Parallel Programming

Parallel Histogram

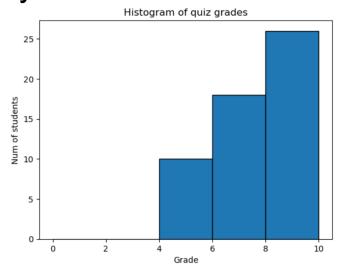
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Overview

- The "histogram" task
- Sequential implementation
- Parallel implementation
 - A naïve kernel
 - A kernel using atomic operation
 - A fast kernel using atomic operation

Histogram

- A method for extracting notable features and patterns from large data sets
 rút trích đặc trưng
 - Feature extraction for object recognition in images
 - Fraud detection in credit card transactions phát hiện giao dịch lỗi
 - Correlating heavenly object movements in astrophysics
- Basic histograms: for each element in the data set, use the value to identify a "bin counter" to increment



The "histogram" task

Input

0	5	4	0
7	1	1	2
2	3	0	6

Histogram 4 bin

Bin	Index	0	1	2	3	
	Range	[0, 1]	[2, 3]	[4, 5]	[6, 7]	
	Count	5	3	2	2	

The "histogram" task

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0	5	4	0
7	1	1	2
2	3	0	6

Histogram 8 bin

Din	Index	0	1	2	3	4	5	6	7
Bin	Range	0	1	2	3	4	5	6	7
Count									

Histogram 4 bin

Bin	Index	0	1	2	3	
	Range	[0, 1]	[2, 3]	[4, 5]	[6, 7]	
Count		5	3	2	2	

The "histogram" task

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0	5	4	0
7	1	1	2
2	3	0	6

Histogram 8 bin

Bin	Index	0	1	2	3	4	5	6	7
	Range	0	1	2	3	4	5	6	7
Count		3	2	2	1	1	1	1	1

Histogram 4 bin

Bin	Index	0	1	2	3	
	Range	[0, 1]	[2, 3]	[4, 5]	[6, 7]	
	Count	5	3	2	2	

Sequential implementation

Compute bin which in[i] falls into
In the rest of this lecture, we'll assume: in[i] ≥ 0 and computeBin(in[i]) = in[i]

Parallel implementation: a naïve kernel

- Let each thread take care of one element in the input array
- Thread will identify the bin of its element and increase the bin counter by 1

```
__global__ void computeHistOnDevice1(int *in, int *hist, int n, int nBins)
{
    int i = blockIdx.x * blockDim.x + threadIdx.x;

    Need synchronize globally, but we can only synchronize within each block ③
    One simple solution: from host, call cudaMemset function ⑤

    if (i < n)
        hist[in[i]] += 1;
}

Read-Modify-Write (RMW); what if more than one threads RMW the same location? ⑥
```

Parallel implementation: a naïve kernel

 A kernel using <u>atomic operation</u>: solve the conflict problem when many threads RMW the same location

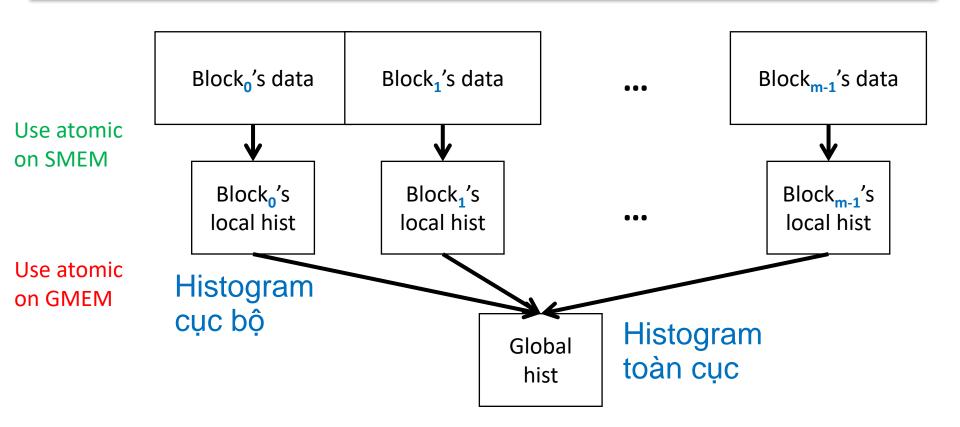
```
__global__ void computeHistOnDevice2(int *in, int *hist, int n, int nBins)
{
    int i = blockIdx.x * blockDim.x + threadIdx.x;
    if (i < n)
        hist[int[i]] += 1;
        atomicAdd(&hist[in[i]], 1);
}</pre>
```

If many threads RMW the same location, they will be **serialized**: threads will RMW in turn

→ Guarantee correctness, but can hurt speed

Coding exercise: file "10-Histogram.cu"

A fast kernel using atomic operation



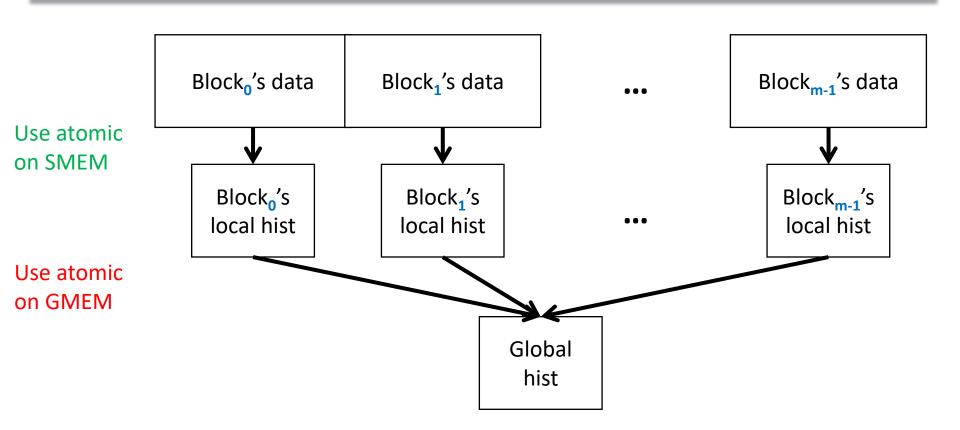
Will it run faster?

Consider a bin; assume in each block's data chunk there are 2 values falling into this bin

The old way costs: ? atomicAdd's on GMEM

This way costs: ?

A fast kernel using atomic operation



Will it run faster?

Consider a bin; assume in each block's data chunk there are 2 values falling into this bin

The old way costs: 2m atomicAdd's on GMEM

This way costs: 2m atomicAdd's on SMEM, 2 atomicAdd's/block and m blocks run in parallel

+ m atomicAdd's on GMEM

Coding exercise: file "10-Histogram.cu"

Reference

- [1] Wen-Mei, W. Hwu, David B. Kirk, and Izzat El Hajj. Programming Massively Parallel Processors: A Hands-on Approach. Morgan Kaufmann, 2022
- [2] Cheng John, Max Grossman, and Ty McKercher. *Professional Cuda C Programming*. John Wiley & Sons, 2014
- [3] Illinois GPU course

https://wiki.illinois.edu/wiki/display/ECE408/ECE408+Home



THE END