

# Types in Passerine

Weekly meeting 2021-04-17  
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# Quick facts

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Y'all know, but: a type system allows a compiler to determine the specific structure of data at each point in the program:

- Strong vs Weak: Casting and comparing data
  - Static vs Dynamic: When types are checked
  - Nominal vs Structural: How types are described
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# Algebraic Data Types

A composite type, i.e. types made up of other types:

- Base types:
  - Integers, strings, floating point numbers, booleans, bytes, etc.
- Product types:
  - Structs or records or tuples
  - Made up of a number of *fields*
- Sum types:
  - Tagged unions
  - Made up of a number of *variants*

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# Mapping ADTs to Passerine

## Product types:

- Tuples  
(indexed by  
number)
- Records  
(indexed by  
name)

## Sum types:

- Enums  
(closed unions)
- Traits  
(open unions)

## Base types:

- Boolean
- Integers
- Reals
- String
- Functions
- Fibers
- Alg. Effs.?
- (More; this for  
now)



# **Part I: Product Types**

# Tuples

## Definition:

- (Label\_0, ..., Label\_N)

## Construction:

- (expression\_0, ..., expression\_N)

## Destruction:

- (pattern\_0, ..., pattern\_N)

## Update:

- (function\_0, ..., function\_N)

# Records

## Definition:

- { field\_0: Label\_0, ... field\_N: Label\_N }

## Construction:

- { field\_0: expression\_0, ... field\_N: expression\_N }

## Destruction:

- { field\_0: pattern\_0, ... field\_N: pattern\_N }

## Update:

- { field\_0: function\_0, ... field\_N: function\_N }

# Aside: Row Polymorphism

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If the record A is a subset of record B, record A is a B.

- `type Foo { a: Int }`
- `type Bar { a: Int, b: String }`

Therefore, Foo is a Bar



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# Part II: Sum Types

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# Enums

## Definition:

- { Label\_0 type\_0, Label\_N type\_N }

## Construction:

- Label expression

## Destruction:

- Label pattern

## Update:

- Label function

# Traits

I'm not sure at all... some options

1. Associated namespaces
2. Dynamic dispatch
3. Multiple dispatch (Julia)
4. Rust traits / Haskell typeclasses
5. ???

# Part III: Traits



# Associated Namespaces

Basic Idea:

- Each type has a record, i.e. an associated namespace
- We can index into that namespace:
  - `Label::field`
- We can define interfaces as records:
  - `type Add T { add: T T -> T }`
- Because of row polymorphism, if `Label` has associated value `Label::add`, `Label` is `Add`.

# Associated Namespaces

Problems:

- How do we define associated namespaces?
  - impl
  - direct
- Diamond problem:
  - Two traits, Couch and Chair both have `::sit`
  - We make an Armchair, both Couch and Chair apply
  - Which `::sit` method is used for Armchair?

# Dynamic / Multiple Dispatch

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Think of this like open match statements:

- We define a 'function' X:
  - Has a number of methods.
  - Each method maps from A -> B
- If we call function(D):
  - Looks up method that takes type D
  - Runs that method against it
- No diamond problem, always call function.

# Dynamic / Multiple Dispatch

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## Issues:

- Only granular on the basis of functions:
    - Can't define a 'Number' that implements add, sub, mul, div, etc. methods.
  - Not clear what method you are calling, or which methods take priority
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# Traits / Typeclasses

- You define a 'trait' X:
  - Has a number of 'members'
  - Each member is a type, function, etc.
- You implement trait X for type Y:
  - All members must be implemented
- Dispatch based on type:
  - If X is implemented for Y and X::member exists:
  - Y.member() will dispatch appropriately.
- Diamond Problem is solved:
  - Use explicit function call syntax:
  - X::member(Y) and Z::member(Y)



# Traits / Typeclasses

Issues:

- With languages like Rust (low-level), annoying to manage &dyn
  - Not an issue for high-level langs like Haskell or Passerine
- Trait objects are a tad confusing
- Have to explicitly implement it for each type
- Introduces additional syntax

Honestly probably the best solution?



# **Part V: Algebraic Effects**

# What is an Algebraic Effect?

- Separates state management from call site:
  - Computations can cause side effects:
    - I/O, FFI, Errors, etc.
  - Algebraic effects define ways to manage for these side effects
- Prior work:
  - Koka
  - Unison (abilities)
  - Eff
- Built on top of continuations:
  - An effect is a sum type that dispatches on continuation.

# Basic example

We define an effect with a number of operations:

- **effect raise** → **has operation raise**

We can use the operation like a function:

- **raise("hello")**

We implement a handler to implement the behaviour of the operation:

- **handler raise(continuation)** → **handle and decide how to continue**

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# Why are they a good fit for Passerine?

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When you think about it, an FFI is an effect handler implemented in another language.

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Passerine is an embedded language, so allows fine-grained control of side effects.

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Natural fit for fibers, and Passerine's concurrency model.

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Would simplify runtime interface upon implementation.

**Discuss!**

- 1. Traits in  $P_n$ ?**
- 2. Alg. Effs. in  $P_n$ ?**

# Moving forward

Easiest ways to contribute:

- Get acquainted with project layout; I'm always open for questions
- Help-wanted labels on GitHub
- Coming to this Weekly Meeting!
- Discussing ideas on Passerine Discord.

I'm a bit burnt out, which is why I haven't been actively hacking on Passerine's codebase.

Also at a bit of a crossroads: Passerine is very simple, and there are a lot of directions to take, we have to make a choice and set a course. (w.r.t. Alg. Effs. and Traits)





**Fin!**

Thanks for tuning in!



## **Next Up:**

Shaw will share  
his work on an  
alternative impl.  
for Passerine!