# ECON 712B - Lecture Notes

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## Lecture 1 - HH Optimization over Finite Time Horizon

• Consumption-savings problem for  $T < \infty$ .

#### **Preferences**

$$\sum_{t=1}^{T} \beta^t u(c_t) \text{ where } 0 < \beta < 1$$

- Assumptions about preferences: u is strictly increasing, strictly concave, and continuously differentiable  $(C^2)$ .
- We also often assume that utility is bounded  $(|u(c)| < K \forall c)$  and Inada conditions hold  $(\lim_{c\to 0} u'(c) = +\infty \text{ and } \lim_{c\to +\infty} u'(c) = 0)$ .
- Think log utility.

#### Constraints

- R is gross interest rate (R = 1 + r where r is net interest rate).
- $\{y_t\}_{t=0}^T$  is labor income.
- $x_0$  is initial wealth.
- Intertemporal budget constraint:

$$\sum_{t=0}^{T} \left(\frac{1}{R}\right)^{t} c_{t} \leq \sum_{t=0}^{T} \left(\frac{1}{R}\right)^{t} y_{t} + x_{0}$$

- $s_t$  is savings.
- Flow budget constraint implies law of motion for wealth:

$$s_t = x_t - c_t + y_t \implies x_{t+1} = Rs_t = R(x_t - c_t + y_t)$$

• Intertemporal budget constraint embeds assumption that HH can borrow any amount at rate r as long as budget constraint is satisfied.

## **HH Problem**

$$\max_{\{c_t\}_{t=0}^T} \sum_{t=1}^T \beta^t u(c_t) \text{ s.t. } \sum_{t=0}^T \left(\frac{1}{R}\right)^t c_t \le \sum_{t=0}^T \left(\frac{1}{R}\right)^t y_t + x_0$$

- Where  $x_0 \ge 0$  and  $y_t \ge 0 \forall t$ .
- Since  $u(c_t)$  continuous,  $U_T(c) = \sum_{t=1}^T \beta^t u(c_t)$  is continuous on  $c = \{c_0, ..., c_t\}$ .

• Given  $\{y_t\}_{t=0}^T$ , the feasible set (below) is a compact subset of  $\mathbb{R}^T$  by Heine-Borel.

$$0 \le \sum_{t=0}^{T} \left(\frac{1}{R}\right)^{t} c_{t} \le \sum_{t=0}^{T} \left(\frac{1}{R}\right)^{t} y_{t} + x_{0}$$

- Existence: Close, bounded subset of  $\mathbb{R}^T_+$  are compact  $\implies \exists c^*$  that maximizes U(c) subject to budget constraint (Weinstrass).
- Uniqueness: Since the feasible set is convex and U(c) is strictly concave (because  $u(c^*)$  is strictly concave),  $\exists$  unique  $c^*$  that maximizes U(c) subject to budget constraint.
- The Lagrangian is:

$$\mathcal{L} = \sum_{t=1}^{T} \beta^t u(c_t) + \lambda \left[ \sum_{t=0}^{T} \left( \frac{1}{R} \right)^t (y_t - c_t) + x_0 \right]$$

## **Recursive Formulation**

- We use backward induction.
- At T, choose  $x_T$  given  $\max u(c_T)$  s.t.  $c_T \le x_T + y_T$ .  $x_{T+1} = 0 \implies c_T = x_T + y_T$ . So,  $V_T(x) = u(x + y_T)$ .
- At T-1, choose  $x_{T_1}$ . From law of motion,  $x_T = Rs_{T-1} = R(x_{T-1} c_{T-1} + y_{T-1})$ .

$$\max_{c_{T-1}} \left\{ u(c_{T-1}) + \beta u(R(x_{T-1} - c_{T-1} + y_{T-1}) + y_T) \right\} \text{ s.t. } 0 \le c_{T-1} \le \min\{\bar{c}(x) : c_T = 0\}$$

• For simplicity,  $y_t = y > 0 \ \forall t$ , so  $\bar{c}(x) = \frac{R(x+y)+y}{R}$ . (HH cannot borrow so much in period T-1 that they can't pay it all back in period T).

$$V_{T-1}(x) = \max_{c} \{u(c) + \beta V_T(R(x-c+y))\}$$
 s.t.  $0 \le c \le \bar{c}(x)$ 

- c(x) is continuous.
- Policy function  $\{c_t\} \implies c_t(x)$ .
- Thus, the Bellman equation is

$$V_t(x) = \max_{0 \le c \le \bar{c}} \{ u(c) + \beta V_{t+1} (R(x - c + y)) \}$$

• A Bellman equation is the key result of the theory of dynamic programming.

## Theorem of the Maximum

- Lower hemi-continuous (lhc):  $\forall y \in \Gamma(x), \{x_n\} \to x, \exists \{y_n\} \text{ s.t. } y_n \to y \text{ and } y_n \in \Gamma(x_n).$
- Upper hemi-continuous (uhc):  $\{x_n\} \to x$ , all  $\{y_n\}$  s.t.  $y_n \in \Gamma(x), y \in \Gamma(x)$ .

## Suppose:

- $x \in X \subseteq \mathbb{R}^L$
- $y \in Y \subset \mathbb{R}^M$
- $f: X \times Y \to \mathbb{R}$  is continuous.

•  $\Gamma: X \rightrightarrows Y$  compact-valued and continuous (uhc and lhc).

Then:

- $h(x) = \max_{y \in \Gamma(x)} f(x, y)$  is continuous.
- $G(x) = \{y \in \Gamma(x) : f(x,y) = h(x)\}$  is nonempty, compact-valued, and uhc.
- In the HH problem, the theorem of the maximum implies that  $V_T(x)$  is continuous.

## Principle of Optimality

- Let  $\{c_t^*\}_{t=0}^T$  solve sequence problem above. with initial  $x_0 \ge 0$ . From law of motion  $x_{t+1} = R(x_t c_t + y_t)$ , we can derive  $\{x_t^*\}$ .
- For arbitrary dates  $0 \le a < b \le T 1$ , let  $x_a^*$  and  $x_{b+1}^*$  be optimal states at those dates. Then the solution to the subproblem is still  $\{c_t^*\}_a^b$ :

$$\max_{\{c_t\}_{t=a}^b} \sum_{t=a}^b \beta^{t-1} u(c_t) \text{ s.t. } x_{t+1} = R(x_t - c_t + y_t) \text{ and given } x_a^*, x_{b+1}^*$$

## Back to Consumer Probem - FOC and Euler Equations

• In period T-1,

$$V_{T-1}(x) = \max_{0 < c < \bar{c}} \{ u(c) + \beta V_T (R(x - c + y)) \}$$

- Inada conditions rule out corner solutions  $\implies c^*$  is interior.
- Strictly concave, differentiable function and convex choice set  $\implies$  FOC determine optimal choice.
- FOC:  $u'(c) = \beta RV'_T(R(x-c+y)).$
- Because  $V_T(x) = u(x+y) \implies V_T'(x) = u'(x+y)$ , FOC  $\implies u'(c) = \beta R u'(R(x-c+y)+y)$  (Euler Equation).
- Using policy function  $c_{t-1} = c(x), u'(c_{T-1}) = \beta Ru'(c_T) \implies$

$$u'(c_t) = \beta R u'(c_{t+1})$$

- Since u is strictly concave  $(u'' < 0) \implies$  consumption smoothing.
- $\beta R = 1 \implies u'(c_t) = u'(c_{t+1}) \implies c_t = \bar{c}$ .
- $\beta R < 1 \implies u'(c_t) < u'(c_{t+1}) \implies c_t > c_{t+1}$ .
- $\beta R > 1 \implies u'(c_t) > u'(c_{t+1}) \implies c_t < c_{t+1}$ .

## Lecture 2 - HH Optimization over Infinity Time Horizon

We'll consider the consumption-saving problem from last class but over an infinite horizon:

$$\max_{\{c_t\}_{t=0}^{\infty}} \sum_{t=0}^{\infty} \beta^t u(c_t) \text{ s.t. } x_{t+1} = R(x_t - c_t + y), c_t \ge 0, x_0 \text{ given.}$$

 $c_0, c_1, c_2, \dots$ 

#### Recursive Formulation of Consumption-Savings Problem

$$V(x_0) = \max_{\{c_t\}_{t=0}^{\infty}} \sum_{t=0}^{\infty} \beta^t u(c_t) \text{ s.t. } x_{t+1} = R(x_t - c_t + y)$$

- Since u is bounded and  $0 < \beta < 1 \implies |V(x_0)| < \infty$ .
- These notes is all about justifying that we can break apart the maximum and apply it to today's value and the rest of time so that the consumer problem becomes:

$$V(x_0) = \max_{\{c_t\}_{t=1}^{\infty}} \left\{ u(c_0) + \sum_{t=0}^{\infty} \beta^t u(c_t) \right\}$$

$$= \max_{c_0} \left\{ u(c_0) + \beta \max_{\{c_t\}_{t=1}^{\infty}} \sum_{s=0}^{\infty} \beta^s u(c_{s+1}) \right\}$$

$$= \max_{c_0} \{ u(c_0) + \beta V(x_1) \}$$
s.t.  $x_1 = R(x_0 + y - c_0)$ 

• Under what conditions, can we find a solution to the sequence problem by using this recursive formulation?

## Sequence Problem (SP)

- Consider more general notation following Stokey, Lucas, and Prescott.
- Let F be the objective function that we're maximizing.  $x_t$  is savings; it is the state variable. We can reformulate F to be the utility function on consumption from above.

$$V^*(x_0) = \sup_{\{x_{t+1}\}} \sum_{t=0}^{\infty} \beta^t F(x_t, x_{t+1}) \text{ s.t. } x_{t+1} \in \Gamma(x_t), x_0 \text{ given.}$$

- $\Gamma(x_t)$  is the feasible correspondence.
- $x_t \in X$
- $A = \operatorname{graph}\Gamma = \{(x, y) \in X \times X : y \in \Gamma(x)\}\$
- $F:A\to\mathbb{R}$
- $\Pi(x_0) = \{\{x_t\}_{t=0}^{\infty} : x_{t+1} \in \Gamma(x_t)\}$  is the set of feasible plans.

#### **Recursive Formulation**

- The Functional Equation (FE) (aka the Bellman equation) is  $V(x) = \sup_{y \in \Gamma(x)} \{F(x,y) + \beta V(y)\}.$
- Today's question is when is  $V^*$  equal to V. If so, we've made forward progress because solving for a function is easier than solving for an infinite series.
- What do we know about  $V^*$  and V?
- To make things easier, assume  $|V^*(x_0)| < \infty$  and  $|V(x_0)| < \infty$ .
- Define  $U(x) := \sum_{t=0}^{\infty} \beta^t F(x_t, x_{t+1})$  and  $U(x) := \sum_{t=0}^n \beta^t F(x_t, x_{t+1})$  for some  $x \in \Pi(x_0)$ .
- Thus, we know that the solution to the SP  $V^*(x_0) \geq U(x) \ \forall x \in \Pi(x_0)$ .
- In addition, we know that for any  $\varepsilon > 0$ , the solution to the SP  $V^*(x_0) \leq U(x) + \varepsilon$  for some  $x \in \Pi(x_0)$ .
- Similarly, we know that the FE  $V(x_0) \ge F(x_0, y) + \beta V(y) \ \forall y \in \Gamma(x_0) \ \text{and} \ V(x_0) \le F(x_0, y) + \beta V(y) + \varepsilon$  for some  $y \in \Gamma(x_0)$ .

#### Theorem:

- Suppose  $\Gamma(x)$  is nonempty  $\forall x \in X$ .
- Suppose  $\forall x_0 \in X$  and  $x \in \Pi(x_0)$ ,  $\lim_{n \to \infty} \sum_{t=0}^n \beta^t F(x_t, x_{t+1})$  exists.
- Then if V solves (FE) and
- $\lim_{n\to\infty} \beta^n V(x_n) = 0 \ \forall x \in \Pi(x_0), \ \forall x_0 \ (a tail boundedness restriction)$
- Then  $V = V^*$  (i.e., V solves the sequence problem).

#### Proof:

- Suppose  $|V| < \infty$  and  $|V^*| < \infty$  (not necessarily, but convienent).
- Say we have a solution V of the Bellman equation (FE).
- By definition, we know that  $\forall x \in \Pi(x_0)$ :

$$V(x_0) \ge F(x_0, x_1) + \beta V(x_1)$$

$$\ge F(x_0, x_1) + \beta F(x_1, x_2) + \beta^2 V(x_2)$$
...
$$\ge U_n(x) + \beta^{n+1} V(x_{n+1})$$

• Take  $\lim_{n\to\infty}$ :

$$V(x_0) \ge U(x) = \sum_{t=0}^{\infty} \beta^t F(x_t, x_{t+1}) \ \forall x \in \Pi(x_0)$$

- [We're halfway done; we showed that V dominates for all feasible plans, and we now need to show that we can get within  $\varepsilon$  away with some feasible plan.]
- Having a solution to the FE also implies that  $\exists x_1 \in \Gamma(x_0), x_2 \in \Gamma(x_1), \dots$  such that

$$V(x_t) \leq F(x_t, x_{t+1}) + V(x_{t+1}) + \delta_{t+1} \ \forall t$$

• Where  $\{\delta_t\}$  s.t.  $\sum_{t=0}^{\infty} \beta^t \delta_t \leq \varepsilon/2$ .

$$\implies V(x_0) \le F(x_0, x_1) + V(x_1) + \delta_1$$
...
$$\le \sum_{t=0}^{n} \beta^t F(x_t, x_{t+1}) + \beta^{n+1} V(x_{n+1}) + \sum_{t=1}^{n+1} \beta^{t-1} \delta_t$$

- For sufficiently large n, we can bound  $\beta^{n+1}V(x_{n+1})$  by  $\varepsilon/2$ .
- By construction  $\sum_{t=1}^{n+1} \beta^{t-1} \delta_t \leq \sum_{t=0}^{\infty} \beta^t \delta_t \leq \varepsilon/2$ .
- Taking the limit as  $n \to \infty$ :

$$V(x_0) \le \sum_{t=0}^{\infty} \beta^t F(x_t, x_{t+1}) + \varepsilon$$

## **Background on Contractions**

• An  $T: S \to S$  on  $(S, \mu)$  metric space is a contraction if  $\mu(Tx, Ty) \leq \beta \mu(x, y) \ \forall x, y \in S$ .

Contraction Mapping Theorem:

- If  $(S, \mu)$  is complete and T is a contraction, then T has a unique fixed point, Tv = v.
- And for any  $v_0 \in S$ ,  $\mu(T^n v_0, v) \leq \beta^n \mu(v_0, v) \iff \lim_{n \to \infty} T^n v_0 = v$ .

### Corollary:

- Suppose  $S' \subseteq S$ , S' closed.
- Then  $T(S') \subseteq S' \implies v \in S'$ .
- Moreover, if  $T(S') \subseteq S'' \subset S'$ , then  $v \in S''$ .
- The last part of this corollary is helpful if we're thinking about strictly increasing functions because we need weak inequalities to get a closed subset.

## Blackwell Sufficient Conditions:

- $X \subseteq \mathbb{R}^n$
- B(X) is the set of bounded functions  $f: X \to \mathbb{R}$ .
- Let  $||f|| = \sup_{x \in X} |f(x)| \implies \mu(f,g) = ||f-g||$ . Thus, this metric space is complete.
- Suppose  $T: B(X) \to B(X)$  satisfies monotonicity and discounting.
  - Monotonicity:  $f, g \in B(X), f(x) \leq g(x) \ \forall x \in X$ , then  $Tf(x) \leq Tg(x) \ \forall x \in X$ .
  - Discounting:  $\exists \beta \in (0,1)$  s.t.  $T(f+a)(x) \leq Tf(x) + \beta a$  where a > 0 and  $f \in B(X)$ .
- Then T is a contraction with modulus  $\beta$ .

## Applying Blackwell to Bellman Equations

- Recall we have only made some weak assumptions (nonempty and tail boundedness). Now, we're added a few more assumptions about the feasible set  $(\Gamma 1)$  and the objective function (F1):
- ( $\Gamma$ 1):  $X \in \mathbb{R}^{\ell}$  is convex.  $\Gamma: X \to X$  nonempty, compact-valued, and continuous (uhc and lhc). [These assumptions are similar to those in the theorem of the maximum.]
- (F1):  $F: A \to \mathbb{R}$  is bounded and continuous with  $0 < \beta < 1$  (recall  $A = \operatorname{graph}\Gamma$ ).
- Define C(X) as the metric space of bounded continuous functions on X with  $||f|| = \sup_{x \in X} |f(x)|$ . [This metric space is complete.]
- Define the Bellman operator as  $(Tf)(x) = \max_{y \in \Gamma(x)} \{F(x,y) + \beta f(y)\}$  for  $f \in C(X)$ .
- With T defined this way, the Bellman equation is the fixed point of the Bellman operator: Tv = v.

## Theorem:

- Under  $(\Gamma 1)$  and (F 1).
- The Bellman operator  $T: C(X) \to C(X)$  is a contraction, hence T has a unique fixed point (i.e.,  $v \in C(X)$  s.t. v = Tv) and  $v_0 \in C(X)$ ,  $||T^n v_0 v|| \le \beta^n ||v_0 v||$ .
- Moreover, the optimal policy correspondence  $G(x) = \{y \in \Gamma(x) : V(x) = F(x,y) + \beta V(y)\}$  is compact-valued and uhc (conclusions from the Theorem of the Maximum).

## Proof:

- The Bellman operator is in fact an operator.  $T: C \to C$  by the Theorem of the Maximum.
- $f, g \in C(X), f \leq g$ .
- T is monotone:

$$(Tf)(x) = \max_{y \in \Gamma(x)} \{F(x,y) + \beta f(y)\} \leq \max_{y \in \Gamma(x)} \{F(x,y) + \beta g(y)\} = (Tg)(x)$$

• T discounts:

$$T(f(x)+a) = \max_{y \in \Gamma(x)} \left\{ F(x,y) + \beta(f(y)+a) \right\} = \max_{y \in \Gamma(x)} \left\{ F(x,y) + \beta f(y) \right\} + \beta a = Tf(x) + \beta a$$

• By Blackwell, T is a contraction. By CMT, the Bellman equation has a unique solution.