

Chapter 6

Relations

Discrete Structures for Computing

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Course outcomes

Course learning outcomes	
L.O.1	Understanding of logic and discrete structures L.O.1.1 – Describe definition of propositional and predicate logic L.O.1.2 – Define basic discrete structures: set, mapping, graphs
L.O.2	Represent and model practical problems with discrete structures L.O.2.1 – Logically describe some problems arising in Computing L.O.2.2 – Use proving methods: direct, contrapositive, induction L.O.2.3 – Explain problem modeling using discrete structures
L.O.3	Understanding of basic probability and random variables L.O.3.1 – Define basic probability theory L.O.3.2 – Explain discrete random variables
L.O.4	Compute quantities of discrete structures and probabilities L.O.4.1 – Operate (compute/ optimize) on discrete structures L.O.4.2 – Compute probabilities of various events, conditional ones, Bayes theorem

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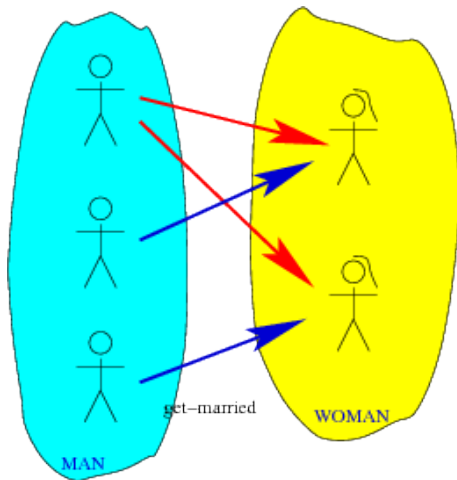
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Function?





Definition

Let A and B be sets. A **binary relation** (*quan hệ hai ngôi*) from a set A to a set B is a set

$$R \subseteq A \times B$$

- Notations:

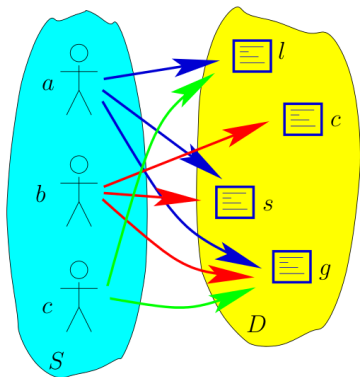
$$(a, b) \in R \longleftrightarrow aRb$$

- **n-ary relations?**

Example

Example

Let $A = \{a, b, c\}$ be the set of students, $B = \{l, c, s, d\}$ be the set of the available optional courses. We can have relation R that consists of pairs (a, b) , where a is a student enrolled in course b .



$$R = \{(a, l), (a, s), (a, g), (b, c), (b, s), (b, g), (c, l), (c, g)\}$$

R	l	c	s	g
a	x		x	x
b		x	x	x
c	x			x



Functions as Relations

- Is a function a relation?

- **Yes!**

- $f : A \rightarrow B$

$$R = \{(a, b) \mid b = f(a)\}$$



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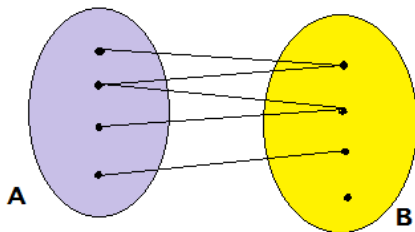
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Functions as Relations

- Is a relation a function?
- **No**



- Relations are a **generalization** of functions



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Relations on a Set

Definition

A **relation on the set A** is a relation from A to A .

Example

Let A be the set $\{1, 2, 3, 4\}$. Which ordered pairs are in the relation $R = \{(a, b) \mid a \text{ divides } b\}$ (a là ước số của b)?

Solution:

$$R = \{(1, 1), (1, 2), (1, 3), (1, 4), (2, 2), (2, 4), (3, 3), (4, 4)\}$$

R	1	2	3	4
1	x	x	x	x
2		x		x
3			x	
4				x



Properties of Relations



Reflexive (phản xạ)	$xRx, \forall x \in A$
Symmetric (đối xứng)	$xRy \rightarrow yRx, \forall x, y \in A$
Antisymmetric (phản đối xứng)	$(xRy \wedge yRx) \rightarrow x = y, \forall x, y \in A$
Transitive (bắc cầu)	$(xRy \wedge yRz) \rightarrow xRz, \forall x, y, z \in A$

Example

Example

Consider the following relations on $\{1, 2, 3, 4\}$:

$$R_1 = \{(1, 1), (1, 2), (2, 1), (2, 2), (3, 4), (4, 1), (4, 4)\},$$

$$R_2 = \{(1, 1), (1, 2), (2, 1)\},$$

$$R_3 = \{(1, 1), (1, 2), (1, 4), (2, 1), (2, 2), (3, 3), (4, 1), (4, 4)\},$$

$$R_4 = \{(2, 1), (3, 1), (3, 2), (4, 1), (4, 2), (4, 3)\},$$

$$R_5 = \{(3, 4)\}$$

Solution:

- Reflexive: R_3
- Symmetric: R_2, R_3
- Antisymmetric: R_4, R_5
- Transitive: R_4, R_5



Example

Example

What is the properties of the **divides** (ước số) relation on the set of positive integers?

Solution:

- $\forall a \in \mathbb{Z}^+, a \mid a$: **reflexive**
- $1 \mid 2$, but $2 \nmid 1$: **not symmetric**
- $\exists a, b \in \mathbb{Z}^+, (a \mid b) \wedge (b \mid a) \rightarrow a = b$: **antisymmetric**
- $a \mid b \Rightarrow \exists k \in \mathbb{Z}^+, b = ak; b \mid c \Rightarrow \exists l \in \mathbb{Z}^+, c = bl$. Hence, $c = a(kl) \Rightarrow a \mid c$: **transitive**



Example

Example

What are the properties of these relations on the set of integers:

$$R_1 = \{(a, b) \mid a \leq b\}$$

$$R_2 = \{(a, b) \mid a > b\}$$

$$R_3 = \{(a, b) \mid a = b \text{ or } a = -b\}$$



Combining Relations

Because relations from A to B are **subsets** of $A \times B$, two relations from A to B can be combined in any way two sets can be combined.

Example

Let $A = \{1, 2, 3\}$ and $B = \{1, 2, 3, 4\}$. List the combinations of relations $R_1 = \{(1, 1), (2, 2), (3, 3)\}$ and $R_2 = \{(1, 1), (1, 2), (1, 3), (1, 4)\}$.

Solution: $R_1 \cup R_2$, $R_1 \cap R_2$, $R_1 - R_2$ and $R_2 - R_1$.

Example

Let A and B be the set of all students and the set of all courses at school, respectively. Suppose $R_1 = \{(a, b) \mid a \text{ has taken the course } b\}$ and $R_2 = \{(a, b) \mid a \text{ requires course } b \text{ to graduate}\}$. What are the relations $R_1 \cup R_2$, $R_1 \cap R_2$, $R_1 \oplus R_2$, $R_1 - R_2$, $R_2 - R_1$?



Composition of Relations



Definition

Let R be **relations** from A to B and S be from B to C . Then the **composite** (*hợp thành*) of S and R is

$$S \circ R = \{(a, c) \in A \times C \mid \exists b \in B (aRb \wedge bSc)\}$$

Example

$$R = \{(0, 0), (0, 3), (1, 2), (0, 1)\}$$

$$S = \{(0, 0), (1, 0), (2, 1), (3, 1)\}$$

$$S \circ R = \{(0, 0), (0, 1), (1, 1)\}$$



Definition

Let R be a relation on the set A . The **powers** (*lũy thừa*) $R^n, n = 1, 2, 3, \dots$ are defined recursively by

$$R^1 = R \quad \text{and} \quad R^{n+1} = R^n \circ R.$$

Example

Let $R = \{(1, 1), (2, 1), (3, 2), (4, 3)\}$. Find the powers $R^n, n = 2, 3, 4, \dots$

Solution:

$$R^2 = \{(1, 1), (2, 1), (3, 1), (4, 2)\}$$

$$R^3 = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

$$R^4 = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

\dots

Representing Relations Using Matrices

Definition

Suppose R is a relation from $A = \{a_1, a_2, \dots, a_m\}$ to $B = \{b_1, b_2, \dots, b_n\}$, R can be represented by the **matrix** $\mathbf{M}_R = [m_{ij}]$, where

$$m_{ij} = \begin{cases} 1 & \text{if } (a_i, b_j) \in R \\ 0 & \text{if } (a_i, b_j) \notin R \end{cases}$$

Example

R is relation from $A = \{1, 2, 3\}$ to $B = \{1, 2\}$. Let $R = \{(2, 1), (3, 1), (3, 2)\}$, the matrix for R is

$$\mathbf{M}_R = \begin{bmatrix} 0 & 0 \\ 1 & 0 \\ 1 & 1 \end{bmatrix}$$

Determine whether the relation has certain properties (reflexive, symmetric, antisymmetric,...)



Representing Relations Using Digraphs



Definition

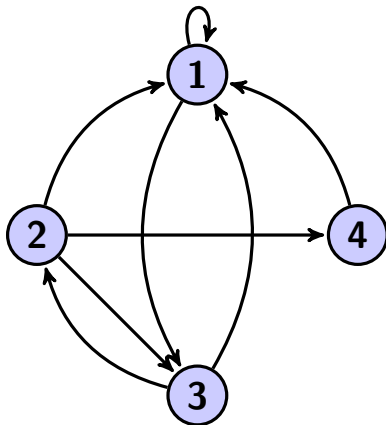
Suppose R is a relation in $A = \{a_1, a_2, \dots, a_m\}$, R can be represented by the **digraph** (đồ thị có hướng) $G = (V, E)$, where

$$V = A$$
$$(a_i, a_j) \in E \text{ if } (a_i, a_j) \in R$$

Example

Given a relation on $A = \{1, 2, 3, 4\}$,
 $R = \{(1, 1), (1, 3), (2, 1), (2, 3), (2, 4), (3, 1), (3, 2), (4, 1)\}$
Draw corresponding digraph.

Resulting digraph



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Definition

The **closure** (*bao đóng*) of relation R with respect to **property** P is the relation S that

- i. **contains** R
- ii. **has** property P
- iii. is **contained in any** relation satisfying (i) and (ii).

S is the “smallest” relation satisfying (i) & (ii)

Reflexive Closure



Example

Let $R = \{(a, b), (a, c), (b, d), (d, c)\}$

The **reflexive closure** of R

$\{(a, b), (a, c), (b, d), (d, c), (a, a), (b, b), (c, c), (d, d)\}$

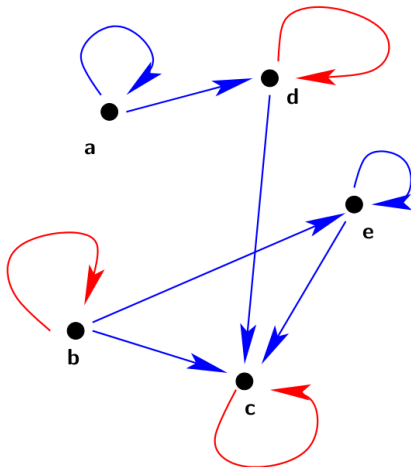
$$R \cup \Delta$$

where

$$\Delta = \{(a, a) \mid a \in A\}$$

diagonal relation (*quan hệ đường chéo*).

Reflexive Closure



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Symmetric Closure



Example

Let $R = \{(a, b), (a, c), (b, d), (c, a), (d, e)\}$

The **symmetric closure** of R

$\{(a, b), (a, c), (b, d), (c, a), (d, e), (b, a), (d, b), (e, d)\}$

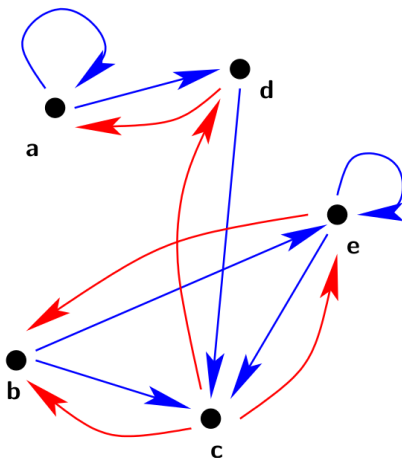
$$R \cup R^{-1}$$

where

$$R^{-1} = \{(b, a) \mid (a, b) \in R\}$$

inverse relation (*quan hệ ngược*).

Symmetric Closure



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Transitive Closure



Example

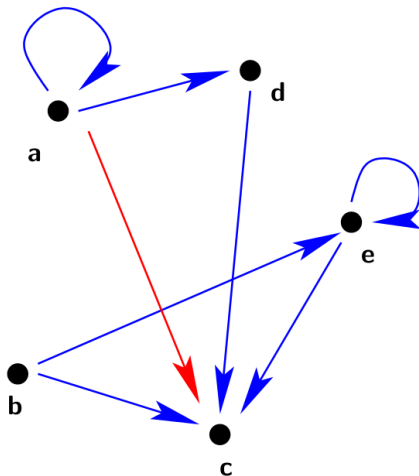
Let $R = \{(a, b), (a, c), (b, d), (d, e)\}$

The **transitive closure** of R

$\{(a, b), (a, c), (b, d), (d, e), (a, d), (b, e), (a, e)\}$

$$\bigcup_{n=1}^{\infty} R^n$$

Transitive Closure



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Equivalence Relations

Definition

A relation on a set A is called an **equivalence relation** (*quan hệ tương đương*) if it is **reflexive**, **symmetric** and **transitive**.

Example (1)

The relation $R = \{(a, b) | a \text{ and } b \text{ are in the same provinces}\}$ is an equivalence relation. a is **equivalent** to b and vice versa, denoted $a \sim b$.

Example (2)

$$R = \{(a, b) \mid a = b \vee a = -b\}$$

R is an equivalence relation.

Example (3)

$$R = \{(x, y) \mid |x - y| < 1\}$$

Is R an equivalence relation?



Example

Example (Congruence Modulo m - Đồng dư modulo m)

Let m be a positive integer with $m > 1$. Show that the relation

$$R = \{(a, b) \mid a \equiv b \pmod{m}\}$$

is an equivalence relation on the set of integers.

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Equivalence Classes

Definition

Let R be an **equivalence relation** on the set A . The set of all elements that are related to an element a of A is called the **equivalence class** (*lớp tương đương*) of a , denoted by

$$[a]_R = \{s \mid (a, s) \in R\}$$

Example

The equivalence class of “Thủ Đức” for the equivalence relation “in the same provinces” is { “Thủ Đức”, “Gò Vấp”, “Bình Thạnh”, “Quận 10”, ... }



Example

Example

What are the equivalence classes of 0, 1, 2, 3 for congruence modulo 4?

Solution:

$$[0]_4 = \{\dots, -8, -4, 0, 4, 8, \dots\}$$

$$[1]_4 = \{\dots, -7, -3, 1, 5, 9, \dots\}$$

$$[2]_4 = \{\dots, -6, -2, 2, 6, 10, \dots\}$$

$$[3]_4 = \{\dots, -5, -1, 3, 7, 11, \dots\}$$



Equivalence Relations and Partitions



Theorem

Let R be an equivalence relation on a set A . These statements for elements a and b of A are equivalent:

- i aRb
- ii $[a] = [b]$
- iii $[a] \cap [b] \neq \emptyset$

Example 1

Example

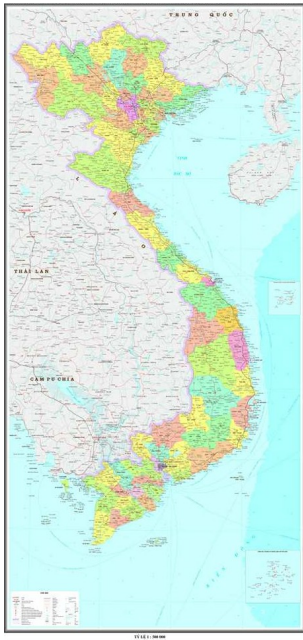
Suppose that $S = \{1, 2, 3, 4, 5, 6\}$. The collection of sets $A_1 = \{1, 2, 3\}$, $A_2 = \{4, 5\}$, and $A_3 = \{6\}$ forms a partition of S , because these sets are disjoint and their union is S .

The equivalence classes of an equivalence relation R on a set S form a **partition** of S .

Every partition of a set can be used to form an **equivalence relation**.



Example 2



Example

Divides set of all cities and towns in Vietnam into set of 64 provinces. We know that:

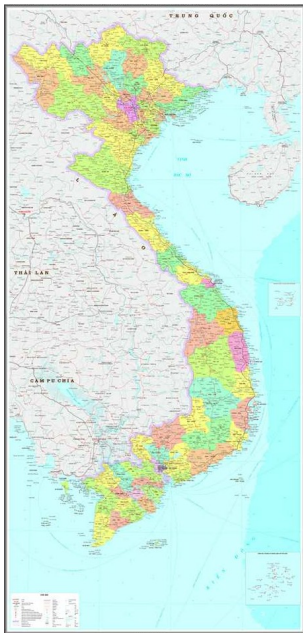
- there are no provinces with no cities or towns
- no city is in more than one province
- every city is accounted for

Definition

A **partition** of a Vietnam is a collection of non-overlapping non-empty subsets of Vietnam (provinces) that, together, make up all of Vietnam.



Relation in a Partition



- We divided based on relation

$$R = \{(a, b) | a \text{ and } b \text{ are in the same provinces}\}$$

- “Thủ Đức” is related (equivalent) to “Gò Vấp”
- “Đà Lạt” is **not** related (not equivalent) to “Long Xuyên”



Partial Order Relations

- Order words such that x comes before y in the dictionary
- Schedule projects such that x must be completed before y
- Order set of integers, where $x < y$

Definition

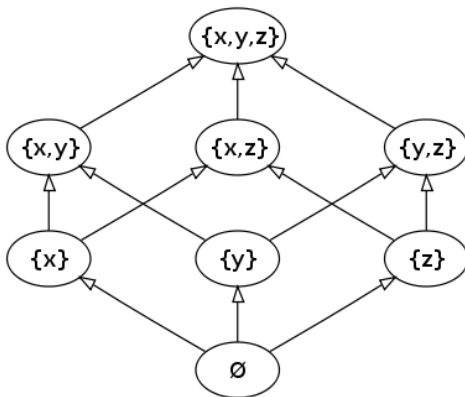
A relation R on a set S is called a **partial ordering** (có thứ tự bộ phận) if it is **reflexive**, **antisymmetric** and **transitive**. A set S together with a partial ordering R is called a partially ordered set, or **poset** (tập có thứ tự bộ phận), and is denoted by (S, R) or (S, \preceq) .

Example

- (\mathbb{Z}, \geq) is a poset
- Let S a set, $(P(S), \subseteq)$ is a poset



Example



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Totally Order Relations

Example

In the poset $(\mathbb{Z}^+, |)$, 3 and 9 are **comparable** (so sánh được), because $3 \mid 9$, but 5 and 7 are not, because $5 \nmid 7$ and $7 \nmid 5$.

→ That's why we call it **partially** ordering.

Definition

If (S, \preceq) is a poset and every two elements of S are comparable, S is called a **totally ordered** (có thứ tự toàn phần). A totally ordered set is also called a **chain** (dây xích).

Example

The poset (\mathbb{Z}, \leq) is totally ordered.



Maximal & Minimal Elements

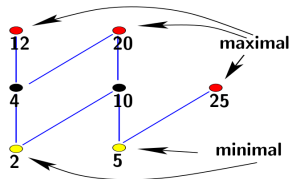


Definition

- a is **maximal** (*cực đại*) in the poset (S, \preceq) if there is no $b \in S$ such that $a \prec b$.
- a is **minimal** (*cực tiểu*) in the poset (S, \preceq) if there is no $b \in S$ such that $b \prec a$.

Example

Which elements of the poset $(\{2, 4, 5, 10, 12, 20, 25\}, |)$ are minimal and maximal?



Greatest Element & Least Element



Definition

- a is the **greatest element** (*lớn nhất*) of the poset (S, \preceq) if $b \preceq a$ for all $b \in S$.
- a is the **least element** (*nhỏ nhất*) of the poset (S, \preceq) if $a \preceq b$ for all $b \in S$.

The greatest and least element are **unique** if it exists.

Example

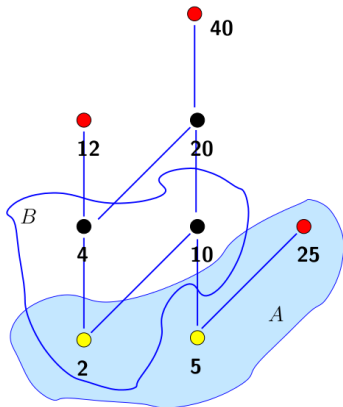
Let S be a set. In the poset $(P(S), \subseteq)$, the least element is \emptyset and the greatest element is S .

Upper Bound & Lower Bound

Definition

Let $A \subseteq (S, \preccurlyeq)$.

- If u is an element of S such that $a \preccurlyeq u$ for all elements $a \in A$, then u is called an **upper bound** (*cận trên*) of A .
- If l is an element of S such that $l \preccurlyeq a$ for all elements $a \in A$, then l is called a **lower bound** (*cận dưới*) of A .



Example

- Subset A does **not** have upper bound and lower bound.
- The upper bound of B are 20, 40 and the lower bound is 2.

