Program Structures and Algorithms Spring 2022 Assignment 3 (WQUPC)

Name: Vraj Himanshu Reshamdalal

(NUID): 002927484

Task:

Step 1:

- (a) Implement height-weighted Quick Union with Path Compression. For this, you will flesh out the class UF_HWQUPC. All you have to do is to fill in the sections marked with // TO BE IMPLEMENTED ... // ...END IMPLEMENTATION.
- (b) Check that the unit tests for this class all work. You must show "green" test results in your submission (screenshot is OK).

Step 2:

Using your implementation of UF_HWQUPC, develop a UF ("union-find") client that takes an integer value n from the command line to determine the number of "sites." Then generates random pairs of integers between 0 and n-1, calling connected() to determine if they are connected and union() if not. Loop until all sites are connected then print the number of connections generated. Package your program as a static method count() that takes n as the argument and returns the number of connections; and a main() that takes n from the command line, calls count() and prints the returned value. If you prefer, you can create a main program that doesn't require any input and runs the experiment for a fixed set of n values. Show evidence of your run(s).

Step 3:

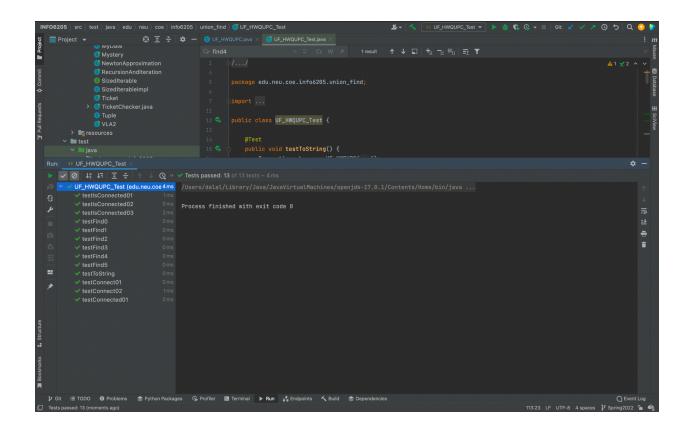
Determine the relationship between the number of objects (n) and the number of pairs (m) generated to accomplish this (i.e. to reduce the number of components from n to 1). Justify your conclusion in terms of your observations and what you think might be going on.

Code Stubs:

Part 1: Completed

```
public int find(int p) {
    validate(p);
   int root = p;
    // FIXME
    while(root != getParent(root)) {
        if(this.pathCompression == true) {
            this.doPathCompression(root);
        }
        root = getParent(root);
    // END
    return root;
```

```
private void doPathCompression(int i) {
    // FIXME update parent to value of grandparent
    parent[i] = parent[parent[i]];
    // END
private void mergeComponents(int i, int j) {
    // FIXME make shorter root point to taller one
    int t1 = find(i);
    int t2 = find(j);
   if(t1 == t2) {
        return;
    }
    if(height[t1] < height[t2]) {</pre>
        parent[t1] = t2;
        height[t2] = height[t2] + height[t1];
    } else {
        parent[t2] = t1;
        height[t1] = height[t1] + height[t2];
    }
    // END
```



Part2: Completed

```
| Notice | Project | Proje
```

Part3: Completed

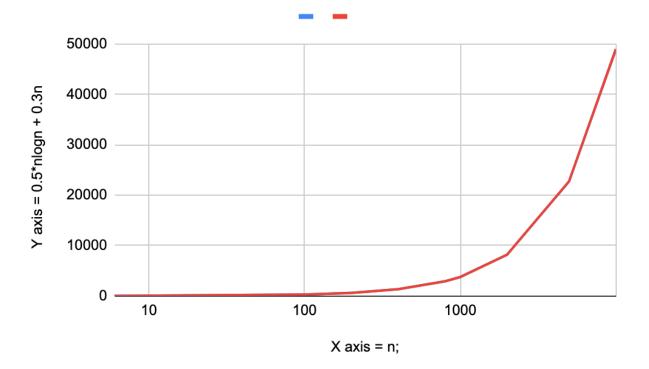
Part 3 output:

```
### Project * © I ÷ 0 - Uniformation of Historic Sendon Para | Project Para | Disput Sendon Para | Project Para | Disput Sendon Para | Project Para | Disput Sendon Disput Sendon Disput Sendon Disput Sendon Disput Sendon Disput Sendon Disput
```

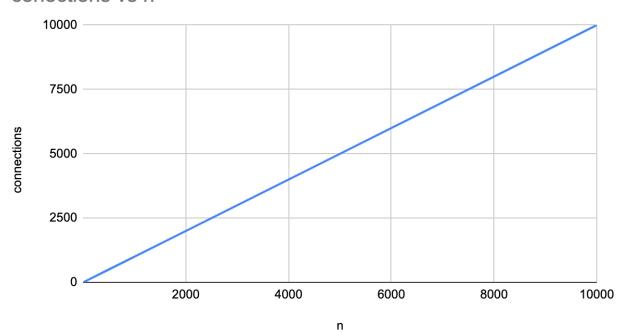
Evidence/graph:

connections	n	0.5*nlogn + 0.29n	m
5	6	7.115278408	8
99	100	259.2585093	261
199	200	587.8317367	590
399	400	1314.292909	1319
799	800	2905.844691	2907
999	1000	3743.877639	3737
1999	2000	8180.90246	8189
4999	5000	22742.98298	22680
9999	10000	48951.70186	48958

Graph



conections vs n



Relationship conclusion:

Based on the calculations and the table shown above we can conclude that, m=number of pairs generated connections=number of connections c = constant

Connections =
$$n-1$$

 $m = 0.5*nlogn + 0.3n OR m = 0.5*nlogn + c$

Justification:

As per the Erdös-Renyi model - the number of pairs generated to get one component is $\sim 1/2N$ In N.

References:

- 1. PSA text book Algorithhms 4th Edition by Robert Sedgewick, Kevin Wayne
- 2. Slides from class
- 3. https://en.wikipedia.org/wiki/Erd%C5%91s%E2%80%93R%C3%A9nyi_model