

We can verify the average length by considering the upper and lower bounds of the path length possible. For example, for data size of $N=2048$, the best case height is $\log_3(2048) = 6.94$ and worst case height is $\log_2(2048) = 11$. So, our observation, i.e. 8.05714 is correct.

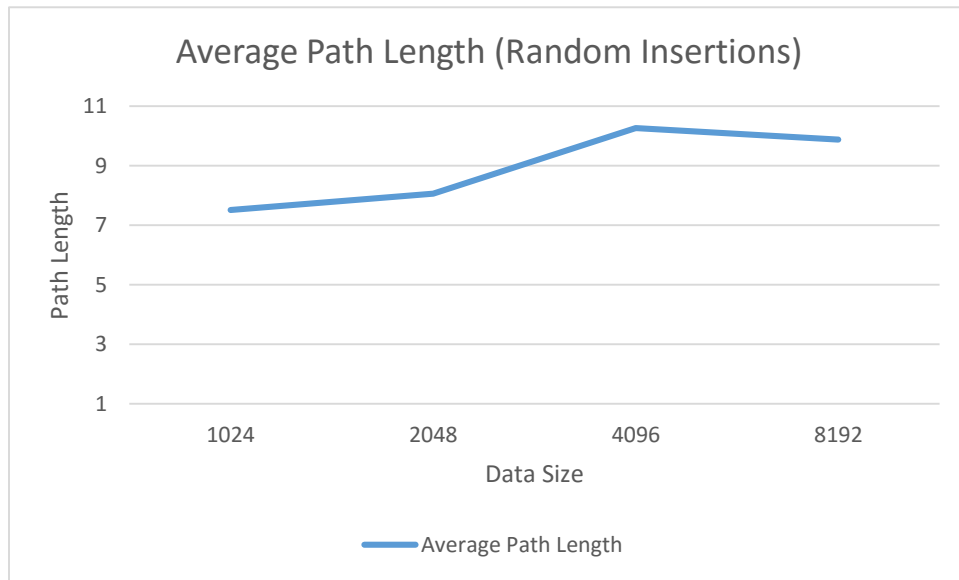


Figure 2 – Plot of Average Path Length for Random Insertions

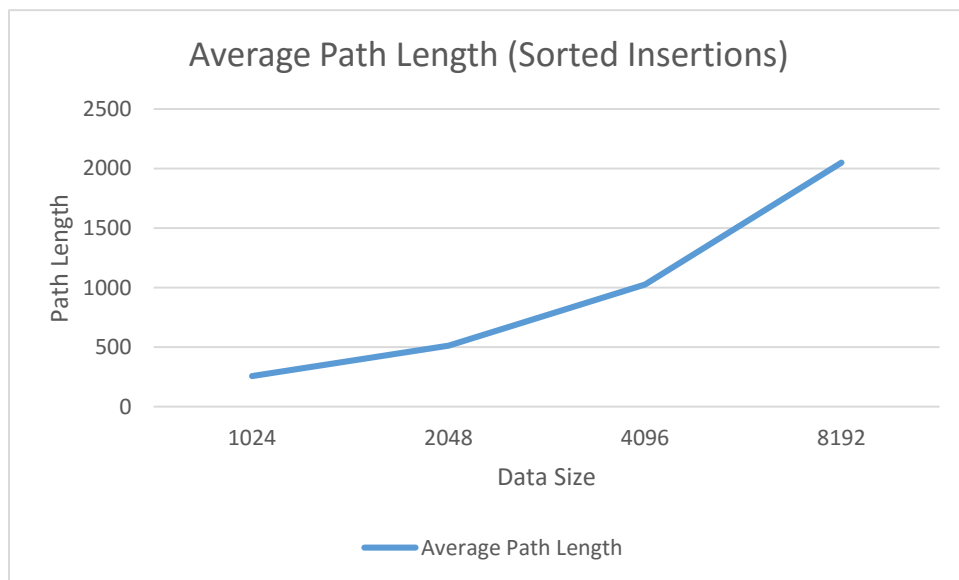


Figure 3 – Plot of Average Path Length for Sorted Insertions

Hypothesis =>

$$P = c * N^b$$

For Random Insertions,

Here, $P1 = 8.05714$ and $P2 = 7.50805$

$N1 = 2048$ and $N2 = 1024$

Therefore, $P1/P2 = 1.073 = (2048/1024)^b$

$b = 1.0731$ and $c = 0.00225$

Therefore, $P = 0.00225 * N^{1.0731}$

Similarly, for Sorted Insertions,

We get, $b = 0.998$ and $c = 0.25$

Therefore, $P = 0.25 * N^{0.998}$

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Q3. Write a program that computes the percentage of red nodes in a given red-black tree. Test program by running at least 100 trials of the experiment of increasing N random keys into an initially empty tree for $N=10^4$, 10^5 and 10^6 and formulate a hypothesis.

=> Red Black Tree implementation includes insertion of node and coloring it Red and then fixing the violation that may arise due to this insertion and moving towards the root while doing so. Tests were run against increasing data size of 10^4 , 10^5 , and 10^6 . Following is the average percentage of red nodes calculated as (# of Red Nodes) / (# of Red Nodes + # of Black Nodes).

Data Size	Average Percentage of Red Nodes
10000	0.484877
100000	0.486498
1000000	0.486583

Figure 4 – Table showing average percentage of red nodes

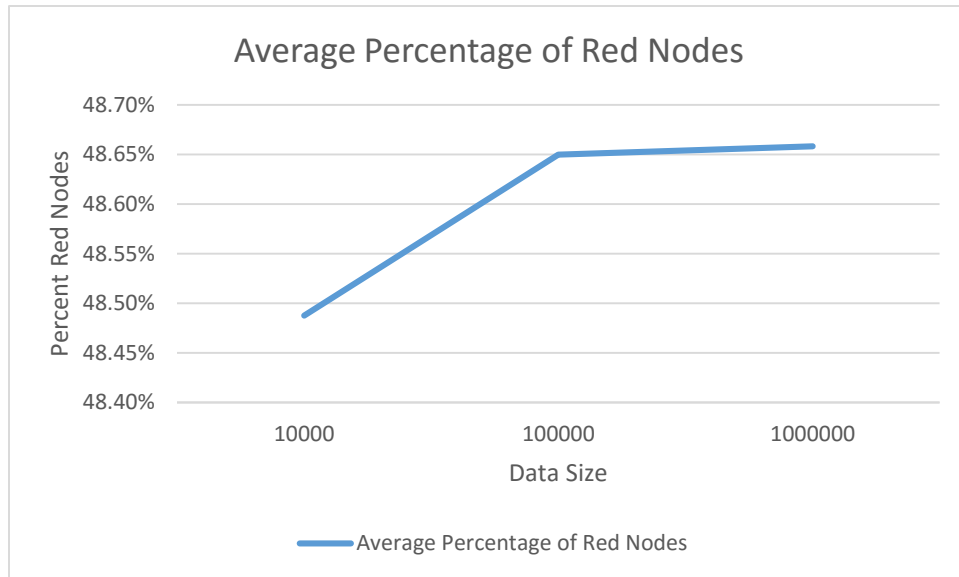


Figure 5 – Plot of Average Percentage of Red Nodes against Data Size

We can see that the average percent of red nodes remains *more or less constant in the range of 48.4% to 48.7%*.

Hypothesis =>

$$P = c * N^b$$

Here, $P_1 = 0.486498$ and $P_2 = 0.484877$

$N_1 = 100000$ and $N_2 = 10000$

Therefore, $P_1/P_2 = 1.003 = (100000/10000)^b$

$b = 0$ and $c = 0.48$

Therefore, **$P = 0.48$** (constant)

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Q4. Run empirical studies to compute the average and standard deviation of the average length of a path to a random node (internal path length divided by tree size) in a red-black BST built by insertion of N random keys into an initially empty tree, for N from 1 to 10,000. Do at least 1,000 trials for each size.

=> Following is the table that shows average and standard deviation of path length for different data sizes between 1 to 10,000 computed by executing 1000 trials. We can see that the average keeps on increasing (rapidly in the beginning and slowly afterwards) as expected with data size, but the standard deviation decreases slowly and settles eventually.

Data Size	Average	Standard Deviation
4	1	0
8	1.696875	0.084856
16	2.488563	0.053846
32	3.363313	0.046502
64	4.310659	0.055133
128	5.279415	0.045988
256	6.271172	0.039916
512	7.276039	0.038935
1024	8.287318	0.036058
2048	9.302272	0.034241
4096	10.31876	0.033071
8192	11.33704	0.033035

Figure 6 – Table showing average and standard deviation of path length

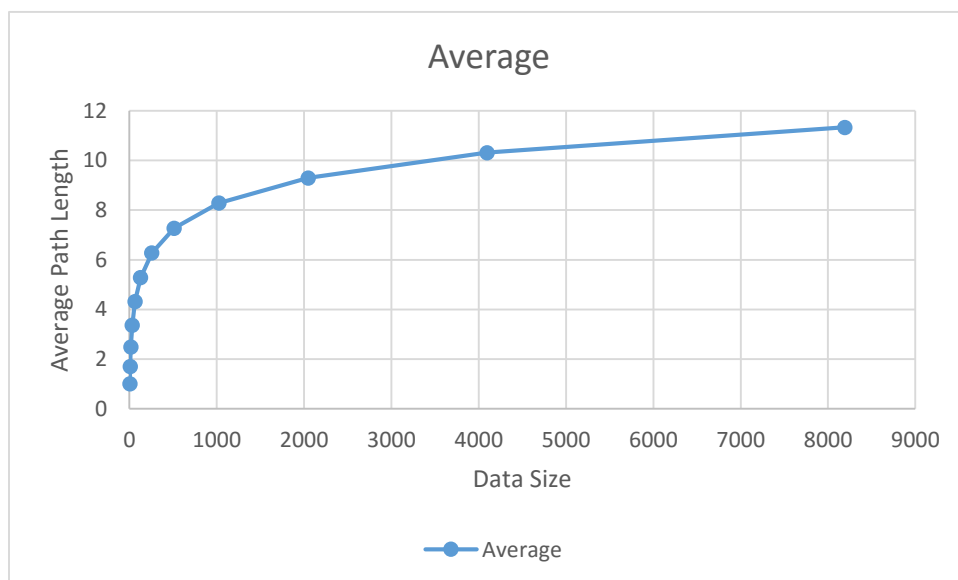


Figure 7 – Scatter plot - the trend average path length follows with increase in data size

Hypothesis => for average of path length

$$P = c * N^b$$

Here, $P_1 = 10.31876$ and $P_2 = 9.302272$

$N_1 = 4096$ and $N_2 = 2048$

Therefore, $P_1/P_2 = 1.1093 = (4096 / 2048)^b$

$b = 0.1496$ and $c = 2.9732$

Therefore, $P = 2.9732 * N^{0.1496}$ (constant)

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