Performance and Optimization

CSC 230 : C and Software Tools

NC State Department of Computer

Science

Topics for Today

- Leftover Topic: Environment Variables
- Code coverage
- Optimization Techniques
 - Timing program execution
 - Profiling
 - Optimizing by Hand
 - Compiler support for optimization

Environment Variables

- Your execution environment supports variables
 - We can access these from the shell

```
$ echo $HOME
/home/jerry
$ HOME=/home/linda
$ echo $HOME
/home/linda
```

- The shell uses these a lot for configuration and to customize behavior
 - HOME, PATH, SHELL, PS1, EDITOR, ...
- But, these aren't just available to the shell, all programs have environment variables.

Using the Environment

- The standard library lets us get/set these variables
 - We need our good friend, stdlib.h
 - Then, we can gen environment variable values

```
char *str = getenv( "HOME" );
printf( "HOME = $s\n", str );
```

- And, we can change them

 These changes are local to the program, they go away when it exits.

Testing Testing

- How do you know if your tests are adequate?
 - Well, there are different types of tests you can perform
 - Unit testing, integration testing, system testing, regression testing, acceptance testing, beta testing
 - black-box vs white box
 - Each type intended to expose a different cause or type of bug.
- For acceptance testing, we could make sure we have a test or two for every (functional) requirement
- For black box testing, we can test output on typical and boundary cases
 - To make sure we cover the range of possible inputs

Code Coverage

- For white-box testing, we can make sure we cover the code itself.
- Ideally, all the source code should get a chance to run under test
 - Every statement executed
 - Every branch direction of every conditional
 - Every sub-expression of every compound conditional

These are details, but they could hide errors.

Tools for Code Coverage

- The compiler will help us measure this
- Compile with the -fprofile-arcs and -ftestcoverage flags

```
gcc -Wall -std=c99 -fprofile-arcs \
  -ftest-coverage bubble1.c -o bubble1
```

Run your program all you like

```
./bubble1 < small_input
./bubble1 < some_other_input
```

Using gcov

• Notice, you get some extra files.

```
eos$ ls
bubble1 bubble1.c
bubble1.gcno bubble1.gcda

written during
execution

written by the
compiler
```

Run gcov to extract the coverage information

```
eos$ gcov bubble1.c
File 'bubble1.c'
Lines executed:66.67% of 36
bubble1.c:creating 'bubble1.c.gcov'
```

Reading gcov Output

 In addition to the summary, gcov writes a text file giving coverage details.

```
42:
             for ( int i=0; !done && i<len; i++ ) {
   -: 43:
               // If we make it through without ...
   1: 44: done = true;
   -: 45:
   5: 46:
               for ( int j=0; j<len-i-1; j++ )
                 if ( list[j ] > list[j+1] ) {
   4: 47:
                   swap(list + j, list + j + 1);
#####: 48:
#####: 49:
                   done = false;
   -: 50:
   -: 51:
   1:
        52:}
                       number of times executed ... line number
```

Reading gcov Output

gcov has some codes for reporting execution count

5	I was executed 5 times
-	I could never be executed (this doesn't count against coverage)
#####	I was never executed (but I could have been)

 Let's see if we can get code coverage up to 100 %

Getting More from gcov

- By default, gcov tells you about statements executed
- It can also tell you how branches went

```
eos$ gcov -b bubble1.c
File 'bubble1.c'
Lines executed:66.67% of 36
Branches executed:100.00% of 18
Taken at least once:77.78% of 18
Calls executed:61.54% of 13
bubble1.c:creating 'bubble1.c.gcov'
```

Reading gcov Output

 With the -b option, gcov reports branching directions for each part of each conditional.

```
1: 40: bool done = false;
-: 41:
2: 42: for ( int i=0; !done && i<len; i++ ) {
branch 0 taken 50% (fallthrough) Which way did !done
branch 1 taken 50%
branch 2 taken 100%
branch 3 taken 0% (fallthrough)
-: 43: // If we make it through ...
1: 44: done = true;
```

Performance and Optimization

- Performance depends on
 - Coding Style
 - Choice of programming language
 - Compiler and Options

-Choice of Algorithm and Data Structure

Code Profiling

 We can measure our program's performance using the time command.

```
time ./program program-options
```

- This reports how long it took the program to execute, CPU time consumed, etc.
- It's fairly coarse grained, we'd like to know more about where our program is spending its time.
- Code profiling tools let us do that.

Code Profiling

- gprof is a tool for profiling during execution
 - It counts the number of times each function is called
 - How much time we spend in each function
- To use gprof, we need to compile with the -pg option:

```
gcc -pg -Wall -std=c99 program.c -o program
```

Generating Profile Information

Then, you can run your program as normal.

```
./program program-options
```

This also generates some profiling output

```
$ ls
gmon.out myProgram myProgram.c
```

Where did this file come from?

Examining Profile Output

- The gmon.out file is in a binary format.
- But, the gprof command will print it in a way that's human readable.

```
$ gprof ./myProgram
<lots of output>
```

 Maybe it would be better to send this output to a file.

```
$ gprof ./myProgram > perf.txt
```

Fun With Time and Gprof

- Let's run bubble4.c against the big_input.txt input file.
 - This file is large enough to take a few seconds to bubble sort.
- First, let's time the execution
- Then, let's profile:
 - Compile with profiling support
 - Run to collect profiling results
 - Examine results with gprof

DIY Optimizations: Code Motion

- You can optimize code yourself
 - By moving duplicate expression evaluation outside loops

```
for ( int i = 0; i < n; i++ )
  for ( int j = 0; j < n; j++ )
    a[ n * i + j ] = f() + g(i) + h(j);</pre>
```



```
int c1 = f();
for ( int i = 0; i < n; i++ ){
  int c2 = c1 + g(i);
  int ri = n * i;
  for ( int j = 0; j < n; j++ )
    a[ ri + j ] = c2 + h(j);
}</pre>
```

DIY Optimizations: Strength Reduction

```
c1 = f();
for ( int i = 0; i < n; i++ ){
  c2 = c1 + g(i);
  int ri = n * i;
  for ( int j = 0; j < n; j++ )
    a[ ri + j ] = c2 + h(j);
}</pre>
```

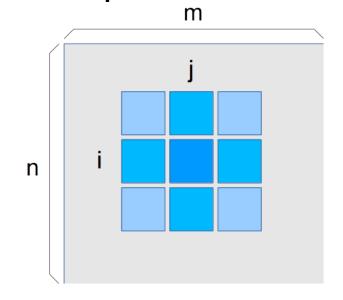


```
c1 = f();
*p = a;
for ( int i = 0; i < n; i++ ){
  c2 = c1 + g(i);
  for ( int j = 0; j < n; j++ )
    *p++ = c2 + h(j);
}</pre>
```

DIY Optimization: Expression Reuse

Consider this neighborhood computation:

```
sum = a[ (i-1)*m + j ];
sum += a[ (i+1)*m + j ];
sum += a[ i*m + j-1 ];
sum += a[ i*m + j+1 ];
```





```
int ctr = i*m + j;
sum = a[ ctr - m ];
sum += a[ ctr + m ];
sum += a[ ctr - 1 ];
sum += a[ ctr + 1 ];
```

DIY Optimization: Inlining

```
double prod( int i, int j )
{
  double sum = 0.0;
  for ( int k = 0; k<n; k++)
     sum += a[i][k] * b[k][j];
  return sum;
}</pre>
```

```
for (i=0; i<n; i++)
  for (j=0; j<n; j++)
  c[i][j] = prod(i, j);</pre>
```

DIY Optimization: Inlining

```
double prod( int i, int j )
 double sum = 0.0;
  for ( int k = 0; k < n; k++)
    sum += a[i][k] * b[k][j];
  return sum;
for (i=0; i<n; i++)
  for (j=0; j< n; j++) {
    double sum = 0.0;
    for (k=0; k< n; k++)
       sum += a[i][k] * b[k][j];
    c[i][j] = sum;
```

Inlining Help from the Language

- C can inline your functions for you
 - Introduced in C99

```
inline double prod( int i, int j )
{
  double sum = 0.0;
  for ( int k = 0; k<n; k++)
     sum += a[i][k] * b[k][j];
  return sum;
}</pre>
```

```
for (i=0; i<n; i++)
  for (j=0; j<n; j++)
  c[i][j] = prod(i, j);</pre>
```

Help from the Compiler

- Many of these refinements can be done automatically by the compiler
- It can try to optimize for execution speed
 - Options -O, -O2 and -O3, in order of increasing optimization.
- Increasing optimization level can:
 - Increase code size
 - Possibly decreasing performance
- But, the compiler can't do everything

DIY Optimization: Reordering Tests

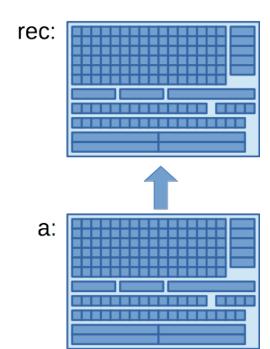
Place more frequent branches earlier

DIY Optimization: Pass by Reference

- C lets us use value semantics with most types
 - ... maybe it's too easy.
 - Passing large objects by value can be expensive.

```
int f( struct BigRecord rec )
{
  int sum = rec.x + rec.y;
  return sum;
}
```

```
struct BigRecord a;
...
v = f( a );
```



DIY Optimization: Pass by Reference

- Pass by reference/address can reduce overhead copying large objects
- Const lets the function say what it plans to do with the object.

```
int f( const struct BigRecord *rec )
{
  int sum = rec->x + rec->y;
  return sum;
}
```

```
struct BigRecord a;
...
v = f( &a );
```

