

# Concepts and Models of Parallel and Datacentric Programming

**Distributed Shared Memory** 

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#### **Outline**

- Organization
- Foundations
- Shared Memory
- 3. GPU Programming
- Bulk-Synchronous Parallelism
- Message Passing
- 6. Distributed Shared Memory
- 7. Parallel Algorithms
- 8. Parallel I/O
- 9. MapReduce
- 10. Apache Spark

- a. PGAS Foundations
- b. DASH Overview
- Distributed Data Structured
- d. DASH Algorithms
- e. Tasking







## **PGAS** Foundations

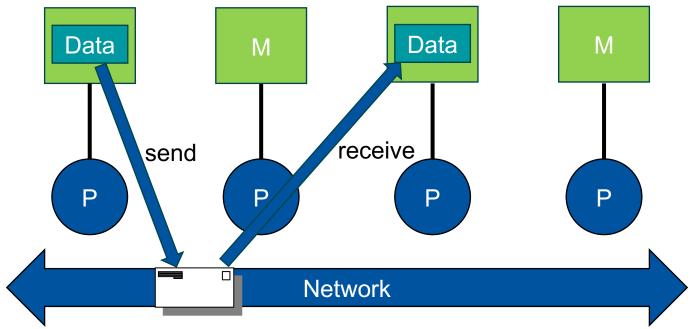






#### **Parallel Architectures**

- Distributed Memory
  - Each processing element (P) has its separate main memory (M)



Data exchange is achieved through message passing over the network

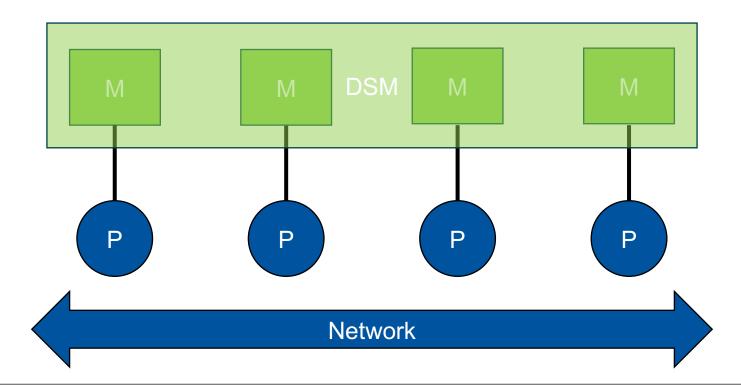






#### **Parallel Architectures**

- Distributed Shared Memory
  - All processing elements (P) have direct access to their main memory (M)
  - Each processing element (P) has access to the shared memory (DSM)



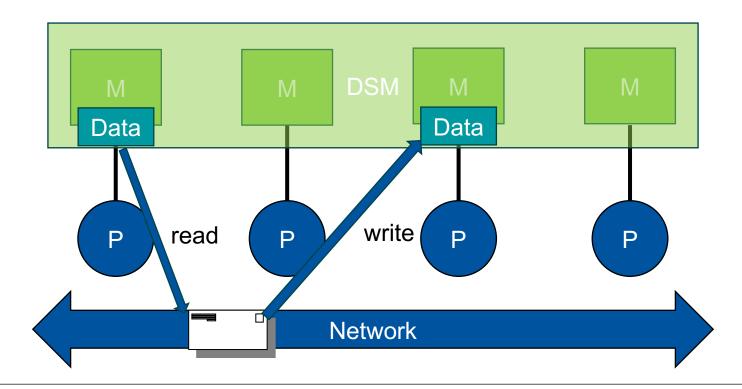






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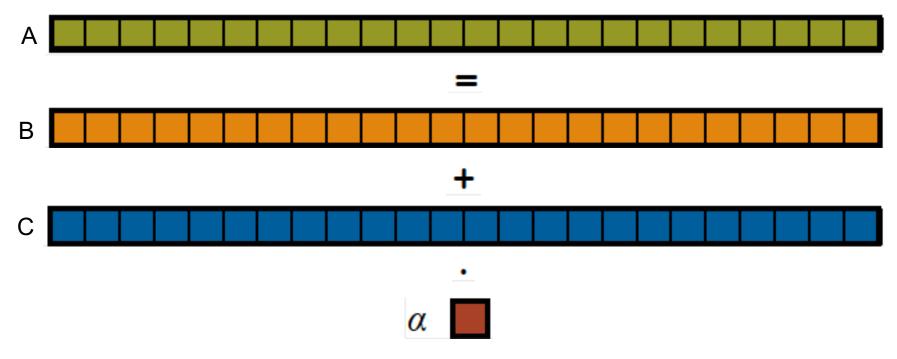








- Given: m-element vectors A, B, C
- Compute:  $\forall_{i \in 1...m}$ :  $A_i = B_i + \alpha * C_i$
- In pictures:



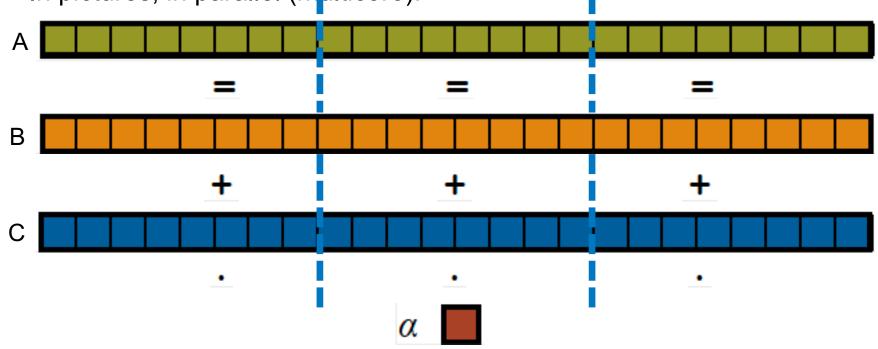






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• In pictures, in parallel (multicore):



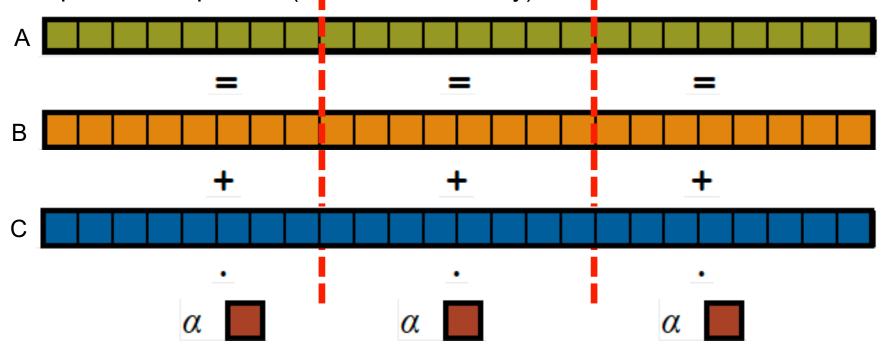






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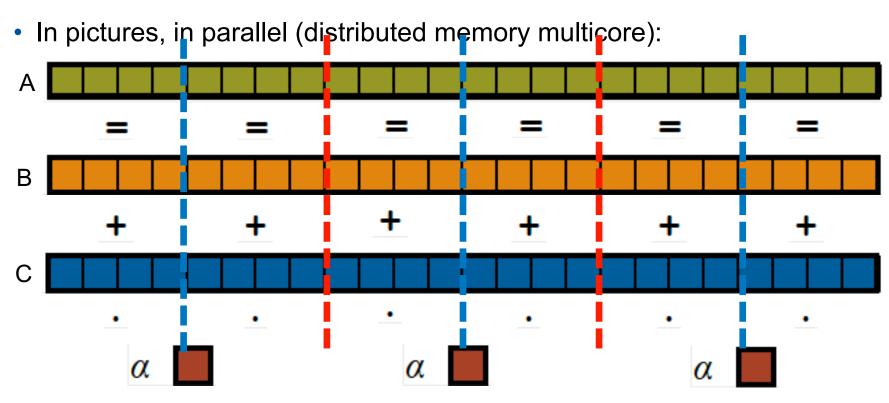
In pictures, in parallel (distributed memory):







- Given: m-element vectors A, B, C
- Compute:  $\forall_{i \in 1...m}$ :  $A_i = B_i + \alpha * C_i$









### **HPC STREAM Triad: MPI + OpenMP**

```
#include <hpcc.h>
#ifdef OPENMP
#include <omp.h>
#endif
static int VectorSize;
static double *a, *b, *c;
int HPCC StarStream(HPCC Params| *params) {
  int myRank, commSize;
  int rv, errCount;
 MPI Comm comm = MPI COMM WORLD;
 MPI Comm size ( comm, &commSize );
 MPI Comm rank (comm, &myRank);
  rv = HPCC_Stream( params, 0 == myRank);
 MPI Reduce ( &rv, &errCount, 1, MPI INT, MPI SUM,
   0, comm );
  return errCount;
int HPCC Stream(HPCC Params *params, int doIO) {
  register int j;
  double scalar;
  VectorSize = HPCC LocalVectorSize( params, 3,
   sizeof(double), 0 );
  a = HPCC XMALLOC( double, VectorSize );
 b = HPCC XMALLOC( double, VectorSize );
  c = HPCC XMALLOC( double, VectorSize );
```

```
if (!a || !b || !c) {
    if (c) HPCC free(c);
    if (b) HPCC free(b);
    if (a) HPCC free(a);
    if (doIO) {
      fprintf(outFile, "Failed to allocate memory (%d).
   \n", VectorSize );
      fclose( outFile );
    return 1;
#ifdef OPENMP
#pragma omp parallel for
#endif
 for (j=0; j<VectorSize; j++) {
   b[j] = 2.0;
    c[j] = 0.0;
  scalar = 3.0;
#ifdef OPENMP
#pragma omp parallel for
#endif
  for (j=0; j<VectorSize; j++)</pre>
    a[i] = b[i] + scalar * c[i];
  HPCC free(c);
  HPCC free(b);
 HPCC free(a);
```







- Shared Memory, e.g., Threading
  - Support dynamic, fine-grain parallelism
  - Considered simpler, more like traditional programming
    - If you want to access something, simply name it

- But:
  - Limited scalability
  - Bugs can be subtle, difficult to track down







- Distributed Memory, e.g., Message-Passing with MPI
  - More constrained model, can only access local data
  - Runs on all large-scale parallel platforms
    - And often can achieve near-optimal performance
  - Can serve as a strong foundation for high-level models

#### But:

- Communication must be used to get copies of remote data
  - Reveals (too) much about how to transfer, not what to transfer
- Couples data transfer and synchronization
- Has classes of bugs of its own







- Hybrid, e.g., MPI + Threading
  - Popular in the context of HPC (Threading: OpenMP) for highest performance
  - Division of labor: each handles what it does best
  - Overheads are amortized across processor cores

- But:
  - Requires multiple notations to express single logical parallel algorithm
  - Distinct semantics of MPI and Threading / OpenMP
  - May hold surprises in terms of unexpected side effects...







- Partitioned Global Address Space
  - (or: partitioned global namespace)
  - Can come as a language (extension) or API
  - Support a shared namespace on distributed memory ...
    - Permit any parallel task to access any lexically visible variable
  - ... with a sense of ownership
    - Variables have well-defined location, local accesses are cheaper than remote ones

	partitioned sh	nared name-/a	ddress space	
private	private	private	private	private
space 0	space 1	space 2	space 3	space 4







## Comparison

Name	Memory Model	Programming Model	Execution Model	Data Structures	Communica- tion
MPI	Distributed Memory	SPMD	Cooperating Executables	Manually fragmented	APIs
OpenMP	Shared Memory	Global-view parallelism	Threading	Shared memory arrays	-
BSP	Distributed Memory	SPMD	SPMD	Distributed variables, Co-arrays	Explicit and Implicit
PGAS	Partitioned Global Address Space	SPMD	SPMD	Co-arrays	Implicit

- In this lecture, we present the DASH model
  - Asynchronous Partitioned Global Address Space
  - Implementation developed in the context of DFG SPPEXA prog





