

# EE910: Digital Communication Systems-I

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May 30, 2022



## Lecture #8A: Maximum Likelihood Sequence Estimation: Viterbi Algorithm



## Maximum Likelihood Sequence Estimation

- Any system with memory can be represented by its state diagram or trellis diagram.
- Viterbi algorithm is a computational efficient way to find maximum likelihood sequence estimate.
- We will take an example of a simple error correcting code having four states to illustrate how maximum likelihood sequence estimation works.



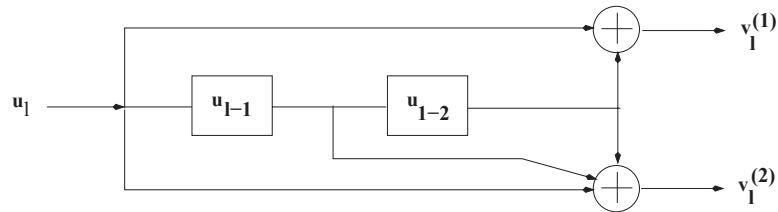
## Viterbi Decoding

- We will consider an example where each information bit is encoded into two coded bits using the error correcting code that has memory two.
- We will assume that the data is transmitted over a binary symmetric channel.



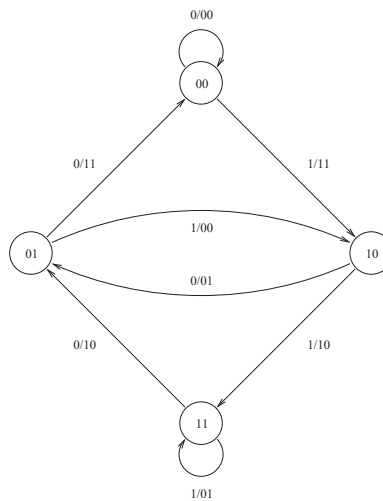
## Viterbi Decoding

Example:



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## Viterbi Decoding



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## Viterbi decoding

On BSC:

- Let the information sequence of length  $L$

$$\mathbf{u} = (u_0, u_1, \dots, u_I, \dots, u_{L-1})$$

is encoded into code sequence of length  $N \triangleq L$

$$\mathbf{v} = (\mathbf{v}_0, \mathbf{v}_1, \dots, \mathbf{v}_I, \dots, \mathbf{v}_{L-1})$$

- If the code sequence  $\mathbf{v}$  is transmitted over a channel, let the received sequence is,

$$\mathbf{r} = (\mathbf{r}_0, \mathbf{r}_1, \dots, \mathbf{r}_I, \dots, \mathbf{r}_{L-1}),$$

where the  $I^{th}$  received block is

$$\mathbf{r}_I = (r_I^{(1)}, r_I^{(2)}).$$

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## Viterbi decoding

On BSC:

- A maximum likelihood decoder finds a path through the trellis that maximizes the path conditional probability

$$P(\mathbf{r}|\mathbf{v}) = \prod_{I=0}^L P(\mathbf{r}_I|\mathbf{v}_I)$$

where the branch conditional probability

$$P(\mathbf{r}_I|\mathbf{v}_I) = \prod_{I=1}^2 P(r_I^{(i)}|v_I^{(i)})$$

- The bit conditional probabilities  $P(r_I^{(i)}|v_I^{(i)})$  are the channel transition probabilities.

Navigation icons: back, forward, search, etc.

## Viterbi decoding

On BSC:

- Maximizing  $P(\mathbf{r}|\mathbf{v})$  is equivalent to maximizing

$$M(\mathbf{r}|\mathbf{v}) \triangleq \log P(\mathbf{r}|\mathbf{v})$$

- $M(\mathbf{r}|\mathbf{v})$  is called the path metric.



## Viterbi decoding

- We have

$$\begin{aligned} M(\mathbf{r}|\mathbf{v}) &= \sum_{l=0}^L \log P(\mathbf{r}_l|\mathbf{v}_l) \\ &= \sum_{l=0}^L M(\mathbf{r}_l|\mathbf{v}_l), \quad (\text{branch metrics}) \end{aligned}$$

$$\begin{aligned} M(\mathbf{r}_l|\mathbf{v}_l) &= \sum_{i=1}^2 \log P(r_l^{(i)}|v_l^{(i)}) \\ &= \sum_{i=1}^2 M(r_l|v_l), \quad (\text{bit metrics}) \end{aligned}$$



## Viterbi decoding

- The partial path metric for the first  $j$  branches of a path  $\mathbf{v}$  is given by

$$M([\mathbf{r}|\mathbf{v}]_j) = \sum_{l=0}^{j-1} M(\mathbf{r}_l|\mathbf{v}_l)$$

- For BSC, the maximum likelihood decoder decodes the received sequence  $\mathbf{r}$  into code sequence  $\mathbf{v}$  that minimizes the Hamming distance  $d(\mathbf{r}, \mathbf{v})$
- The Viterbi algorithm is a computationally efficient method of finding the path through the trellis with the best metric.



## Viterbi decoding

- The Viterbi decoder proceeds through the trellis level by level in search of the path with the best metric.
- At each level, the decoder compares the metric of all partial paths entering each state.
- The decoder stores the partial path entering each state with the best metric (survivor path) and eliminates all other partial paths.
- For  $m \leq l \leq L$ , there are total  $2^m$  survivors.
- To bring the code to zero state is known as termination process.
- The number of survivors decrease during the termination process, until there is only one survivor left.
- The surviving path is the maximum likelihood path.



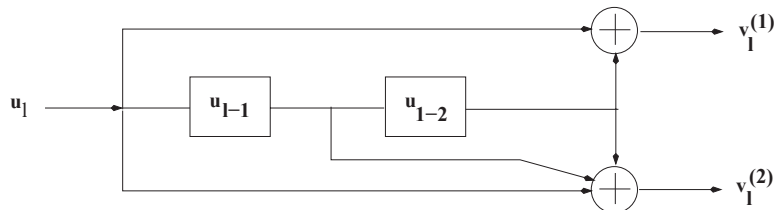
## Viterbi decoding

- Step 1: Starting at level  $l = m$  in the trellis, compute the partial metric for the single path entering each  $m^{th}$  level state. Store the survivor path and its metric for each state.
- Step 2: Increase time  $l$  by one. Compute the partial metric for all the paths entering at the  $(l + 1)^{th}$  level state by adding the branch metric entering that state to the metric of the connecting survivor path at the previous  $l^{th}$  level state. Store the survivor path and its metric for each state.
- Step 3: Repeat Step 2 until you are at the end of the trellis.

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## Viterbi Decoding

Example:

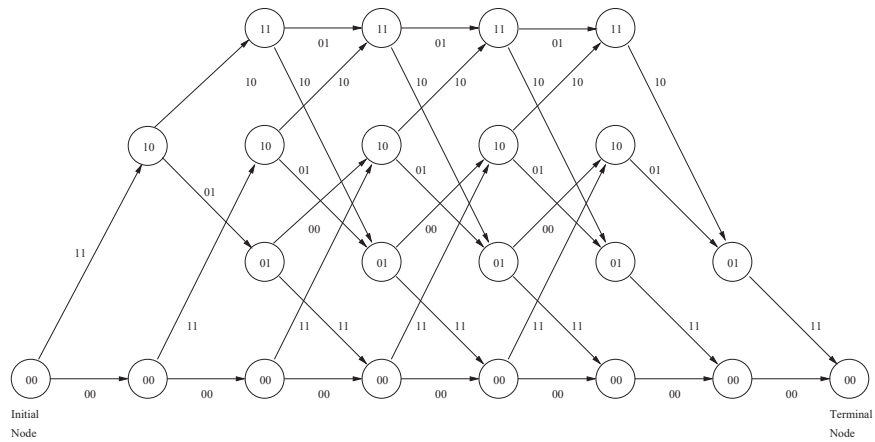


- This (2,1,2) convolutional code with  $L = 7$  (including termination bits) is used on a BSC. The received sequence is

$$\mathbf{r} = (01, 11, 10, 10, 00, 11, 10)$$

Navigation icons: back, forward, search, etc.

## Viterbi Decoding



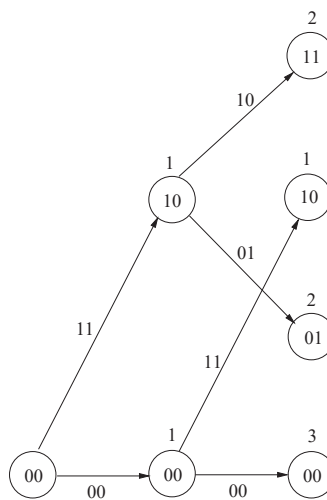
Trellis diagram of (2, 1, 2) convolutional code.

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## Viterbi Decoding



Level 2

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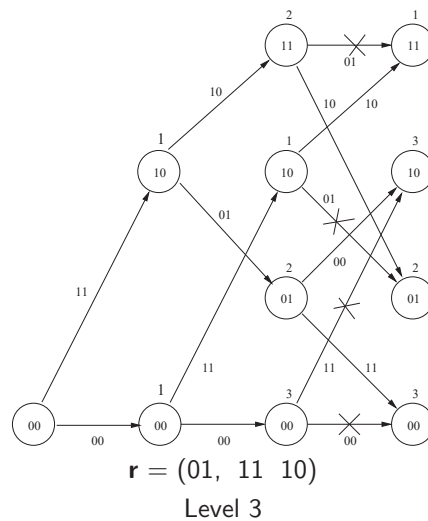
## Viterbi Decoding

$$\mathbf{r} = (01, 11)$$

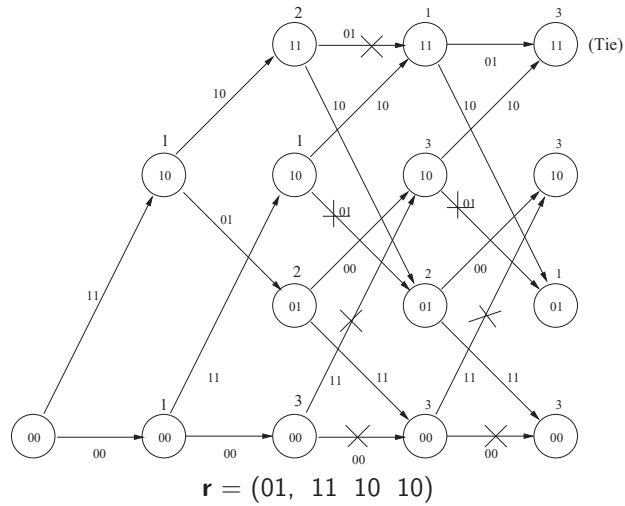
$\mathbf{v}_1 = (00, 00)$	$d(\mathbf{v}_1, \mathbf{r}) = 3$
$\mathbf{v}_2 = (00, 11)$	$d(\mathbf{v}_2, \mathbf{r}) = 1$
$\mathbf{v}_3 = (11, 01)$	$d(\mathbf{v}_3, \mathbf{r}) = 2$
$\mathbf{v}_4 = (11, 10)$	$d(\mathbf{v}_4, \mathbf{r}) = 2$



## Viterbi Decoding



## Viterbi Decoding

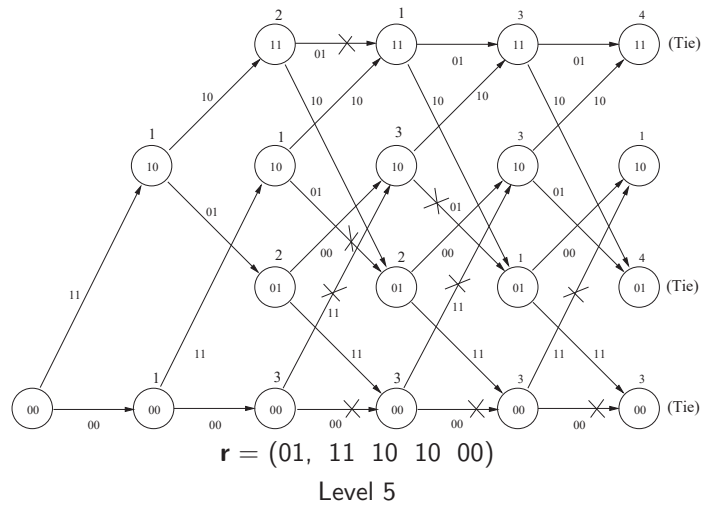


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## Viterbi Decoding

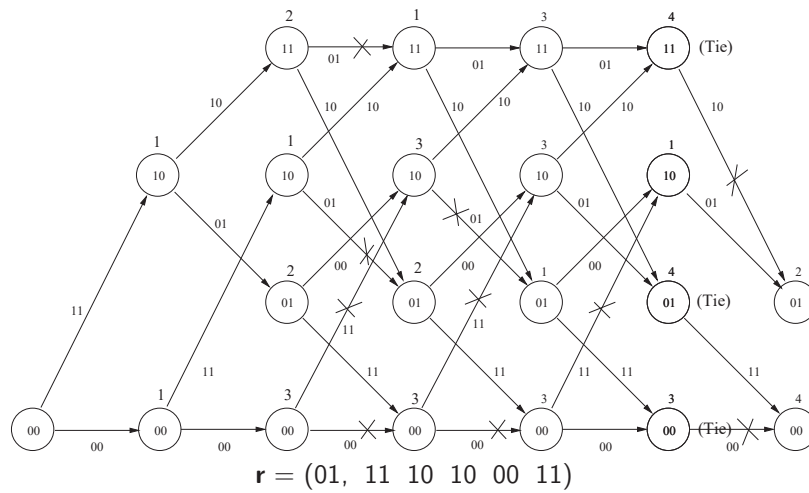


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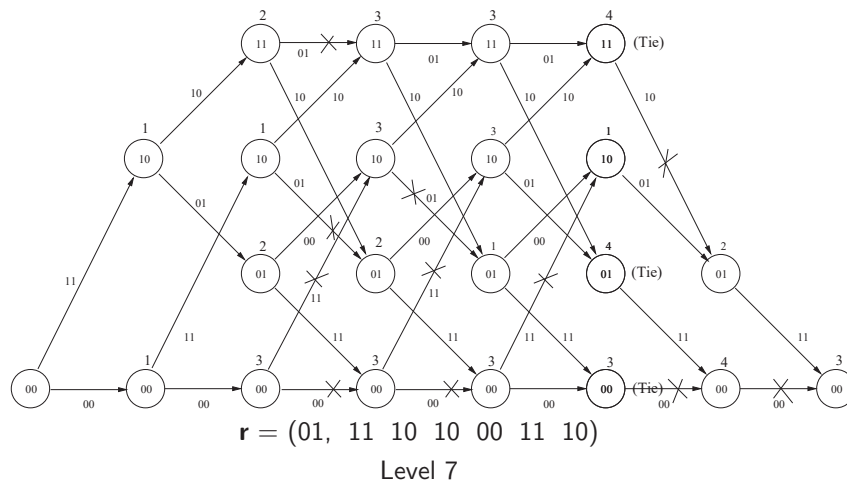


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# Viterbi Decoding

