





• La nivel de bit (Bit level)

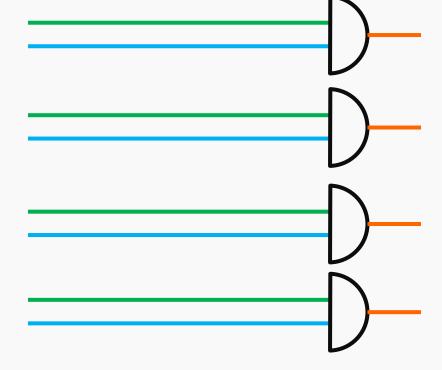
• La nivel de instrucțiune (Instruction level)

La nivel de task (Task parallelism)



- Bit level
 - C = A & B
- Instruction level

Task parallelism





Bit level

Instruction level

$$C[] = A[] + B[]$$

Task parallelism

Adunarea a doi vectori



Bit level

Instruction level

Task parallelism

load Z		A = B + C
store H	•	store A
H=H*100		G[1] = G[2] + G[3]
H++		store G[1]
store H		G[4] = G[5] + G[6]

- Multi-Tasking (pot comunica şi procesele)
- Multi-Threading



Bit level

Instruction level

- Task parallelism
 - Multi-Tasking (pot comunica şi procesele)
 - Multi-Threading
 Cum pot comunica două procese?





Notații pseudocod

```
co S1 || S2 || ... || Sn oc

Ex.1:

x=0; y=0;

co x=x+1 || y=y+1 oc

z=x+y;
```



Notații pseudocod

```
co [cuantificator] {Sj}
Ex. 2:
    co [j=1 to n] {a[j]=0;
b[j]=0;}
```



POSIX threads



POSIX threads

```
pthread t thread;
pthread create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
pthread join(thread, NULL);
```



Compilare pthread

gcc –o executabil cod.c –lpthread –lrt

#include<pthread.h>
#include<semaphore.h>



```
pthread t thread;
pthread create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
                     Acest element reprezintă thread-ul.
                     Este un thread handle
pthread join(thread, NULL);
```



```
pthread t thread;
pthread create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
                 Prin acest parametru am putea să facem
                 recomandări sistemului de operare.
                 Ex: să folosească anumite core-uri.
```

pthread join(thread, NULL);



```
pthread t thread;
pthread create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
               Când se crează thread-ul va porni de
               la funcția dată ca parametru.
```

pthread_join(thread, NULL);



```
pthread t thread;
pthread create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
               Aşa trimitem date thread-ului
pthread_join(thread, NULL);
```



```
pthread t thread;
pthread create(&thread, NULL, threadFunction, arg);
void * threadFunction(void* arg)
                Astfel se pot extrage date din
                thread
pthread join(thread, NULL);
```





```
int main(int argc, char **argv)
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
 return 0;
```



```
int main(int argc, char **argv)
 int i, P=2;
 pthread_t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread_join(tid[i], NULL);
 return 0;
```



```
int main(int argc, char **argv)
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
 return 0;
```



```
int main(int argc, char **argv)
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
 return 0;
```



```
int main(int argc, char **argv)
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
 return 0;
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                             a = a + 1;
                                             return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
  return 0;
```



void* threadFunction(void *arg)

```
int main(int argc, char **argv)
                                             a = a + 1;
                                             return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
  return 0;
```



void* threadFunction(void *arg)

```
int main(int argc, char **argv)
                                             a = a + 1;
                                             return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
  return 0;
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                               a = a + 1;
                                               return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread_join(tid[i], NULL);
                                           void* threadFunction(void *arg)
                                            a = a + 1;
  return 0;
                                            return NULL;
```



void* threadFunction(void *arg)

```
int main(int argc, char **argv)
                                      a=a+1;
                                      return NULL;
 int i, P=2;
 Câte thread-uri avem în execuție?
  pthread create(&(tid[i]), NULL, threadFunction, NULL); <
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
                                  void* threadFunction(void *arg)
                                   a = a + 1:
 return 0;
                                   return NULL;
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                   a=a+1;
                                   return NULL;
 Câte thread-uri avem în execuție?
for(i = 0; i < P; i++) {
  pthread_create(&(tid[i]), NULL, PreadFunction, NULL);
                          main
 for(i = 0; i < P; i++) {
  pthread_join(tid[i]_NULL): thread_function(void *arg)
 return 0;
                 threadFunction
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                               a = a + 1;
                                               return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread_create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread join(tid[i], NULL);
                                          void* threadFunction(void *arg)
                                            a=a+1;
  return 0;
                                            return NULL;
```



void* threadFunction(void *arg)

```
int main(int argc, char **argv)
                                               a = a + 1;
                                               return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread join(tid[i], NULL);
                                          void* threadFunction(void *arg)
                                            a = a + 1;
  return 0;
                                            return NULL;
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                               a = a + 1;
                                                return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread join(tid[i], NULL);
                                           void* threadFunction(void *arg)
                                            a = a + 1;
  return 0;
                                            return NULL;
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                               a = a + 1;
                                                return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread join(tid[i], NULL);
                                           void* threadFunction(void *arg)
                                            a = a + 1;
  return 0;
                                            return NULL;
```



```
Thread-ul main int main (int argo char **argv)

Thread-ul main
așteaptă ca thread-ul
     cu handler-ul tid[i] să
  termine execuția
```



```
void* threadFunction(void *arg)
int main(int argc, char **argv)
                                               a = a + 1;
                                               return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread join(tid[i], NULL);
                                          void* threadFunction(void *arg)
                                            a = a + 1;
  return 0;
                                            return NULL;
```



```
int main(int argc, char **argv)
                                               a = a + 1;
                                               return NULL;
 int i, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
    pthread join(tid[i], NULL);
                                          void* threadFunction(void *arg)
                                            a = a + 1;
  return 0;
                                            return NULL;
```



```
int main(int argc, char **argv)
                                             a = a + 1;
                                             return NULL;
 int I, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
  return 0;
```



```
int main(int argc, char **argv)
                                             a = a + 1;
                                             return NULL;
 int I, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
  return 0;
```



```
int main(int argc, char **argv)
                                             a = a + 1;
                                             return NULL:
 int I, P=2;
 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
  return 0;
```



```
int main(int argc, char **argv)
 int I, P=2;
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 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
 return 0;
```



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int main(int argc, char **argv)
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   pthread join(tid[i], NULL);
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 for(i = 0; i < P; i++) {
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 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
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 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread_join(tid[i], NULL);
 return 0;
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```
int main(int argc, char **argv)
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 for(i = 0; i < P; i++) {
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   pthread_join(tid[i], NULL);
 return 0;
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```
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 pthread t tid[P];
 for(i = 0; i < P; i++) {
   pthread create(&(tid[i]), NULL, threadFunction, NULL);
 for(i = 0; i < P; i++) {
   pthread join(tid[i], NULL);
 return 0;
```





Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

Care este valoare lui a?



$$a = 0$$

Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

Care este valoare lui a?



$$a = 0$$

Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

Care este valoare lui a?

4



$$a = 0$$

Thread 1

$$a = a + 2$$

Thread 2

$$a = a + 2$$

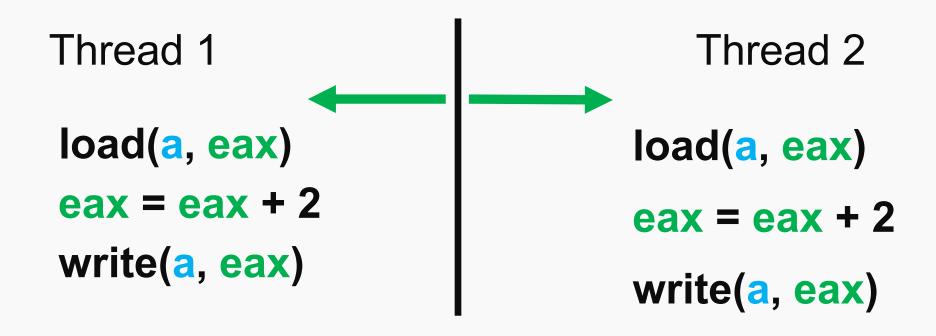
Care este valoare lui a?

4 **ŞI** 2



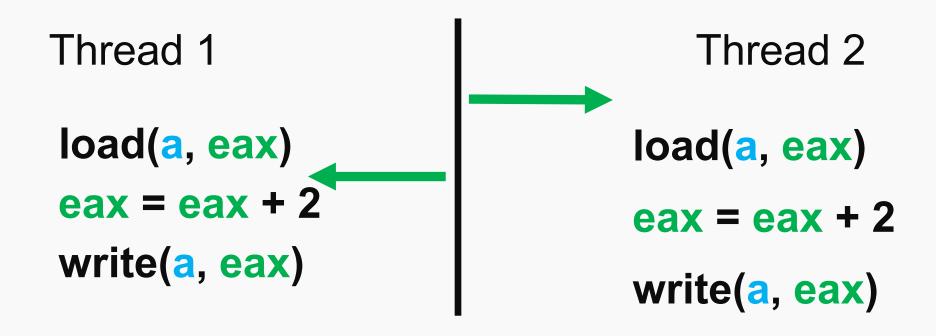


$$a = 0$$





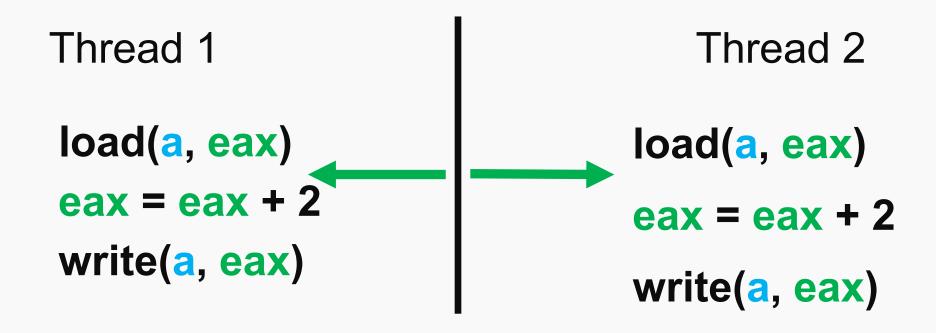
$$a = 0$$



$$eax = 0$$



$$a = 0$$

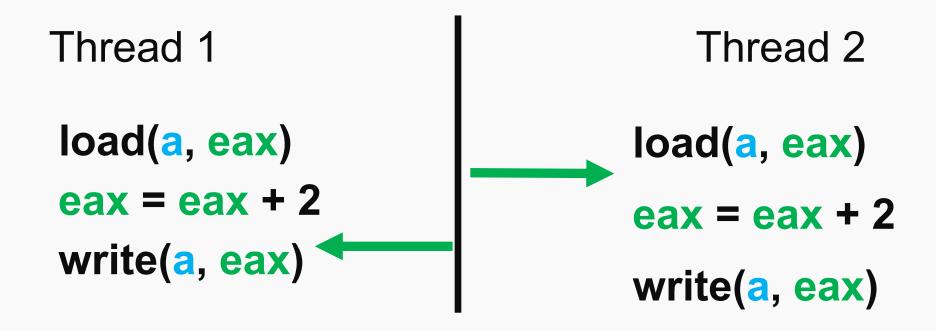


$$eax = 0$$

$$eax = 0$$



$$a = 0$$

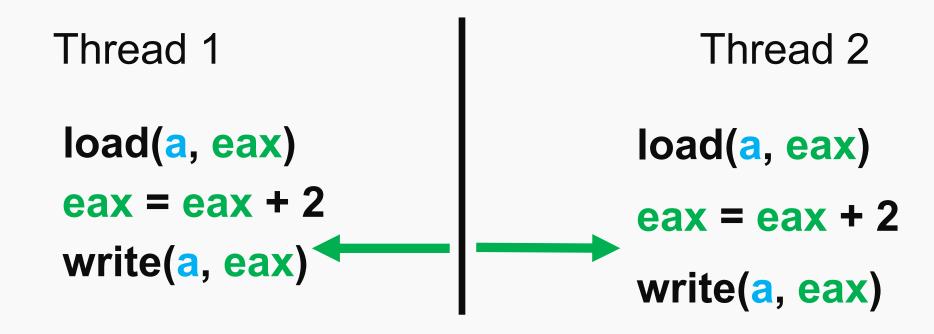


$$eax = 2$$

$$eax = 0$$



$$a = 0$$

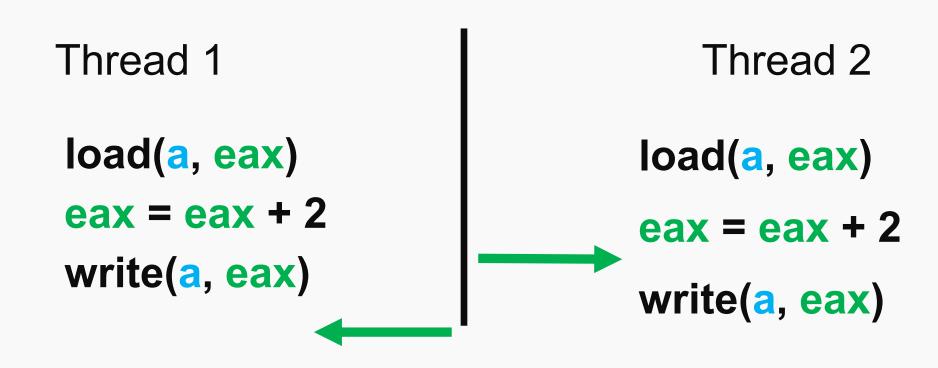


$$eax = 2$$

$$eax = 2$$



$$a = 2$$



$$eax = 2$$

$$eax = 2$$



$$a = 2$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

Thread 2

load(a, eax)

eax = eax + 2

write(a, eax)

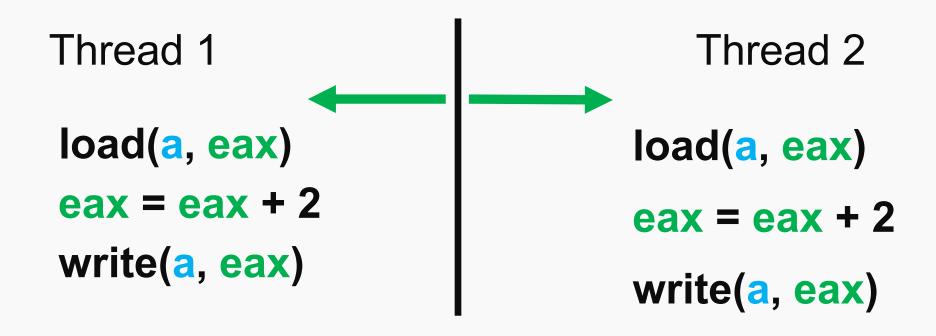
$$eax = 2$$

$$eax = 2$$



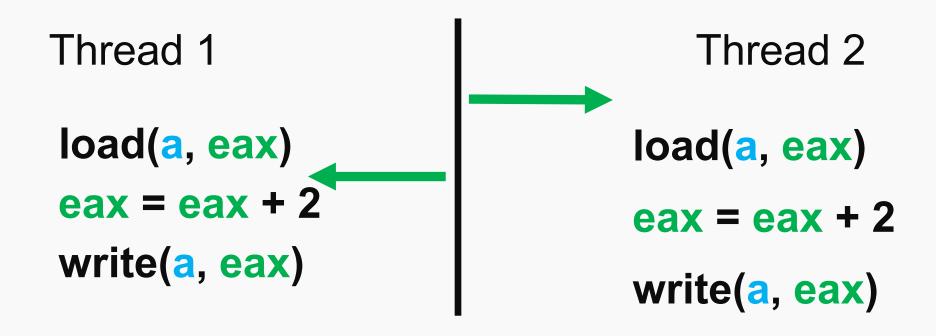


$$a = 0$$





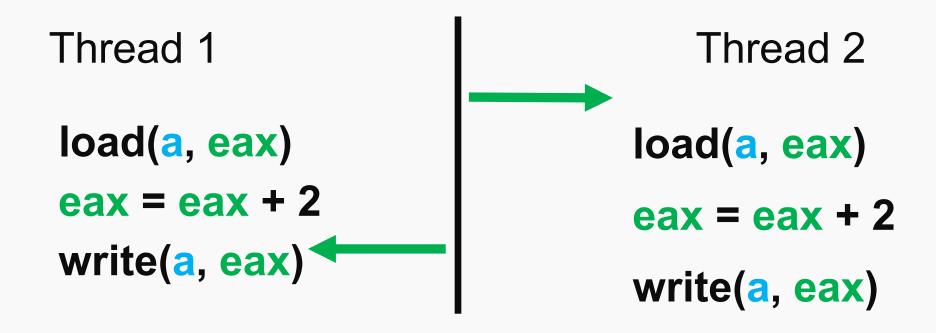
$$a = 0$$



$$eax = 0$$

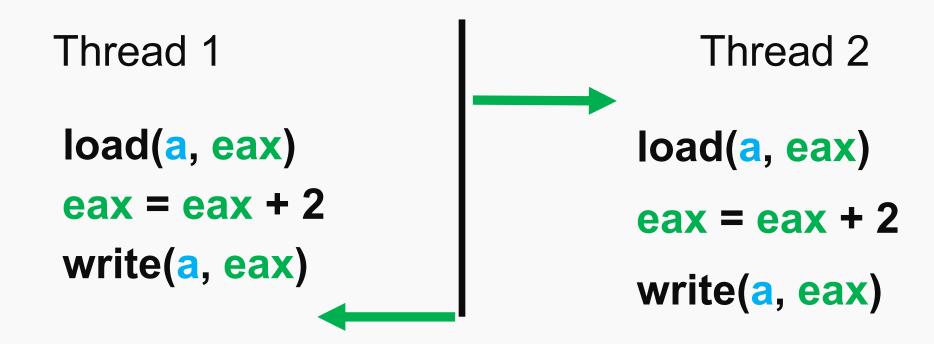


$$a = 0$$



$$eax = 2$$

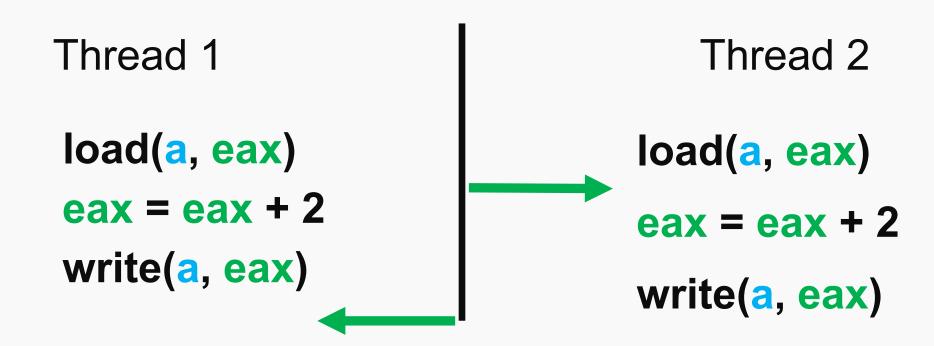




$$eax = 2$$



$$a = 2$$

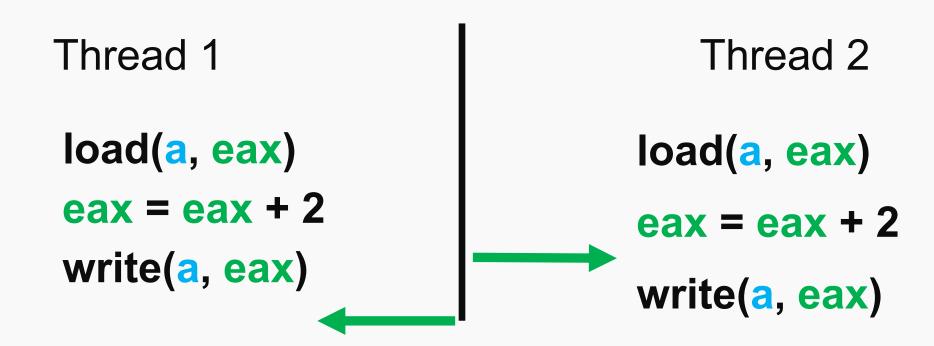


$$eax = 2$$

$$eax = 2$$



$$a = 2$$



$$eax = 2$$

$$eax = 4$$



$$a = 4$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

Thread 2

load(a, eax)

$$eax = eax + 2$$

write(a, eax)

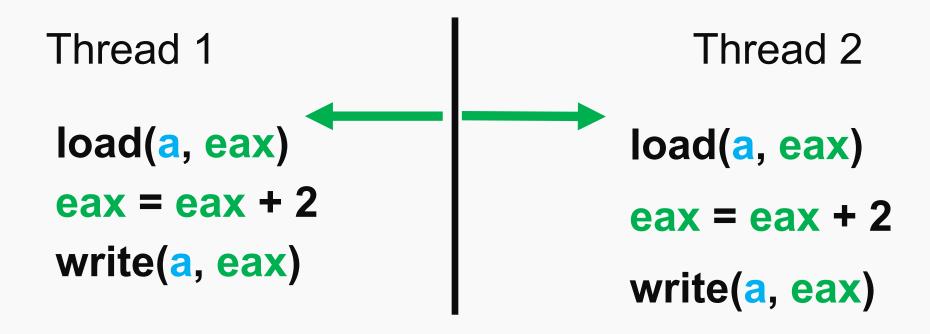
$$eax = 2$$

$$eax = 4$$



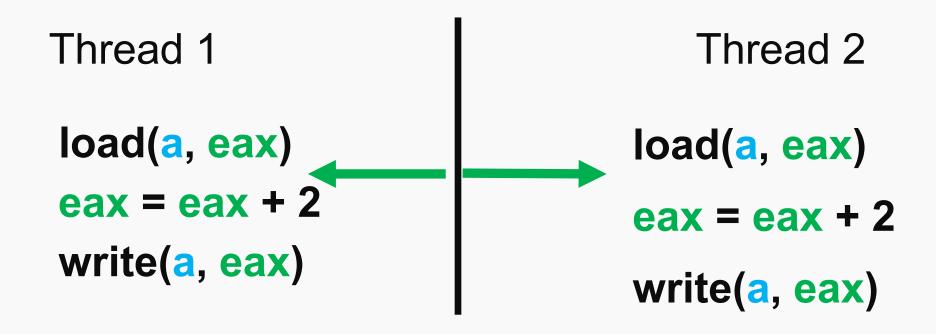


$$a = 0$$





$$a = 0$$

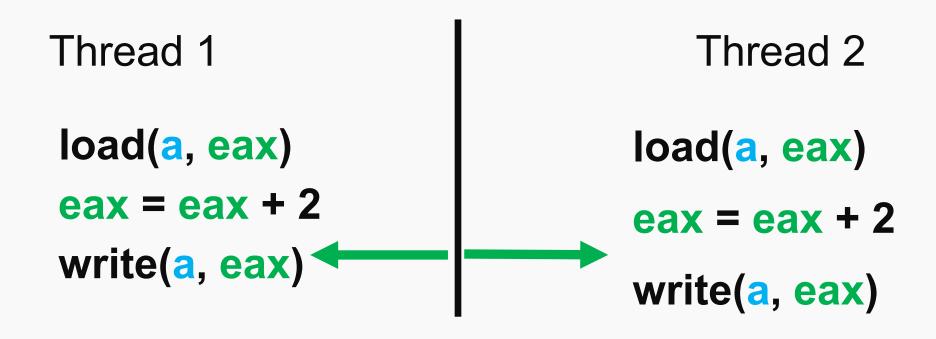


$$eax = 0$$

$$eax = 0$$



$$a = 0$$



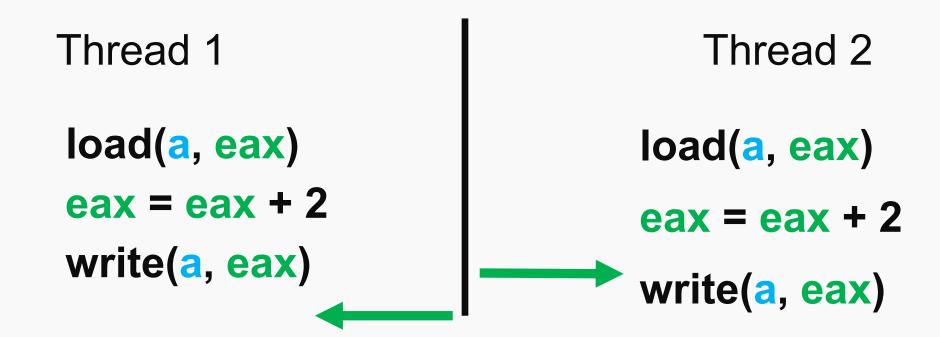
$$eax = 2$$

$$eax = 2$$



Race condition

$$a = 2$$



$$eax = 2$$

$$eax = 2$$



Race condition

$$a = 2$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

Thread 2

load(a, eax)

$$eax = eax + 2$$

write(a, eax)

$$eax = 2$$

$$eax = 2$$



Race condition

$$a = 2$$

Thread 1

load(a, eax)

eax = eax + 2

write(a, eax)

Thread 2

load(a, eax)

eax = eax + 2

write(a, eax)

$$eax = 2$$

$$eax = 2$$





Primitive de sincronizare Synchronization primitives

Atomics

Atomice

•Semaphore



Semafoare

- Binary semaphore (Mutex)
- Critical section

Barrier



Bariere



Fie variabile de 64 biţi pe un procesor 64 biţi C = A + B

load(A, eax)

load(B, ebx)

eax = eax + ebx

write(C, eax)



Fie variabile de 64 biţi pe un procesor 32 biţi C = A + B

```
load(A[0], eax)
load(B[0], ebx)
eax = eax + ebx
write(C[0], eax)
load(A[1], eax)
load(B[1], ebx)
eax = eax + ebx
write(C[1], eax)
```



Fie variabile de 64 biţi pe un procesor 32 biţi C = A + B

```
load(A[0], eax)
load(B[0], ebx)
eax = eax + ebx
write(C[0], eax)
load(A[1], eax)
load(B[1], ebx)
eax = eax + ebx
write(C[1], eax)
```

Putem avea doar jumătate de C modificat



Fie variabile de 64 biţi pe un procesor 32 biţi C = A + B

```
load(A[0], eax)
load(B[0], ebx)
eax = eax + ebx
write(C[0], eax)
load(A[1], eax)
load(B[1], ebx)
eax = eax + ebx
write(C[1], eax)
```

Atomicitatea asigură că C va fi vizibil doar complet modificat, sau complet nemodificat.



Excludere mutuală Mutual exclusion - mutex



Mutual exclusion - Dekker's solution

EWD35 - 1

EWD35.htm

About the sequentiality of process descriptions.

Over de sequentialiteit van procesbeschrijvingen.

Het is niet ongebruikelijk, wanneer een spreker zijn een inleiding. Omdat ik mij hier mogelijk richt tot een g minder vertrouwd is met de problematiek, die ik wil aansn die ik zal moeten gebruiken, wilde ik in dit geval ter in inleidingen houden, nl. een om de achtergrond van de problem tweede, om U een gevoel te geven voor de aard van de wij tegen het lijf zullen lopen.



Theodorus (Dirk) J. Dekker



Mutual exclusion – Dijsktra's Solution

Solution of a Problem in Concurrent Programming Control

E. W. Dijkstra Technological University, Eindhoven, The Netherlands

A number of mainly independent sequential-cyclic processes with restricted means of communication with each other can be made in such a way that at any moment one and only one of them is engaged in the "critical section" of its cycle.

Introduction

Given in this paper is a solution to a problem for which, to the knowledge of the author, has been an open question since at least 1962, irrespective of the solvability. The paper consists of three parts: the problem, the solution, and the proof. Although the setting of the problem might seem somewhat academic at first, the author trusts that anyone familiar with the logical problems that arise in

computer can only request one one-way message at a time. And only this w this problem is f The Solution The common "Boo The integer will only be set h by the others. It well outside the mentioned set to The program "integer j; Li0: b[i] := falseLi1: if $k \neq i$ the Li2: begin c[i]: Li3: if b[k] then go to Lil endelse



Dekker's Solution

```
wants to enter[0] = true
while wants to enter[1] {
      if turn != 0 {
            wants to enter[0] = false
            while turn != 0 { // busy wait }
            wants to enter[0] = true
// critical section ...
turn = 1
wants to enter[0] = false
```



Dijsktra's Solution

```
b[i] = false
while(sw[i]) {
         sw[i] = false
         if (k!=i) {
                   c[i] = true
                   if(b[k])
                            k = i
                   sw[i] = true
         } else {
                   c[i] = false
                  for(j=0;j<N;j++)
                            if(j!=i && !c[j])
                                      sw[i] = true
//critical
b[i] = true
```

lock() - P

unlock() - V

c[i] = true



Peterson's Solution

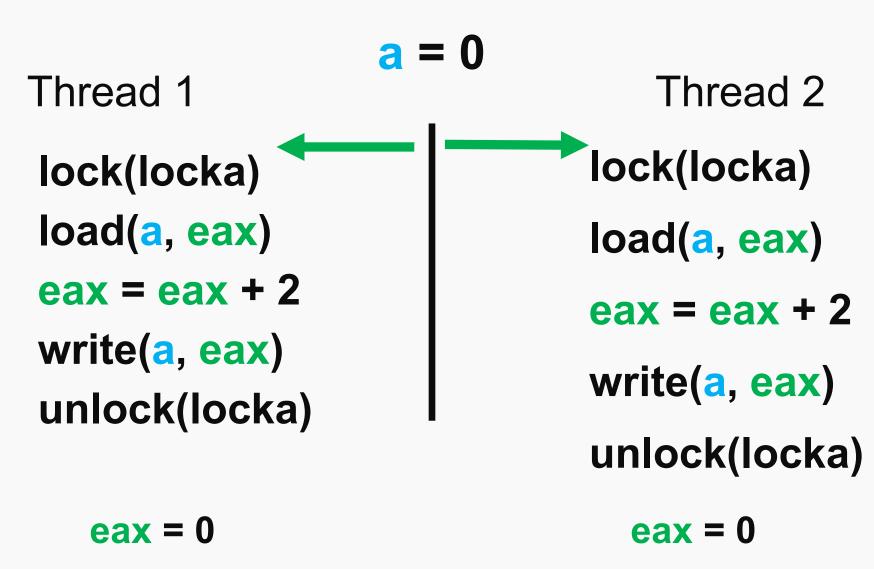
```
flag[0] = true;
turn = 1;
while (flag[1] && turn == 1) { // busy wait }
// critical section ...
flag[0] = false;
```



Hardware assisted Solution

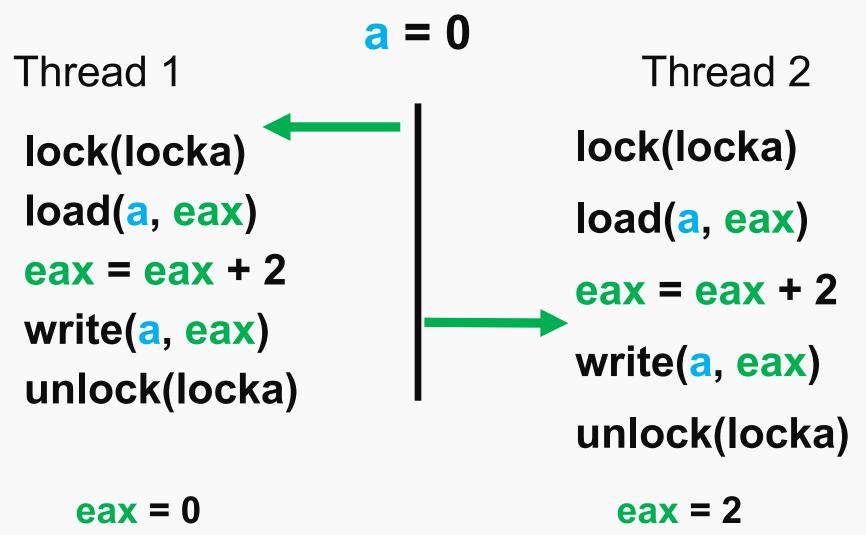
```
while (test_and_set(lock));
// critical section
lock = 0;
```



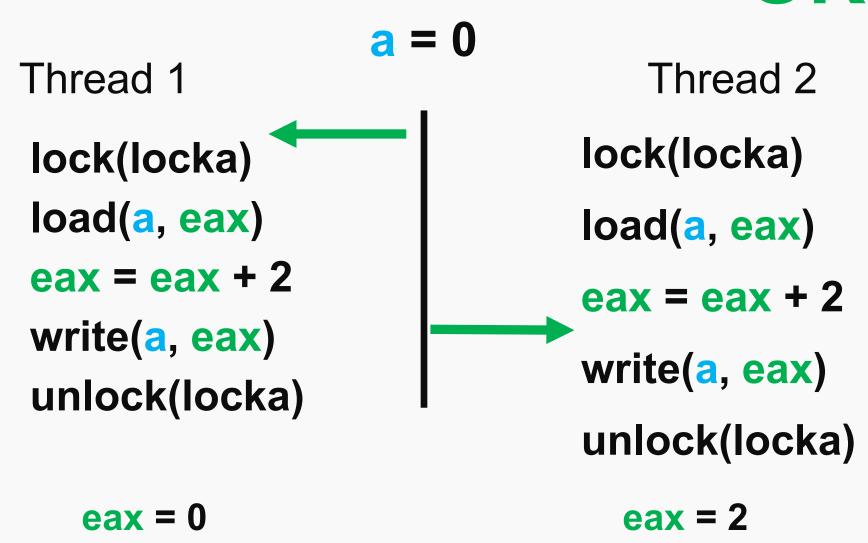
















a = 2

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 0

Thread 2

lock(locka)

load(a, eax)

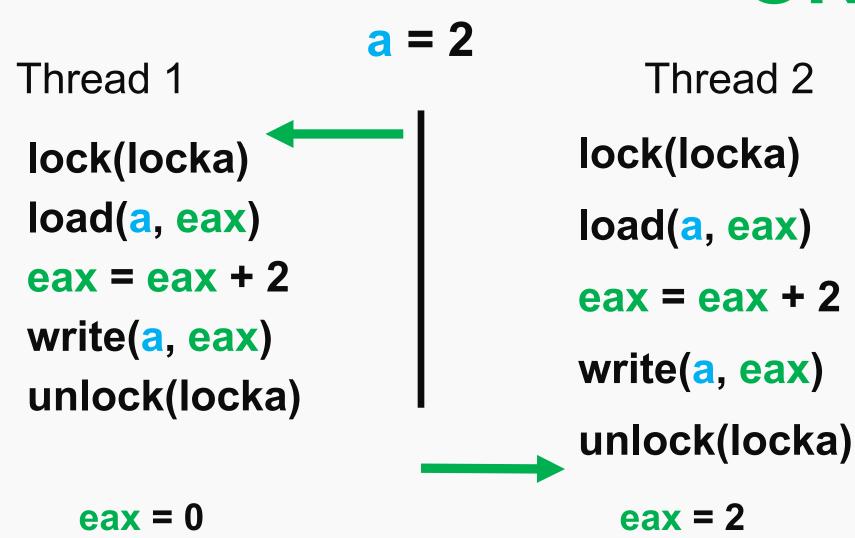
eax = eax + 2

write(a, eax)

unlock(locka)

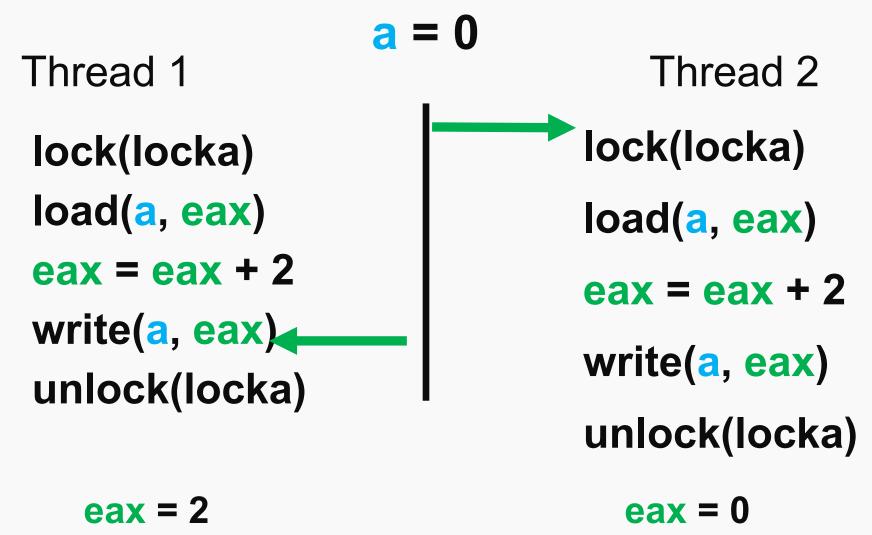
eax = 2



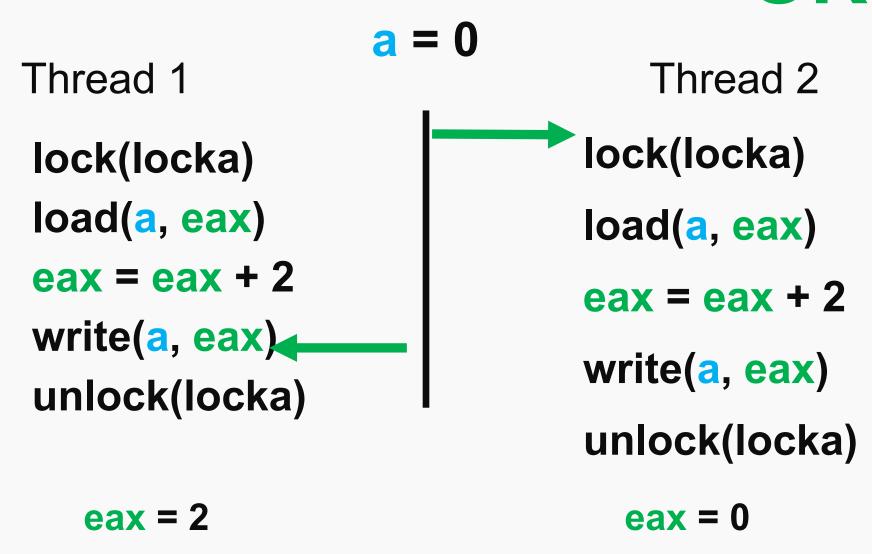
















Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

a = 2

Thread 2

lock(locka)

load(a, eax)

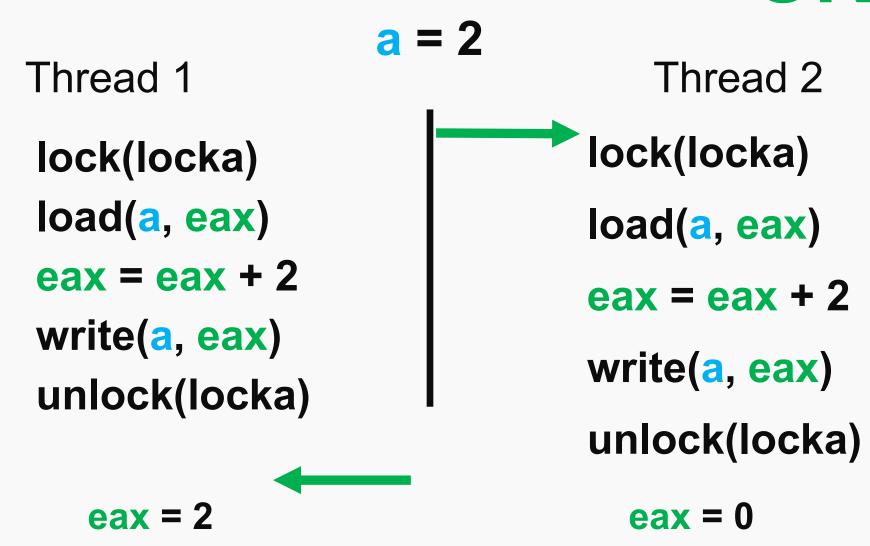
eax = eax + 2

write(a, eax)

unlock(locka)

eax = 0









$$a = 4$$

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2



a = 4

Thread 1

lock(locka)

load(a, eax)

eax = eax + 2

write(a, eax)

unlock(locka)

eax = 2

Thread 2

lock(locka)

load(a, eax)

eax = eax + 2

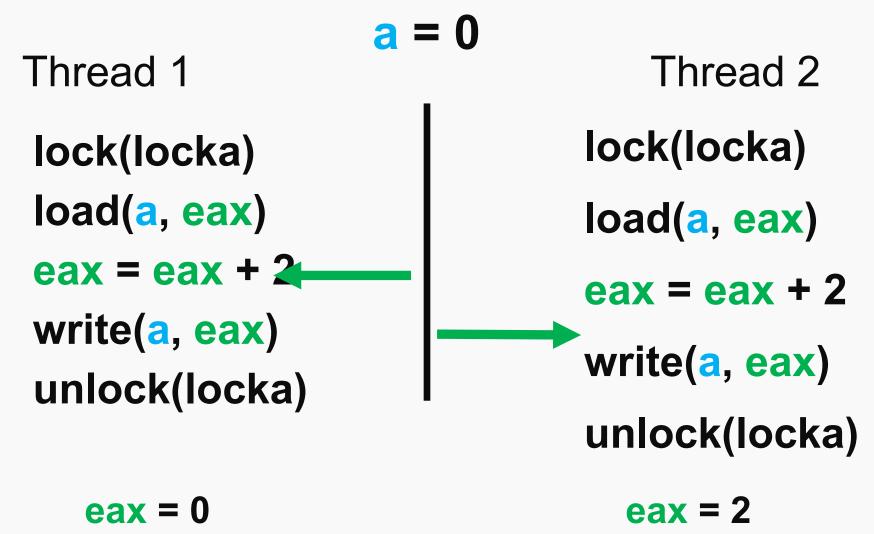
write(a, eax)

unlock(locka)

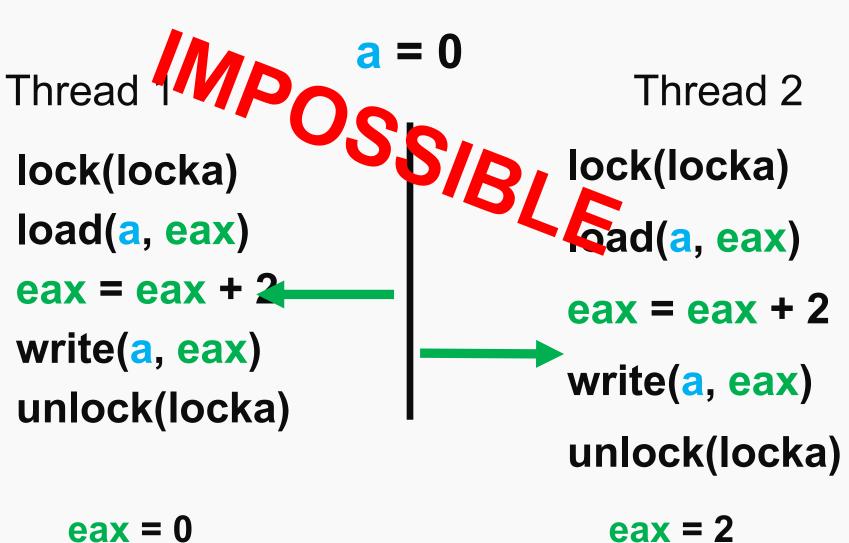
eax = 2













ÎN MAIN

Înainte de a porni thread-urile

```
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL);
```



```
pthread_mutex_t mutex;
pthread_mutex_init(&mutex, NULL);
```

Poate fi folosit să anunțe că acest mutex e împărțit mai multor procese



```
pthread_mutex_lock(&mutex);
load(a, eax)
eax = eax + 2
write(a, eax)
pthread_mutex_unlock(&mutex);
```



ÎN MAIN

După ce au terminat thread-urile

pthread_mutex_destroy(&mutex);





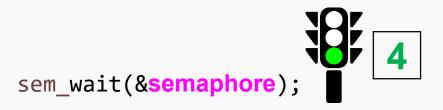
Semaphore

ÎN MAIN

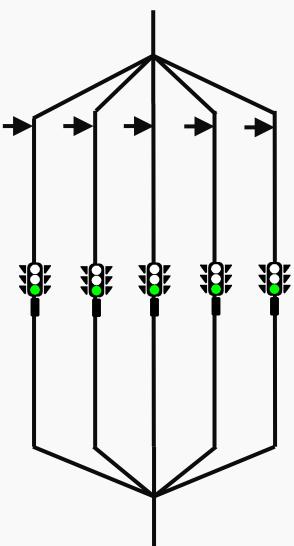
```
Înainte de a porni thread-urile
           gcc –o executabil cod.c –lpthread –lrt
           #include<pthread.h>
           #include<semaphore.h>
sem t semaphore;
int semaphore value= 4;
sem init(& semaphore, 0, semaphore value);
```



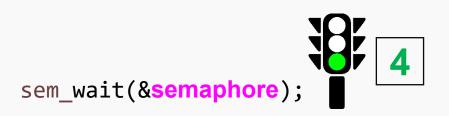
Semaphore

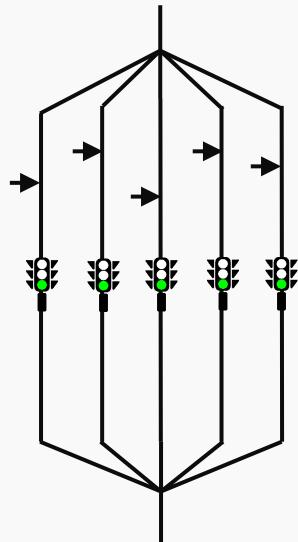


P() sau Proberen - Dijsktra

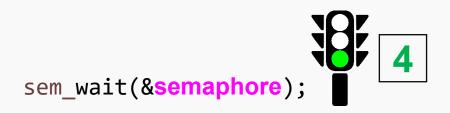


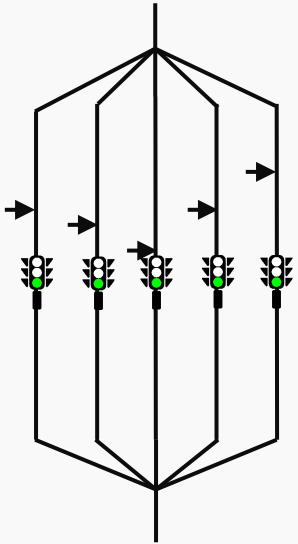






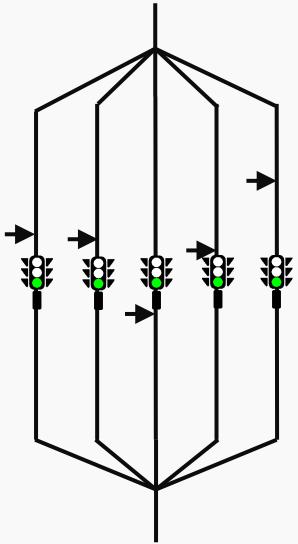








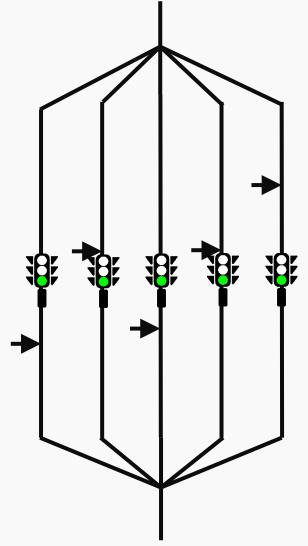




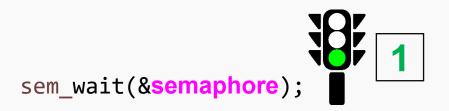


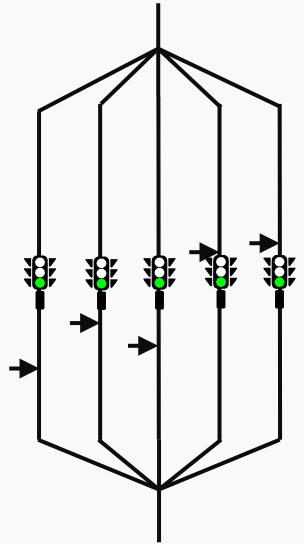
Nu contează că doua thread-uri au ajuns simultan la semafor, acesta este protejat, la fel ca un mutex.

```
sem_wait(&semaphore);
```

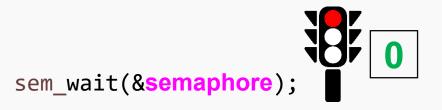


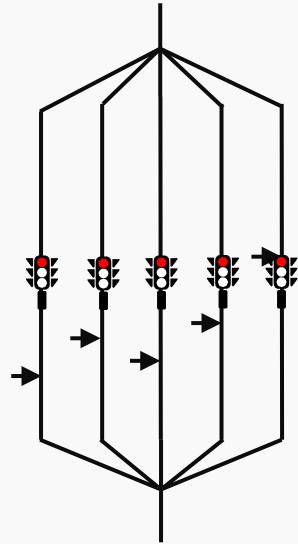






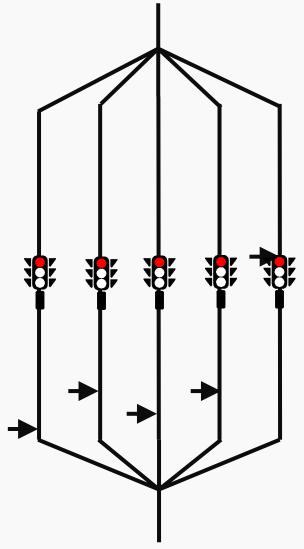






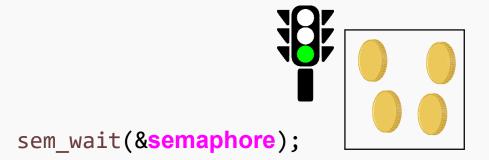


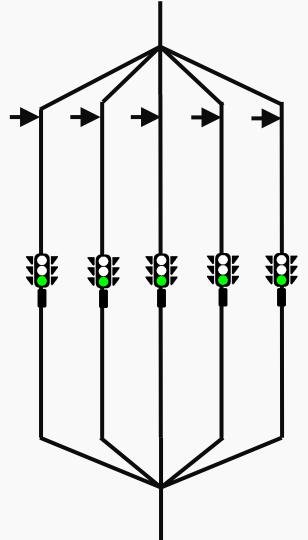






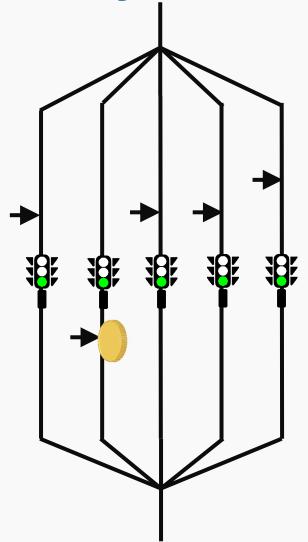
Un semafor are un set de token-uri.





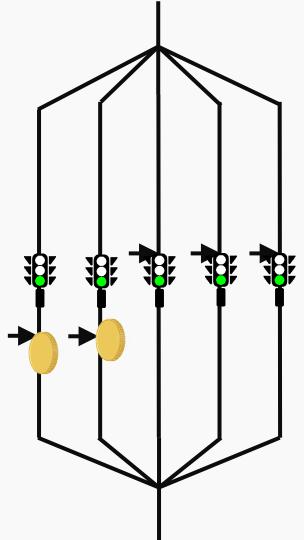






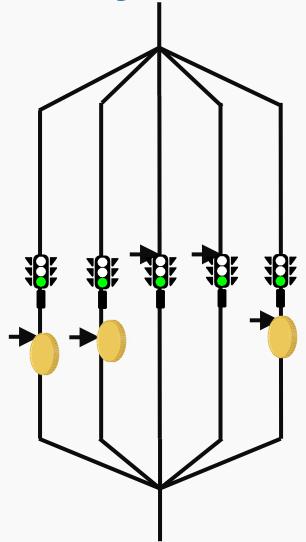






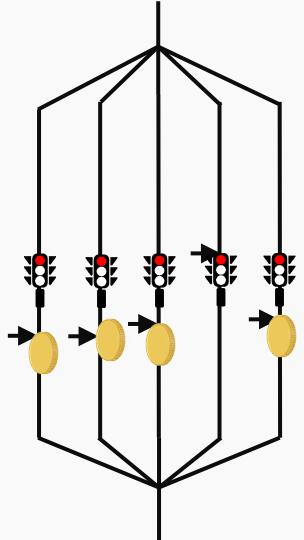




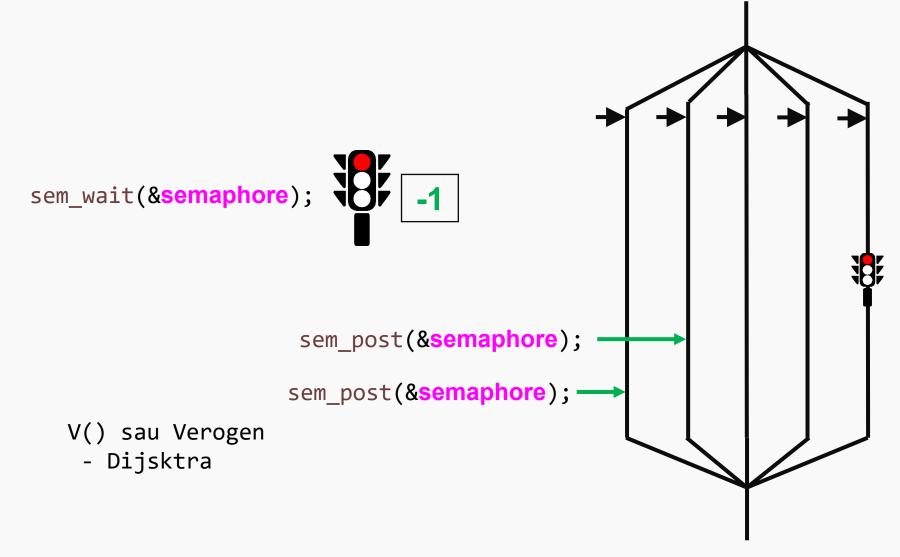




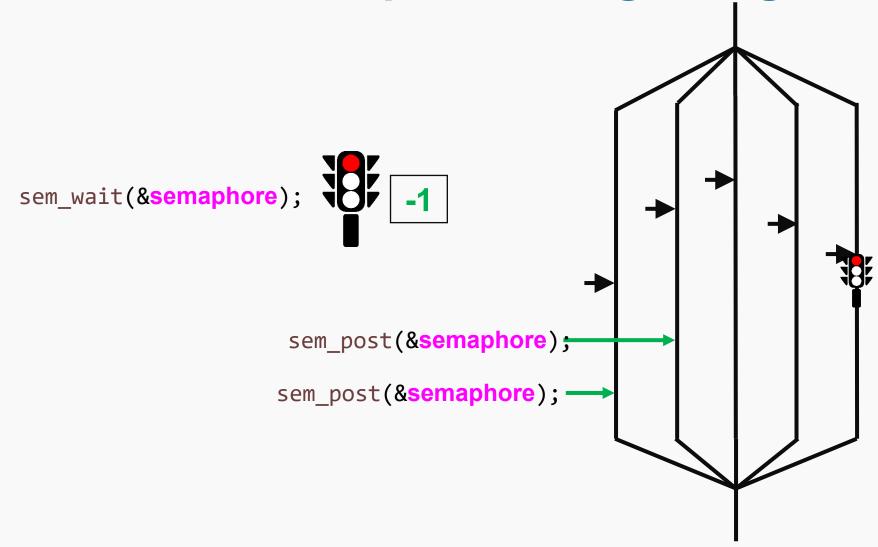




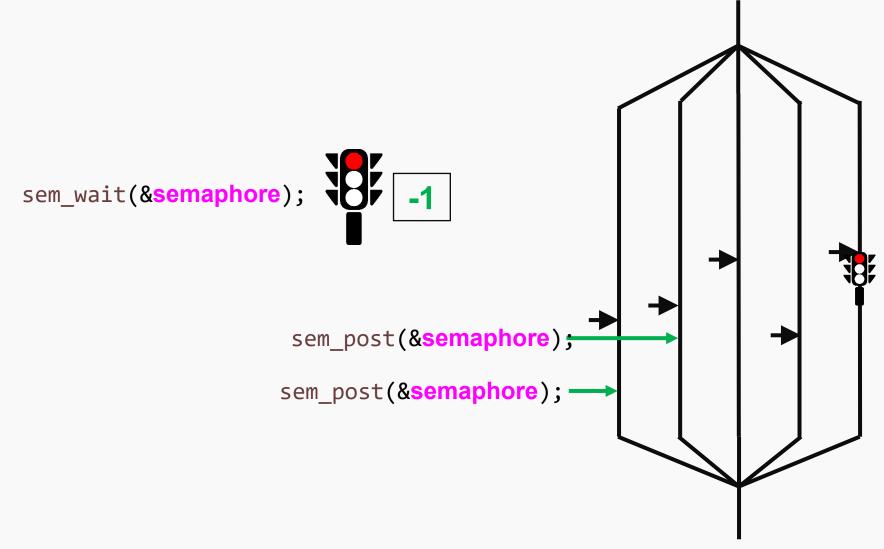




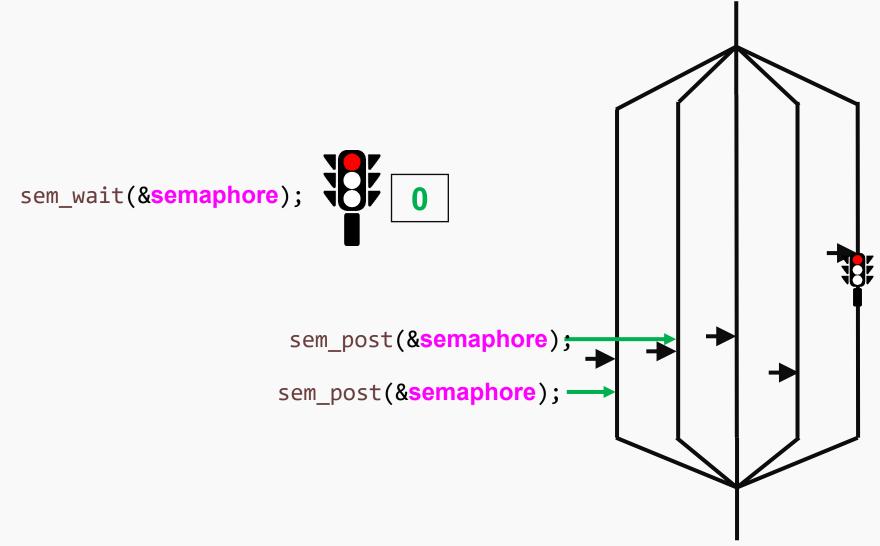




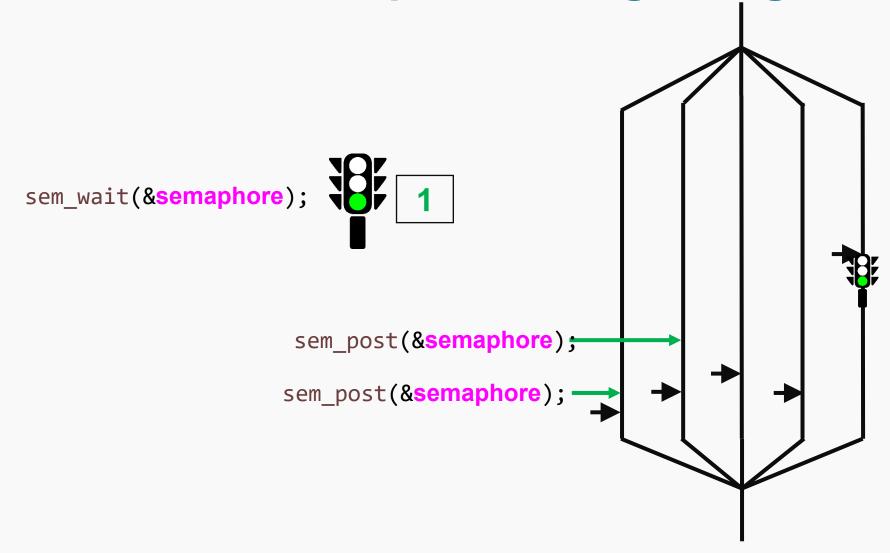




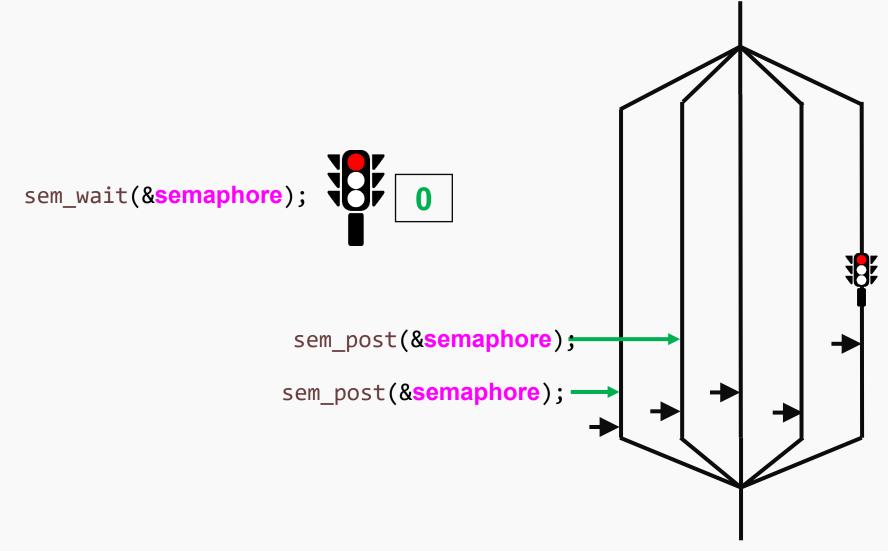














ÎN MAIN

După ce au terminat thread-urile

sem_destroy(& semaphore);







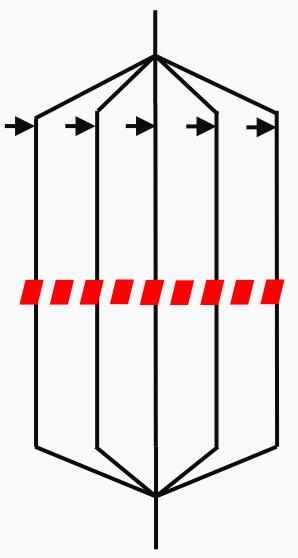
ÎN MAIN

Înainte de a porni thread-urile

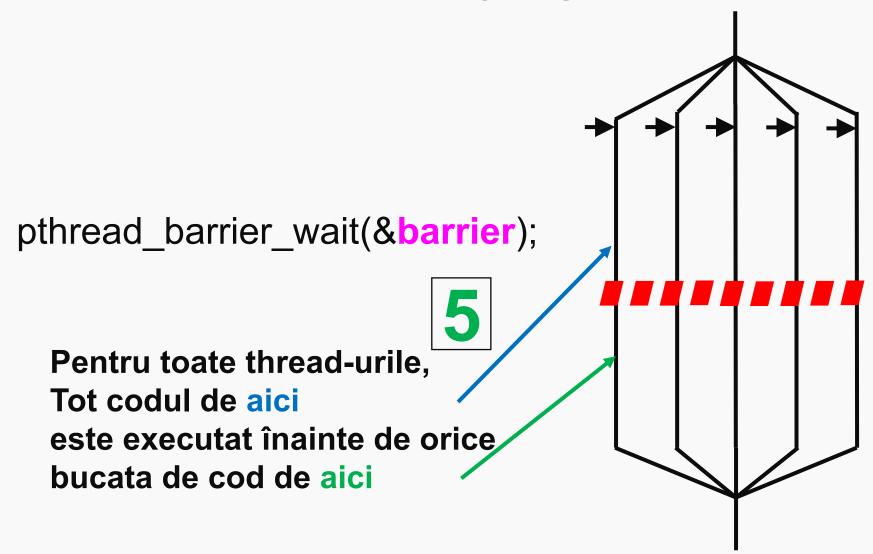
```
pthread_barrier_t barrier;
int num_threads = 5;
pthread_barrier_init(&barrier, NULL, num_threads);
```





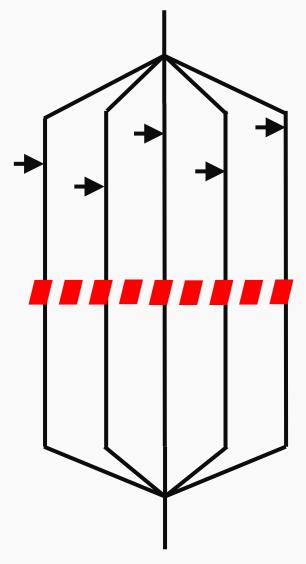






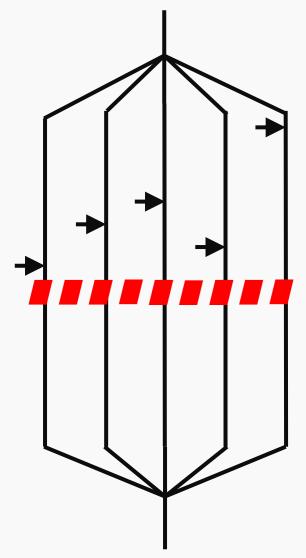






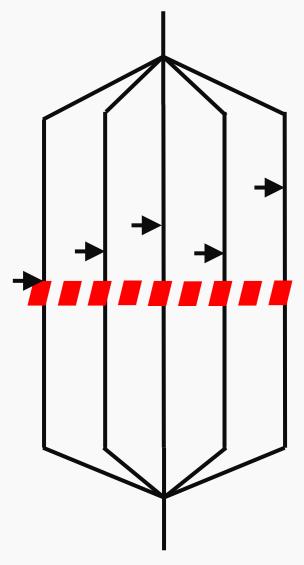






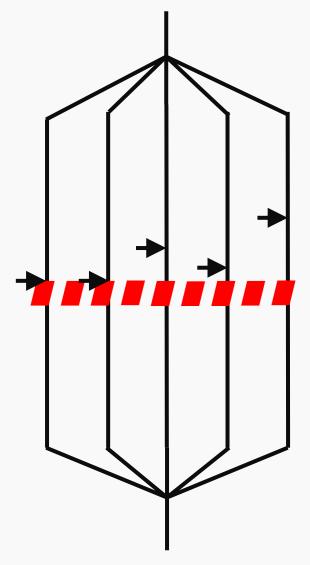








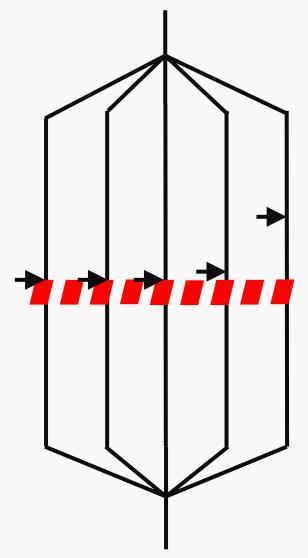






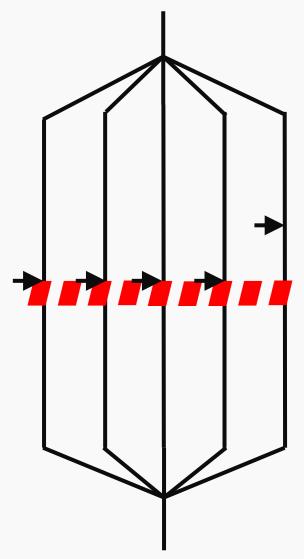
pthread_barrier_wait(&barrier);

2



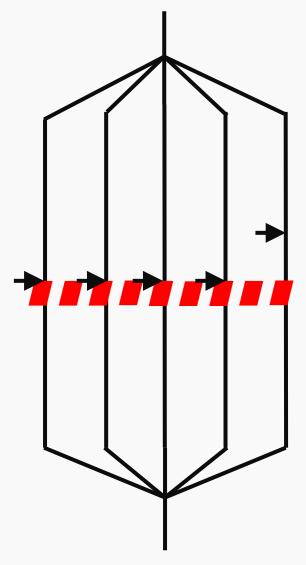






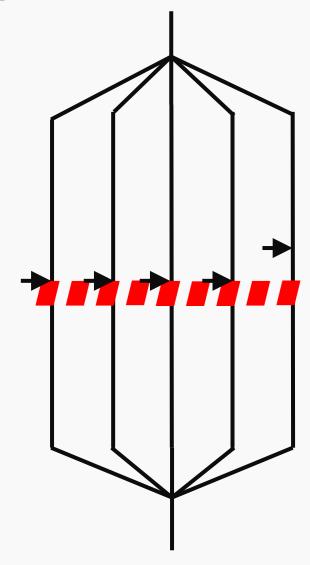






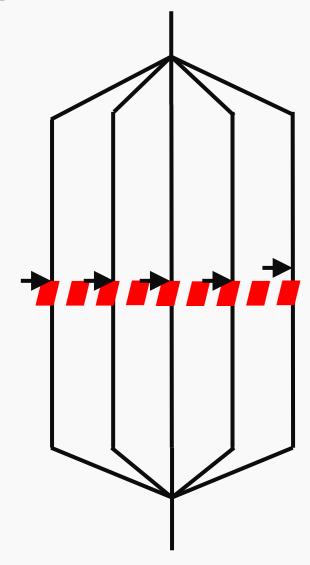






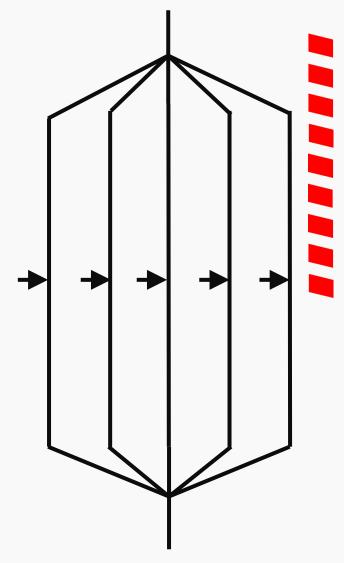






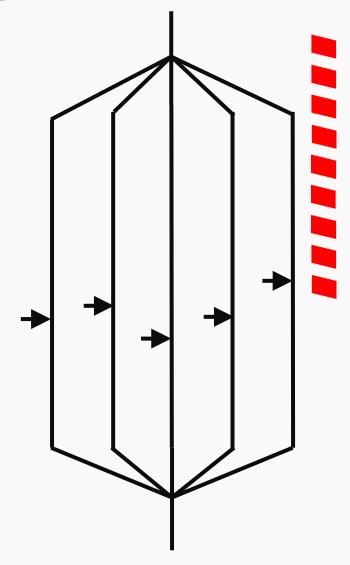












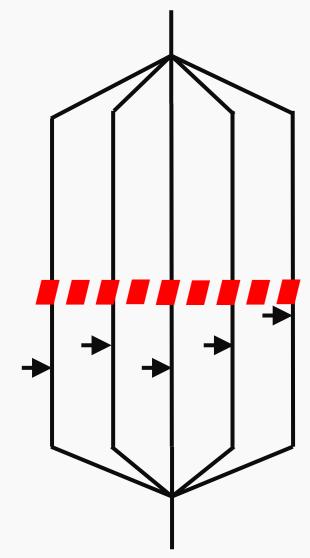


Cum știe o barieră când să se reseteze?

O soluție ar fi: Reusable Barrier in <u>The Little Book of Semaphores</u> <u>By Allen B. Downey</u>









ÎN MAIN

După ce au terminat thread-urile

pthread_barrier_destroy(&barrier);





Unele probleme nu pot fi paralelizate

Calculating the hash of a hash of a hash ... of a string.

Deep First Search

Huffman decoding

Outer loops of most simulations

P complete problems



Paralelizare prin împărțirea problemei

Sunt o serie de probleme care sunt extrem de uşor paralelizabile.

Embarrassingly parallel



Multiplicare unui vector cu un scalar

* 3

27 18 27 12 6 21 18 15 18 3



Toate calculele pot fi efectuate în același timp

* 3

27 18 **27** 12 **6 21 18 15 18 3**



Câte elemente sunt?





Câte elemente sunt?



Cristian Chilipirea – Arhitecturi Paralele



Câte elemente sunt? N

9 6 9 4 2 7 6 5 6

Cristian Chilipirea – Arhitecturi Paralele



Dar câte elemente de procesare?





Dar câte elemente de procesare?

9 6 9 4 2 7 6 5 6



Dar câte thread-uri?

9 6 9 4 2 7 6 5 6



Dar câte thread-uri?



În majoritatea cazurilor obținem performanță maximă când numărul de thread-uri este egal cu numărul de elemente de procesare, sau core-uri.



Cum este P față de N?

9 6 9 4 2 7 6 5 6



P << N

9 6 9 4 2 7 6 5 6



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

.



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

. 1

Thread 1



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

.

Thread 1



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

.

Thread 1



Caz concret: P = 2 Cum împărțim? Putem și random

9 6 9 4 2 7 6 5 6

1

Thread 1



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

Thread 1



Caz concret: P = 2 Cum împărțim?

. .

Este utilă?

Thread 1



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

. 1

Ce ne dorim?

Thread 1



Caz concret: P = 2 Cum împărțim?

9 6 9 4 2 7 6 5 6

. 1

Ce ne dorim?

Thread 1

Thread 2

Aproximativ același număr elemente



Aproximativ N/P elemente

9 6 9 4 2 7 6 5 6

Thread 1



Aproximativ N/P elemente

9 6 9 4 2 7 6 5 6 ... 1

Dacă N nu se divide perfect la P?

Thread 1



Aproximativ N/P elemente Dacă N nu se divide perfect la P?

1

6

4 2 7

9 6 9

Thread 1

8

4 9 2

5 6 3



floor(N/P) elemente floor(15/2) = 7

1

6

4 2 7

9 6 9

Thread 1

8

4 9 2

5 6 3



ceil(N/P) elemente ceil(15/2) = 8

6 5

4 2 7

9 6 9

Thread 1

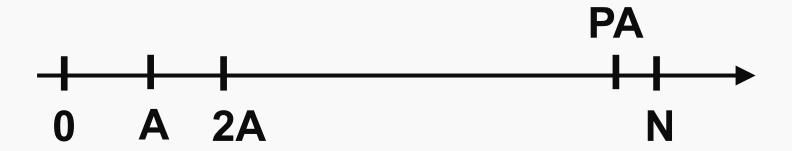
8 1

4 9 2

6 3

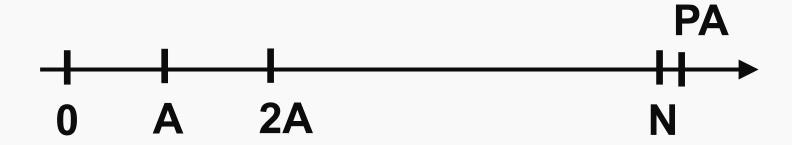


$$A = floor(N/P)$$





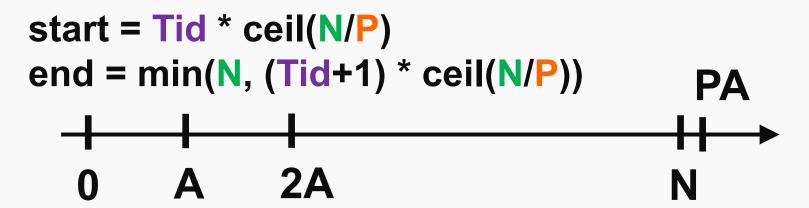
$$A = ceil(N/P)$$





Formule elegante:

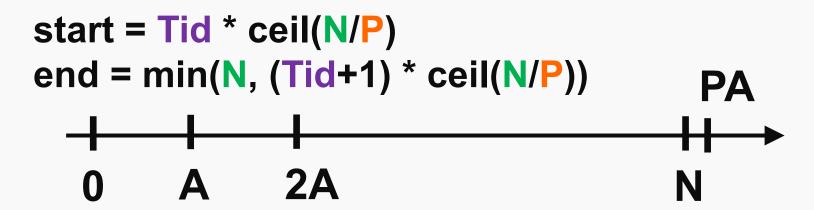
Tid este identificator de thread, are valori de la 0 la P





Formule elegante:

Tid este identificator de thread, are valori de la 0 la P



```
Funcţionează şi:

start = round(Tid * N/P)

end = round((Tid+1) * N/P) De ce?
```