 < n \leq 30 \end{aligned}$$