

# STRING MODULES, NON-KISSING COMPLEXES AND WHATNOT (TEMPORARY TITLE)

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ABSTRACT. This is a work in progress.

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## 1. INTRODUCTION

**Conventions.** If  $Q$  is a quiver and  $k$  is a field, we denote by  $kQ$  the path algebra of  $Q$ . Arrows of a quiver  $Q$  are composed from right to left: if  $1 \xrightarrow{a} 2 \xrightarrow{b} 3$  is a quiver, then  $ba$  is a path while  $ab$  is not. Modules over an algebra are assumed to be *left* modules. By these conventions, a module over  $kQ$  is equivalent to a representation of  $Q$ .

## 2. STRING MODULES

**2.1. String algebras.** String algebras are a class of algebras defined by generators and particularly nice relations. The study of their representations goes back to [4]; we will follow the general framework established in [3].

Let  $k$  be an algebraically closed field, and let  $Q$  be a quiver. While this is not necessary for the theory, it suffices our purposes to assume that  $Q$  is finite.

**Definition 2.1.** A *string algebra* is an algebra of the form  $kQ/I$ , where  $I$  is an ideal generated by paths of length at least two in  $Q$  and such that  $kQ/I$  is finite-dimensional, and satisfying the following properties:

- (1) for any vertex  $i$  of  $Q$ , there are at most two arrows starting at  $i$  and at most two arrows ending in  $i$ ;
- (2) for any arrow  $\beta$  of  $Q$ , there is at most one arrow  $\gamma$  such that  $\beta\gamma \notin I$ ;

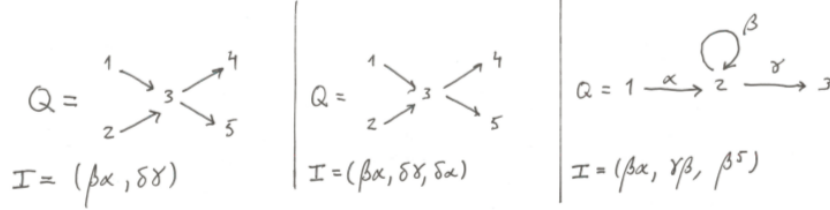
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- (3) for any arrow  $\beta$  of  $Q$ , there is at most one arrow  $\alpha$  such that  $\alpha\beta \notin I$ .

**Example 2.2.**



**Definition 2.3** ([?]). A string algebra  $kQ/I$  is a *gentle algebra* if

- (1)  $I$  is generated by paths of length exactly two;
- (2) for any arrow  $\beta$  of  $Q$ , there is at most one arrow  $\gamma$  such that  $t(\gamma) = s(\beta)$  and  $\beta\gamma \in I$ ;
- (3) for any arrow  $\beta$  of  $Q$ , there is at most one arrow  $\alpha$  such that  $s(\alpha) = t(\beta)$  and  $\alpha\beta \in I$ .

**Example 2.4.** In Example 2.2, only the first algebra is gentle.

String algebras enjoy a particularly nice representation theory: they are tame algebras (as proved in [?]) and their indecomposable representations are completely classified. We will make heavy use of this classification, which we recall in the next section

**2.2. String and band modules.** For any arrow  $\beta$  of any quiver  $Q$ , define a formal inverse  $\beta^{-1}$  with the properties that  $s(\beta^{-1}) = t(\beta)$ ,  $t(\beta^{-1}) = s(\beta)$ ,  $\beta^{-1}\beta = e_{s(\beta)}$  and  $\beta\beta^{-1} = e_{t(\beta)}$ , where  $e_i$  is the path of length zero starting and ending at the vertex  $i$ . Furthermore, let  $(\beta^{-1})^{-1}$ ; thus  $(-)^{-1}$  is an involution on the set of arrows and formal inverses of arrows of  $Q$ . We extend this involution to the set of all paths involving arrows or their formal inverses, by further setting  $(e_i)^{-1} = e_i$  for all vertices  $i$  of  $Q$ .

**Definition 2.5.** Let  $A = kQ/I$  be a string algebra.

- (1) A *string* for  $A$  is a word of the form

$$w = \beta_1^{\varepsilon_1} \beta_2^{\varepsilon_2} \cdots \beta_n^{\varepsilon_n},$$

where

- each  $\beta_i$  is an arrow in  $Q$  and each  $\varepsilon_i$  is an element of  $\{1, -1\}$ ;
- for each  $i \in \{1, 2, \dots, n-1\}$ ,  $t(\beta_{i+1}^{\varepsilon_{i+1}}) = s(\beta_i^{\varepsilon_i})$ ;
- if  $p$  is a path in  $Q$  such that  $p$  or  $p^{-1}$  appears as a subword of  $w$ , then  $p$  is not an element of  $I$ ; and
- $w$  is reduced, in the sense that no subword of the form  $\beta\beta^{-1}$  or  $\beta^{-1}\beta$  appears in  $w$ , with  $\beta$  an arrow of  $Q$ .

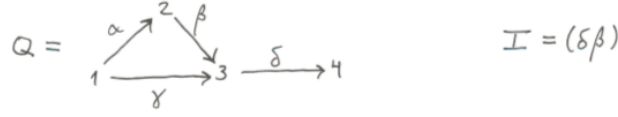
The integer  $n$  is called the *length* of the string  $w$ .

For each vertex  $i$  of  $Q$ , there is also a *string of length zero*, denoted by  $e_i$ , that starts and ends at  $i$ .

- (2) A *band* for  $A$  is a string  $b$  of length at least one such that

- $s(b) = t(b)$ ;
- all powers of  $b$  are strings; and
- $b$  is not itself a power of a strictly smaller string.

**Example 2.6.** For the algebra



some strings are  $e_1, e_2, e_3, e_4, \alpha^\pm, \beta^\pm, \gamma^\pm, \delta^\pm, \delta\gamma, \gamma^{-1}\beta, \dots$

The only bands are  $\gamma^{-1}\beta\alpha, \beta\alpha\gamma^{-1}, \alpha\gamma^{-1}\beta$  and their inverses.

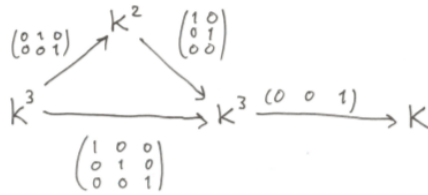
**Definition 2.7.** Let  $w = \beta_1^{\varepsilon_1} \beta_2^{\varepsilon_2} \dots \beta_n^{\varepsilon_n}$  be a string for a string algebra  $A = kQ/I$ . The *string module*  $M(w)$  is the  $A$ -module defined as a representation of  $Q$  as follows:

- Let  $i_0 = t(\beta_1^{\varepsilon_1})$ , and  $i_m = s(\beta_m^{\varepsilon_m})$  for each  $m \in \{1, \dots, n\}$ .
- For each vertex  $j$  of  $Q$ , let  $M(w)_j$  be the vector space with basis given by  $\{v_m \mid i_m = j\}$ .
- For each arrow  $\gamma$  of  $Q$ , the linear map  $M(w)_\gamma : M(w)_{s(\gamma)} \rightarrow M(w)_{t(\gamma)}$  is defined on the basis of  $M(w)_{s(\gamma)}$  by

$$M(w)_\gamma(v_m) = \begin{cases} v_{m-1} & \text{if } \beta_m = \gamma \text{ and } \varepsilon_m = 1; \\ v_{m+1} & \text{if } \beta_{m+1} = \gamma \text{ and } \varepsilon_{m+1} = -1; \\ 0 & \text{else.} \end{cases}$$

It follows from the definition that for any string  $w$ , the string modules  $M(w)$  and  $M(w^{-1})$  are isomorphic.

**Example 2.8.** For the algebra given in Example 2.6, let  $w = \delta\gamma\alpha^{-1}\beta^{-1}\gamma\alpha^{-1}\beta^{-1}\gamma$ . Then  $M(w)$  is given by



**Notation 2.9.** For convenience, even though 0 is not a string by our definition, we will define  $M(0)$  to be the zero module.

**Definition 2.10.** Let  $b = \beta_1^{\varepsilon_1} \beta_2^{\varepsilon_2} \dots \beta_n^{\varepsilon_n}$  be a band for a string algebra  $A = kQ/I$ . Moreover, let  $\lambda \in k^*$ , and let  $d \in \mathbb{N}_{>0}$ . The *band module*  $M(b, \lambda, d)$  is defined as follows:

- For each vertex  $j$  of  $Q$ , the vector space  $M(b, \lambda, d)_j$  is equal to  $k^d$ .
- For each arrow  $\gamma$  of  $Q$  different from  $\beta_1$ , the linear map  $M(b, \lambda, d)_\gamma$  is the identity.

- The linear map  $M(b, \lambda, d)_{\beta_1}$  is equal to  $J_d(\lambda)^{\varepsilon_1}$ , where

$$J_d(\lambda) = \begin{pmatrix} \lambda & 0 & 0 & \cdots & 0 \\ 1 & \lambda & 0 & \cdots & 0 \\ 0 & 1 & \lambda & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & \lambda \end{pmatrix}$$

is the  $d \times d$  Jordan block of type  $\lambda$ .

It follows from the definition that the band modules  $M(b, \lambda, d)$  and  $M(b^{-1}, \lambda^{-1}, d)$  are isomorphic. Moreover, if  $b$  and  $b'$  are two bands that are *cyclically equivalent*, in the sense that one is obtained from the other by cyclically permuting the arrows that constitute it, then we also have that  $M(b, \lambda, d)$  and  $M(b', \lambda, d)$  are isomorphic.

**Example 2.11.** For the algebra defined in Example 2.6, let  $b = \gamma^{-1}\beta\alpha$ . Then  $b$  is a band, and the following are band modules:

$$M(b, \lambda, 1) = \begin{array}{ccccc} & & k & & \\ & \nearrow 1 & & \nwarrow 1 & \\ k & & & & k & \xrightarrow{0} & 0 \\ & \xrightarrow{\lambda^{-1}} & & & \end{array}$$

$$M(b, \lambda, 3) = \begin{array}{ccccc} & & k^3 & & \\ & \nearrow \text{Id} & & \nwarrow \text{Id} & \\ k^3 & & & & k^3 & \xrightarrow{0} & 0 \\ & \xrightarrow{\begin{pmatrix} \lambda^{-1} & 0 & 0 \\ 1 & \lambda^{-1} & 0 \\ 0 & 1 & \lambda^{-1} \end{pmatrix}} & & \end{array}$$

**Theorem 2.12** ([3], p.161). *Let  $A$  be a string algebra. Then the string and band modules over  $A$  form a complete list of indecomposable  $A$ -modules, up to isomorphism. Moreover,*

- *A string module is never isomorphic to a band module.*
- *Two string modules  $M(w)$  and  $M(w')$  are isomorphic if and only if  $w' = w^{\pm 1}$ .*
- *Two band modules  $M(b, \lambda, d)$  and  $M(b', \lambda', d')$  are isomorphic if and only if  $d = d'$ , and either*
  - $b$  is cyclically equivalent to  $b'$ , and  $\lambda = \lambda'$ , or*
  - $b^{-1}$  is cyclically equivalent to  $b'$ , and  $\lambda^{-1} = \lambda'$ .*

**2.3. Auslander-Reiten translation of string modules.** The following notation, still taken from [3], will be useful for dealing with string modules.

For a given string  $w = \beta_1^{\varepsilon_1} \beta_2^{\varepsilon_2} \cdots \beta_n^{\varepsilon_n}$ , we draw  $w$  as follows:

- draw all arrows  $\beta_1, \dots, \beta_n$  from left to right;
- draw all arrows pointing downwards.

Moreover, the strings  $e_i$  of length zero are depicted simply as  $i$ .

This is best understood on an example.

**Example 2.13.** For the algebra defined in Example 2.6, the string  $w = \delta\gamma\alpha^{-1}\beta^{-1}\gamma\alpha^{-1}\beta^{-1}\gamma$  is drawn as follows:



There are some immediate advantages for this depiction of strings. For instance, the socle of  $M(w)$  is the direct sum of the simple modules corresponding to the local minima in the picture of  $w$ ; similarly, the top of  $M(w)$  is the direct sum of the simple modules corresponding to the local maxima.

**Definition 2.14.** (1) A string  $w$  *starts (or ends) on a peak* if there is no arrow  $\beta$  such that  $w\beta$  (or  $\beta^{-1}w$ , respectively) is a string.  
 (2) A string  $w$  *starts (or ends) in a deep* if there is no arrow  $\beta$  such that  $w\beta^{-1}$  (or  $\beta w$ , respectively) is a string.

**Remark 2.15.** Definition 2.14 is best understood by drawing strings. Starting (or ending) on a peak means that one cannot add an arrow at the start (or the end) of  $w$  such that the starting point (or ending point) of the new string would be higher in the picture than that of  $w$ . The same applies when replacing “on a peak” by “in a deep” and “higher” by “lower”.

**Example 2.16.** The string in Example 2.13 starts and ends on a peak, and also ends in a deep.

**Definition 2.17.** (1) Let  $w$  be a string that does not start on a peak. Let  $\beta_0, \beta_1, \dots, \beta_r$  be arrows such that  $w_h := w\beta_0\beta_1^{-1} \dots \beta_r^{-1}$  starts in a deep. We say that  $w_h$  is obtained from  $w$  by *adding a hook* at the start of  $w$ .  
 (2) Let  $w$  be a string that does not end on a peak. Let  $\beta_0, \beta_1, \dots, \beta_r$  be arrows such that  ${}_hw := \beta_r\beta_{r-1} \dots \beta_1\beta_0^{-1}w$  ends in a deep. We say that  ${}_hw$  is obtained from  $w$  by *adding a hook* at the end of  $w$ .

**Remark 2.18.** (1) Unless  $w$  has length zero, the strings  ${}_hw$  and  $w_h$ , when they are defined, are uniquely determined. If  $w$  has length zero, then there may be two possible choices for  ${}_hw$  and  $w_h$ .  
 (2) It is possible that  $r = 0$ .

The dual of a hook is a cohook.

**Definition 2.19.** (1) Let  $w$  be a string that does not start in a deep. Let  $\beta_0, \beta_1, \dots, \beta_r$  be arrows such that  $w_c := w\beta_0^{-1}\beta_1 \dots \beta_r$  starts on a peak. We say that  $w_c$  is obtained from  $w$  by *adding a cohook* at the start of  $w$ .  
 (2) Let  $w$  be a string that does not end in a deep. Let  $\beta_0, \beta_1, \dots, \beta_r$  be arrows such that  ${}_cw := \beta_r^{-1}\beta_{r-1}^{-1} \dots \beta_1^{-1}\beta_0w$  ends on a peak. We say that  ${}_cw$  is obtained from  $w$  by *adding a cohook* at the end of  $w$ .

**Remark 2.20.** (1) As when adding hooks, if  $w$  has length zero, then the strings  ${}_cw$  and  $w_c$ , when they are defined, are uniquely determined. If  $w$  has length zero, then there may be two possible choices for  ${}_cw$  and  $w_c$ .

- (2) It is possible that  $r = 0$ .

**Example 2.21.** In pictures, hooks always look like



while cohooks look like



**Example 2.22.** In Example 2.6, the string



does not start on a deep. Thus



**Example 2.23.** Still in Example 2.6, the string  $w = \delta$  does not start on a peak, and we have



**Lemma 2.24** (and Definition). *Let  $w$  be a string for a string algebra  $A = kQ/I$ .*

- (1) *If  $w$  starts on a peak and is not a path in  $Q$ , then there exists a unique string  $w'$  such that  $w'$  does not start in a deep, and  $w = w'_c$ . We denote  $w'$  by  $w_{c^{-1}}$ , and we say that is obtained from  $w$  by removing a cohook at the start of  $w$ . If  $w$  is a path in  $Q$ , then we put  $w_{c^{-1}} = 0$ .*
- (2) *If  $w$  ends on a peak and is not the inverse of a path in  $Q$ , then there exists a unique string  $w'$  such that  $w'$  does not end in a deep, and  $w = {}_c w'$ . We denote  $w'$  by  ${}_{c^{-1}}w$ , and we say that is obtained from  $w$  by removing a*

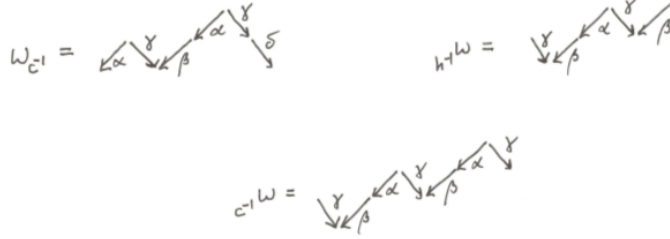
cohook at the end of  $w$ . If  $w$  is the inverse of a path in  $Q$ , then we put  $c^{-1}w = 0$ .

- (3) If  $w$  starts in a deep and is not the inverse of a path in  $Q$ , then there exists a unique string  $w'$  such that  $w'$  does not start on a peak, and  $w = w'_h$ . We denote  $w'$  by  $w_{h^{-1}}$ , and we say that is obtained from  $w$  by removing a hhook at the start of  $w$ . If  $w$  is the inverse of a path in  $Q$ , then we put  $w_{h^{-1}} = 0$ .
- (4) If  $w$  ends in a deep and is not a path in  $Q$ , then there exists a unique string  $w'$  such that  $w'$  does not end on a peak, and  $w = {}_h w'$ . We denote  $w'$  by  ${}_{h^{-1}}w$ , and we say that is obtained from  $w$  by removing a hook at the end of  $w$ . If  $w$  is a path in  $Q$ , then we put  ${}_{h^{-1}}w = 0$ .

**Example 2.25.** For the algebra defined in Example 2.6 and the string



we have



**Theorem 2.26** ([3]). Let  $A = kQ/I$  be a string algebra, and let  $w$  be a string. Then  $\tau(M(w)) = M(w')$ , where  $w'$  is obtained by following the table below.

	$w$ does not start in a deep	$w$ starts in a deep
$w$ does not end in a deep	$cw_c$	$({}_c w)_{h^{-1}}$
$w$ ends in a deep	${}_{h^{-1}}(w_c)$	${}_{h^{-1}}w_{h^{-1}}$

Dually,  $\tau^{-1}(M(w)) = M(w'')$ , where  $w''$  is given in the table below.

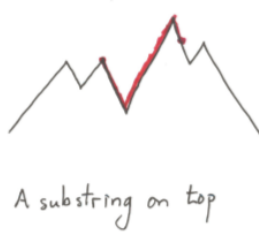
	$w$ does not start on a peak	$w$ starts on a peak
$w$ does not end on a peak	${}_h w_h$	$({}_h w)_{c^{-1}}$
$w$ ends on a peak	$c^{-1}(w_h)$	$c^{-1}w_{c^{-1}}$

Finally, applying  $\tau$  or  $\tau^{-1}$  to a band module gives the same band module.

#### 2.4. Morphisms between string modules.

**Definition 2.27.** Let  $A = kQ/I$  be a string algebra, and let  $w = \beta_1^{\varepsilon_1} \cdots \beta_n^{\varepsilon_n}$  be a string for  $A$ .

- (1) A *substring* of  $w$  is a string of the form  $\beta_i^{\varepsilon_i} \cdots \beta_j^{\varepsilon_j}$ , with  $1 \leq i \leq j \leq n$ . A string  $e_i$  of length zero is also considered to be a substring of  $w$  if one of the  $\beta_i$  starts or ends in  $i$ .
- (2) A substring  $w' = \beta_i^{\varepsilon_i} \cdots \beta_j^{\varepsilon_j}$  is said to be *on top* of  $w$  if
  - $i = 1$  or  $\varepsilon_{i-1} = 1$ , and
  - $j = n$  or  $\varepsilon_{j+1} = -1$ .
 If  $w'$  is of length zero, we say that it is on top of  $w$  if it corresponds to a local maximum in the picture of  $w$ .
- (3) A substring  $w' = \beta_i^{\varepsilon_i} \cdots \beta_j^{\varepsilon_j}$  is said to be *at the bottom* of  $w$  if
  - $i = 1$  or  $\varepsilon_{i-1} = -1$ , and
  - $j = n$  or  $\varepsilon_{j+1} = 1$ .
 If  $w'$  is of length zero, we say that it is on top of  $w$  if it corresponds to a local minimum in the picture of  $w$ .



**Example 2.28.**

**Proposition 2.29.** *Let  $w$  and  $w'$  be two strings for a string algebra  $A$ . Then the dimension of the  $k$ -vector space  $\text{Hom}_A(M(w), M(w'))$  is equal to*

$$\sum_{\xi \text{ substring of } w} (\# \text{ of substrings equal to } \xi \text{ or } \xi^{-1} \text{ on top of } w) \cdot (\# \text{ of substrings equal to } \xi \text{ or } \xi^{-1} \text{ at the bottom of } w)$$

PROOF. Reference?

□

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