Parallel Programming

Parallel Execution in CUDA

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Review of previous lecture

CUDA allows us to organize grid and block as 1D or 2D or 3D

	Compute Capability													
Technical Specifications	5	5.2	5	6	6.1	6.2	7	7.2	7.5	8	8.6	8.7	8.9	9
Maximum number of resident grids per device (Concurrent Kernel Execution)	32		16	128	32	16	128	16	128					
Maximum dimensionality of grid of thread blocks	3													
Maximum x -dimension of a grid of thread blocks	2^{31}-1													
Maximum y- or z-dimension of a grid of thread blocks	65535													
Maximum dimensionality of thread block	3													
Maximum x- or y-dimensionality of a block	1024													
Maximum z-dimension of a block	64													
Maximum number of threads per block	1024													
Warp size	32													
Maximum number of resident blocks per SM	32								16	32	16		24	32
Maximum number of resident warps per SM	64								32	64	48			64
Maximum number of resident threads per SM	2048								1024	2048	1536			2048
Number of 32-bit registers per SM	64 K													
Maximum number of 32-bit registers per thread block	64 K													
Maximum number of 32-bit registers per thread	255													
Maximum amount of shared memory per SM	64 KB	96 KB	64 KB		96 KB	64 KB	96 KB		64 KB	164 KB	100 KB	164 KB	100 KB	228 KB
Maximum amount of shared memory per thread block	48 KB						96 KB	96 KB	64 KB	163 KB	99 KB	163 KB	99 KB	227 KB
Number of shared memory banks	32													
Maximum amount of local memory per thread	512 KB													
Constant memory size	64 KB													
Cache working set per SM for constant memory	8 KB			4 KB	8 KB									

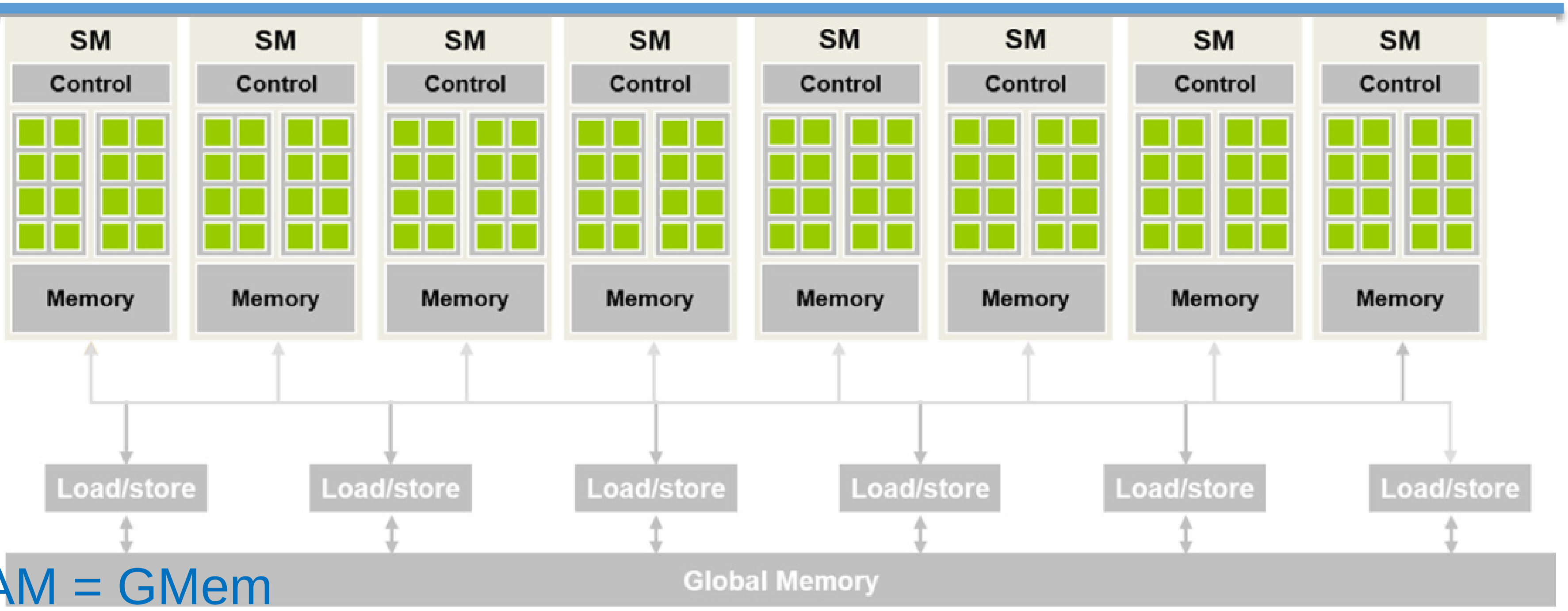
Today: level 2 CUDA

Kiến trúc tính toán của GPU

Aspects of the **GPU compute architecture** that are essential for CUDA C programmers.

- A high-level, simplified view of the **compute architecture** and explore the concepts of flexible *resource assignment*, *scheduling of blocks*, and *occupancy*.
- *Thread scheduling*, *latency tolerance*, *control divergence*, and *synchronization*.
- **API functions** that can be used to query the resources that are available in the GPU and the tools to help estimate the occupancy of the GPU when executing a kernel.

Architecture of a modern GPU



DRAM = GMem

Bộ đa xử lý luồng

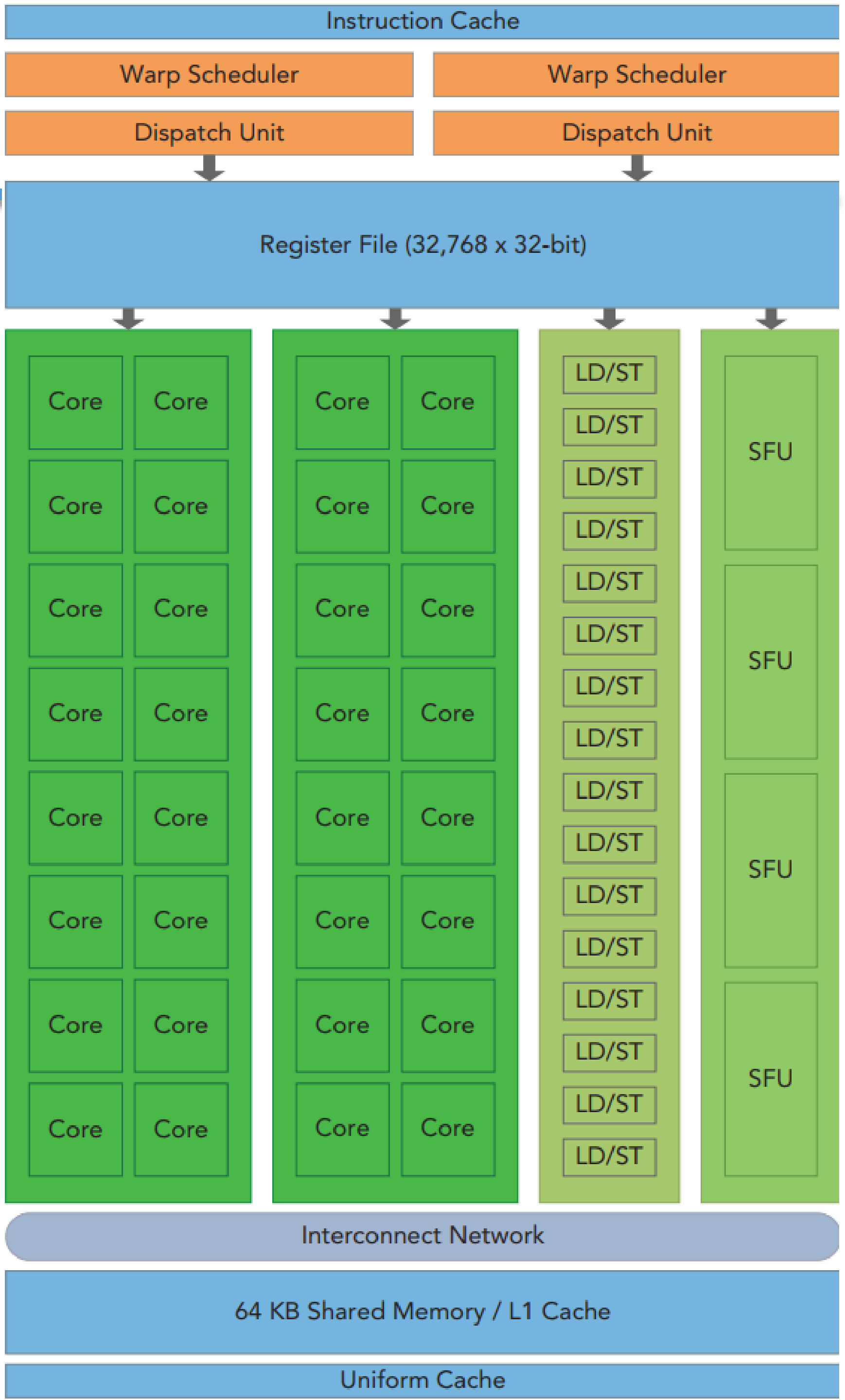
- GPU consists of SMs – Streaming Multiprocessors
  - Each SM consists of SPs – Streaming Processors (or CUDA cores)
  - The SMs have special on-chip memory 1 CUDAcore = 1 SProcessor
- E.g.: the Ampere A100 GPU has 108 SMs with 64 cores each, totaling 6912 cores in the entire GPU

108 SM * 64 core/SM = 6912 core

More about SM

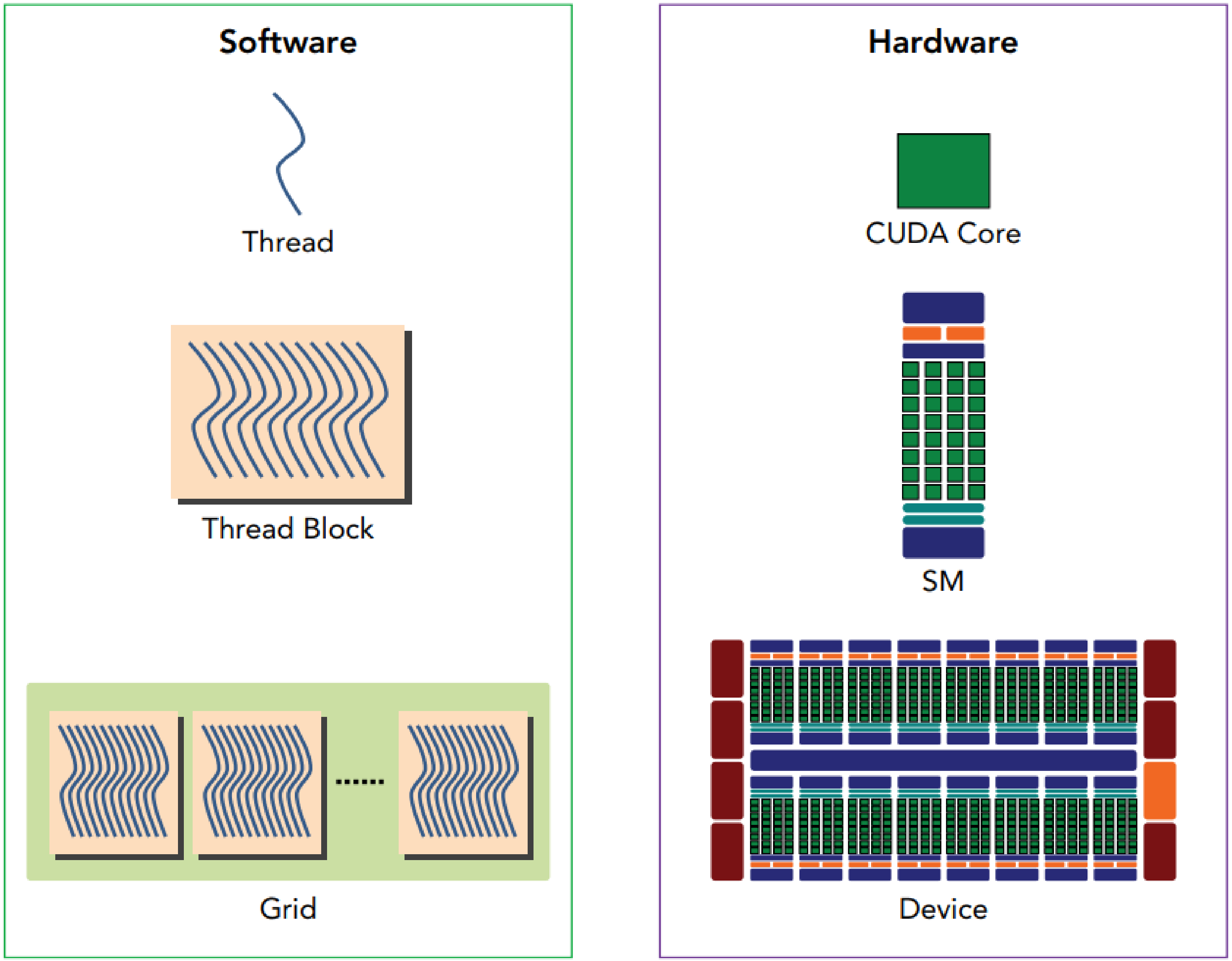
- CUDA Cores
- Shared Memory/L1 Cache
- Register File
- Load/Store Units đọc/ghi bộ nhớ
- Special Function Units SFU
- Warp Scheduler

SFU: tính toán đặc biệt (vd: log, v.v.)
Warp Scheduler: Thực hiện lập lịch
Dispatch Unit: Điều khiển core theo lịch



Fermi SM
Source: Professional CUDA C programming

Architecture of a modern GPU



Parallel execution – SM level

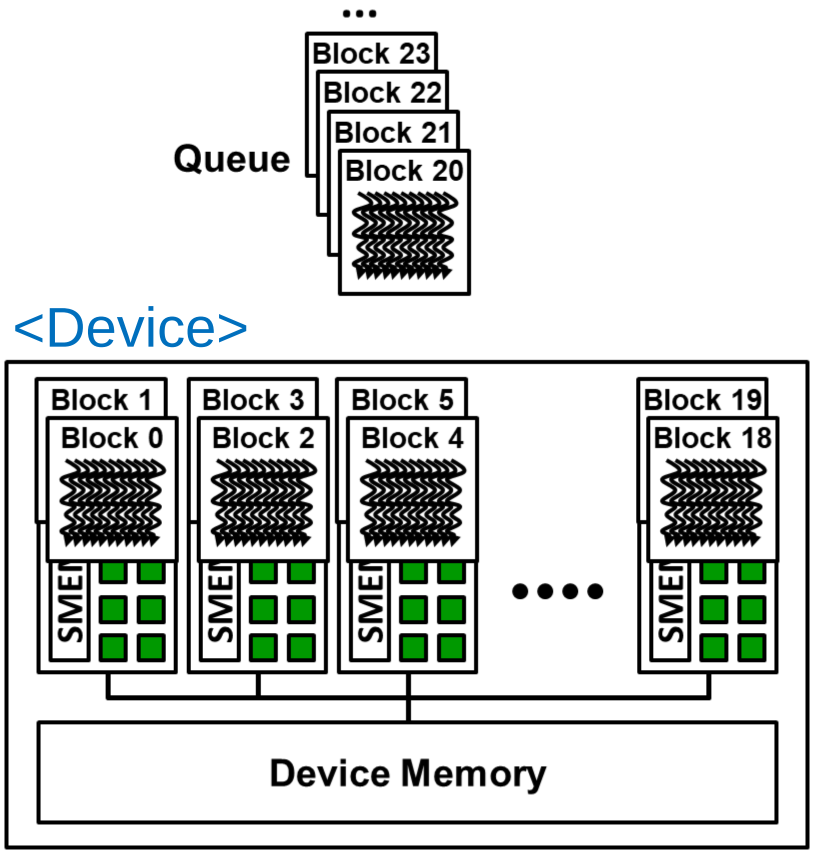
- CUDA virtualizes GPU hardware architecture ảo hóa
  - Block = virtual SM
  - Thread = virtual SP = CUDA core
- When host calls a kernel function, a grid of blocks will be created and each block (virtual SM) will be assigned to a real SM for execution
  - Each SM can contain > 1 block to execute
    - It depends on SM resource limitations and resources each block needs
    - E.g., SM needs resources (registers) to keep track of indexes of blocks and threads as well as their execution state, SM 2.x resources can afford at most 8 blocks and 1536 threads

      If block size is 512 → SM 2.x can contain 3 blocks

      3 blocks * 512 threads/block = 1536 threads

Parallel execution – SM level

- CUDA virtualizes GPU hardware architecture
  - Block = virtual SM
  - Thread = virtual SP
- When host calls a kernel function, a grid of blocks will be created and each block (virtual SM) will be assigned to a real SM for execution
  - Each SM can contain > 1 block to execute
  - Blocks which have not been assigned to SMs will wait in a queue
  - When a block finishes its execution, a block from the queue will be assigned to the available slot in SM



gán theo thứ tự bất kỳ, ta ko kiểm soát được

☐ *Note: blocks can be assigned to SMs in an arbitrary order*

Parallel execution – SM level

Such parallel execution in CUDA:

- Helps achieve scalability ☺

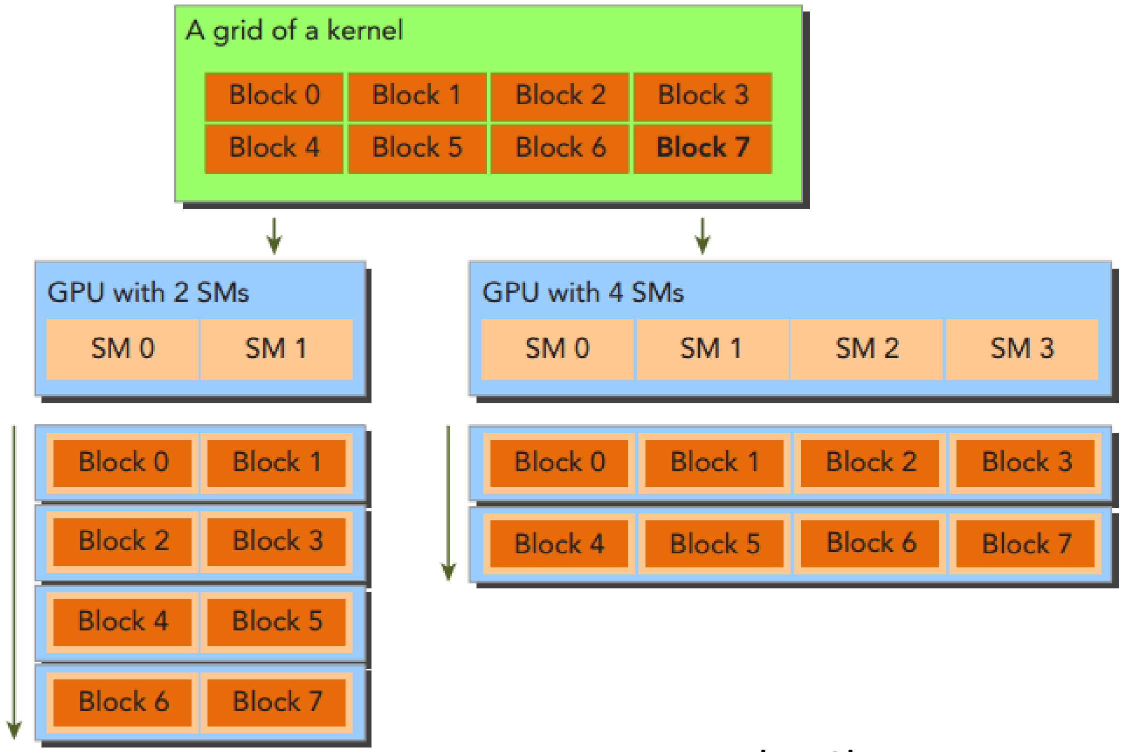


Image source: John Cheng et al. Professional CUDA C Programming. 2014

Parallel execution – SM level

Such parallel execution in CUDA:

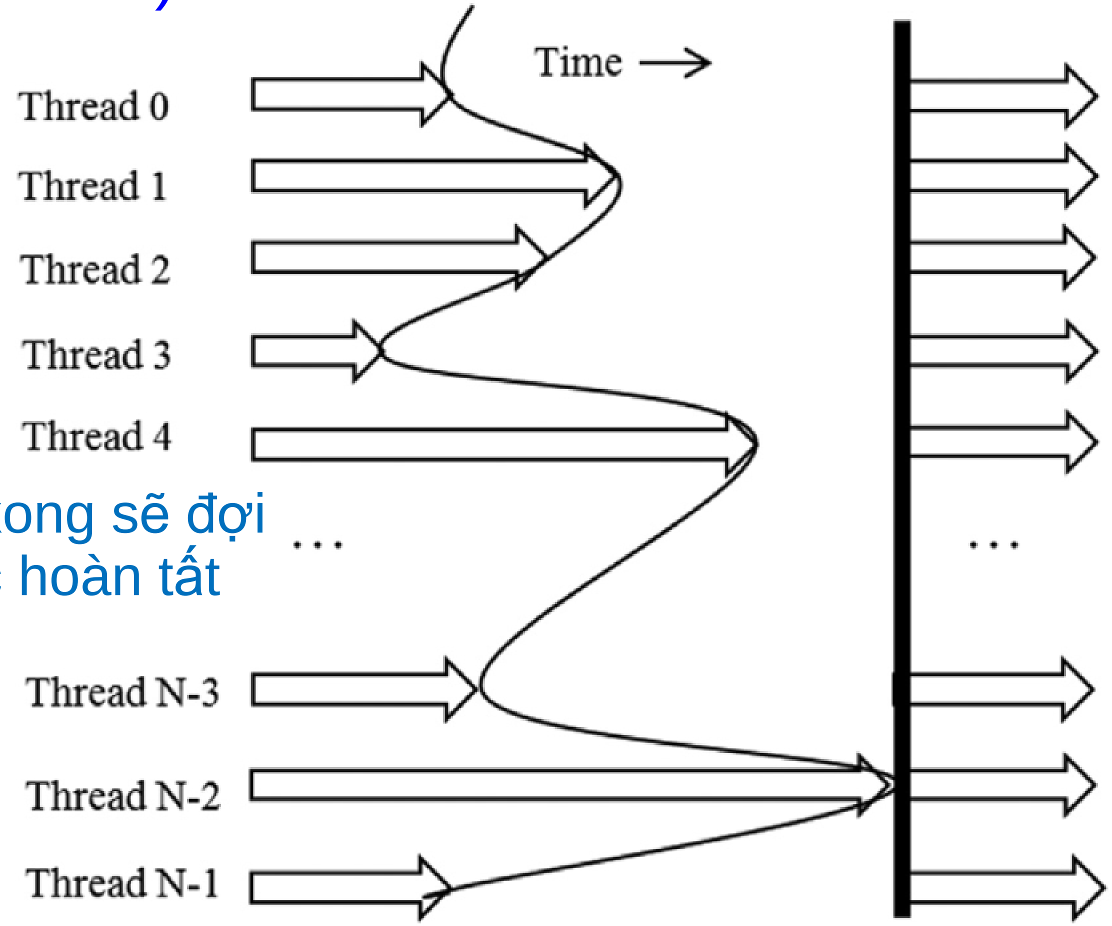
- Helps achieve scalability ☺
- But requires blocks to be independent of each other → threads in different blocks cannot cooperate (synchronize) with each other ☹
  - Assume thread a in block A wants to use a result from thread b in block B,

    and GPU resources can only afford executing one block at a time and currently block A is being executed

    But block B only can be executed when block A is done ☹
  - Threads in the same block can cooperate with each other by using `__syncthreads()` CUDA command

Parallel execution – SM level

An example execution of barrier synchronization (__syncthreads)



Các thread đã xong sẽ đợi các thread khác hoàn tất

Parallel execution – SM level

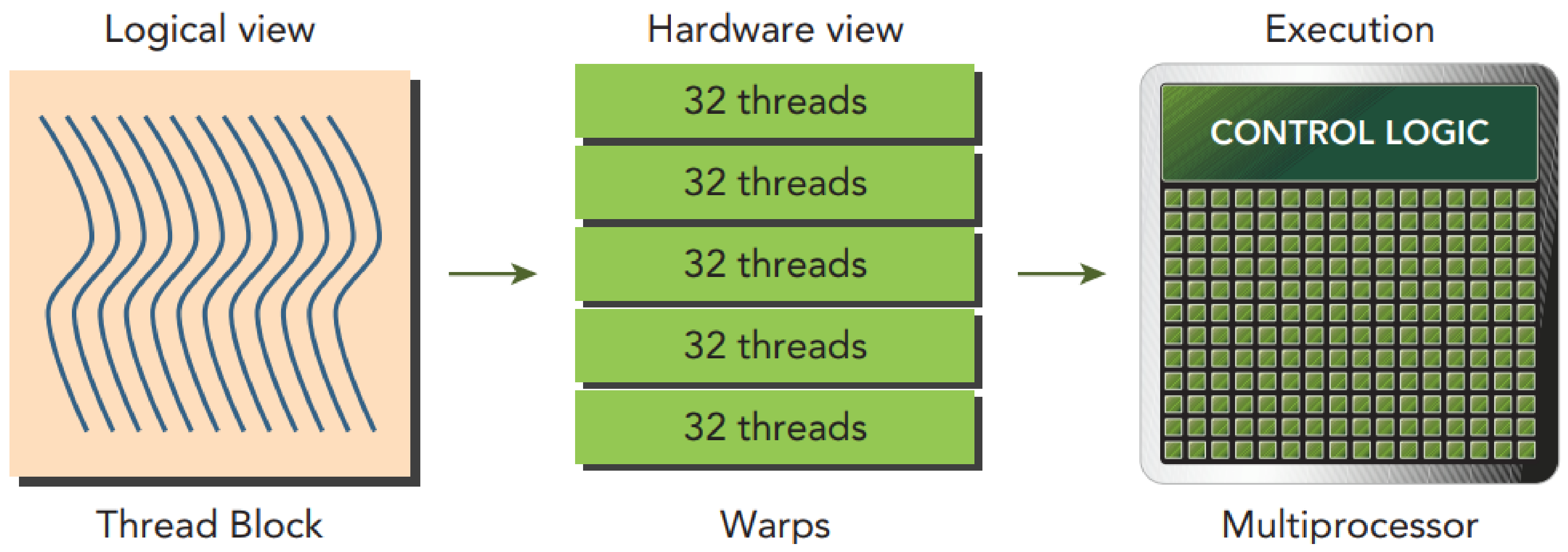
__syncthreads() it must be executed by all threads in a block

```
__global__ void incorrect_barrier_example(int n){
    //...
    if (threadIdx.x % 2 == 0){
        //...
        __syncthreads();
    }
    else{
        //...
        __syncthreads();
    }
}
```

- Not all threads in a block are guaranteed to execute either of the barriers ➔ undefined execution behavior

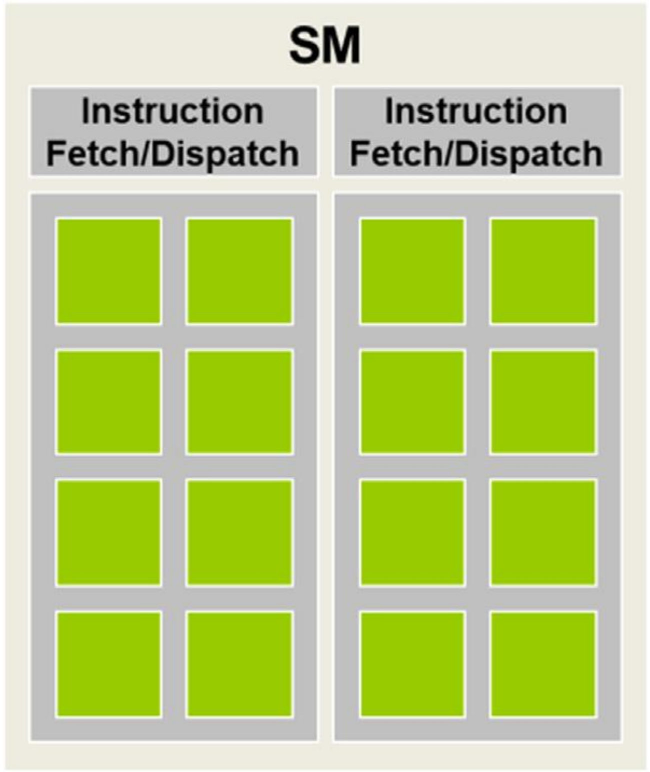
Parallel execution – Warp

- In SM, with each block, system don’t manage and execute each individual thread but a group of 32 threads - called a warp



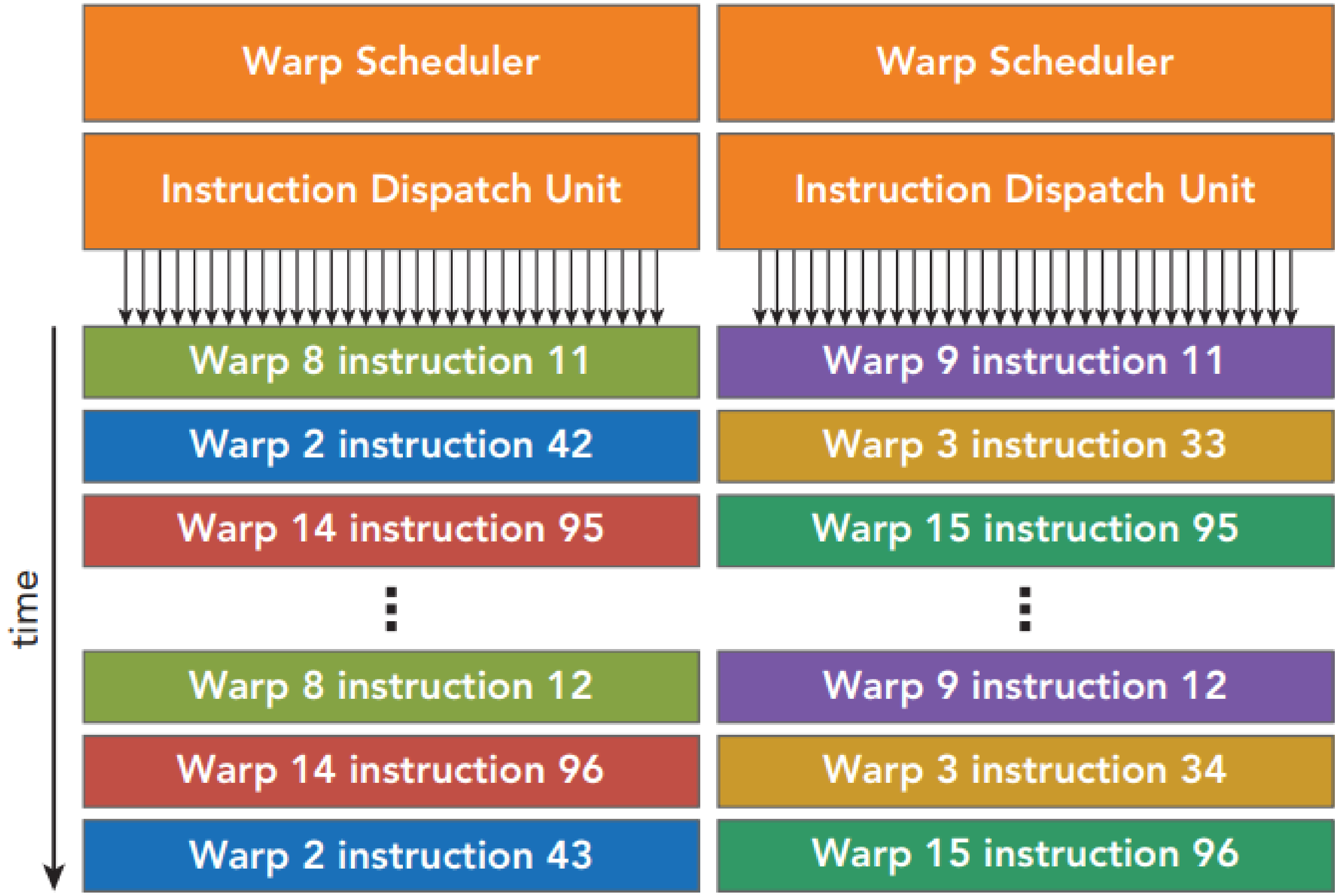
Parallel execution – Warp

- One instruction is executed for all threads in a warp (each thread has its own data
  - This execution model is called SIMT (Single Instruction Multiple Threads)
- The benefit of this execution model?
  - Help simplify hardware: less resources for control and more resources for arithmetic throughput



All threads in a warp execute the same instruction when selected

Parallel execution – Warp



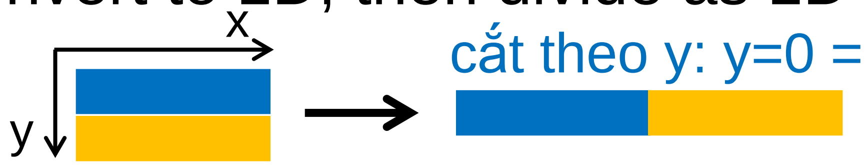
Parallel execution – Warp

How is a block in SM divided into warps?

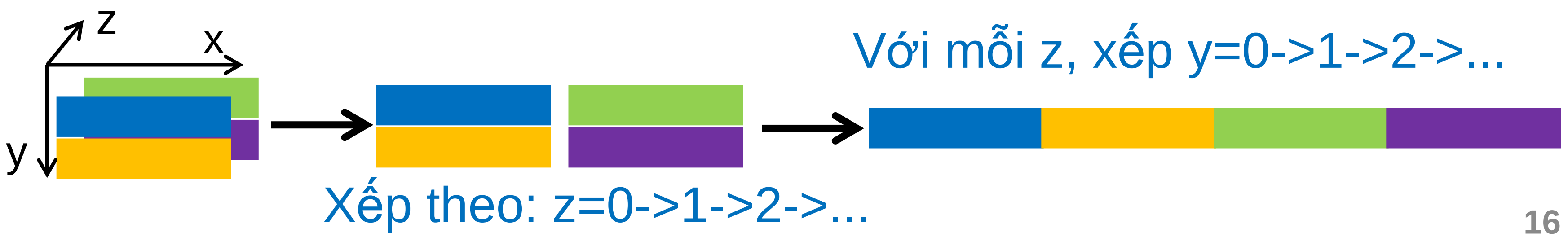
- **Block 1D:** 32 consecutive threads form 1 warp (warp 1: threads 0-31, warp 2: threads 32-63, …)

  If block size is not a multiple of 32 then the last warp will be padded extra threads so that its size is 32, these threads are useless but still consume resources

- **Block 2D:** convert to 1D, then divide as 1D



- **Block 3D:** convert to 1D, then divide as 1D



Parallel execution – SM-inside level

- What if threads in a warp cannot execute the same instruction? Hiện tượng rẽ nhánh do thread-Id

➔ Warp divergence

  - Correctness? OK
  - Speed? Hmm…

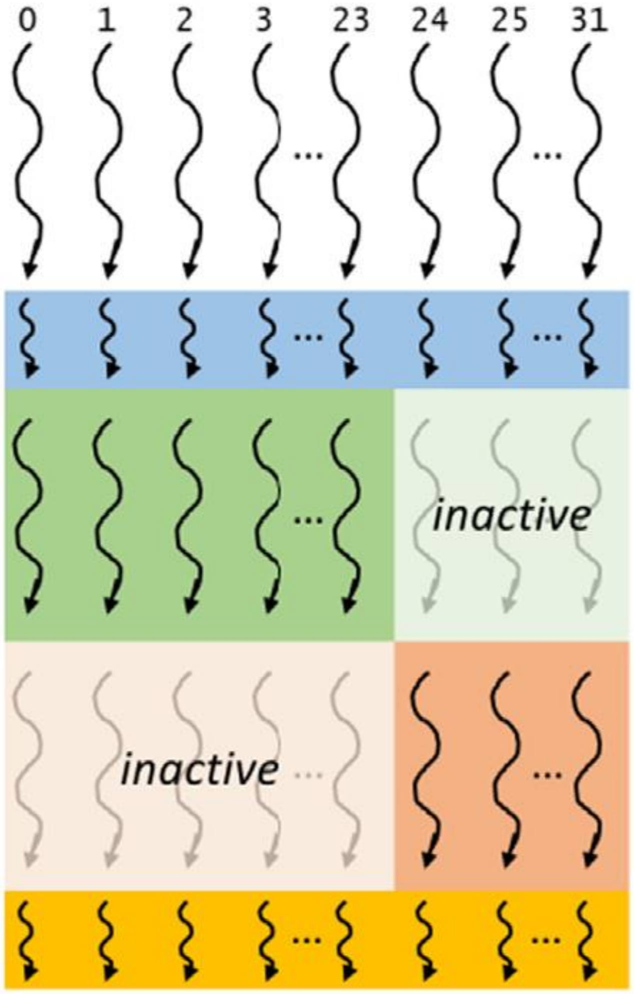
- GPUs use **predicated execution**
  - Each thread computes a yes/no answer for each path
  - Multiple paths taken by threads in a warp are executed serially!
- If all threads in a warp need a barrier synchronization mechanism: __syncwarp()

Parallel execution – SM-inside level

Warp divergence example: branching

ALL THREADS EXECUTE BOTH PATHS
(results kept only when predicate is true for thread)

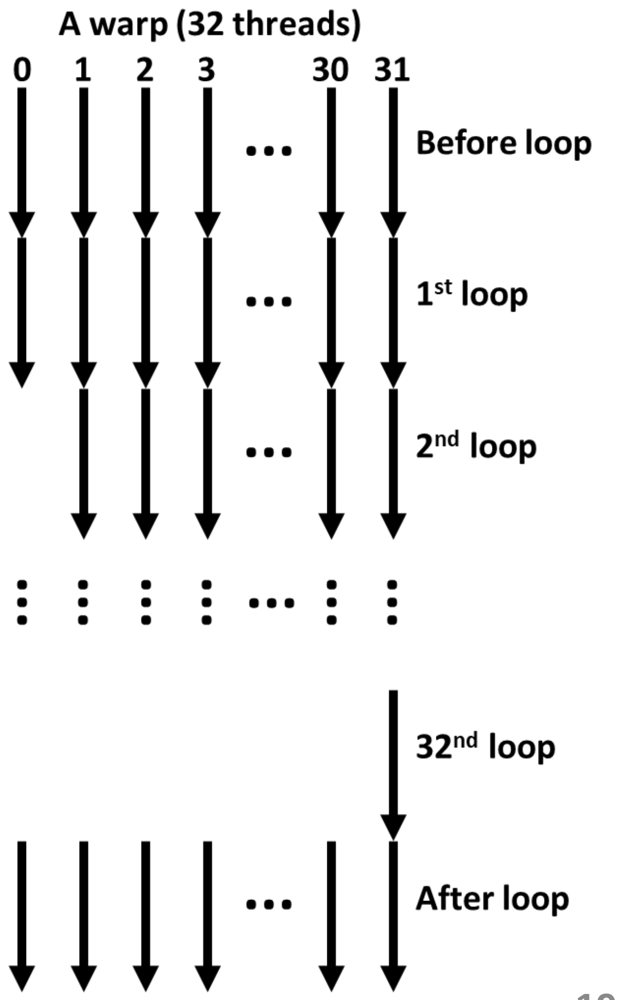
```
if(threadIdx.x < 24) {

    A

} else {

    B

}
C
```



Parallel execution – SM-inside level

Warp divergence example: looping

```
...
for (int i = 0; i <= threadIdx.x; i++)
{
    ...
}
...
```



Avoiding Branch Divergence

- Make branch granularity a multiple of warp size.

```
__global__ void divergence1(float* A, float* B, float* C, int n){
    int i = blockDim.x * blockIdx.x + threadIdx.x;
    if (i < n){
        if (i % 2)
            C[i] = A[i] * B[i];
        else
            C[i] = A[i] / B[i];
    }
}
```

```
__global__ void divergence2(float* A, float* B, float* C, int n)
{
    int i = blockDim.x * blockIdx.x + threadIdx.x;
    if (i < n) {
        if ( (i / 32) % 2 == 0)
            C[i] = A[i] * B[i];
        else
            C[i] = A[i] / B[i];
    }
}
```

Example: warp divergence

- Task: adding 2 matrixes Xem lại cách chia warp 2D -> 1D
  - Matrix size 1000×1000
  - Each thread computes an element in the result matrix
  - Block size 32×32
- How many diverged warps?

A. 0

[✓ B. 1000]

C. 1024

D. 2000

E. I don’t know

Example: warp divergence

Task: adding 2 matrixes

- Matrix size 1000×1000
- Each thread computes an element in the result matrix
- Block size 32×32 => 32 warp / block

Execution time of a diverged warp vs non-diverged warp? (we don't consider non-diverged warps in which all 32 threads fail the if condition and do nothing)

A. Faster

B. Slower

[✓ C. Equal]

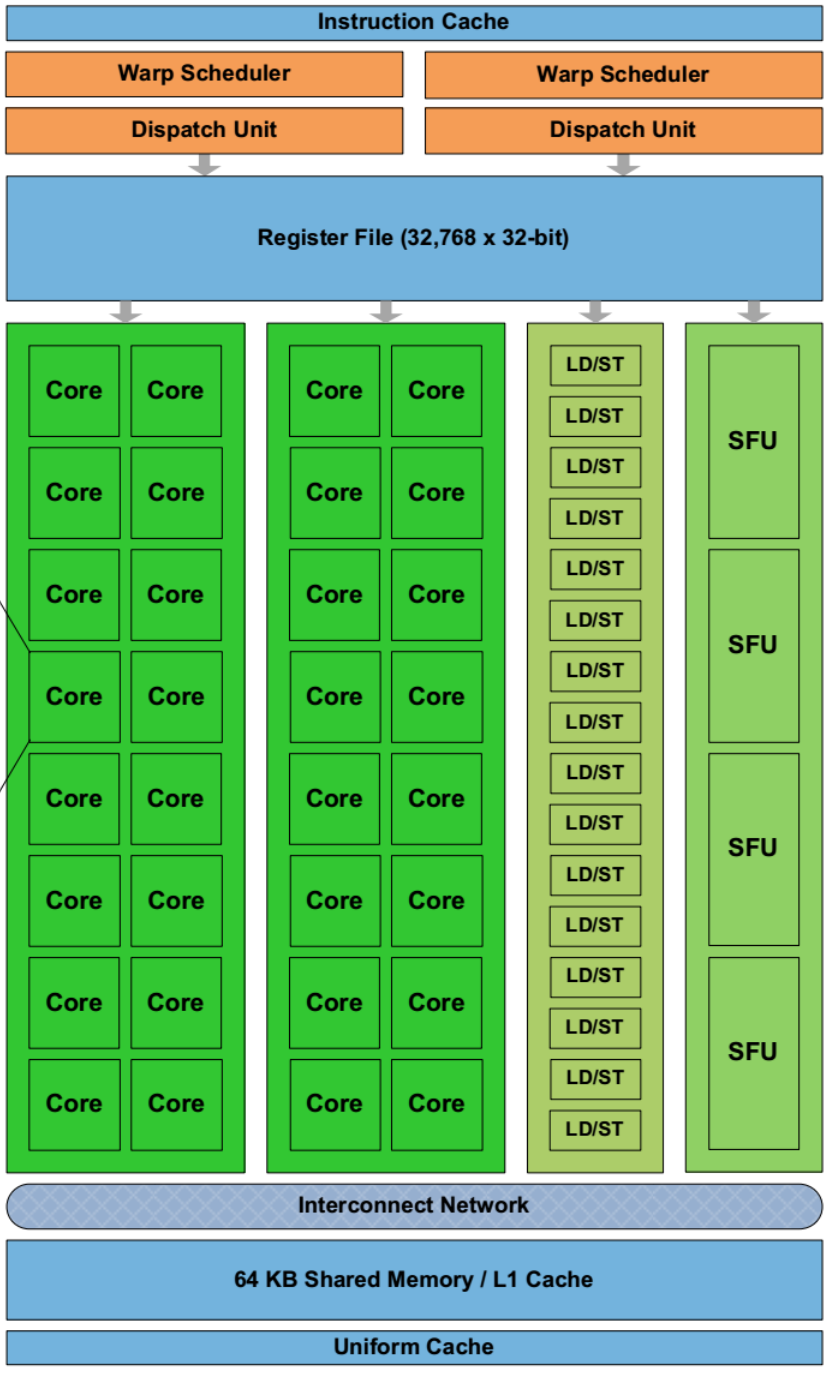
D. I don't know

Warp scheduling & latency tolerance

- In SM, are warps executed in parallel?
  - Not totally so

    E.g., Fermi SM (2.x) can contain at most 48 warps (1536 threads), but it has only 32 cores
- So, in SM how are warps executed exactly?
- Why does SM contain many warps / threads compared to its execution resources (cores)?

Image source: NVIDIA. Fermi white paper



Fermi Streaming Multiprocessor (SM)

Warp scheduling & latency tolerance

Before continuing, let's review instruction pipeline

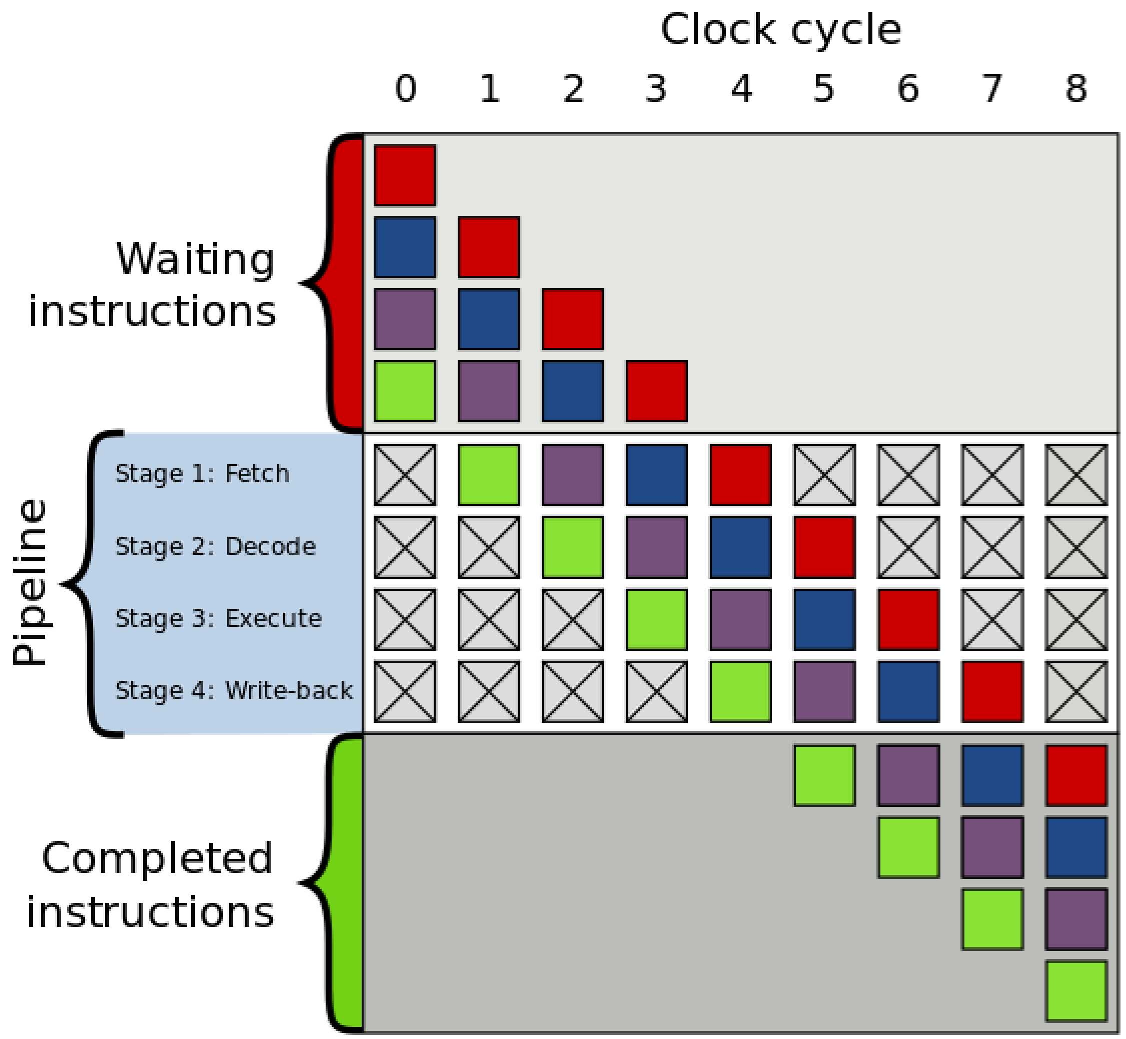
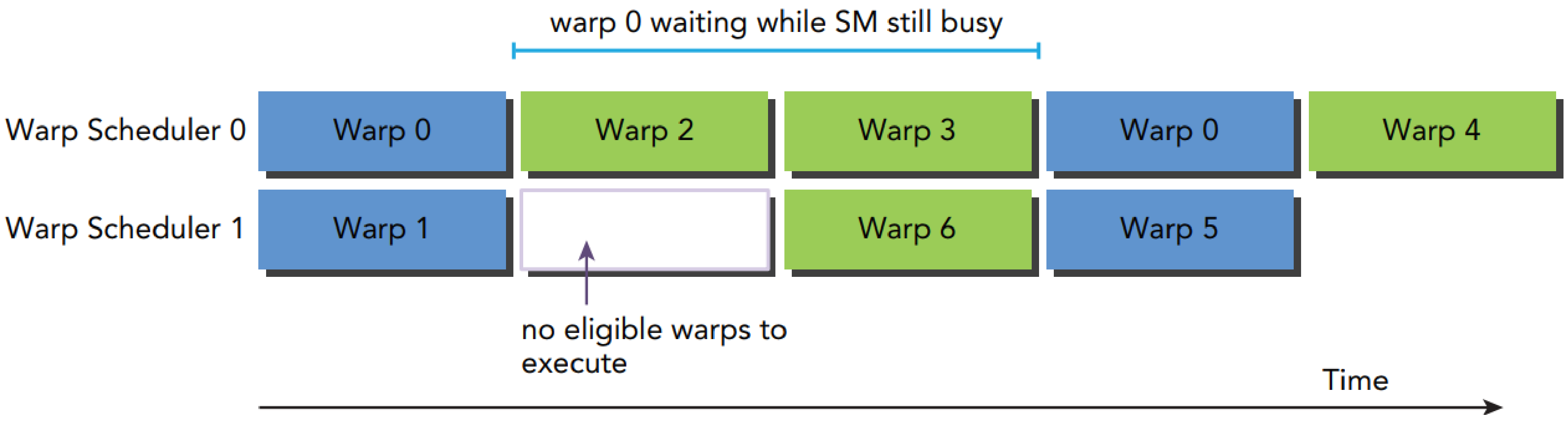


Image source: https://en.wikipedia.org/wiki/Instruction_pipelining

Warp scheduling & latency tolerance

- **Latency**: the number of clock cycles between an instruction being issued and being completed
- **Full compute resource utilization** is achieved when *all warp schedulers have an eligible warp at every clock cycle*.
- The latency of each instruction can be hidden by issuing other instructions in other resident warps



Warp scheduling & latency tolerance

- Two instruction types:
  - Arithmetic instructions: 10-20 cycles
  - Memory instructions: 400-800 cycles for global memory accesses
- If an instruction has latency of n clock cycles, then scheduler will need ~n ready instructions (coming from the same warp or other warps) to “hide” latency, keep pipelines full. (latency tolerance or latency hiding)

➔ desirable for an SM to have many more threads assigned to it than can be simultaneously supported with its execution resources to maximize the chance of finding a warp that is ready to execute at any point in time

Resource partitioning

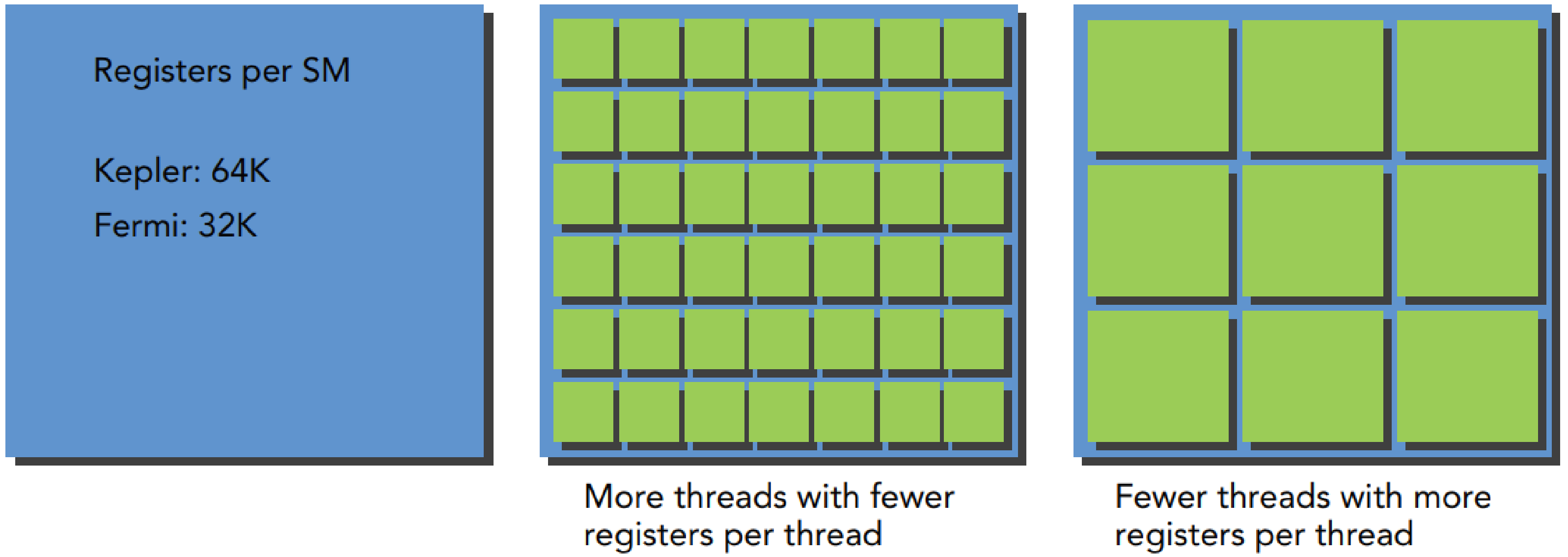
- For each SM:
  - The more warp (threads), the better (to hide latency)
  - Maximum: 2048 Threads/SM = 64 warps/SM (*)
  - Not always can reach this maximum due to resource availability

$$\text{Occupancy} = \frac{\text{\#warps SM contains}}{\text{\# max warps SM can contain}}$$

Resource partitioning

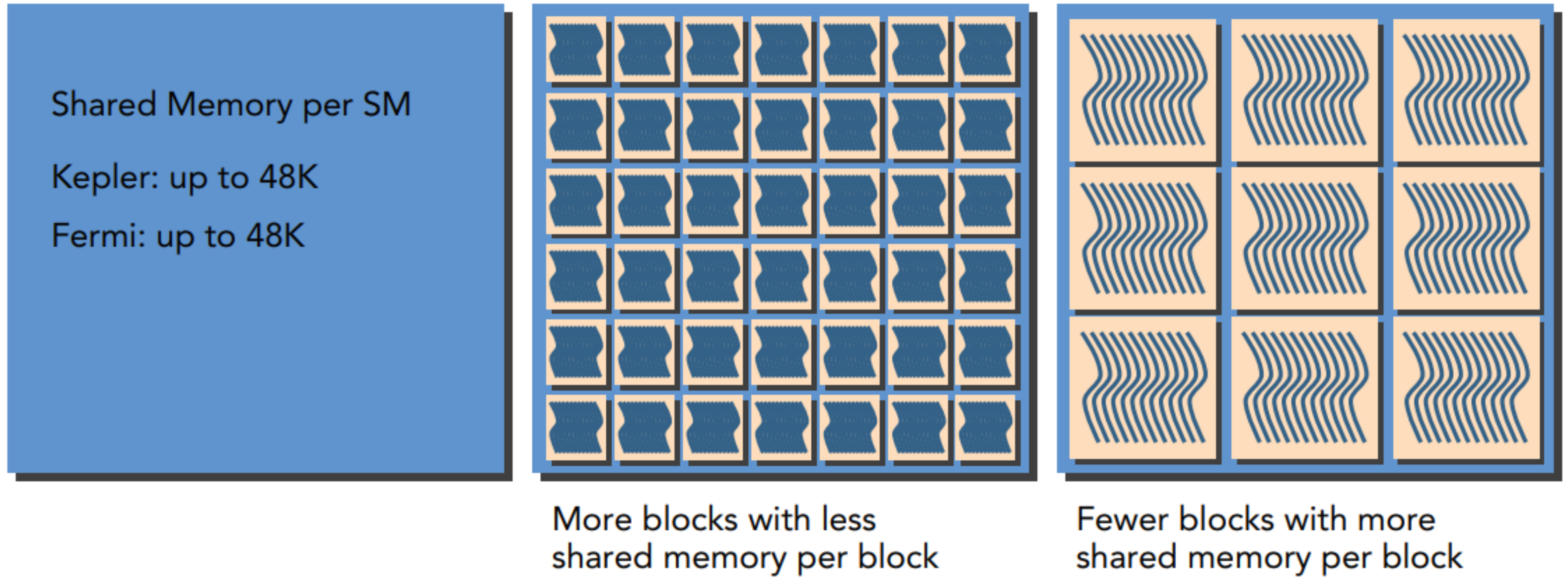
- The execution of a warp mainly consists of the following resources:
  - Program counters
  - Registers
  - Shared memory
- The number of blocks and warps that can simultaneously reside on an SM depends on the number of registers and amount of shared memory available on the SM and required by the kernel.

Resource partitioning



More threads with fewer registers per thread

Fewer threads with more registers per thread



More blocks with less shared memory per block

Fewer blocks with more shared memory per block

Resource partitioning - Example

- A100 GPU: 32 blocks/SM, 64 warps (2048 threads)/ SM, 1024 threads/block, and 65,536 registers/SM
- BlockSize = 32 ➔ ? Theads/SM ➔ Occupancy = ? 50%
  1024
- BlockSize = 768 ➔ ? Theads/SM ➔ Occupancy = ? 75%
  1536
- Kernel uses 64 registers per thread ➔ ? Theads/SM ➔ Occupancy = ? 50%
  1024
- 31 registers/thread, 512 threads/block ➔ ? 2048 Theads/SM ➔ Occupancy = ? 100%
- 33 registers/thread, 512 threads/block ➔ ? 1536 Theads/SM ➔ Occupancy = ? 75%

Resource partitioning - Example

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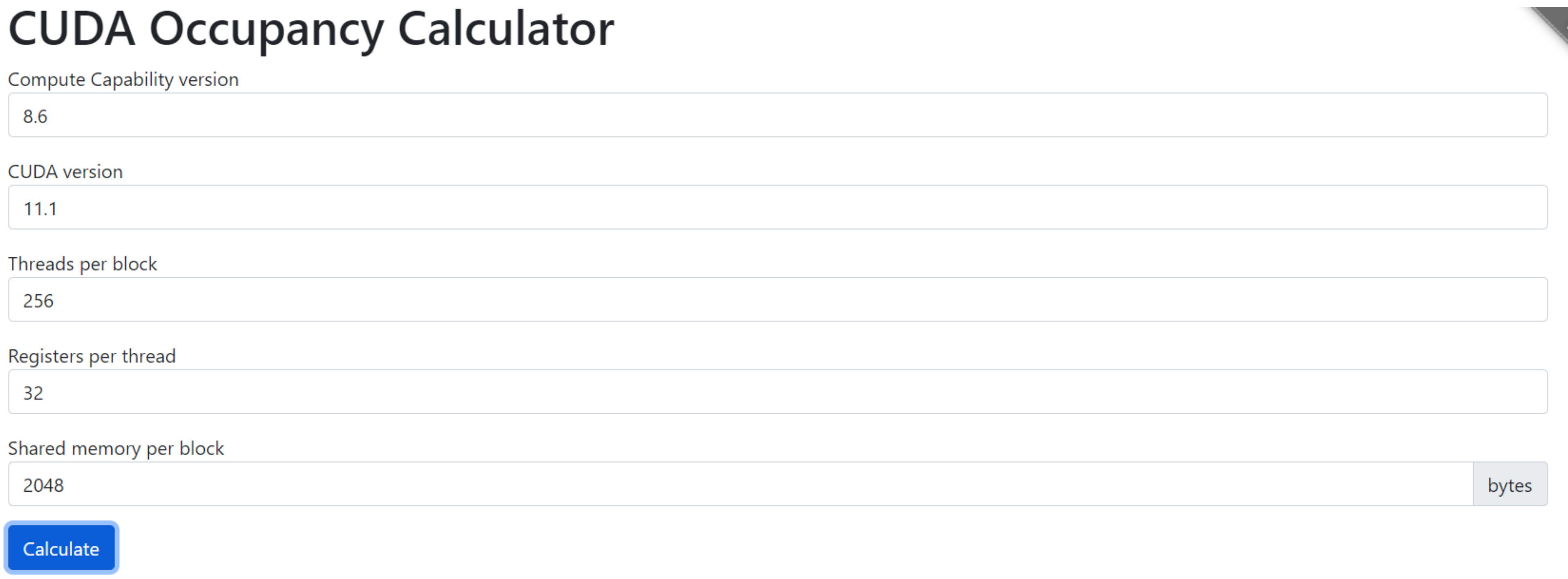
- Kernel uses 64 registers per thread ➔ 1024 Theads/SM ➔ Occupancy = 50%
- 31 registers/thread, 512 threads/block ➔ 2048 Theads/SM ➔ Occupancy = 100%
- 33 registers/thread, 512 threads/block ➔ $\overset{1536}{1538}$ Theads/SM ➔ Occupancy = 75%

Resource partitioning

- Accurate determination of the number of threads running in each SM can be difficult
- Tool: CUDA Occupancy Calculator



CUDA Occupancy Calculator

Compute Capability version	8.6
CUDA version	11.1
Threads per block	256
Registers per thread	32
Shared memory per block	2048 bytes

Example: # needed warps for Kepler SM

- Each SM has 4 warps schedulers

  cần ít nhất

  → need 4+ warps / SM

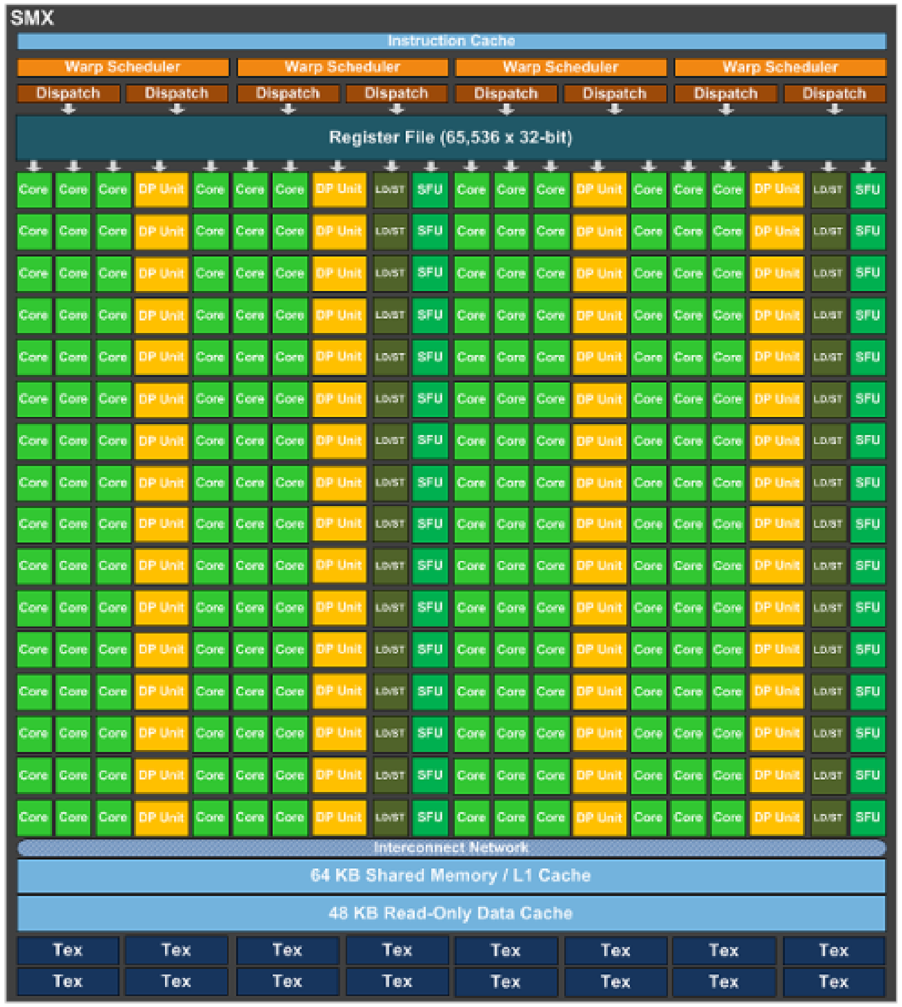
  In practice, need many more than 4 to hide latency (Kepler SM can contain up to 64 warps)

- With programs whose performance is limited by computation throughput (latency of a computation instruction: 10+ clock cycles)
  - Without ILP (Instruction Level Parallelism – adjacent independent instructions in a warp):

    need 4 schedulers × 10+ cycles

    = 40+ warps / SM
  - With ILP: may need fewer warps

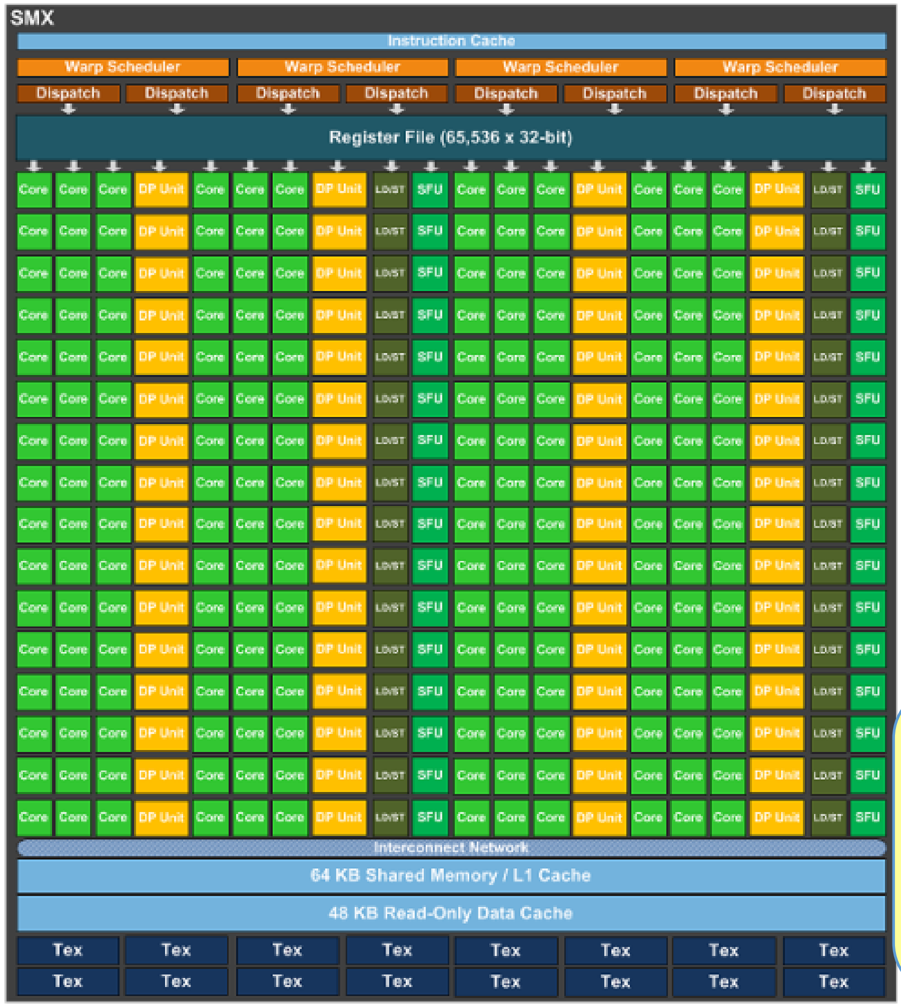
More about cores
Example: Kepler SM


SMX
Instruction Cache
Warp Scheduler
Dispatch
Register File (65,536 x 32-bit)
Core
DP Unit
LD/ST
SFU
Interconnect Network
64 KB Shared Memory / L1 Cache
48 KB Read-Only Data Cache
Tex

- **192 fp32 lanes (cores)**
  - fp32 math
  - Simple int32 math (add,min,etc.)
- **64 fp64 lanes**
- **32 SFU lanes**
  - Int32 multiplies, etc.
  - Transcendentals
- **32 LD/ST lanes**
  - GMEM, SMEM, LMEM accesses
- **16 TEX lanes**
  - Texture access
  - Read-only GMEM access

Image source: Paulius Micikevicius. Performance Optimization, GTC2013

More about cores
Example: Kepler SM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- **64 fp64 lanes**
- **32 SFU lanes**
  - Int32 multiplies, etc.
  - Transcendentals
- **32 LD/ST lanes**

NVIDIA "core" refers to fp32 core

# fp32 cores > # fp64 cores
→ should use 32-bit float when possible

Image source: Paulius Micikevicius. Performance Optimization, GTC2013

Guide for block size selection

bội số của 32

- Block size should be a multiple of 32 (warp size)
- Block size should be selected so that SM has enough warps to hide latency and make use of available resources
  - Occupancy measure: the ratio of # warps SM contains to # max warps SM can contain

    Block size should be selected so that occupancy is high

    Example: assume SM can contain up to 8 blocks and 1536 threads (48 warps); what block size should we pick: 64, 256, 1024 (256 => O=100%)
  - Not necessary: 100% occupancy = max performance
    - Only need enough warp to hide latency and make use of available resources
    - If warps have ILP, we will need fewer warps
    - …

Reference

- [1] Wen-Mei, W. Hwu, David B. Kirk, and Izzat El Hajj. *Programming Massively Parallel Processors: A Hands-on Approach*. Morgan Kaufmann, 2022
- [2] Cheng John, Max Grossman, and Ty McKercher. *Professional Cuda C Programming*. John Wiley & Sons, 2014

THE END