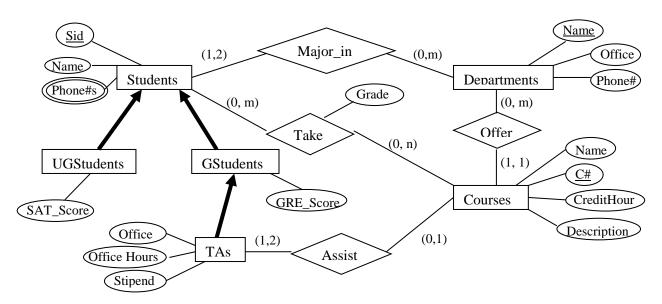
Sample Database Midterm Exam (60 min., closed book and closed notes)

- 1. (6 points) Explain the difference between a procedural query language and a non-procedural query language.
- 2. (9 points) Let E1 and E2 be two entity sets. List three factors that support the creation of a super entity type for E1 and E2 (i.e., E1 and E2 become sub entity types of the super entity type).
- 3. (9 points) Answer "true" or "false" to each of the following statements (no justification is required):
 - a. Given a relation R with a set of tuples T and let A denote an attribute of R. If all values under A in T are always different, then A must be a candidate key of R.
 - b. Suppose X and Y are two different subsets of the attributes of relation R. If X is the primary key of R and X is a subset of Y, then Y must be a superkey of R.
 - c. If A is a foreign key of relation R referencing the primary key B of relation S, then the number of distinct values under A (excluding null value) in R cannot exceed the number of distinct values under B in S.

Answer (Circle one for each question): (a) true false; (b) true false; (c) true false

4. (30 points) Transform the following ER diagram into relations. Use Method 1 to transform the IS-A hierarchy. You need to indicate the key and foreign key(s) of each relation.



5. (7 points) Let R(A, B, C) be a relation schema and $F = \{A \rightarrow B, B \rightarrow C, B \rightarrow A\}$ be a set of functional dependencies on R. Is the decomposition of R into R1(A, B) and R2(A, C) dependency-preserving? Justify your answer.

- 6. (9 points) Answer "Yes" or "No" to each of the following questions (no justification is required).
 - a. Let A be an attribute of relation R. If $A \rightarrow (attr(R) \{A\})$ is true, where attr(R) is the set of all attributes of R. Is A a candidate key of R?
 - b. Let A, B and C be attributes of a relation. Is there redundancy in $\{A \rightarrow BC, B \rightarrow AC\}$?
 - c. Suppose attribute A is a candidate key of relation R and X is a subset of attributes of R. Suppose further that $X \rightarrow A$. Must X be a candidate key of R?

Answer (Circle one for each question): (a) Yes No; (b) Yes No; (c) Yes No

- 7. Given a relation schema R(A, B, C, D, E, G) and a set of FDs $F = \{AB \rightarrow C, A \rightarrow DE, B \rightarrow BC, D \rightarrow CE\}$.
 - a. (6 points) Find the minimal cover Fmin of F. (Note that all trivial functional dependencies such as A → A are redundant.)
 - b. (6 points) Find all candidate keys of R. Show the dependency graph, Vni and Voi.
 - c. (6 points) Is the schema in 3NF? You must justify your answer.
 - d. (6 points) Apply Algorithm LLJ-DPD-3NF to decompose R regardless of whether R is already in 3NF.
- 8. (6 points) Let $R(A_1, ..., A_n)$ be a relation and X be a (sub)set of attributes of R. Argue that if $X \to A_1A_2...A_n$, then X is a superkey of R. (Note that our Lecture Notes have a theorem that says that if X can determine all attributes of a relation, then X is a superkey of the relation. You are not allowed to use this theorem to answer this question. This question actually asks you to prove the correctness of this theorem.) (Hints: Use the definition of functional dependency, the Unique Row Rule, and contradiction to make the argument.)