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## 1 计算几何

### 1.1 二维计算几何基本操作

```

1 const double PI = 3.14159265358979323846264338327950288;
2 double arcsin(const double &a) { return a <= -1.0 ? -PI / 2 : (a >= 1.0 ? PI / 2 : asin(a)); }
3 double arccos(const double &a) { return a <= -1.0 ? PI : (a >= 1.0 ? 0 : acos(a)); }
4 struct point {
5     double x, y; // `something omitted`
6     point rot(const double &a) const { // `counterclockwise`
7         return point(x * cos(a) - y * sin(a), x * sin(a) + y * cos(a));}
8     point rot90() const { // `counterclockwise`
9         return point(-y, x);}
10    point project(const point &p1, const point &p2) const {
11        const point &q = *this;
12        return p1 + (p2 - p1) * (dot(p2 - p1, q - p1) / (p2 - p1).norm());}
13    bool onSeg(const point &a, const point &b) const { // `a, b inclusive`
14        const point &c = *this;
15        return sign(dot(a - c, b - c)) <= 0 && sign(dot(b - a, c - a)) == 0;}
16    double distLP(const point &p1, const point &p2) const { // `dist from *this to line p1→p2`
17        const point &q = *this;
18        return fabs(dot(p2 - p1, q - p1)) / (p2 - p1).len(); }
19    double distSP(const point &p1, const point &p2) const { // `dist from *this to segment [p1, p2]`
20        const point &q = *this;
21        if (dot(p2 - p1, q - p1) < EPS) return (q - p1).len();
22        if (dot(p1 - p2, q - p2) < EPS) return (q - p2).len();
23        return distLP(p1, p2); }
24    bool inAngle(const point &p1, const point &p2) const { // `det(p1, p2) $ge$ 0`
25        const point &q = *this; return det(p1, q) > -EPS && det(p2, q) < EPS;}
26 };
27 bool lineIntersect(const point &a, const point &b, const point &c, const point &d, point &e) {
28     double s1 = det(c - a, d - a), s2 = det(d - b, c - b);
29     if (!sign(s1 + s2)) return false;
30     e = (b - a) * (s1 / (s1 + s2)) + a; return true; }
31 int segIntersectCheck(const point &a, const point &b, const point &c, const point &d, point &o) {
32     static double s1, s2, s3, s4; static int iCnt;
33     int d1 = sign(s1 = det(b - a, c - a)), d2 = sign(s2 = det(b - a, d - a));
34     int d3 = sign(s3 = det(d - c, a - c)), d4 = sign(s4 = det(d - c, b - c));
35     if ((d1 ^ d2) == -2 && (d3 ^ d4) == -2) {
36         o = (c * s2 - d * s1) / (s2 - s1); return true; }
37     iCnt = 0;
38     if (d1 == 0 && c.onSeg(a, b)) o = c, ++iCnt;
39     if (d2 == 0 && d.onSeg(a, b)) o = d, ++iCnt;
40     if (d3 == 0 && a.onSeg(c, d)) o = a, ++iCnt;
41     if (d4 == 0 && b.onSeg(c, d)) o = b, ++iCnt;
42     return iCnt ? 2 : 0; // `不相交返回0, 严格相交返回1, 非严格相交返回2`
43 }
44 struct circle {
45     point o; double r, rSqure;
46     bool inside(const point &a) { // `非严格`
47         return (a - o).len() < r + EPS; }
48     bool contain(const circle &b) const { // `非严格`
49         return sign(b.r + (o - b.o).len() - r) <= 0; }
50     bool disjunct(const circle &b) const { // `非严格`
51         return sign(b.r + r - (o - b.o).len()) <= 0; }
52     int isCL(const point &p1, const point &p2, point &a, point &b) const {
53         double x = dot(p1 - o, p2 - p1), y = (p2 - p1).norm();
54         double d = x * x - y * ((p1 - o).norm() - rSqure);
55         if (d < -EPS) return 0; if (d < 0) d = 0;
56         point q1 = p1 - (p2 - p1) * (x / y);
57         point q2 = (p2 - p1) * (sqrt(d) / y);
58         a = q1 - q2; b = q1 + q2;
59         return q2.len() < EPS ? 1 : 2; }
60     int tanCP(const point &p, point &a, point &b) const { // `返回切点, 注意可能与 $p$ 重合`
61         double x = (p - o).norm(), d = x - rSqure;
62         if (d < -EPS) return 0; if (d < 0) d = 0;
63         point q1 = (p - o) * (rSqure / x);
64         point q2 = ((p - o) * (-r * sqrt(d) / x)).rot90();
65         a = o + (q1 - q2); b = o + (q1 + q2);
66         return q2.len() < EPS ? 1 : 2; }};
67 bool checkCrossCS(const circle &cir, const point &p1, const point &p2) { // `非严格`

```

```

68     const point &c = cir.o;
69     const double &r = cir.r;
70     return c.distSP(p1, p2) < r + EPS && (r < (c - p1).len() + EPS || r < (c - p2).len() + EPS); }
71 bool checkCrossCC(const circle &cir1, const circle &cir2) { // `非严格`
72     const double &r1 = cir1.r, &r2 = cir2.r, d = (cir1.o - cir2.o).len();
73     return d < r1 + r2 + EPS && fabs(r1 - r2) < d + EPS; }
74 int isCC(const circle &cir1, const circle &cir2, point &a, point &b) {
75     const point &c1 = cir1.o, &c2 = cir2.o;
76     double x = (c1 - c2).norm(), y = ((cir1.rSqure - cir2.rSqure) / x + 1) / 2;
77     double d = cir1.rSqure / x - y * y;
78     if (d < -EPS) return 0; if (d < 0) d = 0;
79     point q1 = c1 + (c2 - c1) * y;
80     point q2 = ((c2 - c1) * sqrt(d)).rot90();
81     a = q1 - q2; b = q1 + q2;
82     return q2.len() < EPS ? 1 : 2; }
83 vector<pair<point, point>> tanCC(const circle &cir1, const circle &cir2) {
84     // `注意: 如果只有三条切线, 即 $s1 = 1, s2 = 1$, 返回的切线可能重复, 切点没有问题`
85     vector<pair<point, point>> list;
86     if (cir1.contain(cir2) || cir2.contain(cir1)) return list;
87     const point &c1 = cir1.o, &c2 = cir2.o;
88     double r1 = cir1.r, r2 = cir2.r;
89     point p, a1, b1, a2, b2; int s1, s2;
90     if (sign(r1 - r2) == 0) {
91         p = c2 - c1; p = (p * (r1 / p.len())).rot90();
92         list.push_back(make_pair(c1 + p, c2 + p)); list.push_back(make_pair(c1 - p, c2 - p));
93     } else {
94         p = (c2 * r1 - c1 * r2) / (r1 - r2);
95         s1 = cir1.tanCP(p, a1, b1); s2 = cir2.tanCP(p, a2, b2);
96         if (s1 >= 1 && s2 >= 1) {
97             list.push_back(make_pair(a1, a2)); list.push_back(make_pair(b1, b2)); } }
98     p = (c1 * r2 + c2 * r1) / (r1 + r2);
99     s1 = cir1.tanCP(p, a1, b1); s2 = cir2.tanCP(p, a2, b2);
100    if (s1 >= 1 && s2 >= 1) {
101        list.push_back(make_pair(a1, a2)); list.push_back(make_pair(b1, b2)); }
102    return list; }
103 bool distConvexPin(const point &p1, const point &p2, const point &p3, const point &p4, const point &q) {
104     point o12 = (p1 - p2).rot90(), o23 = (p2 - p3).rot90(), o34 = (p3 - p4).rot90();
105     return (q - p1).inAngle(o12, o23) || (q - p3).inAngle(o23, o34)
106         || ((q - p2).inAngle(o23, p3 - p2) && (q - p3).inAngle(p2 - p3, o23)); }
107 }
108 double distConvexP(int n, point ps[], const point &q) { // `外部点到多边形各顶点的距离`
109     int left = 0, right = n;
110     while (right - left > 1) {
111         int mid = (left + right) / 2;
112         if (distConvexPin(ps[(left + n - 1) % n], ps[left], ps[mid], ps[(mid + 1) % n], q)) right = mid;
113         else left = mid; }
114     return q.distSP(ps[left], ps[right % n]); }
115 double areaCT(const circle &cir, point pa, point pb) {
116     pa = pa - cir.o; pb = pb - cir.o; double R = cir.r;
117     if (pa.len() < pb.len()) swap(pa, pb);
118     if (pb.len() < EPS) return 0; point pc = pb - pa;
119     double a = pa.len(), b = pb.len(), c = pc.len(), S, h, theta;
120     double cosB = dot(pb, pc) / b / c, B = acos(cosB);
121     double cosC = dot(pa, pb) / a / b, C = acos(cosC);
122     if (b > R) {
123         S = C * 0.5 * R * R; h = b * a * sin(C) / c;
124         if (h < R && B < PI * 0.5)
125             S += acos(h / R) * R * R - h * sqrt(R * R - h * h);
126     } else if (a > R) {
127         theta = PI - B - asin(sin(B) / R * b);
128         S = 0.5 * b * R * sin(theta) + (C - theta) * 0.5 * R * R;
129     } else S = 0.5 * sin(C) * b * a;
130     return S; }
131 circle minCircle(const point &a, const point &b) {
132     return circle((a + b) * 0.5, (b - a).len() * 0.5); }
133 circle minCircle(const point &a, const point &b, const point &c) { // `钝角三角形没有被考虑`
134     double a2((b - c).norm()), b2((a - c).norm()), c2((a - b).norm());
135     if (b2 + c2 <= a2 + EPS) return minCircle(b, c);
136     if (a2 + c2 <= b2 + EPS) return minCircle(a, c);
137     if (a2 + b2 <= c2 + EPS) return minCircle(a, b);
138     double A = 2.0 * (a.x - b.x), B = 2.0 * (a.y - b.y);
139     double D = 2.0 * (a.x - c.x), E = 2.0 * (a.y - c.y);
140     double C = a.norm() - b.norm(), F = a.norm() - c.norm();

```

```

141 point p((C * E - B * F) / (A * E - B * D), (A * F - C * D) / (A * E - B * D));
142 return circle(p, (p - a).len()); }
143 circle minCircle(point P[], int N) { // '1-based'
144 if (N == 1) return circle(P[1], 0.0);
145 random_shuffle(P + 1, P + N + 1); circle O = minCircle(P[1], P[2]);
146 Rep(i, 1, N) if(!O.inside(P[i])) { O = minCircle(P[1], P[i]);
147 Foru(j, 1, i) if(!O.inside(P[j])) { O = minCircle(P[1], P[j]);
148 Foru(k, 1, j) if(!O.inside(P[k])) O = minCircle(P[i], P[j], P[k]); }
149 } return O; }

```

## 1.2 圆的面积模板

```

1 struct Event { point p; double alpha; int add; // '构造函数省略'
2 bool operator < (const Event &other) const { return alpha < other.alpha; } };
3 void circleCover(circle *c, int N, double *area) { // '$area[k]$: 至少被覆盖$ks$次'
4 static bool overlap[MAXN][MAXN], g[MAXN][MAXN];
5 Rep(i, 0, N + 1) area[i] = 0.0; Rep(j, 1, N) Rep(j, 1, N) overlap[i][j] = c[i].contain(c[j]);
6 Rep(i, 1, N) Rep(j, 1, N) g[i][j] = !(overlap[i][j] || overlap[j][i] || c[i].disjunct(c[j]));
7 Rep(i, 1, N) { static Event events[MAXN * 2 + 1]; int totE = 0, cnt = 1;
8 Rep(j, 1, N) if (j != i && overlap[j][i]) ++cnt;
9 Rep(j, 1, N) if (j != i && g[i][j]) {
10 circle &a = c[i], &b = c[j]; double l = (a.o - b.o).norm();
11 double s = ((a.r - b.r) * (a.r + b.r) / l + 1) * 0.5;
12 double t = sqrt(-(l - sqrt(a.r - b.r)) * (l - sqrt(a.r + b.r))) / (l * l * 4.0);
13 point dir = b.o - a.o, nDir = point(-dir.y, dir.x);
14 point aa = a.o + dir * s + nDir * t;
15 point bb = a.o + dir * s - nDir * t;
16 double A = atan2(aa.y - a.o.y, aa.x - a.o.x);
17 double B = atan2(bb.y - a.o.y, bb.x - a.o.x);
18 events[totE++] = Event(bb, B, 1); events[totE++] = Event(aa, A, -1); if (B > A) ++cnt;
19 } if (totE == 0) { area[cnt] += PI * c[i].rSquare; continue; }
20 sort(events, events + totE); events[totE] = events[0];
21 Foru(j, 0, totE) {
22 cnt += events[j].add; area[cnt] += 0.5 * det(events[j].p, events[j + 1].p);
23 double theta = events[j + 1].alpha - events[j].alpha; if (theta < 0) theta += 2.0 * PI;
24 area[cnt] += 0.5 * c[i].rSquare * (theta - sin(theta));
25 }}}

```

## 1.3 多边形相关

```

1 struct Polygon { // stored in [0, n)
2 int n; point ps[MAXN];
3 Polygon cut(const point &a, const point &b) {
4 static Polygon res; static point o; res.n = 0;
5 for (int i = 0; i < n; ++i) {
6 int s1 = sign(det(ps[i] - a, b - a));
7 int s2 = sign(det(ps[(i + 1) % n] - a, b - a));
8 if (s1 <= 0) res.ps[res.n++] = ps[i];
9 if (s1 * s2 < 0) {
10 lineIntersect(a, b, ps[i], ps[(i + 1) % n], o);
11 res.ps[res.n++] = o;
12 }
13 } return res;
14 }
15 bool contain(const point &p) const { // 1 if on border or inner, 0 if outer
16 static point A, B; int res = 0;
17 for (int i = 0; i < n; ++i) {
18 A = ps[i]; B = ps[(i + 1) % n];
19 if (p.onSeg(A, B)) return 1;
20 if (sign(A.y - B.y) <= 0) swap(A, B);
21 if (sign(p.y - A.y) > 0) continue;
22 if (sign(p.y - B.y) <= 0) continue;
23 res += (int)(sign(det(B - p, A - p)) > 0);
24 } return res & 1;
25 }
26 #define qs(x) (ps[x] - ps[0])
27 bool convexContain(point p) const { // 'counter-clockwise'
28 point q = qs(n - 1); p = p - ps[0];
29 if (!p.inAngle(qs(1), q)) return false;

```

```

30 int L = 0, R = n - 1;
31 while (L + 1 < R) { int M((L + R) >> 1);
32 if (p.inAngle(qs(M), q)) L = M; else R = M;
33 } if (L == 0) return false; point l(qs(L)), r(qs(R));
34 return sign( fabs(det(l, p)) + fabs(det(p, r)) + fabs(det(r - l, p - l)) - det(l, r) ) == 0;
35 }
36 #undef qs
37 double isPLAtan2(const point &a, const point &b) {
38 double k = (b - a).alpha(); if (k < 0) k += 2 * PI;
39 return k;
40 }
41 point isPL_Get(const point &a, const point &b, const point &s1, const point &s2) {
42 double k1 = det(b - a, s1 - a), k2 = det(b - a, s2 - a);
43 if (sign(k1) == 0) return s1;
44 if (sign(k2) == 0) return s2;
45 return (s1 * k2 - s2 * k1) / (k2 - k1);
46 }
47 int isPL_Dic(const point &a, const point &b, int l, int r) {
48 int s = (det(b - a, ps[l] - a) < 0) ? -1 : 1;
49 while (l <= r) {
50 int mid = (l + r) / 2;
51 if (det(b - a, ps[mid] - a) * s <= 0) r = mid - 1;
52 else l = mid + 1;
53 }
54 return r + 1;
55 }
56 int isPL_Find(double k, double w[]) {
57 if (k <= w[0] || k > w[n - 1]) return 0;
58 int l = 0, r = n - 1, mid;
59 while (l <= r) {
60 mid = (l + r) / 2;
61 if (w[mid] >= k) r = mid - 1;
62 else l = mid + 1;
63 } return r + 1;
64 }
65 bool isPL(const point &a, const point &b, point &cp1, point &cp2) { // '$O(\log N)$'
66 static double w[MAXN * 2]; // 'pay attention to the array size'
67 for (int i = 0; i <= n; ++i) ps[i + n] = ps[i];
68 for (int i = 0; i < n; ++i) w[i + n] = isPLAtan2(ps[i], ps[i + 1]);
69 int i = isPL_Find(isPLAtan2(a, b), w);
70 int j = isPL_Find(isPLAtan2(b, a), w);
71 double k1 = det(b - a, ps[i] - a), k2 = det(b - a, ps[j] - a);
72 if (sign(k1) * sign(k2) > 0) return false; // 'no intersection'
73 if (sign(k1) == 0 || sign(k2) == 0) { // 'intersect with a point or a line in the convex'
74 if (sign(k1) == 0) {
75 if (sign(det(b - a, ps[i + 1] - a)) == 0) cp1 = ps[i], cp2 = ps[i + 1];
76 else cp1 = cp2 = ps[i];
77 return true;
78 }
79 if (sign(k2) == 0) {
80 if (sign(det(b - a, ps[j + 1] - a)) == 0) cp1 = ps[j], cp2 = ps[j + 1];
81 else cp1 = cp2 = ps[j];
82 }
83 return true;
84 }
85 if (i > j) swap(i, j);
86 int x = isPL_Dic(a, b, i, j), y = isPL_Dic(a, b, j, i + n);
87 cp1 = isPL_Get(a, b, ps[x - 1], ps[x]);
88 cp2 = isPL_Get(a, b, ps[y - 1], ps[y]);
89 return true;
90 }
91 double getI(const point &O) const {
92 if (n <= 2) return 0;
93 point G(0.0, 0.0);
94 double S = 0.0, I = 0.0;
95 for (int i = 0; i < n; ++i) {
96 const point &x = ps[i], &y = ps[(i + 1) % n];
97 double d = det(x, y);
98 G = G + (x + y) * d / 3.0;
99 S += d;
100 } G = G / S;
101 for (int i = 0; i < n; ++i) {
102 point x = ps[i] - G, y = ps[(i + 1) % n] - G;
103 I += fabs(det(x, y)) * (x.norm() + dot(x, y) + y.norm());
104 }

```

```

105     return I = I / 12.0 + fabs(S * 0.5) * (0 - G).norm();
106 }
107 };

```

## 1.4 直线与凸包求交点

```

1 int isPL(point a, point b, vector<point> &res) { // 点逆时针给出, 无三点共线
2     static double theta[MAXN];
3     for (int i = 0; i < n; ++i) theta[i] = (list[(i + 1) % n] - list[i]).atan2();
4     double delta = theta[0];
5     for (int i = 0; i < n; ++i) theta[i] = normalize(theta[i] - delta);
6     int x = lower_bound(theta, theta + n, normalize((b - a).atan2() - delta)) - theta;
7     int y = lower_bound(theta, theta + n, normalize((a - b).atan2() - delta)) - theta;
8     for (int k = 0; k <= 1; ++k, swap(a, b), swap(x, y)) {
9         if (y < x) y += n;
10        int l = x, r = y, m;
11        while (l + 1 < r) {
12            if (sign(det(b - a, list[(m = (l + r) / 2) % n] - a))) < 0) l = m;
13            else r = m;
14        }
15        l %= n, r %= n;
16        if (sign(det(b - a, list[r] - list[l])) == 0) {
17            if (sign(det(b - a, list[l] - a)) == 0)
18                return -1; // 直线与 $(list[l], list[r])$ 重合
19        }
20        else {
21            point p; lineIntersect(list[l], list[r], a, b, p);
22            if (p.onSeg(list[l], list[r]))
23                res.push_back(p);
24        }
25    }
26    return res.size();
27 }

```

## 1.5 半平面交

```

1 struct Border {
2     point p1, p2; double alpha;
3     Border() : p1(), p2(), alpha(0.0) {}
4     Border(const point &a, const point &b): p1(a), p2(b), alpha( atan2(p2.y - p1.y, p2.x - p1.x) ) {}
5     bool operator == (const Border &b) const { return sign(alpha - b.alpha) == 0; }
6     bool operator < (const Border &b) const {
7         int c = sign(alpha - b.alpha); if (c != 0) return c > 0;
8         return sign(det(b.p2 - b.p1, p1 - b.p1)) >= 0;
9     }
10 };
11 point isBorder(const Border &a, const Border &b) { // a and b should not be parallel
12     point is; lineIntersect(a.p1, a.p2, b.p1, b.p2, is); return is;
13 }
14 bool checkBorder(const Border &a, const Border &b, const Border &me) {
15     point is; lineIntersect(a.p1, a.p2, b.p1, b.p2, is);
16     return sign(det(me.p2 - me.p1, is - me.p1)) > 0;
17 }
18 double HPI(int N, Border border[]) {
19     static Border que[MAXN * 2 + 1]; static point ps[MAXN];
20     int head = 0, tail = 0, cnt = 0; // [head, tail)
21     sort(border, border + N); N = unique(border, border + N) - border;
22     for (int i = 0; i < N; ++i) {
23         Border &cur = border[i];
24         while (head + 1 < tail && !checkBorder(que[tail - 2], que[tail - 1], cur)) --tail;
25         while (head + 1 < tail && !checkBorder(que[head], que[head + 1], cur)) ++head;
26         que[tail++] = cur;
27     } while (head + 1 < tail && !checkBorder(que[tail - 2], que[tail - 1], que[head])) --tail;
28     while (head + 1 < tail && !checkBorder(que[head], que[head + 1], que[tail - 1])) ++head;
29     if (tail - head <= 2) return 0.0;
30     Foru(i, head, tail) ps[cnt++] = isBorder(que[i], que[(i + 1 == tail) ? (head) : (i + 1)]);
31     double area = 0; Foru(i, 0, cnt) area += det(ps[i], ps[(i + 1) % cnt]);
32     return fabs(area * 0.5); // or (-area * 0.5)
33 }

```

## 1.6 最大面积空凸包

```

1 inline bool toUpRight(const point &a, const point &b) {
2     int c = sign(b.y - a.y); if (c > 0) return true;
3     return c == 0 && sign(b.x - a.x) > 0;
4 }
5 inline bool cmpByPolarAngle(const point &a, const point &b) { // counter-clockwise, shorter first if they share
6     // the same polar angle
7     int c = sign(det(a, b)); if (c != 0) return c > 0;
8     return sign(b.len() - a.len()) > 0;
9 }
10 double maxEmptyConvexHull(int N, point p[]) {
11     static double dp[MAXN][MAXN];
12     static point vec[MAXN];
13     static int seq[MAXN]; // empty triangles formed with $(0, 0), vec[0], vec[seq[i]]$
14     double ans = 0.0;
15     Rep(0, 1, N) {
16         int totVec = 0;
17         Rep(i, 1, N) if (toUpRight(p[0], p[i])) vec[++totVec] = p[i] - p[0];
18         sort(vec + 1, vec + totVec + 1, cmpByPolarAngle);
19         Rep(i, 1, totVec) Rep(j, 1, totVec) dp[i][j] = 0.0;
20         Rep(k, 2, totVec) {
21             int i = k - 1;
22             while (i > 0 && sign( det(vec[k], vec[i]) ) == 0) --i;
23             int totSeq = 0;
24             for (int j; i > 0; i = j) {
25                 seq[++totSeq] = i;
26                 for (j = i - 1; j > 0 && sign(det(vec[i] - vec[k], vec[j] - vec[k])) > 0; --j);
27                 double v = det(vec[i], vec[k]) * 0.5;
28                 if (j > 0) v += dp[i][j];
29                 dp[k][i] = v;
30                 cMax(ans, v);
31             } for (int i = totSeq - 1; i >= 1; --i) cMax( dp[k][ seq[i] ], dp[k][seq[i + 1]] );
32         }
33     } return ans;
34 }

```

## 1.7 最近点对

```

1 int N; point p[maxn];
2 bool cmpByX(const point &a, const point &b) { return sign(a.x - b.x) < 0; }
3 bool cmpByY(const int &a, const int &b) { return p[a].y < p[b].y; }
4 double minimalDistance(point *c, int n, int *ys) {
5     double ret = 1e+20;
6     if (n < 20) {
7         Foru(i, 0, n) Foru(j, i + 1, n) cMin(ret, (c[i] - c[j]).len());
8         sort(ys, ys + n, cmpByY); return ret;
9     } static int mergeTo[maxn];
10     int mid = n / 2; double xmid = c[mid].x;
11     ret = min(minimalDistance(c, mid, ys), minimalDistance(c + mid, n - mid, ys + mid));
12     merge(ys, ys + mid, ys + mid, ys + n, mergeTo, cmpByY);
13     copy(mergeTo, mergeTo + n, ys);
14     Foru(i, 0, n) {
15         while (i < n && sign(fabs(p[ys[i]].x - xmid) - ret) > 0) ++i;
16         int cnt = 0;
17         Foru(j, i + 1, n)
18             if (sign(p[ys[j]].y - p[ys[i]].y - ret) > 0) break;
19             else if (sign(fabs(p[ys[j]].x - xmid) - ret) <= 0) {
20                 ret = min(ret, (p[ys[i]] - p[ys[j]]).len());
21                 if (++cnt >= 10) break;
22             }
23     } return ret;
24 }
25 double work() {
26     sort(p, p + n, cmpByX); Foru(i, 0, n) ys[i] = i; return minimalDistance(p, n, ys);
27 }

```

## 1.8 凸包与点集直径

```

1 vector<point> convexHull(int n, point ps[]) { // `counter-clockwise, strict`
2     static point qs[MAXN * 2];
3     sort(ps, ps + n, cmpByXY);
4     if (n <= 2) return vector<point>(ps, ps + n);
5     int k = 0;
6     for (int i = 0; i < n; qs[k++] = ps[i++])
7         while (k > 1 && det(qs[k-1] - qs[k-2], ps[i] - qs[k-1]) < EPS) k--;
8     for (int i = n-2, t = k; i >= 0; qs[k++] = ps[i--])
9         while (k > t && det(qs[k-1] - qs[k-2], ps[i] - qs[k-1]) < EPS) k--;
10    return vector<point>(qs, qs + k);
11 }
12 double convexDiameter(int n, point ps[]) {
13     if (n < 2) return 0; if (n == 2) return (ps[1] - ps[0]).len();
14     double k, ans = 0;
15     for (int x = 0, y = 1, nx, ny; x < n; ++x) {
16         for (nx = (x == n-1) ? 0 : (x+1); ; y = ny) {
17             ny = (y == n-1) ? 0 : (y+1);
18             if (sign(k = det(ps[nx] - ps[x], ps[ny] - ps[y])) <= 0) break;
19             ans = max(ans, (ps[x] - ps[ny]).len());
20             if (sign(k) == 0) ans = max(ans, (ps[x] - ps[ny]).len());
21         }
22     }
23 }

```

## 1.9 Farmland

```

1 struct node { int begin[MAXN], *end; } a[MAXN]; // 按对 $p[i]$ 的极角的 atan2 值排序
2 bool check(int n, point p[], int b1, int b2, bool vis[MAXN][MAXN]) {
3     static pii l[MAXN * 2 + 1]; static bool used[MAXN];
4     int tp(0), *k, p, p1, p2; double area(0.0);
5     for (l[0] = pii(b1, b2); ; ) {
6         vis[p1 = l[tp].first][p2 = l[tp].second] = true;
7         area += det(p[p1], p[p2]);
8         for (k = a[p2].begin; k != a[p2].end; ++k) if (*k == p1) break;
9         k = (k == a[p2].begin) ? (a[p2].end-1) : (k-1);
10        if ((l[++tp] = pii(p2, *k)) == l[0]) break;
11    } if (sign(area) < 0 || tp < 3) return false;
12    Rep(i, 1, n) used[i] = false;
13    for (int i = 0; i < tp; ++i) if (used[p = l[i].first]) return false; else used[p] = true;
14    return true; // `a face with tp vertices`
15 }
16 int countFaces(int n, point p[]) {
17     static bool vis[MAXN][MAXN]; int ans = 0;
18     Rep(x, 1, n) Rep(y, 1, n) vis[x][y] = false;
19     Rep(x, 1, n) for (int *itr = a[x].begin; itr != a[x].end; ++itr) if (!vis[x][*itr])
20         if (check(n, p, x, *itr, vis)) ++ans;
21     return ans;
22 }

```

## 1.10 Voronoi 图

不能有重点, 点数应当不小于 2

```

1 #define Oi(e) ((e)->oi)
2 #define Dt(e) ((e)->dt)
3 #define On(e) ((e)->on)
4 #define Op(e) ((e)->op)
5 #define Dn(e) ((e)->dn)
6 #define Dp(e) ((e)->dp)
7 #define Other(e, p) ((e)->oi == p ? (e)->dt : (e)->oi)
8 #define Next(e, p) ((e)->oi == p ? (e)->on : (e)->dn)
9 #define Prev(e, p) ((e)->oi == p ? (e)->op : (e)->dp)
10 #define V(p1, p2, u, v) (u = p2->x - p1->x, v = p2->y - p1->y)
11 #define C2(u1, v1, u2, v2) (u1 * v2 - v1 * u2)
12 #define C3(p1, p2, p3) ((p2->x - p1->x) * (p3->y - p1->y) - (p2->y - p1->y) * (p3->x - p1->x))
13 #define Dot(u1, v1, u2, v2) (u1 * u2 + v1 * v2)
14 #define dis(a,b) (sqrt((a->x - b->x) * (a->x - b->x) + (a->y - b->y) * (a->y - b->y)))
15 const int maxn = 110024;
16 const int aix = 4;
17 const double eps = 1e-7;
18 int n, M, k;

```

```

19 struct gEdge {
20     int u, v; double w;
21     bool operator <(const gEdge &e1) const { return w < e1.w - eps; }
22 } E[aix * maxn], MST[maxn];
23 struct point {
24     double x, y; int index; edge *in;
25     bool operator <(const point &p1) const { return x < p1.x - eps || (abs(x - p1.x) <= eps && y < p1.y - eps); }
26 };
27 struct edge { point *oi, *dt; edge *on, *op, *dn, *dp; };
28
29 point p[maxn], *Q[maxn];
30 edge mem[aix * maxn], *elist[aix * maxn];
31 int nfree;
32 void Alloc_memory() { nfree = aix * n; edge *e = mem; for (int i = 0; i < nfree; i++) elist[i] = e++; }
33 void Splice(edge *a, edge *b, point *v) {
34     edge *next;
35     if (Oi(a) == v) next = On(a), On(a) = b; else next = Dn(a), Dn(a) = b;
36     if (Oi(next) == v) Op(next) = b; else Dp(next) = b;
37     if (Oi(b) == v) On(b) = next, Op(b) = a; else Dn(b) = next, Dp(b) = a;
38 }
39 edge *Make_edge(point *u, point *v) {
40     edge *e = elist[---nfree];
41     e->on = e->op = e->dn = e->dp = e; e->oi = u; e->dt = v;
42     if (!u->in) u->in = e;
43     if (!v->in) v->in = e;
44     return e;
45 }
46 edge *Join(edge *a, point *u, edge *b, point *v, int side) {
47     edge *e = Make_edge(u, v);
48     if (side == 1) {
49         if (Oi(a) == u) Splice(Op(a), e, u);
50         else Splice(Dp(a), e, u);
51         Splice(b, e, v);
52     } else {
53         Splice(a, e, u);
54         if (Oi(b) == v) Splice(Op(b), e, v);
55         else Splice(Dp(b), e, v);
56     }
57     return e;
58 }
59 void Remove(edge *e) {
60     point *u = Oi(e), *v = Dt(e);
61     if (u->in == e) u->in = e->on;
62     if (v->in == e) v->in = e->dn;
63     if (Oi(e->on) == u) e->on->op = e->op; else e->on->dp = e->op;
64     if (Oi(e->op) == u) e->op->on = e->on; else e->op->dn = e->on;
65     if (Oi(e->dn) == v) e->dn->op = e->dp; else e->dn->dp = e->dp;
66     if (Oi(e->dp) == v) e->dp->on = e->dn; else e->dp->dn = e->dn;
67     elist[nfree++] = e;
68 }
69 void Low_tangent(edge *e_l, point *o_l, edge *e_r, point *o_r, edge **l_low, point **OL, edge **r_low, point **OR) {
70     }
71     for (point *d_l = Other(e_l, o_l), *d_r = Other(e_r, o_r); ; )
72         if (C3(o_l, o_r, d_l) < -eps) e_l = Prev(e_l, d_l), o_l = d_l, d_l = Other(e_l, o_l);
73         else if (C3(o_l, o_r, d_r) < -eps) e_r = Next(e_r, d_r), o_r = d_r, d_r = Other(e_r, o_r);
74         else break;
75     *OL = o_l, *OR = o_r; *l_low = e_l, *r_low = e_r;
76 }
77 void Merge(edge *lr, point *s, edge *rl, point *u, edge **tangent) {
78     double l1, l2, l3, l4, r1, r2, r3, r4, cot_L, cot_R, u1, v1, u2, v2, n1, cot_n, P1, cot_P;
79     point *O, *D, *OR, *OL; edge *B, *L, *R;
80     Low_tangent(lr, s, rl, u, &L, &OL, &R, &OR);
81     for (*tangent = B = Join(L, OL, R, OR, 0), O = OL, D = OR; ; ) {
82         edge *El = Next(B, O), *Er = Prev(B, D), *next, *prev;
83         point *l = Other(El, O), *r = Other(Er, D);
84         V(l, O, l1, l2); V(l, D, l3, l4); V(r, O, r1, r2); V(r, D, r3, r4);
85         double cl = C2(l1, l2, l3, l4), cr = C2(r1, r2, r3, r4);
86         bool BL = cl > eps, BR = cr > eps;
87         if (!BL && !BR) break;
88         if (BL) {
89             double dl = Dot(l1, l2, l3, l4);
90             for (cot_L = dl / cl; ; Remove(El), El = next, cot_L = cot_n) {
91                 next = Next(El, O); V(Other(next, O), O, u1, v1); V(Other(next, O), D, u2, v2);
92                 n1 = C2(u1, v1, u2, v2); if (!(n1 > eps)) break;
93                 cot_n = Dot(u1, v1, u2, v2) / n1;
94                 if (cot_n > cot_L) break;

```

```

93     }
94   } if (BR) {
95     double dr = Dot(r1, r2, r3, r4);
96     for (cot_R = dr / cr; ; Remove(Er), Er = prev, cot_R = cot_P) {
97       prev = Prev(Er, D); V(Other(prev, D), 0, u1, v1); V(Other(prev, D), D, u2, v2);
98       P1 = C2(u1, v1, u2, v2); if (!(P1 > eps)) break;
99       cot_P = Dot(u1, v1, u2, v2) / P1;
100      if (cot_P > cot_R) break;
101    }
102    } l = Other(El, 0); r = Other(Er, D);
103    if (!BL || (BL && BR && cot_R < cot_L)) B = Join(B, 0, Er, r, 0), D = r;
104    else B = Join(El, l, B, D, 0), O = l;
105  }
106 }
107 void Divide(int s, int t, edge **L, edge **R) {
108   edge *a, *b, *c, *ll, *lr, *rl, *rr, *tangent;
109   int n = t - s + 1;
110   if (n == 2) *L = *R = Make_edge(Q[s], Q[t]);
111   else if (n == 3) {
112     a = Make_edge(Q[s], Q[s + 1]), b = Make_edge(Q[s + 1], Q[t]);
113     Splice(a, b, Q[s + 1]);
114     double v = C3(Q[s], Q[s + 1], Q[t]);
115     if (v > eps) c = Join(a, Q[s], b, Q[t], 0), *L = a, *R = b;
116     else if (v < -eps) c = Join(a, Q[s], b, Q[t], 1), *L = c, *R = c;
117     else *L = a, *R = b;
118   } else if (n > 3) {
119     int split = (s + t) / 2;
120     Divide(s, split, &ll, &lr); Divide(split + 1, t, &rl, &rr);
121     Merge(lr, Q[split], rl, Q[split + 1], &tangent);
122     if (Oi(tangent) == Q[s]) ll = tangent;
123     if (Di(tangent) == Q[t]) rr = tangent;
124     *L = ll; *R = rr;
125   }
126 }
127 void Make_Graph() {
128   edge *start, *e; point *u, *v;
129   for (int i = 0; i < n; i++) {
130     start = e = (u = &p[i]) -> in;
131     do { v = Other(e, u);
132         if (u < v) E[M++] .u = (u - p, v - p, dis(u, v)); // M < aix * maxn
133       } while ((e = Next(e, u)) != start);
134   }
135 }
136 int b[maxn];
137 int Find(int x) { while (x != b[x]) { b[x] = b[b[x]]; x = b[x]; } return x; }
138 void Kruskal() {
139   memset(b, 0, sizeof(b)); sort(E, E + M);
140   for (int i = 0; i < n; i++) b[i] = i;
141   for (int i = 0, kk = 0; i < M && kk < n - 1; i++) {
142     int m1 = Find(E[i].u), m2 = Find(E[i].v);
143     if (m1 != m2) b[m1] = m2, MST[kk++] = E[i];
144   }
145 }
146 void solve() {
147   scanf("%d", &n);
148   for (int i = 0; i < n; i++) scanf("%lf%lf", &p[i].x, &p[i].y), p[i].index = i, p[i].in = NULL;
149   Alloc_memory(); sort(p, p + n);
150   for (int i = 0; i < n; i++) Q[i] = p + i;
151   edge *L, *R; Divide(0, n - 1, &L, &R);
152   M = 0; Make_Graph(); Kruskal();
153 }
154 int main() { solve(); return 0; }

```

### 1.11 四边形双费马点

```

1  typedef complex<double> Tpoint;
2  const double eps = 1e-8;
3  const double sqrt3 = sqrt(3.0);
4  bool cmp(const Tpoint &a, const Tpoint &b) {
5    return a.real() < b.real() - eps || (a.real() < b.real() + eps && a.imag() < b.imag());
6  }
7  Tpoint rotate(const Tpoint &a, const Tpoint &b, const Tpoint &c) {
8    Tpoint d = b - a; d = Tpoint(-d.imag(), d.real());

```

```

9    if (Sign(cross(a, b, c)) == Sign(cross(a, b, a + d))) d *= -1.0;
10   return unit(d);
11 }
12 Tpoint p[10], a[10], b[10];
13 int N, T;
14 double totlen(const Tpoint &p, const Tpoint &a, const Tpoint &b, const Tpoint &c) {
15   return abs(p - a) + abs(p - b) + abs(p - c);
16 }
17 double fermat(const Tpoint &x, const Tpoint &y, const Tpoint &z, Tpoint &cp) {
18   a[0] = a[3] = x; a[1] = a[4] = y; a[2] = a[5] = z;
19   double len = 1e100, len2;
20   for (int i = 0; i < 3; i++) {
21     len2 = totlen(a[i], x, y, z);
22     if (len2 < len) len = len2, cp = a[i];
23   }
24   for (int i = 0; i < 3; i++) {
25     b[i] = rotate(a[i + 1], a[i], a[i + 2]);
26     b[i] = (a[i + 1] + a[i]) / 2.0 + b[i] * (abs(a[i + 1] - a[i]) * sqrt3 / 2.0);
27   }
28   b[3] = b[0];
29   Tpoint cp2 = intersect(b[0], a[2], b[1], a[3]);
30   len2 = totlen(cp2, x, y, z);
31   if (len2 < len) len = len2, cp = cp2;
32   return len;
33 }
34 double getans(const Tpoint &a) {
35   double len = 0; for (int i = 0; i < N; i++) len += abs(a - p[i]);
36   return len;
37 }
38 double mindist(const Tpoint &p, const Tpoint &a, const Tpoint &b, const Tpoint &c, const Tpoint &d) {
39   return min(min(abs(p - a), abs(p - b)), min(abs(p - c), abs(p - d)));
40 }
41 int main() {
42   N = 4;
43   for (cin >> T; T; T--) {
44     double ret = 1e100, len_cur, len_before, len1, len2, len;
45     Tpoint cp, cp1, cp2;
46     Foru(i, 0, N) cin >> p[i];
47     Foru(i, 0, N) ret = min(ret, getans(p[i]));
48     Foru(i, 1, N) Foru(j, 1, N) if (j != i) Foru(k, 1, N) if (k != i && k != j) {
49       cMin(ret, abs(p[0] - p[i]) + abs(p[j] - p[k])
50         + min(min(abs(p[0] - p[j]), abs(p[0] - p[k])),
51             min(abs(p[i] - p[j]), abs(p[i] - p[k]))
52         ));
53       ret = min(ret, getans(intersect(p[0], p[i], p[j], p[k])));
54     }
55     Foru(i, 0, N) Foru(j, i + 1, N) Foru(k, j + 1, N) {
56       double len = fermat(p[i], p[j], p[k], cp);
57       ret = min(ret, len + mindist(p[6 - i - j - k], p[i], p[j], p[k], cp));
58     }
59     sort(p, p + N, cmp);
60     for (int i = 1; i < N; i++) {
61       cp1 = (p[0] + p[i]) / 2.0;
62       int j, k;
63       for (j = 1; j < N && j == i; j++);
64       for (k = 6 - i - j, len_before = 1e100; ; ) {
65         len1 = fermat(cp1, p[j], p[k], cp2);
66         len1 = fermat(cp2, p[0], p[i], cp1);
67         len = len1 + abs(cp2 - p[j]) + abs(cp2 - p[k]);
68         if (len < len_before - (1e-6)) len_before = len;
69         else break;
70       } ret = min(ret, len_before);
71     } printf("%.4f\n", ret);
72   }
73   return 0;
74 }

```

### 1.12 三角形和四边形的费马点

- 费马点: 距几个顶点距离之和最小的点
- 三角形:



- 若每个角都小于  $120^\circ$ : 以每条边向外作正三角形, 得到  $\triangle ABF$ ,  $\triangle BCD$ ,  $\triangle CAE$ , 连接  $AD$ ,  $BE$ ,  $CF$ , 三线必共点于费马点. 该点对三边的张角必然是  $120^\circ$ , 也必然是三个三角形外接圆的交点

- 否则费马点一定是那个大于等于  $120^\circ$  的顶角

- 四边形:

- 在凸四边形中, 费马点为对角线的交点
- 在凹四边形中, 费马点位凹顶点

### 1.13 三维计算几何基本操作

```

1 struct point { double x, y, z; // something omitted
2   friend point det(const point &a, const point &b) {
3     return point(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z, a.x * b.y - a.y * b.x);
4   }
5   friend double mix(const point &a, const point &b, const point &c) {
6     return a.x * b.y * c.z + a.y * b.z * c.x + a.z * b.x * c.y - a.z * b.y * c.x - a.x * b.z * c.y - a.y * b.x * c.z;
7   }
8   double distLP(const point &p1, const point &p2) const {
9     return det(p2 - p1, *this - p1).len() / (p2 - p1).len();
10  }
11  double distFP(const point &p1, const point &p2, const point &p3) const {
12    point n = det(p2 - p1, p3 - p1); return fabs( dot(n, *this - p1) / n.len() );
13  }
14 };
15 double distLL(const point &p1, const point &p2, const point &q1, const point &q2) {
16   point p = q1 - p1, u = p2 - p1, v = q2 - q1;
17   double d = u.norm() * v.norm() - dot(u, v) * dot(u, v);
18   if (sign(d) == 0) return p1.distLP(q1, q2);
19   double s = (dot(p, u) * v.norm() - dot(p, v) * dot(u, v)) / d;
20   return (p1 + u * s).distLP(q1, q2);
21 }
22 double distSS(const point &p1, const point &p2, const point &q1, const point &q2) {
23   point p = q1 - p1, u = p2 - p1, v = q2 - q1;
24   double d = u.norm() * v.norm() - dot(u, v) * dot(u, v);
25   if (sign(d) == 0) return min( min((p1 - q1).len(), (p1 - q2).len()),
26                                min((p2 - q1).len(), (p2 - q2).len()));
27   double s1 = (dot(p, u) * v.norm() - dot(p, v) * dot(u, v)) / d;
28   double s2 = (dot(p, v) * u.norm() - dot(p, u) * dot(u, v)) / d;
29   if (s1 < 0.0) s1 = 0.0; if (s1 > 1.0) s1 = 1.0;
30   if (s2 < 0.0) s2 = 0.0; if (s2 > 1.0) s2 = 1.0;
31   point r1 = p1 + u * s1; point r2 = q1 + v * s2;
32   return (r1 - r2).len();
33 }
34 bool isFL(const point &p, const point &o, const point &q1, const point &q2, point &res) {
35   double a = dot(o, q2 - p), b = dot(o, q1 - p), d = a - b;
36   if (sign(d) == 0) return false;
37   res = (q1 * a - q2 * b) / d;
38   return true;
39 }
40 bool isFF(const point &p1, const point &o1, const point &p2, const point &o2, point &a, point &b) {
41   point e = det(o1, o2), v = det(o1, e);
42   double d = dot(o2, v); if (sign(d) == 0) return false;
43   point q = p1 + v * (dot(o2, p2 - p1) / d);
44   a = q; b = q + e;
45   return true;
46 }

```

### 1.14 凸多面体切割

```

1 vector<vector<point>> > convexCut(const vector<vector<point>> &pss, const point &p, const point &o) {
2   vector<vector<point>> > res;
3   vector<point> sec;
4   for (unsigned itr = 0, size = pss.size(); itr < size; ++itr) {
5     const vector<point> &ps = pss[itr];
6     int n = ps.size();
7     vector<point> qs;

```

```

8     bool dif = false;
9     for (int i = 0; i < n; ++i) {
10      int d1 = sign( dot(o, ps[i] - p) );
11      int d2 = sign( dot(o, ps[(i + 1) % n] - p) );
12      if (d1 <= 0) qs.push_back(ps[i]);
13      if (d1 * d2 < 0) {
14        point q;
15        isFL(p, o, ps[i], ps[(i + 1) % n], q); // must return true
16        qs.push_back(q);
17        sec.push_back(q);
18      }
19      if (d1 == 0) sec.push_back(ps[i]);
20      else dif = true;
21      dif |= dot(o, det(ps[(i + 1) % n] - ps[i], ps[(i + 2) % n] - ps[i])) < -EPS;
22    }
23    if (!qs.empty() && dif)
24      res.insert(res.end(), qs.begin(), qs.end());
25  }
26  if (!sec.empty()) {
27    vector<point> tmp( convexHull2D(sec, o) );
28    res.insert(res.end(), tmp.begin(), tmp.end());
29  }
30  return res;
31 }
32
33 vector<vector<point>> > initConvex() {
34   vector<vector<point>> > pss(6, vector<point>(4));
35   pss[0][0] = pss[1][0] = pss[2][0] = point(-INF, -INF, -INF);
36   pss[0][3] = pss[1][1] = pss[5][2] = point(-INF, -INF, INF);
37   pss[0][1] = pss[2][3] = pss[4][2] = point(-INF, INF, -INF);
38   pss[0][2] = pss[5][3] = pss[4][1] = point(-INF, INF, INF);
39   pss[1][3] = pss[2][1] = pss[3][2] = point( INF, -INF, -INF);
40   pss[1][2] = pss[5][1] = pss[3][3] = point( INF, -INF, INF);
41   pss[2][2] = pss[4][3] = pss[3][1] = point( INF, INF, -INF);
42   pss[5][0] = pss[4][0] = pss[3][0] = point( INF, INF, INF);
43   return pss;
44 }

```

### 1.15 三维凸包

不能有重点

```

1 namespace ConvexHull3D {
2   #define volume(a, b, c, d) (mix(ps[b] - ps[a], ps[c] - ps[a], ps[d] - ps[a]))
3   vector<Facet> getHull(int n, point ps[]) {
4     static int mark[MAXN][MAXN], a, b, c;
5     int stamp = 0;
6     bool exist = false;
7     vector<Facet> facet;
8     random_shuffle(ps, ps + n);
9     for (int i = 2; i < n && !exist; i++) {
10      point ndir = det(ps[0] - ps[i], ps[1] - ps[i]);
11      if (ndir.len() < EPS) continue;
12      swap(ps[i], ps[2]);
13      for (int j = i + 1; j < n && !exist; j++)
14        if (sign(volume(0, 1, 2, j)) != 0) {
15          exist = true;
16          swap(ps[j], ps[3]);
17          facet.push_back(Facet(0, 1, 2));
18          facet.push_back(Facet(0, 2, 1));
19        }
20    }
21    if (!exist) return ConvexHull2D(n, ps);
22    for (int i = 0; i < n; ++i)
23      for (int j = 0; j < n; ++j)
24        mark[i][j] = 0;
25    stamp = 0;
26    for (int v = 3; v < n; ++v) {
27      vector<Facet> tmp;
28      ++stamp;
29      for (unsigned i = 0; i < facet.size(); i++) {
30        a = facet[i].a;
31        b = facet[i].b;
32        c = facet[i].c;

```

```

33         if (sign(volume(v, a, b, c)) < 0)
34             mark[a][b] = mark[a][c] =
35             mark[b][a] = mark[b][c] =
36             mark[c][a] = mark[c][b] = stamp;
37         else tmp.push_back(facet[i]);
38     } facet = tmp;
39     for (unsigned i = 0; i < tmp.size(); i++) {
40         a = facet[i].a; b = facet[i].b; c = facet[i].c;
41         if (mark[a][b] == stamp) facet.push_back(Facet(b, a, v));
42         if (mark[b][c] == stamp) facet.push_back(Facet(c, b, v));
43         if (mark[c][a] == stamp) facet.push_back(Facet(a, c, v));
44     }
45     } return facet;
46 }
47 #undef volume
48 }
49 namespace Gravity {
50     using ConvexHull3D::Facet;
51     point findG(point ps[], const vector<Facet> &facet) {
52         double ws = 0; point res(0.0, 0.0, 0.0), o = ps[ facet[0].a ];
53         for (int i = 0, size = facet.size(); i < size; ++i) {
54             const point &a = ps[ facet[i].a ], &b = ps[ facet[i].b ], &c = ps[ facet[i].c ];
55             point p = (a + b + c + o) * 0.25;
56             double w = mix(a - o, b - o, c - o);
57             ws += w;
58             res = res + p * w;
59         } res = res / ws;
60         return res;
61     }
62 }

```

## 1.16 球面点表面点距离

```

1 double distOnBall(double lati1, double longi1, double lati2, double longi2, double R) {
2     lati1 *= PI / 180; longi1 *= PI / 180;
3     lati2 *= PI / 180; longi2 *= PI / 180;
4     double x1 = cos(lati1) * sin(longi1);
5     double y1 = cos(lati1) * cos(longi1);
6     double z1 = sin(lati1);
7     double x2 = cos(lati2) * sin(longi2);
8     double y2 = cos(lati2) * cos(longi2);
9     double z2 = sin(lati2);
10    double theta = acos(x1 * x2 + y1 * y2 + z1 * z2);
11    return R * theta;
12 }

```

## 1.17 长方体表面点距离

```

1 int r;
2 void turn(int i, int j, int x, int y, int z, int x0, int y0, int L, int W, int H) {
3     if (z == 0) r = min(r, x * x + y * y);
4     else {
5         if (i >= 0 && i < 2) turn(i + 1, j, x0 + L + z, y, x0 + L - x, x0 + L, y0, H, W, L);
6         if (j >= 0 && j < 2) turn(i, j + 1, x, y0 + W + z, y0 + W - y, x0, y0 + W, L, H, W);
7         if (i <= 0 && i > -2) turn(i - 1, j, x0 - z, y, x - x0, x0 - H, y0, H, W, L);
8         if (j <= 0 && j > -2) turn(i, j - 1, x, y0 - z, y - y0, x0, y0 - H, L, H, W);
9     }
10 }
11 int calc(int L, int H, int W, int x1, int y1, int z1, int x2, int y2, int z2) {
12     if (z1 != 0 && z1 != H)
13         if (y1 == 0 || y1 == W) swap(y1, z1), swap(y2, z2), swap(W, H);
14         else swap(x1, z1), swap(x2, z2), swap(L, H);
15     if (z1 == H) z1 = 0, z2 = H - z2;
16     r = INF; turn(0, 0, x2 - x1, y2 - y1, z2, -x1, -y1, L, W, H);
17     return r;
18 }

```

## 1.18 最小覆盖球

```

1 int outCnt; point out[4], res; double radius;
2 void ball() {
3     static point q[3];
4     static double m[3][3], sol[3], L[3], det;
5     int i, j; res = point(0.0, 0.0, 0.0); radius = 0.0;
6     switch (outCnt) {
7         case 1: res = out[0]; break;
8         case 2: res = (out[0] + out[1]) * 0.5; radius = (res - out[0]).norm();
9             break;
10        case 3:
11            q[0] = out[1] - out[0]; q[1] = out[2] - out[0];
12            for (i = 0; i < 2; ++i) for (j = 0; j < 2; ++j)
13                m[i][j] = dot(q[i], q[j]) * 2.0;
14            for (i = 0; i < 2; ++i) sol[i] = dot(q[i], q[i]);
15            det = m[0][0] * m[1][1] - m[0][1] * m[1][0];
16            if (sign(det) == 0) return;
17            L[0] = (sol[0] * m[1][1] - sol[1] * m[0][1]) / det;
18            L[1] = (sol[1] * m[0][0] - sol[0] * m[1][0]) / det;
19            res = out[0] + q[0] * L[0] + q[1] * L[1];
20            radius = (res - out[0]).norm();
21            break;
22        case 4:
23            q[0] = out[1] - out[0]; q[1] = out[2] - out[0]; q[2] = out[3] - out[0];
24            for (i = 0; i < 3; ++i) for (j = 0; j < 3; ++j) m[i][j] = dot(q[i], q[j]) * 2;
25            for (i = 0; i < 3; ++i) sol[i] = dot(q[i], q[i]);
26            det = m[0][0] * m[1][1] * m[2][2] + m[0][1] * m[1][2] * m[2][0]
27                + m[0][2] * m[2][1] * m[1][0] - m[0][2] * m[1][1] * m[2][0]
28                - m[0][1] * m[1][0] * m[2][2] - m[0][0] * m[1][2] * m[2][1];
29            if (sign(det) == 0) return;
30            for (j = 0; j < 3; ++j) { for (i = 0; i < 3; ++i) m[i][j] = sol[i];
31                L[j] = (m[0][0] * m[1][1] * m[2][2] + m[0][1] * m[1][2] * m[2][0]
32                    + m[0][2] * m[2][1] * m[1][0] - m[0][2] * m[1][1] * m[2][0]
33                    - m[0][1] * m[1][0] * m[2][2] - m[0][0] * m[1][2] * m[2][1]) / det;
34                for (i = 0; i < 3; ++i) m[i][j] = dot(q[i], q[j]) * 2;
35            } res = out[0];
36            for (i = 0; i < 3; ++i) res += q[i] * L[i]; radius = (res - out[0]).norm();
37        }
38    }
39    void minball(int n, point pt[]) {
40        ball();
41        if (outCnt < 4) for (int i = 0; i < n; ++i)
42            if ((res - pt[i]).norm() > +radius + EPS) {
43                out[outCnt] = pt[i]; ++outCnt; minball(i, pt); --outCnt;
44                if (i > 0) {
45                    point Tt = pt[i];
46                    memmove(&pt[1], &pt[0], sizeof(point) * i);
47                    pt[0] = Tt;
48                }
49            }
50    }
51    pair<point, double> main(int npoint, point pt[]) { // 0-based
52        random_shuffle(pt, pt + npoint); radius = -1;
53        for (int i = 0; i < npoint; i++) { if ((res - pt[i]).norm() > EPS + radius) {
54            outCnt = 1; out[0] = pt[i]; minball(i, pt); } }
55        return make_pair(res, sqrt(radius));
56    }

```

## 1.19 三维向量操作矩阵

- 绕单位向量  $u = (u_x, u_y, u_z)$  右手方向旋转  $\theta$  度的矩阵:

$$\begin{bmatrix} \cos \theta + u_z^2(1 - \cos \theta) & u_x u_y(1 - \cos \theta) - u_z \sin \theta & u_x u_z(1 - \cos \theta) + u_y \sin \theta \\ u_y u_x(1 - \cos \theta) + u_z \sin \theta & \cos \theta + u_y^2(1 - \cos \theta) & u_y u_z(1 - \cos \theta) - u_x \sin \theta \\ u_z u_x(1 - \cos \theta) - u_y \sin \theta & u_z u_y(1 - \cos \theta) + u_x \sin \theta & \cos \theta + u_z^2(1 - \cos \theta) \end{bmatrix}$$

$$= \cos \theta I + \sin \theta \begin{bmatrix} 0 & -u_z & u_y \\ u_z & 0 & -u_x \\ -u_y & u_x & 0 \end{bmatrix} + (1 - \cos \theta) \begin{bmatrix} u_x^2 & u_x u_y & u_x u_z \\ u_y u_x & u_y^2 & u_y u_z \\ u_z u_x & u_z u_y & u_z^2 \end{bmatrix}$$

- 点  $a$  绕单位向量  $u = (u_x, u_y, u_z)$  右手方向旋转  $\theta$  度的对应点为  $a' = a \cos \theta + (u \times a) \sin \theta + (u \otimes u) a (1 - \cos \theta)$



- 关于向量  $v$  作对称变换的矩阵  $H = I - 2\frac{vv^T}{v^T v}$ ,
- 点  $a$  对称点:  $a' = a - 2\frac{v^T a}{v^T v} \cdot v$

1.20 立体角

对于任意一个四面体  $OABC$ , 从  $O$  点观察  $\triangle ABC$  的立体角  $\tan \frac{\Omega}{2} = \frac{\text{mix}(\vec{a}, \vec{b}, \vec{c})}{|a||b||c|+(\vec{a} \cdot \vec{b})|c|+(\vec{a} \cdot \vec{c})|b|+(\vec{b} \cdot \vec{c})|a|}$ .

2 数据结构

2.1 动态凸包 (只支持插入)

```
1 #define x first // `upperHull $\leftarrow (x, y)$`
2 #define y second // `lowerHull $\leftarrow (x, -y)$`
3 typedef map<int, int> mii;
4 typedef map<int, int>::iterator mit;
5 struct point { point(const mit &p): x(p->first), y(p->second) {} };
6 inline bool checkInside(mii &a, const point &p) { // `border inclusive`
7     int x = p.x, y = p.y; mit p1 = a.lower_bound(x);
8     if (p1 == a.end()) return false; if (p1->x == x) return y <= p1->y;
9     if (p1 == a.begin()) return false; mit p2(p1--);
10    return sign(det(p - point(p1), point(p2) - p)) >= 0;
11 } inline void addPoint(mii &a, const point &p) { // `no collinear points`
12     int x = p.x, y = p.y; mit pnt = a.insert(make_pair(x, y)).first, p1, p2;
13     for (pnt->y = y; ; a.erase(p2)) {
14         p1 = pnt; if (++p1 == a.end()) break;
15         p2 = p1; if (++p1 == a.end()) break;
16         if (det(point(p2) - p, point(p1) - p) < 0) break;
17     } for ( ; ; a.erase(p2)) {
18         if ((p1 = pnt) == a.begin()) break;
19         if (--p1 == a.begin()) break; p2 = p1--;
20         if (det(point(p2) - p, point(p1) - p) > 0) break;
21     }
22 }
```

2.2 Rope 用法

```
1 #include <ext/rope>
2 using __gnu_cxx::crope; using __gnu_cxx::rope;
3 a = b.substr(from, len); // `${from, from + len}$`
4 a = b.substr(from); // `${from, from}$`
5 b.c_str(); // `might lead to memory leaks`
6 b.delete_c_str(); // `delete the c\{}str that created before`
7 a.insert(p, str); // `insert str before position $p$`
8 a.erase(i, n); // `erase ${i, i + n}$`
```

2.3 Treap

```
1 struct node { int key, prio, size; node *ch[2]; } base[MAXN], *top, *root, *null, nil;
2 typedef node *tree;
3 tree newNode(int key) {
4     static int seed = 3312;
5     top->key = key; top->prio = seed = int(seed * 48271LL % 2147483647);
6     top->size = 1; top->ch[0] = top->ch[1] = null; return top++;
7 }
8 void Rotate(tree &x, int d) {
9     tree y = x->ch[!d]; x->ch[!d] = y->ch[d]; y->ch[d] = x; y->size = x->size;
10    x->size = x->ch[0]->size + 1 + x->ch[1]->size; x = y;
11 }
12 void Insert(tree &t, int key) {
13     if (t == null) t = newNode(key);
14     else { int d = t->key < key; Insert(t->ch[d], key); ++t->size;
```

```
15         if (t->ch[d]->prio < t->prio) Rotate(t, !d);
16     }
17 }
18 void Delete(tree &t, int key) {
19     if (t->key != key) { Delete(t->ch[t->key < key], key); --t->size; }
20     else if (t->ch[0] == null) t = t->ch[1];
21     else if (t->ch[1] == null) t = t->ch[0];
22     else { int d = t->ch[0]->prio < t->ch[1]->prio;
23         Rotate(t, d); Delete(t->ch[d], key); --t->size;
24     }
25 }
```

2.4 可持久化 Treap

```
1 inline bool randomByKey(int a, int b) {
2     static long long seed = 1;
3     return (seed = seed * 48271 % 2147483647) * (a + b) < 2147483647LL * a;
4 }
5 tree merge(tree x, tree y) {
6     if (x == null) return y; if (y == null) return x;
7     tree t = NULL;
8     if (randomByKey(x->size, y->size)) t = newNode(x), t->r = merge(x->r, y);
9     else t = newNode(y), t->l = merge(x, y->l);
10    update(t); return t;
11 }
12 void splitByKey(tree t, int k, tree &l, tree &r) { // `${-\infty, k} [k, +\infty)$`
13     if (t == null) l = r = null;
14     else if (t->key < k) l = newNode(t), splitByKey(t->r, k, l->r, r), update(l);
15     else r = newNode(t), splitByKey(t->l, k, l, r->l), update(r);
16 }
17 void splitBySize(tree t, int k, tree &l, tree &r) { // `${1, k} [k, +\infty)$`
18     static int s; if (t == null) l = r = null;
19     else if ((s = t->l->size + 1) < k) l = newNode(t), splitBySize(t->r, k - s, l->r, r), update(l);
20     else r = newNode(t), splitBySize(t->l, k, l, r->l), update(r);
21 }
```

2.5 左偏树

```
1 tree merge(tree a, tree b) {
2     if (a == null) return b;
3     if (b == null) return a;
4     if (a->key > b->key) swap(a, b);
5     a->r = merge(a->r, b);
6     a->r->fa = a;
7     if (a->lc->dist < a->r->dist) swap(a->lc, a->r);
8     a->dist = a->r->dist + 1;
9     return a;
10 }
11 void erase(tree t) {
12     tree x = t->fa, y = merge(t->lc, t->rc);
13     if (y != null) y->fa = x;
14     if (x == null) root = y;
15     else
16         for ((x->lc == t ? x->lc : x->rc) = y; x != null; y = x, x = x->fa) {
17             if (x->lc->dist < x->rc->dist) swap(x->lc, x->rc);
18             if (x->rc->dist + 1 == x->dist) return;
19             x->dist = x->rc->dist + 1;
20         }
21 }
```

2.6 Link-Cut Tree

```
1 struct node { int rev; node *pre, *ch[2]; } base[MAXN], nil, *null;
2 typedef node *tree;
3 #define isRoot(x) (x->pre->ch[0] != x && x->pre->ch[1] != x)
4 #define isRight(x) (x->pre->ch[1] == x)
5 inline void MakeRev(tree t) { if (t != null) { t->rev ^= 1; swap(t->ch[0], t->ch[1]); } }
```

```

6 inline void PushDown(tree t) { if (t->rev) { MakeRev(t->ch[0]); MakeRev(t->ch[1]); t->rev = 0; } }
7 inline void Rotate(tree x) {
8     tree y = x->pre; PushDown(y); PushDown(x);
9     int d = isRight(x);
10    if (!isRoot(y)) y->pre->ch[isRight(y)] = x; x->pre = y->pre;
11    if ((y->ch[d] = x->ch[d]) != null) y->ch[d]->pre = y;
12    x->ch[!d] = y; y->pre = x; Update(y);
13 }
14 inline void Splay(tree x) {
15     PushDown(x); for (tree y; !isRoot(x); Rotate(x)) {
16         y = x->pre; if (!isRoot(y)) Rotate(isRight(x) != isRight(y) ? x : y);
17     } Update(x);
18 }
19 inline void Splay(tree x, tree to) {
20     PushDown(x); for (tree y; (y = x->pre) != to; Rotate(x)) if (y->pre != to)
21         Rotate(isRight(x) != isRight(y) ? x : y);
22     Update(x);
23 }
24 inline tree Access(tree t) {
25     tree last = null; for (; t != null; last = t, t = t->pre) Splay(t), t->ch[1] = last, Update(t);
26     return last;
27 }
28 inline void MakeRoot(tree t) { Access(t); Splay(t); MakeRev(t); }
29 inline tree FindRoot(tree t) { Access(t); Splay(t); tree last = null;
30     for (; t != null; last = t, t = t->ch[0]) PushDown(t); Splay(last); return last;
31 }
32 inline void Join(tree x, tree y) { MakeRoot(y); y->pre = x; }
33 inline void Cut(tree t) { Access(t); Splay(t); t->ch[0]->pre = null; t->ch[0] = null; Update(t); }
34 inline void Cut(tree x, tree y) {
35     tree upper = (Access(x), Access(y));
36     if (upper == x) { Splay(x); y->pre = null; x->ch[1] = null; Update(x); }
37     else if (upper == y) { Access(x); Splay(y); x->pre = null; y->ch[1] = null; Update(y); }
38     else assert(0); // 'impossible to happen'
39 }
40 inline int Query(tree a, tree b) { // 'query the cost in path a <=> b, lca inclusive'
41     Access(a); tree c = Access(b); // c is lca
42     int v1 = c->ch[1]->maxCost; Access(a);
43     int v2 = c->ch[1]->maxCost;
44     return max(max(v1, v2), c->cost);
45 }
46 void Init() {
47     null = &nil; null->ch[0] = null->ch[1] = null->pre = null; null->rev = 0;
48     Rep(i, 1, N) { node &n = base[i]; n.rev = 0; n.pre = n.ch[0] = n.ch[1] = null; }
49 }

```

## 3 字符串相关

### 3.1 Manacher

```

1 char t[1001]; // 要处理的字符串
2 char s[1001 * 2]; // 中间插入特殊字符以后的
3 int Z[1001 * 2], L, R; // Gusfield's Algorithm
4 // 由a往左, 由b往右, 对称地做字符匹配
5 int match(int a, int b) {
6     int i = 0;
7     while (a - i >= 0 && b + i < n && s[a - i] == s[b + i]) i++;
8     return i;
9 }
10 void longest_palindromic_substring()
11 {
12     int N = strlen(t);
13     /* 在 t 中插入特殊字符, 存放到 s */ memset(s, '.', N*2+1);
14     for (int i = 0; i < N; ++i) s[i * 2 + 1] = t[i];
15     N = N * 2 + 1;
16     Z[0] = 1; L = R = 0;
17     for (int i = 1; i < N; ++i) {
18         int ii = L - (i - L), n = R + 1 - i; // i的映射位
19         if (i > R) Z[i] = match(i, i), L = i, R = i + Z[i] - 1;
20         else if (Z[ii] == n) Z[i] = n + match(i - n, i + n), L = i, R = i + Z[i] - 1;
21         else Z[i] = min(Z[ii], n);

```

```

22     }
23     // 寻找最长回文子串的長度。
24     int n = 0, p = 0;
25     for (int i = 0; i < N; ++i) if (Z[i] > n) n = Z[p = i];
26     // 記得去掉特殊字元。
27     cout << "最长回文子串的長度是" << (n - 1) / 2;
28     // 印出最长回文子串, 記得別印特殊字元。
29     for (int i = p - Z[p] + 1; i <= p + Z[p] - 1; ++i) if (i & 1) cout << s[i];
30 }

```

### 3.2 KMP

$next[i] = \max\{len | A[0 \dots len - 1] = A \text{ 的第 } i \text{ 位向前或后的长度为 } len \text{ 的串}\}$   
 $ext[i] = \max\{len | A[0 \dots len - 1] = B \text{ 的第 } i \text{ 位向前或后的长度为 } len \text{ 的串}\}$

```

1 void KMP(char *a, int la, char *b, int lb, int *next, int *ext) {
2     —a; —b; —next; —ext;
3     for (int i = 2, j = next[1] = 0; i <= la; i++) {
4         while (j && a[j + 1] != a[i]) j = next[j]; if (a[j + 1] == a[i]) ++j; next[i] = j;
5     } for (int i = 1, j = 0; i <= lb; ++i) {
6         while (j && a[j + 1] != b[i]) j = next[j]; if (a[j + 1] == b[i]) ++j; ext[i] = j;
7         if (j == la) j = next[j];
8     }
9 } void ExKMP(char *a, int la, char *b, int lb, int *next, int *ext) {
10     next[0] = la; for (int &j = next[1] = 0; j + 1 < la && a[j] == a[j + 1]; ++j);
11     for (int i = 2, k = 1; i < la; ++i) {
12         int p = k + next[k], l = next[i - k]; if (l < p - i) next[i] = l;
13         else for (int &j = next[k = i] = max(0, p - i); i + j < la && a[j] == a[i + j]; ++j);
14     } for (int &j = ext[0] = 0; j < la && j < lb && a[j] == b[j]; ++j);
15     for (int i = 1, k = 0; i < lb; ++i) {
16         int p = k + ext[k], l = next[i - k]; if (l < p - i) ext[i] = l;
17         else for (int &j = ext[k = i] = max(0, p - i); j < la && i + j < lb && a[j] == b[i + j]; ++j);
18     }
19 }

```

### 3.3 Aho-Corasick 自动机

```

1 void construct() {
2     static tree Q[MAX_NODE]; int head = 0, tail = 0;
3     for (root->fail = root, Q[++tail] = root; head < tail; ) {
4         tree x = Q[++head];
5         // if (x->fail->danger) x->danger = true;
6         Rep(d, 0, sigma - 1) if (!x->next[d])
7             x->next[d] = (x == root) ? (root) : (x->fail->next[d]);
8         else {
9             x->next[d]->fail = (x == root) ? (root) : (x->fail->next[d]);
10            Q[++tail] = x->next[d];
11        }
12    }
13 }

```

### 3.4 后缀自动机

```

1 struct SAM {
2     int in[Maxn * 2 + 1][Sigma], fa[Maxn * 2 + 1], max[Maxn * 2 + 1], tot, last;
3     void init(int n) {
4         tot = last = 0;
5         for (int i = 0; i <= 2 * n + 1; ++i)
6             memset(in[i], -1, sizeof in[i]), fa[i] = -1;
7     }
8     void add(int x) {
9         int v = last; ++tot, last = tot, max[last] = max[v] + 1;
10        while (v != -1 && in[v][x] == -1) in[v][x] = last, v = fa[v];
11        if (v == -1) { fa[last] = 0; return; }
12        int p = in[v][x];
13        if (max[p] == max[v] + 1) fa[last] = p;
14        else {

```

```

15     int np = ++tot;
16     max[np] = max[v] + 1; fa[np] = fa[p], fa[p] = np, fa[last] = np;
17     while(v != -1 && in[v][x] == p) in[v][x] = np, v = fa[v];
18     memcpy(in[np], in[p], sizeof in[p]);
19 }

```

### 3.5 后缀数组

待排序的字符串放在  $r[0 \dots n-1]$  中, 最大值小于  $m$ .

$r[0 \dots n-2] > 0, r[n-1] = 0$ .

结果放在  $sa[0 \dots n-1]$ .

```

1 namespace SuffixArrayDoubling {
2     int wa[MAXN], wb[MAXN], wv[MAXN], ws[MAXN];
3     int cmp(int *r, int a, int b, int l) { return r[a] == r[b] && r[a+l] == r[b+l]; }
4     void da(int *r, int *sa, int n, int m) {
5         int i, j, p, *x = wa, *y = wb, *t;
6         for (i = 0; i < m; i++) ws[i] = 0;
7         for (i = 0; i < n; i++) ws[x[i]] = r[i]++;
8         for (i = 1; i < m; i++) ws[i] += ws[i-1];
9         for (i = n-1; i >= 0; i--) sa[ws[x[i]]] = i;
10        for (j = 1, p = 1; p < n; j *= 2, m = p) {
11            for (p = 0, i = n-j; i < n; i++) y[p++] = i;
12            for (i = 0; i < n; i++) if (sa[i] >= j) y[p++] = sa[i] - j;
13            for (i = 0; i < n; i++) wv[i] = x[y[i]];
14            for (i = 0; i < m; i++) ws[i] = 0;
15            for (i = 0; i < n; i++) ws[wv[i]]++;
16            for (i = 1; i < m; i++) ws[i] += ws[i-1];
17            for (i = n-1; i >= 0; i--) sa[ws[wv[i]]] = y[i];
18            for (t = x, x = y, y = t, p = 1, x[sa[0]] = 0, i = 1; i < n; i++)
19                x[sa[i]] = cmp(y, sa[i-1], sa[i], j) ? p-1 : p++;
20        }
21    }
22    namespace SuffixArrayDC3 { // `r 与 sa 大小需 3 倍`
23        #define F(x) ((x) / 3 + ((x) % 3 == 1 ? 0 : tb))
24        #define G(x) ((x) < tb ? (x) * 3 + 1 : ((x) - tb) * 3 + 2)
25        int wa[MAXN], wb[MAXN], wv[MAXN], ws[MAXN];
26        int c0(int *r, int a, int b) {
27            return r[a] == r[b] && r[a+1] == r[b+1] && r[a+2] == r[b+2];
28        }
29        int c12(int k, int *r, int a, int b) {
30            if (k == 2) return r[a] < r[b] || (r[a] == r[b] && c12(1, r, a+1, b+1));
31            else return r[a] < r[b] || (r[a] == r[b] && wv[a+1] < wv[b+1]);
32        }
33        void sort(int *r, int *a, int *b, int n, int m) {
34            for (int i = 0; i < n; i++) wv[i] = r[a[i]];
35            for (int i = 0; i < m; i++) ws[i] = 0;
36            for (int i = 0; i < n; i++) ws[wv[i]]++;
37            for (int i = 1; i < m; i++) ws[i] += ws[i-1];
38            for (int i = n-1; i >= 0; i--) b[ws[wv[i]]] = a[i];
39        }
40        void dc3(int *r, int *sa, int n, int m) {
41            int i, j, *rn = r + n, *san = sa + n, ta = 0, tb = (n+1) / 3, tbc = 0, p;
42            r[n] = r[n+1] = 0;
43            for (i = 0; i < n; i++) if (i % 3 != 0) wa[tbc++] = i;
44            sort(r+2, wa, wb, tbc, m);
45            sort(r+1, wb, wa, tbc, m);
46            sort(r, wa, wb, tbc, m);
47            for (p = 1, rn[F(wb[0])] = 0, i = 1; i < tbc; i++)
48                rn[F(wb[i])] = c0(r, wb[i-1], wb[i]) ? p-1 : p++;
49            if (p < tbc) dc3(rn, san, tbc, p);
50            else for (i = 0; i < tbc; i++) san[rn[i]] = i;
51            for (i = 0; i < tbc; i++) if (san[i] < tb) wb[ta++] = san[i] * 3;
52            if (n % 3 == 1) wb[ta++] = n-1;
53            sort(r, wb, wa, ta, m);
54            for (i = 0; i < tbc; i++) wv[wb[i]] = G(san[i]);
55            for (i = 0, j = 0, p = 0; i < ta && j < tbc; p++)
56                sa[p] = c12(wb[j] % 3, r, wa[i], wb[j]) ? wa[i++] : wb[j++];
57            for (; i < ta; p++) sa[p] = wa[i++];
58            for (; j < tbc; p++) sa[p] = wb[j++];
59        }
60    }

```

```

61 }
62 namespace CalcHeight {
63     int rank[MAXN], height[MAXN];
64     void calheight(int *r, int *sa, int n) {
65         int i, j, k = 0; for (i = 1; i <= n; i++) rank[sa[i]] = i;
66         for (i = 0; i < n; height[rank[i++]] = k)
67             for (k ? k-- : 0, j = sa[rank[i]-1]; r[i+k] == r[j+k]; k++);
68     }
69     void init(int len)
70     {
71         for(int i = 0; i <= len + 10; ++i)
72             rank[i] = height[i] = 0;
73     }
74 }
75 //Sample
76 int r[MAXN]; char s[MAXN];
77 int main()
78 {
79     int len;
80     scanf("%s", s);
81     len = strlen(s);
82     for(int i = 0; i < len; ++i) r[i] = s[i] - 'a' + 1;
83     r[len] = 0;
84     SuffixArrayDoubling::da(r, sa, n+1, 30);
85     CalcHeight::calheight(r, sa, n);
86     //Then the value of sa[0~len-1] is 1 ~ n, so init RMQ carefully(1-n not 0-n-1)
87     return 0;
88 }
89 }

```

### 3.6 环串最小表示

```

1 int minimalRepresentation(int N, char *s) { // s must be double-sized and 0-based
2     int i, j, k, l; for (i = 0; i < N; ++i) s[i+N] = s[i]; s[N+N] = 0;
3     for (i = 0, j = 1; j < N; ) {
4         for (k = 0; k < N && s[i+k] == s[j+k]; ++k);
5         if (k >= N) break; if (s[i+k] < s[j+k]) j += k+1;
6         else l = i+k, i = j, j = max(l, j)+1;
7     } return i; // [i, i+N) is the minimal representation
8 }

```

### 3.7 回文自动机

```

1 #include <cstdlib>
2 #include <cstdio>
3 #include <cstring>
4 #include <algorithm>
5
6 const int C = 26, N = 100000, S = N + 2 + C;
7 char string[N+2];
8 int s, length[S], suffix[S], go[S][C];
9 int extend(int p, int i) {
10     while (string[i-1-length[p]] != string[i]) {p = suffix[p];}
11     int q = suffix[p];
12     while (string[i-1-length[q]] != string[i]) {q = suffix[q];}
13     int c = string[i] - 'a';
14     int pp = go[p][c]; int qq = go[q][c];
15     if (pp == -1) {
16         length[pp = go[p][c] = s++] = length[p] + 2;
17         suffix[pp] = qq;
18         memset(go[pp], -1, sizeof(go[pp]));
19     } return pp;
20 }
21 int main() {
22     int tests; scanf("%d", &tests);
23     for (int t = 1; t <= tests; ++t) {
24         printf("Case %d: ", t);
25         for (int i = 0; i < C+2; ++i) {
26             suffix[i] = 1; length[i] = std::min(i-1, 1);
27             memset(go[i], -1, sizeof(go[i]));
28         }

```

```

29     suffix[0] = suffix[1] = 0;
30     for (int i = 0; i < C; ++i)
31         go[0][i] = 2 + i;
32     s = C + 2; string[0] = '#';
33     scanf("%s", string + 1);
34     int n = strlen(string + 1); int p = 0;
35     for (int i = 1; i <= n; ++i)
36         p = extend(p, i);
37     int result = s - (C + 2); std::sort(string + 1, string + n + 1);
38     result += std::unique(string + 1, string + n + 1) - string - 1;
39     printf("%d\n", result);
40 }
41 return 0;
42 }

```

## 4 图论

### 4.1 Dominator Tree

```

1  #include <cstdio>
2  #include <cstdlib>
3  #include <cstring>
4  #include <iostream>
5  #include <algorithm>
6  #include <vector>
7
8  using namespace std;
9
10 const int oo = 1073741819;
11
12 const int Maxn = 200000;
13 const int Maxm = 200000;
14
15 vector<int> g[Maxn];
16
17 //idom[i] is the dominator of i, node id — 1 based(1 ~ n), n is the source
18 class DominatorTree
19 {
20 public:
21     int tail[4][Maxm], n, m;
22     int Next[4][Maxm], sora[4][Maxm];
23     int ss[4], top, w_time;
24     int rel[Maxn], semi[Maxn], b[Maxn], idom[Maxn], best[Maxn], st[Maxn], pre[Maxn];
25     void origin()
26     {
27         for (int e = 0; e <= 3; e++) ss[e] = n;
28         for (int i = 1; i <= n; i++) {
29             for (int e = 0; e <= 3; e++)
30                 tail[e][i] = i, Next[e][i] = 0;
31             rel[i] = 0, semi[i] = idom[i] = pre[i] = 0, best[i] = i;
32             b[i] = i;
33         }
34         rel[0] = oo;
35     }
36     void link(int e, int x, int y)
37     {
38         ++ss[e], Next[e][tail[e][x]] = ss[e], tail[e][x] = ss[e], sora[e][ss[e]] = y, Next[e][ss[e]] = 0;
39     }
40     void dfs(int x, int y)
41     {
42         ++w_time, rel[x] = w_time;
43         st[++top] = x, pre[x] = y;
44         for (int i = x, ne; Next[0][i];) {
45             i = Next[0][i], ne = sora[0][i];
46             if (!rel[ne]) dfs(ne, x);
47         }
48     }
49     int find(int x)
50     {
51         int y = b[x];
52         if (b[x] != x) b[x] = find(b[x]);
53         if (rel[semi[best[y]]] < rel[semi[best[x]]])
54             best[x] = best[y];
55         return b[x];
56     }

```

```

57 //n — number of vertex, m — number of edges, e — edge set
58 void init(int _n, int _m, const vector<pair<int, int> > &e)
59 {
60     n = _n, m = _m;
61     origin();
62     for (int i = 0; i < m; i++) {
63         link(0, e[i].first, e[i].second);
64         link(1, e[i].second, e[i].first);
65     }
66     w_time = 0, top = 0;
67     dfs(n, 0);
68 }
69
70 void work()
71 {
72     for (int i = top; i >= 1; i--) {
73         int ne = st[i];
74         for (int j = ne, na; Next[1][j];) {
75             j = Next[1][j], na = sora[1][j];
76             if (!rel[na]) continue;
77             int y;
78             if (rel[na] > rel[ne]) {
79                 find(na);
80                 y = semi[best[na]];
81             }
82             else y = na;
83             if (rel[y] < rel[semi[ne]]) semi[ne] = y;
84         }
85         if (ne != n) link(2, semi[ne], ne);
86         for (int j = ne, na; Next[2][j];) {
87             j = Next[2][j], na = sora[2][j];
88             find(na);
89             int y = best[na];
90             if (semi[y] == semi[na]) idom[na] = semi[na];
91             else idom[na] = y;
92         }
93         for (int j = ne, na; Next[0][j];) {
94             j = Next[0][j], na = sora[0][j];
95             if (pre[na] == ne) {
96                 na = find(na);
97                 b[na] = ne;
98             }
99         }
100     }
101     for (int i = 2; i <= top; i++) {
102         int ne = st[i];
103         if (idom[ne] != semi[ne]) idom[ne] = idom[idom[ne]];
104         link(3, idom[ne], ne);
105     }
106 }
107 }dom;

```

### 4.2 带花树

```

1 namespace Blossom {
2     int n, head, tail, S, T, lca;
3     int match[MAXN], Q[MAXN], pred[MAXN], label[MAXN], inq[MAXN], inb[MAXN];
4     vector<int> link[MAXN];
5     inline void push(int x) { Q[tail++] = x; inq[x] = true; }
6     int findCommonAncestor(int x, int y) {
7         static bool inPath[MAXN]; for (int i = 0; i < n; ++i) inPath[i] = 0;
8         for (; x = pred[ match[x] ]) { x = label[x]; inPath[x] = true; if (x == S) break; }
9         for (; y = pred[ match[y] ]) { y = label[y]; if (inPath[y]) break; } return y;
10    }
11    void resetTrace(int x, int lca) {
12        while (label[x] != lca) { int y = match[x]; inb[ label[x] ] = inb[ label[y] ] = true;
13            x = pred[y]; if (label[x] != lca) pred[x] = y; }
14    void blossomContract(int x, int y) {
15        lca = findCommonAncestor(x, y);
16        Foru(i, 0, n) inb[i] = 0; resetTrace(x, lca); resetTrace(y, lca);
17        if (label[x] != lca) pred[x] = y; if (label[y] != lca) pred[y] = x;
18        Foru(i, 0, n) if (inb[ label[i] ]) { label[i] = lca; if (!inq[i]) push(i); }
19    }
20    bool findAugmentingPath() {

```

```

21 Foru(i, 0, n) pred[i] = -1, label[i] = i, inq[i] = 0;
22 int x, y, z; head = tail = 0;
23 for (push(S); head < tail; ) for (int i = (int)link[x = Q[head++]].size() - 1; i >= 0; --i) {
24     y = link[x][i]; if (label[x] == label[y] || x == match[y]) continue;
25     if (y == S || (match[y] >= 0 && pred[ match[y] ] >= 0)) blossomContract(x, y);
26     else if (pred[y] == -1) {
27         pred[y] = x; if (match[y] >= 0) push(match[y]);
28     } else {
29         for (x = y; x >= 0; x = z) {
30             y = pred[x], z = match[y]; match[x] = y, match[y] = x;
31         } return true; }}} return false;
32 }
33 int findMaxMatching() {
34     int ans = 0; Foru(i, 0, n) match[i] = -1;
35     for (S = 0; S < n; ++S) if (match[S] == -1) if (findAugmentingPath()) ++ans;
36     return ans;
37 }
38 }

```

### 4.3 KM

```

1 int N, Tcnt, w[MAXN][MAXN], slack[MAXN];
2 int lx[MAXN], linkx[MAXN], visy[MAXN], ly[MAXN], linky[MAXN], visx[MAXN]; // 初值全为0
3 bool DFS(int x) { visx[x] = Tcnt;
4     Rep(y, 1, N) if(visy[y] != Tcnt) { int t = lx[x] + ly[y] - w[x][y];
5         if (t == 0) { visy[y] = Tcnt;
6             if (!linky[y] || DFS(linky[y])) { linkx[x] = y; linky[y] = x; return true; }
7         } else cMin(slack[y], t);
8     } return false;
9 } void KM() {
10     Tcnt = 0; Rep(x, 1, N) Rep(y, 1, N) cMax(lx[x], w[x][y]);
11     Rep(S, 1, N) { Rep(i, 1, N) slack[i] = INF;
12         for (++Tcnt; !DFS(S); ++Tcnt) { int d = INF;
13             Rep(y, 1, N) if(visy[y] != Tcnt) cMin(d, slack[y]);
14             Rep(x, 1, N) if(visx[x] == Tcnt) lx[x] -= d;
15             Rep(y, 1, N) if(visy[y] == Tcnt) ly[y] += d; else slack[y] -= d;
16         }
17     }
18 }

```

### 4.4 双连通分量

```

1 #include <iostream>
2 #include <cstdio>
3 #include <cstring>
4 #include <cstdlib>
5 #include <vector>
6
7 using namespace std;
8
9 const int MAXN = 100000 + 10;
10
11 int dfn[MAXN], low[MAXN], bccno[MAXN], dfn_clock, bcc_cnt, Top;
12 vector<int> G[MAXN], bcc[MAXN];
13 pair<int, int> stk[MAXN];
14 bool iscut[MAXN];
15 int n, m;
16
17 void dfs(int p, int fa) {
18     low[p] = dfn[p] = ++dfn_clock;
19     int child = 0;
20     for (int i = 0; i < G[p].size(); ++i) {
21         int v = G[p][i];
22         if (!dfn[v]) {
23             stk[++Top] = make_pair(p, v);
24             dfs(v, p);
25             child++;
26             low[p] = min(low[p], low[v]);
27             if (low[v] >= dfn[p]) {
28                 iscut[p] = 1;
29                 ++bcc_cnt;
30                 bcc[bcc_cnt].clear();

```

```

31         for (;;) {
32             pair<int, int> x = stk[Top];
33             --Top;
34             if (bccno[x.first] != bcc_cnt) {
35                 bccno[x.first] = bcc_cnt;
36                 bcc[bcc_cnt].push_back(x.first);
37             }
38             if (bccno[x.second] != bcc_cnt) {
39                 bccno[x.second] = bcc_cnt;
40                 bcc[bcc_cnt].push_back(x.second);
41             }
42             if (x.first == p && x.second == v)
43                 break;
44         }
45     }
46 } else
47     if (dfn[v] < dfn[p] && v != fa) {
48         stk[++Top] = make_pair(p, v);
49         low[p] = min(low[p], dfn[v]);
50     }
51 }
52 if (fa < 0 && child == 1) iscut[p] = 0;
53 }
54
55 void find_bcc(int n) {
56     for (int i = 1; i <= n; ++i) dfn[i] = 0;
57     for (int i = 1; i <= n; ++i) iscut[i] = 0;
58     for (int i = 1; i <= n; ++i) bccno[i] = 0;
59     dfn_clock = bcc_cnt = 0;
60     for (int i = 1; i <= n; ++i)
61         if (!dfn[i])
62             dfs(i, -1);
63 }
64
65 int main() {
66     scanf("%d%d", &n, &m);
67     for (int a, b, i = 1; i <= m; ++i) {
68         scanf("%d%d", &a, &b);
69         G[a].push_back(b);
70         G[b].push_back(a);
71     }
72     find_bcc(n);
73     return 0;
74 }
75
76 }

```

### 4.5 强连通分量

```

1 #include <iostream>
2 #include <cstdio>
3 #include <cstring>
4 #include <cstdlib>
5 #include <vector>
6 #include <algorithm>
7
8 using namespace std;
9
10 const int MAXN = 100000 + 10;
11
12 vector<int> G[MAXN];
13 int n, m;
14 int dfn[MAXN], low[MAXN], stk[MAXN], Top, scc_cnt, sccno[MAXN], dfn_clock;
15
16 void dfs(int p) {
17     dfn[p] = low[p] = ++dfn_clock;
18     stk[++Top] = p;
19     for (int i = 0; i < (int)G[p].size(); ++i) {
20         int v = G[p][i];
21         if (!dfn[v]) {
22             dfs(v);
23             low[p] = min(low[p], low[v]);
24         }
25         else if (!sccno[v])
26             low[p] = min(low[p], dfn[v]);
27     }
28     if (low[p] == dfn[p]) {

```

```

29     scc_cnt++;
30     for (;;) {
31         int x = stk[Top];
32         --Top;
33         sccno[x] = scc_cnt;
34         if (x == p) break;
35     }
36 }
37 }
38 void find_scc(int n) {
39     dfn_clock = scc_cnt = 0;
40     for (int i = 1; i <= n; ++i) sccno[i] = 0;
41     for (int i = 1; i <= n; ++i) dfn[i] = low[i] = 0;
42     for (int i = 1; i <= n; ++i)
43         if (!dfn[i])
44             dfs(i);
45 }
46 }

```

## 4.6 2-SAT 与 Kosaraju

注意 Kosaraju 需要建反图

```

1 namespace SCC {
2     int code[MAXN * 2], seq[MAXN * 2], sCnt;
3     void DFS_1(int x) { code[x] = 1;
4         for (edge e(fir[x]); e; e = e->next) if (code[e->to] == -1) DFS_1(e->to);
5         seq[++sCnt] = x;
6     } void DFS_2(int x) { code[x] = sCnt;
7         for (edge e(fir2[x]); e; e = e->next) if (code[e->to] == -1) DFS_2(e->to); }
8     void SCC(int N) {
9         sCnt = 0; for (int i = 1; i <= N; ++i) code[i] = -1;
10        for (int i = 1; i <= N; ++i) if (code[i] == -1) DFS_1(i);
11        sCnt = 0; for (int i = 1; i <= N; ++i) code[i] = -1;
12        for (int i = N; i >= 1; --i) if (code[seq[i]] == -1) {
13            ++sCnt; DFS_2(seq[i]); }
14    }
15    // true   - 2i - 1
16    // false  - 2i
17    bool TwoSat() { SCC::SCC(N + N);
18        // if code[2i - 1] == code[2i]: no solution
19        // if code[2i - 1] > code[2i]: i selected. else i not selected
20    }

```

## 4.7 全局最小割 Stoer-Wagner

```

1 int minCut(int N, int G[MAXN][MAXN]) { // 0-based
2     static int weight[MAXN], used[MAXN]; int ans = INT_MAX;
3     while (N > 1) {
4         for (int i = 0; i < N; ++i) used[i] = false; used[0] = true;
5         for (int i = 0; i < N; ++i) weight[i] = G[i][0];
6         int S = -1, T = 0;
7         for (int _r = 2; _r <= N; ++_r) { // N - 1 selections
8             int x = -1;
9             for (int i = 0; i < N; ++i) if (!used[i])
10                if (x == -1 || weight[i] > weight[x]) x = i;
11            for (int i = 0; i < N; ++i) weight[i] += G[x][i];
12            S = T; T = x; used[x] = true;
13        } ans = min(ans, weight[T]);
14        for (int i = 0; i < N; ++i) G[i][S] += G[i][T], G[S][i] += G[i][T];
15        G[S][S] = 0; --N;
16        for (int i = 0; i <= N; ++i) swap(G[i][T], G[i][N]);
17        for (int i = 0; i < N; ++i) swap(G[T][i], G[N][i]);
18    } return ans;
19 }

```

## 4.8 重口味费用流

```

1 struct MinCostFlow {
2     int e[M], succ[M], last[N], val[M], cost[M], sum, dis[N], visit[N], slack[N];
3     int source, target, totFlow, totCost;
4     void init(int n) {
5         for (int i = 1; i <= n; ++i) last[i] = 0;
6         sum = 1;
7     }
8     void add(int a, int b, int c, int d) {
9         e[++sum] = b, succ[sum] = last[a], last[a] = sum;
10        e[++sum] = a, succ[sum] = last[b], last[b] = sum;
11        val[sum - 1] = c, val[sum] = 0;
12        cost[sum - 1] = d, cost[sum] = -d;
13    }
14    int modlable() {
15        int delta = INF;
16        for (int i = 1; i <= target; ++i) {
17            if (!visit[i] && slack[i] < delta) delta = slack[i];
18            slack[i] = INF;
19        }
20        if (delta == INF) return 1;
21        for (int i = 1; i <= target; ++i) if (visit[i])
22            dis[i] += delta;
23        return 0;
24    }
25    int dfs(int x, int flow) {
26        if (x == target) {
27            totFlow += flow;
28            totCost += flow * (dis[source] - dis[target]);
29            return flow;
30        }
31        visit[x] = 1; int left = flow;
32        for (int i = last[x]; i; i = succ[i]) {
33            if (val[i] > 0 && !visit[e[i]]) {
34                int y = e[i];
35                if (dis[y] + cost[i] == dis[x]) {
36                    int delta = dfs(y, min(left, val[i]));
37                    val[i] -= delta, val[i ^ 1] += delta;
38                    left -= delta;
39                    if (!left) {
40                        visit[x] = 0;
41                        return flow;
42                    }
43                } else slack[y] = min(slack[y], dis[y] + cost[i] - dis[x]);
44            }
45        }
46        return flow - left;
47    }
48    pair<int, int> minCost() {
49        totFlow = 0, totCost = 0;
50        fill(dis + 1, dis + target + 1, 0);
51        do {
52            do {
53                fill(visit + 1, visit + target + 1, 0);
54            } while (dfs(source, INF));
55        } while (!modlable());
56        return make_pair(totFlow, totCost);
57    }
58 } mcf;
59 }

```

## 4.9 欧拉路

```

1 vector<int> eulerianWalk(int N, int S) {
2     static int res[MAXN], stack[MAXN]; static edge cur[MAXN];
3     int rcnt = 0, top = 0, x; for (int i = 1; i <= N; ++i) cur[i] = fir[i];
4     for (stack[top++] = S; top; ) {
5         for (x = stack[--top]; ; ) {
6             edge &e = cur[x]; if (e == NULL) break;
7             stack[top++] = x; x = e->to; e = e->next;
8             // 对于无向图需要删掉反向边
9         } res[rcnt++] = x;
10    } reverse(res, res + rcnt); return vector<int>(res, res + rcnt);
11 }

```



## 4.10 稳定婚姻

```

1 namespace StableMatching {
2     int pairM[MAXN], pairW[MAXN], p[MAXN];
3     // init: pairM[0...n-1] = pairW[0...n-1] = -1, p[0...n-1] = 0
4     void stableMatching(int n, int orderM[MAXN][MAXN], int preferW[MAXN][MAXN]) {
5         for (int i = 0; i < n; i++) while (pairM[i] < 0) {
6             int w = orderM[i][p[i]++], m = pairW[w];
7             if (m == -1) pairM[i] = w, pairW[w] = i;
8             else if (preferW[w][i] < preferW[w][m])
9                 pairM[m] = -1, pairM[i] = w, pairW[w] = i, i = m;
10        }
11    }
12 }

```

## 4.11 最大团搜索

```

1 namespace MaxClique { // 1-based
2     int g[MAXN][MAXN], len[MAXN], list[MAXN][MAXN], mc[MAXN], ans, found;
3     void DFS(int size) {
4         if (len[size] == 0) { if (size > ans) ans = size, found = true; return; }
5         for (int k = 0; k < len[size] && !found; ++k) {
6             if (size + len[size] - k <= ans) break;
7             int i = list[size][k]; if (size + mc[i] <= ans) break;
8             for (int j = k + 1, len[size + 1] = 0; j < len[size]; ++j) if (g[i][list[size][j]])
9                 list[size + 1][len[size + 1]++] = list[size][j];
10            DFS(size + 1);
11        }
12    }
13    int work(int n) {
14        mc[n] = ans = 1; for (int i = n - 1; i >= 1; --i) { found = false; len[1] = 0;
15            for (int j = i + 1; j <= n; ++j) if (g[i][j]) list[1][len[1]++] = j;
16            DFS(1); mc[i] = ans;
17        } return ans;
18    }
19 }

```

## 4.12 极大团计数

```

1 namespace MaxCliqueCounting {
2     int n, ans;
3     int ne[MAXN], ce[MAXN];
4     int g[MAXN][MAXN], list[MAXN][MAXN];
5     void dfs(int size) {
6         int i, j, k, t, cnt, best = 0;
7         bool bb;
8         if (ne[size] == ce[size]) {
9             if (ce[size] == 0) ++ans;
10            return;
11        }
12        for (t = 0, i = 1; i <= ne[size]; ++i) {
13            for (cnt = 0, j = ne[size] + 1; j <= ce[size]; ++j)
14                if (!g[list[size][i]][list[size][j]]) ++cnt;
15            if (t == 0 || cnt < best) t = i, best = cnt;
16        }
17        if (t && best <= 0) return;
18        for (k = ne[size] + 1; k <= ce[size]; ++k) {
19            if (t > 0) {
20                for (i = k; i <= ce[size]; ++i)
21                    if (!g[list[size][t]][list[size][i]]) break;
22                swap(list[size][k], list[size][i]);
23            }
24            i = list[size][k];
25            ne[size + 1] = ce[size + 1] = 0;
26            for (j = 1; j < k; ++j)

```

```

32            if (g[i][list[size][j]])
33                list[size + 1][++ne[size + 1]] = list[size][j];
34            for (ce[size + 1] = ne[size + 1], j = k + 1; j <= ce[size]; ++j)
35                if (g[i][list[size][j]])
36                    list[size + 1][++ce[size + 1]] = list[size][j];
37            dfs(size + 1);
38            ++ne[size];
39            --best;
40            for (j = k + 1, cnt = 0; j <= ce[size]; ++j)
41                if (!g[i][list[size][j]]) ++cnt;
42            if (t == 0 || cnt < best) t = k, best = cnt;
43            if (t && best <= 0) break;
44        }
45    }
46    void work() {
47        int i;
48        ne[0] = 0;
49        ce[0] = 0;
50        for (i = 1; i <= n; ++i)
51            list[0][++ce[0]] = i;
52        ans = 0;
53        dfs(0);
54    }
55 }

```

## 4.13 最小树形图

```

1 namespace EdmondsAlgorithm { // O(ElogE + V^2) !!! 0-based !!!
2     struct enode { int from, c, key, delta, dep; enode *ch[2], *next;
3     } ebase[maxm], *etop, *fir[maxn], nil, *null, *inEdge[maxn], *chs[maxn];
4     typedef enode *edge; typedef enode *tree;
5     int n, m, setFa[maxn], deg[maxn], que[maxn];
6     inline void pushDown(tree x) { if (x->delta) {
7         x->ch[0]->key += x->delta; x->ch[0]->delta += x->delta;
8         x->ch[1]->key += x->delta; x->ch[1]->delta += x->delta; x->delta = 0;
9     } }
10    tree merge(tree x, tree y) {
11        if (x == null) return y; if (y == null) return x;
12        if (x->key > y->key) swap(x, y); pushDown(x); x->ch[1] = merge(x->ch[1], y);
13        if (x->ch[0]->dep < x->ch[1]->dep) swap(x->ch[0], x->ch[1]);
14        x->dep = x->ch[1]->dep + 1; return x;
15    }
16    void addEdge(int u, int v, int w) {
17        etop->from = u; etop->c = etop->key = w; etop->delta = etop->dep = 0;
18        etop->next = fir[v]; etop->ch[0] = etop->ch[1] = null;
19        fir[v] = etop; inEdge[v] = merge(inEdge[v], etop);
20    }
21    void deleteMin(tree &r) { pushDown(r); r = merge(r->ch[0], r->ch[1]); }
22    int findSet(int x) { return setFa[x] == x ? x : setFa[x] = findSet(setFa[x]); }
23    void clear(int V, int E) {
24        null = &nil; null->ch[0] = null->ch[1] = null; null->dep = -1;
25        n = V; m = E; etop = ebase; Foru(i, 0, V) fir[i] = NULL; Foru(i, 0, V) inEdge[i] = null;
26    }
27    int solve(int root) { int res = 0, head, tail;
28        for (int i = 0; i < n; ++i) setFa[i] = i;
29        for ( ; ; ) { memset(deg, 0, sizeof(int) * n); chs[root] = inEdge[root];
30            for (int i = 0; i < n; ++i) if (i != root && setFa[i] == i) {
31                while (findSet(inEdge[i]->from) == findSet(i)) deleteMin(inEdge[i]);
32                ++deg[ findSet(chs[i] = inEdge[i]->from) ];
33            }
34            for (int i = head = tail = 0; i < n; ++i)
35                if (i != root && setFa[i] == i && deg[i] == 0) que[tail++] = i;
36            while (head < tail) {
37                int x = findSet(chs[que[head++]]->from);
38                if (deg[x] == 0) que[tail++] = x;
39            } bool found = false;
40            for (int i = 0; i < n; ++i) if (i != root && setFa[i] == i && deg[i] > 0) {
41                int j = i; tree temp = null; found = true;
42                do {setFa[j] = findSet(chs[j]->from)} = i;

```

```

43         deleteMin(inEdge[j]); res += chs[j]→key;
44         inEdge[j]→key = chs[j]→key; inEdge[j]→delta = chs[j]→key;
45         temp = merge(temp, inEdge[j]);
46     } while (j != i); inEdge[i] = temp;
47 } if (!found) break;
48 } for (int i = 0; i < n; ++i) if (i != root && setFa[i] == i) res += chs[i]→key;
49 return res;
50 }
51 }
52 namespace ChuLiu { // O(V ^ 3) !!! 1-based !!!
53     int n, used[maxn], pass[maxn], eg[maxn], more, que[maxn], g[maxn][maxn];
54     void combine(int id, int &sum) { int tot = 0, from, i, j, k;
55         for (; id != 0 && !pass[id]; id = eg[id]) que[tot++] = id, pass[id] = 1;
56         for (from = 0; from < tot && que[from] != id; from++);
57         if (from == tot) return; more = 1;
58         for (i = from; i < tot; i++) {
59             sum += g[eg[que[i]]][que[i]]; if (i == from) continue;
60             for (j = used[que[i]] = 1; j <= n; j++) if (!used[j])
61                 if (g[que[i]][j] < g[id][j]) g[id][j] = g[que[i]][j];
62         }
63         for (i = 1; i <= n; i++) if (!used[i] && i != id)
64             for (j = from; j < tot; j++) {
65                 k = que[j]; if (g[i][id] > g[i][k] - g[eg[k]][k])
66                     g[i][id] = g[i][k] - g[eg[k]][k];
67             }
68     }
69     void clear(int V) { n = V; Rep(i, 1, V) Rep(j, 1, V) g[i][j] = inf; }
70     int solve(int root) {
71         int i, j, k, sum = 0; memset(used, 0, sizeof(int) * (n + 1));
72         for (more = 1; more; ) {
73             more = 0; memset(eg, 0, sizeof(int) * (n + 1));
74             for (i = 1; i <= n; i++) if (!used[i] && i != root) {
75                 for (j = 1, k = 0; j <= n; j++) if (!used[j] && i != j)
76                     if (k == 0 || g[j][i] < g[k][i]) k = j;
77                 eg[i] = k;
78             } memset(pass, 0, sizeof(int) * (n + 1));
79             for (i = 1; i <= n; i++) if (!used[i] && !pass[i] && i != root)
80                 combine(i, sum);
81         } for (i = 1; i <= n; i++) if (!used[i] && i != root) sum += g[eg[i]][i];
82         return sum;
83     }
84 }

```

#### 4.14 离线动态最小生成树

$O(Q \log^2 Q)$ .  $(qx[i], qy[i])$  表示将编号为  $qx[i]$  的边的权值改为  $qy[i]$ , 删除一条边相当于将其权值改为  $\infty$ , 加入一条边相当于将其权值从  $\infty$  变成某个值.

```

1  const int maxn = 100000 + 5;
2  const int maxm = 1000000 + 5;
3  const int maxq = 1000000 + 5;
4  const int qsize = maxm + 3 * maxq;
5  int n, m, Q, x[qsize], y[qsize], z[qsize], qx[maxq], qy[maxq], a[maxn], *tz;
6  int kx[maxn], ky[maxn], kt, vd[maxn], id[maxm], app[maxm];
7  bool extra[maxn];
8  void init() {
9      scanf("%d%d", &n, &m); for (int i = 0; i < m; i++) scanf("%d%d%d", x + i, y + i, z + i);
10     scanf("%d", &Q); for (int i = 0; i < Q; i++) { scanf("%d%d", qx + i, qy + i); qx[i] = -1; }
11 }
12 int find(int x) {
13     int root = x, next; while (a[root]) root = a[root];
14     while ((next = a[x]) != 0) a[x] = root, x = next; return root;
15 }
16 inline bool cmp(const int &a, const int &b) { return tz[a] < tz[b]; }
17 void solve(int *qx, int *qy, int Q, int n, int *x, int *y, int *z, int m, long long ans) {
18     int ri, rj;
19     if (Q == 1) {
20         for (int i = 1; i <= n; i++) a[i] = 0; z[qx[0]] = qy[0];
21         for (int i = 0; i < m; i++) id[i] = i;
22         tz = z; sort(id, id + m, cmp);
23         for (int i = 0; i < m; i++) {
24             ri = find(x[id[i]]); rj = find(y[id[i]]);
25             if (ri != rj) ans += z[id[i]], a[ri] = rj;

```

```

26         } printf("%I64d\n", ans);
27         return;
28     } int tm = kt = 0, n2 = 0, m2 = 0;
29     for (int i = 1; i <= n; i++) a[i] = 0;
30     for (int i = 0; i < Q; i++) {
31         ri = find(x[qx[i]]); rj = find(y[qx[i]]); if (ri != rj) a[ri] = rj;
32     }
33     for (int i = 0; i < m; i++) extra[i] = true;
34     for (int i = 0; i < Q; i++) extra[qx[i]] = false;
35     for (int i = 0; i < m; i++) if (extra[i]) id[tm++] = i;
36     tz = z; sort(id, id + tm, cmp);
37     for (int i = 0; i < tm; i++) {
38         ri = find(x[id[i]]); rj = find(y[id[i]]);
39         if (ri != rj)
40             a[ri] = rj, ans += z[id[i]], kx[kt] = x[id[i]], ky[kt] = y[id[i]], kt++;
41     }
42     for (int i = 1; i <= n; i++) a[i] = 0;
43     for (int i = 0; i < kt; i++) a[find(kx[i])] = find(ky[i]);
44     for (int i = 1; i <= n; i++) if (a[i] == 0) vd[i] = ++n2;
45     for (int i = 1; i <= n; i++) if (a[i] != 0) vd[i] = vd[find(i)];
46     int *Nx = x + m, *Ny = y + m, *Nz = z + m;
47     for (int i = 0; i < m; i++) app[i] = -1;
48     for (int i = 0; i < Q; i++)
49         if (app[qx[i]] == -1)
50             Nx[m2] = vd[x[qx[i]]], Ny[m2] = vd[y[qx[i]]], Nz[m2] = z[qx[i]], app[qx[i]] = m2, m2++;
51     for (int i = 0; i < Q; i++) {
52         z[qx[i]] = qy[i];
53         qx[i] = app[qx[i]];
54     }
55     for (int i = 1; i <= n2; i++) a[i] = 0;
56     for (int i = 0; i < tm; i++) {
57         ri = find(vd[x[id[i]]]); rj = find(vd[y[id[i]]]);
58         if (ri != rj)
59             a[ri] = rj, Nx[m2] = vd[x[id[i]]], Ny[m2] = vd[y[id[i]]], Nz[m2] = z[id[i]], m2++;
60     }
61     int mid = Q / 2;
62     solve(qx, qy, mid, n2, Nx, Ny, Nz, m2, ans);
63     solve(qx + mid, qy + mid, Q - mid, n2, Nx, Ny, Nz, m2, ans);
64 }
65 void work() { if (Q) solve(qx, qy, Q, n, x, y, z, m, 0); }
66 int main() { init(); work(); return 0; }

```

#### 4.15 弦图

- 任何一个弦图都至少有一个单纯点, 不是完全图的弦图至少有两个不相邻的单纯点.

- 设第  $i$  个点在弦图的完美消除序列第  $p(i)$  个. 令  $N(v) = \{w | w \text{ 与 } v \text{ 相邻且 } p(w) > p(v)\}$  弦图的极大团一定是  $v \cup N(v)$  的形式.

- 弦图最多有  $n$  个极大团.

- 设  $next(v)$  表示  $N(v)$  中最前的点. 令  $w^*$  表示所有满足  $A \in B$  的  $w$  中最后的一个点. 判断  $v \cup N(v)$  是否为极大团, 只需判断是否存在一个  $w$ , 满足  $Next(w) = v$  且  $|N(v)| + 1 \leq |N(w)|$  即可.

- 最小染色: 完美消除序列从后往前依次给每个点染色, 给每个点染上可以染的最小的颜色. (团数 = 色数)

- 最大独立集: 完美消除序列从前往后能选就选.

- 最小团覆盖: 设最大独立集为  $\{p_1, p_2, \dots, p_t\}$ , 则  $\{p_1 \cup N(p_1), \dots, p_t \cup N(p_t)\}$  为最小团覆盖. (最大独立集数 = 最小团覆盖数)

```

1  class Chordal { // 1-Based, G is the Graph, must be sorted before call Check_Chordal
2  public: // Construct will sort it automatically
3      int v[Maxn], id[Maxn]; bool inseq[Maxn]; priority_queue<pair<int, int> > pq;
4      vector<int> Construct_Perfect_Elimination_Sequence(vector<int> *G, int n) { // O(m + n log n)
5          vector<int> seq(n + 1, 0);
6          for (int i = 0; i <= n; ++i) inseq[i] = false, sort(G[i].begin(), G[i].end()), v[i] = 0;
7          int cur = n; pair<int, int> Mx; while(!pq.empty()) pq.pop(); pq.push(make_pair(0, 1));

```

```

8   for (int i = n; i >= 1; --i) {
9       while (!pq.empty() && (Mx = pq.top(), inseq[Mx.second] || Mx.first != v[Mx.second])) pq.pop();
10      id[Mx.second] = cur;
11      int x = seq[cur--] = Mx.second, sz = (int)G[Mx.second].size(); inseq[x] = true;
12      for (int j = 0; j < sz; ++j) {
13          int y = G[x][j]; if(!inseq[y]) pq.push(make_pair(++v[y], y));
14      }
15  } return seq;
16
17  bool Check_Chordal(vector<int> *G, vector<int> &seq, int n) { // 0(n + mlogn), plz gen seq first
18      bool isChordal = true;
19      for (int i = n - 1; i >= 1 && isChordal; --i) {
20          int x = seq[i], sz, y = -1;
21          if ((sz = (int)G[x].size()) == 0) continue;
22          for(int j = 0; j < sz; ++j) {
23              if (id[G[x][j]] < i) continue;
24              if (y == -1 || id[y] > id[G[x][j]]) y = G[x][j];
25          } if (y == -1) continue;
26          for (int j = 0; j < sz; ++j) {
27              int y1 = G[x][j]; if (id[y1] < i) continue;
28              if (y1 == y || binary_search(G[y].begin(), G[y].end(), y1)) continue;
29              isChordal = false; break;
30          }
31      } return isChordal;
32  }
33  };

```

#### 4.16 K 短路 (允许重复)

```

1  #define for_each(it, v) for (vector<Edge*>::iterator it = (v).begin(); it != (v).end(); ++it)
2  const int MAX_N = 10000, MAX_M = 50000, MAX_K = 10000, INF = 1000000000;
3  struct Edge { int from, to, weight; };
4  struct HeapNode { Edge* edge; int depth; HeapNode* child[4]; }; // child[0..1] for heap G, child[2..3] for heap
   out edge
5  int n, m, k, s, t; Edge* edge[MAX_M];
6  int dist[MAX_N]; Edge* prev[MAX_N];
7  vector<Edge*> graph[MAX_N]; vector<Edge*> graphR[MAX_N];
8  HeapNode* nullNode; HeapNode* heapTop[MAX_N];
9
10 HeapNode* createHeap(HeapNode* curNode, HeapNode* newNode) {
11     if (curNode == nullNode) return newNode; HeapNode* rootNode = new HeapNode;
12     memcpy(rootNode, curNode, sizeof(HeapNode));
13     if (newNode->edge->weight < curNode->edge->weight) {
14         rootNode->edge = newNode->edge; rootNode->child[2] = newNode->child[2]; rootNode->child[3] = newNode->
   child[3];
15         newNode->edge = curNode->edge; newNode->child[2] = curNode->child[2]; newNode->child[3] = curNode->child
   [3];
16     } if (rootNode->child[0]->depth < rootNode->child[1]->depth) rootNode->child[0] = createHeap(rootNode->child
   [0], newNode);
17     else rootNode->child[1] = createHeap(rootNode->child[1], newNode);
18     rootNode->depth = max(rootNode->child[0]->depth, rootNode->child[1]->depth) + 1;
19     return rootNode;
20
21 }
22 bool heapNodeMoreThan(HeapNode* node1, HeapNode* node2) { return node1->edge->weight > node2->edge->weight; }
23
24 int main() {
25     scanf("%d%d%d", &n, &m, &k); scanf("%d%d", &s, &t); s--, t--;
26     while (m--) { Edge* newEdge = new Edge;
27         int i, j, w; scanf("%d%d%d", &i, &j, &w);
28         i--, j--; newEdge->from = i; newEdge->to = j; newEdge->weight = w;
29         graph[i].push_back(newEdge); graphR[j].push_back(newEdge);
30     }
31     //Dijkstra
32     queue<int> dfsOrder; memset(dist, -1, sizeof(dist));
33     typedef pair<int, pair<int, Edge*>> > DijkstraQueueItem;
34     priority_queue<DijkstraQueueItem, vector<DijkstraQueueItem>, greater<DijkstraQueueItem>> dq;
35     dq.push(make_pair(0, make_pair(t, (Edge*) NULL)));
36     while (!dq.empty()) {
37         int d = dq.top().first; int i = dq.top().second.first;
38         Edge* edge = dq.top().second.second; dq.pop();
39         if (dist[i] != -1) continue;
40         dist[i] = d; prev[i] = edge; dfsOrder.push(i);

```

```

41     for_each(it, graphR[i]) dq.push(make_pair(d + (*it)->weight, make_pair((*it)->from, *it)));
42 }
43 //Create edge heap
44 nullNode = new HeapNode; nullNode->depth = 0; nullNode->edge = new Edge; nullNode->edge->weight = INF;
45 fill(nullNode->child, nullNode->child + 4, nullNode);
46 while (!dfsOrder.empty()) {
47     int i = dfsOrder.front(); dfsOrder.pop();
48     if (prev[i] == NULL) heapTop[i] = nullNode;
49     else heapTop[i] = heapTop[prev[i]->to];
50     vector<HeapNode*> heapNodeList;
51     for_each(it, graph[i]) { int j = (*it)->to; if (dist[j] == -1) continue;
52         (*it)->weight += dist[j] - dist[i]; if (prev[i] != *it) {
53             HeapNode* curNode = new HeapNode;
54             fill(curNode->child, curNode->child + 4, nullNode);
55             curNode->depth = 1; curNode->edge = *it;
56             heapNodeList.push_back(curNode);
57         }
58     } if (!heapNodeList.empty()) { //Create heap out
59         make_heap(heapNodeList.begin(), heapNodeList.end(), heapNodeMoreThan);
60         int size = heapNodeList.size();
61         for (int p = 0; p < size; p++) {
62             heapNodeList[p]->child[2] = 2 * p + 1 < size ? heapNodeList[2 * p + 1] : nullNode;
63             heapNodeList[p]->child[3] = 2 * p + 2 < size ? heapNodeList[2 * p + 2] : nullNode;
64             heapTop[i] = createHeap(heapTop[i], heapNodeList.front());
65         }
66     } //Walk on DAG
67     typedef pair<long long, HeapNode*> DAGQueueItem;
68     priority_queue<DAGQueueItem, vector<DAGQueueItem>, greater<DAGQueueItem>> aq;
69     if (dist[s] == -1) printf("NO\n");
70     else { printf("%d\n", dist[s]);
71         if (heapTop[s] != nullNode) aq.push(make_pair(dist[s] + heapTop[s]->edge->weight, heapTop[s]));
72     } k--; while (k--) {
73         if (aq.empty()) { printf("NO\n"); continue; }
74         long long d = aq.top().first; HeapNode* curNode = aq.top().second; aq.pop();
75         printf("%I64d\n", d);
76         if (heapTop[curNode->edge->to] != nullNode)
77             aq.push(make_pair(d + heapTop[curNode->edge->to]->edge->weight, heapTop[curNode->edge->to]));
78         for (int i = 0; i < 4; i++) if (curNode->child[i] != nullNode)
79             aq.push(make_pair(d + curNode->edge->weight + curNode->child[i]->edge->weight, curNode->child[i]));
80     } return 0;
81 }

```

#### 4.17 K 短路 (不允许重复)

```

1  int Num[10005][205], Path[10005][205], dev[10005], from[10005], value[10005], dist[205], Next[205], Graph
   [205][205];
2  int N, M, K, s, t, tot, cnt; bool forbid[205], hasNext[10005][205];
3  struct cmp {
4      bool operator()(const int &a, const int &b) {
5          int *i, *j; if (value[a] != value[b]) return value[a] > value[b];
6          for (i = Path[a], j = Path[b]; (*i) == (*j); i++, j++);
7          return (*i) > (*j);
8      }
9  };
10 void Check(int idx, int st, int *path, int &res) {
11     int i, j; for (i = 0; i < N; i++) dist[i] = 1000000000, Next[i] = t;
12     dist[t] = 0; forbid[t] = true; j = t;
13     for ( ; ; ) {
14         for (i = 0; i < N; i++) if (!forbid[i] && (i != st || !hasNext[idx][j]) && (dist[j] + Graph[i][j] < dist[
   i] || (dist[j] + Graph[i][j] == dist[i] && j < Next[i])))
15             Next[i] = j, dist[i] = dist[j] + Graph[i][j];
16         j = -1; for (i = 0; i < N; i++) if (!forbid[i] && (j == -1 || dist[i] < dist[j])) j = i;
17         if (j == -1) break; forbid[j] = 1; if (j == st) break;
18     } res += dist[st]; for (i = st; i != t; i = Next[i], path++) (*path) = i; (*path) = i;
19 }
20 int main() {
21     int i, j, k, l;
22     while (scanf("%d%d%d%d", &N, &M, &K, &s, &t) && N) {
23         priority_queue<int, vector<int>, cmp> Q;
24         for (i = 0; i < N; i++) for (j = 0; j < N; j++) Graph[i][j] = 1000000000;
25         for (i = 0; i < M; i++) { scanf("%d%d%d", &j, &k, &l); Graph[j - 1][k - 1] = 1; }
26         s--; t--;

```

```
27     memset(forbid, false, sizeof(forbid)); memset(hasNext[0], false, sizeof(hasNext[0]));
28     Check(0, s, Path[0], value[0]); dev[0] = 0; from[0] = 0; Num[0][0] = 0; Q.push(0);
29     cnt = 1; tot = 1;
30     for (i = 0; i < K; i++) {
31         if (Q.empty()) break; l = Q.top(); Q.pop();
32         for (j = 0; j <= dev[l]; j++) Num[l][j] = Num[from[l]][j];
33         for (; Path[l][j] != t; j++) {
34             memset(hasNext[tot], false, sizeof(hasNext[tot])); Num[l][j] = tot++;
35             } for (j = 0; Path[l][j] != t; j++) hasNext[Num[l][j]][Path[l][j] + 1] = true;
36             for (j = dev[l]; Path[l][j] != t; j++) {
37                 memset(forbid, false, sizeof(forbid)); value[cnt] = 0;
38                 for (k = 0; k < j; k++) {
39                     forbid[Path[l][k]] = true;
40                     Path[cnt][k] = Path[l][k];
41                     value[cnt] += Graph[Path[l][k]][Path[l][k + 1]];
42                 } Check(Num[l][j], Path[l][j], &Path[cnt][j], value[cnt]);
43                 if (value[cnt] > 2000000) continue;
44                 dev[cnt] = j; from[cnt] = l; Q.push(cnt); cnt++;
45             }
46         }
47         if (i < K || value[l] > 2000000) printf("None\n");
48         else {
49             for (i = 0; Path[l][i] != t; i++) printf("%d-", Path[l][i] + 1);
50             printf("%d\n", t + 1);
51         }
52     } return 0;
53 }
```

4.18 小知识

- 平面图: 一定存在一个度小于等于 5 的点.  $E \leq 3V - 6$ . 欧拉公式:  $V + F - E = 1 +$  连通块数
- 图连通度:
  - 1.  $k$ - 连通 ( $k$ -connected): 对于任意一对结点都至少存在结点各不相同的  $k$  条路
  - 2. 点连通度 ( $vertex\ connectivity$ ): 把图变成非连通图所需删除的最少点数
  - 3. Whitney 定理: 一个图是  $k$ - 连通的当且仅当它的点连通度至少为  $k$
- Lindstroem-Gessel-Viennot Lemma: 给定一个图的  $n$  个起点和  $n$  个终点, 令  $A_{ij}$  = 第  $i$  个起点到第  $j$  个终点的路径条数, 则从起点到终点的不相交路径条数为  $det(A)$
- 欧拉回路与树形图的联系: 对于出度等于入度的连通图  $s(G) = t_i(G) \prod_{j=1}^n (d^+(v_j) - 1)!$
- 密度子图: 给定无向图, 选取点集及其导出子图, 最大化  $W_e + P_v$  (点权可负).
  - $(S, u) = U, (u, T) = U - 2P_u - D_u, (u, v) = (v, u) = W_e$
  - $ans = \frac{Un - C[S, T]}{2}$ , 解集为  $S - \{s\}$
- 最大权闭合图: 选  $a$  则  $a$  的后继必须被选
  - $P_u > 0, (S, u) = P_u, P_u < 0, (u, T) = -P_u$
  - $ans = \sum_{P_u > 0} P_u - C[S, T]$ , 解集为  $S - \{s\}$
- 判定边是否属于最小割:
  - 可能属于最小割:  $(u, v)$  不属于同一 SCC
  - 一定在所有最小割中:  $(u, v)$  不属于同一 SCC, 且  $S, u$  在同一 SCC,  $u, T$  在同一 SCC
- 图同构 Hash:  $F_t(i) = (F_{t-1}(i) \times A + \sum_{i \rightarrow j} F_{t-1}(j) \times B + \sum_{j \leftarrow i} F_{t-1}(j) \times C + D \times (i = a)) \pmod P$ , 枚举点  $a$ , 迭代  $K$  次后求得的  $F_k(a)$  就是  $a$  点所对应的 Hash 值.

5 数学

5.1 博弈论相关

1. Anti-SG:  
规则与 Nim 基本相同, 取最后一个的输。  
先手必胜当且仅当:  
(1) 所有堆的石子数都为 1 且游戏的 SG 值为 0;  
(2) 有些堆的石子数大于 1 且游戏的 SG 值不为 0。
2. SJ 定理:  
对于任意一个 Anti-SG 游戏, 如果我们规定当局面中, 所有的单一游戏的 SG 值为 0 时, 游戏结束, 则先手必胜当且仅当:  
(1) 游戏的 SG 函数不为 0 且游戏中某个单一游戏的 SG 函数大于 1;  
(2) 游戏的 SG 函数为 0 且游戏中没有单一游戏的 SG 函数大于 1。
3. Multi-SG 游戏:  
可以将一堆石子分成多堆。
4. Every-SG 游戏:  
每一个可以移动的棋子都要移动。  
对于我们可以赢的单一游戏, 我们一定要拿到这一场游戏的胜利。  
只需要考虑如何让我们必胜的游戏尽可能长的玩下去, 对手相反。  
于是就来一个 DP,  
 $step[v] = 0;$  ( $v$  为终止状态)  
 $step[v] = \maxstep[u] + 1;$  ( $sg[v]>0, sg[u]=0$ )  
 $step[v] = \minstep[u] + 1;$  ( $sg[v]=0$ )
5. 翻硬币游戏:  
N 枚硬币排成一排, 有的正面朝上, 有的反面朝上。游戏者根据某些约束翻硬币 (如: 每次只能翻一或两枚, 或者每次只能翻连续的几枚), 但他所翻动的硬币中, 最右边的必须是从正面翻到反面。谁不能翻谁输。  
结论: 局面的 SG 值为局面中每个正面朝上的棋子单一存在时的 SG 值的异或和。可用数学归纳法证明。
6. 无向树删边游戏:  
规则如下:  
给出一个有 N 个点的树, 有一个点作为树的根节点。游戏者轮流从树中删去边, 删去一条边后, 不与根节点相连的部分将被移走。谁无路可走谁输。  
结论:  
叶子节点的 SG 值为 0; 中间节点的 SG 值为它的所有子节点的 SG 值加 1 后的异或和。是用数学归纳法证明。
7. Christmas Game(PKU3710):  
题目大意:  
有 N 个局部联通的图。Harry 和 Sally 轮流从图中删边, 删去一条边后, 不与根节点相连的部分将被移走。Sally 为先手。图是通过从基础树中加一些边得到的。所有形成的环保证不共用边, 且只与基础树有一个公共点。谁无路可走谁输。环的处理成为了解题的关键。  
性质:  
(1) 对于长度为奇数的环, 去掉其中任意一个边之后, 剩下的两个链长度同奇偶, 抑或之后的 SG 值不可能为奇数, 所以它的 SG 值为 1;  
(2) 对于长度为偶数的环, 去掉其中任意一个边之后, 剩下的两个链长度异奇偶, 抑或之后的 SG 值不可能为 0, 所以它的 SG 值为 0; 所以我们可以去掉所有的偶环, 将所有的奇环变为长短为 1 的链。  
这样的话, 我们已经将这题改造成了上一节的模型。
8. 无向图的删边游戏:  
我们将 Christmas Game 这道题进行一步拓展——去掉对环的限制条件, 这个模型应该怎样处理?  
无向图的删边游戏:  
一个无向联通图, 有一个点作为图的根。游戏者轮流从图中删去边, 删去一条边后, 不与根节点相连的部分将被移走。谁无路可走谁输。  
结论:  
对无向图做如下改动: 将图中的任意一个偶环缩成一个新点, 任意一个奇环缩成一个新点加一个新边; 所有连

到原先环上的边全部改为与新点相连。这样的改动不会影响图的 SG 值。

9. Staircase nim:  
楼梯从地面由下向上编号为 0 到 n。游戏者在每次操作时可以将楼梯  $j(1 \leq j \leq n)$  上的任意多但至少一个硬币移动到楼梯  $j-1$  上。将最后一枚硬币移至地上的人获胜。  
结论:  
设该游戏 Sg 函数为奇数格棋子数的 Xor 和 S。  
如果  $S=0$ , 则先手必败, 否则必胜。

5.2 单纯形 Cpp

max {cx|Ax ≤ b, x ≥ 0}

```
1 const int MAXN = 11000, MAXM = 1100;
2 // `here MAXN is the MAX number of conditions, MAXM is the MAX number of vars`
3
4 int avali[MAXN], avacnt;
5 double A[MAXN][MAXM];
6 double b[MAXN], c[MAXN];
7 double* simplex(int n, int m) {
8 // `here n is the number of conditions, m is the number of vars`
9     m++;
10    int r = n, s = m - 1;
11    static double D[MAXN + 2][MAXM + 1];
12    static int ix[MAXN + MAXM];
13    for (int i = 0; i < n + m; i++) ix[i] = i;
14    for (int i = 0; i < n; i++) {
15        for (int j = 0; j < m - 1; j++) D[i][j] = -A[i][j];
16        D[i][m - 1] = 1;
17        D[i][m] = b[i];
18        if (D[r][m] > D[i][m]) r = i;
19    }
20    for (int j = 0; j < m - 1; j++) D[n][j] = c[j];
21    D[n + 1][m - 1] = -1;
22    for (double d; ; ) {
23        if (r < n) {
24            int t = ix[s]; ix[s] = ix[r + m]; ix[r + m] = t;
25            D[r][s] = 1.0 / D[r][s];
26            for (int j = 0; j <= m; j++) if (j != s) D[r][j] *= -D[r][s];
27            avacnt = 0;
28            for (int i = 0; i <= m; ++i)
29                if (fabs(D[r][i]) > EPS)
30                    avali[avacnt++] = i;
31            for (int i = 0; i <= n + 1; i++) if (i != r) {
32                if (fabs(D[i][s]) < EPS) continue;
33                double *cur1 = D[i], *cur2 = D[r], tmp = D[i][s];
34                //for (int j = 0; j <= m; j++) if (j != s) cur1[j] += cur2[j] * tmp;
35                for (int j = 0; j < avacnt; ++j) if (avali[j] != s) cur1[avali[j]] += cur2[avali[j]] * tmp;
36                D[i][s] *= D[r][s];
37            }
38        }
39        r = -1; s = -1;
40        for (int j = 0; j < m; j++) if (s < 0 || ix[s] > ix[j]) {
41            if (D[n + 1][j] > EPS || D[n + 1][j] > -EPS && D[n][j] > EPS) s = j;
42        }
43        if (s < 0) break;
44        for (int i = 0; i < n; i++) if (D[i][s] < -EPS) {
45            if (r < 0 || (d = D[r][m] / D[r][s] - D[i][m] / D[i][s]) < -EPS
46                || d < EPS && ix[r + m] > ix[i + m])
47                r = i;
48        }
49        if (r < 0) return null; // `非有界`
50    }
51    if (D[n + 1][m] < -EPS) return null; // `无法执行`
52    static double x[MAXN - 1];
53    for (int i = m; i < n + m; i++) if (ix[i] < m - 1) x[ix[i]] = D[i - m][m];
54    return x; // `值为 $D[n][m]$`
55 }
```

5.3 自适应辛普森

```
1 double area(const double &left, const double &right) {
2     double mid = (left + right) / 2;
3     return (right - left) * (calc(left) + 4 * calc(mid) + calc(right)) / 6;
4 }
5
6 double simpson(const double &left, const double &right,
7               const double &eps, const double &area_sum) {
8     double mid = (left + right) / 2;
9     double area_left = area(left, mid);
10    double area_right = area(mid, right);
11    double area_total = area_left + area_right;
12    if (std::abs(area_total - area_sum) < 15 * eps) {
13        return area_total + (area_total - area_sum) / 15;
14    }
15    return simpson(left, mid, eps / 2, area_left)
16           + simpson(mid, right, eps / 2, area_right);
17 }
18
19 double simpson(const double &left, const double &right, const double &eps) {
20     return simpson(left, right, eps, area(left, right));
21 }
```

5.4 FFT

```
1 namespace FFT {
2     #define mul(a, b) (Complex(a.x * b.x - a.y * b.y, a.x * b.y + a.y * b.x))
3     struct Complex { }; // `something omitted`
4     void FFT(Complex P[], int n, int oper) {
5         for (int i = 1, j = 0; i < n - 1; i++) {
6             for (int s = n; j ^= s >= 1, ~j & s; );
7             if (i < j) swap(P[i], P[j]);
8         }
9         for (int d = 0; (1 << d) < n; d++) {
10            int m = 1 << d, m2 = m * 2;
11            double p0 = PI / m * oper;
12            Complex unit_p0(cos(p0), sin(p0));
13            for (int i = 0; i < n; i += m2) {
14                Complex unit(1.0, 0.0);
15                for (int j = 0; j < m; j++) {
16                    Complex &P1 = P[i + j + m], &P2 = P[i + j];
17                    Complex t = mul(unit, P1);
18                    P1 = Complex(P2.x - t.x, P2.y - t.y);
19                    P2 = Complex(P2.x + t.x, P2.y - t.y);
20                    unit = mul(unit, unit_p0);
21                }
22            }
23        }
24        vector<int> doFFT(const vector<int> &a, const vector<int> &b) {
25            vector<int> ret(max(0, (int) a.size() + (int) b.size() - 1), 0);
26            static Complex A[MAXB], B[MAXB], C[MAXB];
27            int len = 1; while (len < (int)ret.size()) len *= 2;
28            for (int i = 0; i < len; i++) A[i] = i < (int)a.size() ? a[i] : 0;
29            for (int i = 0; i < len; i++) B[i] = i < (int)b.size() ? b[i] : 0;
30            FFT(A, len, 1); FFT(B, len, 1);
31            for (int i = 0; i < len; i++) C[i] = mul(A[i], B[i]);
32            FFT(C, len, -1);
33            for (int i = 0; i < (int)ret.size(); i++)
34                ret[i] = (int) (C[i].x / len + 0.5);
35            return ret;
36        }
37    }
```

5.5 整数 FFT

```
1 namespace FFT {
2     // `替代方案: $23068673( = 11 * 2 ^ {21} + 1)$, 原根为 $3$`
3     const int MOD = 786433, PRIMITIVE_ROOT = 10; // `$3 * 2 ^ {18} + 1$`
4     const int MAXB = 1 << 20;
5     int getMod(int downLimit) { // `或者现场自己找一个MOD`
6         for (int c = 3; ; ++c) { int t = (c << 21) | 1;
```



```

7         if (t >= downLimit && isPrime(t)) return t;
8     }}
9     int modInv(int a) { return a <= 1 ? a : (long long) (MOD - MOD / a) * modInv(MOD % a) % MOD; }
10    void NTT(int P[], int n, int oper) {
11        for (int i = 1, j = 0; i < n - 1; i++) {
12            for (int s = n; j ^= s >= 1, ~j & s;);
13            if (i < j) swap(P[i], P[j]);
14        }
15        for (int d = 0; (1 << d) < n; d++) {
16            int m = 1 << d, m2 = m * 2;
17            long long unit_p0 = powMod(PRIMITIVE_ROOT, (MOD - 1) / m2);
18            if (oper < 0) unit_p0 = modInv(unit_p0);
19            for (int i = 0; i < n; i += m2) {
20                long long unit = 1;
21                for (int j = 0; j < m; j++) {
22                    int &P1 = P[i + j + m], &P2 = P[i + j];
23                    int t = unit * P1 % MOD;
24                    P1 = (P2 - t + MOD) % MOD; P2 = (P2 + t) % MOD;
25                    unit = unit * unit_p0 % MOD;
26                }
27            }
28            vector<int> mul(const vector<int> &a, const vector<int> &b) {
29                vector<int> ret(max(0, (int) a.size() + (int) b.size() - 1), 0);
30                static int A[MAXB], B[MAXB], C[MAXB];
31                int len = 1; while (len < (int)ret.size()) len <= 1;
32                for (int i = 0; i < len; i++) A[i] = i < (int)a.size() ? a[i] : 0;
33                for (int i = 0; i < len; i++) B[i] = i < (int)b.size() ? b[i] : 0;
34                NTT(A, len, 1); NTT(B, len, 1);
35                for (int i = 0; i < len; i++) C[i] = (long long) A[i] * B[i] % MOD;
36                NTT(C, len, -1); for (int i = 0, inv = modInv(len); i < (int)ret.size(); i++) ret[i] = (long long) C[i] *
37                    inv % MOD;
38                return ret;
39            }
40        }
41    }
42 }

```

## 5.6 扩展欧几里得

$$ax + by = g = \gcd(x, y)$$

```

1 void exgcd(LL x, LL y, LL &a0, LL &b0, LL &g) {
2     LL a1 = b0 = 0, b1 = a0 = 1, t;
3     while (y != 0) {
4         t = a0 - x / y * a1, a0 = a1, a1 = t;
5         t = b0 - x / y * b1, b0 = b1, b1 = t;
6         t = x % y, x = y, y = t;
7     } if (x < 0) a0 = -a0, b0 = -b0, x = -x;
8     g = x;
9 }

```

## 5.7 线性同余方程

- 中国剩余定理: 设  $m_1, m_2, \dots, m_k$  两两互素, 则同余方程组  $x \equiv a_i \pmod{m_i}$  for  $i = 1, 2, \dots, k$  在  $[0, M = m_1 m_2 \dots m_k)$  内有唯一解. 记  $M_i = M / m_i$ , 找出  $p_i$  使得  $M_i p_i \equiv 1 \pmod{m_i}$ , 记  $e_i = M_i p_i$ , 则  $x \equiv e_1 a_1 + e_2 a_2 + \dots + e_k a_k \pmod{M}$

- 多变元线性同余方程组: 方程的形式为  $a_1 x_1 + a_2 x_2 + \dots + a_n x_n + b \equiv 0 \pmod{m}$ , 令  $d = (a_1, a_2, \dots, a_n, m)$ , 有解的充要条件是  $d | b$ , 解的个数为  $m^{n-1} d$

## 5.8 Miller-Rabin 素性测试

```

1 bool test(LL n, int base) {
2     LL m = n - 1, ret = 0; int s = 0;
3     for (; m % 2 == 0; ++s) m >>= 1; ret = pow_mod(base, m, n);
4     if (ret == 1 || ret == n - 1) return true;
5     for (--s; s >= 0; --s) {
6         ret = multiply_mod(ret, ret, n); if (ret == n - 1) return true;
7     } return false;
8 }

```

```

9 LL special[7] = {
10     1373653LL,          25326001LL,
11     3215031751LL,       25000000000LL,
12     2152302898747LL,    3474749660383LL, 341550071728321LL};
13 /*
14  * n < 2047                test[] = {2}
15  * n < 1,373,653          test[] = {2, 3}
16  * n < 9,080,191          test[] = {31, 73}
17  * n < 25,326,001         test[] = {2, 3, 5}
18  * n < 4,759,123,141      test[] = {2, 7, 61}
19  * n < 1,122,004,669,633  test[] = {2, 13, 23, 1662803}
20  * n < 2,152,302,898,747  test[] = {2, 3, 5, 7, 11}
21  * n < 3,474,749,660,383  test[] = {2, 3, 5, 7, 11, 13}
22  * n < 341,550,071,728,321 test[] = {2, 3, 5, 7, 11, 13, 17}
23  * n < 3,825,123,056,546,413,051 test[] = {2, 3, 5, 7, 11, 13, 17, 19, 23}
24 */
25 bool is_prime(LL n) {
26     if (n < 2) return false;
27     if (n < 4) return true;
28     if (!test(n, 2) || !test(n, 3)) return false;
29     if (n < special[0]) return true;
30     if (!test(n, 5)) return false;
31     if (n < special[1]) return true;
32     if (!test(n, 7)) return false;
33     if (n == special[2]) return false;
34     if (n < special[3]) return true;
35     if (!test(n, 11)) return false;
36     if (n < special[4]) return true;
37     if (!test(n, 13)) return false;
38     if (n < special[5]) return true;
39     if (!test(n, 17)) return false;
40     if (n < special[6]) return true;
41     return test(n, 19) && test(n, 23) && test(n, 29) && test(n, 31) && test(n, 37);
42 }

```

## 5.9 PollardRho

```

1 LL pollardRho(LL n, LL seed) {
2     LL x, y, head = 1, tail = 2; x = y = random() % (n - 1) + 1;
3     for (; ; ) {
4         x = addMod(multiplyMod(x, x, n), seed, n);
5         if (x == y) return n; LL d = gcd(myAbs(x - y), n);
6         if (1 < d && d < n) return d;
7         if (++head == tail) y = x, tail <= 1;
8     } vector<LL> divisors;
9     void factorize(LL n) { // 需要保证 n > 1
10        if (isPrime(n)) divisors.push_back(n);
11        else { LL d = n;
12            while (d >= n) d = pollardRho(n, random() % (n - 1) + 1);
13            factorize(n / d); factorize(d);
14        }
15    }

```

## 5.10 多项式求根

```

1 const double error = 1e-12;
2 const double inf = 1e+12;
3 int n; double a[10], x[10];
4 double f(double a[], int n, double x) {
5     double tmp = 1, sum = 0;
6     for (int i = 0; i <= n; i++) sum = sum + a[i] * tmp, tmp = tmp * x;
7     return sum;
8 }
9 double binary(double l, double r, double a[], int n) {
10    int sl = sign(f(a, n, l)), sr = sign(f(a, n, r));
11    if (sl == 0) return l; if (sr == 0) return r;
12    if (sl * sr > 0) return inf;
13    while (r - l > error) {
14        double mid = (l + r) / 2;
15        int ss = sign(f(a, n, mid));
16        if (ss == 0) return mid;
17        if (ss * sl > 0) l = mid; else r = mid;
18    }

```



```
18     } return l;
19 }
20 void solve(int n, double a[], double x[], int &nx) {
21     if (n == 1) { x[1] = -a[0] / a[1]; nx = 1; return; }
22     double da[10], dx[10]; int ndx;
23     for (int i = n; i >= 1; i--) da[i - 1] = a[i] * i;
24     solve(n - 1, da, dx, ndx); nx = 0;
25     if (ndx == 0) {
26         double tmp = binary(-infi, infi, a, n);
27         if (tmp < infi) x[++nx] = tmp; return;
28     } double tmp = binary(-infi, dx[1], a, n);
29     if (tmp < infi) x[++nx] = tmp;
30     for (int i = 1; i <= ndx - 1; i++) {
31         tmp = binary(dx[i], dx[i + 1], a, n);
32         if (tmp < infi) x[++nx] = tmp;
33     } tmp = binary(dx[ndx], infi, a, n);
34     if (tmp < infi) x[++nx] = tmp;
35 }
36 int main() {
37     scanf("%d", &n);
38     for (int i = n; i >= 0; i--) scanf("%lf", &a[i]);
39     int nx; solve(n, a, x, nx);
40     for (int i = 1; i <= nx; i++) printf("%.6f\n", x[i]);
41     return 0;
42 }
```

5.11 线性递推

for  $a_{i+n} = (\sum_{i=0}^{n-1} k_j a_{i+j}) + d$ ,  $a_m = (\sum_{i=0}^{n-1} c_i a_i) + c_n d$

```
1 vector<int> recFormula(int n, int k[], int m) {
2     vector<int> c(n + 1, 0);
3     if (m < n) c[m] = 1;
4     else {
5         static int a[MAX_K * 2 + 1];
6         vector<int> b = recFormula(n, k, m >> 1);
7         for (int i = 0; i < n + n; ++i) a[i] = 0;
8         int s = m & 1;
9         for (int i = 0; i < n; i++) {
10             for (int j = 0; j < n; j++) a[i + j + s] += b[i] * b[j];
11             c[n] += b[i];
12         } c[n] = (c[n] + 1) * b[n];
13         for (int i = n * 2 - 1; i >= n; i--) {
14             int add = a[i]; if (add == 0) continue;
15             for (int j = 0; j < n; j++) a[i - n + j] += k[j] * add;
16             c[n] += add;
17         } for (int i = 0; i < n; ++i) c[i] = a[i];
18     } return c;
19 }
```

5.12 离散对数

$A^x \equiv B \pmod C$ , 对非质数  $C$  也适用.

```
1 int modLog(int A, int B, int C) {
2     static pii baby[MAX_SQRT_C + 11];
3     int d = 0; LL k = 1, D = 1; B %= C;
4     for (int i = 0; i < 100; ++i, k = k * A % C) // '$[0, \log C]$',
5         if (k == B) return i;
6     for (int g; ; ++d) {
7         g = gcd(A, C); if (g == 1) break;
8         if (B % g != 0) return -1;
9         B /= g; C /= g; D = (A / g * D) % C;
10    } int m = (int) ceil(sqrt((double) C)); k = 1;
11    for (int i = 0; i <= m; ++i, k = k * A % C) baby[i] = pii(k, i);
12    sort(baby, baby + m + 1); // [0, m]
13    int n = unique(baby, baby + m + 1, equalFirst) - baby, am = powMod(A, m, C);
14    for (int i = 0; i <= m; ++i) {
15        LL e, x, y; exgcd(D, C, x, y, e); e = x * B % C;
16        if (e < 0) e += C;
17        if (e >= 0) {
```

```
18         int k = lower_bound(baby, baby + n, pii(e, -1)) - baby;
19         if (baby[k].first == e) return i * m + baby[k].second + d;
20     } D = D * am % C;
21 } return -1;
22 }
```

5.13 平方剩余

- Legendre Symbol: 对奇质数  $p$ ,  $(\frac{a}{p}) = \begin{cases} 1 & \text{是平方剩余} \\ -1 & \text{是非平方剩余} \\ 0 & a \equiv 0 \pmod p \end{cases} = a^{\frac{p-1}{2}} \pmod p$

- 若  $p$  是奇质数,  $(\frac{-1}{p}) = 1$  当且仅当  $p \equiv 1 \pmod 4$

- 若  $p$  是奇质数,  $(\frac{2}{p}) = 1$  当且仅当  $p \equiv \pm 1 \pmod 8$

- 若  $p, q$  是奇素数且互质,  $(\frac{p}{q})(\frac{q}{p}) = (-1)^{\frac{p-1}{2} \times \frac{q-1}{2}}$

- Jacobi Symbol: 对奇数  $n = p_1^{\alpha_1} p_2^{\alpha_2} \cdots p_k^{\alpha_k}$ ,  $(\frac{a}{n}) = (\frac{a}{p_1})^{\alpha_1} (\frac{a}{p_2})^{\alpha_2} \cdots (\frac{a}{p_k})^{\alpha_k}$

- Jacobi Symbol 为  $-1$  则一定不是平方剩余, 所有平方剩余的 Jacobi Symbol 都是 1, 但 1 不一定是平方剩余

$ax^2 + bx + c \equiv 0 \pmod p$ , 其中  $a \not\equiv 0 \pmod p$ , 且  $p$  是质数

```
1 inline int normalize(LL a, int P) { a %= P; return a < 0 ? a + P : a; }
2 vector<int> QuadraticResidue(LL a, LL b, LL c, int P) {
3     int h, t; LL r1, r2, delta, pb = 0;
4     a = normalize(a, P); b = normalize(b, P); c = normalize(c, P);
5     if (P == 2) { vector<int> res;
6         if (c % P == 0) res.push_back(0);
7         if ((a + b + c) % P == 0) res.push_back(1);
8         return res;
9     } delta = b * rev(a + a, P) % P;
10    a = normalize(-c * rev(a, P) + delta * delta, P);
11    if (powMod(a, P / 2, P) + 1 == P) return vector<int>(0);
12    for (t = 0, h = P / 2; h % 2 == 0; ++t, h /= 2);
13    r1 = powMod(a, h / 2, P);
14    if (t > 0) { do b = random() % (P - 2) + 2;
15        while (powMod(b, P / 2, P) + 1 != P); }
16    for (int i = 1; i <= t; ++i) {
17        LL d = r1 * r1 % P * a % P;
18        for (int j = 1; j <= t - i; ++j) d = d * d % P;
19        if (d + 1 == P) r1 = r1 * pb % P; pb = pb * pb % P;
20    } r1 = a * r1 % P; r2 = P - r1;
21    r1 = normalize(r1 - delta, P); r2 = normalize(r2 - delta, P);
22    if (r1 > r2) swap(r1, r2); vector<int> res(1, r1);
23    if (r1 != r2) res.push_back(r2);
24    return res;
25 }
```

5.14 N 次剩余

- 若  $p$  为奇质数,  $a$  为  $p$  的  $n$  次剩余的充要条件是  $a^{\frac{p-1}{(a, p-1)}} \equiv 1 \pmod p$ .

$x^N \equiv a \pmod p$ , 其中  $p$  是质数

```
1 vector<int> solve(int p, int N, int a) {
2     if ((a %= p) == 0) return vector<int>(1, 0);
3     int g = findPrimitiveRoot(p), m = modLog(g, a, p); // g ^ m = a (mod p)
4     if (m == -1) return vector<int>(0);
5     LL B = p - 1, x, y, d; exgcd(N, B, x, y, d);
6     if (m % d != 0) return vector<int>(0);
7     vector<int> ret; x = (x * (m / d) % B + B) % B; // g ^ B mod p = g ^ (p - 1) mod p = 1
```

```
8     for (int i = 0, delta = B / d; i < d; ++i) {
9         x = (x + delta) % B; ret.push_back((int)powMod(g, x, p));
10    } sort(ret.begin(), ret.end());
11    ret.resize(unique(ret.begin(), ret.end()) - ret.begin());
12    return ret;
13 }
```

5.15 Pell 方程

$$\begin{pmatrix} x_k \\ y_k \end{pmatrix} = \begin{pmatrix} x_1 & dy_1 \\ y_1 & x_1 \end{pmatrix}^{k-1} \begin{pmatrix} x_1 \\ y_1 \end{pmatrix}$$

```
1 pair<ULL, ULL> Pell(int n) {
2     static ULL p[50] = {0, 1}, q[50] = {1, 0}, g[50] = {0, 0}, h[50] = {0, 1}, a[50];
3     ULL t = a[2] = Sqrt(n);
4     for (int i = 2; ; ++i) {
5         g[i] = -g[i-1] + a[i] * h[i-1];
6         h[i] = (n - g[i] * g[i]) / h[i-1];
7         a[i+1] = (g[i] + t) / h[i];
8         p[i] = a[i] * p[i-1] + p[i-2];
9         q[i] = a[i] * q[i-1] + q[i-2];
10        if (p[i] * p[i] - n * q[i] * q[i] == 1) return make_pair(p[i], q[i]);
11    } return make_pair(-1, -1);
12 }
```

5.16 小知识

- 勾股数: 设正整数  $n$  的质因数分解为  $n = \prod p_i^{a_i}$ , 则  $x^2 + y^2 = n$  有整数解的充要条件是  $n$  中不存在形如  $p_i \equiv 3 \pmod{4}$  且指数  $a_i$  为奇数的质因数  $p_i$ .  $(\frac{a-b}{2})^2 + ab = (\frac{a+b}{2})^2$ .
- 素勾股数: 若  $m$  和  $n$  互质, 而且  $m$  和  $n$  中有一个是偶数, 则  $a = m^2 - n^2, b = 2mn, c = m^2 + n^2$ , 则  $a, b, c$  是素勾股数.
- Stirling 公式:  $n! \approx \sqrt{2\pi n}(\frac{n}{e})^n$
- Pick 定理: 简单多边形, 不自交, 顶点如果全是整点. 则: 严格在多边形内部的整点数 +  $\frac{1}{2}$ 在边上的整点数 - 1 = 面积
- Mersenne 素数:  $p$  是素数且  $2^p - 1$  的数是素数. (10000 以内的  $p$  有: 2, 3, 5, 7, 13, 17, 19, 31, 61, 89, 107, 127, 521, 607, 1279, 2203, 2281, 3217, 4253, 4423, 9689, 9941)
- 序列差分表: 差分表的第 0 条对角线确定原序列. 设原序列为  $h_i$ , 第 0 条对角线为  $c_0, c_1, \dots, c_p, 0, 0, \dots$  有这样两个公式:  $h_n = \binom{n}{0}c_0 + \binom{n}{1}c_1 + \dots + \binom{n}{p}c_p, \sum_{k=0}^n h_k = \binom{n+1}{1}c_0 + \binom{n+1}{2}c_2 + \dots + \binom{n+1}{p+1}c_p$
- GCD:  $\gcd(2^a - 1, 2^b - 1) = 2^{\gcd(a,b)} - 1$
- Fermat 分解算法: 从  $t = \sqrt{n}$  开始, 依次检查  $t^2 - n, (t+1)^2 - n, (t+2)^2 - n, \dots$ , 直到出现一个平方数  $y$ , 由于  $t^2 - y^2 = n$ , 因此分解得  $n = (t-y)(t+y)$ . 显然, 当两个因数很接近时这个方法能很快找到结果, 但如果遇到一个素数, 则需要检查  $\frac{n+1}{2} - \sqrt{n}$  个整数
- 牛顿迭代:  $x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}$
- 球与盒子的动人故事: ( $n$  个球,  $m$  个盒子,  $S$  为第二类斯特林数)
  - 1. 球同, 盒同, 无空: dp
  - 2. 球同, 盒同, 可空: dp
  - 3. 球同, 盒不同, 无空:  $\binom{n-1}{m-1}$

- 4. 球同, 盒不同, 可空:  $\binom{n+m-1}{n-1}$
- 5. 球不同, 盒同, 无空:  $S(n, m)$
- 6. 球不同, 盒同, 可空:  $\sum_{k=1}^m S(n, k)$
- 7. 球不同, 盒不同, 无空:  $m!S(n, m)$
- 8. 球不同, 盒不同, 可空:  $m^n$

- 组合数奇偶性: 若  $(n \& m) = m$ , 则  $\binom{n}{m}$  为奇数, 否则为偶数

- 格雷码  $G(x) = x \otimes (x \gg 1)$

- Fibonacci 数:
  - $F_0 = F_1 = 1, F_i = F_{i-1} + F_{i-2}, F_{-i} = (-1)^{i-1}F_i$
  - $F_i = \frac{1}{\sqrt{5}}((\frac{1+\sqrt{5}}{2})^n - (\frac{1-\sqrt{5}}{2})^n)$
  - $\gcd(F_n, F_m) = F_{\gcd(n,m)}$
  - $F_{i+1}F_i - F_i^2 = (-1)^i$
  - $F_{n+k} = F_kF_{n+1} + F_{k-1}F_n$

- 第一类 Stirling 数:  $[n_k]$  代表第一类无符号 Stirling 数, 代表将  $n$  阶置换群中有  $k$  个环的置换个数;  $s(n, k)$  代表有符号型,  $s(n, k) = (-1)^{n-k} [n_k]$ .

- $(x)^{(n)} = \sum_{k=0}^n [n_k]x^k, (x)_n = \sum_{k=0}^n s(n, k)x^k$
- $[n_k] = n[n_{k-1}^{n-1}] + [n_{k-1}^{n-1}], [0] = 1, [n] = [n_n^0] = 0$
- $[n_{n-2}^2] = \frac{1}{4}(3n-1)\binom{n}{3}, [n_{n-3}^3] = \binom{n}{2}\binom{n}{4}$
- $\sum_{k=0}^a [n_k] = n! - \sum_{k=0}^n [n_{k+a+1}^n]$
- $\sum_{p=k}^n [n_p^p] \binom{p}{k} = [n_{k+1}^{n+1}]$

- 第二类 Stirling 数:  $\{n_k\} = S(n, k)$  代表  $n$  个不同的球, 放到  $k$  个相同的盒子里, 盒子非空.

- $\{n_k\} = \frac{1}{k!} \sum_{j=0}^k (-1)^j \binom{k}{j} (k-j)^n$
- $\{n_k^{n+1}\} = k\{n_k^n\} + \{n_{k-1}^n\}, \{0_0^0\} = 1, \{0_n^0\} = \{n_n^0\} = 0$
- 奇偶性:  $(n-k) \& \frac{k-1}{2} == 0$

- Bell 数:  $B_n$  代表将  $n$  个元素划分成若干个非空集合的方案数

- $B_0 = B_1 = 1, B_n = \sum_{k=0}^{n-1} \binom{n-1}{k} B_k$
- $B_n = \sum_{k=0}^n \{n_k\}$
- Bell 三角形:  $a_{1,1} = 1, a_{n,1} = a_{n-1,n-1}, a_{n,m} = a_{n,m-1} + a_{n-1,m-1}, B_n = a_{n,1}$

- 对质数  $p$ ,  $B_{n+p} \equiv B_n + B_{n+1} \pmod p$
- 对质数  $p$ ,  $B_{n+p^m} \equiv mB_n + B_{n+1} \pmod p$
- 对质数  $p$ , 模的周期一定是  $\frac{p^p-1}{p-1}$  的约数,  $p \leq 101$  时就是这个值
- 从  $B_0$  开始, 前几项是 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, 115975...

• Bernoulli 数

- $B_0 = 1, B_1 = \frac{1}{2}, B_2 = \frac{1}{6}, B_4 = -\frac{1}{30}, B_6 = \frac{1}{42}, B_8 = B_4, B_{10} = \frac{5}{66}$
- $\sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}$
- $B_m = 1 - \sum_{k=0}^{m-1} \binom{m}{k} \frac{B_k}{m-k+1}$

- 完全数:  $x$  是偶完全数等价于  $x = 2^{n-1}(2^n - 1)$ , 且  $2^n - 1$  是质数.

6 其他

6.1 Extended LIS

```
1 int G[MAXN][MAXN];
2 void insertYoung(int v) {
3     for (int x = 1, y = INT_MAX; ; ++x) {
4         Down(y, *G[x]); while (y > 0 && G[x][y] >= v) --y;
5         if (++y > *G[x]) { ++*G[x]; G[x][y] = v; break; }
6         else swap(G[x][y], v);
7     }
8 }
9 int solve(int N, int seq[]) {
10     Rep(i, 1, N) *G[i] = 0;
11     Rep(i, 1, N) insertYoung(seq[i]);
12     printf("%d\n", *G[1] + *G[2]);
13     return 0;
14 }
```

6.2 生成 nCk

```
1 void nCk(int n, int k) {
2     for (int comb = (1 << k) - 1; comb < (1 << n); ) {
3         int x = comb & -comb, y = comb + x;
4         comb = (((comb & ~y) / x) >> 1) | y;
5     }
6 }
```

6.3 nextPermutation

```
1 boolean nextPermutation(int[] is) {
2     int n = is.length;
3     for (int i = n - 1; i > 0; i--) {
4         if (is[i - 1] < is[i]) {
5             int j = n; while (is[i - 1] >= is[--j]);
6             swap(is, i - 1, j); // swap is[i - 1], is[j]
7             rev(is, i, n); // reverse is[i, n]
8             return true;
9         }
10    } rev(is, 0, n);
11    return false;
12 }
```

6.4 Josephus 数与逆 Josephus 数

```
1 int josephus(int n, int m, int k) { int x = -1;
2     for (int i = n - k + 1; i <= n; i++) x = (x + m) % i; return x;
3 }
4 int invJosephus(int n, int m, int x) {
5     for (int i = n; ; i--) { if (x == i) return n - i; x = (x - m % i + i) % i; }
6 }
```

6.5 直线下的整点个数

求  $\sum_{i=0}^{n-1} \lfloor \frac{a+bi}{m} \rfloor$

```
1 LL count(LL n, LL a, LL b, LL m) {
2     if (b == 0) return n * (a / m);
3     if (a >= m) return n * (a / m) + count(n, a % m, b, m);
4     if (b >= m) return (n - 1) * n / 2 * (b / m) + count(n, a, b % m, m);
5     return count((a + b * n) / m, (a + b * n) % m, m, b);
6 }
```

6.6 Java 多项式

```
1 class Polynomial {
2     final static Polynomial ZERO = new Polynomial(new int[] { 0 });
3     final static Polynomial ONE = new Polynomial(new int[] { 1 });
4     final static Polynomial X = new Polynomial(new int[] { 0, 1 });
5     int[] coef;
6     static Polynomial valueOf(int val) { return new Polynomial(new int[] { val }); }
7     Polynomial(int[] coef) { this.coef = Arrays.copyOf(coef, coef.length); }
8     Polynomial add(Polynomial o, int mod); // omitted
9     Polynomial subtract(Polynomial o, int mod); // omitted
10    Polynomial multiply(Polynomial o, int mod); // omitted
11    Polynomial scale(int o, int mod); // omitted
12    public String toString() {
13        int n = coef.length; String ret = "";
14        for (int i = n - 1; i > 0; --i) if (coef[i] != 0)
15            ret += coef[i] + "x^" + i + "+";
16        return ret + coef[0];
17    }
18    static Polynomial lagrangeInterpolation(int[] x, int[] y, int mod) {
19        int n = x.length; Polynomial ret = Polynomial.ZERO;
20        for (int i = 0; i < n; ++i) {
21            Polynomial poly = Polynomial.valueOf(y[i]);
22            for (int j = 0; j < n; ++j) if (i != j) {
23                poly = poly.multiply(
24                    Polynomial.X.subtract(Polynomial.valueOf(x[j]), mod), mod);
25                poly = poly.scale(powMod(x[i] - x[j] + mod, mod - 2, mod), mod);
26            } ret = ret.add(poly, mod);
27        } return ret;
28    }
29 }
```

6.7 long long 乘法取模

```
1 LL multiplyMod(LL a, LL b, LL P) { // 需要保证 a 和 b 非负
2     LL t = (a * b - LL((long double)a / P * b + 1e-3) * P) % P;
3     return t < 0 : t + P : t;
4 }
```

6.8 重复覆盖

```

1 namespace DLX {
2     struct node { int x, y; node *l, *r, *u, *d; } base[MAX * MAX], *top, *head;
3     typedef node *link;
4     int row, col, nGE, ans, stamp, cntc[MAX], vis[MAX];
5     vector<link> eachRow[MAX], eachCol[MAX];
6     inline void addElement(int x, int y) {
7         top->x = x; top->y = y; top->l = top->r = top->u = top->d = NULL;
8         eachRow[x].push_back(top); eachCol[y].push_back(top++);
9     }
10    void init(int _row, int _col, int _nGE) {
11        row = _row; col = _col; nGE = _nGE; top = base; stamp = 0;
12        for (int i = 0; i <= col; ++i) vis[i] = 0;
13        for (int i = 0; i <= row; ++i) eachRow[i].clear();
14        for (int i = 0; i <= col; ++i) eachCol[i].clear();
15        for (int i = 0; i <= col; ++i) addElement(0, i);
16        head = eachCol[0].front();
17    }
18    void build() {
19        for (int i = 0; i <= row; ++i) {
20            vector<link> &v = eachRow[i];
21            sort(v.begin(), v.end(), cmpByY);
22            int s = v.size();
23            for (int j = 0; j < s; ++j) {
24                link l = v[j], r = v[(j + 1) % s]; l->r = r; r->l = l;
25            }
26        }
27        for (int i = 0; i <= col; ++i) {
28            vector<link> &v = eachCol[i];
29            sort(v.begin(), v.end(), cmpByX);
30            int s = v.size();
31            for (int j = 0; j < s; ++j) {
32                link u = v[j], d = v[(j + 1) % s]; u->d = d; d->u = u;
33            }
34        }
35        for (int i = 0; i <= col; ++i) cntc[i] = (int) eachCol[i].size() - 1;
36    }
37    void removeExact(link c) {
38        c->l->r = c->r; c->r->l = c->l;
39        for (link i = c->d; i != c; i = i->d)
40            for (link j = i->r; j != i; j = j->r) {
41                j->d->u = j->u; j->u->d = j->d; --cntc[j->y];
42            }
43    }
44    void resumeExact(link c) {
45        for (link i = c->u; i != c; i = i->u)
46            for (link j = i->l; j != i; j = j->l) {
47                j->d->u = j; j->u->d = j; ++cntc[j->y];
48            }
49        c->l->r = c; c->r->l = c;
50    }
51    void removeRepeat(link c) {
52        for (link i = c->d; i != c; i = i->d) {
53            i->l->r = i->r; i->r->l = i->l;
54        }
55    }
56    void resumeRepeat(link c) {
57        for (link i = c->u; i != c; i = i->u) {
58            i->l->r = i; i->r->l = i;
59        }
60    }
61    int calCH() {
62        int y, res = 0; ++stamp;
63        for (link c = head->r; (y = c->y) <= row && c != head; c = c->r) {
64            if (vis[y] != stamp) {
65                vis[y] = stamp; ++res;
66                for (link i = c->d; i != c; i = i->d)
67                    for (link j = i->r; j != i; j = j->r) vis[j->y] = stamp;
68            }
69        }
70        return res;
71    }
72    void DFS(int dep) { if (dep + calCH() >= ans) return;
73        if (head->r->y > nGE || head->r == head) {
74            if (ans > dep) ans = dep; return;
75            link c = NULL;
76            for (link i = head->r; i->y <= nGE && i != head; i = i->r)
77                if (!c || cntc[i->y] < cntc[c->y]) c = i;
78        }
79    }
80    int solve() { build(); ans = INF; DFS(0); return ans; }
81 }

```

```

76     for (link i = c->d; i != c; i = i->d) {
77         removeRepeat(i);
78         for (link j = i->r; j != i; j = j->r) if (j->y <= nGE) removeRepeat(j);
79         for (link j = i->r; j != i; j = j->r) if (j->y > nGE) removeExact(base + j->y);
80         DFS(dep + 1);
81         for (link j = i->l; j != i; j = j->l) if (j->y > nGE) resumeExact(base + j->y);
82         for (link j = i->l; j != i; j = j->l) if (j->y <= nGE) resumeRepeat(j);
83         resumeRepeat(i);
84     }
85 }
86 int solve() { build(); ans = INF; DFS(0); return ans; }
87 }

```

## 6.9 星期几判定

```

1 int getDay(int y, int m, int d) {
2     if (m <= 2) m += 12, y--;
3     if (y < 1752 || (y == 1752 && m < 9) || (y == 1752 && m == 9 && d < 3))
4         return (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 + 5) % 7 + 1;
5     return (d + 2 * m + 3 * (m + 1) / 5 + y + y / 4 - y / 100 + y / 400) % 7 + 1;
6 }

```

## 6.10 LCSequence Fast

```

1 ULL *a, *b, *s, c, d;
2 for (i = 0, a = appear[(int)B[k]], b = row[max(k - 1, 0)], s = X; i < bitSetLen; ++i)
3     *s++ = *a++ | *b++; // `X = row[i - 1] or appear[ B[i] ]`
4 for (i = 0, a = dp, c = d = 0; i < bitSetLen; ++a, c = d, ++i)
5     d = *a >> 63, *a = ~((*a << 1) + c); // `row[i] = -(row[i] << 1) + 1`
6 for (i = 0, a = dp, b = X, c = 0; i < bitSetLen; ++a, ++b, ++i)
7     d = *b + c, c = (*a >= -d) * a += d; // `row[i] += X`
8 for (i = 0, a = dp, b = X; i < bitSetLen; ++a, ++b, ++i)
9     *a = (*a ^ *b) & *b; // `row[i] = X and (row[i] xor X)`

```

## 6.11 C Split

```

1 for (char *tok = strtok(ins, delimiters); tok; tok = strtok(NULL, delimiters))
2     puts(tok); // '会破坏原字符串ins'

```

## 6.12 builtin 系列

- int \_\_builtin\_ffs (unsigned int x) 返回 x 的最后一位 1 的是从后向前第几位, 比如 7368( 1110011001000) 返回 4.
- int \_\_builtin\_clz (unsigned int x) 返回前导的 0 的个数.
- int \_\_builtin\_ctz (unsigned int x) 返回后面的 0 的个数, 和 \_\_builtin\_clz 相对.
- int \_\_builtin\_popcount (unsigned int x) 返回二进制表示中 1 的个数.
- int \_\_builtin\_parity (unsigned int x) 返回 x 的奇偶校验位, 也就是 x 的 1 的个数模 2 的结果.

## 7 Templates

### 7.1 Eclipse 配置

Exec=env UBUNTU\_MENUPROXY= /opt/eclipse/eclipse  
 preference general keys 把 word completion 设置成 alt+c, 把 content assistant 设置成 alt + /

## 7.2 C++

```
1 #pragma comment(linker, "/STACK:10240000")
2 #include <stdio>
3 #include <stdlib>
4 #include <string>
5 #include <iostream>
6 #include <algorithm>
7 #define Rep(i, a, b) for(int i = (a); i <= (b); ++i)
8 #define Foru(i, a, b) for(int i = (a); i < (b); ++i)
9 using namespace std;
10 typedef long long LL;
11 typedef pair<int, int> pii;
12 namespace BufferedReader {
13     char buff[MAX_BUFFER + 5], *ptr = buff, c; bool flag;
14     bool nextChar(char &c) {
15         if ( (c = *ptr++) == 0 ) {
16             int tmp = fread(buff, 1, MAX_BUFFER, stdin);
17             buff[tmp] = 0; if (tmp == 0) return false;
18             ptr = buff; c = *ptr++;
19         } return true;
20     }
21     bool nextUnsignedInt(unsigned int &x) {
22         for (;;) { if (!nextChar(c)) return false; if ('0' <= c && c <= '9') break; }
23         for (x = c - '0'; nextChar(c); x = x * 10 + c - '0') if (c < '0' || c > '9') break;
24         return true;
25     }
26     bool nextInt(int &x) {
27         for (;;) { if (!nextChar(c)) return false; if (c == '-' || ('0' <= c && c <= '9')) break; }
28         for ((c == '-') ? (x = 0, flag = true) : (x = c - '0', flag = false); nextChar(c); x = x * 10 + c - '0')
29             if (c < '0' || c > '9') break;
30         if (flag) x = -x; return true;
31     }
32 };
33 #endif
```

## 7.3 Java

```
1 import java.io.*;
2 import java.util.*;
3 import java.math.*;
4
5 public class Main {
6     public void solve() {}
7     public void run() {
8         tokenizer = null; out = new PrintWriter(System.out);
9         in = new BufferedReader(new InputStreamReader(System.in));
10        solve();
11        out.close();
12    }
13    public static void main(String[] args) {
14        new Main().run();
15    }
16    public StringTokenizer tokenizer;
17    public BufferedReader in;
18    public PrintWriter out;
19    public String next() {
20        while (tokenizer == null || !tokenizer.hasMoreTokens()) {
21            try { tokenizer = new StringTokenizer(in.readLine()); }
22            catch (IOException e) { throw new RuntimeException(e); }
23        } return tokenizer.nextToken();
24    }
25 }
```

## 7.4 gcc 配置

在.bashrc 中加入 export CXXFLAGS="-Wall -Wconversion -Wextra -g3"