

General

→ setxkbmap ch -variant de_nodeadkeys

make

CXX = g++

CFLAGS = -Wall -O3 -std=c++11

test: test.cpp

\$(CXX) \$(CFLAGS) -o \$@ \$*

clean:

rm -f test

#PCSEI HS16

Threads

compile

g++ -std=c++11 -pthread

import

#include <thread>

launch func

std::thread t(foo, arg1);

launch λ

std::thread (λ) { // do something };

mutex:

import

#include <mutex>

declare

std::mutex mtx;

lock

mtx.lock();

⋮

mtx.unlock();

lock_guard

std::lock_guard<std::mutex> l(mtx);

→ locked when initialized

→ unlocked when destroyed

unique_lock

std::unique_lock<std::mutex> l(mtx, &id);

→ locks work with multiple locks (is dead lock)

l.unlock();, l.lock(); for manual

Example

int nthreads = 4;

int nstep = N/nthreads;

std::vector<std::thread> threads(nthreads);

for (int t = 0; t < nthreads; t++) {

threads[t] = thread([t, t] {

for (int i = t*nstep; i < (t+1)*nstep; i++) {

z[i] = x[i] + y[i];

};

}

for (std::thread& t : threads) {

t.join();

}

MPI

compile

mpicc, mpic++

import

#include <mpi.h>

run

mpirun -np 4 ./a.out

info

MPI_Comm_rank (MPI_COMM_WORLD, &rank);

MPI_Comm_size (MPI_COMM_WORLD, &size);

Example

int main (int argc, char **argv) {

MPI_Init (&argc, &argv);

int rank;

MPI_Comm_rank (MPI_COMM_WORLD, &rank);

if (rank == 0) { // master

MPI_Status status;

char txt[100];

MPI_Recv (txt, 100, MPI_CHAR, 1, 42, &status);

MPI_COMM_WORLD, &status);

std::cout << txt << '\n';

} else { // worker

std::string text = "Bla";

MPI_Send (const_cast<char*> (text.c_str()),

text.size()+1, MPI_CHAR, 0, 42,

MPI_COMM_WORLD);

} MPI_Finalize();

}

Open MP

compile

g++ -fopenmp -std=c++11

import

#include <omp.h>

info

omp_get_thread_num() → id

omp_get_num_threads() → #threads

#PCSEI #S16

F. Vogel

Parallel Scaling
Andals Law:

$$S(n) = \frac{1}{(1-p) + \frac{p}{n}}$$

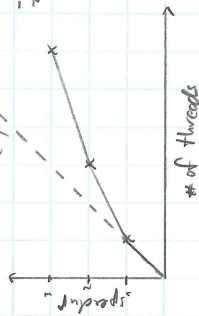
Gustafson's Law:

$$S(n) = 1 - p + n \cdot p \quad (\text{weak scaling})$$

Strong scaling:

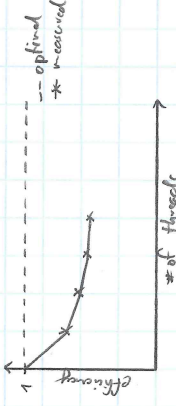
$$S(n) = \frac{T(1)}{T(n)} \quad E(n) = \frac{S(n)}{n}$$

-- optimal
 * measured



Weak scaling:

$$E(n) = \frac{T(1)}{T(n)} \quad \text{fixed problem size per thread}$$



Utk: Problem size for N^2 solver is $N^2 \log N$

Roofline Model

Operational Intensity [Flop/Byte]

$$r = \frac{\text{\# of Flops}}{\text{\# of Bytes}}$$

Example: $RHS_{i,j} = C_n \cdot (\hat{q}_{i,j} + \hat{q}_{i,j-1}^* + \hat{q}_{i,j+1}^* + \hat{q}_{i,j+2}^* - 4\hat{q}_{i,j})$
 $\rightarrow 6 \text{ Flops (4 ADD + 2 MUL)}$

\rightarrow Mem. Access:

\rightarrow No cache: 5 read + 1 write = 6

\rightarrow Infinite cache: 1 read + 1 write = 2

\rightarrow Op. Int.

\rightarrow No cache: $\frac{6.48}{6 \text{ Flops}} = 0.25 \text{ Flop/B}$

\rightarrow Infinite cache: $\frac{6 \text{ Flops}}{2.4 \text{ B}} = 0.75 \text{ Flop/B}$

Non-ideal performance

\rightarrow peak = processor clock/sec \cdot vector size \cdot instructions/clock \cdot # cores

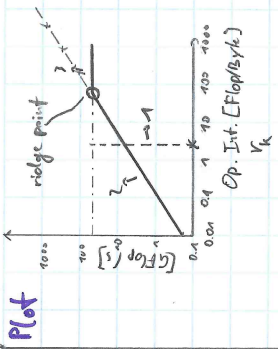
\rightarrow break = memory clock/sec \cdot channel size \cdot # channels / [GB/s/ops]

Bugs:

- \rightarrow Threads: Pos: thread id by value, not by reference
- \rightarrow SIMD: Race condition by interleaving of loads
- \rightarrow SIMD: Undefined behavior by loading not allocated memory
- \rightarrow OMP: implicit barrier @ end of for, all threads have to call it. Fix: schedule(dynamic) nowait
- \rightarrow MPI: Don't update asynchronous send buffer
- \rightarrow Threads: Not passing allocated variables
- \rightarrow Threads: Race condition by not locking result double
- \rightarrow Threads: Result reduction before join is called in each thread
- \rightarrow Threads: Master thread: not releasing lock before calling join can result in deadlock
- \rightarrow MPI: Ordering needed to prevent deadlock (replacing with asynchronous works too)

- 1: Locality
- 2: Communication
- 3: Computation

limit by comp --- $f_{peak} = f_{lim}$
 limit by OMP --- $r_k \cdot break = f_{lim}$



Plot