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## 1 计算几何

## 1.1 二维计算几何基本操作

```
1 const double PI = 3.14159265358979323846264338327950288;
 2 double arcSin(const double &a) {
        return a \le -1.0? -PI / 2: (a >= 1.0? PI / 2: asin(a)):
 4 }
 5 double arcCos(const double &a) {
       return a \le -1.0? PI : (a \ge 1.0 ? 0 : a\cos(a);
7 }
 8 struct point {
        double x, y; // `something omitted'
        point rot(const double &a) const { // `counter-clockwise`
           return point(x * cos(a) - y * sin(a), x * sin(a) + y * cos(a));
12
13
       point rot90() const { // `counter-clockwise`
           return point(-y, x);
14
15
       point project(const point &p1, const point &p2) const {
16
17
           const point &q = *this;
           return p1 + (p2 - p1) * (dot(p2 - p1, q - p1) / (p2 - p1).norm());
18
19
       bool onSeg(const point &a, const point &b) const { // `a, b inclusive`
20
           const point &c = *this:
21
22
           return sign(dot(a - c, b - c)) \leq 0 \& sign(det(b - a, c - a)) == 0;
23
       double distLP(const point &p1, const point &p2) const { // `dist from *this to line p1->p2`
24
           const point &q = *this;
25
           return fabs(det(p2 - p1, q - p1)) / (p2 - p1).len();
26
27
       double distSP(const point &p1, const point &p2) const { // `dist from *this to segment [p1, p2]`
28
           const point &q = *this;
29
30
           if (dot(p2 - p1, q - p1) < EPS) return (q - p1).len();
           if (dot(p1-p2, q-p2) < EPS) return (q-p2).len();
31
32
           return distLP(p1, p2);
33
34
       bool inAngle(const point &p1, const point &p2) const { // `det(p1, p2) $\quad qe$ 0`
35
           const point &q = *this; return det(p1, q) > -EPS && det(p2, q) < EPS;</pre>
36
38 bool lineIntersect(const point &a, const point &b, const point &c, const point &d, point &e) {
       double s1 = det(c - a, d - a);
       double s2 = det(d - b, c - b);
       if (!sign(s1 + s2)) return false;
       e = (b - a) * (s1 / (s1 + s2)) + a;
43
44 }
45 int segIntersectCheck(const point &a, const point &b, const point &c, const point &d, point &o) {
       static double s1, s2, s3, s4;
       static int iCnt;
       int d1 = sign(s1 = det(b - a, c - a));
       int d2 = sign(s2 = det(b - a, d - a));
       int d3 = sign(s3 = det(d - c, a - c));
```

```
int d4 = sign(s4 = det(d - c, b - c));
        if ((d1 \wedge d2) == -2 &\& (d3 \wedge d4) == -2) {
            o = (c * s2 - d * s1) / (s2 - s1);
            return true;
54
 55
        iCnt = 0:
        if (d1 == 0 \&\& c.onSeg(a, b)) o = c, ++iCnt;
        if (d2 == 0 \&\& d.onSeg(a, b)) o = d, ++iCnt;
        if (d3 == 0 \&\& a.onSeq(c, d)) o = a, ++iCnt;
        if (d4 == 0 \&\& b.onSeg(c, d)) o = b, ++iCnt;
        return iCnt ? 2 : 0; // `不相交返回0, 严格相交返回1, 非严格相交返回2
61
62 }
63 struct circle {
64
        point o;
        double r, rSqure;
        bool inside(const point &a) { // 洋严格
67
            return (a - o).len() < r + EPS;
68
        bool contain(const circle &b) const { // 非严格
            return sign(b.r + (o - b.o).len() - r) \le 0;
 70
71
        bool disjunct(const circle &b) const { // 注严格
 72
 73
            return sign(b.r + r - (o - b.o).len()) \le 0;
74
        int isCL(const point &p1, const point &p2, point &a, point &b) const {
 75
            double x = dot(p1 - o, p2 - p1), y = (p2 - p1).norm();
 76
            double d = x * x - y * ((p1 - o).norm() - rSqure);
 77
           if (d < -EPS) return 0;
 78
           if (d < 0) d = 0;
 80
            point q1 = p1 - (p2 - p1) * (x / y);
            point q2 = (p2 - p1) * (sqrt(d) / y);
 81
 82
            a = q1 - q2; b = q1 + q2;
            return q2.len() < EPS ? 1 : 2;
 83
84
        int tanCP(const point &p, point &p, point &b) const { // `返回切点, 注意可能与 $p$ 重合
 85
            double x = (p - o).norm(), d = x - rSqure;
 86
           if (d < —EPS) return 0;
 87
 88
           if (d < 0) d = 0;
            point q1 = (p - o) * (rSqure / x);
 89
            point q2 = ((p - o) * (-r * sqrt(d) / x)).rot90();
 91
            a = o + (q1 - q2); b = o + (q1 + q2);
            return q2.len() < EPS ? 1 : 2;
 92
 93
      }
 95 bool checkCrossCS(const circle &cir, const point &p1, const point &p2) { // 羊严格
        const point &c = cir.o;
        const double &r = cir.r;
        return c.distSP(p1, p2) < r + EPS
           && (r < (c - p1).len() + EPS | | r < (c - p2).len() + EPS);
 99
100 }
101 bool checkCrossCC(const circle &cir1, const circle &cir2) { // 非严格`
        const double &r1 = cir1.r, &r2 = cir2.r, d = (cir1.o - cir2.o).len();
        return d < r1 + r2 + EPS && fabs(r1 - r2) < d + EPS;
```

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```
104 }
int isCC(const circle &cir1, const circle &cir2, point &a, point &b) {
        const point &c1 = cir1.o, &c2 = cir2.o;
106
        double x = (c1 - c2).norm(), y = ((cir1.rSqure - cir2.rSqure) / x + 1) / 2;
        double d = cir1.rSqure / x - y * y;
108
        if (d < -EPS) return 0:
110
        if (d < 0) d = 0;
        point q1 = c1 + (c2 - c1) * y;
        point q2 = ((c2 - c1) * sqrt(d)).rot90();
112
113
        a = q1 - q2; b = q1 + q2;
114
        return q2.len() < EPS ? 1 : 2;
115 }
116 vector<pair<point, point> > tanCC(const circle &cir1, const circle &cir2) {
         、注意: 如果只有三条切线, 即 $s1 = 1, s2 = 1$, 返回的切线可能重复, 切点没有问题
117 /
        vector<pair<point, point> > list;
118
        if (cir1.contain(cir2) || cir2.contain(cir1)) return list;
119
        const point &c1 = cir1.o, &c2 = cir2.o;
120
121
        double r1 = cir1.r, r2 = cir2.r;
        point p, a1, b1, a2, b2;
        int s1, s2;
123
        if (sign(r1 - r2) == 0) {
124
            p = c2 - c1:
125
126
            p = (p * (r1 / p.len())).rot90();
            list.push_back(make_pair(c1 + p, c2 + p));
            list.push_back(make_pair(c1 - p, c2 - p));
128
        } else {
            p = (c2 * r1 - c1 * r2) / (r1 - r2);
130
131
            s1 = cir1.tanCP(p, a1, b1);
            s2 = cir2.tanCP(p, a2, b2);
132
            if (s1 >= 1 \&\& s2 >= 1) {
134
                list.push_back(make_pair(a1, a2));
135
                list.push_back(make_pair(b1, b2));
            }
136
137
        p = (c1 * r2 + c2 * r1) / (r1 + r2):
138
139
        s1 = cir1.tanCP(p, a1, b1);
        s2 = cir2.tanCP(p, a2, b2);
140
        if (s1 >= 1 \&\& s2 >= 1) {
141
            list.push_back(make_pair(a1, a2));
142
            list.push_back(make_pair(b1, b2));
143
144
145
        return list;
146 }
147 bool distConvexPIn(const point &p1, const point &p2, const point &p3, const point &p4, const point &q) {
        point o12 = (p1 - p2).rot90(), o23 = (p2 - p3).rot90(), o34 = (p3 - p4).rot90();
        return (q - p1).inAngle(o12, o23) || (q - p3).inAngle(o23, o34)
149
            II ((q - p2).inAngle(o23, p3 - p2) && (q - p3).inAngle(p2 - p3, o23));
150
151 }
152 double distConvexP(int n, point ps[], const point &q) { // 外部点到多边形的距离
        int left = 0, right = n;
154
        while (right - left > 1) {
            int mid = (left + right) / 2;
155
156
            if (distConvexPIn(ps[(left + n - 1) % n], ps[left], ps[mid], ps[(mid + 1) % n], q))
```

```
157
                right = mid;
158
            else left = mid;
        return q.distSP(ps[left], ps[right % n]);
160
161
162 double areaCT(const circle &cir, point pa, point pb) {
        pa = pa - cir.o; pb = pb - cir.o;
        double R = cir.r;
164
165
        if (pa.len() < pb.len()) swap(pa, pb);</pre>
166
        if (pb.len() < EPS) return 0;
        point pc = pb - pa;
        double a = pa.len(), b = pb.len(), c = pc.len(), S, h, theta;
168
169
        double cosB = dot(pb, pc) / b / c, B = acos(cosB):
        double cosC = dot(pa, pb) / a / b, C = acos(cosC);
        if (b > R) {
           S = C * 0.5 * R * R;
172
173
           h = b * a * sin(C) / c;
174
            if (h < R && B < PI * 0.5)
               S = a\cos(h / R) * R * R - h * sqrt(R * R - h * h);
        } else if (a > R) {
176
            theta = PI - B - asin(sin(B) / R * b);
            S = 0.5 * b * R * sin(theta) + (C - theta) * 0.5 * R * R:
178
179
        else S = 0.5 * sin(C) * b * a;
        return S:
180
181
182 circle minCircle(const point &a, const point &b) {
        return circle((a + b) * 0.5, (b - a).len() * 0.5);
184 }
185 circle minCircle(const point &a, const point &b, const point &c) { // `钝角三角形没有被考虑
        double a2( (b-c).norm() ), b2( (a-c).norm() ), c2( (a-b).norm() );
        if (b2 + c2 <= a2 + EPS) return minCircle(b, c);
        if (a2 + c2 <= b2 + EPS) return minCircle(a, c);
        if (a2 + b2 <= c2 + EPS) return minCircle(a, b);</pre>
189
        double A = 2.0 * (a.x - b.x), B = 2.0 * (a.y - b.y);
190
        double D = 2.0 * (a.x - c.x), E = 2.0 * (a.y - c.y);
191
192
        double C = a.norm() - b.norm(), F = a.norm() - c.norm();
        point p((C * E - B * F) / (A * E - B * D), (A * F - C * D) / (A * E - B * D));
193
        return circle(p, (p - a).len());
194
195
196 circle minCircle(point P□, int N) { // `1—based
        if (N == 1) return circle(P[1], 0.0);
        random_shuffle(P + 1, P + N + 1); circle 0 = minCircle(P[1], P[2]);
198
        Rep(i, 1, N) if(!0.inside(P[i])) { 0 = minCircle(P[1], P[i]);}
199
            Foru(j, 1, i) if(!0.inside(P[j])) { 0 = minCircle(P[i], P[j]);
200
201
                Foru(k, 1, j) if(!0.inside(P[k])) 0 = minCircle(P[i], P[j], P[k]); }
202
        } return 0;
203
```

#### 1.2 二维计算几何基本操作

```
1 const double PI = 3.14159265358979323846264338327950288;
2 double arcSin(const double &a) {
3    return a <= -1.0 ? -PI / 2 : (a >= 1.0 ? PI / 2 : asin(a));
```

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```
5 double arcCos(const double &a) {
        return a \leftarrow -1.0? PI : (a >= 1.0 ? 0 : acos(a));
 7 }
 8 struct point {
        double x, y; // `something omitted`
       point rot(const double &a) const { // `counter-clockwise`
11
            return point(x * cos(a) - y * sin(a), x * sin(a) + y * cos(a));
12
13
       point rot90() const { // `counter-clockwise`
14
            return point(-y, x);
15
       }
16
       point project(const point &p1, const point &p2) const {
17
            const point &q = *this;
           return p1 + (p2 - p1) * (dot(p2 - p1, q - p1) / (p2 - p1).norm());
19
20
       bool onSeq(const point &a, const point &b) const { // `a, b inclusive`
           const point &c = *this;
21
            return sign(dot(a - c, b - c)) \leftarrow 0 && sign(det(b - a, c - a)) == 0;
22
23
24
       double distLP(const point &p1, const point &p2) const { // `dist from *this to line p1->p2`
25
            const point &a = *this:
26
            return fabs(det(p2 - p1, q - p1)) / (p2 - p1).len();
27
       double distSP(const point &p1, const point &p2) const { // `dist from *this to segment [p1, p2]`
28
            const point &q = *this;
29
           if (dot(p2 - p1, q - p1) < EPS) return (q - p1).len();
30
           if (dot(p1 - p2, q - p2) < EPS) return (q - p2).len();
           return distLP(p1, p2);
32
33
       }
       bool inAngle(const point &p1, const point &p2) const { // `det(p1, p2) $\quad qe$ 0`
34
            const point &q = *this; return det(p1, q) > -EPS && det(p2, q) < EPS;</pre>
36
37 };
38 bool lineIntersect(const point &a. const point &b. const point &c. const point &d. point &e) {
       double s1 = det(c - a, d - a);
       double s2 = det(d - b, c - b);
       if (!sign(s1 + s2)) return false;
42
       e = (b - a) * (s1 / (s1 + s2)) + a;
       return true:
45 int segIntersectCheck(const point &a, const point &b, const point &c, const point &d, point &o) {
       static double s1, s2, s3, s4;
       static int iCnt;
47
       int d1 = sign(s1 = det(b - a, c - a));
       int d2 = sign(s2 = det(b - a, d - a));
       int d3 = sign(s3 = det(d - c, a - c));
50
51
       int d4 = sian(s4 = det(d - c, b - c)):
       if ((d1 \wedge d2) == -2 & (d3 \wedge d4) == -2) {
           o = (c * s2 - d * s1) / (s2 - s1);
53
54
           return true;
55
       iCnt = 0;
```

```
if (d1 == 0 \&\& c.onSeg(a, b)) o = c, ++iCnt;
        if (d2 == 0 \&\& d.onSeg(a, b)) o = d, ++iCnt;
        if (d3 == 0 \&\& a.onSeg(c, d)) o = a, ++iCnt;
        if (d4 == 0 \&\& b.onSeq(c, d)) o = b, ++iCnt;
        return iCnt ? 2 : 0; // `不相交返回0, 严格相交返回1, 非严格相交返回2
62 }
 63 struct circle {
        point o;
64
65
        double r, rSqure;
        bool inside(const point &a) { // `非严格`
            return (a - o).len() < r + EPS;
67
68
 69
        bool contain(const circle &b) const { // 消严格
 70
            return sign(b.r + (o - b.o).len() - r) \ll 0;
 71
        bool disjunct(const circle &b) const { // 非严格
 72
73
            return sign(b.r + r - (o - b.o).len()) \le 0;
74
 75
        int isCL(const point &p1, const point &p2, point &a, point &b) const {
            double x = dot(p1 - o, p2 - p1), y = (p2 - p1).norm();
 76
 77
            double d = x * x - y * ((p1 - o).norm() - rSqure);
            if (d < -EPS) return 0:
 78
 79
           if (d < 0) d = 0;
            point q1 = p1 - (p2 - p1) * (x / y);
            point q2 = (p2 - p1) * (sqrt(d) / y);
 81
            a = q1 - q2; b = q1 + q2;
 82
            return q2.len() < EPS ? 1 : 2;
 83
 84
        int tanCP(const point &p, point &a, point &b) const { // `返回切点, 注意可能与 $p$ 重合
 85
 86
            double x = (p - o).norm(), d = x - rSqure;
            if (d < -EPS) return 0;
 87
 88
           if (d < 0) d = 0;
            point q1 = (p - o) * (rSqure / x);
 89
            point q2 = ((p - o) * (-r * sqrt(d) / x)).rot90();
 91
            a = o + (q1 - q2); b = o + (q1 + q2);
 92
            return q2.len() < EPS ? 1 : 2;
 93
94 };
 95 bool checkCrossCS(const circle &cir, const point &p1, const point &p2) { // 注严格
        const point &c = cir.o:
        const double &r = cir.r;
        return c.distSP(p1, p2) < r + EPS
 98
99
           && (r < (c - p1).len() + EPS | | r < (c - p2).len() + EPS);
101 bool checkCrossCC(const circle &cir1, const circle &cir2) { // 非严格`
        const double &r1 = cir1.r, &r2 = cir2.r, d = (cir1.o - cir2.o).len();
        return d < r1 + r2 + EPS && fabs(r1 - r2) < d + EPS;
104 }
int isCC(const circle &cir1, const circle &cir2, point &a, point &b) {
        const point &c1 = cir1.o, &c2 = cir2.o;
        double x = (c1 - c2).norm(), y = ((cir1.rSqure - cir2.rSqure) / x + 1) / 2;
        double d = cir1.rSqure / x - y * y;
        if (d < -EPS) return 0;
```

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```
if (d < 0) d = 0;
        point q1 = c1 + (c2 - c1) * y;
112
        point q2 = ((c2 - c1) * sqrt(d)).rot90();
        a = q1 - q2; b = q1 + q2;
113
        return q2.len() < EPS ? 1 : 2;
114
115 }
116 vector<pair<point, point> > tanCC(const circle &cir1, const circle &cir2) {
         `注意: 如果只有三条切线, 即 $s1 = 1, s2 = 1$, 返回的切线可能重复, 切点没有问题
         vector<pair<point, point> > list;
118
119
        if (cir1.contain(cir2) | | cir2.contain(cir1)) return list;
120
        const point &c1 = cir1.o, &c2 = cir2.o;
        double r1 = cir1.r, r2 = cir2.r;
121
122
        point p, a1, b1, a2, b2;
123
        int s1, s2;
        if (sign(r1 - r2) == 0) {
124
125
            p = c2 - c1;
            p = (p * (r1 / p.len())).rot90();
126
            list.push_back(make_pair(c1 + p, c2 + p));
            list.push_back(make_pair(c1 - p, c2 - p));
        } else {
            p = (c2 * r1 - c1 * r2) / (r1 - r2);
130
            s1 = cir1.tanCP(p, a1, b1);
131
132
            s2 = cir2.tanCP(p, a2, b2);
            if (s1 >= 1 \&\& s2 >= 1) {
133
                list.push_back(make_pair(a1, a2));
134
                list.push_back(make_pair(b1, b2));
135
136
            }
        p = (c1 * r2 + c2 * r1) / (r1 + r2);
138
        s1 = cir1.tanCP(p, a1, b1);
139
140
        s2 = cir2.tanCP(p, a2, b2);
141
        if (s1 >= 1 \&\& s2 >= 1) {
            list.push_back(make_pair(a1, a2));
142
            list.push_back(make_pair(b1, b2));
143
        }
144
145
        return list;
146 }
147 bool distConvexPIn(const point &p1, const point &p2, const point &p3, const point &p4, const point &q) {
        point o12 = (p1 - p2).rot90(), o23 = (p2 - p3).rot90(), o34 = (p3 - p4).rot90();
148
        return (q - p1).inAngle(o12, o23) || (q - p3).inAngle(o23, o34)
149
150
            II ((q - p2).inAngle(o23, p3 - p2) \&& (q - p3).inAngle(p2 - p3, o23));
151 }
152 double distConvexP(int n, point ps[], const point &q) { // 外部点到多边开的距离
        int left = 0, right = n;
        while (right - left > 1) {
            int mid = (left + right) / 2;
155
            if (distConvexPIn(ps[(left + n - 1) % n], ps[left], ps[mid], ps[(mid + 1) % n], q))
156
                riaht = mid:
158
            else left = mid;
159
160
        return q.distSP(ps[left], ps[right % n]);
161 }
162 double areaCT(const circle &cir, point pa, point pb) {
```

```
pa = pa - cir.o; pb = pb - cir.o;
        double R = cir.r;
164
165
        if (pa.len() < pb.len()) swap(pa, pb);</pre>
        if (pb.len() < EPS) return 0;</pre>
166
        point pc = pb - pa;
        double a = pa.len(), b = pb.len(), c = pc.len(), S, h, theta;
169
        double cosB = dot(pb, pc) / b / c, B = acos(cosB);
        double cosC = dot(pa, pb) / a / b, C = acos(cosC);
171
        if (b > R) {
172
           S = C * 0.5 * R * R;
173
            h = b * a * sin(C) / c;
174
           if (h < R && B < PI * 0.5)
175
                S = a\cos(h / R) * R * R - h * sart(R * R - h * h):
176
        } else if (a > R) {
            theta = PI - B - asin(sin(B) / R * b);
            S = 0.5 * b * R * sin(theta) + (C - theta) * 0.5 * R * R;
178
        else S = 0.5 * sin(C) * b * a;
179
180
        return S;
181
182 circle minCircle(const point &a, const point &b) {
        return circle((a + b) * 0.5, (b - a).len() * 0.5);
183
184
185 circle minCircle(const point &a, const point &b, const point &c) { // `钝角三角形没有被考虑'
        double a2( (b - c).norm() ), b2( (a - c).norm() ), c2( (a - b).norm() );
        if (b2 + c2 <= a2 + EPS) return minCircle(b, c);
        if (a2 + c2 <= b2 + EPS) return minCircle(a, c);
188
        if (a2 + b2 <= c2 + EPS) return minCircle(a, b);
        double A = 2.0 * (a.x - b.x), B = 2.0 * (a.y - b.y);
        double D = 2.0 * (a.x - c.x), E = 2.0 * (a.y - c.y);
191
        double C = a.norm() - b.norm(), F = a.norm() - c.norm();
        point p((C * E - B * F) / (A * E - B * D), (A * F - C * D) / (A * E - B * D));
193
        return circle(p, (p - a).len());
195 }
196 circle minCircle(point P[], int N) { // `1—based`
        if (N == 1) return circle(P[1], 0.0);
197
198
        random_shuffle(P + 1, P + N + 1); circle 0 = minCircle(P[1], P[2]);
        Rep(i, 1, N) if(!0.inside(P[i])) { 0 = minCircle(P[1], P[i]);}
199
            Foru(j, 1, i) if(!0.inside(P[j])) { 0 = minCircle(P[i], P[j]);
200
                Foru(k, 1, j) if(!0.inside(P[k])) 0 = minCircle(P[i], P[j], P[k]); }
201
202
        } return 0;
203
```

## 1.3 圆的面积模板

```
struct Event { point p; double alpha; int add; // 本語文的略。
bool operator < (const Event &other) const { return alpha < other.alpha; } };

void circleKCover(circle *c, int N, double *area) { // `$area[k]$ : 至少被覆盖$k$次

static bool overlap[MAXN][MAXN], g[MAXN][MAXN];

Rep(i, 0, N + 1) area[i] = 0.0; Rep(i, 1, N) Rep(j, 1, N) overlap[i][j] = c[i].contain(c[j]);

Rep(i, 1, N) Rep(j, 1, N) g[i][j] = !(overlap[i][j] || overlap[j][i] || c[i].disjunct(c[j]));

Rep(i, 1, N) { static Event events[MAXN * 2 + 1]; int totE = 0, cnt = 1;

Rep(j, 1, N) if (j != i && overlap[j][i]) ++cnt;

Rep(j, 1, N) if (j != i && overlap[j][j]) {
```

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```
circle &a = c[i], &b = c[j]; double l = (a.o - b.o).norm();
               double s = ((a.r - b.r) * (a.r + b.r) / l + 1) * 0.5;
               double t = sqrt(-(1 - sqr(a.r - b.r)) * (1 - sqr(a.r + b.r)) / (1 * 1 * 4.0));
12
               point dir = b.o - a.o, nDir = point(-dir.y, dir.x);
13
               point aa = a.o + dir * s + nDir * t;
               point bb = a.o + dir * s - nDir * t;
               double A = atan2(aa.y - a.o.y, aa.x - a.o.x);
16
17
               double B = atan2(bb.y - a.o.y, bb.x - a.o.x);
               events[totE++] = Event(bb, B, 1); events[totE++] = Event(aa, A, -1); if (B > A) ++cnt;
18
19
           } if (totE == 0) { area[cnt] += PI * c[i].rSquare; continue; }
20
           sort(events, events + totE); events[totE] = events[0];
21
           Foru(j, 0, totE) {
22
               cnt += events[j].add; area[cnt] += 0.5 * det(events[j].p, events[j + 1].p);
               double theta = events[j + 1].alpha — events[j].alpha; if (theta < 0) theta += 2.0 * PI;
               area[cnt] += 0.5 * c[i].rSquare * (theta - sin(theta));
25 }}}
```

## 1.4 多边形相关

```
1 struct Polygon { // stored in [0, n)
       int n; point ps[MAXN];
       Polygon cut(const point &a, const point &b) {
           static Polygon res; static point o; res.n = 0;
           for (int i = 0; i < n; ++i) {
               int s1 = sign(det(ps[i] - a, b - a));
               int s2 = sign(det(ps[(i + 1) % n] - a, b - a));
               if (s1 \le 0) res.ps[res.n++] = ps[i];
               if (s1 * s2 < 0) {
                   lineIntersect(a, b, ps[i], ps[(i + 1) % n], o);
                   res.ps[res.n++] = o;
12
13
           } return res;
14
       bool contain(const point &p) const { // 1 if on border or inner, 0 if outter
15
16
           static point A, B; int res = 0;
17
           for (int i = 0; i < n; ++i) {
18
               A = ps[i]; B = ps[(i + 1) % n];
19
               if (p.onSeg(A, B)) return 1;
               if (sign(A.y - B.y) \le 0) swap(A, B);
               if (sign(p.y - A.y) > 0) continue;
21
               if (sign(p.y - B.y) \le 0) continue;
22
23
               res += (int)(sign(det(B - p, A - p)) > 0);
24
           } return res & 1;
25
       #define qs(x) (ps[x] - ps[0])
26
27
       bool convexContain(point p) const { // `counter-clockwise`
28
           point q = qs(n-1); p = p - ps[0];
           if (!p.inAngle(qs(1), q)) return false;
29
           int L = 0, R = n - 1;
30
           while (L + 1 < R) { int M((L + R) >> 1);
31
32
               if (p.inAngle(qs(M), q)) L = M; else R = M;
           } if (L == 0) return false; point l(qs(L)), r(qs(R));
33
           return sign( fabs(det(l, p)) + fabs(det(p, r)) + fabs(det(r - l, p - l)) - det(l, r) ) == 0;
34
```

```
35
       #undef as
36
       double isPLAtan2(const point &a, const point &b) {
37
           double k = (b - a).alpha(); if (k < 0) k += 2 * PI;
38
           return k:
39
40
41
       point isPL_Get(const point &a, const point &b, const point &s1, const point &s2) {
42
           double k1 = det(b - a, s1 - a), k2 = det(b - a, s2 - a);
           if (sign(k1) == 0) return s1;
43
44
           if (sign(k2) == 0) return s2;
           return (s1 * k2 - s2 * k1) / (k2 - k1);
45
46
       int isPL_Dic(const point &a, const point &b, int l, int r) {
47
48
           int s = (det(b - a, ps[l] - a) < 0) ? -1 : 1;
49
           while (l \ll r) {
               int mid = (l + r) / 2;
50
51
               if (\det(b-a, ps[mid]-a) * s \le 0) r = mid-1;
               else l = mid + 1;
52
53
          }
           return r + 1;
54
55
       int isPL Find(double k. double w[]) {
56
57
           if (k \leftarrow w[0] \mid k > w[n-1]) return 0;
           int l = 0, r = n - 1, mid;
58
           while (l \ll r) {
59
               mid = (l + r) / 2;
60
               if (w[mid] >= k) r = mid - 1;
61
               else l = mid + 1;
62
           } return r + 1;
63
64
       bool isPL(const point &a, const point &b, point &cp1, point &cp2) { // `$0 (log N)$`
65
           static double w[MAXN * 2]; // `pay attention to the array size`
66
           for (int i = 0; i \le n; ++i) ps[i + n] = ps[i];
67
68
           for (int i = 0; i < n; ++i) w[i] = w[i + n] = isPLAtan2(ps[i], ps[i + 1]);
           int i = isPL Find(isPLAtan2(a, b), w):
69
70
           int j = isPL_Find(isPLAtan2(b, a), w);
           double k1 = det(b - a, ps[i] - a), k2 = det(b - a, ps[j] - a);
71
72
           if (sign(k1) * sign(k2) > 0) return false; // `no intersection`
73
           if (sign(k1) == 0 \mid | sign(k2) == 0)  // `intersect with a point or a line in the convex`
               if (sign(k1) == 0) {
74
75
                   if (sign(det(b-a, ps[i+1]-a)) == 0) cp1 = ps[i], cp2 = ps[i+1];
76
                   else cp1 = cp2 = ps[i];
77
                   return true;
78
79
               if (sign(k2) == 0) {
                   if (sign(det(b-a, ps[j+1]-a)) == 0) cp1 = ps[j], cp2 = ps[j+1];
80
                   else cp1 = cp2 = ps[j];
81
82
              }
83
               return true;
84
85
           if (i > j) swap(i, j);
86
           int x = isPL_Dic(a, b, i, j), y = isPL_Dic(a, b, j, i + n);
87
           cp1 = isPL\_Get(a, b, ps[x - 1], ps[x]);
```

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```
cp2 = isPL\_Get(a, b, ps[y-1], ps[y]);
            return true;
90
        double getI(const point &0) const {
91
            if (n <= 2) return 0;
92
            point G(0.0, 0.0);
94
            double S = 0.0, I = 0.0;
            for (int i = 0; i < n; ++i) {
95
                const point &x = ps[i], &y = ps[(i + 1) \% n];
97
                double d = det(x, y);
98
                G = G + (x + y) * d / 3.0;
                S += d;
99
100
            rac{1}{3} G = G / S:
            for (int i = 0; i < n; ++i) {
                point x = ps[i] - G, y = ps[(i + 1) % n] - G;
                I \leftarrow fabs(det(x, y)) * (x.norm() + dot(x, y) + y.norm());
103
104
            return I = I / 12.0 + fabs(S * 0.5) * (0 - G).norm();
105
107 };
```

#### 1.5 直线与凸包求交点

```
ı int isPL(point a, point b, vector<point> &res) { // `点逆时针给出,无三点共线
       static double theta[MAXN];
       for (int i = 0; i < n; ++i) theta[i] = (list[(i + 1) % n] - list[i]).atan2();
       double delta = theta[0];
       for (int i = 0; i < n; ++i) theta[i] = normalize(theta[i] - delta);</pre>
       int x = lower_bound(theta, theta + n, normalize((b - a).atan2() - delta)) - theta;
       int y = lower_bound(theta, theta + n, normalize((a - b).atan2() - delta)) - theta;
       for (int k = 0; k \le 1; ++k, swap(a, b), swap(x, y)) {
           if (y < x) y += n;
           int l = x, r = y, m;
           while (l + 1 < r) {
11
               if (sign(det(b-a, list[(m = (l + r) / 2) % n] - a)) < 0) l = m;
               else r = m:
          }
14
15
          1 %= n, r %= n;
           if (sign(det(b-a, list[r]-list[l])) == 0) {
16
               if (sign(det(b-a, list[l]-a)) == 0)
               return—l; // `直线与 $(list[l], list[r])$ 重合`
19
          }
           else {
20
21
               point p; lineIntersect(list[l], list[r], a, b, p);
               if (p.onSeg(list[l], list[r]))
               res.push_back(p);
23
          }
24
25
       return res.size();
27 }
```

#### 1.6 半平面交

```
1 struct Border {
       point p1, p2; double alpha;
       Border(): p1(), p2(), alpha(0.0) {}
       Border(const point &a, const point &b): p1(a), p2(b), alpha(atan2(p2.y-p1.y, p2.x-p1.x)) {}
       bool operator == (const Border &b) const { return sign(alpha - b.alpha) == 0; }
       bool operator < (const Border &b) const {</pre>
           int c = sign(alpha - b.alpha); if (c != 0) return c > 0;
           return sign(det(b.p2 - b.p1, p1 - b.p1)) >= 0;
11 point isBorder(const Border &a, const Border &b) { // a and b should not be parallel
       point is; lineIntersect(a.p1, a.p2, b.p1, b.p2, is); return is;
12
13 }
14 bool checkBorder(const Border &a, const Border &b, const Border &me) {
       point is; lineIntersect(a.p1, a.p2, b.p1, b.p2, is);
       return sign(det(me.p2 - me.p1, is - me.p1)) > 0;
16
17 }
18 double HPI(int N, Border border□) {
       static Border que[MAXN * 2 + 1]; static point ps[MAXN];
       int head = 0, tail = 0, cnt = 0; // [head, tail)
21
       sort(border, border + N); N = unique(border, border + N) - border;
       for (int i = 0; i < N; ++i) {
22
23
           Border &cur = border[i];
           while (head + 1 < tail && !checkBorder(que[tail - 2], que[tail - 1], cur)) —tail;
24
25
           while (head + 1 < tail && !checkBorder(que[head], que[head + 1], cur)) ++head;</pre>
26
           que[tail++] = cur;
       } while (head + 1 < tail && !checkBorder(que[tail - 2], que[tail - 1], que[head])) —tail;
       while (head + 1 < tail \& !checkBorder(que[head], que[head + 1], que[tail - 1])) ++head;
28
29
       if (tail - head <= 2) return 0.0:
       Foru(i, head, tail) ps[cnt++] = isBorder(que[i], que[(i + 1 == tail) ? (head) : (i + 1)]);
30
       double area = 0; Foru(i, 0, cnt) area += det(ps[i], ps[(i + 1) \% cnt]);
31
       return fabs(area * 0.5); // or (—area * 0.5)
32
33 }
```

## 1.7 最大面积空凸包

```
inline bool toUpRight(const point &a, const point &b) {
       int c = sign(b.y - a.y); if (c > 0) return true;
       return c == 0 \& sign(b.x - a.x) > 0;
3
4 }
5 inline bool cmpByPolarAngle(const point &a, const point &b) { // `counter-clockwise, shorter first if they share
         the same polar angle
       int c = sign(det(a, b)); if (c != 0) return c > 0;
       return sign(b.len() - a.len()) > 0;
8 }
9 double maxEmptyConvexHull(int N, point p□) {
       static double dp[MAXN][MAXN];
       static point vec[MAXN];
       static int seq[MAXN]; // `empty triangles formed with $(0, 0), vec[o], vec[ seq[i] ]$
       double ans = 0.0;
13
```

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```
Rep(o, 1, N) {
           int totVec = 0;
16
           Rep(i, 1, N) if (toUpRight(p[o], p[i])) vec[++totVec] = p[i] - p[o];
           sort(vec + 1, vec + totVec + 1, cmpByPolarAngle);
17
           Rep(i, 1, totVec) Rep(j, 1, totVec) dp[i][j] = 0.0;
           Rep(k, 2, totVec) {
               int i = k - 1;
               while (i > 0 && sign( det(vec[k], vec[i]) ) == 0) —i;
22
               int totSeq = 0;
23
               for (int j; i > 0; i = j) {
24
                   seq[++totSeq] = i;
                   for (j = i - 1; j > 0 \& sign(det(vec[i] - vec[k], vec[j] - vec[k])) > 0; ---j);
                   double v = det(vec[i], vec[k]) * 0.5;
                   if (j > 0) \lor += dp[i][j];
                   dp[k][i] = v;
                   cMax(ans, v);
               } for (int i = totSeq -1; i >= 1; -i) cMax( dp[k][ seq[i] ], dp[k][seq[i + 1]] );
30
31
      } return ans;
33 }
```

#### 1.8 最近点对

```
1 int N; point p[maxn];
 2 bool cmpByX(const point &a, const point &b) { return sign(a.x - b.x) < 0; }
 3 bool cmpByY(const int &a, const int &b) { return p[a].y < p[b].y; }</pre>
 4 double minimalDistance(point *c, int n, int *ys) {
       double ret = 1e+20;
       if (n < 20) {
           Foru(i, 0, n) Foru(j, i + 1, n) cMin(ret, (c[i] - c[j]).len());
           sort(ys, ys + n, cmpByY); return ret;
       } static int mergeTo[maxn];
       int mid = n / 2; double xmid = c[mid].x;
       ret = min(minimalDistance(c, mid, ys), minimalDistance(c + mid, n - mid, ys + mid));
       merge(ys, ys + mid, ys + mid, ys + n, mergeTo, cmpByY);
       copy(mergeTo, mergeTo + n, ys);
14
       Foru(i, 0, n) {
15
           while (i < n && sign(fabs(p[ys[i]].x - xmid) - ret) > 0) ++i;
16
           int cnt = 0;
17
           Foru(j, i + 1, n)
               if (sign(p[ys[j]].y - p[ys[i]].y - ret) > 0) break;
               else if (sign(fabs(p[ys[j]].x - xmid) - ret) \le 0) {
19
                   ret = min(ret, (p[ys[i]] - p[ys[j]]).len());
                   if (++cnt >= 10) break;
       } return ret;
23
24 }
25 double work() {
       sort(p, p + n, cmpByX); Foru(i, 0, n) ys[i] = i; return minimalDistance(p, n, ys);
27 }
```

## 1.9 凸包与点集直径

```
1 vector<point> convexHull(int n, point ps[]) { // `counter-clockwise, strict`
       static point qs[MAXN * 2];
       sort(ps, ps + n, cmpByXY);
       if (n <= 2) return vector<point>(ps, ps + n);
       int k = 0;
       for (int i = 0; i < n; qs[k++] = ps[i++])
           while (k > 1 \& det(qs[k-1] - qs[k-2], ps[i] - qs[k-1]) < EPS) - k;
       for (int i = n - 2, t = k; i >= 0; qs[k++] = ps[i--])
           while (k > t \& det(qs[k-1] - qs[k-2], ps[i] - qs[k-1]) < EPS) - k;
10
       return vector<point>(qs, qs + k);
11 }
12 double convexDiameter(int n, point ps[]) {
       if (n < 2) return 0; if (n == 2) return (ps[1] - ps[0]).len();
       double k, ans = 0;
15
       for (int x = 0, y = 1, nx, ny; x < n; ++x) {
           for(nx = (x == n - 1) ? (0) : (x + 1); ; y = ny) {
16
17
              ny = (y == n - 1) ? (0) : (y + 1);
              if ( sign(k = det(ps[nx] - ps[x], ps[ny] - ps[y])) \le 0) break;
18
19
          } ans = max(ans, (ps[x] - ps[y]).len());
           if (sign(k) == 0) ans = max(ans, (ps[x] - ps[ny]).len());
20
21
       } return ans;
22 }
```

#### 1.10 Farmland

```
1 struct node { int begin[MAXN], *end; } a[MAXN]; // `按对 $p[i]$ 的极角的 atan2 值排字
2 bool check(int n, point p[], int b1, int b2, bool vis[MAXN][MAXN]) {
       static pii l[MAXN * 2 + 1]; static bool used[MAXN];
       int tp(0), *k, p, p1, p2; double area(0.0);
       for (l[0] = pii(b1, b2); ; ) {
          vis[p1 = l[tp].first][p2 = l[tp].second] = true;
           area += det(p[p1], p[p2]);
           for (k = a[p2].begin; k != a[p2].end; ++k) if (*k == p1) break;
          k = (k == a[p2].begin) ? (a[p2].end - 1) : (k - 1);
           if ((l[++tp] = pii(p2, *k)) == l[0]) break;
10
       } if (sign(area) < 0 || tp < 3) return false;</pre>
12
       Rep(i, 1, n) used[i] = false;
       for (int i = 0; i < tp; ++i) if (used[p = l[i].first]) return false; else used[p] = true;</pre>
13
       return true; // `a face with tp vertices`
14
15 }
16 int countFaces(int n, point p□) {
17
       static bool vis[MAXN][MAXN]; int ans = 0;
       Rep(x, 1, n) Rep(y, 1, n) vis[x][y] = false;
18
19
       Rep(x, 1, n) for (int *itr = a[x].begin; itr != a[x].end; ++itr) if (!vis[x][*itr])
           if (check(n, p, x, *itr, vis)) ++ans;
20
21
       return ans;
22 }
```

## 1.11 Voronoi 图

不能有重点, 点数应当不小于 2

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```
if (0i(b) == v) Splice(0p(b), e, v);
 1 #define 0i(e) ((e)→oi)
                                                                                                                                                         else Splice(Dp(b), e, v);
 2 #define Dt(e) ((e)->dt)
                                                                                                                                            55
 3 #define On(e) ((e)→on)
                                                                                                                                                    } return e;
                                                                                                                                            56
 4 #define Op(e) ((e)->op)
                                                                                                                                            57 }
 5 #define Dn(e) ((e)->dn)
                                                                                                                                            58 void Remove(edge *e) {
 6 #define Dp(e) ((e)->dp)
                                                                                                                                                    point *u = 0i(e). *v = Dt(e):
 7 #define Other(e, p) ((e)->oi == p ? (e)->oi)
                                                                                                                                                    if (u \rightarrow in == e) u \rightarrow in = e \rightarrow on;
 8 #define Next(e, p) ((e)\rightarrowoi == p ? (e)\rightarrowon : (e)\rightarrowdn)
                                                                                                                                                    if (v \rightarrow in == e) v \rightarrow in = e \rightarrow dn;
 9 #define Prev(e, p) ((e)\rightarrowoi == p ? (e)\rightarrowop : (e)\rightarrowdp)
                                                                                                                                                    if (0i(e\rightarrow on) == u) e\rightarrow on\rightarrow op = e\rightarrow op; else e\rightarrow on\rightarrow dp = e\rightarrow op;
10 #define V(p1, p2, u, v) (u = p2 \rightarrow x - p1 \rightarrow x, v = p2 \rightarrow y - p1 \rightarrow y)
                                                                                                                                                    if (0i(e \rightarrow op) == u) e \rightarrow op \rightarrow on = e \rightarrow on; else e \rightarrow op \rightarrow dn = e \rightarrow on;
11 #define C2(u1, v1, u2, v2) (u1 * v2 - v1 * u2)
                                                                                                                                                    if (0i(e\rightarrow dn) == v) e\rightarrow dn\rightarrow op = e\rightarrow dp; else e\rightarrow dn\rightarrow dp = e\rightarrow dp;
12 #define C3(p1, p2, p3) ((p2\tox - p1\tox) * (p3\toy - p1\toy) - (p2\toy - p1\toy) * (p3\tox - p1\tox))
                                                                                                                                                    if (0i(e\rightarrow dp) == v) e\rightarrow dp\rightarrow on = e\rightarrow dn; else e\rightarrow dp\rightarrow dn = e\rightarrow dn;
                                                                                                                                            65
13 #define Dot(u1, v1, u2, v2) (u1 * u2 + v1 * v2)
                                                                                                                                            66
                                                                                                                                                    elistInfree++1 = e:
14 #define dis(a,b) (sqrt( (a\rightarrow x-b\rightarrow x)*(a\rightarrow x-b\rightarrow x)+(a\rightarrow y-b\rightarrow y)*(a\rightarrow y-b\rightarrow y) ))
                                                                                                                                            67 }
15 const int maxn = 110024;
                                                                                                                                            68 void Low_tangent(edge *e_l, point *o_l, edge *e_r, point *o_r, edge **l_low, point **OL, edge **r_low, point **OR
16 const int aix = 4:
                                                                                                                                                      } (
17 const double eps = 1e-7;
                                                                                                                                                    for (point *d_l = 0ther(e_l, o_l), *d_r = 0ther(e_r, o_r); ; )
18 int n, M, k;
                                                                                                                                                         if (C3(o_1, o_r, d_1) < -eps)  e_1 = Prev(e_1, d_1), o_1 = d_1, d_1 = 0ther(e_1, o_1);
19 struct gEdge {
                                                                                                                                                         else if (C3(o_1, o_r, d_r) < -eps) e_r = Next(e_r, d_r), o_r = d_r, d_r = 0ther(e_r, o_r);
                                                                                                                                                         else break:
        int u, v; double w;
                                                                                                                                            72
        bool operator <(const gEdge &e1) const { return w < e1.w - eps; }</pre>
                                                                                                                                            73
                                                                                                                                                    *OL = o_l, *OR = o_r; *l_low = e_l, *r_low = e_r;
22 } E[aix * maxn], MST[maxn]:
                                                                                                                                            74 }
23 struct point {
                                                                                                                                            75 void Merge(edge *lr, point *s, edge *rl, point *u, edge **tangent) {
        double x, y; int index; edge *in;
                                                                                                                                                    double l1, l2, l3, l4, r1, r2, r3, r4, cot_L, cot_R, u1, v1, u2, v2, n1, cot_n, P1, cot_P;
        bool operator <(const point &p1) const { return x < p1.x - eps | | (abs(x - p1.x) <= eps && y < p1.y - eps); }
                                                                                                                                                    point *0, *D, *OR, *OL; edge *B, *L, *R;
                                                                                                                                                    Low_tangent(lr, s, rl, u, &L, &OL, &R, &OR);
26 };
                                                                                                                                            78
27 struct edge { point *oi, *dt; edge *on, *op, *dn, *dp; };
                                                                                                                                                    for (*tangent = B = Join(L, OL, R, OR, \emptyset), 0 = OL, D = OR; ; ) {
                                                                                                                                            79
                                                                                                                                                         edge *El = Next(B, 0), *Er = Prev(B, D), *next, *prev;
28
                                                                                                                                            80
                                                                                                                                                         point *l = Other(El, 0), *r = Other(Er, D);
29 point p[maxn], *Q[maxn];
                                                                                                                                            81
30 edge mem[aix * maxn], *elist[aix * maxn];
                                                                                                                                            82
                                                                                                                                                         V(1, 0, 11, 12); V(1, D, 13, 14); V(r, 0, r1, r2); V(r, D, r3, r4);
                                                                                                                                                         double cl = C2(l1, l2, l3, l4), cr = C2(r1, r2, r3, r4);
                                                                                                                                            83
32 void Alloc_memory() { nfree = aix * n; edge *e = mem; for (int i = 0; i < nfree; i++) elist[i] = e++; }
                                                                                                                                            84
                                                                                                                                                         bool BL = cl > eps, BR = cr > eps;
33 void Splice(edge *a, edge *b, point *v) {
                                                                                                                                                        if (!BL && !BR) break;
                                                                                                                                            85
        edge *next;
                                                                                                                                            86
                                                                                                                                                         if (BL) {
        if (0i(a) == v) next = 0n(a), 0n(a) = b; else next = 0n(a), 0n(a) = b;
                                                                                                                                                              double dl = Dot(l1, l2, l3, l4):
                                                                                                                                            87
        if (0i(next) == v) Op(next) = b; else Dp(next) = b;
                                                                                                                                            88
                                                                                                                                                              for (cot_L = dl / cl; ; Remove(El), El = next, cot_L = cot_n) {
37
        if (0i(b) == v) On(b) = next, Op(b) = a; else Dn(b) = next, Dp(b) = a;
                                                                                                                                                                  next = Next(El, 0); V(0ther(next, 0), 0, u1, v1); V(0ther(next, 0), D, u2, v2);
                                                                                                                                            89
38 }
                                                                                                                                            90
                                                                                                                                                                  n1 = C2(u1, v1, u2, v2); if (!(n1 > eps)) break;
39 edge *Make_edge(point *u, point *v) {
                                                                                                                                                                  cot_n = Dot(u1, v1, u2, v2) / n1;
                                                                                                                                            91
        edge *e = elist[--nfree];
                                                                                                                                                                  if (cot_n > cot_L) break;
                                                                                                                                                             }
        e \rightarrow on = e \rightarrow op = e \rightarrow dn = e \rightarrow dp = e; e \rightarrow oi = u; e \rightarrow dt = v;
                                                                                                                                            93
                                                                                                                                                        } if (BR) {
        if (!u \rightarrow in) u \rightarrow in = e;
                                                                                                                                            94
        if (!v \rightarrow in) v \rightarrow in = e;
                                                                                                                                            95
                                                                                                                                                             double dr = Dot(r1, r2, r3, r4);
                                                                                                                                                             for (cot_R = dr / cr; ; Remove(Er), Er = prev, cot_R = cot_P) {
44
        return e;
                                                                                                                                            96
45 }
                                                                                                                                            97
                                                                                                                                                                  prev = Prev(Er, D); V(0ther(prev, D), 0, u1, v1); V(0ther(prev, D), D, u2, v2);
                                                                                                                                                                  P1 = C2(u1, v1, u2, v2); if (!(P1 > eps)) break;
46 edge *Join(edge *a, point *u, edge *b, point *v, int side) {
                                                                                                                                            98
                                                                                                                                                                  cot_P = Dot(u1, v1, u2, v2) / P1;
        edge *e = Make_edge(u, v);
                                                                                                                                            99
        if (side == 1) {
                                                                                                                                           100
                                                                                                                                                                  if (cot P > cot R) break:
            if (0i(a) == u) Splice(0p(a), e, u);
49
                                                                                                                                           101
             else Splice(Dp(a), e, u);
                                                                                                                                                        } l = Other(El, 0); r = Other(Er, D);
50
                                                                                                                                           102
                                                                                                                                                        if (!BL || (BL && BR && cot_R < cot_L)) B = Join(B, 0, Er, r, 0), D = r;
51
             Splice(b, e, v);
52
       } else {
                                                                                                                                           104
                                                                                                                                                         else B = Join(El, 1, B, D, 0), 0 = 1;
             Splice(a, e, u);
                                                                                                                                           105
```

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```
106 }
107 void Divide(int s, int t, edge **L, edge **R) {
        edge *a, *b, *c, *ll, *lr, *rl, *rr, *tangent;
108
        int n = t - s + 1;
109
        if (n == 2) *L = *R = Make_edge(Q[s], Q[t]);
110
        else if (n == 3) {
111
            a = Make\_edge(Q[s], Q[s + 1]), b = Make\_edge(Q[s + 1], Q[t]);
112
            Splice(a, b, Q[s + 1]);
113
            double v = C3(0\lceil s\rceil, 0\lceil s + 1\rceil, 0\lceil t\rceil);
114
115
            if (v > eps) c = Join(a, 0[s], b, 0[t], 0), *L = a, *R = b;
116
            else if (v \leftarrow eps) c = Join(a, Q[s], b, Q[t], 1), *L = c, *R = c;
            else *L = a, *R = b;
118
        } else if (n > 3) {
119
            int split = (s + t) / 2;
            Divide(s, split, &ll, &lr); Divide(split + 1, t, &rl, &rr);
120
            Merge(lr, Q[split], rl, Q[split + 1], &tangent);
121
122
            if (0i(tangent) == 0[s]) ll = tangent;
            if (Dt(tangent) == Q[t]) rr = tangent;
123
             *L = 11; *R = rr;
124
125
126 }
127 void Make_Graph() {
128
        edge *start, *e; point *u, *v;
129
        for (int i = 0; i < n; i++) {
            start = e = (u = &p[i]) \rightarrow in;
130
            do\{ v = 0 ther(e, u);
131
                 if (u < v) E[M++].u = (u - p, v - p, dis(u, v)); // M < aix * maxn
132
            } while ((e = Next(e, u)) != start);
133
134
135 }
136 int b[maxn];
int Find(int x) { while (x != b[x]) \{ b[x] = b[b[x]]; x = b[x]; \} return x; }
138 void Kruskal() {
        memset(b, 0, sizeof(b)); sort(E, E + M);
139
        for (int i = 0: i < n: i++) b\lceil i \rceil = i:
140
141
        for (int i = 0, kk = 0; i < M && kk < n - 1; i++) {
            int m1 = Find(E[i].u), m2 = Find(E[i].v);
142
            if (m1 != m2) b[m1] = m2, MST[kk++] = E[i];
143
144
145 }
146 void solve() {
147
        scanf("%d", &n);
        for (int i = 0; i < n; i++) scanf("%lf%lf", &p[i].x, &p[i].y), p[i].index = i, p[i].in = NULL;
148
        Alloc_memory(); sort(p, p + n);
149
        for (int i = 0; i < n; i++) Q[i] = p + i;
        edge *L, *R; Divide(0, n - 1, &L, &R);
151
152
        M = 0; Make_Graph(); Kruskal();
153 }
154 int main() { solve(); return 0; }
```

## 1.12 四边形双费马点

1 typedef complex<double> Tpoint;

```
2 const double eps = 1e-8;
 3 const double sqrt3 = sqrt(3.0);
 4 bool cmp(const Tpoint &a, const Tpoint &b) {
       return a.real() < b.real() - eps || (a.real() < b.real() + eps && a.imaq() < b.imaq());
 6 }
 7 Tpoint rotate(const Tpoint &a, const Tpoint &b, const Tpoint &c) {
       Tpoint d = b - a; d = Tpoint(-d.imag(), d.real());
       if (Sign(cross(a, b, c)) == Sign(cross(a, b, a + d))) d *= -1.0;
10
       return unit(d);
11 }
12 Tpoint p[10], a[10], b[10];
13 int N, T;
14 double totlen(const Tpoint &p. const Tpoint &a. const Tpoint &b. const Tpoint &c) {
15
       return abs(p-a) + abs(p-b) + abs(p-c);
16 }
17 double fermat(const Tpoint &x, const Tpoint &y, const Tpoint &z, Tpoint &cp) {
       a[0] = a[3] = x; a[1] = a[4] = y; a[2] = a[5] = z;
       double len = 1e100, len2;
19
       for (int i = 0; i < 3; i++) {
           len2 = totlen(a[i], x, y, z);
21
22
           if (len2 < len) len = len2, cp = a[i];
23
24
       for (int i = 0; i < 3; i++) {
           b[i] = rotate(a[i + 1], a[i], a[i + 2]);
25
           b[i] = (a[i + 1] + a[i]) / 2.0 + b[i] * (abs(a[i + 1] - a[i]) * sqrt3 / 2.0);
26
27
       b[3] = b[0];
28
       Tpoint cp2 = intersect(b[0], a[2], b[1], a[3]);
       len2 = totlen(cp2, x, y, z);
31
       if (len2 < len) len = len2, cp = cp2;
32
       return len;
33 }
34 double getans(const Tpoint &a) {
35
       double len = 0; for (int i = 0; i < N; i++) len += abs(a - p[i]);
36
37 }
38 double mindist(const Tpoint &p, const Tpoint &a, const Tpoint &b, const Tpoint &c, const Tpoint &d) {
       return min( min(abs(p-a), abs(p-b)), min(abs(p-c), abs(p-d)));
40 }
41 int main() {
42
       N = 4;
43
       for (cin >> T; T; T—) {
44
           double ret = 1e100, len_cur, len_before, len1, len2, len;
45
           Tpoint cp, cp1, cp2;
           Foru(i, 0, N) cin >> p[i];
46
           Foru(i, 0, N) ret = min(ret, getans(p[i]));
47
48
           Foru(i, 1, N) Foru(j, 1, N) if (j != i) Foru(k, 1, N) if (k != i && k != j) {
49
               cMin(ret, abs(p[0] - p[i]) + abs(p[j] - p[k])
                       + min( min(abs(p[0] - p[j]), abs(p[0] - p[k])),
                              min(abs(p[i] - p[j]), abs(p[i] - p[k]))
51
52
               ret = min(ret, getans(intersect(p[0], p[i], p[i], p[k])));
53
54
```

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```
Foru(i, 0, N) Foru(j, i + 1, N) Foru(k, j + 1, N) {
               double len = fermat(p[i], p[j], p[k], cp);
               ret = min(ret, len + mindist(p[6-i-j-k], p[i], p[j], p[k], cp));
           }
           sort(p, p + N, cmp);
           for (int i = 1; i < N; i++) {
               cp1 = (p[0] + p[i]) / 2.0;
               int j, k;
62
               for (j = 1; j < N \&\& j == i; j++);
               for (k = 6 - i - j, len_before = 1e100; ;) {
                   len1 = fermat(cp1, p[j], p[k], cp2);
                   len1 = fermat(cp2, p[0], p[i], cp1);
                   len = len1 + abs(cp2 - p[j]) + abs(cp2 - p[k]);
                   if (len < len_before - (1e-6)) len_before = len;</pre>
                   else break;
               } ret = min(ret, len_before);
70
71
           } printf("%.4f\n", ret);
72
       return 0;
73
74 }
```

## 1.13 三角形和四边形的费马点

- 费马点: 距几个顶点距离之和最小的点
- 三角形:
  - 若每个角都小于  $120^\circ$ : 以每条边向外作正三角形,得到  $\Delta ABF$ ,  $\Delta BCD$ ,  $\Delta CAE$ ,连接 AD, BE, CF, 三线必共点于费马点. 该点对三边的张角必然是  $120^\circ$ ,也必然是三个三角形外接圆的交点
  - 否则费马点一定是那个大干等于 120° 的顶角
- 四边形:
  - 在凸四边形中, 费马点为对角线的交点
  - 在凹四边形中, 费马点位凹顶点

#### 1.14 三维计算几何基本操作

```
15 double distLL(const point &p1, const point &p2, const point &q1, const point &q2) {
       point p = q1 - p1, u = p2 - p1, v = q2 - q1;
       double d = u.norm() * v.norm() - dot(u, v) * dot(u, v);
17
       if (sign(d) == 0) return p1.distLP(q1, q2);
       double s = (dot(p, u) * v.norm() - dot(p, v) * dot(u, v)) / d;
       return (p1 + u * s).distLP(q1, q2);
20
21 }
22 double distSS(const point &p1, const point &p2, const point &q1, const point &q2) {
       point p = q1 - p1, u = p2 - p1, v = q2 - q1;
       double d = u.norm() * v.norm() - dot(u, v) * dot(u, v);
       if (sign(d) == 0) return min(min((p1 - q1).len(), (p1 - q2).len()),
25
26
                                     min((p2 - q1).len(), (p2 - q2).len()));
       double s1 = (dot(p, u) * v.norm() - dot(p, v) * dot(u, v)) / d;
28
       double s2 = (dot(p, v) * u.norm() - dot(p, u) * dot(u, v)) / d;
       if (s1 < 0.0) s1 = 0.0; if (s1 > 1.0) s1 = 1.0;
30
       if (s2 < 0.0) s2 = 0.0; if (s2 > 1.0) s2 = 1.0;
       point r1 = p1 + u * s1; point r2 = q1 + v * s2;
       return (r1 - r2).len();
33 }
34 bool isFL(const point &p, const point &o, const point &q1, const point &q2, point &res) {
       double a = dot(o, q2 - p), b = dot(o, q1 - p), d = a - b;
       if (sign(d) == 0) return false;
       res = (q1 * a - q2 * b) / d;
38
       return true;
39 }
40 bool isFF(const point &p1, const point &o1, const point &p2, const point &o2, point &a, point &b) {
       point e = det(o1, o2), v = det(o1, e);
       double d = dot(o2, v); if (sign(d) == 0) return false;
43
       point q = p1 + v * (dot(o2, p2 - p1) / d);
44
       a = a; b = a + e;
45
       return true;
46 }
```

## 1.15 凸多面体切割

```
1 vector<vector<point> > convexCut(const vector<vector<point> > &pss, const point &p, const point &o) {
       vector<vector<point> > res;
       vector<point> sec;
       for (unsigned itr = 0, size = pss.size(); itr < size; ++itr) {</pre>
           const vector<point> &ps = pss[itr];
           int n = ps.size();
           vector<point> qs;
           bool dif = false;
           for (int i = 0; i < n; ++i) {
10
               int d1 = sign( dot(o, ps[i] - p) );
11
               int d2 = sign( dot(o, ps[(i + 1) \% n] - p) );
               if (d1 \le 0) qs.push_back(ps[i]);
12
13
               if (d1 * d2 < 0) {
14
15
                   isFL(p, o, ps[i], ps[(i + 1) % n], q); // must return true
16
                   qs.push_back(q);
17
                    sec.push_back(q);
```

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```
if (d1 == 0) sec.push_back(ps[i]);
               else dif = true;
               dif l = dot(o, det(ps[(i + 1) % n] - ps[i], ps[(i + 2) % n] - ps[i])) < -EPS;
21
22
           if (!qs.empty() && dif)
23
               res.insert(res.end(), qs.begin(), qs.end());
^{24}
25
       if (!sec.empty()) {
26
27
           vector<point> tmp( convexHull2D(sec, o) );
28
           res.insert(res.end(), tmp.begin(), tmp.end());
       }
29
30
       return res:
31 }
33 vector<vector<point> > initConvex() {
34
       vector<vector<point> > pss(6, vector<point>(4));
       pss[0][0] = pss[1][0] = pss[2][0] = point(-INF, -INF, -INF);
35
       pss[0][3] = pss[1][1] = pss[5][2] = point(-INF, -INF, INF);
       pss[0][1] = pss[2][3] = pss[4][2] = point(-INF, INF, -INF);
38
       pss[0][2] = pss[5][3] = pss[4][1] = point(-INF, INF, INF);
       pss[1][3] = pss[2][1] = pss[3][2] = point(INF, -INF, -INF);
39
       pss[1][2] = pss[5][1] = pss[3][3] = point(INF, -INF, INF);
       pss[2][2] = pss[4][3] = pss[3][1] = point(INF, INF, -INF);
41
       pss[5][0] = pss[4][0] = pss[3][0] = point(INF, INF);
43
       return pss;
44 }
```

## 1.16 三维凸包

#### 不能有重点

```
1 namespace ConvexHull3D {
       #define volume(a, b, c, d) (mix(ps[b] - ps[a], ps[c] - ps[a], ps[d] - ps[a]))
       vector<Facet> getHull(int n, point ps∏) {
           static int mark[MAXN][MAXN], a, b, c;
           int stamp = 0;
           bool exist = false;
           vector<Facet> facet;
           random\_shuffle(ps, ps + n);
           for (int i = 2; i < n && !exist; i++) {
               point ndir = det(ps[0] - ps[i], ps[1] - ps[i]);
11
               if (ndir.len() < EPS) continue;</pre>
               swap(ps[i], ps[2]);
13
               for (int j = i + 1; j < n && !exist; j++)
                   if (sign(volume(0, 1, 2, j)) != 0) {
14
                       exist = true:
15
                        swap(ps[i], ps[3]);
                       facet.push_back(Facet(0, 1, 2));
                       facet.push_back(Facet(0, 2, 1));
19
           if (!exist) return ConvexHull2D(n, ps);
21
22
           for (int i = 0; i < n; ++i)
```

```
for (int j = 0; j < n; ++j)
23
                   mark[i][j] = 0;
24
           stamp = 0;
25
26
           for (int v = 3; v < n; ++v) {
               vector<Facet> tmp;
27
28
                ++stamp:
               for (unsigned i = 0; i < facet.size(); i++) {</pre>
29
30
                   a = facet[i].a;
                   b = facet[i].b;
31
32
                   c = facet[i].c;
33
                   if (sign(volume(v, a, b, c)) < 0)
                       mark[a][b] = mark[a][c] =
34
35
                       mark[b][a] = mark[b][c] =
36
                       mark[c][a] = mark[c][b] = stamp;
37
                   else tmp.push_back(facet[i]);
               } facet = tmp;
38
39
               for (unsigned i = 0; i < tmp.size(); i++) {
                   a = facet[i].a; b = facet[i].b; c = facet[i].c;
40
                   if (mark[a][b] == stamp) facet.push_back(Facet(b, a, v));
41
                   if (mark[b][c] == stamp) facet.push_back(Facet(c, b, v));
42
43
                   if (mark[c][a] == stamp) facet.push_back(Facet(a, c, v));
44
               }
45
           } return facet;
46
       #undef volume
47
48 }
49 namespace Gravity {
       using ConvexHull3D::Facet;
       point findG(point ps[], const vector<Facet> &facet) {
51
52
           double ws = 0; point res(0.0, 0.0, 0.0), o = ps[facet[0].a];
           for (int i = 0, size = facet.size(); i < size; ++i) {</pre>
53
54
               const point &a = ps[ facet[i].a ], &b = ps[ facet[i].b ], &c = ps[ facet[i].c ];
               point p = (a + b + c + o) * 0.25;
55
56
               double w = mix(a - o, b - o, c - o);
57
               WS += W:
58
               res = res + p * w;
           } res = res / ws;
59
60
           return res;
61
62 }
```

## 1.17 球面点表面点距离

```
double distOnBall(double lati1, double longi1, double lati2, double longi2, double R) {
    lati1 *= PI / 180; longi1 *= PI / 180;
    lati2 *= PI / 180; longi2 *= PI / 180;

double x1 = cos(lati1) * sin(longi1);

double y1 = cos(lati1) * cos(longi1);

double z1 = sin(lati1);

double x2 = cos(lati2) * sin(longi2);

double y2 = cos(lati2) * cos(longi2);

double z2 = sin(lati2);

double theta = acos(x1 * x2 + y1 * y2 + z1 * z2);
```

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```
return R * theta;
return R * theta;
```

## 1.18 长方体表面点距离

## 1.19 最小覆盖球

```
int outCnt; point out[4], res; double radius;
2 void ball() {
       static point q[3];
       static double m[3][3], sol[3], L[3], det;
       int i, j; res = point(0.0, 0.0, 0.0); radius = 0.0;
       switch (outCnt) {
       case 1: res = out[0]; break;
       case 2: res = (out[0] + out[1]) * 0.5; radius = (res - out[0]).norm();
            break:
10
            q[0] = out[1] - out[0]; q[1] = out[2] - out[0];
            for (i = 0; i < 2; ++i) for (j = 0; j < 2; ++j)
                m[i][j] = dot(q[i], q[j]) * 2.0;
13
14
            for (i = 0; i < 2; ++i) sol[i] = dot(q[i], q[i]);
            det = m[0][0] * m[1][1] - m[0][1] * m[1][0];
            if (sign(det) == 0) return;
            L[0] = (sol[0] * m[1][1] - sol[1] * m[0][1]) / det;
            L[1] = (sol[1] * m[0][0] - sol[0] * m[1][0]) / det;
            res = out\lceil 0 \rceil + q \lceil 0 \rceil * L \lceil 0 \rceil + q \lceil 1 \rceil * L \lceil 1 \rceil;
            radius = (res - out[0]).norm();
            break;
21
22
23
            q[0] = out[1] - out[0]; q[1] = out[2] - out[0]; q[2] = out[3] - out[0];
            for (i = 0; i < 3; ++i) for (j = 0; j < 3; ++j) m[i][j] = dot(q[i], q[j]) * 2;
24
            for (i = 0; i < 3; ++i) sol[i] = dot(q[i], q[i]);
```

```
det = m[0][0] * m[1][1] * m[2][2] + m[0][1] * m[1][2] * m[2][0]
27
               + m[0][2] * m[2][1] * m[1][0] - m[0][2] * m[1][1] * m[2][0]
              -m[0][1] * m[1][0] * m[2][2] - m[0][0] * m[1][2] * m[2][1];
29
           if (sign(det) == 0) return;
           for (j = 0; j < 3; ++j) { for (i = 0; i < 3; ++i) m[i][j] = sol[i];
30
              L[j] = (m[0][0] * m[1][1] * m[2][2] + m[0][1] * m[1][2] * m[2][0]
                     + m[0][2] * m[2][1] * m[1][0] - m[0][2] * m[1][1] * m[2][0]
32
                    -m[0][1] * m[1][0] * m[2][2] - m[0][0] * m[1][2] * m[2][1]) / det;
33
               for (i = 0; i < 3; ++i) m[i][j] = dot(q[i], q[j]) * 2;
34
35
           } res = out[0];
           for (i = 0; i < 3; ++i) res += q[i] * L[i]; radius = (res - out[0]).norm();
37
38 }
39 void minball(int n, point pt□) {
       ball();
       if (outCnt < 4) for (int i = 0; i < n; ++i)
42
           if ((res - pt[i]).norm() > +radius + EPS) {
               out[outCnt] = pt[i]; ++outCnt; minball(i, pt); --outCnt;
43
              if (i > 0) {
                   point Tt = pt[i];
45
                   memmove(&pt[1], &pt[0], sizeof(point) * i);
46
                   pt[0] = Tt;
50 }
51 pair<point, double> main(int npoint, point pt[]) { // 0—based
       random_shuffle(pt, pt + npoint); radius = -1;
       for (int i = 0; i < npoint; i++) { if ((res - pt[i]).norm() > EPS + radius) {
           outCnt = 1; out[0] = pt[i]; minball(i, pt); } }
55
       return make_pair(res, sqrt(radius));
56 }
```

## 1.20 三维向量操作矩阵

• 绕单位向量  $u = (u_x, u_y, u_z)$  右手方向旋转  $\theta$  度的矩阵:

```
\begin{bmatrix} \cos\theta + u_x^2(1 - \cos\theta) & u_x u_y(1 - \cos\theta) - u_z \sin\theta & u_x u_z(1 - \cos\theta) + u_y \sin\theta \\ u_y u_x(1 - \cos\theta) + u_z \sin\theta & \cos\theta + u_y^2(1 - \cos\theta) & u_y u_z(1 - \cos\theta) - u_x \sin\theta \\ u_z u_x(1 - \cos\theta) - u_y \sin\theta & u_z u_y(1 - \cos\theta) + u_x \sin\theta & \cos\theta + u_z^2(1 - \cos\theta) \end{bmatrix}
= \cos\theta I + \sin\theta \begin{bmatrix} 0 & -u_z & u_y \\ u_z & 0 & -u_x \\ -u_y & u_x & 0 \end{bmatrix} + (1 - \cos\theta) \begin{bmatrix} u_x^2 & u_x u_y & u_x u_z \\ u_y u_x & u_y^2 & u_y u_z \\ u_z u_x & u_z u_y & u_z^2 \end{bmatrix}
```

- 点 a 绕单位向量  $u=(u_x,u_y,u_z)$  右手方向旋转  $\theta$  度的对应点为  $a'=a\cos\theta+(u\times a)\sin\theta+(u\otimes u)a(1-\cos\theta)$
- 关于向量 v 作对称变换的矩阵  $H = I 2 \frac{vv^T}{v^T v}$
- 点 a 对称点:  $a' = a 2\frac{v^T a}{v^T v} \cdot v$

# 1.21 立体角

对于任意一个四面体 OABC,从 O 点观察  $\Delta ABC$  的立体角  $\tan \frac{\Omega}{2} = \frac{\min(\overrightarrow{a}, \overrightarrow{b}, \overrightarrow{c})}{|a||b||c|+(\overrightarrow{a} \cdot \overrightarrow{b})|c|+(\overrightarrow{a} \cdot \overrightarrow{c})|b|+(\overrightarrow{b} \cdot \overrightarrow{c})|a|}$ .

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## 2 数据结构

## 2.1 动态凸包 (只支持插入)

```
1 #define x first // `upperHull $\leftarrow (x, y)$
 2 #define y second // `lowerHull $\leftarrow (x, -y)$`
 3 typedef map<int, int> mii;
 4 typedef map<int, int>::iterator mit;
 5 struct point { point(const mit &p): x(p→first), y(p→second) {} };
 6 inline bool checkInside(mii &a, const point &p) { // `border inclusive`
       int x = p.x, y = p.y; mit p1 = a.lower_bound(x);
       if (p1 == a.end()) return false; if (p1\rightarrowx == x) return y <= p1\rightarrowy;
       if (p1 == a.begin()) return false; mit p2(p1—);
       return sign(det(p - point(p1), point(p2) - p)) >= 0;
11 } inline void addPoint(mii &a, const point &p) { // `no collinear points`
       int x = p.x, y = p.y; mit pnt = a.insert(make_pair(x, y)).first, p1, p2;
       for (pnt \rightarrow y = y; ; a.erase(p2)) {
13
           p1 = pnt; if (++p1 == a.end()) break;
14
           p2 = p1; if (++p1 == a.end()) break;
           if (det(point(p2) - p, point(p1) - p) < 0) break;
       } for ( ; ; a.erase(p2)) {
           if ((p1 = pnt) == a.begin()) break;
           if (--p1 == a.begin()) break; p2 = p1--;
19
           if (\det(point(p2) - p, point(p1) - p) > 0) break;
21
22 }
```

## 2.2 Rope 用法

## 2.3 Treap

```
1 struct node { int key, prio, size; node *ch[2]; } base[MAXN], *top, *root, *null, nil;
2 typedef node *tree;
3 tree newNode(int key) {
4    static int seed = 3312;
5    top—key = key; top—prio = seed = int(seed * 48271LL % 2147483647);
6    top—size = 1; top—ch[0] = top—ch[1] = null; return top++;
7 }
8 void Rotate(tree &x, int d) {
9    tree y = x—ch[!d]; x—ch[!d] = y—ch[d]; y—ch[d] = x; y—size = x—size;
10    x—size = x—ch[0]—size + 1 + x—ch[1]—size; x = y;
11 }
```

```
12 void Insert(tree &t, int key) {
         if (t == null) t = newNode(key);
         else { int d = t \rightarrow key < key; Insert(t \rightarrow ch[d], key); ++t \rightarrow size;
              if (t->ch[d]->prio < t->prio) Rotate(t, !d);
15
16
17 }
18 void Delete(tree &t, int key) {
         if (t->key != key) { Delete(t->ch[t->key < key], key); --t->size; }
         else if (t\rightarrow ch[0] == null) t = t\rightarrow ch[1];
20
21
         else if (t\rightarrow ch[1] == null) t = t\rightarrow ch[0];
         else { int d = t \rightarrow ch[0] \rightarrow prio < t \rightarrow ch[1] \rightarrow prio;
23
              Rotate(t, d); Delete(t\rightarrowch[d], key); —t\rightarrowsize;
24
25 }
```

## 2.4 可持久化 Treap

```
inline bool randomBySize(int a, int b) {
         static long long seed = 1;
         return (seed = seed * 48271 % 2147483647) * (a + b) < 2147483647LL * a;
 4 }
 5 tree merge(tree x, tree y) {
         if (x == null) return y; if (y == null) return x;
         tree t = NULL;
         if (randomBySize(x\rightarrow size, y\rightarrow size)) t = newNode(x), t\rightarrow r = merge(x\rightarrow r, y);
         else t = newNode(y), t \rightarrow l = merge(x, y \rightarrow l);
10
         update(t); return t;
11 }
12 void splitByKey(tree t, int k, tree &l, tree &r) { // `$[—\infty, k) [k, +infty)$`
         if (t == null) l = r = null;
14
         else if (t\rightarrow key < k) l = newNode(t), splitByKey(t\rightarrow r, k, l\rightarrow r, r), update(l);
                                  r = \text{newNode}(t), \text{splitByKey}(t \rightarrow l, k, l, r \rightarrow l), \text{update}(r);
15
16 }
17 void splitBySize(tree t, int k, tree &l, tree &r) { // `$[1, k) [k, +\infty)$`
         static int s; if (t == null) l = r = null;
         else if ((s = t\rightarrow l\rightarrow size + 1) < k) l = newNode(t), splitBySize(t\rightarrow r, k - s, l\rightarrow r, r), update(l);
20
                                                    r = \text{newNode}(t), \text{splitBySize}(t \rightarrow l, k, l, r \rightarrow l), \text{update}(r);
21 }
```

# 2.5 左偏树

```
1 tree merge(tree a, tree b) {
2     if (a == null) return b;
3     if (b == null) return a;
4     if (a \to key > b \to key) swap(a, b);
5     a \to rc = merge(a \to rc, b);
6     a \to rc \to fa = a;
7     if (a \to lc \to dist < a \to rc \to dist) swap(a \to lc, a \to rc);
8     a \to dist = a \to rc \to dist + 1;
9     return a;
10 }</pre>
```

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#### 2.6 Link-Cut Tree

```
1 struct node { int rev; node *pre, *ch[2]; } base[MAXN], nil, *null;
 2 typedef node *tree;
 3 #define isRoot(x) (x\rightarrow pre\rightarrow ch[0] != x && x\rightarrow pre\rightarrow ch[1] != x)
 4 #define isRight(x) (x->pre->ch[1] == x)
 5 inline void MakeRev(tree t) { if (t != null) { t\rightarrow rev ^= 1; swap(t\rightarrow ch[0], t\rightarrow ch[1]); } }
 6 inline void PushDown(tree t) { if (t->rev) { MakeRev(t->ch[0]); MakeRev(t->ch[1]); t->rev = 0; } }
 7 inline void Rotate(tree x) {
        tree y = x \rightarrow pre; PushDown(y); PushDown(x);
        int d = isRight(x);
 9
        if (!isRoot(y)) y->pre->ch[isRight(y)] = x; x->pre = y->pre;
        if ((y\rightarrow ch[d] = x\rightarrow ch[!d]) != null) y\rightarrow ch[d]\rightarrow pre = y;
        x\rightarrow ch[!d] = y; y\rightarrow pre = x; Update(y);
12
13 }
14 inline void Splay(tree x) {
        PushDown(x); for (tree y; !isRoot(x); Rotate(x)) {
             y = x \rightarrow pre; if (!isRoot(y)) Rotate(isRight(x) != isRight(y) ? x : y);
17
        } Update(x);
18 }
19 inline void Splay(tree x, tree to) {
        PushDown(x); for (tree y; (y = x \rightarrow pre) != to; Rotate(x)) if (y \rightarrow pre != to)
             Rotate(isRight(x) != isRight(y) ? x : y);
        Update(x);
22
23 }
24 inline tree Access(tree t) {
        tree last = null; for (; t != null; last = t, t = t\rightarrowpre) Splay(t),t\rightarrowch[1] = last, Update(t);
        return last;
27 }
28 inline void MakeRoot(tree t) { Access(t); Splay(t); MakeRev(t); }
29 inline tree FindRoot(tree t) { Access(t); Splay(t); tree last = null;
30
         for ( ; t != null; last = t, t = t->ch[0]) PushDown(t); Splay(last); return last;
31 }
32 inline void Join(tree x, tree y) { MakeRoot(y); y→pre = x; }
33 inline void Cut(tree t) {Access(t); Splay(t); t->ch[0]->pre = null; t->ch[0] = null; Update(t);}
34 inline void Cut(tree x, tree y) {
        tree upper = (Access(x), Access(y));
        if (upper == x) { Splay(x); y\rightarrowpre = null; x\rightarrowch[1] = null; Update(x); }
        else if (upper == y) { Access(x); Splay(y); x \rightarrow pre = null; y \rightarrow ch[1] = null; Update(y); }
        else assert(0); // `impossible to happen`
39 }
```

## 2.7 树链剖分

```
1 #include <cstdio>
2 #include <cstdlib>
3 #include <cstring>
 4 #include <cmath>
5 #include <iostream>
6 #include <fstream>
7 #include <algorithm>
8 #include <vector>
9 #include <string>
10 #define lson l,mid,rt<<1</pre>
#define rson mid+1,r,rt<<1|1</pre>
12
13 using namespace std;
14
15 const int MAX = 111111;
16 typedef long long LL;
17 typedef vector<int>::iterator iter;
18 struct qry_node {
       int u,v,w;
20 }qrys[MAX];
21 struct tree_node {
       LL sum;
       LL mark;
23
24 }tree[MAX*4];
25 vector<int> ori[MAX];
26 int pre[MAX], size[MAX], heavy[MAX], deep[MAX], f[MAX][20];
27 int num[MAX],block[MAX],pathHead[MAX],ind = 0;
29 void insert(int u,int v)
30 {
31
       ori[u].push_back(v);
       ori[v].push_back(u);
32
33 }
34
35 void prepare_split(int u,int pre)
36 {
37
       int tmp = 0;
38
       pre[u] = pre;
       for (iter it = ori[u].begin(); it != ori[u].end(); ++it) {
40
           int v = (*it);
```

```
if (v != pre) {
                prepare_split(v,u);
                if (size[v] > tmp) {
43
44
                   tmp = size[v];
                   heavy[u] = v;
                size[u] += size[v];
49
50
51
       size[u]++;
52 }
53
54 void split(int u,int bel)
       block[u] = num[u] = ++ind;
57
       pathHead[u] = bel;
       if (heavy[u]) split(heavy[u],bel);
58
       block[u] = max(block[u],block[heavy[u]]);
       for (iter it = ori[u].begin(); it != ori[u].end(); ++ it) {
           int v = (*it);
61
           if (v != pre[u] && heavy[u] != v) {
62
63
                split(v,v);
               block[u] = max(block[u],block[v]);
64
65
66
67 }
69 void push_up(int l,int r,int rt)
70 {
       if (l != r) tree[rt].sum = tree[rt<<1].sum + tree[(rt<<1)+1].sum;
71
72 }
73 void push_down(int l,int r,int rt)
74 {
       if (tree[rt].mark != 0 && l != r) {
75
76
           int mid = (l + r) \gg 1;
           tree[rt << 1].mark += tree[rt].mark;</pre>
77
           tree[rt << 1 | 1].mark += tree[rt].mark;</pre>
           tree[rt << 1].sum += (mid - l + 1) * tree[rt].mark;</pre>
           tree[rt \ll 1 \mid 1].sum += (r - mid) * tree[rt].mark;
           tree[rt].mark = 0;
82
83 }
84
85 void build(int l,int r,int rt)
       tree[rt].sum = tree[rt].mark = 0;
87
       if (l == r) return;
88
       int mid = (l+r)>>1;
       build(lson);
91
       build(rson);
93 void upd(int l,int r,int rt,int a,int b,int c)
```

```
94 {
 95
        push_down(l,r,rt);
        int tmp = tree[rt].sum;
 96
        if (a \le 1 \&\& b \ge r) \{
 97
            tree[rt].sum += (r - l + 1) * c;
 98
            tree[rt].mark += c;
100
            return;
101
        int mid = (l + r) \gg 1;
102
103
        if (a \leftarrow mid) upd(lson,a,b,c);
104
        if (b > mid) upd(rson,a,b,c);
105
        push_up(l,r,rt);
106 }
107 LL gry(int l,int r,int rt,int a,int b)
108 {
        push_down(l,r,rt);
109
        if (a \le 1 \&\& b \ge r) {
            return tree[rt].sum;
112
        int mid = (l + r) \gg 1;
113
114
        LL ret = 0;
        if (a <= mid) ret += qry(lson,a,b);</pre>
115
116
        if (b > mid) ret += qry(rson, a, b);
        return ret;
117
118 }
119
120 void lca_prepare(int u)
121 {
        for (iter it = ori[u].begin();it != ori[u].end(); ++it) {
123
            int v = (*it);
            if (v != pre[u]) {
124
125
                 deep[v] = deep[u]+1;
                 f[v][0] = u;
126
127
                 for (int tmp = u, dep = 0; tmp; f[v][dep+1] = f[tmp][dep], tmp = f[tmp][dep], dep++);
128
                lca_prepare(v);
129
130
131 }
132
133 int get_lca(int u,int v)
134 {
135
        int lose = abs(deep[u] - deep[v]), pos = 0;
136
        if (deep[u] < deep[v]) swap(u,v);</pre>
        while (lose) {
137
138
            if (lose & 1) u = f[u][pos];
139
            pos ++;
140
            lose >>= 1;
141
        }
142
        pos = 0;
        while (u != v) {
143
144
            if (f[u][pos] != f[v][pos] || (f[u][pos] == f[v][pos] && !pos)) {
                u = f[u][pos];
145
146
                v = f[v][pos];
```

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```
else {
149
                pos---;
150
151
152
153
        return u;
154 }
155
156 int n,m;
157
158 int main()
159 {
160
        freopen("tree.in", "r", stdin);
161
        freopen("tree.out", "w", stdout);
        ios::sync_with_stdio(false);
162
163
164
        cin >> n;
165
        for (int i = 1; i < n; ++i) {
            int a,b;
166
167
            cin >> a >> b;
            a ++ ,b ++ ;
168
169
            insert(a,b);
170
        memset(pre,0,sizeof(pre));
171
        memset(size,0,sizeof(size));
172
        prepare_split(1,1);
173
        split(1,1);
174
        lca_prepare(1);
175
176
        build(1,n,1);
177
        cin >> m;
178
        for (int i = 1; i <= m; ++i) {
            string c;
179
            cin >> c;
180
            if (c[0] == 'A') {
181
182
                int u,v,w,lca;
                cin >> u >> v >> w;
183
                u++, v++;
184
                lca = get_lca(u,v);
185
                while (pathHead[u] != pathHead[lca]) {
186
187
                    upd(1,n,1,num[pathHead[u]],num[u],w);
                    u = pre[pathHead[u]];
188
                }upd(1,n,1,num[lca],num[u],w);
189
                while (pathHead[v] != pathHead[lca]) {
190
191
                    upd(1,n,1,num[pathHead[v]],num[v],w);
                    v = pre[pathHead[v]];
192
                }upd(1,n,1,num[lca],num[v],w);
193
194
                upd(1,n,1,num[lca],num[lca],-w);
195
            else {
196
197
                int u;
                cin >> u;
198
199
                u++;
```

147

pos++;

# 2.8 Splay 维护数列

```
1 #include <iostream>
 2 #include <cstdio>
 3 #include <cstdlib>
 4 #define keyTree root→ch[1]→ch[0]
 6 using namespace std;
 8 \text{ const int } N = 500000;
 9 const int INF = 1001;
10 int n, m, a[N];
11
12 int max(int x, int y, int z)
13 {
        return max(x, max(y, z));
14
15 }
16
17 struct node {
        int key, maxL, maxR, maxSum, sum, same, size;
        bool rev;
19
        node *pre, *ch[2];
20
        inline void reverse(){
21
22
             if (size == 0) return;
             rev ^= 1;
23
24
             swap(ch[0], ch[1]);
             swap(maxL, maxR);
25
26
27
        inline void saming(int x){
             if (size == 0) return;
28
29
             key = same = x;
30
             maxL = maxR = maxSum = sum = x * size;
31
             if (x < 0)
                  maxL = maxR = maxSum = x;
32
33
        inline void push_up(){
34
35
             sum = ch[0] \rightarrow sum + ch[1] \rightarrow sum + key;
             size = ch[0] \rightarrow size + ch[1] \rightarrow size + 1;
36
             maxL = max(ch[0] \rightarrow maxL, ch[0] \rightarrow sum+key, ch[0] \rightarrow sum+key+ch[1] \rightarrow maxL);
37
             maxR = max(ch[1] \rightarrow maxR, ch[1] \rightarrow sum+key, ch[1] \rightarrow sum+key+ch[0] \rightarrow maxR);
38
             maxSum = max(ch[0] \rightarrow maxSum, max(ch[1] \rightarrow maxSum, max(ch[0] \rightarrow maxR+key, max(ch[1] \rightarrow maxL+key, max(ch[0] \rightarrow maxR+key+key)
39
           ch[1]—>maxL,key)))));
40
        inline void push_down(){
41
42
             if (rev){
                  ch[0]-->reverse();
43
                  ch[1]-->reverse();
44
```

```
rev = 0;
                                                                                                                                                      99
              if (same != INF){
                                                                                                                                                     100
                   ch[0]—>saming(same);
                   ch[1]—>saming(same);
                                                                                                                                                     102
                                                                                                                                                     103
              same = INF;
                                                                                                                                                     104
52
53 };
                                                                                                                                                     106
54
                                                                                                                                                     107
55 class splayTree{
                                                                                                                                                     108
56 public:
                                                                                                                                                     109
                                                                                                                                                                    x \rightarrow ch[c] = y;
57
         node *root, *null:
                                                                                                                                                     110
                                                                                                                                                                    y \rightarrow pre = x;
         node buf[N]; // 内存池
                                                                                                                                                     111
                                                                                                                                                                    if (y == root)
         int top; // 内存池使用量
                                                                                                                                                     112
                                                                                                                                                                         root = x;
         node *stk[N]; // 内存回收
                                                                                                                                                     113
                                                                                                                                                                    y→push_up();
61
         int cnt; // 内存回收量
                                                                                                                                                     114
         int num;
62
63
         int pos, tot, c, pop;
                                                                                                                                                     116
         inline void erase(node *x){
64
                                                                                                                                                     117
65
              x \rightarrow size = x \rightarrow sum = x \rightarrow maxL = x \rightarrow maxR = x \rightarrow maxSum = 0;
                                                                                                                                                     118
66
67
         inline node *newNode(int value){
                                                                                                                                                     120
                                                                                                                                                                              break;
              node *x;
                                                                                                                                                     121
                                                                                                                                                                        }
              if (cnt) x = stk[cnt-];
                                                                                                                                                     122
69
              else x = \&buf[top++];
70
              x\rightarrow key = x\rightarrow maxL = x\rightarrow maxR = x\rightarrow maxSum = x\rightarrow sum = value;
71
                                                                                                                                                     124
              x \rightarrow size = 1, x \rightarrow rev = 0;
                                                                                                                                                     125
              x \rightarrow pre = x \rightarrow ch[0] = x \rightarrow ch[1] = null;
                                                                                                                                                                         else
                                                                                                                                                     126
              x \rightarrow same = INF;
75
              return x;
                                                                                                                                                     128
                                                                                                                                                     129
                                                                                                                                                                    x \rightarrow push_up();
         inline void init(){
                                                                                                                                                     130
              top = cnt = 0;
78
                                                                                                                                                     131
              num = n;
                                                                                                                                                    132
              null = newNode(—INF);
                                                                                                                                                     133
              null \rightarrow size = 0, null \rightarrow sum = 0;
                                                                                                                                                     134
              root = newNode(—INF);
                                                                                                                                                     135
              root \rightarrow sum = 0;
                                                                                                                                                     136
              root \rightarrow ch[1] = newNode(-INF);
                                                                                                                                                     137
              root->ch[1]-->pre = root;
                                                                                                                                                     138
              root->ch[1]->sum =0;
                                                                                                                                                     139
                                                                                                                                                     140
                                                                                                                                                                    splay(t, x);
         inline node *build(int l,int r){
                                                                                                                                                     141
88
              if (l>r) return null;
                                                                                                                                                     142
              int mid = (l+r) \gg 1;
                                                                                                                                                     143
              node *x = newNode(a[mid]);
91
                                                                                                                                                     144
92
              x \rightarrow ch[0] = build(l,mid-1);
                                                                                                                                                     145
              x \rightarrow ch[1] = build(mid+1,r);
                                                                                                                                                     146
              if (x\rightarrow ch[0] != null) x\rightarrow ch[0]\rightarrow pre = x;
                                                                                                                                                     147
              if (x\rightarrow ch[1] != null) x\rightarrow ch[1]\rightarrow pre = x;
                                                                                                                                                     148
              x->push_up();
                                                                                                                                                                    num += tot;
                                                                                                                                                     149
              return x;
                                                                                                                                                     150
```

```
98
         inline void rotate(node *x,int c){
               node *y = x \rightarrow pre;
              y→push_down();
               x->push_down();
              y \rightarrow ch[!c] = x \rightarrow ch[c];
              if (y\rightarrow ch[!c] != null)
                   y\rightarrow ch[!c]\rightarrow pre = y;
              x \rightarrow pre = y \rightarrow pre;
              if (x->pre != null)
                    x \rightarrow pre \rightarrow ch[y == x \rightarrow pre \rightarrow ch[1]] = x;
         inline void splay(node *x, node *g){
              x->push_down();
               while (x \rightarrow pre != g){
                   if (x \rightarrow pre \rightarrow pre == g){
                         rotate(x, x == x \rightarrow pre \rightarrow ch[0]);
                   node *y = x \rightarrow pre, *z = y \rightarrow pre;
                    int f = (y == z \rightarrow ch[0]);
                   if (x == y \rightarrow ch[f])
                          rotate(x, !f), rotate(x, f);
                         rotate(y, f), rotate(x, f);
         inline void select(node *x, int k){
              node *t = root;
               while (true) {
                    t→push_down();
                    int tmp = t->ch[0]->size;
                    if (tmp == k) break;
                   if (tmp < k) k = tmp + 1, t = t \rightarrow ch[1];
                    else t = t \rightarrow ch[0];
         inline void recycle(node *x){
              if (x->ch[0] != null) recycle(x->ch[0]);
              stk[++cnt] = x;
              if (x\rightarrow ch[1] != null) recycle(x\rightarrow ch[1]);
         inline void insert(){
              scanf("%d%d", &pos, &tot);
               for (int i = 1; i \leftarrow tot; ++i)
```

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```
151
                 scanf("%d", &a[i]);
            select(null, pos);
152
            select(root, pos+1);
153
            keyTree = build(1, tot);
154
            keyTree->pre = root->ch[1];
155
             splay(keyTree, null);
156
157
        inline void del(){
158
            scanf("%d%d", &pos, &tot);
159
160
             select(null, pos-1);
161
            select(root, pos+tot);
            if (keyTree != null){
163
                 num -= keyTree->size;
164
                 recycle(keyTree);
165
                 root \rightarrow ch[1] \rightarrow ch[0] = null;
                 root->ch[1]->push_up();
166
167
                 root→push_up();
168
             splay(root->ch[1], null);
170
        inline void make_same(){
171
            scanf("%d%d%d", &pos, &tot, &c);
172
173
            select(null, pos-1);
            select(root, pos+tot);
174
            if (keyTree != null){
175
                 keyTree→saming(c);
176
                 splay(keyTree, null);
177
178
        inline void reverse(){
180
            scanf("%d%d", &pos, &tot);
181
182
            select(null, pos-1);
            select(root, pos+tot);
183
            if (keyTree != null){
184
                 keyTree→reverse();
185
186
                 splay(keyTree, null);
            }
187
188
        inline void max_sum(){
189
             printf("%d\n", root->maxSum);
190
191
192
        inline void get_sum(){
            scanf("%d%d", &pos, &tot);
193
            select(null, pos-1);
194
195
            select(root, pos+tot);
            if (keyTree != null){
196
                 printf("%d\n", keyTree->sum);
197
                 splay(keyTree, null);
198
199
                 keyTree->push_down();
            } else printf("0\n");
201
202 }spt;
203
```

```
204 int main(){
        scanf("%d%d", &n, &m);
205
        for (int i = 1; i <= n; ++i)
206
            scanf("%d", &a[i]);
207
        spt.init();
208
        spt.keyTree = spt.build(1,n);
209
        spt.keyTree->pre = spt.root->ch[1];
210
211
        spt.splay(spt.keyTree, spt.null);
        char op[30];
212
213
        for (int i = 1; i <= m; ++i) {
214
            scanf("%s", op);
215
            switch (op[0]){
216
            case 'I': spt.insert(); break;
217
            case 'D': spt.del(); break;
218
            case 'R': spt.reverse(); break;
            case 'G': spt.get_sum(); break;
219
220
            case 'S': spt.make_same(); break;
            case 'M':
221
222
                if (op[2] == 'X') spt.max_sum();
                else spt.make_same(); break;
223
224
225
226
        return 0;
227 }
```

## 3 字符串相关

#### 3.1 Manacher

```
// 要处理的字符串
1 char t[1001];
2 char s[1001 * 2]; // 中间插入特殊字符以后的
3 int Z[1001 * 2], L, R; // Gusfield's Algorithm
5 // 由6往左, 由6往右, 对称地做字符匹配
6 int match(int a, int b)
7 {
      int i = 0;
      while (a - i) = 0 \& b + i < n \& s[a - i] = s[b + i]) i++;
      return i;
10
11 }
12
13 void longest_palindromic_substring()
14 {
      int N = strlen(t);
15
16
      // 在 t 中插入特殊字符, 存放到 s
17
      memset(s, '.', N*2+1);
18
      for (int i=0; i<N; ++i) s[i*2+1] = t[i];
19
      N = N*2+1;
20
21
      // modified Gusfield's lgorithm
22
      Z[0] = 1;
23
```

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```
L = R = 0;
       for (int i=1; i<N; ++i)
26
              int ii = L - (i - L); // i的映射位置
27
              int n = R + 1 - i;
28
              if (i > R)
30
                 {
31
                     Z[i] = match(i, i);
32
                     L = i;
34
                     R = i + Z[i] - 1;
              else if (Z[ii] == n)
37
                 {
                     Z[i] = n + match(i-n, i+n);
                     L = i;
39
                     R = i + Z[i] - 1;
40
41
              else
                 Z[i] = min(Z[ii], n);
43
44
45
       // 尋找最長迴文子字串的長度。
46
      int n = 0, p = 0;
47
      for (int i=0; i<N; ++i)
48
          if (Z[i] > n)
49
              n = Z[p = i];
50
      // 記得去掉薪菜字元。
52
      cout << "最長迴文子字串的長度是" << (n—1) / 2;
53
54
      // 印出最長迴文子字串, 記得別印特殊字元。
55
      for (int i=p-Z[p]+1; i <=p+Z[p]-1; ++i)
56
57
          if (i & 1)
58
              cout << s[i];</pre>
59 }
```

#### 3.2 最大回文正方形

```
1 #include <cstdio>
2 #include <cstdlib>
3 #include <cstring>
4 #include <cmath>
5 #include <iostream>
6 #include <fstream>
7 #include <algorithm>
8 #include <string>
9

10 using namespace std;
11
12 const int MAX = 711;
13 int n, m, tmp, t;
14 string line[MAX], row[MAX], col[MAX];
```

```
int row_z[MAX][MAX * 2], col_z[MAX][MAX * 2];
int row_palinlen[MAX][MAX][2], col_palinlen[MAX][MAX][2];
17 int maxlen, sx, sy, ex, ey;
19 void match(int start, int &z, string &s,int len)
20 {
       while (start - z \ge 0 \&\& start + z < len \&\& s[start - z] == s[start + z])
^{21}
22
23 }
24
25 void calc_z(string s, int *z)
26 {
27
       static int mid, right, len, reflect_i, suply_pos;
28
29
       z[0] = 1;
       len = s.size();
30
31
       mid = right = 0;
32
       for (int i = 1; i < len; ++i) {
          reflect_i = 2 * mid - i;
34
35
           suply_pos = i + z[reflect_i] - 1;
36
          if (i > right) {
37
              match(i, z[i] = 0, s, len);
              mid = i, right = i + z[i] - 1;
38
          } else if (suply_pos == right) {
39
              match(i, z[i] = z[reflect_i], s, len);
40
               mid = i, right = i + z[i] - 1;
41
42
          } else {
               z[i] = min(z[reflect_i], right - i + 1);
43
44
45
46 }
47
48 int main()
49 {
50
       ios::sync_with_stdio(false);
51
52
       cin >> n >> m;
53
       for (int i = 0; i < n; ++i) {
          row[i] = "#";
54
55
           cin >> line[i];
56
           for (int j = 0; j < m; ++j) {
57
            row[i] += line[i][j];
               row[i] += '#';
58
59
          }
           calc_z(row[i], row_z[i]);
60
61
           for (int j = 0, cnt = 1; j < m; ++j, cnt += 2) {
62
              row_palinlen[i][j][1] = row_z[i][cnt] - 1;
63
               row_palinlen[i][j][0] = row_z[i][cnt + 1] - 1;
          }
64
65
66
       for (int i = 0; i < m; ++i) {
           col[i] = "#";
```

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```
for (int j = 0; j < n; ++j) {
                 col[i] += line[j][i];
                 col[i] += "#";
 71
            calc_z(col[i], col_z[i]);
 72
             for (int j = 0, cnt = 1; j < n; ++j, cnt += 2) {
                 col_palinlen[j][i][1] = col_z[i][cnt] - 1;
                 col_palinlen[j][i][0] = col_z[i][cnt + 1] - 1;
 75
 76
 77
 78
        tmp = min(n,m);
 79
        maxlen = 1:
 80
 81
        sx = sy = ex = ey = 0;
        for (int i = 0; i < n; ++i)
 83
 84
            for (int j = 0; j < m; ++j) {
 85
                 int k,c;
                 t = tmp;
                 for (k = 1, c = 0; k \le t; k += 2, c++) {
                     if (i - c < 0 \mid i \mid i + c >= n \mid i \mid j - c < 0 \mid i \mid j + c >= m) break;
                     t = min(t, row_palinlen[i - c][j][1]);
                     t = min(t, row_palinlen[i + c][j][1]);
                    t = min(t, col_palinlen[i][j - c][1]);
                     t = min(t, col_palinlen[i][j + c][1]);
                     if (k > t) break;
                     if (t <= maxlen) break;
                 if (k-2 > maxlen) {
                     c---;
 98
                     maxlen = k - 2;
                     sx = i - c; ex = i + c;
                     sy = j - c; ey = j + c;
102
103
                 t = tmp;
                 for (k = 2, c = 0; k \le t; k += 2, c++) {
104
                     if (i - c < 0 | i + c + 1 >= n | i - c < 0 | i + c + 1 >= m) break;
                     t = min(t, row_palinlen[i - c][j][0]);
106
                    t = min(t, row_palinlen[i + c + 1][j][0]);
                     t = min(t, col_palinlen[i][j - c][0]);
                     t = min(t, col_palinlen[i][j + c + 1][0]);
109
                     if (k > t) break;
                     if (t <= maxlen) break;
                 if (k-2 > maxlen) {
113
114
                     c---;
                     maxlen = k - 2:
115
116
                     sx = i - c; ex = i + c + 1;
                     sy = j - c; ey = j + c + 1;
118
119
        cout << sx + 1 << " " << sy + 1 << " "
```

#### 3.3 KMP

```
next[i] = \max\{len|A[0...len-1] = A的第 i 位向前或后的长度为 len 的串}
  ext[i] = max\{len|A[0...len-1] = B的第 i 位向前或后的长度为 len 的串}
void KMP(char *a, int la, char *b, int lb, int *next, int *ext) {
     -a; -b; -next; -ext;
      for (int i = 2, j = next[1] = 0; i \le la; i++) {
          while (j && a[j + 1] != a[i]) j = next[j]; if (a[j + 1] == a[i]) ++j; next[i] = j;
      } for (int i = 1, j = 0; i \le lb; ++i) {
          while (j && a[j + 1] != b[i]) j = next[j]; if (a[j + 1] == b[i]) ++j; ext[i] = j;
          if (j == la) j = next[j];
9 } void ExKMP(char *a, int la, char *b, int lb, int *next, int *ext) {
      next[0] = la; for (int &j = next[1] = 0; j + 1 < la && a[j] == a[j + 1]; ++j);
      for (int i = 2, k = 1; i < la; ++i) {
11
12
          int p = k + next[k], l = next[i - k]; if (l ;
          else for (int \&j = next[k = i] = max(0, p - i); i + j < la && a[j] == a[i + j]; ++j);
13
      } for (int &j = ext[0] = 0; j < la && j < lb && a[j] == b[j]; ++j);
      for (int i = 1, k = 0; i < lb; ++i) {
          int p = k + ext[k], l = next[i - k]; if (l ;
16
17
          else for (int \&j = ext[k = i] = max(0, p - i); j < la &\& i + j < lb && a[j] == b[i + j]; ++j);
18
19 }
```

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## 3.4 Aho-Corasick 自动机

```
void construct() {
    static tree Q[MAX_NODE]; int head = 0, tail = 0;
    for (root > fail = root, Q[++tail] = root; head < tail; ) {
        tree x = Q[++head];
        if (x > fail > danger) x > danger = true;
        Rep(d, 0, sigma - 1) if (!x > next[d]);
            x > next[d] = (x == root) ? (root) : (x > fail > next[d]);
        else {
            x > next[d] > fail = (x == root) ? (root) : (x > fail > next[d]);
            Q[++tail] = x > next[d];
        }
    }
}
```

# 3.5 后缀自动机

```
1 struct SAM {
2    int in[Maxn * 2 + 1][Sigma], fa[Maxn * 2 + 1], max[Maxn * 2 + 1], tot, last;
3    void init(int n) {
```

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```
tot = last = 0;
           for(int i = 0; i \le 2 * n + 1; ++i)
               memset(in[i], -1, sizeof in[i]), fa[i] = -1;
       void add(int x) {
           int v = last; ++tot, last = tot, max[last] = max[v] + 1;
           while(v != -1 \& in[v][x] == -1) in[v][x] = last, v = fa[v];
11
           if(v == -1) \{ fa[last] = 0; return; \}
           int p = in[v][x];
12
           if(max[p] == max[v] + 1) fa[last] = p;
           else {
15
               int np = ++tot;
               max[np] = max[v] + 1; fa[np] = fa[p], fa[p] = np, fa[last] = np;
               while(v := -1 \&\& in[v][x] == p) in[v][x] = np, v = fa[v];
               memcpy(in[np], in[p], sizeof in[p]);
19 }}};
```

```
3.6 后缀数组-1
待排序的字符串放在 r[0...n-1] 中, 最大值小于 m.
   r[0...n-2] > 0, r[n-1] = 0.
   结果放在 sa[0...n-1].
 1 namespace SuffixArrayDoubling {
       int wa[MAXN], wb[MAXN], wv[MAXN], ws[MAXN];
       int cmp(int *r, int a, int b, int l) {
           return r[a] == r[b] \&\& r[a + 1] == r[b + 1];
       void da(int *r, int *sa, int n, int m) {
           int i, j, p, *x = wa, *y = wb, *t;
           for (i = 0; i < m; i++) ws[i] = 0;
           for (i = 0; i < n; i++) ws[x[i] = r[i]]++;
           for (i = 1; i < m; i++) ws[i] += ws[i-1];
           for (i = n - 1; i \ge 0; i \longrightarrow) sa[-ws[x[i]]] = i;
           for (j = 1, p = 1; p < n; j *= 2, m = p) {
13
               for (p = 0, i = n - j; i < n; i++) y[p++] = i;
14
               for (i = 0; i < n; i++) if (sa[i] >= j) y[p++] = sa[i] - j;
               for (i = 0; i < n; i++) wv[i] = x[y[i]];
               for (i = 0; i < m; i++) ws[i] = 0;
               for (i = 0; i < n; i++) ws[wv[i]]++;
               for (i = 1; i < m; i++) ws[i] += ws[i-1];
               for (i = n - 1; i \ge 0; i \longrightarrow) sa[\longrightarrow ws[wv[i]]] = y[i];
19
               for (t = x, x = y, y = t, p = 1, x[sa[0]] = 0, i = 1; i < n; i++)
                   x[sa[i]] = cmp(y, sa[i-1], sa[i], j) ? p-1 : p++;
22
23
24 }
25 namespace SuffixArrayDC3 { // `r 与 sa 大小需 3 倍
       #define F(x) ((x) / 3 + ((x) % 3 == 1 ? 0 : tb))
       #define G(x) ((x) < tb ? (x) * 3 + 1 : ((x) - tb) * 3 + 2)
       int wa[MAXN], wb[MAXN], wv[MAXN], ws[MAXN];
       int c0(int *r, int a, int b) {
29
           return r[a] == r[b] \&\& r[a + 1] == r[b + 1] \&\& r[a + 2] == r[b + 2];
```

```
int c12(int k, int *r, int a, int b) {
32
           if (k == 2) return r[a] < r[b] || (r[a] == r[b] && c12(1, r, a + 1, b + 1));
                        return r[a] < r[b] \mid | (r[a] == r[b] && wv[a + 1] < wv[b + 1]);
34
35
       void sort(int *r. int *a. int *b. int n. int m) {
           for (int i = 0; i < n; i++) wv[i] = r[a[i]];
37
38
           for (int i = 0; i < m; i++) ws[i] = 0;
           for (int i = 0; i < n; i++) ws[wv[i]]++;
           for (int i = 1; i < m; i++) ws[i] += ws[i-1];
           for (int i = n - 1; i \ge 0; i \longrightarrow) b \sqsubseteq ws[wv[i]]] = a[i];
41
42
43
       void dc3(int *r. int *sa. int n. int m) {
44
           int i, j, *rn = r + n, *san = sa + n, ta = 0, tb = (n + 1) / 3, tbc = 0, p;
45
           r\lceil n\rceil = r\lceil n + 1\rceil = 0;
           for (i = 0; i < n; i++) if (i % 3 != 0) wa[tbc++] = i;
46
47
           sort(r + 2, wa, wb, tbc, m);
           sort(r + 1, wb, wa, tbc, m);
           sort(r, wa, wb, tbc, m);
           for (p = 1, rn[F(wb[0])] = 0, i = 1; i < tbc; i++)
50
51
               rn[F(wb[i])] = c0(r, wb[i-1], wb[i]) ? p-1 : p++;
52
           if (p < tbc) dc3(rn, san, tbc, p):
53
           else for (i = 0; i < tbc; i++) san[rn[i]] = i;
           for (i = 0; i < tbc; i++) if (san[i] < tb) wb[ta++] = san[i] * 3;
54
           if (n \% 3 == 1) wb[ta++] = n-1;
55
           sort(r, wb, wa, ta, m);
56
           for (i = 0; i < tbc; i++) wv[wb[i] = G(san[i])] = i;
57
           for (i = 0, j = 0, p = 0; i < ta && j < tbc; p++)
58
59
               sa[p] = c12(wb[j] % 3, r, wa[i], wb[j]) ? wa[i++] : wb[j++];
60
           for (; i < ta; p++) sa[p] = wa[i++];
           for (; j < tbc; p++) sa[p] = wb[j++];
61
62
       #undef F
63
64
       #undef G
65 }
66 namespace CalcHeight {
       int rank[MAXN], height[MAXN];
68
       void calheight(int *r, int *sa, int n) {
69
           int i, j, k = 0;
           for (i = 1; i \le n; i++) rank[sa[i]] = i;
           for (i = 0; i < n; height[rank[i++]] = k)
               for (k ? k - : 0, j = sa[rank[i] - 1]; r[i + k] == r[j + k]; k++);
72
73
74 }
```

## 3.7 后缀数组-2

```
1 namespace SuffixArrayDoubling {
2    int wa[MAXN], wb[MAXN], wv[MAXN];
3    int cmp(int *r, int a, int b, int l) { return r[a] == r[b] && r[a + l] == r[b + l]; }
4    void da(int *r, int *sa, int n, int m) {
5        int i, j, p, *x = wa, *y = wb, *t;
6        for (i = 0; i < m; i++) ws[i] = 0;</pre>
```

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```
for (i = 0; i < n; i++) ws[x[i] = r[i]]++;
            for (i = 1; i < m; i++) ws[i] += ws[i-1];
            for (i = n - 1; i \ge 0; i \longrightarrow) sa[-ws[x[i]]] = i;
            for (j = 1, p = 1; p < n; j *= 2, m = p) {
                for (p = 0, i = n - j; i < n; i++) y[p++] = i;
                for (i = 0; i < n; i++) if (sa[i] >= j) y[p++] = sa[i] - j;
                for (i = 0; i < n; i++) wv[i] = x[y[i]];
14
                for (i = 0; i < m; i++) ws[i] = 0;
                for (i = 0; i < n; i++) ws[wv[i]]++;
15
16
                for (i = 1; i < m; i++) ws[i] += ws[i-1];
                for (i = n - 1; i \ge 0; i \longrightarrow) sa[\longrightarrow ws[wv[i]]] = y[i];
17
                for (t = x, x = y, y = t, p = 1, x[sa[0]] = 0, i = 1; i < n; i++)
18
19
                    x[sa[i]] = cmp(y, sa[i-1], sa[i], j) ? p-1 : p++;
20 }}}
21 namespace CalcHeight {
        int rank[MAXN], height[MAXN];
23
        void calheight(int *r, int *sa, int n) {
           int i, j, k = 0; for (i = 1; i \le n; i++) rank[sa[i]] = i;
24
25
            for (i = 0; i < n; height[rank[i++]] = k)
                for (k ? k - : 0, j = sa[rank[i] - 1]; r[i + k] == r[j + k]; k++);
26
27 }
       void init(int len)
28
       {
29
            for(int i = 0; i \le len + 10; ++i)
30
                rank[i] = height[i] = 0;
31
32
33 }
34 //Sample
35 int r[MAXN]; char s[MAXN];
36 int main()
37 {
       int len;
       scanf("%s", s);
39
40
       len = strlen(s);
       for(int i = 0; i < len; ++i) r[i] = s[i] - 'a' + 1;
41
       r[len] = 0;
       SuffixArrayDoubling::da(r, sa, n + 1, 30);
43
44
       CalcHeight::calheight(r, sa, n);
45
       //Then the value of sa[0\sim len-1] is 1\sim n, so init RMO carefully(1\sim n not 0\sim n-1)
46
        return 0;
47 }
48 }
```

## 3.8 环串最小表示

```
int minimalRepresentation(int N, char *s) { // s must be double—sized and 0—based
int i, j, k, l; for (i = 0; i < N; ++i) s[i + N] = s[i]; s[N + N] = 0;

for (i = 0, j = 1; j < N; ) {
    for (k = 0; k < N && s[i + k] == s[j + k]; ++k);
    if (k >= N) break; if (s[i + k] < s[j + k]) j += k + 1;
    else l = i + k, i = j, j = max(l, j) + 1;
    return i; // [i, i + N) is the minimal representation
}</pre>
```

## 3.9 回文自动机

```
1 #include <cstdlib>
2 #include <cstdio>
3 #include <cstring>
4 #include <algorithm>
5
6 const int C = 26;
7 const int N = 100000;
8 const int S = N + 2 + C;
10 char string[N + 2];
int s, length[S], suffix[S], go[S][C];
13 int extend(int p, int i)
14 {
15
       while (string[i - 1 - length[p]] != string[i]) {
16
           p = suffix[p];
17
18
       int q = suffix[p];
       while (string[i - 1 - length[q]] != string[i]) {
19
           q = suffix[q];
20
21
22
       int c = string[i] - 'a';
23
       int pp = go[p][c];
       int qq = go[q][c];
25
       if (pp == -1) {
           length[pp = go[p][c] = s ++] = length[p] + 2;
26
27
           suffix[pp] = qq;
28
           memset(go[pp], -1, sizeof(go[pp]));
29
30
       return pp;
31 }
32
33 int main()
34 {
35
       int tests;
36
       scanf("%d", &tests);
       for (int t = 1; t <= tests; ++ t) {
           printf("Case #%d: ", t);
38
           for (int i = 0; i < C + 2; ++ i) {
39
               suffix[i] = 1;
40
41
               length[i] = std::min(i - 1, 1);
               memset(go[i], -1, sizeof(go[i]));
42
43
44
           suffix[0] = suffix[1] = 0;
45
           for (int i = 0; i < C; ++ i) {
               go[0][i] = 2 + i;
46
47
           }
           s = C + 2;
48
49
           string[0] = '#';
           scanf("%s", string + 1);
50
51
           int n = strlen(string + 1);
```

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## 4 图论

#### 4.1 Dominator Tree

```
1 #include <cstdio>
2 #include <cstdlib>
 3 #include <cstring>
 4 #include <iostream>
 5 #include <algorithm>
 6 #include <vector>
 8 using namespace std;
10 const int oo = 1073741819;
11
12 const int Maxn = 200000;
13 const int Maxm = 200000;
14
15 vector<int> g[Maxn];
17 //idom[i] is the dominator of i, node id — 1 based(1 \sim n), n is the source
18 class DominatorTree
19 {
20 public:
21
       int tail[4][Maxm], n, m;
       int Next[4][Maxm], sora[4][Maxm];
       int ss[4], top, w_time;
24
       int rel[Maxn], semi[Maxn], b[Maxn], idom[Maxn], best[Maxn], st[Maxn], pre[Maxn];
       void origin()
25
26
       {
27
           for (int e = 0; e <= 3; e++) ss[e] = n;
           for (int i = 1; i <= n; i++) {
28
               for (int e = 0; e <= 3; e++)
29
                   tail[e][i] = i, Next[e][i] = 0;
               rel[i] = 0, semi[i] = idom[i] = pre[i] = 0, best[i] = i;
31
               b[i] = i;
32
           }
33
34
           rel[0] = oo;
35
       void link(int e, int x, int y)
```

```
++ss[e], Next[e][tail[e][x]] = ss[e], tail[e][x] = ss[e], sora[e][ss[e]] = y, Next[e][ss[e]] = 0;
38
39
       void dfs(int x, int y)
40
41
           ++w_time, rel[x] = w_time;
           st[++top] = x, pre[x] = y;
43
44
           for (int i = x, ne; Next[0][i];) {
               i = Next[0][i], ne = sora[0][i];
45
46
               if (!rel[ne]) dfs(ne, x);
47
           }
48
49
       int find(int x)
50
51
           int y = b[x];
           if (b[x] != x) b[x] = find(b[x]);
52
53
           if (rel[semi[best[y]]]<rel[semi[best[x]]])</pre>
               best[x] = best[y];
54
55
           return b[x];
56
57
       //n — number of vertex, m — number of edges, e — edge set
58
       void init(int _n, int _m, const vector<pair<int, int> > &e)
       {
59
60
           n = _n, m = _m;
           origin();
61
           for (int i = 0; i < m; i++) {
62
              link(0, e[i].first, e[i].second);
63
               link(1, e[i].second, e[i].first);
64
65
66
           w_{time} = 0, top = 0;
           dfs(n, 0);
67
68
69
70
       void work()
71
72
           for (int i = top; i >= 1; i \longrightarrow) {
               int ne = st[i];
73
74
               for (int j = ne, na; Next[1][j];) {
75
                   j = Next[1][j], na = sora[1][j];
                   if (!rel[na]) continue;
76
77
                   int y;
                   if (rel[na]>rel[ne]) {
78
79
                       find(na);
                       y = semi[best[na]];
80
81
                   else y = na;
82
                   if (rel[y]<rel[semi[ne]]) semi[ne] = y;</pre>
83
84
               }
               if (ne != n) link(2, semi[ne], ne);
85
               for (int j = ne, na; Next[2][j];) {
86
                   j = Next[2][j], na = sora[2][j];
87
                   find(na);
88
                   int y = best[na];
```

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```
if (semi[y] == semi[na]) idom[na] = semi[na];
                    else idom[na] = y;
 92
                for (int j = ne, na; Next[0][j];) {
 93
                    j = Next[0][j], na = sora[0][j];
 94
                    if (pre[na] == ne) {
                        na = find(na);
                        b[na] = ne;
 97
100
            for (int i = 2; i <= top; i++) {
101
102
                int ne = st[i];
103
                if (idom[ne] != semi[ne]) idom[ne] = idom[idom[ne]];
                link(3, idom[ne], ne);
105
106
107 }dom;
```

#### 4.2 树 Hash

```
1 #include <cstdio>
 2 #include <cstdlib>
3 #include <cstrina>
 4 #include <cmath>
5 #include <iostream>
 6 #include <algorithm>
 7 #include <vector>
 8 #include <map>
 9 #include <queue>
10
11 using namespace std;
12
13 const int mm = 1051697, p = 1e9 + 9, q = 1e9 + 7;
14 const int N = 100000 + 10;
15 vector<int> vec[N];
int n, size[N], mark[N], deg[N], father[N];
17 long long f[N], g[N], rtp[N], rtq[N];
18 map<pair<long long, long long>, int> mp;
20 struct Node {
21
       int a, b, v;
       Node() {}
       Node(int _a, int _b, int _v) {
23
           a = _a, b = _b, v = _v;
24
25
26
       bool operator < (const Node &rhs) const {
           if (a == rhs.a)
27
               return b < rhs.b;</pre>
28
           return a < rhs.a;</pre>
30
31 };
32
```

```
33 struct HashNode {
34
       int pos;
       long long val1, val2;
35
       HashNode() {}
36
       HashNode(int _pos, long long _val1, long long _val2) {
37
38
           pos = _pos;
           val1 = _val1;
39
40
           val2 = _val2;
41
42
       bool operator < (const HashNode &rhs) const {
           if (val1 == rhs.val1)
43
               return val2 < rhs.val2;</pre>
44
45
           return val1 < rhs.val1:
46
47 };
48
49 void hashwork(int u)
50 {
       vector<Node> data;
       size[u] = 1;
52
53
       for (int i = 0; i < (int)vec[u].size(); ++i) {</pre>
           int v = vec[u][i];
54
55
           hashwork(v);
           data.push_back(Node(f[v], g[v], size[v]));
56
           size[u] += size[v];
57
58
59
       data.push_back(Node(1, 1, size[u]));
       sort(data.begin(), data.end());
60
61
62
       int len = 0;
       f[u] = 1;
63
64
       for (int i = 0; i < (int)data.size(); ++i) {</pre>
           f[u] = ((f[u] * data[i].a) % p * rtp[len]) % p;
65
66
           g[u] = ((g[u] * data[i].b) % q + rtq[len]) % q;
           len += data[i].v;
67
68
69 }
70
71 int main()
72 {
       ios::sync_with_stdio(false);
       rtp[0] = rtq[0] = 1;
74
75
       for (int i = 1; i < N; ++i) {
           rtp[i] = (rtp[i - 1] * mm) % p;
76
77
           rtq[i] = (rtq[i-1] * mm) % q;
78
79
80
       queue<int> que;
81
       cin >> n;
       for (int v = 2; v <= n; ++v) {
83
           int u;
84
           cin >> u;
85
           vec[u].push_back(v);
```

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```
father[v] = u;
            deg[u]++;
 88
        memset(size, 0, sizeof(size));
 89
 90
        memset(f, 0, sizeof(f));
        memset(g, 0, sizeof(g));
 91
        for (int i = 1; i <= n; ++i)
            if (deg[i] == 0)
 93
 94
                que.push(i);
 95
        while (!que.empty()) {
            int u = que.front();
            //cout << u << endl;</pre>
            que.pop();
            deg[father[u]]—;
            if (deg[father[u]] == 0) que.push(father[u]);
            vector<Node> data;
103
            size[u] = 1;
104
            for (int i = 0; i < (int)vec[u].size(); ++i) {
                int v = vec[u][i];
106
                //hashwork(v);
                data.push_back(Node(f[v], g[v], size[v]));
108
                size[u] += size[v];
110
            data.push_back(Node(1, 1, size[u]));
112
            sort(data.begin(), data.end());
113
114
            int len = 0;
            f[u] = 1;
116
            for (int i = 0; i < (int)data.size(); ++i) {</pre>
117
                f[u] = ((f[u] * data[i].a) % p * rtp[len]) % p;
                g[u] = ((g[u] * data[i].b) % q + rtq[len]) % q;
118
                len += data[i].v;
119
120
121
122
        //hashwork(1);
123
124
125
          vector<HashNode> ans;
          for (int i = 1; i <= n; ++i) {
126
127
          ans.push_back(HashNode(i, f[i], g[i]));
128
          sort(ans.begin(), ans.end());
129
          int tot = 0;
131
          for (int i = 0, j; i < (int)ans.size(); i = j) {
132
          for (j = i; j < (int)ans.size() \&& (ans[j].val1 == ans[i].val1 && ans[j].val2 == ans[i].val2); ++j)
133
134
                mark[ans[j].pos] = tot;
        */
136
137
        int tot = 0;
138
        for (int i = 1; i <= n; ++i) {
```

```
139
            pair<long long, long long> pr = make_pair(f[i], g[i]);
140
            if (mp.count(pr) == 0) {
141
                mp[pr] = ++tot;
                mark[i] = tot;
142
143
            } else {
                mark[i] = mp[pr];
144
145
146
147
        for (int i = 1; i <= n; ++i) {
148
            cout << mark[i];
149
            if (i == n) cout << endl;</pre>
            else cout << " ";
152
        return 0;
153 }
```

## 4.3 带花树

```
1 namespace Blossom {
       int n, head, tail, S, T, lca;
       int match[MAXN], O[MAXN], pred[MAXN], label[MAXN], ing[MAXN], inb[MAXN];
       vector<int> link[MAXN];
       inline void push(int x) { Q[tail++] = x; inq[x] = true; }
       int findCommonAncestor(int x, int y) {
           static bool inPath[MAXN]; for (int i = 0; i < n; ++i) inPath[i] = 0;
           for (;; x = pred[match[x]]) { x = label[x]; inPath[x] = true; if (x == S) break; }
           for ( ; ; y = pred[ match[y] ]) { y = label[y]; if (inPath[y]) break; } return y;
10
11
       void resetTrace(int x, int lca) {
           while (label[x] != lca) { int y = match[x]; inb[ label[x] ] = inb[ label[y] ] = true;
12
               x = pred[y]; if (label[x] != lca) pred[x] = y; }}
13
       void blossomContract(int x, int y) {
14
15
           lca = findCommonAncestor(x, y);
           Foru(i, 0, n) inb[i] = 0; resetTrace(x, lca); resetTrace(y, lca);
16
           if (label[x] != lca) pred[x] = y; if (label[y] != lca) pred[y] = x;
           Foru(i, 0, n) if (inb[ label[i] ]) { label[i] = lca; if (!inq[i]) push(i); }
18
19
20
       bool findAugmentingPath() {
           Foru(i, 0, n) pred[i] = -1, label[i] = i, inq[i] = 0;
21
22
           int x, y, z; head = tail = 0;
           for (push(S); head < tail; ) for (int i = (int)link[x = Q[head++]].size() - 1; i >= 0; --i) {
23
               y = link[x][i]; if (label[x] == label[y] | | x == match[y]) continue;
24
               if (y == S \mid | (match[y] >= 0 \&\& pred[match[y]] >= 0)) blossomContract(x, y);
25
26
               else if (pred[y] == -1) {
                   pred[y] = x; if (match[y] >= 0) push(match[y]);
27
28
29
                       for (x = y; x >= 0; x = z) {
30
                       y = pred[x], z = match[y]; match[x] = y, match[y] = x;
                   } return true; }}} return false;
31
32
33
       int findMaxMatching() {
34
           int ans = 0; Foru(i, 0, n) match[i] = -1;
35
           for (S = 0; S < n; ++S) if (match[S] == -1) if (findAugmentingPath()) ++ans;
```

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```
4 TZN#
```

return ans;

#### 4.4 KM

37

38 }

```
int N, Tcnt, w[MAXN][MAXN], slack[MAXN];
___ int_lx[MAXN], linkx[MAXN], visy[MAXN], ly[MAXN], linky[MAXN], visx[MAXN]; // `初值全为0
 3 bool DFS(int x) { visx[x] = Tcnt;
       Rep(y, 1, N) if(visy[y] != Tcnt) { int t = lx[x] + ly[y] - w[x][y];
           if (t == 0) \{ visy[y] = Tcnt;
               if (!linky[y] | | DFS(linky[y])) \{ linkx[x] = y; linky[y] = x; return true; \}
           } else cMin(slack[y], t);
       } return false;
9 } void KM() {
       Tcnt = 0; Rep(x, 1, N) Rep(y, 1, N) cMax(lx[x], w[x][y]);
       Rep(S, 1, N) \{ Rep(i, 1, N) \ slack[i] = INF; 
11
           for (++Tcnt; !DFS(S); ++Tcnt) { int d = INF;
12
               Rep(y, 1, N) if(visy[y] != Tcnt) cMin(d, slack[y]);
13
               Rep(x, 1, N) if(visx[x] == Tcnt) lx[x] -= d;
14
15
               Rep(y, 1, N) if(visy[y] == Tcnt) ly[y] += d; else slack[y] -= d;
17
18 }
```

## 4.5 双连通分量

```
1 #include <iostream>
2 #include <cstdio>
3 #include <cstring>
 4 #include <cstdlib>
5 #include <vector>
7 using namespace std;
9 const int MAXN = 100000 + 10;
int dfn[MAXN], low[MAXN], bccno[MAXN], dfn_clock, bcc_cnt, Top;
12 vector <int> G[MAXN], bcc[MAXN];
13 pair <int, int> stk[MAXN];
14 bool iscut[MAXN];
15 int n, m;
17 void dfs(int p, int fa) {
       low[p] = dfn[p] = ++dfn_clock;
       int child = 0;
       for (int i = 0; i < G[p].size(); ++i) {
20
           int v = G[p][i];
21
22
           if (!dfn[v]) {
               stk[++Top] = make_pair(p, v);
23
24
               dfs(v, p);
```

```
25
               child++;
               low[p] = min(low[p], low[v]);
26
27
               if (low[v] >= dfn[p]) {
                   iscut[p] = 1;
28
                   ++bcc_cnt;
29
                   bcc[bcc_cnt].clear();
                   for (;;) {
31
                       pair <int, int> x = stk[Top];
32
                       —Тор;
33
34
                       if (bccno[x.first] != bcc_cnt) {
35
                           bccno[x.first] = bcc_cnt;
                           bcc[bcc_cnt].push_back(x.first);
36
37
38
                       if (bccno[x.second] != bcc_cnt) {
39
                           bccno[x.second] = bcc_cnt;
                           bcc[bcc_cnt].push_back(x.second);
40
41
                       if (x.first == p && x.second == v)
42
43
                           break;
44
45
46
47
           else
               if (dfn[v] < dfn[p] && v != fa) {</pre>
48
                   stk[++Top] = make_pair(p, v);
49
                   low[p] = min(low[p], dfn[v]);
50
               }
51
52
       if (fa < 0 && child == 1) iscut[p] = 0;
53
54 }
55
56 void find_bcc(int n) {
       for (int i = 1; i \le n; ++i) dfn[i] = 0;
       for (int i = 1; i \le n; ++i) iscut[i] = 0;
       for (int i = 1; i \le n; ++i) bccno[i] = 0;
59
       dfn_clock = bcc_cnt = 0;
       for (int i = 1; i <= n; ++i)
61
62
           if (!dfn[i])
63
               dfs(i, -1);
64 }
65
66 int main() {
67
       scanf("%d%d", &n, &m);
       for (int a, b, i = 1; i <= m; ++i) {
68
           scanf("%d%d", &a, &b);
           G[a].push_back(b);
70
           G[b].push_back(a);
71
72
73
       find_bcc(n);
74
75
       return 0;
76 }
```

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## 4.6 强连通分量

```
1 #include <iostream>
 2 #include <cstdio>
 3 #include <cstring>
 4 #include <cstdlib>
 5 #include <vector>
 6 #include <algorithm>
 8 using namespace std;
10 const int MAXN = 100000 + 10;
12 vector <int> G[MAXN];
14 int dfn[MAXN], low[MAXN], stk[MAXN], Top, scc_cnt, sccno[MAXN], dfn_clock;
15
16 void dfs(int p) {
       dfn[p] = low[p] = ++dfn_clock;
       stk[++Top] = p;
       for (int i = 0; i < (int)G[p].size(); ++i) {
           int v = G[p][i];
20
           if (!dfn[v]) {
21
               dfs(v);
23
               low[p] = min(low[p], low[v]);
           }
24
25
           else if (!sccno[v])
               low[p] = min(low[p], dfn[v]);
27
       if (low[p] == dfn[p]) {
28
           scc_cnt++;
29
           for (;;) {
30
               int x = stk[Top];
              —Top;
               sccno[x] = scc_cnt;
33
               if (x == p) break;
34
36
37 }
39 void find_scc(int n) {
       dfn_clock = scc_cnt = 0;
       for (int i = 1; i <= n; ++i) sccno[i] = 0;
       for (int i = 1; i \le n; ++i) dfn[i] = low[i] = 0;
       for (int i = 1; i <= n; ++i)
43
           if (!dfn[i])
               dfs(i);
45
```

# 4.7 2-SAT 与 Kosaraju

注意 Kosaraju 需要建反图

```
1 namespace SCC {
       int code[MAXN * 2], seq[MAXN * 2], sCnt;
       void DFS_1(int x) { code[x] = 1;
            for (edge e(fir[x]); e; e = e\rightarrownext) if (code[e\rightarrowto] == -1) DFS_1(e\rightarrowto);
            seq[++sCnt] = x;
       } void DFS_2(int x) { code[x] = sCnt;
            for (edge e(fir2[x]); e; e = e\rightarrownext) if (code[e\rightarrowto] == -1) DFS_2(e\rightarrowto); }
       void SCC(int N) {
           sCnt = 0; for (int i = 1; i <= N; ++i) code[i] = -1;
10
            for (int i = 1; i \leftarrow N; ++i) if (code[i] == -1) DFS_1(i);
11
            sCnt = 0; for (int i = 1; i <= N; ++i) code[i] = -1;
            for (int i = N; i >= 1; —i) if (code[seq[i]] == -1) {
12
13
                ++sCnt; DFS_2(seq[i]); }
14
15 \}// true - 2i - 1
16 // false — 2i
17 bool TwoSat() { SCC::SCC(N + N);
       // if code[2i - 1] = code[2i]: no solution
       // if code[2i - 1] > code[2i]: i selected. else i not selected
20 }
```

## 4.8 全局最小割 Stoer-Wagner

```
int minCut(int N, int G[MAXN][MAXN]) { // 0—based
       static int weight[MAXN], used[MAXN]; int ans = INT_MAX;
       while (N > 1) {
           for (int i = 0; i < N; ++i) used[i] = false; used[0] = true;
           for (int i = 0; i < N; ++i) weight[i] = G[i][0];
           int S = -1, T = 0;
           for (int _r = 2; _r \leftarrow N; _{++}r) { //N-1 selections
              int x = -1:
               for (int i = 0; i < N; ++i) if (!used[i])
                   if (x == -1 \mid l \mid weight[i] > weight[x]) x = i;
11
               for (int i = 0; i < N; ++i) weight[i] += G[x][i];
               S = T; T = x; used[x] = true;
12
13
          } ans = min(ans, weight[T]);
           for (int i = 0; i < N; ++i) G[i][S] += G[i][T], G[S][i] += G[i][T];
14
15
           G[S][S] = 0; --N;
           for (int i = 0; i \le N; ++i) swap(G[i][T], G[i][N]);
16
17
           for (int i = 0; i < N; ++i) swap(G[T][i], G[N][i]);
18
       } return ans;
19 }
```

# 4.9 Hopcroft-Karp

```
1 int N, M, level[MAXN], matchX[MAXN], matchY[MAXN];
2 bool used[MAXN];
3 bool DFS(int x) {
4     used[x] = true; for (edge e(fir[x]); e; e = e = next) {
5         int y = e => to, z = matchY[y];
6     if (z == -1 || (!used[z] && level[x] < level[z] && DFS(z))) {</pre>
```

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```
matchX[x] = y; matchY[y] = x; return true;
       } return false;
10 }
11 int maxMatch() {
        for (int i = 0; i < N; ++i) used[i] = false;
        for (int i = 0; i < N; ++i) matchX[i] = -1;
14
       for (int i = 0; i < M; ++i) matchY[i] = -1;
       for (int i = 0; i < N; ++i) level[i] = -1;
15
16
       int match = 0, d;
17
       for ( ; ; match += d) {
           static int Q[MAXN * 2 + 1];
18
19
           int head = 0. tail = d = 0:
           for (int x = 0; x < N; ++x) level[x] = -1;
            for (int x = 0; x < N; ++x) if (matchX[x] == -1)
                level[x] = 0, Q[++tail] = x;
22
23
           while (head < tail)
                for (edge e(fir[x = Q[++head]]); e; e = e \rightarrow next) {
24
                    int y = e \rightarrow to, z = matchY[y];
                   if (z != -1 \&\& level[z] < 0) level[z] = level[x] + 1, Q[++tail] = z;
           for (int x = 0; x < N; ++x) used[x] = false;
28
            for (int x = 0; x < N; ++x) if (matchX[x] == -1) if (DFS(x)) ++d;
           if (d == 0) break;
31
       } return match;
32 }
```

#### 4.10 欧拉路

#### 4.11 稳定婚姻

```
1  namespace StableMatching {
2    int pairM[MAXN], pairW[MAXN], p[MAXN];
3    // init: pairM[0...n - 1] = pairW[0...n - 1] = -1, p[0...n - 1] = 0
4    void stableMatching(int n, int orderM[MAXN][MAXN], int preferW[MAXN][MAXN]) {
5        for (int i = 0; i < n; i++) while (pairM[i] < 0) {
6            int w = orderM[i][p[i]++], m = pairW[w];
7            if (m == -1) pairM[i] = w, pairW[w] = i;
8            else if (preferW[w][i] < preferW[w][m])</pre>
```

## 4.12 最大团搜索

```
1 namespace MaxClique { // 1—based
       int g[MAXN][MAXN], len[MAXN], list[MAXN][MAXN], mc[MAXN], ans, found;
       void DFS(int size) {
           if (len[size] == 0) { if (size > ans) ans = size, found = true; return; }
           for (int k = 0; k < len[size] && !found; ++k) {
               if (size + len[size] - k <= ans) break;</pre>
               int i = list[size][k]; if (size + mc[i] <= ans) break;</pre>
               for (int j = k + 1, len[size + 1] = 0; j < len[size]; ++j) if (q[i][list[size][j]])
                   list[size + 1][len[size + 1]++] = list[size][j];
               DFS(size + 1);
11
           }
12
       int work(int n) {
13
14
           mc[n] = ans = 1; for (int i = n - 1; i; --i) { found = false; len[1] = 0;
15
               for (int j = i + 1; j \le n; ++j) if (g[i][j]) list[1][len[1]++] = j;
16
               DFS(1); mc[i] = ans;
17
           } return ans;
18
19 }
```

## 4.13 极大团计数

```
1 namespace MaxCliqueCounting {
       int n, ans;
       int ne[MAXN], ce[MAXN];
       int g[MAXN][MAXN], list[MAXN][MAXN];
       void dfs(int size) {
           int i, j, k, t, cnt, best = 0;
           bool bb;
           if (ne[size] == ce[size]) {
               if (ce[size] == 0)
                   ++ans:
11
               return;
12
13
           for (t = 0, i = 1; i \le ne[size]; ++i) {
               for (cnt = 0, j = ne[size] + 1; j <= ce[size]; ++j)
14
15
                   if (!g[list[size][i]][list[size][j]])
16
                       ++cnt;
               if (t == 0 || cnt < best)
17
                   t = i, best = cnt;
18
19
20
           if (t && best <= 0)
21
               return;
22
           for (k = ne[size] + 1; k \leftarrow ce[size]; ++k) {
```

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```
if (t > 0) {
                   for (i = k; i \leftarrow ce[size]; ++i)
                       if (!g[list[size][t]][list[size][i]])
                           break:
                   swap(list[size][k], list[size][i]);
               i = list[size][k];
               ne[size + 1] = ce[size + 1] = 0;
31
               for (j = 1; j < k; ++j)
                   if (a[i][list[size][i]])
32
                       list[size + 1][++ne[size + 1]] = list[size][j];
               for (ce[size + 1] = ne[size + 1], j = k + 1; j <= ce[size]; ++j)
                   if (g[i][list[size][j]])
                       list[size + 1][++ce[size + 1]] = list[size][j];
               dfs(size + 1);
               ++ne[size];
               -best;
               for (j = k + 1, cnt = 0; j \le ce[size]; ++j)
                   if (!g[i][list[size][j]])
                       ++cnt;
               if (t == 0 || cnt < best)
43
                   t = k, best = cnt;
44
               if (t && best <= 0)
                   break;
47
48
49
       void work() {
           int i;
           ne[0] = 0;
51
52
           ce[0] = 0;
53
           for (i = 1; i \le n; ++i)
54
               list[0][++ce[0]] = i;
           ans = 0;
           dfs(0);
56
57
```

#### 4.14 最小树形图

```
1 namespace EdmondsAlgorithm { // O(ElogE + V^2) !!! 0-based !!!
          struct enode { int from, c, key, delta, dep; enode *ch[2], *next;
         } ebase[maxm], *etop, *fir[maxn], nil, *null, *inEdge[maxn], *chs[maxn];
          typedef enode *edge; typedef enode *tree;
          int n, m, setFa[maxn], deg[maxn], que[maxn];
         inline void pushDown(tree x) { if (x\rightarrow delta) {
               x\rightarrow ch[0]\rightarrow key += x\rightarrow delta; x\rightarrow ch[0]\rightarrow delta += x\rightarrow delta;
               x\rightarrow ch[1]\rightarrow key += x\rightarrow delta; x\rightarrow ch[1]\rightarrow delta += x\rightarrow delta; x\rightarrow delta = 0;
         }}
          tree merge(tree x, tree y) {
10
               if (x == null) return y; if (y == null) return x;
11
12
               if (x\rightarrow key > y\rightarrow key) swap(x, y); pushDown(x); x\rightarrow ch[1] = merge(x\rightarrow ch[1], y);
13
               if (x\rightarrow ch[0]\rightarrow dep < x\rightarrow ch[1]\rightarrow dep) swap(x\rightarrow ch[0], x\rightarrow ch[1]);
14
               x\rightarrow dep = x\rightarrow ch[1] \rightarrow dep + 1; return x;
```

```
15
16
       void addEdge(int u, int v, int w) {
17
            etop \rightarrow from = u; etop \rightarrow c = etop \rightarrow key = w; etop \rightarrow delta = etop \rightarrow dep = 0;
            etop->next = fir[v]; etop->ch[0] = etop->ch[1] = null;
18
            fir[v] = etop; inEdge[v] = merge(inEdge[v], etop++);
19
20
21
        void deleteMin(tree &r) { pushDown(r); r = merge(r \rightarrow ch[0], r \rightarrow ch[1]); }
        int findSet(int x) { return setFa[x] == x ? x : setFa[x] = findSet(setFa[x]); }
22
23
        void clear(int V, int E) {
24
            null = &nil; null \rightarrow ch[0] = null \rightarrow ch[1] = null; null \rightarrow dep = -1;
25
            n = V; m = E; etop = ebase; Foru(i, 0, V) fir[i] = NULL; Foru(i, 0, V) inEdge[i] = null;
26
27
       int solve(int root) { int res = 0, head, tail;
28
            for (int i = 0; i < n; ++i) setFa[i] = i;
            for ( ; ; ) { memset(deg, 0, sizeof(int) * n); chs[root] = inEdge[root];
29
                for (int i = 0; i < n; ++i) if (i != root \&\& setFa[i] == i) {
30
                    while (findSet(inEdge[i]->from) == findSet(i)) deleteMin(inEdge[i]);
31
                    ++deg[ findSet((chs[i] = inEdge[i]) >> from) ];
32
                for (int i = head = tail = 0; i < n; ++i)
34
                    if (i != root && setFa[i] == i && deg[i] == 0) que[tail++] = i;
35
                while (head < tail) {</pre>
36
37
                    int x = findSet(chs[que[head++]]—>from);
                    if (--deg[x] == 0) que[tail++] = x;
38
               } bool found = false;
39
                for (int i = 0; i < n; ++i) if (i != root && setFa[i] == i && deg[i] > 0) {
40
                    int j = i; tree temp = null; found = true;
41
                    do {setFa[j = findSet(chs[j]->from)] = i;
                        deleteMin(inEdge[j]); res += chs[j]->key;
43
                        inEdge[j]->key -= chs[j]->key; inEdge[j]->delta -= chs[j]->key;
44
45
                        temp = merge(temp, inEdge[j]);
                    } while (j != i); inEdge[i] = temp;
                } if (!found) break;
47
           } for (int i = 0; i < n; ++ i) if (i != root \&\& setFa[i] == i) res += chs[i] \rightarrow key;
48
49
            return res:
50
51 }
52 namespace ChuLiu { // O(V ^ 3) !!! 1—based !!!
       int n, used[maxn], pass[maxn], eg[maxn], more, que[maxn], g[maxn][maxn];
53
        void combine(int id, int &sum) { int tot = 0, from, i, j, k;
54
55
            for ( ; id != 0 && !pass[id]; id = eg[id]) que[tot++] = id, pass[id] = 1;
56
            for (from = 0; from < tot && que[from] != id; from++);</pre>
57
            if (from == tot) return; more = 1;
58
            for (i = from; i < tot; i++) {
59
                sum += g[eg[que[i]]][que[i]]; if (i == from) continue;
                for (j = used[que[i]] = 1; j \ll n; j++) if (!used[j])
60
61
                    if (g[que[i]][j] < g[id][j]) g[id][j] = g[que[i]][j];</pre>
62
63
            for (i = 1; i <= n; i++) if (!used[i] && i != id)
                for (j = from; j < tot; j++) {
64
65
                    k = que[j]; if (g[i][id] > g[i][k] - g[eg[k]][k])
                    g[i][id] = g[i][k] - g[eg[k]][k];
66
               }
```

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```
void clear(int V) { n = V; Rep(i, 1, V) Rep(j, 1, V) g[i][j] = inf; }
       int solve(int root) {
70
           int i, j, k, sum = 0; memset(used, 0, sizeof(int) * (n + 1));
71
           for (more = 1; more; ) {
               more = 0; memset(eg, 0, sizeof(int) * (n + 1));
               for (i = 1; i <= n; i++) if (!used[i] && i != root) {
                   for (j = 1, k = 0; j \le n; j++) if (!used[j] \&\& i != j)
75
                       if (k == 0 | | g[j][i] < g[k][i]) | k = j;
                   eq[i] = k;
               } memset(pass, 0, sizeof(int) * (n + 1));
               for (i = 1; i <= n; i++) if (!used[i] && !pass[i] && i != root)
                   combine(i, sum):
           for (i = 1; i <= n; i++) if (!used[i] && i != root) sum += g[eg[i]][i];</pre>
           return sum;
83
84 }
```

#### 4.15 离线动态最小生成树

 $O(Qlog^2Q)$ . (qx[i],qy[i]) 表示将编号为 qx[i] 的边的权值改为 qy[i], 删除一条边相当于将其权值改为  $\infty$ , 加入一条 边相当于将其权值从  $\infty$  变成某个值.

```
1 const int maxn = 100000 + 5;
 2 \text{ const int maxm} = 10000000 + 5:
 3 \text{ const int maxq} = 10000000 + 5;
 4 const int qsize = maxm + 3 * maxq;
 5 int n, m, Q, x[qsize], y[qsize], z[qsize], qx[maxq], qy[maxq], a[maxn], *tz;
6 int kx[maxn], ky[maxn], kt, vd[maxn], id[maxm], app[maxm];
7 bool extra[maxm];
 8 void init() {
       scanf("%d%d", &n, &m); for (int i = 0; i < m; i++) scanf("%d%d%d", x + i, y + i, z + i);
       scanf("%d", \&Q); for (int i = 0; i < 0; i++) { <math>scanf("%d%d", qx + i, qy + i); qx[i]—}; }
11 }
12 int find(int x) {
       int root = x, next; while (a[root]) root = a[root];
       while ((next = a[x]) != 0) a[x] = root, x = next; return root;
15 }
16 inline bool cmp(const int &a, const int &b) { return tz[a] < tz[b]; }</pre>
17 void solve(int *qx, int *qy, int Q, int n, int *x, int *y, int *z, int m, long long ans) {
       int ri, rj;
       if (0 == 1) {
            for (int i = 1; i \le n; i++) a[i] = 0; z[qx[0]] = qy[0];
           for (int i = 0; i < m; i++) id[i] = i;
           tz = z; sort(id, id + m, cmp);
           for (int i = 0; i < m; i++) {
23
               ri = find(x[id[i]]); rj = find(y[id[i]]);
24
                if (ri != rj) ans += z[id[i]], a[ri] = rj;
           } printf("%I64d\n", ans);
           return;
       \} int tm = kt = 0, n2 = 0, m2 = 0;
       for (int i = 1; i \le n; i++) a[i] = 0;
       for (int i = 0; i < 0; i++) {
           ri = find(x[qx[i]]); rj = find(y[qx[i]]); if (ri != rj) a[ri] = rj;
```

```
for (int i = 0; i < m; i++) extra[i] = true;
       for (int i = 0; i < 0; i++) extra[qx[i]] = false;
       for (int i = 0; i < m; i++) if (extra[i]) id[tm++] = i;
       tz = z; sort(id, id + tm, cmp);
       for (int i = 0; i < tm; i++) {
38
          ri = find(x[id[i]]); rj = find(y[id[i]]);
39
           if (ri != rj)
               a[ri] = rj, ans += z[id[i]], kx[kt] = x[id[i]], ky[kt] = y[id[i]], kt++;
40
41
       for (int i = 1; i \le n; i++) a[i] = 0;
43
       for (int i = 0; i < kt; i++) a[find(kx[i])] = find(ky[i]);
       for (int i = 1; i \le n; i++) if (a[i] == 0) vd[i] = ++n2;
       for (int i = 1; i \le n; i++) if (a[i] != 0) vd[i] = vd[find(i)];
       int *Nx = x + m, *Ny = y + m, *Nz = z + m;
       for (int i = 0; i < m; i++) app[i] = -1;
48
       for (int i = 0; i < 0; i++)
           if (app[qx[i]] == -1)
49
               Nx[m2] = vd[x[qx[i]]], Ny[m2] = vd[y[qx[i]]], Nz[m2] = z[qx[i]], app[qx[i]] = m2, m2++;
       for (int i = 0; i < 0; i++) {
51
           z[qx[i]] = qy[i];
           qx[i] = app[qx[i]];
54
       for (int i = 1; i \le n2; i++) a[i] = 0;
55
56
       for (int i = 0; i < tm; i++) {
          ri = find(vd[x[id[i]]]); rj = find(vd[y[id[i]]]);
57
               a[ri] = rj, Nx[m2] = vd[x[id[i]]], Ny[m2] = vd[y[id[i]]], Nz[m2] = z[id[i]], m2++;
59
60
61
       int mid = Q / 2;
       solve(qx, qy, mid, n2, Nx, Ny, Nz, m2, ans);
       solve(qx + mid, qy + mid, Q - mid, n2, Nx, Ny, Nz, m2, ans);
64 }
65 void work() { if (Q) solve(qx, qy, Q, n, x, y, z, m, 0); }
66 int main() { init(); work(); return 0; }
```

## 4.16 弦图

- 任何一个弦图都至少有一个单纯点, 不是完全图的弦图至少有两个不相邻的单纯点.
- 设第 i 个点在弦图的完美消除序列第 p(i) 个. 令  $N(v)=\{w|w \neg v$ 相邻且 $p(w)>p(v)\}$  弦图的极大团一定是  $v\cup N(v)$  的形式.
- 弦图最多有 n 个极大团。
- 设 next(v) 表示 N(v) 中最前的点. 令 w\* 表示所有满足  $A \in B$  的 w 中最后的一个点. 判断  $v \cup N(v)$  是否为极大团, 只需判断是否存在一个 w, 满足 Next(w) = v 且  $|N(v)| + 1 \le |N(w)|$  即可.
- 最小染色: 完美消除序列从后往前依次给每个点染色, 给每个点染上可以染的最小的颜色. (团数 = 色数)
- 最大独立集:完美消除序列从前往后能选就选
- 最小团覆盖: 设最大独立集为  $\{p_1,p_2,\ldots,p_t\}$ , 则  $\{p_1\cup N(p_1),\ldots,p_t\cup N(p_t)\}$  为最小团覆盖. (最大独立集数 = 最小团覆盖数)

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```
1 class Chordal { // 1—Based, G is the Graph, must be sorted before call Check_Chordal
 2 public: // Construct will sort it automatically
       int v[Maxn], id[Maxn]; bool inseq[Maxn]; priority_queue<pair<int, int> > pq;
       vector<int> Construct_Perfect_Elimination_Sequence(vector<int> *G, int n) { // O(m + nlogn)
           vector<int> seq(n + 1, 0);
                                                                                                                            18
           for (int i = 0; i <= n; ++i) inseq[i] = false, sort(G[i].begin(), G[i].end()), v[i] = 0;
           int cur = n; pair<int, int> Mx; while(!pq.empty()) pq.pop(); pq.push(make_pair(0, 1));
                                                                                                                            20
           for (int i = n; i >= 1; --i) {
               while (!pq.empty() && (Mx = pq.top(), inseq[Mx.second] || Mx.first != v[Mx.second])) pq.pop();
               id[Mx.second] = cur;
11
               int x = seq[cur—] = Mx.second, sz = (int)G[Mx.second].size(); inseq[x] = true;
               for (int j = 0; j < sz; ++j) {
                                                                                                                            25
12
13
                   int y = G[x][j]; if(!inseq[y]) pq.push(make_pair(++v[y], y));
                                                                                                                            26
14
                                                                                                                            27
           } return seq;
                                                                                                                            28
16
                                                                                                                            29
17
       bool Check_Chordal(vector<int> *G, vector<int> &seq, int n) { // O(n + mlogn), plz gen seq first
                                                                                                                            30
18
           bool isChordal = true;
                                                                                                                            31
           for (int i = n - 1; i >= 1 && isChordal; --i) {
19
                                                                                                                            32
               int x = seq[i], sz, y = -1;
20
                                                                                                                            33
               if ((sz = (int)G[x].size()) == 0) continue;
21
                                                                                                                            34
               for(int j = 0; j < sz; ++j) {
22
                                                                                                                            35
                   if (id[G[x][j]] < i) continue;</pre>
                                                                                                                            36
                   if (y == -1 || id[y] > id[G[x][j]]) y = G[x][j];
24
                                                                                                                            37
25
               } if (y == -1) continue;
                                                                                                                            38
               for (int j = 0; j < sz; ++j) {
26
                                                                                                                            39
                   int y1 = G[x][j]; if (id[y1] < i) continue;
27
                                                                                                                             40
                   if (y1 == y \mid l \text{ binary_search}(G[y].begin(), G[y].end(), y1)) continue;
                                                                                                                             41
                   isChordal = false; break;
                                                                                                                            42
                                                                                                                            43
31
           } return isChordal;
                                                                                                                            44
32
33 };
                                                                                                                             46
                                                                                                                            47
                                                                                                                             48
```

# 4.17 K 短路 (允许重复)

```
1 #define for_each(it, v) for (vector<Edge*>::iterator it = (v).begin(); it != (v).end(); ++it)
2 const int MAX_N = 10000, MAX_M = 50000, MAX_K = 10000, INF = 10000000000;
3 struct Edge { int from, to, weight; };
4 struct HeapNode { Edge* edge; int depth; HeapNode* child[4]; }; // child[0..1] for heap G, child[2..3] for heap
         out edae
6 int n, m, k, s, t; Edge* edge[MAX_M];
7 int dist[MAX_N]; Edge* prev[MAX_N];
8 vector<Edge*> graph[MAX_N]; vector<Edge*> graphR[MAX_N];
9 HeapNode* nullNode; HeapNode* heapTop[MAX_N];
10
11 HeapNode* createHeap(HeapNode* curNode, HeapNode* newNode) {
      if (curNode == nullNode) return newNode; HeapNode* rootNode = new HeapNode;
       memcpy(rootNode, curNode, sizeof(HeapNode));
13
14
       if (newNode->edge->weight < curNode->edge->weight) {
           rootNode->edge = newNode->edge; rootNode->child[2] = newNode->child[2]; rootNode->child[3] = newNode->
15
         child[3];
```

```
newNode->edge = curNode->edge; newNode->child[2] = curNode->child[2]; newNode->child[3] = curNode->child
                 } if (rootNode\rightarrow child[0]\rightarrow depth < rootNode\rightarrow child[1]\rightarrow depth) rootNode\rightarrow child[0] = createHeap(rootNode\rightarrow child[0]\rightarrow depth) rootNode\rightarrow child[0]\rightarrow depth < rootNode\rightarrow child[0]\rightarrow depth </rootNode\rightarrow child[0]\rightarrow depth < rootNode\rightarrow child[0]\rightarrow child[0]
                  else rootNode->child[1] = createHeap(rootNode->child[1], newNode);
                  rootNode->depth = max(rootNode->child[0]->depth, rootNode->child[1]->depth) + 1;
                  return rootNode;
21 }
22 bool heapNodeMoreThan(HeapNode* node1, HeapNode* node2) { return node1—sedge—sweight > node2—sedge—sweight; }
24 int main() {
                 scanf("%d%d%d", &n, &m, &k); scanf("%d%d", &s, &t); s—, t—;
                 while (m—) { Edge* newEdge = new Edge;
                           int i, j, w; scanf("%d%d%d", &i, &j, &w);
                           i---, j---; newEdge->from = i; newEdge->to = j; newEdge->weight = w;
                            graph[i].push_back(newEdge); graphR[j].push_back(newEdge);
                 //Dijkstra
                  queue<int> dfsOrder; memset(dist, -1, sizeof(dist));
                  typedef pair<int, pair<int, Edge*> > DijkstraQueueItem;
                 priority_queue<DijkstraQueueItem, vector<DijkstraQueueItem>, greater<DijkstraQueueItem> > dq;
                 dq.push(make_pair(0, make_pair(t, (Edge*) NULL)));
                 while (!dq.empty()) {
                           int d = dq.top().first; int i = dq.top().second.first;
                           Edge* edge = dq.top().second.second; dq.pop();
                           if (dist[i] != -1) continue;
                            dist[i] = d; prev[i] = edge; dfsOrder.push(i);
                            for_each(it, graphR[i]) dq.push(make_pair(d + (*it)->weight, make_pair((*it)->from, *it)));
                 }
                 //Create edge heap
                 nullNode = new HeapNode; nullNode->depth = 0; nullNode->edge = new Edge; nullNode->edge->weight = INF;
                 fill(nullNode->child, nullNode->child + 4, nullNode);
                 while (!dfsOrder.empty()) {
                           int i = dfsOrder.front(); dfsOrder.pop();
                           if (prev[i] == NULL) heapTop[i] = nullNode;
 49
                            else heapTop[i] = heapTop[prev[i]->to];
                            vector<HeapNode*> heapNodeList;
50
                            for_each(it, graph[i]) { int j = (*it) \rightarrow to; if (dist[j] == -1) continue;
51
                                      (*it)→weight += dist[i] - dist[i]; if (prev[i] != *it) {
52
                                               HeapNode* curNode = new HeapNode;
53
54
                                               fill(curNode->child, curNode->child + 4, nullNode);
55
                                               curNode->depth = 1; curNode->edge = *it;
56
                                               heapNodeList.push_back(curNode);
57
                           } if (!heapNodeList.empty()) { //Create heap out
58
                                      make_heap(heapNodeList.begin(), heapNodeList.end(), heapNodeMoreThan);
59
60
                                      int size = heapNodeList.size();
61
                                      for (int p = 0; p < size; p++) {
 62
                                               heapNodeList[p] \rightarrow child[2] = 2 * p + 1 < size ? heapNodeList[2 * p + 1] : nullNode;
                                               heapNodeList[p] \rightarrow child[3] = 2 * p + 2 < size ? heapNodeList[2 * p + 2] : nullNode;
 63
64
                                     } heapTop[i] = createHeap(heapTop[i], heapNodeList.front());
65
                 } //Walk on DAG
```

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```
typedef pair<long long, HeapNode*> DAGQueueItem;
       priority_queue<DAGQueueItem, vector<DAGQueueItem>, greater<DAGQueueItem> > aq;
       if (dist[s] == -1) printf("NO\n");
       else { printf("%d\n", dist[s]);
           if (heapTop[s] != nullNode) aq.push(make_pair(dist[s] + heapTop[s]->edge->weight, heapTop[s]));
71
       } k---; while (k---) {
           if (aq.empty()) { printf("NO\n"); continue; }
74
           long long d = aq.top().first; HeapNode* curNode = aq.top().second; aq.pop();
75
           printf("%I64d\n", d);
           if (heapTop[curNode->edge->to] != nullNode)
               aq.push(make_pair(d + heapTop[curNode->edge->to]->edge->weight, heapTop[curNode->edge->to]));
           for (int i = 0; i < 4; i++) if (curNode\rightarrowchild[i] != nullNode)
               aq.push(make_pair(d - curNode->edge->weight + curNode->child[i]->edge->weight, curNode->child[i]));
       } return 0;
4.18 K 短路 (不允许重复)
1 int Num[10005][205], Path[10005][205], dev[10005], from[10005], value[10005], dist[205], Next[205], Graph
         [205][205];
2 int N, M, K, s, t, tot, cnt; bool forbid[205], hasNext[10005][205];
 3 struct cmp {
       bool operator()(const int &a, const int &b) {
           int *i, *j; if (value[a] != value[b]) return value[a] > value[b];
           for (i = Path[a], j = Path[b]; (*i) == (*j); i++, j++);
           return (*i) > (*j);
9 };
10 void Check(int idx, int st, int *path, int &res) {
       int i, j; for (i = 0; i < N; i++) dist[i] = 1000000000, Next[i] = t;
       dist[t] = 0; forbid[t] = true; j = t;
       for (;;) {
         for (i = 0; i < N; i++) if (!forbid[i] && (i != st || !hasNext[idx][j]) && (dist[j] + Graph[i][j] < dist[
         i] \mid | (dist[i] + Graph[i][i] == dist[i] && i < Next[i])))
               Next[i] = j, dist[i] = dist[j] + Graph[i][j];
```

j = -1; for (i = 0; i < N; i++) if (!forbid[i] && (j == -1 || dist[i] < dist[j])) j = i;

} res += dist[st]; for (i = st; i != t; i = Next[i], path++) (\*path) = i; (\*path) = i;

for (i = 0; i < N; i++) for (j = 0; j < N; j++) Graph[i][j] = 10000000000;

for  $(j = 0; j \le dev[l]; j++) Num[l][j] = Num[from[l]][j];$ 

for (i = 0; i < M; i++) { scanf("%d%d%d", &j, &k, &l); Graph[j-1][k-1] = l; }

memset(forbid, false, sizeof(forbid)); memset(hasNext[0], false, sizeof(hasNext[0]));
Check(0, s, Path[0], value[0]); dev[0] = 0; from[0] = 0; Num[0][0] = 0; Q.push(0);

if (j == -1) break; forbid[j] = 1; if (j == st) break;

while (scanf("%d%d%d%d%d", &N, &M, &K, &s, &t) && N) {

if (Q.empty()) break; l = Q.top(); Q.pop();

priority\_queue<int, vector<int>, cmp> Q;

for (; Path[l][j] != t; j++) {

16

18

24

19 }

20 int main() {

int i, j, k, l;

cnt = 1; tot = 1;

for (i = 0; i < K; i++) {

```
memset(hasNext[tot], false, sizeof(hasNext[tot])); Num[l][j] = tot++;
               } for (j = 0; Path[l][j] != t; j++) hasNext[Num[l][j]][Path[l][j + 1]] = true;
35
               for (j = dev[l]; Path[l][j] != t; j++) {
                   memset(forbid, false, sizeof(forbid)); value[cnt] = 0;
37
                   for (k = 0; k < j; k++) {
38
                       forbid[Path[l][k]] = true;
                       Path[cnt][k] = Path[l][k];
40
                       value[cnt] += Graph[Path[l][k]][Path[l][k + 1]];
41
                   } Check(Num[l][j], Path[l][j], &Path[cnt][j], value[cnt]);
42
43
                   if (value[cnt] > 2000000) continue;
44
                   dev[cnt] = j; from[cnt] = l; Q.push(cnt); cnt++;
45
46
47
           if (i < K || value[l] > 2000000) printf("None\n");
48
               for (i = 0; Path[l][i] != t; i++) printf("%d-", Path[l][i] + 1);
49
50
               printf("%d\n", t + 1);
51
       } return 0;
53 }
```

#### 4.19 小知识

- 平面图: 一定存在一个度小干等于 5 的点. E < 3V 6. 欧拉公式: V + F E = 1 + 连通块数
- 图连通度:
  - 1. k— 连通 (k-connected): 对于任意一对结点都至少存在结点各不相同的 k 条路
  - 2. 点连通度 (vertex connectivity): 把图变成非连通图所需删除的最少点数
  - 3. Whitney 定理: 一个图是 k— 连通的当且仅当它的点连通度至少为 k
- Lindstroem-Gessel-Viennot Lemma: 给定一个图的 n 个起点和 n 个终点,令  $A_{ij}=$  第 i 个起点到第 j 个终点的路径条数,则从起点到终点的不相交路径条数为 det(A)
- 欧拉回路与树形图的联系: 对于出度等于入度的连通图  $s(G) = t_i(G) \prod_{i=1}^n (d^+(v_i) 1)!$
- 密度子图: 给定无向图,选取点集及其导出子图,最大化 W<sub>e</sub> + P<sub>v</sub> (点权可负).

- 
$$(S,u) = U$$
,  $(u,T) = U - 2P_u - D_u$ ,  $(u,v) = (v,u) = W_e$   
- ans =  $\frac{Un - C[S,T]}{2}$ , 解集为  $S - \{s\}$ 

$$-P_u > 0$$
,  $(S, u) = P_u$ ,  $P_u < 0$ ,  $(u, T) = -P_u$   
 $-\text{ans} = \sum_{P_u > 0} P_u - C[S, T]$ , 解集为  $S - \{s\}$ 

- 判定边是否属于最小割
  - 可能属于最小割: (u,v) 不属于同一 SCC
  - 一定在所有最小割中: (u,v) 不属于同一 SCC, 且 S,u 在同一 SCC, u,T 在同一 SCC
- 图同构 Hash: F<sub>t</sub>(i) = (F<sub>t-1</sub>(i) × A + ∑<sub>i→j</sub> F<sub>t-1</sub>(j) × B + ∑<sub>j←i</sub> F<sub>t-1</sub>(j) × C + D × (i = a)) (mod P),
   枚举点 a, 迭代 K 次后求得的 F<sub>k</sub>(a) 就是 a 点所对应的 Hash 值.

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## 5 数学

## 5.1 博弈论相关

1. Anti-SG:

规则与 Nim 基本相同, 取最后一个的输。

先手必胜当目仅当:

- (1) 所有堆的石子数都为 1 且游戏的 SG 值为 0;
- (2) 有些堆的石子数大于 1 且游戏的 SG 值不为 0。
- 2. SJ 定理:

对于任意一个 Anti-SG 游戏,如果我们规定当局面中,所有的单一游戏的 SG 值为 0 时,游戏结束,则先手必胜当日仅当:

- (1) 游戏的 SG 函数不为 0 且游戏中某个单一游戏的 SG 函数大于 1;
- (2) 游戏的 SG 函数为 0 且游戏中没有单一游戏的 SG 函数大于 1。
- 3. Multi-SG 游戏:

可以将一堆石子分成多堆.

4. Every-SG 游戏:

每一个可以移动的棋子都要移动.

对于我们可以赢的单一游戏,我们一定要拿到这一场游戏的胜利.

只需要考虑如何让我们必胜的游戏尽可能长的玩下去,对手相反。

于是就来一个 DP,

step[v] = 0; (v 为终止状态)

step[v] = maxstep[u] + 1; (sg[v]>0,sg[u]=0)

step[v] = minstep[u] + 1; (sg[v] = 0)

5. 翻硬币游戏:

N 枚硬币排成一排,有的正面朝上,有的反面朝上。游戏者根据某些约束翻硬币 (如:每次只能翻一或两枚,或者每次只能翻连续的几枚),但他所翻动的硬币中,最右边的必须是从正面翻到反面。谁不能翻谁输。

结论: 局面的 SG 值为局面中每个正面朝上的棋子单一存在时的 SG 值的异或和。可用数学归纳法证明。

6. 无向树删边游戏:

抑则加下

给出一个有 N 个点的树,有一个点作为树的根节点。游戏者轮流从树中删去边,删去一条边后,不与根节点相连的部分将被移走。谁无路可走谁输。

结论:

叶子节点的 SG 值为 0; 中间节点的 SG 值为它的所有子节点的 SG 值加 1 后的异或和。是用数学归纳法证明。

7. Christmas Game(PKU3710):

题目大意:

有 N 个局部联通的图。Harry 和 Sally 轮流从图中删边,删去一条边后,不与根节点相连的部分将被移走。 23 Sally 为先手。图是通过从基础树中加一些边得到的。所有形成的环保证不共用边,且只与基础树有一个公共 24 点。谁无路可走谁输。环的处理成为了解题的关键。 25 Mt 后:

- (1) 对于长度为奇数的环,去掉其中任意一个边之后,剩下的两个链长度同奇偶,抑或之后的 SG 值不可能为奇数,所以它的 SG 值为 1;
- (2) 对于长度为偶数的环,去掉其中任意一个边之后,剩下的两个链长度异奇偶,抑或之后的 SG 值不可能为 29

0, 所以它的 SG 值为 0; 所以我们可以去掉所有的偶环, 将所有的奇环变为长短为 1 的链。这样的话, 我们已经将这道题改造成了上一节的模型。

8. 无向图的删边游戏:

我们将 Christmas Game 这道题进行一步拓展——去掉对环的限制条件,这个模型应该怎样处理? 无向图的删边游戏:

一个无向联通图,有一个点作为图的根。游戏者轮流从图中删去边,删去一条边后,不与根节点相连的部分将被移走。谁无路可走谁输。

结论:

对无向图做如下改动:将图中的任意一个偶环缩成一个新点,任意一个奇环缩成一个新点加一个新边;所有连到原先环上的边全部改为与新点相连。这样的改动不会影响图的 SG 值。

9. Staircase nim:

楼梯从地面由下向上编号为 0 到 n。游戏者在每次操作时可以将楼梯 j(1 <= j <= n) 上的任意多但至少一个硬币移动到楼梯 j(1 <= j <= n) 上的任意多但至少一个硬币移动到楼梯 j(1 <= j <= n) 上的任意多但至少一个硬

结论:

设该游戏 Sg 函数为奇数格棋子数的 Xor 和 S。

如果 S=0,则先手必败,否则必胜。

# 5.2 单纯形 Cpp $\max \{cx | Ax < b, x > 0\}$

```
1 const int MAXN = 11000, MAXM = 1100;
2 // `here MAXN is the MAX number of conditions, MAXM is the MAX number of vars
4 int avali[MAXM], avacnt;
5 double A[MAXN][MAXM]:
6 double b[MAXN], c[MAXM];
7 double* simplex(int n, int m) {
8 // `here n is the number of conditions, m is the number of vars
       m++:
       int r = n, s = m - 1;
10
       static double D[MAXN + 2][MAXM + 1];
       static int ix[MAXN + MAXM];
13
       for (int i = 0; i < n + m; i++) ix[i] = i;
14
       for (int i = 0; i < n; i++) {
           for (int j = 0; j < m-1; j++) D[i][j] = -A[i][j];
           D[i][m-1] = 1;
16
17
           D[i][m] = b[i];
           if (D[r][m] > D[i][m]) r = i;
       for (int j = 0; j < m - 1; j++) D[n][j] = c[j];
       D[n + 1][m - 1] = -1;
21
       for (double d; ; ) {
          if (r < n) {
              int t = ix[s]; ix[s] = ix[r + m]; ix[r + m] = t;
25
              D[r][s] = 1.0 / D[r][s];
              for (int j = 0; j \leftarrow m; j++) if (j != s) D[r][j] *= -D[r][s];
26
               for (int i = 0; i <= m; ++i)
                  if(fabs(D[r][i]) > EPS)
```

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```
avali[avacnt++] = i;
                for (int i = 0; i \le n + 1; i++) if (i != r) {
                    if(fabs(D[i][s]) < EPS) continue;</pre>
                    double *cur1 = D[i], *cur2 = D[r], tmp = D[i][s];
                    //for (int j = 0; j \le m; j++) if (j != s) cur1[j] += cur2[j] * tmp;
                    for(int j = 0; j < avacnt; ++j) if(avali[j] != s) cur1[avali[j]] += cur2[avali[j]] * tmp;</pre>
                    D[i][s] *= D[r][s];
           r = -1; s = -1;
           for (int j = 0; j < m; j++) if (s < 0 || ix[s] > ix[j]) {
                if (D[n + 1][j] > EPS \mid I \mid D[n + 1][j] > -EPS && D[n][j] > EPS) s = j;
42
           if (s < 0) break;
           for (int i = 0; i < n; i++) if (D[i][s] < -EPS) {
                if (r < 0 \mid l \mid (d = D[r][m] / D[r][s] - D[i][m] / D[i][s]) < -EPS
                          II d < EPS && ix[r + m] > ix[i + m])
47
                    r = i;
           if (r < 0) return null; // 津有界
50
       if (D[n + 1][m] <—EPS) return null; // `无法执行`
51
52
        static double x[MAXM - 1];
        for (int i = m; i < n + m; i++) if (ix[i] < m - 1) x[ix[i]] = D[i - m][m];
53
       return x; // `值为 $D[n][m]$`
54
55 }
```

#### 5.3 自适应辛普森

```
1 double area(const double &left, const double &right) {
        double mid = (left + right) / 2;
        return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
 4 }
 6 double simpson(const double &left, const double &right,
                  const double &eps, const double &area_sum) {
       double mid = (left + right) / 2;
 9
       double area_left = area(left, mid);
10
       double area_right = area(mid, right);
11
        double area_total = area_left + area_right;
       if (std::abs(area_total - area_sum) < 15 * eps) {</pre>
13
           return area_total + (area_total - area_sum) / 15;
14
15
       return simpson(left, mid, eps / 2, area_left)
16
            + simpson(mid, right, eps / 2, area_right);
17 }
18
19 double simpson(const double &left, const double &right, const double &eps) {
        return simpson(left, right, eps, area(left, right));
21 }
```

#### 5.4 FFT

```
1 namespace FFT {
       #define mul(a, b) (Complex(a.x * b.x - a.y * b.y, a.x * b.y + a.y * b.x))
       struct Complex {}; // `something omitted`
       void FFT(Complex P□, int n, int oper) {
            for (int i = 1, j = 0; i < n - 1; i++) {
                for (int s = n; j \triangleq s >>= 1, ~j \& s; );
               if (i < j) swap(P[i], P[j]);</pre>
            for (int d = 0; (1 << d) < n; d++) {
                int m = 1 \ll d, m2 = m * 2;
                double p0 = PI / m * oper;
11
12
               Complex unit_p0(cos(p0), sin(p0));
13
                for (int i = 0; i < n; i += m2) {
                    Complex unit(1.0, 0.0);
14
                    for (int j = 0; j < m; j++) {
15
                        Complex &P1 = P[i + j + m], &P2 = P[i + j];
16
                        Complex t = mul(unit, P1);
                        P1 = Complex(P2.x - t.x, P2.y - t.y);
18
19
                        P2 = Complex(P2.x + t.x, P2.y - t.y);
                        unit = mul(unit, unit_p0);
20
21
       }}}}
22
       vector<int> doFFT(const vector<int> &a. const vector<int> &b) {
23
            vector<int> ret(max(0, (int) a.size() + (int) b.size() - 1), 0);
            static Complex A[MAXB], B[MAXB], C[MAXB];
24
25
           int len = 1; while (len < (int)ret.size()) len *= 2;</pre>
           for (int i = 0; i < len; i++) A[i] = i < (int)a.size() ? <math>a[i] : 0;
26
            for (int i = 0; i < len; i++) B[i] = i < (int)b.size() ? <math>b[i] : 0;
27
28
           FFT(A, len, 1); FFT(B, len, 1);
           for (int i = 0; i < len; i++) C[i] = mul(A[i], B[i]);</pre>
29
30
           FFT(C, len, -1);
            for (int i = 0; i < (int)ret.size(); i++)</pre>
31
32
               ret[i] = (int) (C[i].x / len + 0.5);
33
            return ret;
34
35 }
```

## 5.5 整数 FFT

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```
if (i < j) swap(P[i], P[j]);</pre>
           for (int d = 0; (1 << d) < n; d++) {
                int m = 1 \ll d, m2 = m * 2;
16
                long long unit_p0 = powMod(PRIMITIVE\_ROOT, (MOD - 1) / m2);
                if (oper < 0) unit_p0 = modInv(unit_p0);</pre>
                for (int i = 0; i < n; i += m2) {
                   long long unit = 1;
21
                    for (int j = 0; j < m; j++) {
                        int &P1 = P[i + j + m], &P2 = P[i + j];
22
                        int t = unit * P1 % MOD;
                       P1 = (P2 - t + MOD) \% MOD; P2 = (P2 + t) \% MOD;
24
25
                        unit = unit * unit p0 % MOD:
       }}}}
26
        vector<int> mul(const vector<int> &a, const vector<int> &b) {
27
           vector<int> ret(max(0, (int) a.size() + (int) b.size() - 1), 0);
28
           static int A[MAXB], B[MAXB], C[MAXB];
29
           int len = 1; while (len < (int)ret.size()) len <<= 1;</pre>
30
            for (int i = 0; i < len; i++) A[i] = i < (int)a.size() ? <math>a[i] : 0;
            for (int i = 0; i < len; i++) B[i] = i < (int)b.size() ? <math>b[i] : 0;
           NTT(A, len, 1); NTT(B, len, 1);
33
           for (int i = 0; i < len; i++) C[i] = (long long) A[i] * B[i] % MOD;
34
           NTT(C, len, -1); for (int i = 0, inv = modInv(len); i < (int)ret.size(); i++) ret[i] = (long long) C[i] *
          inv % MOD:
           return ret;
37
38 }
```

#### 5.6 扩展欧几里得

## 5.7 线性同余方程

- 中国剩余定理: 设  $m_1, m_2, \cdots, m_k$  两两互素,则同余方程组  $x \equiv a_i \pmod{m_i}$  for  $i=1,2,\cdots,k$  在  $[0,M=m_1m_2\cdots m_k)$  内有唯一解. 记  $M_i=M/m_i$ ,找出  $p_i$  使得  $M_ip_i\equiv 1\pmod{m_i}$ ,记  $e_i=M_ip_i$ ,则  $x\equiv e_1a_1+e_2a_2+\cdots+e_ka_k\pmod{M}$
- 多变元线性同余方程组: 方程的形式为  $a_1x_1+a_2x_2+\cdots+a_nx_n+b\equiv 0\pmod m$ , 令  $d=(a_1,a_2,\cdots,a_n,m)$ , 有解的充要条件是 dlb, 解的个数为  $m^{n-1}d$

#### 5.8 Miller-Rabin 素性测试

```
1 bool test(LL n, int base) {
       LL m = n - 1, ret = 0; int s = 0;
       for (; m \% 2 == 0; ++s) m >>= 1; ret = pow_mod(base, m, n);
       if (ret == 1 | | | ret == n - 1) return true;
       for (-s; s \ge 0; -s) {
           ret = multiply_mod(ret, ret, n); if (ret == n - 1) return true;
       } return false;
8 }
9 LL special[7] = {
       1373653LL,
                          25326001LL,
       3215031751LL.
                          250000000000LL
       2152302898747LL, 3474749660383LL, 341550071728321LL};
13 /*
14 * n < 2047
                                       test[] = {2}
* n < 1,373,653
                                       test[] = \{2, 3\}
* n < 9,080,191
                                       test[] = {31, 73}
* n < 25,326,001
                                       test[] = \{2, 3, 5\}
* n < 4,759,123,141
                                       test[] = \{2, 7, 61\}
* n < 1,122,004,669,633
                                       test = \{2, 13, 23, 1662803\}
20 * n < 2,152,302,898,747
                                       test [] = {2, 3, 5, 7, 11}
* n < 3,474,749,660,383
                                       test[] = \{2, 3, 5, 7, 11, 13\}
* n < 341,550,071,728,321
                                       test[] = {2, 3, 5, 7, 11, 13, 17}
* n < 3,825,123,056,546,413,051
                                       test[] = {2, 3, 5, 7, 11, 13, 17, 19, 23}
24 */
25 bool is_prime(LL n) {
       if (n < 2) return false;
       if (n < 4) return true;
       if (!test(n, 2) || !test(n, 3)) return false;
       if (n < special[0]) return true;</pre>
       if (!test(n, 5)) return false;
       if (n < special[1]) return true;
31
       if (!test(n, 7)) return false;
       if (n == special[2]) return false;
       if (n < special[3]) return true;</pre>
35
       if (!test(n, 11)) return false;
       if (n < special[4]) return true;
       if (!test(n, 13)) return false;
       if (n < special[5]) return true;</pre>
       if (!test(n, 17)) return false;
39
       if (n < special[6]) return true;</pre>
41
       return test(n, 19) && test(n, 23) && test(n, 29) && test(n, 31) && test(n, 37);
42 }
```

#### 5.9 PollardRho

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#### 5.10 多项式求根

```
1 const double error = 1e-12;
 2 const double infi = 1e+12;
 3 int n; double a[10], x[10];
 4 double f(double a□, int n, double x) {
       double tmp = 1, sum = 0;
       for (int i = 0; i \le n; i++) sum = sum + a[i] * tmp, tmp = tmp * x;
       return sum;
8 }
9 double binary(double 1, double r, double a□, int n) {
       int sl = sign(f(a, n, l)), sr = sign(f(a, n, r));
       if (sl == 0) return l; if (sr == 0) return r;
       if (sl * sr > 0) return infi;
       while (r - l > error) {
           double mid = (l + r) / 2;
           int ss = sign(f(a, n, mid));
           if (ss == 0) return mid;
           if (ss * sl > 0) l = mid: else r = mid:
18
       } return l;
19 }
20 void solve(int n, double a□, double x□, int &nx) {
       if (n == 1) \{ x[1] = -a[0] / a[1]; nx = 1; return; \}
       double da\lceil 10 \rceil, dx\lceil 10 \rceil; int ndx;
       for (int i = n; i >= 1; i \longrightarrow) da[i - 1] = a[i] * i;
       solve(n-1, da, dx, ndx); nx = 0;
24
       if (ndx == 0) {
           double tmp = binary(-infi, infi, a, n);
27
           if (tmp < infi) x[++nx] = tmp; return;
       } double tmp = binary(-infi, dx[1], a, n);
       if (tmp < infi) x[++nx] = tmp;
29
       for (int i = 1; i \le ndx - 1; i++) {
30
31
           tmp = binary(dx[i], dx[i + 1], a, n);
           if (tmp < infi) x[++nx] = tmp;
       } tmp = binary(dx[ndx], infi, a, n);
       if (tmp < infi) x[++nx] = tmp;
34
35 }
36 int main() {
       scanf("%d", &n);
       for (int i = n; i \ge 0; i—) scanf("%lf", &a[i]);
       int nx; solve(n, a, x, nx);
       for (int i = 1; i \le nx; i++) printf("%0.6f\n", x[i]);
       return 0;
42 }
```

#### 5.11 线性递推

```
for a_{i+n} = (\sum_{i=0}^{n-1} k_i a_{i+j}) + d, a_m = (\sum_{i=0}^{n-1} c_i a_i) + c_n d
1 vector<int> recFormula(int n, int k□, int m) {
       vector<int> c(n + 1, 0);
       if (m < n) c[m] = 1;
           static int a[MAX_K * 2 + 1];
           vector<int> b = recFormula(n, k, m >> 1);
           for (int i = 0; i < n + n; ++i) a[i] = 0;
           int s = m \& 1:
            for (int i = 0; i < n; i++) {
               for (int j = 0; j < n; j++) a[i + j + s] += b[i] * b[j];
               c[n] += b[i];
11
           c[n] = (c[n] + 1) * b[n];
12
           for (int i = n * 2 - 1; i >= n; i \longrightarrow) {
13
               int add = a[i]; if (add == 0) continue;
14
               for (int j = 0; j < n; j++) a[i - n + j] += k[j] * add;
15
               c[n] += add;
16
           } for (int i = 0; i < n; ++i) c[i] = a[i];
17
18
       } return c;
19 }
```

## 5.12 原根

原根 g: g 是模 n 简化剩余系构成的乘法群的生成元. 模 n 有原根的充要条件是  $n=2,4,p^n,2p^n,$  其中 p 是奇质数, n 是正整数

```
1 vector<int> findPrimitiveRoot(int N) {
       if (N \le 4) return vector(1, \max(1, N-1));
       static int factor[100];
       int phi = N, totF = 0;
       { // `check no solution and calculate phi`
          int M = N, k = 0;
          if (~M & 1) M >>= 1, phi >>= 1;
          if (~M & 1) return vector<int>(0);
           for (int d = 3; d * d <= M; ++d) if (M % d == 0) {
              if (++k > 1) return vector<int>(0);
11
               for (phi -= phi / d; M % d == 0; M /= d);
12
          } if (M > 1) {
               if (++k > 1) return vector<int>(0): phi -= phi / M:
13
14
       } { // `factorize phi`
15
          int M = phi;
           for (int d = 2; d * d <= M; ++d) if (M % d == 0) {
17
              for (; M % d == 0; M /= d); factor[++totF] = d;
18
19
          if (M > 1) factor[++totF] = M;
       } vector<int> ans;
20
       for (int g = 2; g <= N; ++g) if (Gcd(g, N) == 1) {
          bool good = true;
22
23
           for (int i = 1; i <= totF && good; ++i)
              if (powMod(g, phi / factor[i], N) == 1) good = false;
24
25
           if (!good) continue;
```

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```
for (int i = 1, gp = g; i <= phi; ++i, gp = (LL)gp * g % N)

if (Gcd(i, phi) == 1) ans.push_back(gp);

break;

sort(ans.begin(), ans.end());

return ans;

1 }</pre>
```

## 5.13 离散对数

 $A^x \equiv B \pmod{C}$ , 对非质数 C 也适用.

```
int modLog(int A, int B, int C) {
       static pii baby[MAX_SQRT_C + 11];
       int d = 0; LL k = 1, D = 1; B %= C;
       for (int i = 0; i < 100; ++i, k = k * A % C) // `$[0, \log C]$
           if (k == B) return i;
       for (int g; ; ++d) {
           g = gcd(A, C); if (g == 1) break;
           if (B % g != 0) return -1;
           B /= g; C /= g; D = (A / g * D) % C;
      } int m = (int) ceil(sqrt((double) C)); k = 1;
       for (int i = 0; i \le m; ++i, k = k * A % C) baby[i] = pii(k, i);
       sort(baby, baby + m + 1); // [0, m]
       int n = unique(baby, baby + m + 1, equalFirst) — baby, am = powMod(A, m, C);
       for (int i = 0; i <= m; ++i) {
          LL e, x, y; exgcd(D, C, x, y, e); e = x * B % C;
          if (e < 0) e += C;
           if (e >= 0) {
               int k = lower\_bound(baby, baby + n, pii(e, -1)) - baby;
               if (baby[k].first == e) return i * m + baby[k].second + d;
           D = D * am % C;
21
       } return -1;
```

## 5.14 平方剩余

- Legrendre Symbol: 对奇质数  $p, \left(\frac{a}{p}\right) = \begin{cases} 1 & \text{是平方剩余} \\ -1 & \text{是非平方剩余} = a^{\frac{p-1}{2}} \bmod p \\ 0 & a \equiv 0 \pmod p \end{cases}$
- 若 p 是奇质数,  $\left(\frac{-1}{p}\right) = 1$  当且仅当  $p \equiv 1 \pmod{4}$
- 若 p 是奇质数,  $(\frac{2}{p}) = 1$  当且仅当  $p \equiv \pm 1 \pmod{8}$
- 若 p,q 是奇素数且互质,  $(\frac{p}{q})(\frac{q}{n}) = (-1)^{\frac{p-1}{2} \times \frac{q-1}{2}}$
- Jacobi Symbol: 对奇数  $n=p_1^{\alpha_1}p_2^{\alpha_2}\cdots p_k^{\alpha_k}, (\frac{a}{n})=(\frac{a}{p_1})^{\alpha_1}(\frac{a}{p_2})^{\alpha_2}\cdots (\frac{a}{p_k})^{\alpha_k}$
- Jacobi Symbol 为 -1 则一定不是平方剩余,所有平方剩余的 Jacobi Symbol 都是 1, 但 1 不一定是平方剩余  $ax^2+bx+c\equiv 0\pmod p$ , 其中  $a\neq 0\pmod p$ , 且 p 是质数

```
1 inline int normalize(LL a, int P) { a \approx P; return a < 0 ? a + P : a; }
 2 vector<int> QuadraticResidue(LL a, LL b, LL c, int P) {
       int h, t; LL r1, r2, delta, pb = 0;
       a = normalize(a, P); b = normalize(b, P); c = normalize(c, P);
       if (P == 2) { vector<int> res;
           if (c % P == 0) res.push_back(0);
           if ((a + b + c) \% P == 0) res.push_back(1);
           return res;
      } delta = b * rev(a + a, P) % P;
       a = normalize(-c * rev(a, P) + delta * delta, P);
       if (powMod(a, P / 2, P) + 1 == P) return vector<int>(0);
       for (t = 0, h = P / 2; h \% 2 == 0; ++t, h /= 2);
       r1 = powMod(a, h / 2, P);
       if (t > 0) { do b = random() % (P - 2) + 2;
           while (powMod(b, P / 2, P) + 1 != P); }
       for (int i = 1; i <= t; ++i) {
17
         LL d = r1 * r1 % P * a % P;
18
           for (int j = 1; j \le t - i; ++j) d = d * d % P;
          if (d + 1 == P) r1 = r1 * pb % P; pb = pb * pb % P;
      r1 = a * r1 % P; r2 = P - r1;
       r1 = normalize(r1 - delta, P); r2 = normalize(r2 - delta, P);
       if (r1 > r2) swap(r1, r2); vector<int> res(1, r1);
       if (r1 != r2) res.push_back(r2);
       return res;
24
25 }
```

## 5.15 N 次剩余

• 若 p 为奇质数, a 为 p 的 n 次剩余的充要条件是  $a^{\frac{p-1}{(a,p-1)}} \equiv 1 \pmod{p}$ .

 $x^N \equiv a \pmod{p}$ , 其中 p 是质数

# 5.16 Pell 方程

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_1 & dy_1 \\ y_1 & x_1 \end{pmatrix}^{k-1} \begin{pmatrix} x_1 \\ y_1 \end{pmatrix}$$
 pair Pell(int n) { 
$$2 \quad \text{static ULL p[50] = \{0, 1\}, q[50] = \{1, 0\}, g[50] = \{0, 0\}, h[50] = \{0, 1\}, a[50]; }$$

```
3  ULL t = a[2] = Sqrt(n);
4  for (int i = 2; ; ++i) {
5     g[i] = -g[i - 1] + a[i] * h[i - 1];
6     h[i] = (n - g[i] * g[i]) / h[i - 1];
7     a[i + 1] = (g[i] + t) / h[i];
8     p[i] = a[i] * p[i - 1] + p[i - 2];
9     q[i] = a[i] * q[i - 1] + q[i - 2];
10     if (p[i] * p[i] - n * a[i] * a[i] == 1) return make_pair(p[i], a[i]);
11     } return make_pair(-1, -1);
12 }
```

#### 5.17 小知识

- 勾股数: 设正整数 n 的质因数分解为  $n = \prod p_i^{a_i}$ , 则  $x^2 + y^2 = n$  有整数解的充要条件是 n 中不存在形如  $p_i \equiv 3 \pmod{4}$  且指数  $a_i$  为奇数的质因数  $p_i$ .  $(\frac{a-b}{2})^2 + ab = (\frac{a+b}{2})^2$ .
- 素勾股数: 若 m 和 n 互质, 而且 m 和 n 中有一个是偶数, 则  $a=m^2-n^2,$  b=2mn,  $c=m^2+n^2,$  则 a、b、c 是素勾股数.
- Stirling 公式:  $n! \approx \sqrt{2\pi n} (\frac{n}{\epsilon})^n$
- Pick 定理: 简单多边形, 不自交, 顶点如果全是整点. 则: 严格在多边形内部的整点数  $+\frac{1}{2}$  在边上的整点数 -1= 面积
- Mersenne 素数: p 是素数且 2<sup>p</sup> 1 的数是素数. (10000 以内的 p 有: 2, 3, 5, 7, 13, 17, 19, 31, 61, 89, 107, 127, 521, 607, 1279, 2203, 2281, 3217, 4253, 4423, 9689, 9941)
- 序列差分表: 差分表的第 0 条对角线确定原序列. 设原序列为  $h_i$ , 第 0 条对角线为  $c_0, c_1, \ldots, c_p, 0, 0, \ldots$  有这样两个公式:  $h_n = \binom{n}{0}c_0 + \binom{n}{1}c_1 + \ldots + \binom{n}{n}c_p$ ,  $\sum_{k=0}^n h_k = \binom{n+1}{1}c_0 + \binom{n+1}{2}c_2 + \ldots + \binom{n+1}{n+1}c_p$
- GCD:  $gcd(2^a 1, 2^b 1) = 2^{gcd(a,b)} 1$
- Fermat 分解算法: 从  $t=\sqrt{n}$  开始, 依次检查  $t^2-n, (t+1)^2-n, (t+2)^2-n, \ldots$ , 直到出现一个平方数 y, 由于  $t^2-y^2=n$ , 因此分解得 n=(t-y)(t+y). 显然, 当两个因数很接近时这个方法能很快找到结果, 但如果遇到一个素数. 则需要检查  $\frac{n+1}{2}-\sqrt{n}$  个整数
- 牛顿迭代:  $x_1 = x_0 \frac{f(x_0)}{f'(x_0)}$
- 球与盒子的动人故事:  $(n \ \bigcirc \ \ )$  水  $\cap$  金子, S 为第二类斯特林数)
  - 1. 球同, 盒同, 无空: dp
  - 2. 球同, 盒同, 可空: dp
  - 3. 球同, 盒不同, 无空:  $\binom{n-1}{n}$
  - 4. 球同, 盒不同, 可空:  $\binom{n+m-1}{n-1}$
  - 5. 球不同, 盒同, 无空: S(n, m)
  - 6. 球不同, 盒同, 可空:  $\sum_{k=1}^{m} S(n,k)$
  - 7. 球不同, 盒不同, 无空: m!S(n,m)
  - 1. 然中间, 血中间, 九工. 111.5(1
  - 8. 球不同, 盒不同, 可空:  $m^n$
- 组合数奇偶性: 若 (n&m)=m, 则  $\binom{n}{m}$  为奇数, 否则为偶数

- 格雷码  $G(x) = x \otimes (x >> 1)$
- Fibonacci 数:

$$-F_0 = F_1 = 1, F_i = F_{i-1} + F_{i-2}, F_{-i} = (-1)^{i-1} F_i$$

$$-F_i = \frac{1}{\sqrt{5}} \left( \left( \frac{1 + \sqrt{5}}{2} \right)^n - \left( \frac{1 - \sqrt{5}}{2} \right)^n \right)$$

$$-\gcd(F_n, F_m) = F_{\gcd(n,m)}$$

$$-F_{i+1} F_i - F_i^2 = (-1)^i$$

$$-F_{n+k} = F_k F_{n+1} + F_{k-1} F_n$$

• 第一类 Stirling 数:  $\binom{n}{k}$  代表第一类无符号 Stirling 数, 代表将 n 阶置换群中有 k 个环的置换个数; s(n,k) 代表有符号型,  $s(n,k)=(-1)^{n-k}\binom{n}{k}$ .

$$-(x)^{(n)} = \sum_{k=0}^{n} {n \brack k} x^{k}, (x)_{n} = \sum_{k=0}^{n} s(n,k) x^{k}$$

$$- {n \brack k} = n {n-1 \brack k} + {n-1 \brack k-1}, {0 \brack 0} = 1, {n \brack 0} = {0 \brack n} = 0$$

$$- {n \brack n-2} = \frac{1}{4} (3n-1) {n \brack 3}, {n \brack n-3} = {n \brack 2} {n \brack 4}$$

$$- \sum_{k=0}^{a} {n \brack k} = n! - \sum_{k=0}^{n} {n \brack k+a+1}$$

$$- \sum_{n=k}^{n} {n \brack p} {n \brack k} = {n+1 \brack k+1}$$

• 第二类 Stirling 数:  $\binom{n}{k} = S(n,k)$  代表 n 个不同的球, 放到 k 个相同的盒子里, 盒子非空.

$$- \left\{ {n \atop k} \right\} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^j {k \choose j} (k-j)^n$$

$$- \left\{ {n+1 \atop k} \right\} = k {n \atop k} + {n \atop k-1}, {0 \atop 0} = 1, {n \atop 0} = {0 \atop n} = 0$$

$$- 奇偶性: (n-k) & \frac{k-1}{2} = 0$$

• Bell 数:  $B_n$  代表将 n 个元素划分成若干个非空集合的方案数

$$-B_0=B_1=1, B_n=\sum_{k=0}^{n-1} \binom{n-1}{k} B_k$$

$$-B_n=\sum_{k=0}^n \binom{n}{k}$$

$$-Bell 三角形: a_{1,1}=1, a_{n,1}=a_{n-1,n-1}, a_{n,m}=a_{n,m-1}+a_{n-1,m-1}, B_n=a_{n,1}$$

$$- 对质数 \ p, B_{n+p}\equiv B_n+B_{n+1} \ (\text{mod }p)$$

$$- 对质数 \ p, B_{n+p^m}\equiv mB_n+B_{n+1} \ (\text{mod }p)$$

$$- 对质数 \ p, 模的周期一定是 \ \frac{p^p-1}{p-1} \ 的约数, \ p\leq 101 \ \text{时就是这个值}$$

• Bernoulli 数

$$-B_0 = 1, B_1 = \frac{1}{2}, B_2 = \frac{1}{6}, B_4 = -\frac{1}{30}, B_6 = \frac{1}{42}, B_8 = B_4, B_{10} = \frac{5}{66}$$

- 从 B<sub>0</sub> 开始, 前几项是 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975 · · ·

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$$-\sum_{k=1}^{n} k^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} n^{m+1-k}$$
$$-B_{m} = 1 - \sum_{k=0}^{m-1} {m \choose k} \frac{B_{k}}{m-k+1}$$

• 完全数: x 是偶完全数等价于  $x = 2^{n-1}(2^n - 1)$ , 且  $2^n - 1$  是质数.

## 6 其他

#### 6.1 Extended LIS

```
1 int G[MAXN][MAXN];
2 void insertYoung(int v) {
3     for (int x = 1, y = INT_MAX; ; ++x) {
4         Down(y, *G[x]); while (y > 0 && G[x][y] >= v) —y;
5         if (++y > *G[x]) { ++*G[x]; G[x][y] = v; break; }
6         else swap(G[x][y], v);
7     }
8     }
9     int solve(int N, int seq[]) {
10         Rep(i, 1, N) *G[i] = 0;
11         Rep(i, 1, N) insertYoung(seq[i]);
12         printf("%d\n", *G[1] + *G[2]);
13         return 0;
14 }
```

#### 6.2 生成 nCk

```
void nCk(int n, int k) {
for (int comb = (1 << k) - 1; comb < (1 << n); ) {
    int x = comb & -comb, y = comb + x;

comb = (((comb & ~y) / x) >> 1) | y;
}
}
```

#### 6.3 nextPermutation

```
1 boolean nextPermutation(int[] is) {
2     int n = is.length;
3     for (int i = n - 1; i > 0; i—) {
4         if (is[i - 1] < is[i]) {
5             int j = n; while (is[i - 1] >= is[—j]);
6             swap(is, i - 1, j); // swap is[i - 1], is[j]
7             rev(is, i, n); // reverse is[i, n)
8             return true;
9         }
10     } rev(is, 0, n);
11     return false;
12 }
```

## 6.4 Josephus 数与逆 Josephus 数

```
1 int josephus(int n, int m, int k) { int x = -1;
2     for (int i = n - k + 1; i <= n; i++) x = (x + m) % i; return x;
3 }
4 int invJosephus(int n, int m, int x) {
5     for (int i = n; ; i---) { if (x == i) return n - i; x = (x - m % i + i) % i; }
6 }</pre>
```

#### 6.5 直线下的整点个数

## 6.6 Java 多项式

```
1 class Polynomial {
       final static Polynomial ZERO = new Polynomial(new int[] { 0 });
       final static Polynomial ONE = new Polynomial(new int[] { 1 });
       final static Polynomial X = new Polynomial(new int[] { 0, 1 });
       int∏ coef;
       static Polynomial valueOf(int val) { return new Polynomial(new int[] { val }); }
       Polynomial(int[] coef) { this.coef = Arrays.copyOf(coef, coef.length); }
       Polynomial add(Polynomial o, int mod); // omitted
       Polynomial subtract(Polynomial o, int mod); // omitted
       Polynomial multiply(Polynomial o, int mod); // omitted
       Polynomial scale(int o, int mod); // omitted
11
12
       public String toString() {
           int n = coef.length; String ret = "";
           for (int i = n - 1; i > 0; —i) if (coef[i] != 0)
14
               ret += coef[i] + "x^" + i + "+";
15
           return ret + coef[0];
16
17
       static Polynomial lagrangeInterpolation(int[] x, int[] y, int mod) {
18
19
           int n = x.length; Polynomial ret = Polynomial.ZERO;
           for (int i = 0; i < n; ++i) {
20
               Polynomial poly = Polynomial.value0f(y[i]);
21
               for (int j = 0; j < n; ++j) if (i != j) {
22
23
                   poly = poly.multiply(
24
                       Polynomial.X.subtract(Polynomial.valueOf(x[j]), mod), mod);
                   poly = poly.scale(powMod(x[i] - x[j] + mod, mod - 2, mod), mod);
               } ret = ret.add(poly, mod);
27
           } return ret;
28
29 }
```

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## 6.7 long long 乘法取模

```
1 LL multiplyMod(LL a, LL b, LL P) { // 需要保证 a 和 b 非负;
2 LL t = (a * b — LL((long double)a / P * b + 1e—3) * P) % P;
3 return t < 0 : t + P : t;
4 }
```

## 6.8 重复覆盖

```
1 namespace DLX {
         struct node { int x, y; node *1, *r, *u, *d; } base[MAX * MAX], *top, *head;
         typedef node *link;
        int row, col, nGE, ans, stamp, cntc[MAX], vis[MAX];
         vector<link> eachRow[MAX], eachCol[MAX];
         inline void addElement(int x, int y) {
             top \rightarrow x = x; top \rightarrow y = y; top \rightarrow l = top \rightarrow r = top \rightarrow u = top \rightarrow d = NULL;
             eachRow[x].push_back(top); eachCol[y].push_back(top++);
10
         void init(int _row, int _col, int _nGE) {
              row = _row; col = _col; nGE = _nGE; top = base; stamp = 0;
11
             for (int i = 0; i \le col; ++i) vis[i] = 0;
12
13
             for (int i = 0; i <= row; ++i) eachRow[i].clear();</pre>
             for (int i = 0; i <= col; ++i) eachCol[i].clear();</pre>
              for (int i = 0; i \le col; ++i) addElement(0, i);
             head = eachCol[0].front();
        }
17
        void build() {
18
             for (int i = 0; i <= row; ++i) {
19
                   vector<link> &v = eachRow[i];
21
                   sort(v.begin(), v.end(), cmpByY);
22
                   int s = v.size();
                   for (int j = 0; j < s; ++j) {
                       link l = v[j], r = v[(j + 1) \% s]; l \rightarrow r = r; r \rightarrow l = l;
             for (int i = 0; i <= col; ++i) {
                   vector<link> &v = eachCol[i];
                   sort(v.begin(), v.end(), cmpByX);
                   int s = v.size();
                   for (int j = 0; j < s; ++j) {
31
                       link u = v[j], d = v[(j + 1) \% s]; u \rightarrow d = d; d \rightarrow u = u;
32
             } for (int i = 0; i \le col; ++i) cntc[i] = (int) eachCol[i].size() - 1;
34
35
        void removeExact(link c) {
36
37
             c \rightarrow l \rightarrow r = c \rightarrow r; c \rightarrow r \rightarrow l = c \rightarrow l;
              for (link i = c \rightarrow d; i != c; i = i \rightarrow d)
38
                   for (link j = i \rightarrow r; j != i; j = j \rightarrow r) {
                       j\rightarrow d\rightarrow u = j\rightarrow u; j\rightarrow u\rightarrow d = j\rightarrow d; ---cntc[j\rightarrow y];
41
        void resumeExact(link c) {
```

```
44
                 for (link i = c \rightarrow u; i != c; i = i \rightarrow u)
                       for (link j = i \rightarrow l; j != i; j = j \rightarrow l) {
45
46
                             j\rightarrow d\rightarrow u = j; j\rightarrow u\rightarrow d = j; ++cntc[j\rightarrow y];
47
                 c \rightarrow l \rightarrow r = c; c \rightarrow r \rightarrow l = c;
48
49
50
          void removeRepeat(link c) {
51
                 for (link i = c \rightarrow d; i != c; i = i \rightarrow d) {
                       i \rightarrow l \rightarrow r = i \rightarrow r; i \rightarrow r \rightarrow l = i \rightarrow l;
52
53
54
           void resumeRepeat(link c) {
55
56
                 for (link i = c \rightarrow u: i != c: i = i \rightarrow u) {
57
                       i \rightarrow l \rightarrow r = i; i \rightarrow r \rightarrow l = i;
58
59
60
          int calcH() {
61
                 int y, res = 0; ++stamp;
                 for (link c = head \rightarrow r; (y = c \rightarrow y) \leftarrow row \&\& c != head; <math>c = c \rightarrow r) {
                      if (vis[y] != stamp) {
63
                            vis[y] = stamp; ++res;
64
                            for (link i = c \rightarrow d; i != c; i = i \rightarrow d)
65
                                  for (link j = i \rightarrow r; j != i; j = j \rightarrow r) vis[j \rightarrow y] = stamp;
                      }
                } return res;
68
69
           void DFS(int dep) { if (dep + calcH() >= ans) return;
70
                if (head \rightarrow r \rightarrow y > nGE | l | head \rightarrow r == head) {
71
                      if (ans > dep) ans = dep; return;
72
73
                } link c = NULL;
                 for (link i = head \rightarrow r; i \rightarrow y \leftarrow nGE \& i != head; i = i \rightarrow r)
74
75
                      if (!c || cntc[i\rightarrowy] < cntc[c\rightarrowy]) c = i;
                 for (link i = c \rightarrow d; i != c; i = i \rightarrow d) {
76
77
                      removeRepeat(i);
78
                      for (link j = i \rightarrow r; j != i; j = j \rightarrow r) if (j \rightarrow y \leftarrow nGE) removeRepeat(j);
79
                       for (link j = i \rightarrow r; j != i; j = j \rightarrow r) if (j \rightarrow y > nGE) removeExact(base + j \rightarrow y);
                      DFS(dep + 1);
80
                       for (link j = i \rightarrow l; j != i; j = j \rightarrow l) if (j \rightarrow y > nGE) resumeExact(base + j \rightarrow y);
81
                       for (link j = i \rightarrow l; j != i; j = j \rightarrow l) if (j \rightarrow y \leftarrow nGE) resumeRepeat(j);
82
                       resumeRepeat(i);
84
85
           int solve() { build(); ans = INF; DFS(0); return ans; }
87 }
```

# 6.9 星期几判定

```
int getDay(int y, int m, int d) {
   if (m <= 2) m += 12, y—;
   if (y < 1752 || (y == 1752 && m < 9) || (y == 1752 && m == 9 && d < 3))
   return (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 + 5) % 7 + 1;
   return (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
   if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
  if (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1
```

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## 6.10 LCSequence Fast

#### 6.11 C Split

```
1 for (char *tok = strtok(ins, delimiters); tok; tok = strtok(NULL, delimiters))
2 puts(tok); // '会破怀京字符串ins'
```

## 6.12 builtin 系列

- int \_\_\_builtin\_ffs (unsigned int x) 返回 x 的最后一位 1 的是从后向前第几位, 比如 7368( 1110011001000) 返回 4.
- int \_\_\_builtin\_clz (unsigned int x) 返回前导的 0 的个数.
- int \_\_\_builtin\_ctz (unsigned int x) 返回后面的 0 个个数, 和 \_\_\_builtin\_clz 相对.
- int builtin popcount (unsigned int x) 返回二进制表示中 1 的个数.
- int \_\_\_builtin\_parity (unsigned int x) 返回 x 的奇偶校验位, 也就是 x 的 1 的个数模 2 的结果.

## 7 Templates

# 7.1 Eclipse 配置

Exec=env UBUNTU\_MENUPROXY= /opt/eclipse/eclipse preference general keys 把 word completion 设置成 alt+c, 把 content assistant 设置成 alt + /

#### 7.2 C++

```
bool nextChar(char &c) {
15
           if ((c = *ptr++) == 0) {
               int tmp = fread(buff, 1, MAX_BUFFER, stdin);
16
               buff[tmp] = 0; if (tmp == 0) return false;
               ptr = buff; c = *ptr++;
18
           } return true;
19
20
21
       bool nextUnsignedInt(unsigned int &x) {
           for (;;){if (!nextChar(c)) return false; if ('0'<=c && c<='9') break;}</pre>
22
23
           for (x=c-'0'; nextChar(c); x = x * 10 + c - '0') if (c < '0' | | c > '9') break;
24
           return true;
25
26
       bool nextInt(int &x) {
27
           for (;;) { if (!nextChar(c)) return false; if (c=='-' || ('0'<=c && c<='9')) break; }
           for ((c=='-')? (x=0,flag=true): (x=c-'0',flag=false); nextChar(c); x=x*10+c-'0')
               if (c<'0' || c>'9') break;
29
30
           if (flag) x=-x; return true;
31
32 };
33 #endif
```

#### 7.3 Java

```
1 import java.io.*;
2 import java.util.*;
3 import java.math.*;
5 public class Main {
       public void solve() {}
       public void run() {
           tokenizer = null; out = new PrintWriter(System.out);
           in = new BufferedReader(new InputStreamReader(System.in));
           solve();
           out.close();
12
13
       public static void main(String[] args) {
           new Main().run();
14
15
       public StringTokenizer tokenizer;
16
       public BufferedReader in;
       public PrintWriter out;
18
19
       public String next() {
           while (tokenizer == null || !tokenizer.hasMoreTokens()) {
               try { tokenizer = new StringTokenizer(in.readLine()); }
21
               catch (IOException e) { throw new RuntimeException(e); }
22
           } return tokenizer.nextToken();
23
24
25 }
```

# 7.4 gcc 配置

在.bashrc 中加入 export CXXFLAGS="-Wall -Wconversion -Wextra -g3"