

Q1.  $G1 : S \rightarrow ABS$

$S \rightarrow AB$

$A \rightarrow Aa$

$A \rightarrow a$

$B \rightarrow bA$

aabaab

$S \rightarrow ABS \rightarrow aABS \rightarrow aaBS \rightarrow aaBAB \rightarrow aabAAB \rightarrow aabaAB \rightarrow aabaaB \rightarrow aabaabA$

Will not be accepted by  $G1$  because none of the non-terminal can't step to an element of the set  $\{ \epsilon \cup \text{epsilon star} \}$  that leads to the end with terminal 'b'

aaaaba

this string will be accepted  $G1$

aabbba

this string will not be accepted by string  $G1$  because the non-terminal can't step to an element of the set  $( \epsilon \cup \text{epsilon star} )$  that lead to the terminal that have a pair 'bb'.

abaaba

this string could be accepted by  $G1$

Q2

Let  $R$  is RegExp

$T$  is Term

$F$  is Factor

$L$  is Letter

$G2 : R \rightarrow T \mid T + R \mid (R) \mid \epsilon$

$T \rightarrow F \mid TF$

$F \rightarrow L \mid F^*$

$L \rightarrow c \mid 1 \mid R$  ( $c$  is an element of epsilon)

$$R \rightarrow T + R \rightarrow T.F + R \rightarrow T.F^* + R$$

$$TR^* + R \rightarrow T(R)^* + R \rightarrow T(T+R)^* + R$$

$$T(T+T+R) + R \rightarrow T(T.F + T+R)^* + R$$

$$T(FF + T + R) + R \rightarrow R \rightarrow T(LF + T + R)^* + R$$

$$T(bF + T + R)^* + R \rightarrow T(bL + T+R)^* + R$$

$$T(ba + T + R)^* + R \rightarrow T(ba + F + R)^* + R$$

$$T(ba + L + R)^* + R \rightarrow T(ba + b + R)^* + R$$

$$T(ba + b + T)^* + R \rightarrow T(ba + b + TF)^* + R$$

$$T(ba + b + FF)^* + R \rightarrow T(ba + b + LF)^* + R$$

$$T(ba + b + LL)^* + R \rightarrow T(ba + b + bL)^* + R$$

$$T(ba + b + bb)^* + R \rightarrow FF^*(ba + b + bb)^* + R$$

$$LF^*(ba + b + bb)^* + R \rightarrow LL^*(ba + b + bb)^* + R$$

$$aL^*(ba + b + bb)^* + R \rightarrow ab^*(ba + b + bb)^* + R$$

$$ab^*(ba + b + bb) + T + R \rightarrow ab^*(ba + b + bb)^* + T + T$$

$$ab^*(ba + b + bb) + TF + T \rightarrow ab^*(ba + b + bb)^* + T + T$$

$$ab^*(ba + b + bb)^* + T L L + T$$

$$ab^*(ba + b + bb)^* + T ba + T$$

$$ab(ba + b + bb)^* + Fba + F$$

$$ab(ba + b + bb)^* + Lba + L$$

$$ab(ba + b + bb)^* + aba + 1$$

Q.3

$$G3 : S \rightarrow X|Y|Z$$

$$X \rightarrow C D$$

$$C \rightarrow aE$$

$$E \rightarrow \text{epsilon} | bE$$

$$D \rightarrow \text{epsilon} | FD$$

$$F \rightarrow G|H|I$$
$$G \rightarrow BA$$
$$H \rightarrow B$$
$$I \rightarrow B$$
$$Y \rightarrow ABA$$
$$Z \rightarrow \text{epsilon}$$
$$A \rightarrow a$$
$$B \rightarrow b$$

Q4

String that are accepted are : b, aba, aabaa, aaabaaa, bb, abba, aabbaa, etc. The CFG will look like:

$$S \rightarrow aSa \mid B$$
$$B \rightarrow bB'$$
$$B' \rightarrow bB' \mid \text{epsilon}$$

To create a CNF we will need to

1. Remove NULL rules
2. Remove unit rules
3. Remove more than 1 terminal on the Rhts
4. Remove more than 2 non-terminal on the Rhts