



Robotic Systems II

Lecture 4: Deterministic Optimal Control & LQR

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Continuous Time:

$$\begin{aligned} \underset{\mathbf{x}(t), \mathbf{u}(t)}{\operatorname{argmin}} \mathcal{J}(\mathbf{x}(t), \mathbf{u}(t)) &= \int_{t_0}^{t_f} \ell(\mathbf{x}(t), \mathbf{u}(t)) dt + \ell_F(\mathbf{x}(t_f)) \\ \text{s.t.} \quad \dot{\mathbf{x}}(t) &= f(\mathbf{x}(t), \mathbf{u}(t)) \end{aligned}$$

where

- $\mathbf{x}(t) \in \mathbb{R}^N$, $\mathbf{u}(t) \in \mathbb{R}^M$ are the state and control trajectories
- $\mathcal{J}(\mathbf{x}(t), \mathbf{u}(t))$ is the “cost function”
- $\ell(\mathbf{x}(t), \mathbf{u}(t))$ is the “stage cost”
- $\ell_F(\mathbf{x}(t_f))$ is the “terminal cost”
- $\dot{\mathbf{x}}(t) = f(\mathbf{x}(t), \mathbf{u}(t))$ are the “dynamics constraints”
- We can potentially add more constraints (e.g. torque limits)

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- Only very few problems can be solved analytically

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- The solution is open loop control trajectories; we know everything!
- Only very few problems can be solved analytically
- Let's discretize!

Discrete Time:

$$\begin{aligned} \underset{\mathbf{x}_{1:K}, \mathbf{u}_{1:K-1}}{\operatorname{argmin}} \quad & \mathcal{J}(\mathbf{x}_{1:K}, \mathbf{u}_{1:K-1}) = \sum_{k=1}^{K-1} \ell(\mathbf{x}_k, \mathbf{u}_k) + \ell_F(\mathbf{x}_K) \\ \text{s.t.} \quad & \mathbf{x}_{k+1} = f_{\text{discrete}}(\mathbf{x}_k, \mathbf{u}_k) \end{aligned}$$

- $\mathbf{x}_k \in \mathbb{R}^N$ and $\mathbf{u}_k \in \mathbb{R}^M$ are vectors
- The is now **“finite dimensional”**
- The solution is still open loop control trajectories; we know everything!
- We usually call $\mathbf{x}_k, \mathbf{u}_k$ **“knot points”**

Pontryagin's Minimum Principle

Pontryagin's Minimum Principle basically refers to the **KKT conditions** for the optimal control problem. We will skip the derivation here and focus on the result:

$$\mathbf{x}_{k+1} = \nabla_{\lambda} H(\mathbf{x}_k, \mathbf{u}_k, \lambda_{k+1})$$

$$\lambda_k = \nabla_{\mathbf{x}} H(\mathbf{x}_k, \mathbf{u}_k, \lambda_{k+1})$$

$$\mathbf{u}_{k+1} = \underset{\mathbf{u}}{\operatorname{argmin}} H(\mathbf{x}_k, \mathbf{u}, \lambda_{k+1}), \text{ s.t. } \mathbf{u} \in \mathcal{U}$$

$$\lambda_K = \frac{\partial \ell_F}{\partial \mathbf{x}_K}$$

where $H(\mathbf{x}_k, \mathbf{u}_k, \lambda_{k+1}) = \ell(\mathbf{x}_k, \mathbf{u}_k) + \lambda_{k+1}^T f_{\text{discrete}}(\mathbf{x}_k, \mathbf{u}_k)$.

What is the Linear Quadratic Regulator (LQR) problem?

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$$\underset{\mathbf{x}_{1:K}, \mathbf{u}_{1:K-1}}{\operatorname{argmin}} \mathcal{J}(\mathbf{x}_{1:K}, \mathbf{u}_{1:K-1}) = \sum_{k=1}^{K-1} \left(\frac{1}{2} \mathbf{x}_k^T \mathbf{Q}_k \mathbf{x}_k + \frac{1}{2} \mathbf{u}_k^T \mathbf{R}_k \mathbf{u}_k \right) + \frac{1}{2} \mathbf{x}_K^T \mathbf{Q}_K \mathbf{x}_K$$

$$\text{s.t.} \quad \mathbf{x}_{k+1} = \mathbf{A}_k \mathbf{x}_k + \mathbf{B}_k \mathbf{u}_k$$

$$\mathbf{Q}_k \geq 0$$

$$\mathbf{R}_k > 0$$

- Widely used in many real applications
- The “workhorse” of optimal control
- We know everything about it!
- Infinite variations and extensions!
- **Time Invariant** if: $\mathbf{A}_k = \mathbf{A}, \mathbf{B}_k = \mathbf{B}, \mathbf{Q}_k = \mathbf{Q}, \mathbf{R}_k = \mathbf{R}, \forall k$

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$$\mathbf{x}_{k+1} = \nabla_{\lambda} H(\mathbf{x}_k, \mathbf{u}_k, \lambda_{k+1}) = \mathbf{A}\mathbf{x}_k + \mathbf{B}\mathbf{u}_k$$

$$\lambda_k = \nabla_{\mathbf{x}} H(\mathbf{x}_k, \mathbf{u}_k, \lambda_{k+1}) = \mathbf{Q}\mathbf{x}_k + \mathbf{A}^T \lambda_{k+1}$$

$$\mathbf{u}_{k+1} = \underset{\mathbf{u}}{\operatorname{argmin}} H(\mathbf{x}_k, \mathbf{u}, \lambda_{k+1}) = -\mathbf{R}^{-1} \mathbf{B}^T \lambda_{k+1}$$

$$\lambda_K = \frac{\partial \ell_F}{\partial \mathbf{x}_K} = \mathbf{Q}_K \mathbf{x}_K$$

LQR via Shooting (2)

- 1 We start with a guess of the trajectory of $\mathbf{u}_{1:K-1}$
- 2 Forward pass (rollout) given $\mathbf{u}_{1:K-1}$ and \mathbf{x}_1 to get $\mathbf{x}_{1:K}$
- 3 Backward pass to compute $\lambda_{1:K}$ and $\mathbf{u}_{1:K-1}$ (or better $\Delta\mathbf{u}_{1:K-1}$)
- 4 Forward pass (rollout) with line search on $\Delta\mathbf{u}_{1:K-1}$ to get new $\mathbf{x}_{1:K}$
- 5 Go back to 3 until convergence

Double Integrator:

$$\mathbf{x} = \begin{bmatrix} \mathbf{q} \\ \dot{\mathbf{q}} \end{bmatrix}$$

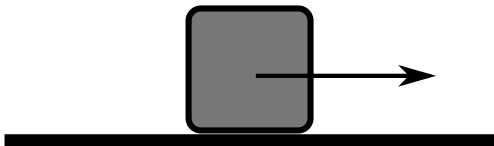
$$\mathbf{u} = K\ddot{\mathbf{q}}$$

where

$$\mathbf{q} = [x] \in \mathbb{R}$$

$$\dot{\mathbf{q}} = [v] = [\dot{x}] \in \mathbb{R}$$

$$\mathbf{u} = [K\ddot{x}] \in \mathbb{R}$$

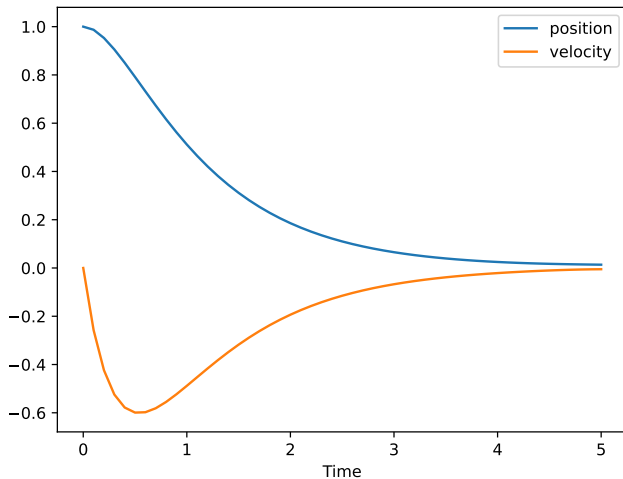


Discrete Dynamics:

$$\mathbf{x}_{k+1} = f_{\text{discrete}}(\mathbf{x}_k, \mathbf{u}_k)$$

$$= \begin{bmatrix} 1 & dt \\ 0 & 1 \end{bmatrix} \mathbf{x}_k + \begin{bmatrix} \frac{1}{2K} dt^2 \\ \frac{dt}{K} \end{bmatrix} \mathbf{u}_k$$

LQR via Shooting - Code Example



LQR as a Quadratic Programming Problem

Looking at the LQR problem, it really looks like a QP. Can we write it like one?

LQR as a Quadratic Programming Problem

Looking at the LQR problem, it really looks like a QP. Can we write it like one? Yes we can! We assume that \mathbf{x}_1 (initial conditions) is given, and define:

$$\mathbf{z} = \begin{bmatrix} \mathbf{u}_1 \\ \mathbf{x}_2 \\ \mathbf{u}_2 \\ \vdots \\ \mathbf{x}_K \end{bmatrix}, \mathbf{H} = \begin{bmatrix} \mathbf{R}_1 & & & & \\ & \mathbf{Q}_2 & & & \\ & & \mathbf{R}_2 & & \\ & & & \ddots & \\ & & & & \mathbf{Q}_K \end{bmatrix}$$

Now we can define:

$$\begin{aligned} \underset{\mathbf{z}}{\operatorname{argmin}} \mathcal{J}(\mathbf{z}) &= \frac{1}{2} \mathbf{z}^T \mathbf{H} \mathbf{z} \\ \text{s.t.} \quad & \text{“dynamics constraints”} \end{aligned}$$

LQR as a Quadratic Programming Problem (2)

For the dynamics constraints we have:

$$\underbrace{\begin{bmatrix} B_1 & -I & 0 & \dots \\ 0 & A_2 & B_2 & -I & 0 & \dots \\ & & \ddots & & & \\ & & & A_{K-1} & B_{K-1} & -I \end{bmatrix}}_G \begin{bmatrix} u_1 \\ x_2 \\ u_2 \\ \vdots \\ x_K \end{bmatrix} = \underbrace{\begin{bmatrix} -A_1 x_1 \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix}}_d$$

Now we have a full QP:

$$\begin{aligned} \underset{z}{\operatorname{argmin}} \mathcal{J}(z) &= \frac{1}{2} z^T H z \\ \text{s.t.} \quad & Gz - d = 0 \end{aligned}$$

LQR as a Quadratic Programming Problem (2)

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Can we solve this?

LQR as a Quadratic Programming Problem (2)

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Can we solve this? Of course we can! The Lagrangian is:

$$\mathcal{L}(\mathbf{z}, \boldsymbol{\lambda}) = \frac{1}{2} \mathbf{z}^T \mathbf{H} \mathbf{z} + \boldsymbol{\lambda}^T (\mathbf{G} \mathbf{z} - \mathbf{d})$$

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KKT Conditions:

$$\nabla_{\mathbf{z}} \mathcal{L} = \mathbf{H} \mathbf{z} + \mathbf{G}^T \boldsymbol{\lambda} = \mathbf{0}$$

$$\nabla_{\boldsymbol{\lambda}} \mathcal{L} = \mathbf{G} \mathbf{z} - \mathbf{d} = \mathbf{0}$$

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KKT System:

$$\begin{bmatrix} \mathbf{H} & \mathbf{G}^T \\ \mathbf{G} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{z} \\ \boldsymbol{\lambda} \end{bmatrix} = \begin{bmatrix} \mathbf{0} \\ \mathbf{d} \end{bmatrix}$$

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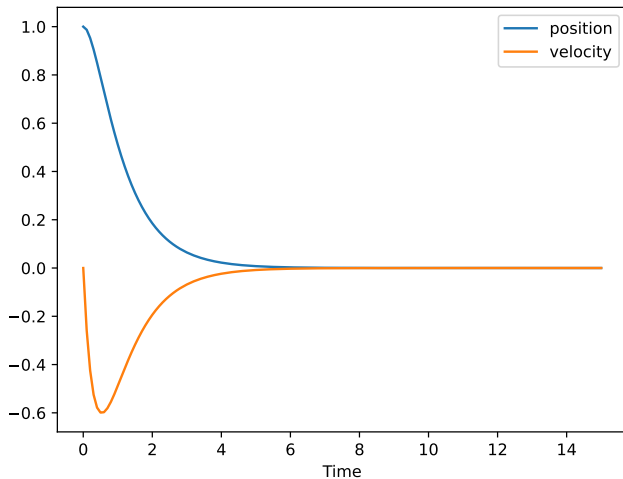
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We have the exact solution by solving ONE linear system!

LQR as a QP - Code Example



Let's write the KKT System in “detail” (for $K = 4$):

$$\begin{bmatrix}
 R_1 & 0 & \cdots & 0 & B_1^T & \cdots & 0 \\
 0 & Q_2 & 0 & \cdots & 0 & -I & A_2^T \\
 \vdots & 0 & R_2 & 0 & \cdots & 0 & B_2^T \\
 & \vdots & 0 & Q_3 & 0 & \cdots & -I \\
 & & \vdots & 0 & R_3 & 0 & \cdots \\
 0 & 0 & & \vdots & 0 & Q_4 & -I
 \end{bmatrix}
 \begin{bmatrix}
 u_1 \\
 x_2 \\
 u_2 \\
 x_3 \\
 u_3 \\
 x_4
 \end{bmatrix}
 =
 \begin{bmatrix}
 0 \\
 0 \\
 0 \\
 0 \\
 0 \\
 0
 \end{bmatrix}$$

$$\begin{bmatrix}
 B_1 & -I & 0 & 0 & \vdots & 0 & \ddots & \vdots \\
 \vdots & A_2 & B_2 & -I & 0 & \vdots & & \\
 0 & \cdots & A_3 & B_3 & -I & \cdots & 0 &
 \end{bmatrix}
 \begin{bmatrix}
 \lambda_2 \\
 \lambda_3 \\
 \lambda_4
 \end{bmatrix}
 =
 \begin{bmatrix}
 -A_1 x_1 \\
 0 \\
 0
 \end{bmatrix}$$

LQR via Riccati Recursion (2)

We first solve the last line of the upper block:

$$\mathbf{Q}_4 \mathbf{x}_4 - \lambda_4 = \mathbf{0} \Rightarrow \lambda_4 = \mathbf{Q}_K \mathbf{x}_4$$

Then we move one line up:

$$\begin{aligned} \mathbf{R}_3 \mathbf{u}_3 + \mathbf{B}_3^T \lambda_4 &= \mathbf{R}_3 \mathbf{u}_3 + \mathbf{B}_3^T \mathbf{Q}_K \mathbf{x}_4 = \mathbf{0} \\ \Rightarrow \mathbf{R}_3 \mathbf{u}_3 + \mathbf{B}_3^T \mathbf{Q}_K (\mathbf{A}_3 \mathbf{x}_3 + \mathbf{B}_3 \mathbf{u}_3) &= \mathbf{0} \\ \Rightarrow \mathbf{u}_3 &= - \underbrace{(\mathbf{R}_3 + \mathbf{B}_3^T \mathbf{Q}_K \mathbf{B}_3)^{-1} \mathbf{B}_3^T \mathbf{Q}_K \mathbf{A}_3}_{\mathbf{K}_3} \mathbf{x}_3 \end{aligned}$$

One more line:

$$\begin{aligned} \mathbf{Q}_3 \mathbf{x}_3 - \lambda_3 + \mathbf{A}_3^T \lambda_4 &= \mathbf{0} \Rightarrow \mathbf{Q}_3 \mathbf{x}_3 - \lambda_3 + \mathbf{A}_3^T \mathbf{Q}_K \mathbf{x}_4 = \mathbf{0} \\ \Rightarrow \mathbf{Q}_3 \mathbf{x}_3 - \lambda_3 + \mathbf{A}_3^T \mathbf{Q}_K (\mathbf{A}_3 \mathbf{x}_3 + \mathbf{B}_3 \mathbf{u}_3) &= \mathbf{0} \\ \Rightarrow \lambda_3 &= \underbrace{\left(\mathbf{Q}_3 + \mathbf{A}_3^T \mathbf{Q}_K (\mathbf{A}_3 - \mathbf{B}_3 \mathbf{K}_3) \right)}_{\mathbf{P}_3} \mathbf{x}_3 \end{aligned}$$

The recursion is now revealed:

$$\mathbf{P}_K = \mathbf{Q}_K$$

$$\mathbf{K}_k = (\mathbf{R}_k + \mathbf{B}_k^T \mathbf{P}_{k+1} \mathbf{B}_k)^{-1} \mathbf{B}_k^T \mathbf{P}_{k+1} \mathbf{A}_k$$

$$\mathbf{P}_k = \mathbf{Q}_k + \mathbf{A}_k^T \mathbf{P}_{k+1} (\mathbf{A}_k - \mathbf{B}_k \mathbf{K}_k)$$

- QP has complexity $O(K^3(N + M)^3)$
- Riccati recursion has complexity $O(K(N + M)^3)$
- The Riccati recursion is the optimal exploitation of the sparsity of the KKT system!

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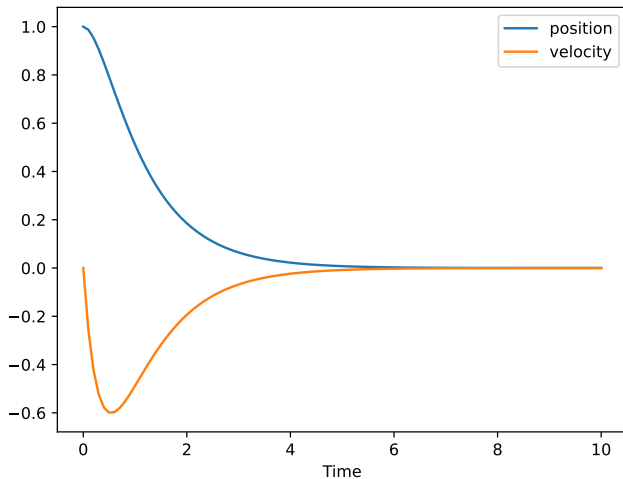
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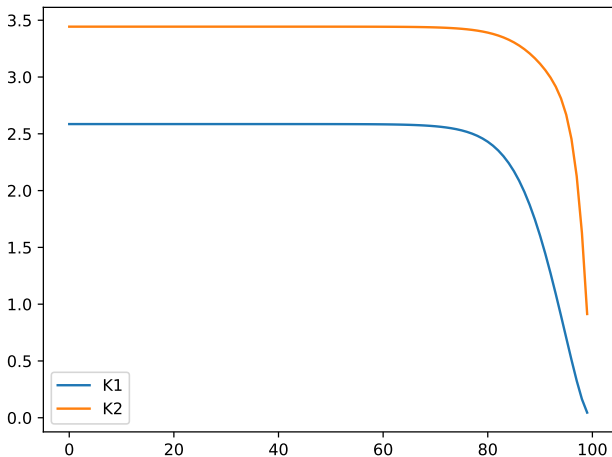
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- QP has complexity $O(K^3(N + M)^3)$
- Riccati recursion has complexity $O(K(N + M)^3)$
- The Riccati recursion is the optimal exploitation of the sparsity of the KKT system!
- **We also get for free a feedback policy: $\mathbf{u} = -\mathbf{K}\mathbf{x}$!**

LQR via Riccati Recursion - Code Example



LQR via Riccati Recursion - Code Example (2)



Infinite Horizon LQR

- Only possible for the time invariant linear system!
- \mathbf{K}_k and \mathbf{P}_k matrices converge to constant values!
- How can we find those values?
 - 1 We do the Riccati recursion long enough until they converge!

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 - 2 Let's have a look at the recursion equations again:

$$K_k = (R_k + B_k^T P_{k+1} B_k)^{-1} B_k^T P_{k+1} A_k$$
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- 3 At the limit, we need $P_k = P_{k+1}$ and $K_k = K_{k+1}$
 - 4 We solve a root finding problem (e.g. with Newton's method)!
- The recursion version is basically the Fixed Point method!

How can we know LQR will always find a solution?

- We get to choose \mathbf{Q}_k and \mathbf{R}_k
- We cannot choose \mathbf{A}_k and \mathbf{B}_k
- What can we do?

How can we know LQR will always find a solution?

- We get to choose Q_k and R_k
- We cannot choose A_k and B_k
- What can we do? **Let's start with the time invariant case:**

$$\begin{aligned}
 \mathbf{x}_K &= \mathbf{A}\mathbf{x}_{K-1} + \mathbf{B}\mathbf{u}_{K-1} \\
 &= \mathbf{A}(\mathbf{A}\mathbf{x}_{K-2} + \mathbf{B}\mathbf{u}_{K-2}) + \mathbf{B}\mathbf{u}_{K-1} \\
 &\quad \vdots \\
 &= \mathbf{A}^K \mathbf{x}_1 + \mathbf{A}^{K-1} \mathbf{B} \mathbf{u}_1 + \mathbf{A}^{K-2} \mathbf{B} \mathbf{u}_2 + \cdots + \mathbf{B} \mathbf{u}_{K-1} \\
 &= \underbrace{\begin{bmatrix} \mathbf{B} & \mathbf{A}^2 \mathbf{B} & \cdots & \mathbf{A}^{K-1} \mathbf{B} \end{bmatrix}}_C \begin{bmatrix} \mathbf{u}_{K-1} \\ \mathbf{u}_{K-2} \\ \vdots \\ \mathbf{u}_1 \end{bmatrix} + \mathbf{A}^K \mathbf{x}_1 = \mathbf{0}
 \end{aligned}$$

- This is equivalent to a least squares problem:

$$\begin{bmatrix} \mathbf{u}_{K-1} \\ \mathbf{u}_{K-2} \\ \vdots \\ \mathbf{u}_1 \end{bmatrix} = [\mathbf{C}^T (\mathbf{C}\mathbf{C}^T)^{-1}] (\mathbf{x}_K - \mathbf{A}^K \mathbf{x}_1)$$

- For $\mathbf{C}\mathbf{C}^T$ to be invertible, we need $\text{rank}(\mathbf{C}) = N$
- \mathbf{C} is usually called the “**Controllability Matrix**”
- We follow the same procedure for the time variant case
- In the time invariant case we can stop at N steps (because of the Cayley-Hamilton theorem)

Thank you

- Any Questions?
- Office Hours:
 - Tue & Thu (09:00-11:00)
 - 24/7 by email (costashatz@upatras.gr, subject: *ECE_RSII_AM*)
- Material and Announcements



Laboratory of Automation & Robotics