



# IMPERATIVE LANGAUGES

February 15<sup>th</sup>



# LAST TIME



★ We:

- Gave an operational semantics for the lambda calculus
- Compared different evaluation orders for the lambda calculus
- Explored laziness in Haskell



# AGENDA



★ We will:

- Look at how big-step semantics guide the implementation of an interpreter
- Look at how to model state in operational semantics
- Show how to use monads to implement effects in pure languages



# IMP SEMANTICS



★ Recall the syntax for IMP, our core imperative language:

$a ::= \mathbb{V} \mid \mathbb{I} \mid a * a \mid a + a$

$b ::= \text{true} \mid \text{false} \mid \text{not } b \mid b \ \&\& \ b$   
 $\mid a == a \mid a < a$

$S ::= \mathbb{V} := a \mid \text{skip} \mid S1; S2 \mid$   
 $\mid \text{if } b \text{ then } S1 \text{ else } S2 \mid \text{while } b \text{ do } S$

★ Semantics for (basic) arithmetic expressions is straightforward:

$$\frac{}{n \Downarrow n} \text{E-CONST}$$

$$\frac{n \Downarrow v_n \quad m \Downarrow v_m}{n * m \Downarrow n * m} \text{E-TIMES}$$

$$\frac{n \Downarrow v_n \quad m \Downarrow v_m}{n + m \Downarrow n + m} \text{E-PLUS}$$



# HASKELL IMPLEMENTATION



★ Has a ‘natural’ implementation in Haskell:

```
data AExp = Const Int | Plus AExp AExp | Times AExp AExp
```

```
eval (Const i) = i
```

```
eval (Plus a1 a2) = (eval a1) + (eval a2)
```

```
eval (Times a1 a2) = (eval a1) * (eval a2)
```

★ How could we add boolean expressions?

```
b ::= true | false | not b | b && b | a == a | a < a
```





# HANDLING STATE



- ★ Core issue is that Variable require state
  - This is what makes IMP an interesting language
  - Intuitively, how do we evaluate the following expression?

$$\frac{}{x + 3 \Downarrow}$$

- How do we represent global state in an evaluation rule?

$$\frac{}{x \Downarrow ?} \quad \text{E-VAR}$$

**I'VE GOT A FEVER**

**$\alpha$   $\Omega$**

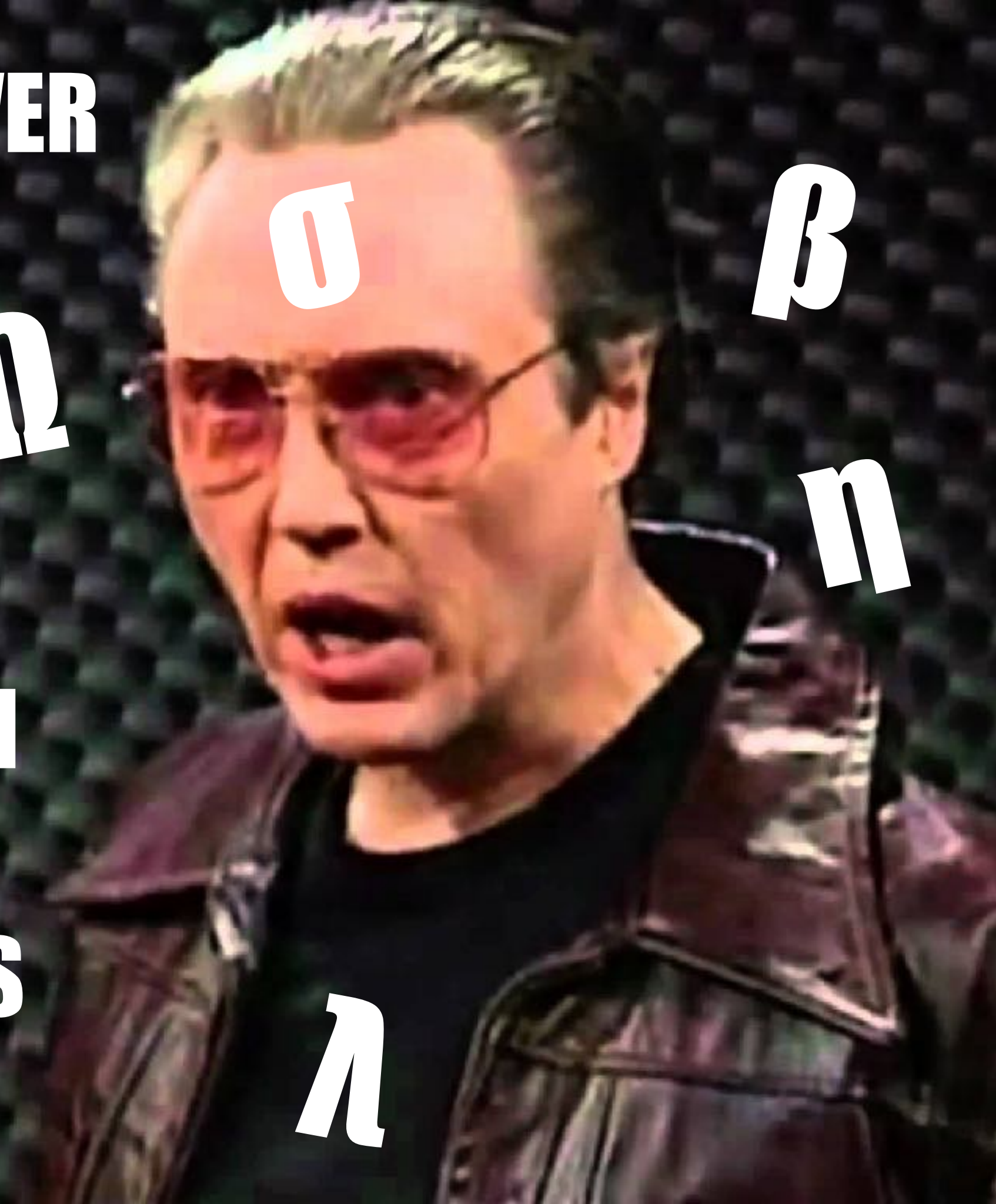
**AND THE ONLY  
PRESCRIPTION  
IS MORE  
GREEK LETTERS**

**$\sigma$**

**$\beta$**

**$\eta$**

**$\lambda$**

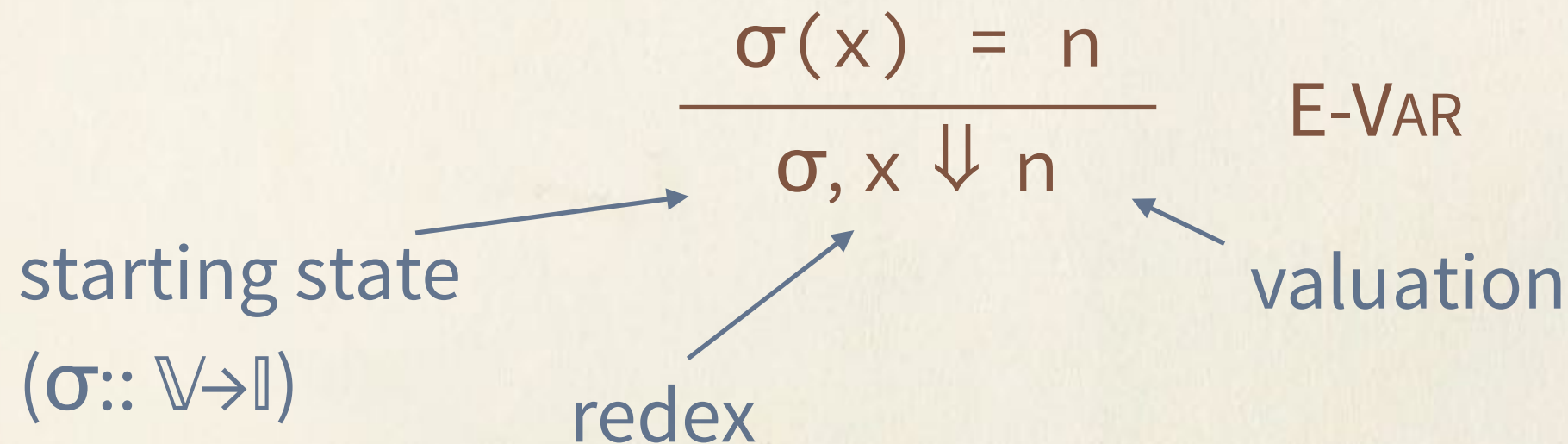




# ADDING STATE



- ★ Idea: evaluation becomes a three-way relation:



- ★ Adding state to the relation lets us use it in the premise
- ★ We have to update other rules as well:

$$\frac{\sigma, x \Downarrow v_x \quad \sigma, y \Downarrow v_y}{\sigma, x+y \Downarrow v_x + v_y} \quad \text{E-VAR}$$





# HASKELL IMPLEMENTATION



## ★ What about statements?

$S ::= \mathbb{V} := a \mid \text{skip} \mid S1; S2 \mid \text{if } b \text{ then } S1 \text{ else } S2 \mid$   
 $\text{while } b \text{ do } S$

- ★ In C or Java, we might just have a global variable for the mapping and we would update it appropriately.
- ★ Haskell is a pure language— there are no mutable variables!
- ★ What's the solution?



# HASKELL IMPLEMENTATION



- ★ In C or Java, we might just have a global variable for the mapping and we would update it appropriately.
  - ★ Haskell is a pure language— there are no mutable variables!
  - ★ What's the solution?
- 
- ★ The key bit was threading the state parameter through in a way that was consistent with the semantics
  - ★ Passing around all these variables was a little cumbersome.



# MONADS



- ★ This idiom is so common that Haskell has a standard method of approaching it:

**MONADS**

- ★ What is a monad:
  - “A monad is just a monoid in the category of endofunctors”  
—Phil Wadler
  - A monad is a means for wrapping values in a context and passing that context around.



# MONADS



## ★ Monad typeclass (basic) definition:

class Monad m where  
 return :: a -> m a      ← Type of Containers or Contexts  
 (>>=) :: m a -> (a -> m b) -> m b      ← put a value into the container  
 ← sequence two contexts together

## ★ The Option Monad:

```
instance Monad Option where  
  return a = Just a  
  Nothing >=> f = Nothing  
  Just a >=> f = f x
```





# MONADS



## ★ The State Monad:

```
data State s a = State (s -> (a, s))
instance Monad (State s) where
    return x = State $ \s -> (x,s)
    (State h) >>= f = State $ \s -> let (a, newState) = h s
                                      in g newState
                                      where (State g) = f a
```

## ★ The Monad laws:

- left identity:  $\text{return } x \gg= f \cong f \ x$
- right identity:  $m \gg= \text{return} \cong m$
- associativity:  $(m \gg= f) \gg= g \cong m \gg= (\lambda x \rightarrow f \ x \gg= g)$



# RECAP



★ We will:

- Look at how big-step semantics guide the implementation of an interpreter
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