CS307 System Practicum Assignment 2

Problem 1

Dining Philosophers Problem

Inferences -

Approaches used were -

Trivial Approach -

Each philosopher attempts to pick up the left fork first and then the second fork. Use an array of mutex locks to simulate forks which each philosopher can hold. Then each philosopher drops the left fork and then right. This is prone to deadlock as it is possible that each of the philosophers picks up the left fork and then every philosopher will be waiting for a right fork which will never become available.

In circular fashion order is

P0 F1 P1 F2 P2 F3 P3 F4 P4 F0

// Here, P stands for Philosopher and F for Fork.

Approach 1 (main.cpp)-

Change the trivial approach so that the odd numbered philosophers pick up the left fork first and the even philosophers pick up the right fork first. This will prevent any deadlocks from occurring as at least one philosopher can always acquire both forks. Starvation also does not appear to be a problem here but can occur if say P0 is waiting for P1 to drop his left fork but after dropping P1 again picks it up. This is highly unlikely because of the random think times introduced in the question. The main problem with this solution is that P4 will eat the most in this case. This is because P0 and P1 will fight for F1 and P2 and P3 will fight for F3 but there will always be both forks available for P4, i.e. F0 and F4 will always be initially available. Thus, this solution gives P4 a slight advantage over others.

Approach 2 (main2.cpp)-

Change the trivial approach so that all philosophers except the last one pick up the right fork first. The last philosopher picks up the left fork first. This will prevent any deadlocks from occurring as at least one philosopher can always acquire both forks. Starvation is also unlikely because of the random times of thinking. So, in our case P0, P1, P2 and P3 pickup the right fork first and P4 picks up the left fork first. Here also, we observe that P0 does the most eating in each case because P4 and P3 will fight for F4 and both F0 and F1 will always be available for P0 initially. Thus, this solution gives P0 a slight advantage over others.

How to fix uneven distribution?

This is highly unlikely to solve in the classical problem. In each solution to this we have to modify the problem by some amount.

- 1. We can introduce a waiter (queue) which keeps hold of how much food each philosopher has eaten and assign forks accordingly.
- 2. We can also modify Approach 2 we used above to cycle the philosopher with the opposite fork picking pattern after some interval of time, say every 2 minutes. This will ensure that each philosopher gets a fair chance and will ensure a fair distribution.

Observations -Approach 1 (main.cpp)out1.txt

Philosopher	Thinking	Eating	1 Fork Acquired	
P0	616	616	616	
P1	634	635	635	
P2	649	649	649	
Р3	673	673	673	
P4	692	693	693	

out2.txt

Philosopher	Thinking	Eating	1 Fork Acquired	
P0	614	614	615	
P1	639	639	639	
P2	648	648	648	
Р3	664		665	
P4	699	699	700	

out3.txt

Philosopher	Thinking	Eating	1 Fork Acquired	
P0	612	612	613	
P1	636	637	637	
P2	648		648	
P3	3 670		671	
P4	698	698	698	

out4.txt

Philosopher	losopher Thinking		1 Fork Acquired	
P0	614	614	614	
P1	640	640	640	
P2	P2 650		650	
P3	P3 664		665	
P4	697		697	

out5.txt

Philosopher	losopher Thinking		1 Fork Acquired	
P0	620	620	620	
P1	636	636	636	
P2	653	654	654	
Р3	661	661	661	
P4	694	695	695	

Approach 2 (main2.cpp)-

out2_1.txt

Philosopher	Thinking	Eating	1 Fork Acquired		
P0	691		691		
P1	680	680	680		
P2	664	665	665		
P3	P3 618		618		
P4	4 658		658		

out2_2.txt

Philosopher	Philosopher Thinking		1 Fork Acquired	
P0	698	698	699	

P1	679	679	679
P2	668	668	668
P3	615	615	615
P4	650	651	651

out2_3.txt

Philosopher	sopher Thinking		1 Fork Acquired	
P0	688	688	688	
P1	671	671	672	
P2	P2 648		648	
P3	P3 595		596	
P4	4 641		641	

out2_4.txt

Philosopher	Thinking	Eating	1 Fork Acquired	
P0	700		700	
P1	665	665	665	
P2	P2 642		642	
P3	P3 599		600	
P4	637		637	

out2_5.txt

Philosopher	Thinking	Eating	1 Fork Acquired	
Р0	692	692	693	
P1	678	678	678	
P2	671	672	672	
Р3	612		613	
P4	656	657	657	

Problem 2

Matrix Multiplication

Observations -

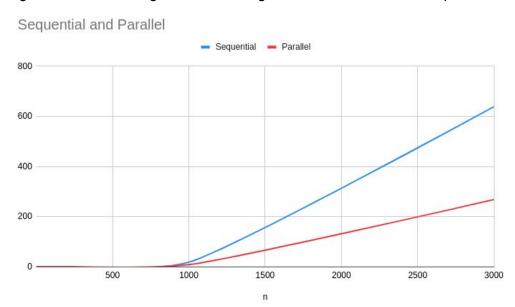
Parallel program runs with 4 threads All times are in seconds

	n=	=1	n=	50	n=2	200	n=1	000	n=3	000
	Sequen tial	Parallel	Sequen tial	Parallel	Sequen tial	Parallel	Sequen tial	Parallel	Sequen tial	Parallel
	0.0000	0.0005 7741	0.0075 0554	0.0069 76	0.1230 2338	0.0690 1864	18.077 2234	8.1692 8457	643.63 5285	267.52 91197
	0.0000	0.0007	0.0072	0.0068	0.1229	0.0706	18.024	8.0096	639.38	266.06
	0082	222	9521	415	6388	0032	07869	2116	174	54095
	0.0000	0.0006	0.0074	0.0052	0.1297	0.0651	17.939	8.0649	632.38	267.00
	0096	7246	6023	7879	082	6734	85098	3118	27295	17897
	0.0000	0.0005	0.0075	0.0069	0.1184	0.0707	18.256	8.1038	635.49	270.26
	0076	5715	4953	5252	7829	6258	51759	7556	6553	24048
	0.0000	0.0005	0.0072	0.0066	0.1244	0.0641	18.282	8.0954	641.73	267.40
	0104	1749	4737	8165	0977	5848	75159	2102	54021	19674
Mean:	0.0000	0.0006	0.0074	0.0065	0.1237	0.0679	18.116	8.0886	638.52	267.65
	00876	09342	11576	46092	16704	41472	08445	26698	63419	21382

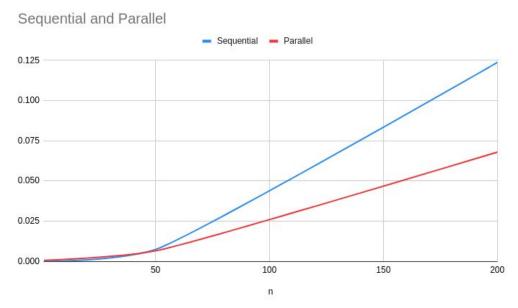
n	Sequential	Parallel
1	0.000000876	0.000609342
50	0.007411576	0.006546092
200	0.123716704	0.067941472
1000	18.11608445	8.088626698
3000	638.5263419	267.6521382

Inferences -

For small matrices (size < 50x50), the sequential program executes faster since the thread creation overhead (time taken to create and join threads) is more than multiplication time. For larger matrices, dividing the work among the cores finishes the multiplication faster.



Graph between time taken and n for n=1 to n=3000



Magnified Above Graph for n=1 to n=200. We can observe that sequential programs run faster for smaller values of n.

Problem 3

Linux Kernel

Kernel compilation and installation

- Linux kernel version 5.9.1
- After downloading and extracting the source code, I used `make gconfig` (similar to `make menuconfig` but provides GUI) to configure what to include in kernel. I removed some drivers like GPU drivers and manufacturer-specific drivers which are not required because it is running inside a VM.
- Then used `make` and `make modules_install` to compile modules, and `sudo make install` to install the compiled kernel.
- Rebooted into new kernel, and installed `linux-headers-5.9.1` package and dependencies from the Debian repository.

Kernel modules

- Module NULLdereference dereferences a NULL pointer. It causes a kernel oops an error that the kernel can recover from and the system can continue working.
- Module panic calls the panic() function defined in the kernel libraries. This simply causes a kernel panic which is a non-recoverable error and freezes the system. The kernel can be configured to reboot in case of a kernel panic, but by default nothing is specified, so the only option in default settings, after a kernel panic is to power off the machine.

Output from sudo dmesg after sudo insmod NULLdereference

```
[ 9847.880983] Loading custom kernel module.
[ 9847.881004] BUG: kernel NULL pointer dereference, address:
00000000000000000
[ 9847.881014] #PF: supervisor read access in kernel mode
[ 9847.881020] #PF: error_code(0x0000) - not-present page
[ 9847.881024] PGD 0 P4D 0
[ 9847.881035] Oops: 0000 [#1] SMP PTI
[ 9847.881046] CPU: 0 PID: 86252 Comm: insmod Tainted: G OE
5.9.1 #1
[ 9847.881052] Hardware name: innotek GmbH VirtualBox/VirtualBox, BIOS
VirtualBox 12/01/2006
[ 9847.881066] RIP: 0010:init_module+0x15/0x2c [NULLdereference]
[ 9847.881074] Code: Bad RIP value.
```

```
9847.881081] RSP: 0018:ffffad7506397c50 EFLAGS: 00010246
9847.881087] RAX: 00000000000001d RBX: 00000000000000 RCX:
00000000000000000
9847.881093] RDX: 000000000000000 RSI: ffff8de3dd418cc0 RDI:
ffff8de3dd418cc0
9847.881097| RBP: ffffad7506397c50 R08: ffff8de3dd418cc0 R09:
00000000000000004
fffffffc08bd000
[ 9847.881107] R13: ffff8de3a79e0070 R14: ffffffffc08bf000 R15:
00000000000000000
[ 9847.881113] FS: 00007f607f79b540(0000) GS:ffff8de3dd400000(0000)
knlGS:000000000000000000
[ 9847.881120] CS: 0010 DS: 0000 ES: 0000 CR0: 0000000080050033
[ 9847.881125] CR2: ffffffffc08bcfeb CR3: 00000000052d4004 CR4:
00000000000706f0
[ 9847.881140] Call Trace:
[ 9847.881160] do one initcall+0x4a/0x1fa
[ 9847.881214] ? kmem cache alloc trace+0x17e/0x2f0
[ 9847.881234] load_module+0x280c/0x2b40
9847.881248] do sys finit module+0xbe/0x120
[ 9847.881255] ? do sys finit module+0xbe/0x120
 9847.881267] x64 sys finit module+0x1a/0x20
[ 9847.881291] do syscall 64+0x38/0x90
[ 9847.881302] entry SYSCALL 64 after hwframe+0x44/0xa9
[ 9847.881309] RIP: 0033:0x7f607f8e089d
[ 9847.881320] Code: 00 c3 66 2e 0f 1f 84 00 00 00 00 00 90 f3 0f 1e fa 48
89 f8 48 89 f7 48 89 d6 48 89 ca 4d 89 c2 4d 89 c8 4c 8b 4c 24 08 0f 05
<48> 3d 01 f0 ff ff 73 01 c3 48 8b 0d c3 f5 0c 00 f7 d8 64 89 01 48
[ 9847.881326] RSP: 002b:00007ffc31fb00a8 EFLAGS: 00000246 ORIG RAX:
0000000000000139
[ 9847.881334] RAX: ffffffffffffffda RBX: 000055c84dc177b0 RCX:
00007f607f8e089d
9847.881339] RDX: 000000000000000 RSI: 000055c84cb38358 RDI:
00000000000000003
 9847.881344] RBP: 0000000000000000 R08: 00000000000000 R09:
00007f607f9b4260
```

```
[ 9847.881349] R10: 000000000000000 R11: 000000000000246 R12:
000055c84cb38358
[ 9847.881354] R13: 00000000000000 R14: 000055c84dc17750 R15:
00000000000000000
[ 9847.881362] Modules linked in: NULLdereference(OE+) btrfs
blake2b generic xor zstd compress raid6 pq ufs qnx4 hfsplus hfs minix ntfs
msdos jfs xfs libcrc32c cpuid vboxvideo(OE) nls iso8859 1 vmwgfx
snd intel8x0 snd ac97 codec ac97 bus snd pcm snd seq midi ttm
snd seq midi event intel rapl msr drm kms helper snd rawmidi
intel rapl common snd seq cec crct10dif pclmul snd seq device
ghash clmulni intel fb sys fops aesni intel snd timer syscopyarea
crypto simd snd sysfillrect cryptd glue helper sysimgblt soundcore rapl
joydev input leds video mac hid serio raw sch fq codel parport pc ppdev lp
drm parport ip tables x tables autofs4 hid generic usbhid hid psmouse
e1000 pata acpi crc32 pclmul ahci libahci i2c piix4
[ 9847.881461] CR2: 0000000000000000
[ 9847.881470] ---[ end trace 543e4fc2a480d1d6 ]---
```