Chapter 10
Trees

Discrete Structures for Computing on January 11, 2017

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Contents

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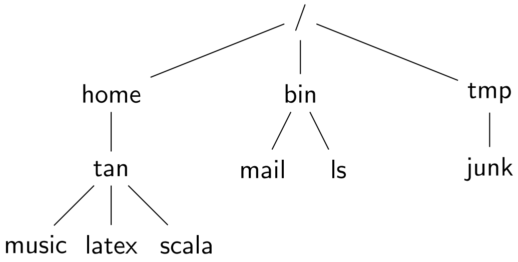
Course outcomes

	Course learning outcomes
L.O.1	Understanding of logic and discrete structures
	L.O.1.1 – Describe definition of propositional and predicate logic L.O.1.2 – Define basic discrete structures: set, mapping, graphs
L.O.2	Represent and model practical problems with discrete structures
	L.O.2.1 – Logically describe some problems arising in Computing L.O.2.2 – Use proving methods: direct, contrapositive, induction L.O.2.3 – Explain problem modeling using discrete structures
L.O.3	Understanding of basic probability and random variables
	L.O.3.1 – Define basic probability theory L.O.3.2 – Explain discrete random variables
L.O.4	Compute quantities of discrete structures and probabilities
	L.O.4.1 – Operate (compute/ optimize) on discrete structures L.O.4.2 – Compute probabilities of various events, conditional ones, Bayes theorem

Contents

Introduction

- Very useful in computer science: search algorithm, game winning strategy, decision making, sorting, ...
- Other disciplines: chemical compounds, family trees, organizational tree, ...



Tree

Definition

A tree (*cây*) is a connected undirected graph with no simple circuits. Consequently, a tree must be a simple graph.

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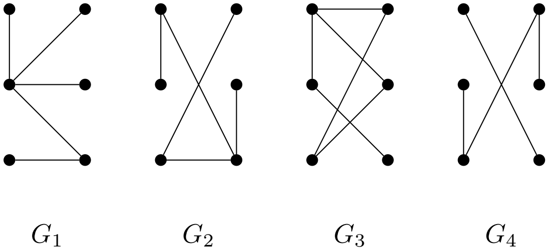

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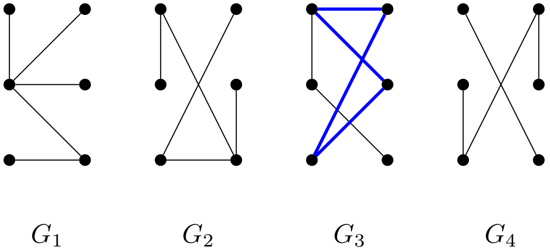

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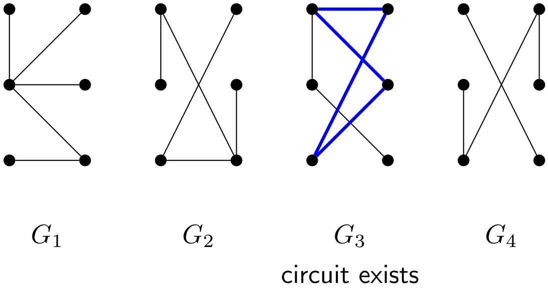

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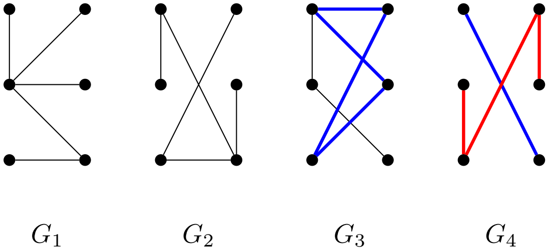

circuit exists

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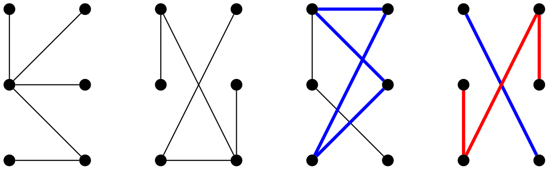

G_1 G_2 G_3 G_4

circuit exists

not connected

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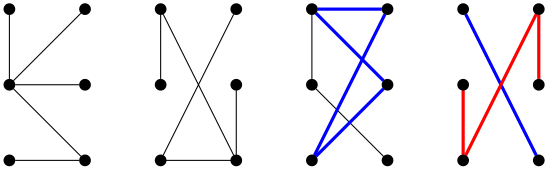

G_1 G_2 G_3 G_4

circuit exists

not connected

Definition

Graphs containing no simple circuits that are not necessarily connected is forest (*rừng*), in which each connected component is a tree.

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Rooted Trees

Definition

A rooted tree (*cây có gốc*) is a tree in which:

- One vertex has been designated as the root and
- Every edge is directed away from the root

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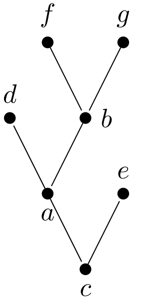

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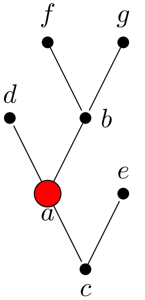

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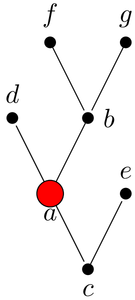

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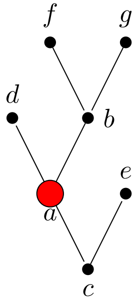

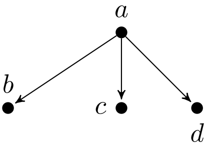

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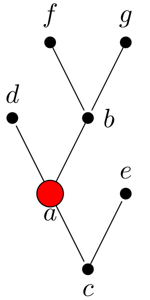

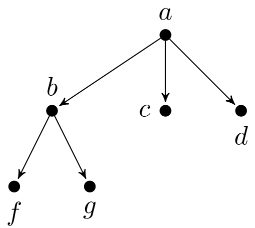

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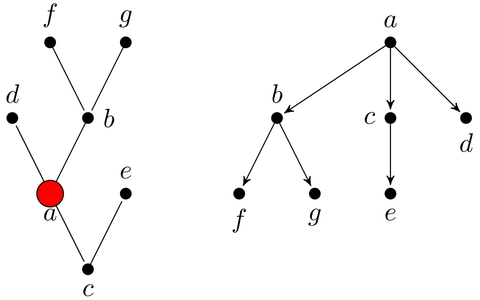

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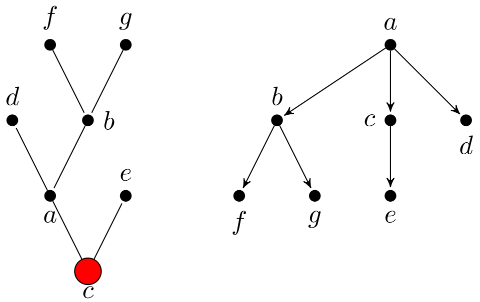

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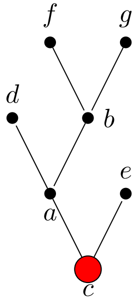

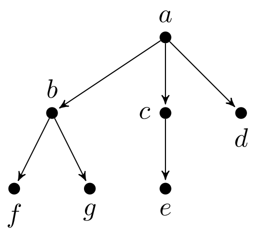

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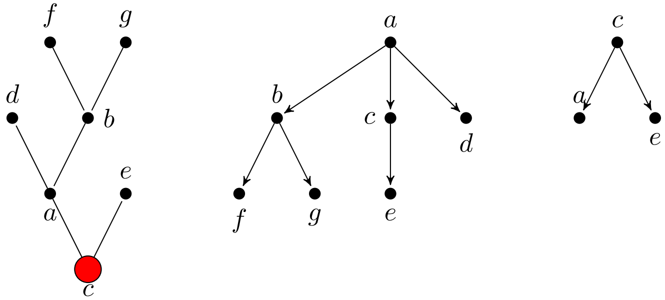

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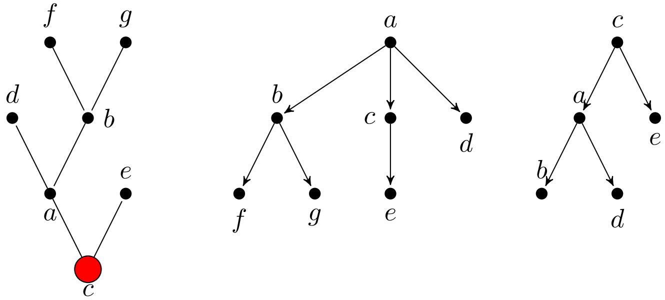

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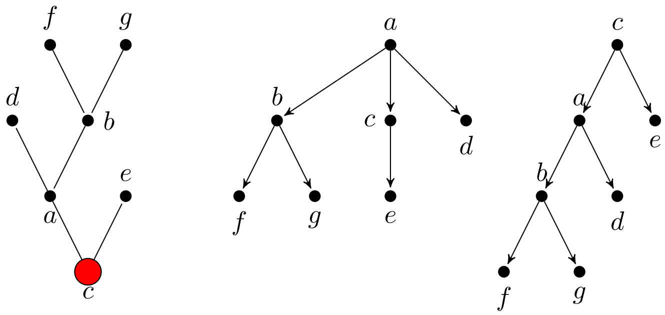

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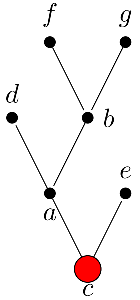
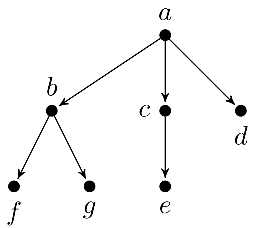
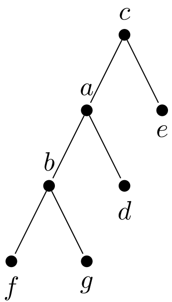


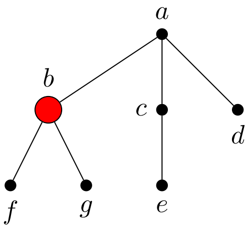




Terminology

Definition

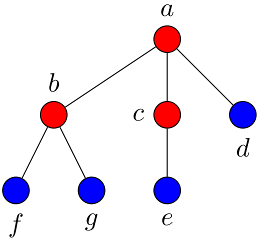
- parent (*cha*) of v is the unique u such that there is a directed edge from u to v
- when u is the parent of v, v is called a child (*con*) of u
- vertices with the same parent are called siblings (*anh em*)
- the ancestors (*tổ tiên*) of a vertex are the vertices in the path from the root to this vertex (excluding the vertex itself)
- descendants (*con cháu*) of a vertex v are those vertices that have v as an ancestor



Terminology

Definition

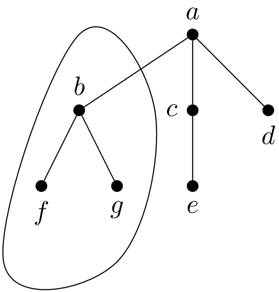
- a vertex of a tree is called a leaf (*lá*) if it has no children
- vertices that have children are called internal vertices (*đỉnh trong*)



Terminology

Definition

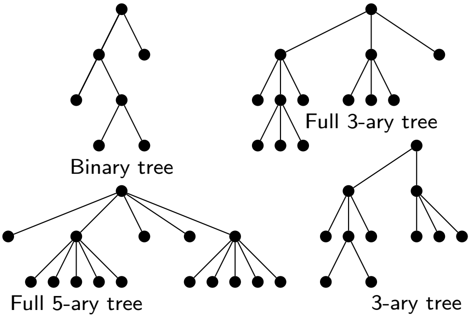
If a is a vertex in a tree, the subtree (*cây con*) with a as its root is the subgraph of the tree consisting of a and its descendants and all edges incident to these descendants.



m-ary tree

Definition

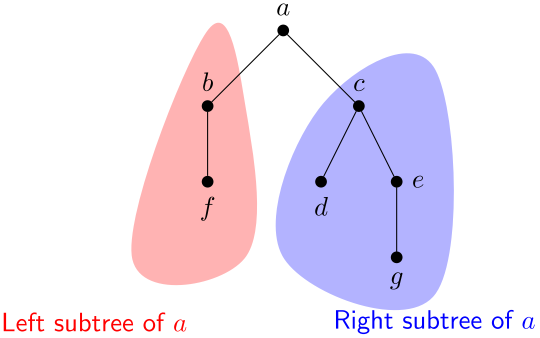
- m-ary tree (*cây m-phân*): at most m children on each internal vertex of a rooted tree.
- **full** m-ary tree (*cây m-phân đầy đủ*): every internal vertex has exactly m children.
- An m-ary tree with $m = 2$ is called a binary tree (*cây nhị phân*).



Ordered Rooted Trees

Definition

- An ordered rooted tree (*cây có gốc có thứ tự*) is a rooted tree where the children of each internal vertex are ordered (e.g. in order from left to right).



Ordered Rooted Trees

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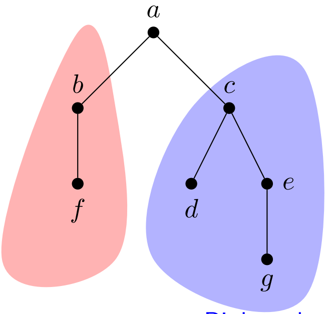
- In an ordered binary tree (*cây nhị phân có thứ tự*), if an internal vertex has two children, the first child is called the left child (*con bên trái*) and the second is called the right child (*con bên phải*).



Left subtree of a Right subtree of a

Properties & Theorems

Theorem

A tree with n vertices has $n - 1$ edges.

Theorem

A full m-ary tree

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A tree with n vertices has $n - 1$ edges.

Theorem

A full m-ary tree

(i) *n vertices has $(n-1)/m$ internal vertices and $[(m-1)n+1]/m$ leaves*

(ii) *i internal vertices has $n = mi + 1$ vertices and $(m-1)i+1$ leaves*

(iii) *ℓ leaves has $n = (m\ell - 1)/(m-1)$ vertices and $(\ell - 1)/(m-1)$ internal vertices*

Example

Example (Chain Letter Game)

- Each person who receives the letter is asked to send it on to four other peoples.
- Some peoples do this, but others do not send any letters.
- How many people have seen the letter, including the first person, if no one receives more than one letter and if the chain letter ends after there have been 100 people who read it but did not send it out ?
- How many people sent out the letter?

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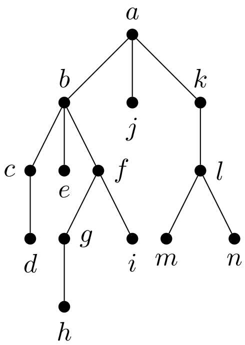
Solution

- *Using 4-ary tree with 100 leaves corresponding to 100 persons who did not send out the letter.*
- $\Longrightarrow n = (ml - 1)/(m - 1) = (4 \times 100 - 1)/(4 - 1) = 133$ *vertices and* $i = n - l = 133 - 100 = 33$ *internal vertices.*

Level and Height

Definition

- The level (*mức*) of a vertex v in a rooted tree is the length of the unique path from the root to this vertex.
- The level of the root is defined to be zero.
- The height (*độ cao*) of a rooted tree is the maximum of the levels of vertices (i.e. the length of the longest path from the root to any vertex).



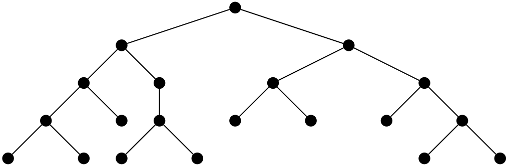
Example

- Level of root $a = 0$, $b, j, k = 1$ and $c, e, f, l = 2 \ldots$
- Because the largest level of any vertex is 4, this tree has height 4.

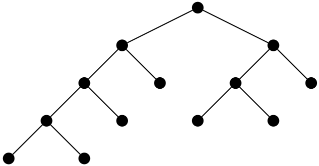
Balanced m-ary Trees

Definition

A rooted m-ary tree of height h is balanced (*cân đối*) if all leaves are at levels h or $h-1$.



T_1



T_2

Balanced m-ary Tree

Theorem

There are at most m^h leaves in an m-ary tree of height h.

Balanced m-ary Tree

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It can be proved by using mathematical induction on the height.

Balanced m-ary Tree

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Corollary

- *If an m-ary tree of height h has ℓ leaves, then $h \geq \lceil \log_m \ell \rceil$.*
- *If the m-ary tree is full and balanced, then $h = \lceil \log_m \ell \rceil$.*

Exercise

Exercise (Chess tournament)

Suppose 1000 people enter a chess tournament. Use a rooted tree model of the tournament to determine how many games must be played to determine a champion. If a player is eliminated after one loss and games are played until only one entrant has not lost. (Assume there are no ties)

Exercise (Isomorphic)

How many different isomers (*đồng phân*) do the following saturated hydrocarbons have ?

- C_3H_8
- C_5H_{12}
- C_6H_{14}

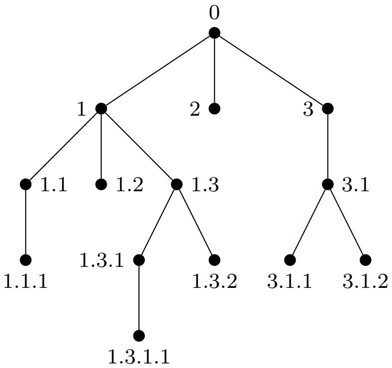
Question

Exercise

- How many vertices and how many leaves does a complete m-ary tree of height h have?
- Show that a full m-ary balanced tree (*cây m-phân hoàn hảo*) of height h has more than m^{h-1} leaves.
- How many edges are there in a forest of t trees containing a total of n vertices?

Labeling Ordered Rooted Trees

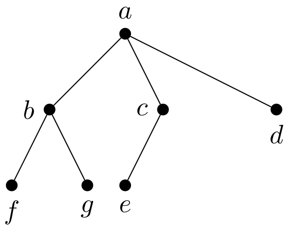
- Ordered rooted trees are often used to store information.
- Need a procedure for visiting each vertex of an ordered rooted tree to access data.
- Ordering and labeling the vertices is important to traverse them in any procedure
- Universal address system (*hệ địa chỉ phổ dụng*)
  $0 < 1 < 1.1 < 1.1.1 < 1.2 < 1.3 < \ldots < 2 < 3 < 3.1 < \ldots$



Traversal Algorithms (Thuật toán duyệt cây)

Preorder Traversal (duyệt tiền thứ tự - NLR)

procedure *preorder*(T: ordered rooted tree)
  r := root of T
  print r
  for each child c of r from left to right
    $T(c)$:= subtree with c as its root
    preorder($T(c)$)



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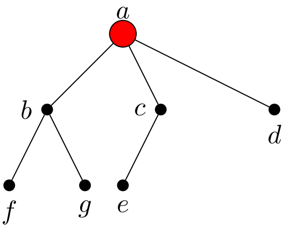

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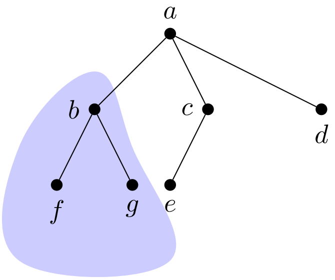

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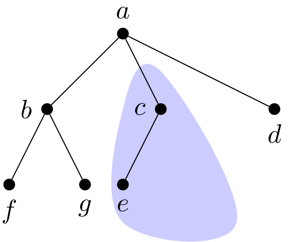

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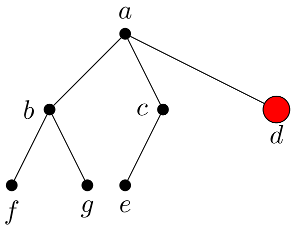

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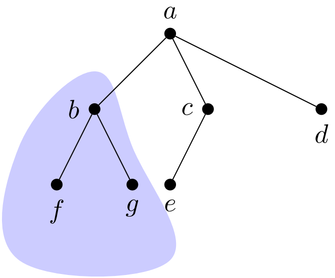

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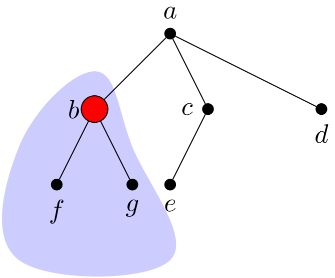

a   b   d

Traversal Algorithms (Thuật toán duyệt cây)

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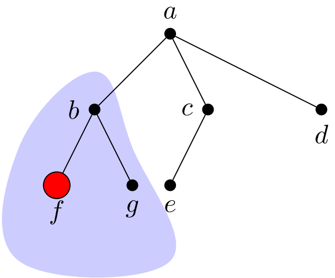

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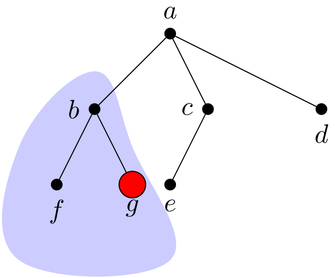

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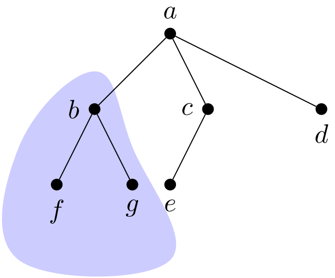

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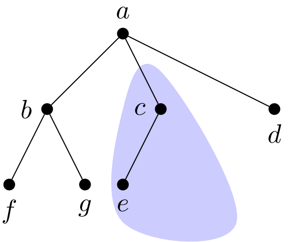

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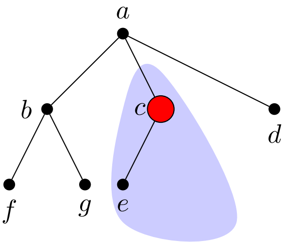

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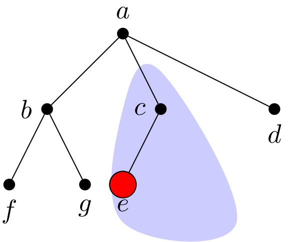

a b f g c e d

Traversal Algorithms

Inorder Traversal (Duyệt trung thứ tự - LNR)

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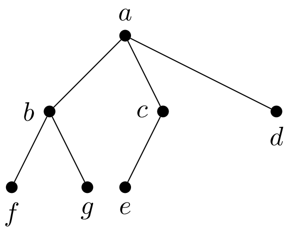

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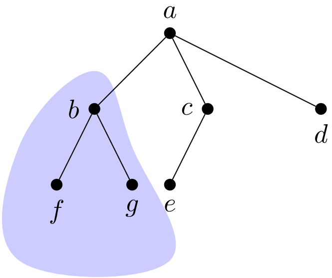

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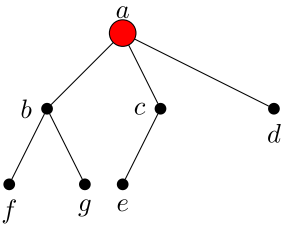

a

Traversal Algorithms

Huynh Tuong Nguyen, Nguyen An Khuong, Vo Thanh Hung

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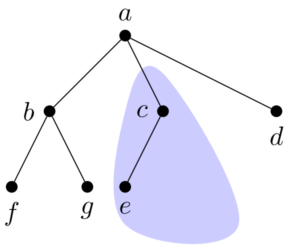

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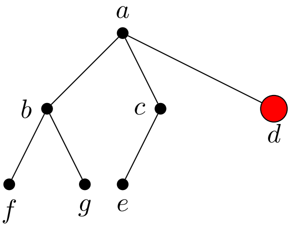

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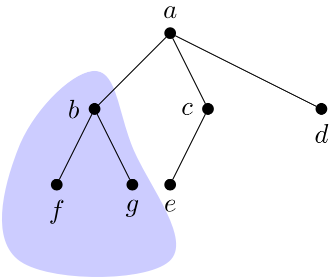

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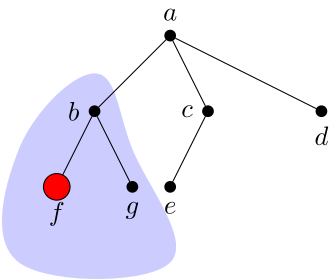

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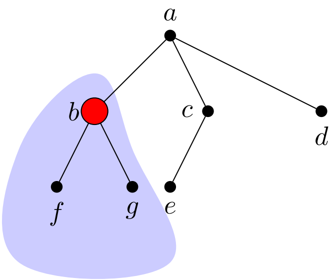

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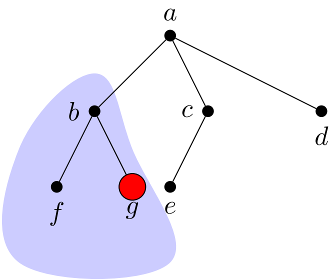

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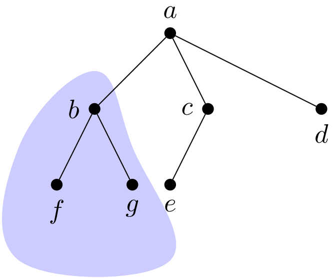

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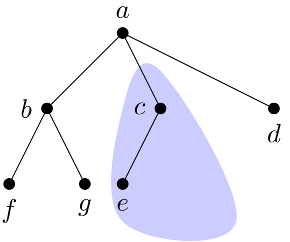

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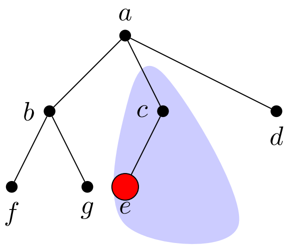

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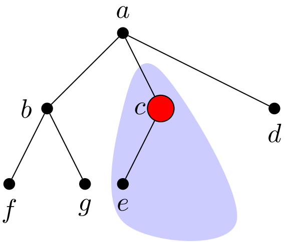

f b g a e c d

Traversal Algorithms

Postorder Traversal (Duyệt hậu thứ tự - LRN)

procedure *postorder*(T: ordered rooted tree)
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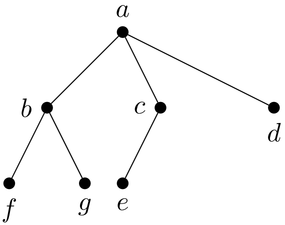

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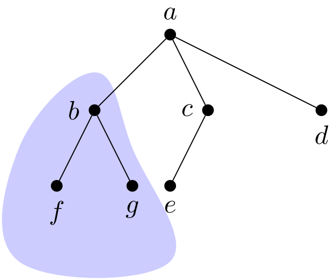

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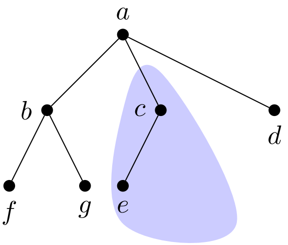

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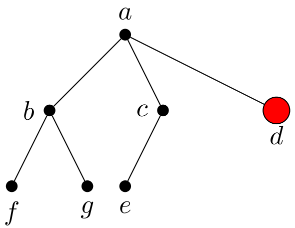

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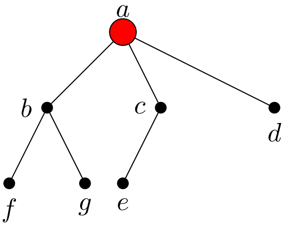

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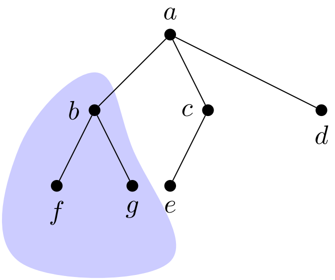

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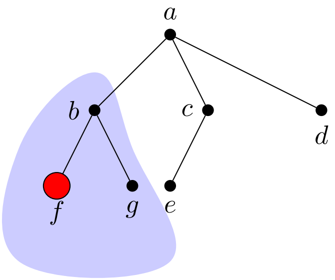

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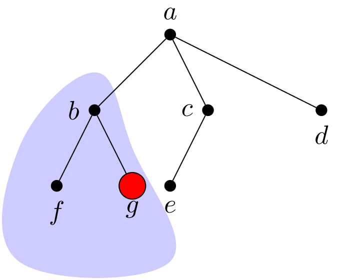

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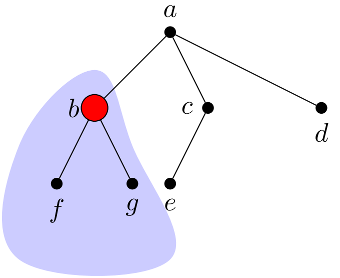

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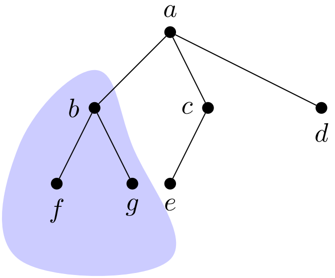

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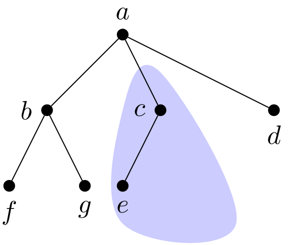

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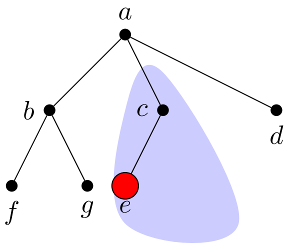

f g b e d a

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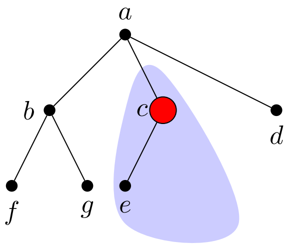
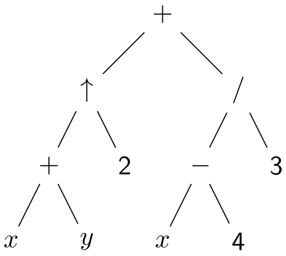


f g b e c d a

Infix, Prefix and Postfix Notations

- Infix (*trung tố*):
  $((x+y)\uparrow 2)+((x-4)/3)$
- Prefix (*tiền tố*):
  $+\ \uparrow\ +\ x\ y\ 2\ /\ -\ x\ 4\ 3$
- Postfix (*hậu tố*):
  $x\ y\ +\ 2\ \uparrow\ x\ 4\ -\ 3\ /\ +$



Exercise

Exercise

Find the ordered rooted tree representing

$$(\neg(p \wedge q) \vee (\neg q \wedge r)) \rightarrow (\neg p \vee \neg r)$$

Then use this rooted tree to find the prefix, postfix and infix forms of this expression

Exercise

Exercise

Find the ordered rooted tree representing

$$(\neg(p \wedge q) \vee (\neg q \wedge r)) \rightarrow (\neg p \vee \neg r)$$

Then use this rooted tree to find the prefix, postfix and infix forms of this expression

Solution

- *Constructing the rooted tree from the bottom up*
- *Preorder traversal creates prefix notation*
  $\rightarrow \vee \neg \wedge p\ q \wedge \neg q\ r \vee \neg p\ \neg r$
- *Postorder traversal creates postfix notation*
  $p\ q \wedge \neg \vee q \neg r \wedge p \neg r \neg \vee \rightarrow$
- *Inorder traversal creates infix notation (with parentheses)*
  $p\ q \neg \vee q \neg \wedge r \rightarrow p \neg \vee r \neg$

Exercise

Exercise

Find postorder traversal of a binary tree with inorder D B H E I A F C J G K and preorder A B D E H I C F G J K.

Exercise

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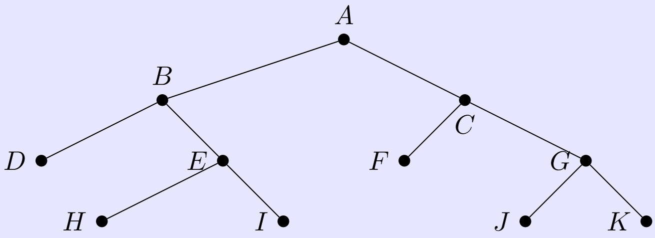
Find postorder traversal of a binary tree with inorder D B H E I A F C J G K and preorder A B D E H I C F G J K.

Solution



Post order: D H I E B F J K G C A.

Exercise

Exercise

Find in-order traversal of a binary tree with pre-order A D E B J C F H I G and post-order E J B D H I F G C A.

Exercise

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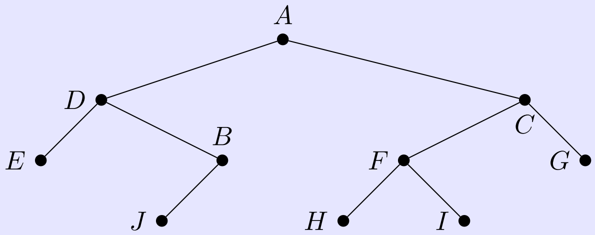
Find in-order traversal of a binary tree with pre-order A D E B J C F H I G and post-order E J B D H I F G C A.

Solution



In-order: E D J B A H F I C G.

Exercise

Exercise

How many different trees are there with the in-order of K E B J C A H G I D F and father-child relations respecting to the alphabet order.

Exercise

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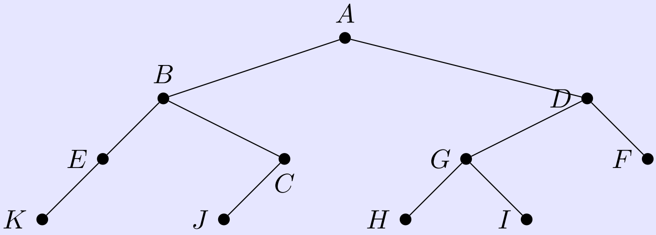
How many different trees are there with the in-order of K E B J C A H G I D F and father-child relations respecting to the alphabet order.

Solution



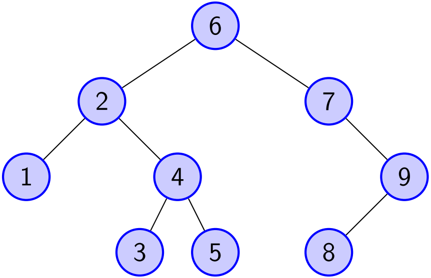
Pre-order: E D J B A H F I C G.

Binary Search Trees

Definition

Binary search tree (*cây tìm kiếm nhị phân* - BST) is a binary tree in which the assigned key of a vertex is:

- larger than the keys of all vertices in its left subtree, and
- smaller than the keys of all vertices in its right subtree.



Adding and Locating an Item in BST

Example

Form a BST for the words *mathematics*, *physics*, *geography*, *zoology*, *meteorology*, *geology*, *psychology*, *chemistry* using alphabetical order.

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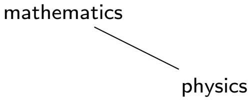

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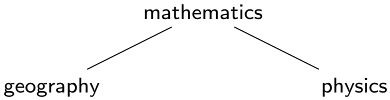

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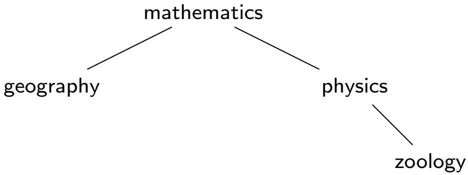

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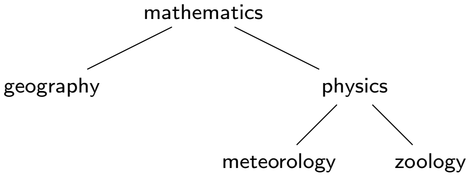

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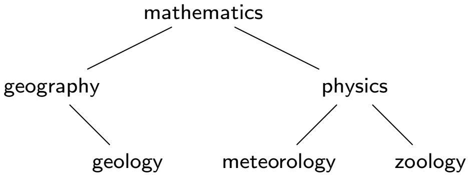

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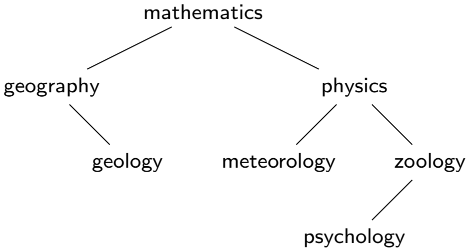

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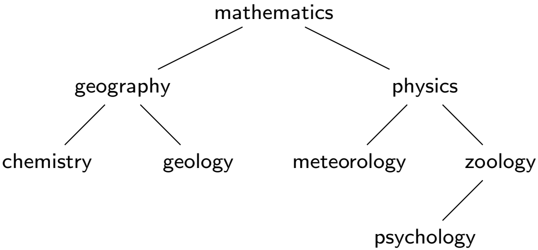

Complexity in searching

$O(\log(n))$ vs. $O(n)$ in linear list

Decision Trees (Cây quyết định)

Example

There are seven coins, all with the same weight, and a counterfeit coin that weighs less than the others. How many weighings are necessary using a balance scale to determine which of the eight coins is the counterfeit one? Give an algorithm for finding this counterfeit coin.

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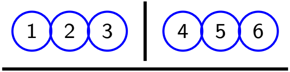

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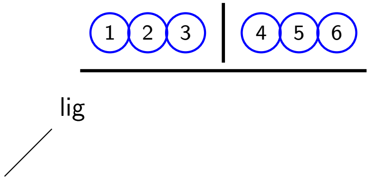

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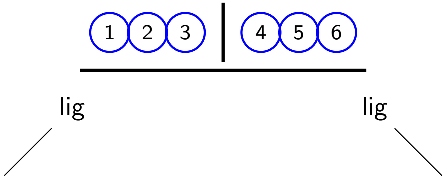

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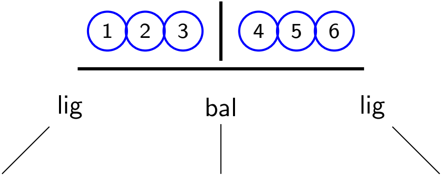

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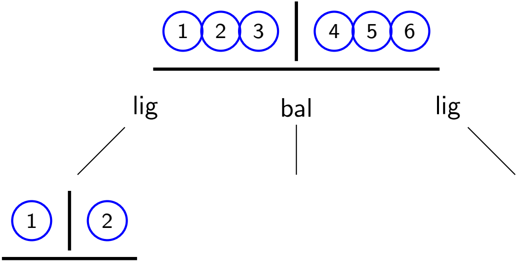

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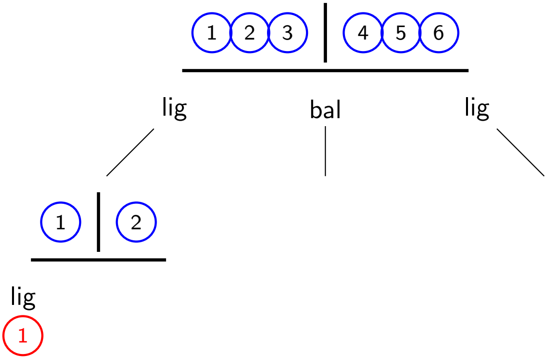

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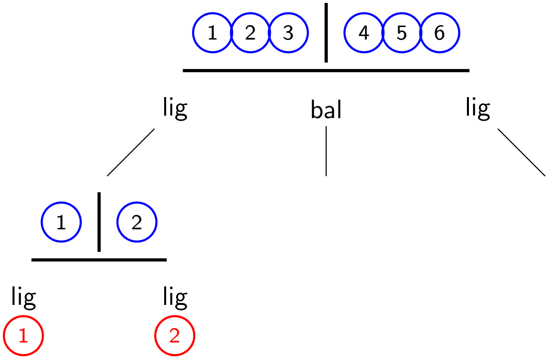

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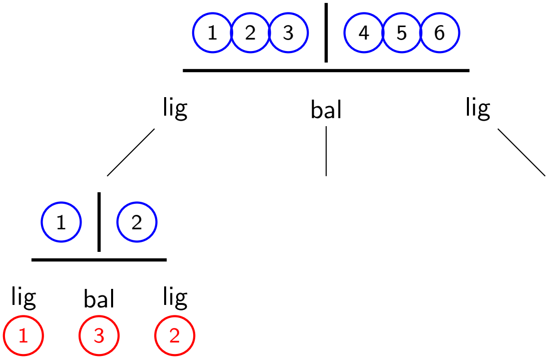

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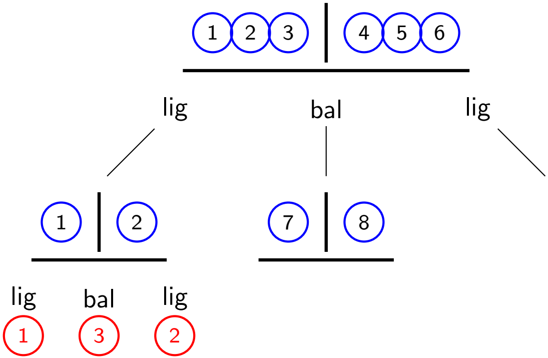

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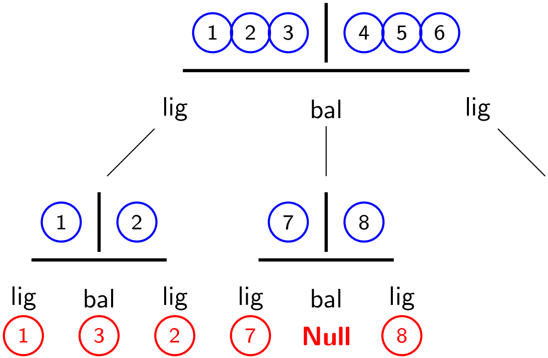

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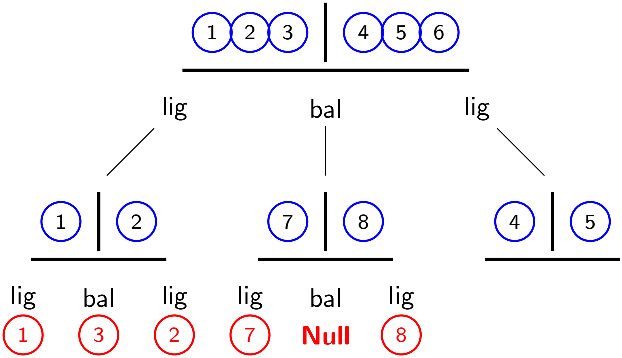

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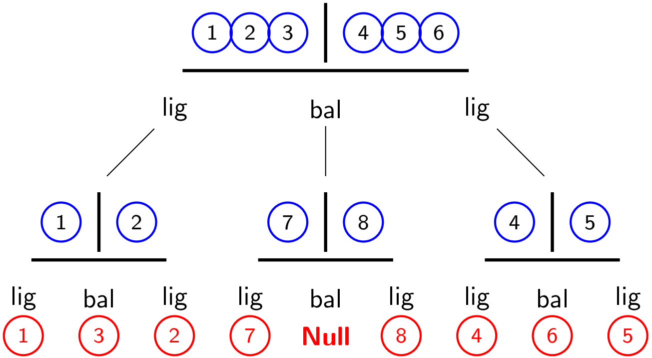

Yet Another Application

Example

If we know that the probability that a person has tuberculosis (TB) is $p(\mathrm{TB}) = 0.0005$.
We also know $p(+|\mathrm{TB}) = 0.999$ and $p(-|\overline{\mathrm{TB}}) = 0.99$.
What is $p(\mathrm{TB}|+)$ and $p(\overline{\mathrm{TB}}|-)$?

Start! •

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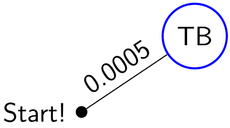

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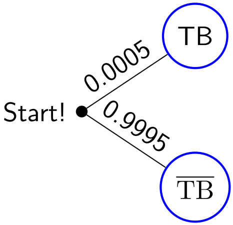

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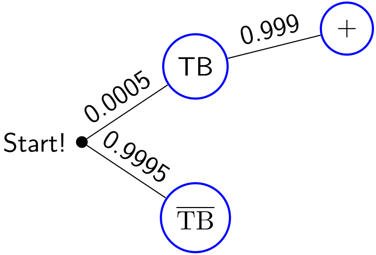

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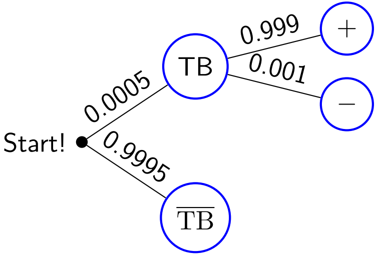

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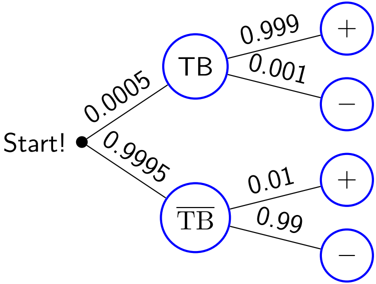

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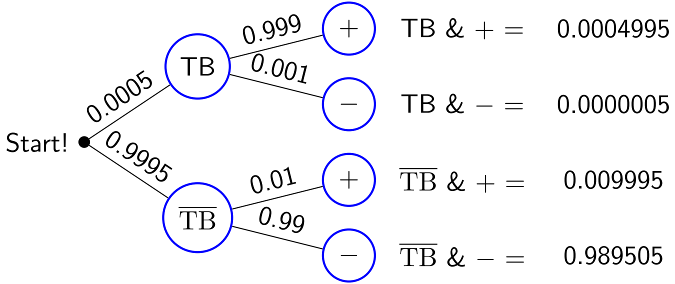

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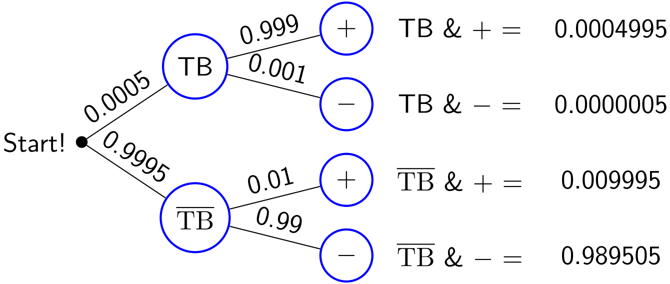

$$p(\mathrm{TB}|+) = \frac{p(\mathrm{TB} \cap +)}{p(+)} = \frac{0.0004995}{0.0004995 + 0.009995} \approx 0.0476$$

Problem

Definition

- A spanning tree (*cây khung*) in a graph G is a subgraph of G that is a tree which contains all vertices of G.

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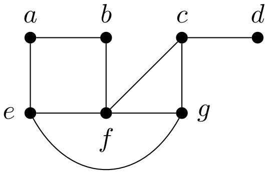

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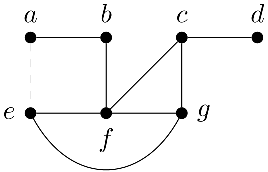

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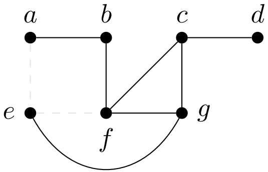

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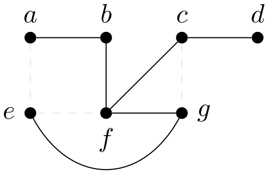
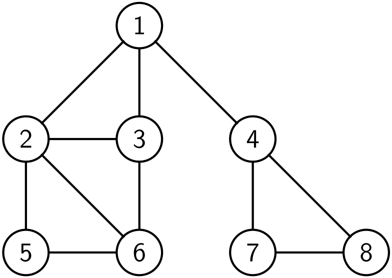


Depth-First Search (Tìm kiếm ưu tiên chiều sâu)



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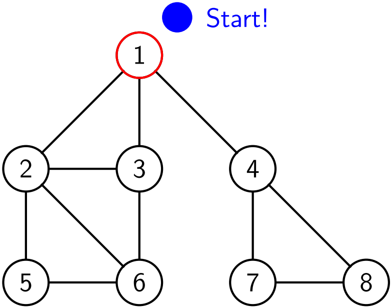

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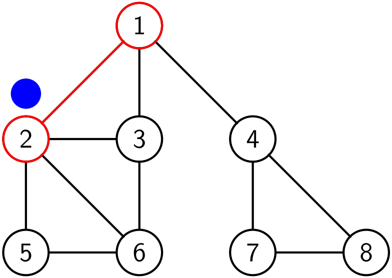

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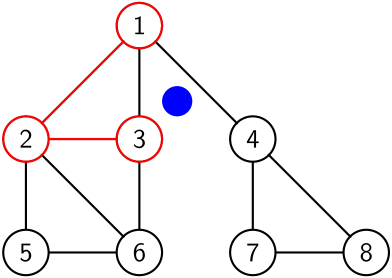

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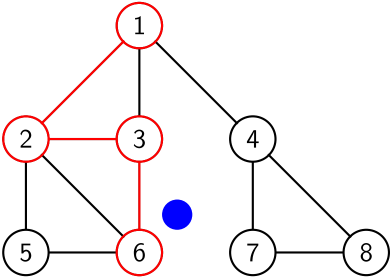

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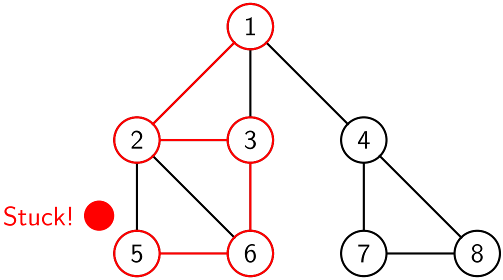

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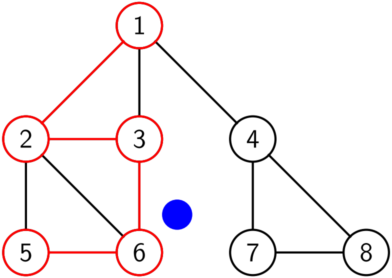

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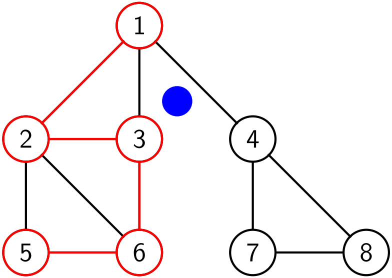

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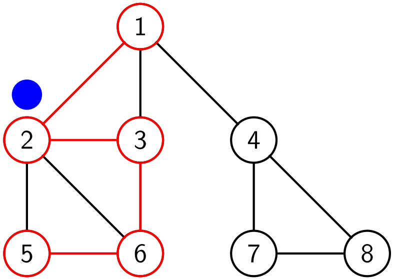

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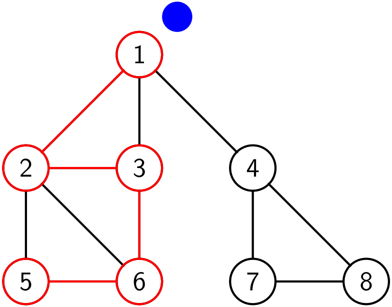

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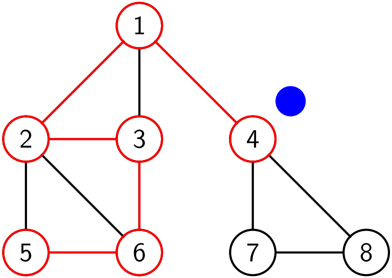

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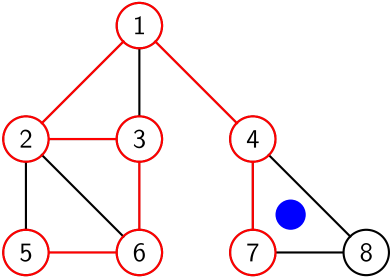

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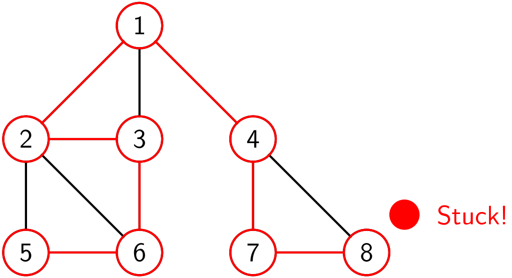

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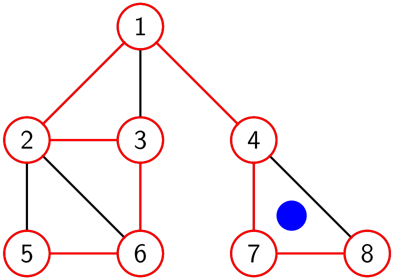

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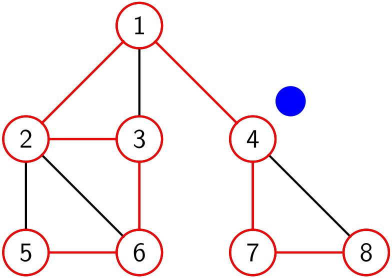
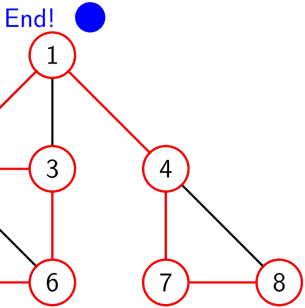
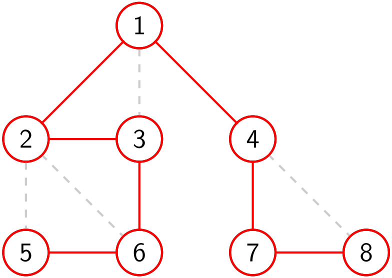


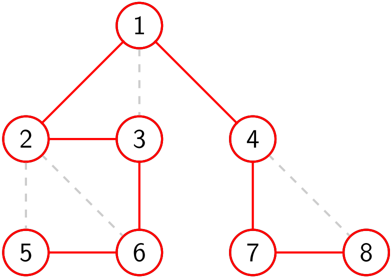
Depth-First Search (Tìm kiếm ưu tiên chiều sâu)



Depth-First Search (Tìm kiếm ưu tiên chiều sâu)



Depth-First Search (Tìm kiếm ưu tiên chiều sâu)



Property

- Go deeper as you can
- Backtrack (*quay lui*) to possible branch when you are stuck.
- $O(e)$ or $O(n^2)$

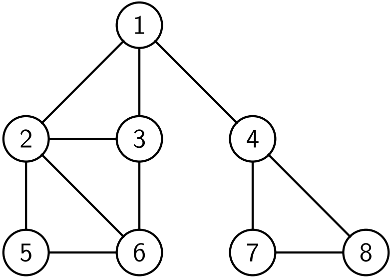
Depth-First Search

Algorithm

procedure *DFS* (G)
    $T :=$ tree consisting only vertex v_1
    visit(v_1)

procedure *visit*$(v$: vertex of $G)$ /* recursive */
    for each vertex w adjacent to v and not in T
        add w and edge $\{v, w\}$ to T
        visit(w)

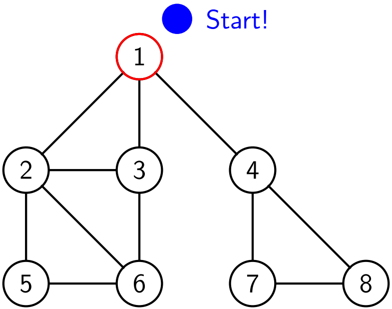
Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



vertex	L
	$\emptyset$

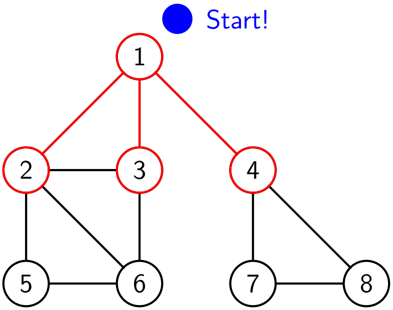
Huynh Tuong Nguyen, Nguyen An Khuong, Vo Thanh Hung

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



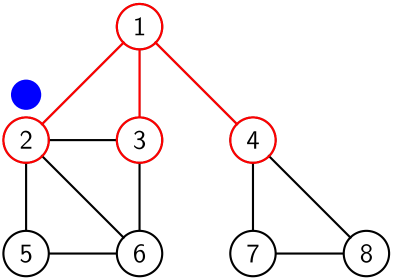
vertex	L
	$\emptyset$
1	

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



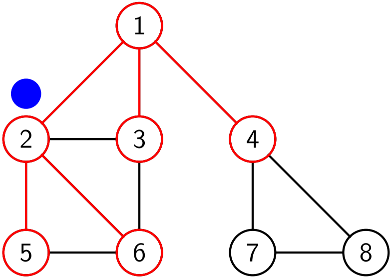
vertex	L
	$\emptyset$
1	2, 3, 4

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



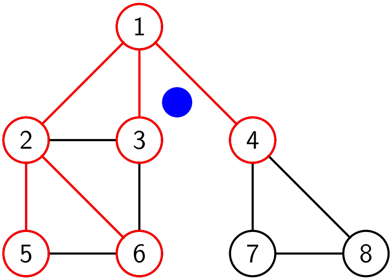
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



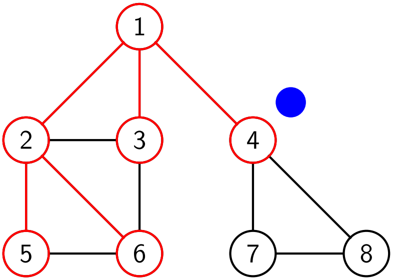
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



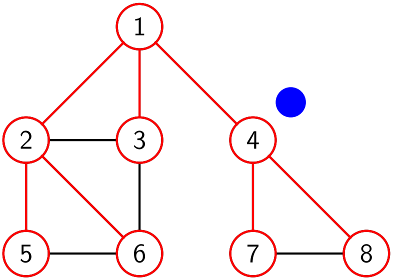
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



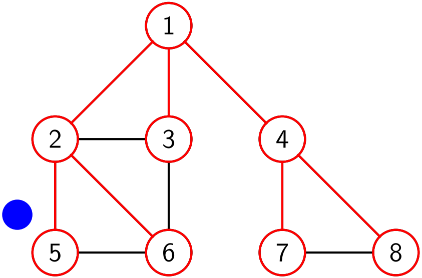
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



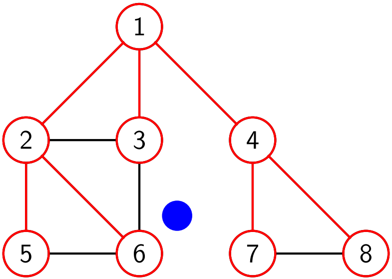
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6, 7, 8

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



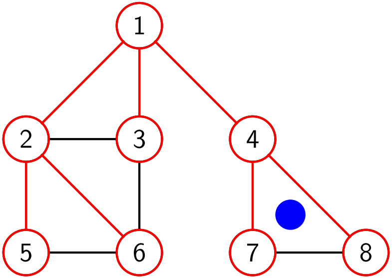
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6, 7, 8
5	6, 7, 8

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



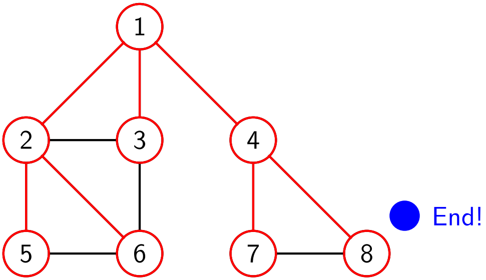
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6, 7, 8
5	6, 7, 8
6	7, 8

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



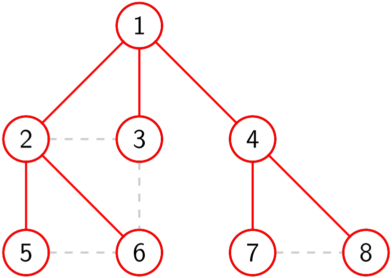
vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6, 7, 8
5	6, 7, 8
6	7, 8
7	8

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6, 7, 8
5	6, 7, 8
6	7, 8
7	8
8	$\emptyset$

Breadth-First Search (Tìm kiếm ưu tiên chiều rộng)



vertex	L
	$\emptyset$
1	2, 3, 4
2	3, 4, 5, 6
3	4, 5, 6
4	5, 6, 7, 8
5	6, 7, 8
6	7, 8
7	8
8	$\emptyset$

Property

- $O(e)$ or $O(n^2)$

Breadth-First Search

Algorithm

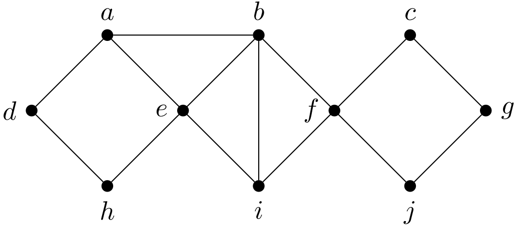
procedure *BFS* (G)

$T :=$ tree consisting only vertex v_1
$L :=$ empty list
put v_1 in the list L of unprocessed vertices
while L is not empty
    remove the first vertex, v, from L
    for each neighbor w of v
        if w is not in L and not in T **then**
            add w to the end of the list L
            add w and edge $\{v, w\}$ to T

Exercise

Exercise

Find spanning tree in the following graphs.



Exercise

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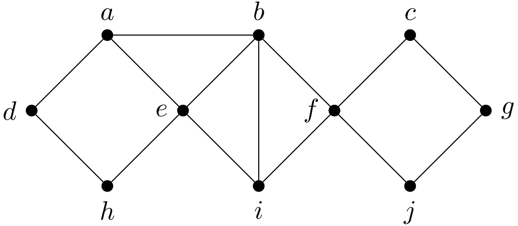

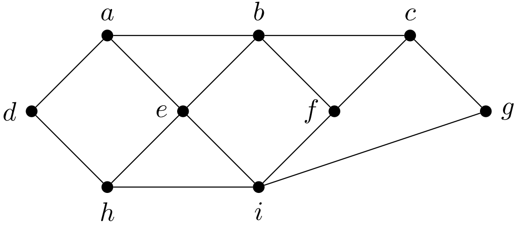

Minimum Spanning Trees

Definition

- A minimum spanning tree (*cây khung nhỏ nhất*) in a connected weighted graph is a spanning tree that has the smallest possible sum of weights of its edges.

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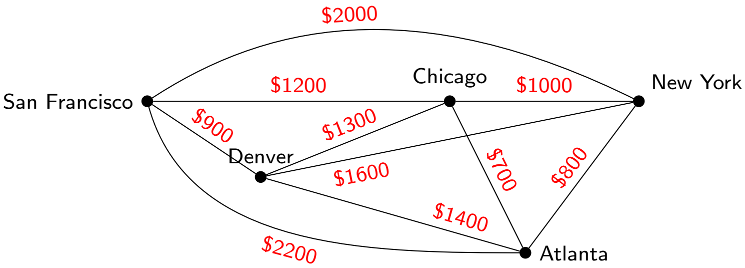

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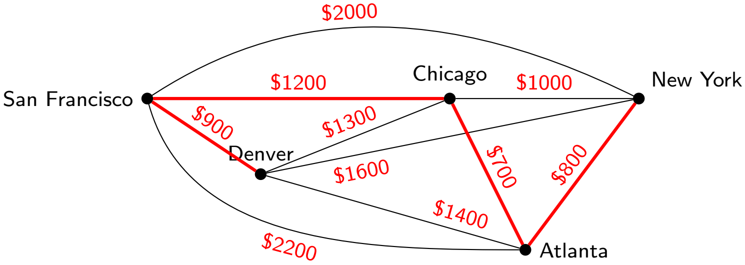


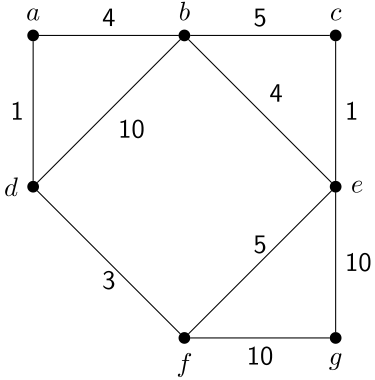
Prim's Algorithm (Nearest-Neighbor)

Prim's Algorithm (1957)

procedure *Prim*(G)
  $T :=$ a minimum-weight edge
  for $i := 1$ to $n - 2$
    $e :=$ an edge of minimum weight incident to a vertex in T and not forming a simple circuit in T if added to T
    $T := T$ with e added
  return T

Prim's Algorithm (Nearest-Neighbor)

- Pick a vertex to start from
- Iteratively absorb smallest edge possible



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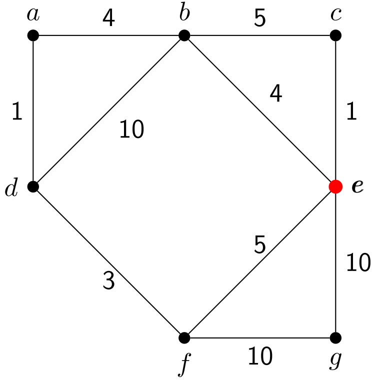

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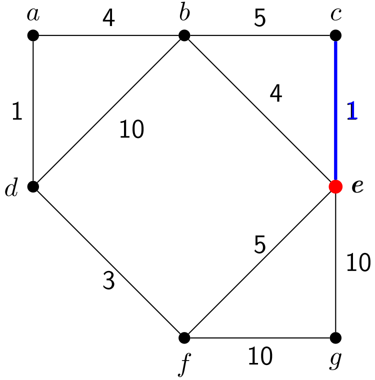

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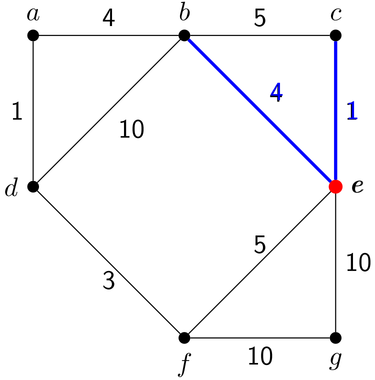

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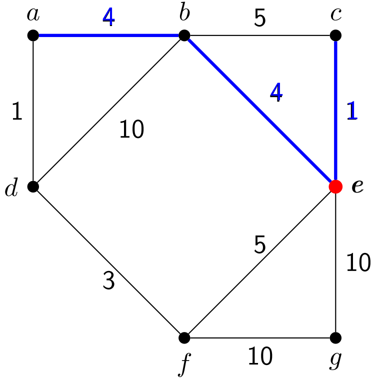

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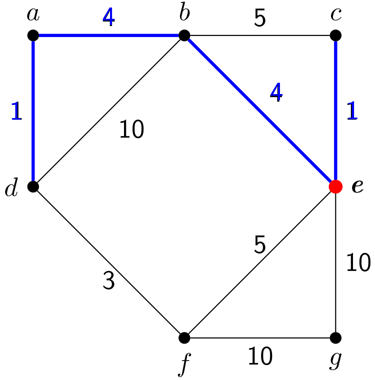

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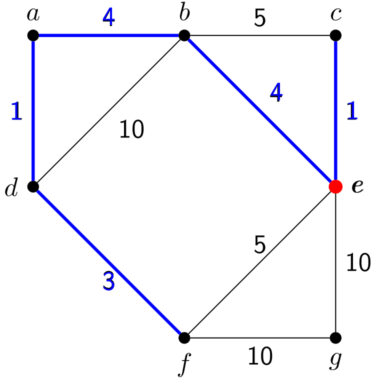

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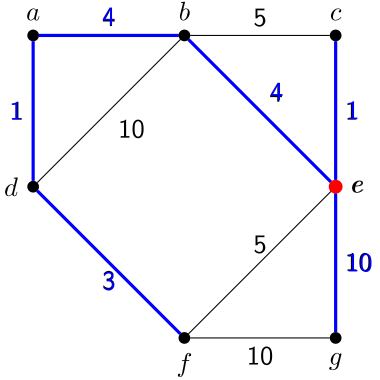


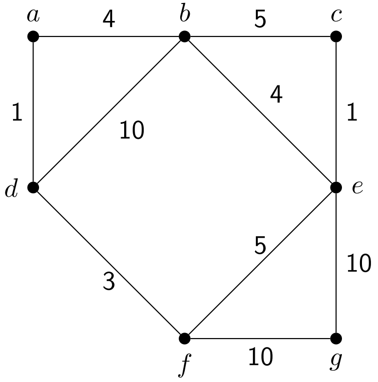
Kruskal's Algorithm (Lightest-Edge)

Kruskal's Algorithm (1958)

procedure *Kruskal*(G)
  $T :=$ empty graph
  for $i := 1$ **to** $n - 1$
    $e :=$ any edge in G with smallest weight that does not form a simple circuit when added to T
    $T := T$ with e added
  return T

Kruskal's Algorithm (Lightest-Edge)

- Iteratively add smallest edge possible



Kruskal's Algorithm (Lightest-Edge)

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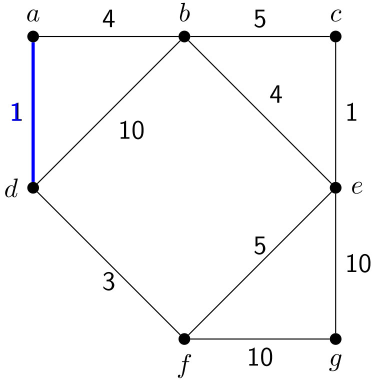

Kruskal's Algorithm (Lightest-Edge)

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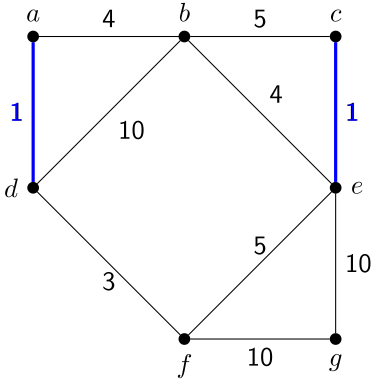

Kruskal's Algorithm (Lightest-Edge)

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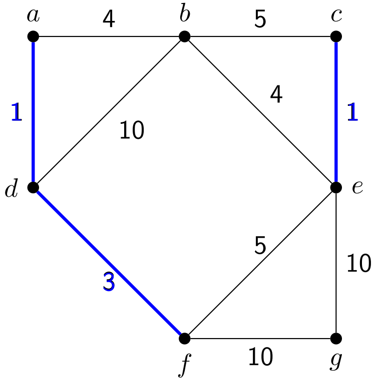

Kruskal's Algorithm (Lightest-Edge)

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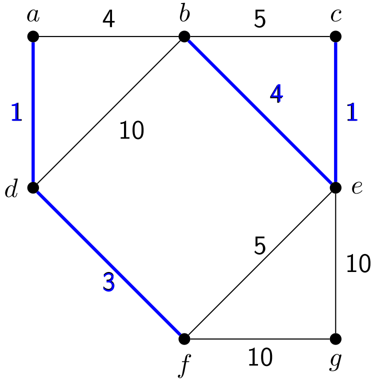

Kruskal's Algorithm (Lightest-Edge)

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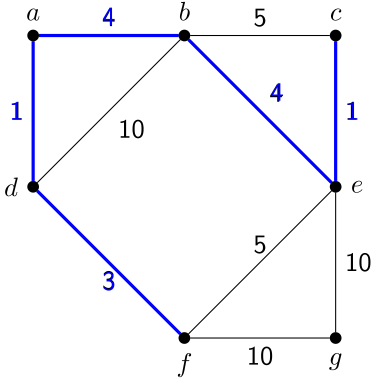

Kruskal's Algorithm (Lightest-Edge)

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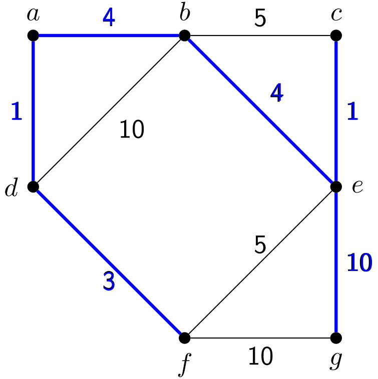

Exercise

Exercise

By using Prim's and Kruskal's algorithm, determine minimum spanning tree in the following graphs.

.

Exercise

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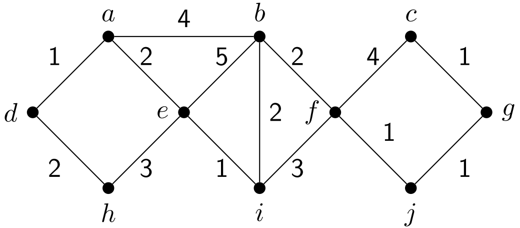
By using Prim's and Kruskal's algorithm, determine minimum spanning tree in the following graphs.

.

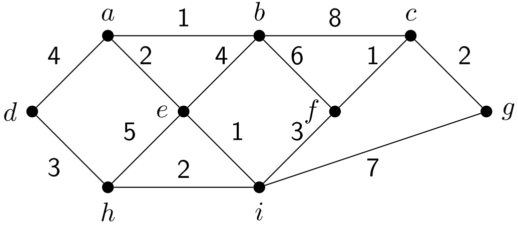


Exercise

Exercise

By using Prim's and Kruskal's algorithm, determine minimum spanning tree in the following graphs.

.

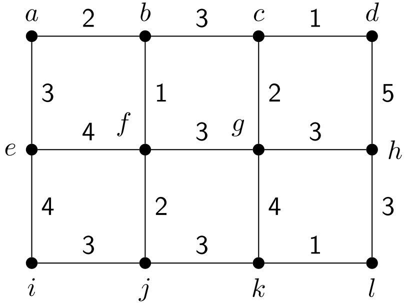


Exercise

Exercise

By using Prim's and Kruskal's algorithm, determine minimum spanning tree in the following graphs.

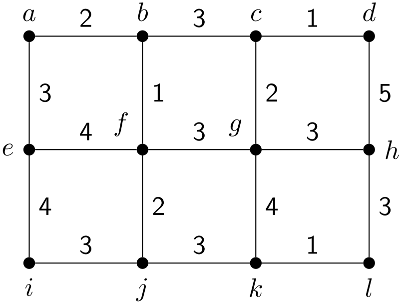
.



Exercise

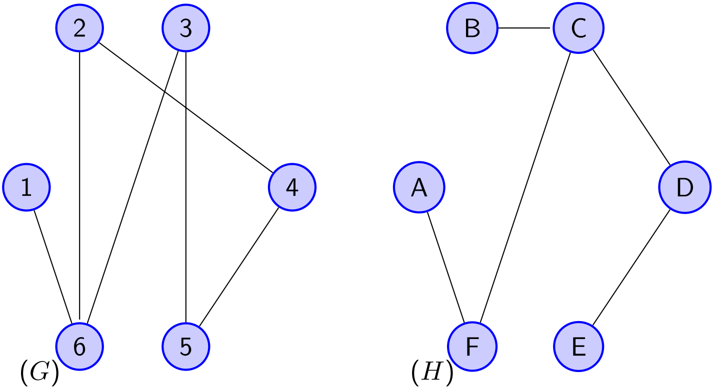
Exercise

By using Prim's and Kruskal's algorithm, determine minimum spanning tree in the following graphs. (and maximum spanning tree (*cây khung cực đại*).



Revision

Cho hai đồ thị G và H như sau:



Chọn phát biểu đúng.

A) G là cây

B) G và H là đẳng cấu

C) Xoá một cạnh trong G thì thu được một cây

D) Xoá một cạnh trong H thì thu được một cây

Revision

What is the value of each of these prefix expressions?

a) - * 2 / 8 4 3

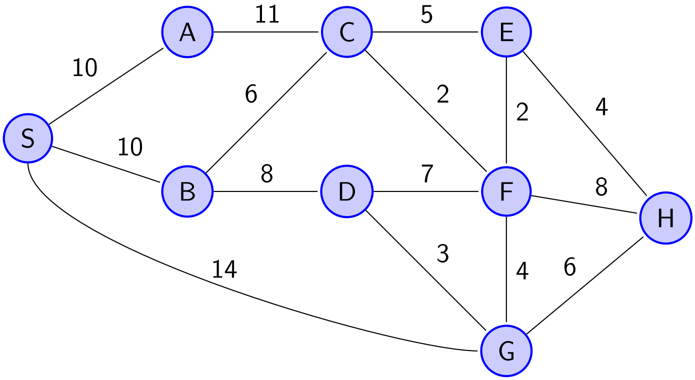
b) $*$ - * 3 3 * 4 2 5

c) $+$ - $*$ 3 2 $+$ 2 3 / 6 - 4 2

d) * $+$ 3 $+$ 3 $*$ 3 $+$ 3 3 3

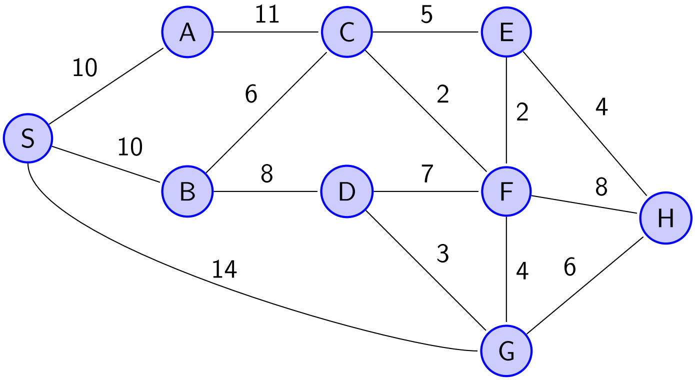
Revision

Xác định cây phủ tối thiểu cho đồ thị như trong hình vẽ dưới (áp dụng hai phương pháp).



Revision

Xác định cây phủ tối thiểu cho đồ thị như trong hình vẽ dưới (áp dụng hai phương pháp).



By using Prim's or Kruskal's algorithm, could we determine a minimum spanning tree in a directed graph?
Explain it.

Revision

Cho $G = (V, E)$ là một đồ thị đơn và vô hướng bất kỳ, có n đỉnh. Định nghĩa đồ thị bù của G là $G^c = (V, F)$ thỏa hai tính chất: $G \cup G^c = K_n$ và $E \cap F = \emptyset$.

Gọi T là cây bao trùm của đồ thị đầy đủ K_6. Số lượng cạnh của đồ thị bù T^c là

A) 5

B) 10

C) 15

D) 20